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TECHNICAL MANUAL

## PROGRAMMING MANUAL

FOR

COMPUTER, GUN DIRECTION, M18 (FADAC)

HEADQUARTERS, DEPARTMENT OF THE ARMY

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## GENERAL DESCRIPTION

### 1.1 PURPOSE

The purpose of this publication is to describe techniques required to program the Computer, Gun Direction, M18 (FADAC).

### 1.2 PHYSICAL AND OPERATIONAL CHARACTERISTICS

FADAC is a general purpose, transistorized digital computer designed for field use. By proper programming, it may be used to perform a variety of tasks. Information flow is serial by bit, parallel by function, allowing up to 12,800 operations per second.

Size.
24 inches by 14 inches by 34 inches.
Weight.
Approximately 210 lbs .
Power.
Three-phase, 4 wire, 400 Hz system, 120 volts line to neutral/208 volts line to line, approximately 750 watts.

Temperature.
$-25^{\circ} \mathrm{F}$ to $125^{\circ} \mathrm{F}$ (external ambient at sea level); with rear cover installed, to $-40^{\circ} \mathrm{F}$. Automatic temperature protection is provided.

Commands.
One command per word; each command contains both address of operand and address of next instruction ( $1+1$ system).

Numbers.
Straight binary for internal operations; automatic conversion to other codes for input-output; two's complement notation for negative numbers.

Word Length.
Thirty-two programmable bits; sign bit and 31 magnitude bits for numerical value.

Memory Type.
Magnetic Disc, 6000 rpm nominal speed of rotation.

Storage Capacity.
Sixty-four channels of 128 words each ( 8192 words) in main memory. Also provided are two 16 word highspeed loops for rapid access, three 1 -word registers for arithmetic operations and control, and one 2 -word register for output display information storage. All channels and loops have one read and one write head, except the 16 -word loops which have an additional read head.

## Pulse Repetition Rate.

Nominal pulse repetition rate, 460 kilopulses per second.

Access Times.
One word-time is 78 to 83 micro-seconds. One disc revolution requires approximately 10 milliseconds .

Input.
Input to the computer is from the manual keyboard or the mechanical tape reader on the control panel, at approximately 10 characters/second, or from external sources, such as: Reproducer, Signal Data, AN/GSQ-64, a high speed photo electric reader designated the Signal Data Reproducer (SDR) (600 characters/sec); another FADAC; magnetic tape ( 4000 char/sec max); other devices. Input may be in 5 -channel Teletype or 8 -channel Fieldata code.

## Output.

FADAC can transmit data to either the Nixie tube visual display unit on the control panel, or through the output plug, to any output device such as a printer, paper tape punch unit, another FADAC, or other Fieldata or Teletype equipment. Output information is in 5-level teletype, 8 level Fieldata or 2-wire Teletype codes, as required. In addition to the standard codes, FADAC can be programmed to generate any 8 -bit code. Output rates depend on the output code selected, as well as the output equipment. Maximum rates are as follows:
(1) $4000 \mathrm{char} / \mathrm{sec}$ in $\mathrm{A} / 6$
(2) 3550 char/sec in A/5 and A/4
(3) $650 \mathrm{char} / \mathrm{sec}$ in octal or BCD

## Additional Features of the Computer.

(1) Parity check on information transfers.
(2) Verify indicator (located on the SDR) on input in Program-Fill mode.
(3) Gating for extemal support equipment to permit logic failure isolation to single printed circuit board.
(4) Marginal test circuit forpreventive maintenance.
(5) Voltage transient warning indicator.
(6) Temperature warning indicator.
(7) High speed (2 bits-at-a-time) multiplication, division, and shifts.
(8) Self-checking for overflow.
(9) Standardized printed circuit boards to minimize logistic problems.

### 1.3 FUNCTIONAL DESCRIPTION

A functional diagram of the FADAC system is shown in Figure 1.1.

## Control Unit.

The computer control unit interprets and processes the programmed instructions. Each instruction word called from memory is processed by the control logic through two non-program-addressable registers, X and I. During execution of the current instruction in the Iregister, the next instruction is stored in the X -register. On completion of execution of the current instruction, the next instruction is sent to the I-register, and the new next instruction is read into the X-register. If the next instruction is found just as execution of the current instruction is completed, the X-register is bypassed, and the instruction is sent directly to the Iregister. (See discussion of Multiplexing, Section 3.1.)

## Arithmetic Unit.

The Arithmetic Unit performs the actual work involved in problet. solution. In addition to the four arithmetic operations, this section can shift and cycle numbers right and left, and assist in operations which make it possible for the computer to make decisions, to accept and store input information, and to format and transmit output information. The Arithmetic Unit is


Arrows show flow of information

Figure 1.1. Functional Diagram of System
comprised of arithmetic logic, input/output logic, and three one-word, program-addressable registers: the Accumulator (A), the lower Accumulator (L), and the Number register (N).

The A-register is the principal arithmetic register in FADAC. It is employed in all arithmetic, logical, input/output, and decision making operations.

The L-register is an extension of the A-register in some operations (as in Multiply and Divide); and functions independently of $A$ in others (as in Zero L, Store L). It also serves as the location counter for input/ output operations.

The N-register holds the second operand for some operations (as in Multiply and Divide). It serves as an input/output buffer, and is used in program control transfer operations to store the transfer instruction, if that instruction is flagged.

The A, L, and N-registers are physically part of the FADAC memory disc.

Memory Unit.
The FADAC memory unit stores information in the form of instruction words and data words. It is comprised of 8192 words of main memory, plus two 16 -word rapid access loops ( R and Q ) and one 2-word display loop (D).

Input.
Data and instructions can be placed in FADAC memory using any of the devices listed in Section 1.2. The Signal Data Reproducer (SDR) is the primary input device.

FADAC output is transmitted to the NIXIE display panel, or through the FADAC plug ( $\mathrm{J}-10$ ) to any external device, as listed in Section 1.2.

### 1.4 DESCRIPTION OF MEMORY

## General.

FADAC memory is a rotating disc, coated with ferrous oxide similar to the coating on conventional magnetic tape. The disc rotates at approximately 6000 rpm under stationary read and write heads. To record information on the disc, an electrical signal generates a magnetic field in a write head. This field produces magnetization on the disc which remains, even when the computer is de-energized, until replaced by new information. The read heads are similar to the write heads, but operate in reverse. A change in magnetization on the disc produces an electrical current in the head. These signals produce the 1 's and 0 's for the binary operation of the computer.

## Memory Layout.

A conceptual layout of the FADAC memory is shown in Figure 1.2. Although the actual position of the heads and the relative positions of the tracks are not exactly as shown, this diagram provides a useful overall pictorial diagram. The disk is designed in channels, each equipped with read and write heads. The sector track is permanently recorded and has a read head only. The purpose of the sector track is to supply data which notifies the control unit which sector is about to pass under the read heads for the other channels. The main memory read heads are all at the same effective sector location at any given point in time.


Figure 1.2. Conceptual Layout of FADAC Memory

Sector Track Contents.
The sector track contains 128 words of factory recorded information, available serially from a read flip-flop. A binary number appears in every word of the sector track. In each successive word the number is increased by 1 , such that the octal numbers 0 to 177 are represented. These numbers serve as reference for numbering the memory sectors. Also recorded on the sector tract are the Teletype (TT) and Fieldata (FD) tape codes, in inverted format, for each of the 16 BCD and Octal characters recognized by FADAC. The four least significant bits of each sector number represent the BCD and Octal translation of the TT and FD tape code for a given character. The tape codes are repeated eight times around the sector track. See Section 2.2.7 for a description of how this information is used during I/O operations and Appendix C for a list of BCD and OCTAL codes.

## Main Memory.

The 64 main memory channels are addressed by the even octal numbers 0 through 136 and 300 through 336. The 128 words or sectors of each channel are addressed by consecutive octal numbers 0 through 177 inclusive. To clarify discussion, the channels in main memory will be referenced throughout this manual as follows:

> Memory I $=$ Channels 0 through 76
> Memory II $=$ Channels 100 through 136
> Memory III $=$ Channels 300 through 336

Hot and Cold Memory.
Main memory is further divided into Working Storage (Hot Memory) in which both the read and write heads are energized, and Permanent Storage (Cold Memory) in
which the read heads ONLY are energized. At programmer's option, the size of working storage may be specified as either 4, 12, or 16 channels, and the FADAC "Hot Channel Select Switch" (see Section 4.5) manually set to energize the write heads of specific Hot Channels, as listed below:

Working
Storage Size

Hot Channels
$\frac{\text { Set }}{\text { (In Octal) }}$

$$
\left\{\begin{array}{c}
70-76 \\
70-76 \\
110-11.6 \\
130-136
\end{array}\right.
$$

$$
40-76
$$

Wi th the SDR connected to FADAC and the Auxiliary Memory Switch on the SDR in the OFF position, all write heads in main memory are energized to allow all of memory to be filled. However, during normal field use, memory is filled, the SDR disconnected, and only the contents of Hot memory can be varied. This design was adopted to assure that the contents of permanent memory are not inadvertently destroyed because of operator error or power failure during field use.

### 1.4.1 PROGRAM-ADDRESSABLE SPECIAL REGISTERS

Description.
The $A, L, N, D, R$, and Q-regi-ters are short,
recirculating loops. Each bit is sensed by the read head and rewritten by the write head. Each word is shifted, bit-by-bit, through the loop. The loop information remains the same unless changed by some arithmetic or control operation. A, L, and N are 1 -word loops whose contents are rewritten each word time. $D$ is a 2-word loop, the contents of which are rewritten two sectors ahead of the read head. Both the $R$ and $Q$-loops have an intermediate read head located in the middle of the loop. This head is not involved in the recirculation process. However, it is used automatically by FADAC to expedite execution of a command which causes the contents of a sector in the R or Q -loop to be read.

The $R$ and $Q$-Loops.
The rapid access loops, $R$ and $Q$, may have octal sector addresses 0 through 177. However, only the last four bits, ( 00 through 17) are recognized. When addressing the R and Q -loops, therefore, only sectors 00 through 17 should be used. Note the contents of each sector of the $R$ and the $Q$-loops are repeated eight times around the $R$ and $Q$-loop tracks. The channel address of the $R$-loop is 142.00 and of the $Q$-loop, 152.00.

## One-Word Loops.

The A, L, and N-registers each effectively occupy a full track on the FADAC memory disc. The contents of each loop are always immediately available under a read head. They are addressed as $170.00,172.00$ and 174.00 respectively.

The D-Register.
When the main memory read heads are over any even numbered sector, the contents of D.O are under a read head; while when the main memory read heads are at any odd numbered sector, the contents of D. 1 are
available, since the D-register also etfectively occupies a full track on the memory disc. Since the read and write heads of the D-loop are spaced two sectors apart both heads are effectively at the same sector of $D$ at any given word time. The D -loop address is 160.000 (D.0) and 160.001 (D.1).

Summary of Register Address.
All register addresses are summarized in Table 1.1.

TABLE 1.1 SUMMARY OF REGISTER ADDRESSES

| Register | Channel <br> Address <br> (Octal) | Sector <br> Address <br> (Octal) | Description |
| :---: | :---: | :---: | :---: |
| R | 142 | 000-017 | 16 words |
| Q | 152 | 000-017 | 16 words |
| D. 0 | 160 | 000 or even sec. | 1 word |
| D. 1 | 160 | 001 or odd sec. | 1 word |
| A | 170 | 000 | A-register (Accumulator) |
| L | 172 | 000 | L-register (Lower |
| N | 174 | 000 | N-register (Number <br> Accumulator) |

### 1.5 CONTROL BUTTONS

The eight control buttons located on the front panel of FADAC are effectively wired directly to the memory locations shown in Table 1.2.

## CONTROL BUTTON

ADDRESS

| Control Buttons | Address (Octal) |
| :--- | :--- |
| Sample Matrix | $000 \quad . \quad 000$ |
| Test | $000 \quad . \quad 001$ |
| Set Up | $000 \quad . \quad 002$ |
| Recall | $000 \quad . \quad 003$ |
| Send | $000 \quad$. |
| Compute | $000 \quad$. |
| Trig | $000 \quad . \quad 005$ |
| Receive | $000 \quad$. |

When a control button is depressed, the flip-flop is set for Memory II, FADAC enters the compute mode, and the instruction stored in the control button location is executed. Computation is usually initiated in this manner. Note that the control button locations may be called in a program and their contents executed under program control.

### 1.6 NUMBER SYSTEMS USED

Number systems applicable to the FADAC computer are the decimal system, the octal system and the binary system. Information going into or coming from the computer is either decimal, octal or alphabetic. However, the computer operates internally with binary numbers only, and all negative numbers are represented in two's complement notation.

## CHAPTER 2

## WORD FORMATS AND INSTRUCTION REPERTOIRE

### 2.1 WORD FORMATS

## General.

Information stored in FADAC memory contains no inherent distinction between numerical data and computer instructions. Both are stored as 36 -bit binary coded numbers, of which 32 bits are programmable, and 4 a.e machine-generated. Of the 4 non-programmable bits, 1 is the parity bit, and 3 are used for synchronization. When a word is read from memory, the state of the computer logic will determine whether the information is to be interpreted as a data word or an instruction.

FADAC Octal Formats.
Each 32 bit FADAC word may be represented by 11 FADAC Octal digits. The first digit is formed from the two most significant bits of the binary word, and has a maximum value of 3 . Each of the remaining digits is formed from three bits of the FADAC word, and has a maximum value of 7 . Because of the short (i.e., two bits only) first digit in the FADAC format, a pure octal number and its FADAC-Octal representation are not the same.

## Data Words in FADAC-Octal.

Information used as an operand is interpreted by the computer as a signed, 31 bit, binary-coded word, with a fixed point between the sign (bit position $O$ ) and the most significant bit. (See Fig. 2.1.) Positive numbers have a zero sign bit, and negative numbers a one. Negative numbers are represented in two's complement form.

a) Bit Position
b) Max. Binary Value
c) Max. Octal Value
d) FADAC Octal Position

Figure 2.1. FADAC Octal Format
Instruction Word Format in FADAC-Octal.
Commands are stored in the memory in the form of one instruction per computer word, represented by 11 FADAC-Octal digits. Each command consists of four separate fields; a flag, a next instruction address, an operation code (op code), and an operand address. (See Fig. 2.2.)

(A) Bit Position
(B) FADAC-Octal Position

Figure 2.2. Instruction Word Format

The flag, which is in the sign position, provides for the modification of a FADAC instruction. The use of the flag is described for each instruction in the FADAC instruction list.

The Next Instruction Address, which tells the computer where to go next in the program sequence, consists of a channel and a sector portion. The channel portion consists of 7 bits. These are represented by 3 octal digits, which are evenly numbered from 000 to 174. The sector portion consists of 7 bits. These are represented by 3 octal digits which are numbered sequentially from 000 to 177 .

The Operation Code, which represents the operation to be performed, usually consists of 6 bits. These are represented by 2 octal digits which are evenly numbered from 00 to 76 . The shift, cycle, I/O, and most special instructions require a special format which is described in the sections covering these commands.

The Operand Address is the address of the number to be operated upon. It consists of a channel and sector portion and is in the same form as the Next Instruction Address.

A total of 15 octal digits seems to be required to represent an instruction word: 1 for the flag, 2 for the operation code, and 6 each for the addresses. Since a FADAC word is limited to only 11 digits, the four extra digits are combined in the following manner:

Let:
$x_{1}$ represent the octal digits in each of the four fields, and $y_{1}$ represent the FADAC Octal digits:

F ADAC Instruction Word

| Flag | Next Instruction |  | Op. Code | Operand Address |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | Channel | Sector |  | Channel | Sector |
| ${ }^{\text {x }} 1$ | $\mathrm{x}_{2} \mathrm{x}_{3} \mathrm{x}_{4}$ | $\mathrm{x}_{5} \mathrm{x}_{6} \mathrm{x}_{7}$ | $\mathrm{x}_{8} \mathrm{x}_{9}$ | $\mathrm{x}_{10} \mathrm{x}_{11} \quad \mathrm{x}_{12}$ | $\mathrm{x}_{13} \mathrm{x}_{14} \mathrm{x}_{15}$ |
| ${ }^{y_{1}}$ | $y_{2} y_{3}$ | $\mathrm{y}_{4} \mathrm{y}_{5}$ | $\mathrm{y}_{6} \mathrm{y}_{7}$ | $\mathrm{y}_{8} \mathrm{y}_{9}$ | ${ }^{\mathrm{y}} 10{ }^{\mathrm{y}}{ }_{11}$ |

Where:

$$
\begin{array}{ll}
y_{1}=x_{1}+x_{2} \\
y_{2}=x_{3} & \frac{y_{7}=x_{9}+x_{10}}{y_{3}=x_{4}+x_{5}} \\
y_{4}=x_{6} & \frac{y_{8}=x_{11}}{y_{9}=x_{12}+x_{13}} \\
y_{6}=x_{8} & y_{10}=x_{14} \\
y_{11}=x_{15}
\end{array}
$$

Note that all channels and operation codes are-evenly numbered. Thus whenever an odd number ends a channel address field, a sector greater than or equal to 100 is called (i.e., $y_{3}=x_{4}+x_{5}$ or $y_{9}=x_{12}+x_{13}$ ). Whenever an odd number ends the op code field, the channel of the operand address is $\geq 100$ (i.e., $y_{7}=x_{9}+x_{10}$ ).

The first digit is derived as follows:

$$
y_{1}=x_{1}+x_{2}
$$

$\mathrm{y}_{1}=0$ when no flag is called and NI channel is $<100$
$y_{1}=1$ when no flag is called and NI channel is $\geq 100$
$\mathrm{y}_{1}=2$ when a flag is called and NI channel is $<100$
$y_{1}=3$ when a flag is called and NI channel is $\geq 100$
To illustrate: Given the instruction "Clear and Add (24) the contents of channel 102 , sector 133 , and go to channel 112, sector 123 ," the word format is as follows:

| Flag | Next Instruction |  | Op. <br> Code | Operand Address |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | Channel | Sector |  | Channel | Sector |
| 0 | 11 | 123 | $\begin{array}{cc} 2 & 4 \\ (C L A) \end{array}$ |  | $133$ |
| 1 | 13 | 23 |  | 03 | 33 |

The FADAC octal word for this command is;
11323250333
Or consider the coding required for the instructions "Subtract" (02) the contents of channel 112, sector 113 , and go to channel 112 , sector $024 . "$

| Flag | Next Instruction |  | Code | Operand Address |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | Channel | Sector |  | Channel | Sector |
| 2 | 112 | 024 | $\left.\left\lvert\, \begin{array}{cc} 0 & 2 \\ (S \cup B * \end{array}\right.\right)$ | 112 | $113$ |
| 3 | 12 | 24 | 03 | 13 | 13 |

Octal

FADAC. Octal

The FADAC octal word for this command is:

In addition to Octal data words, information may be input and/or output in BCD, Alpha-4, Alpha-5, or Alpha6 code. Fig. 2-3 illustrates the data word format for each of these codes.

| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 | 21 | 22 | 23 | 24 | 25 | 26 | 27 | 28 | 29 | 30 | 31 | (a)(b) |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 |  |  |  | 2 |  |  |  |  |  |  |  | 4 |  |  |  | 5 |  |  |  | 6 |  |  |  | 7 |  |  |  | 8 |  |  |  |  |
| i |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  | 4 |  |  |  |  |  |  |  | (c) |
|  |  | 1 |  |  |  |  |  | 2 |  |  |  |  |  | 3 |  |  |  |  |  | 4 |  |  |  |  |  | 5 |  |  |  |  |  | (d) |
| 1 |  | 2 |  |  |  |  |  | 3 |  |  |  |  |  | 4 |  |  |  |  |  | 5 |  |  |  |  |  | 6 |  |  |  |  |  | (e) |

(a) Bit Position
(b) BCD Code $=4$ bits/character; 8 characters/ FADAC word
(c) Alpha-4 Code $=8$ bits/character; 4 characters/ FADAC word (First bit of each character is inverted on output)
(d) Alpha-5 Code $=6$ bits/character; 5 characters/ FADAC word (Bits 0 and 1 are ignored)
(e) Alpha-6 Code $=6$ bits/character; 6 characters/ FADAC word ( 4 most significant bits of first character are stripped on input, machine-generated on output)

Figure 2.3. Data Word Formats in BCD and Alphanumeric Codes

### 2.2 INSTRUCTION REPERTOIRE

General.
All FADAC instructions fall into seven general categories according to function:

Arithmetic
Store and Load
Transfer
Shift \& Cycle
Logic and Special
Discrete Input/Output
Parallel Input/Output
Execution time for each instruction is shown in the Word Time portion of the description. Since one word time is approximately 78 to 83 microseconds, timing may be obtained by multiplying $80 \times 10^{-6}$ by the number of word times required to execute each instruction.

In the instruction descriptions, the following conventions are used:
(1) $M$ represents the operand address of the instruction being executed.
(2) (M)represents the contents of location M.
(3) C represents the channel address of the operand address, or the 2 least significant digits of the 4 digit extended op codes required for shift, cycle, I/O and most special commands.
(4) S represents the sector address of the operand address in all categories except in instructions requiring a special format.
(5) A flagged instruction is represented in mnemonic form by writing an asterisk following the operation code, and in machine code by placing a 1 in bit position O. For example: CLA, CLA*, STO, STO*, TRA, TRA*. The effect of the flag is specified in the description of each instruction.
(6) Unless explicitly mentioned under the description of each command, special registers are
not affected by the execution of a given command.
(7) All numbers are Octal unless otherwise specified, except numbers used to identify bit positions, which are designated in decimal.

### 2.2.1 ARTHMETIC

General.
If overflow occurs upon execution of the following instructions, the computer will automatically flash the Error light indicator on the control panel. The next instruction will be recognized only if it is a TRANSFER ON OVERFLOW (TOV) instruction, in which case the Error light indicator will be turned off. If any other instruction follows, the computer will halt and the Error light will continue to flash.

Rules for the Sign in Arithmetic Operations.
The rules for the sign of an arithmetic operation are the same as those of algebra. A negative zero is never generated by an FADAC arithmetic operation. However, the number 20000000000, which is not interpreted as a negative zero, may be input to FADAC and operated upon.
ADD $\quad 00$ ADD M Word Time $=1$

The algebraic sum of the contents of $M$ and the contents of the A-register replace the contents of the A-register. Overflow can occur. A flag bit of 1 will cause (M) to be put into the N-register. A flag bit of 0 leaves the contents of the N -register unaltered.

## SUBTRACT 02 SUB M Word Time = 1

The word in location $M$ is subtracted from the A-register, and the difference replaces the contents of the A-register. Overflow can occur. A flag bit of 1
will cause the contents of $M$ to be put into the $N$ register. A flag bit of 0 leaves the contents of the N register unaltered.

## MULTIPLY $\quad 20 \mathrm{MPY}$ M $\quad$ Word Time $=22$

The contents of the memory location $M$ are multiplied by the contents of the A-register. The product appears in the $A$ and L-registers unrounded. The most significant 31 bits and the sign are in the A-register. The least significant 31 bits are in the L-register. Because the number in the L-register is the least significant part of the total product, it does not have a sign associated with it. The most significant bit of the 31 bits of the least significant part of the product appears in the sign position of the L-register. The least significant bit position of the L-register always contains a 0 after multiplication, which is not considered part of the product. Note that the contents of location M are stored in the N -register. The flag bit is ignored during execution of this instruction.

$$
\underline{\text { DIVIDE }} \quad 30 \text { DIV } \mathbf{M} \quad \text { Word Time }=22
$$

The 63 bit number appearing in the $A$ and Lregisters is divided by (M). The most significant bit of the portion of the dividend that appears in the Lregister is in the sign position. The least significant bit of the L-register is not considered part of the dividend and must be a zero. The quotient in the Aregister will be affected if the value of this bit is 1 .

If the flag bit is 0 , the A-register will contain a rounded quotient and the L-register will cortain a remainder whose value has no relationship to the rounded quotient.

If the flag bit is 1 , the A-register will contain an unrounded quotient and the L-register will contain
a remainder with the same sign as the divisor. This remainder satisfies the relation:

$$
(\text { Divisor } \times \text { Quotient })+\text { Remainder }=\text { Dividend }
$$

The contents of location $M$ are put into the N register. Overflow can occur if the absolute value of the contents of the A-register is equal to or greater than the value of the contents of location M.

## CLEAR AND ADD 24 CLA M Word Time $=1$

The contents of location $M$ are transferred to the A-register. Location $M$ can refer to any addressable location. A flag bit of 1 transfers the previous contents of the A-register into the L-register. If the flag bit is 0 , the contents of the L-register remain unaltered.

## CLEAR AND SUBSTRACT 26 CLS M Word Time $=1$

The negative of the contents of memory location $M$ (in 2's complement form), is transferred to the Aregister. Location $M$ can refer to any addressable location. A flag bit of 1 transfers the previous contents of the A-register into the L-register. If the flag bit is 0 , the contents of the L-register remain unaltered.

### 2.2.2 STORE \& LOAD

Most store instructions include a capability to write the specified information into the N -register as into location $M$. If $M$ is in Hot memory, the indicated change occurs in $M$ and in the special register. However, if M is in Cold memory, only the N -register is changed. Each instruction description includes information as to whether a flag is required to cause storage into N . Note particularly the special precautions required when an STO or STP command is executed.

## STORE A-REGISTER 50 STA M Word Time $=1$

The contents of the A-register are stored in location $M$. Location $M$ can refer to any addressable location. If the flag bit is 1 , the word in A will be put in the N -register as well as into location M.

NOTE: If M refers to registers A, L or N:
With flag bit $=0$, no storage takes place into M or N .

With flag bit $=1$, no storage takes place into M , but the contents of A are sent to N .

## STORE L 52 STL M Word Time $=1$

The contents of the L-register are stored in location M. M may be any addressable location. If the flag bit is one, the word in L will be put into the N -register as well as into location M.

NOTE: If $M$ refers to registers $A, L$ or $N$ :
With flag bit $=0$, no storage takes place into M or N .

With flag bit $=1$, no storage takes place into M , but the contents of L are sent to N .
STORE N-REGISTER 40 STN M Word Time $=1$
The-contents of the N -register are stored in location M. M may be any addressable location. However, if M refers to the $\mathrm{A}, \mathrm{L}$ or N -register, no storage takes place, and the instruction serves only as a one-word time delay. The flag bit is ignored in this instruction.

STORE D $\quad 42$ STD M Word Time $=1$
The contents of sector $i$ of the 2 -word D-register are stored in location $M$, where $i=00$ if $M$ is even, and $\mathrm{i}=01$ if M is odd. M may be any addressable location. If the flag bit is one, the contents of D.i will be stored in the N -register as well as in M .

NOTE: If $M$ refers to registers $A, L$, or $N$ :
With flag bit $=0$, no storage takes place into M or N

With flag bit $=1$, no storage takes place into $M$, but the contents of the $D$-loop under read head at time of execution are stored in N .

STORE OPERAND ADDRESS 70 STO M Word Time $=1$ if M refers to R or Q .

Word Time $=3$ if $M$ refers to any other location.
(a) The operand address of the word in location $M$ is replaced by the operand address of the word in the A-register (bits 19 through 31). The modified contents of $M$ are stored in the $N$-register. M may be any addressable location.
(b) If M is located in Cold memory and no external input device, such as the SDR, is connected, the contents of $M$ are always unchanged. However, the $N$ register is modified to contain the instruction in M, with its operand address replaced by the operand address currently in the A-register.
(c) If $M$ is located in Memory II or III, and this memory is Hot, the sign bit of the L-register will determine whether ( $M$ ) is modified by an STO or an STP instruction. If $L$ is negative, no modification of ( $M$ ) will occur, although the modified instruction will be stored in $N$. If $L$ is positive, both (M) and ( N ) will be modified by execution of an STO or STP instruction. Therefore, to assure the modification of (M) in a Hot Memory II or Memory III location, the STO or STP could be preceded by a ZEL (Zero L) instruction. If M is in Memory

I, the sign bit in L hàs no effect on the execution of STO or STP commands.
(d) NOTE: If $M$ refers to the A-register, A remains unchanged and the contents of $A$ are stored in the N -register. If M refers to the L -register, the contents of L remain unchanged. However, the N -register is modified to contain the instruction in L, with its operand address (bits 19 thru 31) replaced by the operand address currently in A . If M refers to the N -register, bits 19 thru 31 in A are stored in bits 19 thru 31 of N. The remaining portion of the N -register is unchanged. The flag bit is ignored in this instruction.

STORE PROGRAM ADDRESS 60 STP M Word Time $=1$ if $M$ refers to R or Q .

Word Time $=3$ if $M$ refers to any other location.
(a) The next instruction address of the word in location $M$ is replaced by the next instruction address of the word in the A-register (bit positions 1 thru 13). The modified contents of M are stored in the N -register. $M$ may refer to any addressable location.
(b) If M is located in Cold memory and the SDR is disconnected, the contents of $M$ remain unchanged. However, the N -register is modified to contain the instruction in M, with its NI address replaced by the NI address currently in the A-register (bits 1 thru 13).
(c) See Part (c) under Store Operand Address. The contents of the L-register have the same affect on an STP instruction as on an STO.
(d) If M refers to the A-register, A remains unchanged and the contents of $A$ are stored in $N$. If $M$ refers to the

L-register, $L$ remains unchanged. However, the $N$ register is modified to contain the instruction in $L$, with its NI address replaced by the NI address currently in the A-register (bits 1 thru 13). If $M$ refers to the $N$-register bits 1 thru 13 are stored in bits 1 thru 13 of N . The remaining portion of N remains unchanged. The flag bit is ignored in this instruction.

$$
\text { STORE R } \quad 62 \text { STR M } \quad \text { Word Time }=22
$$

The contents of the 20 words of the R-loop are transferred through the D-register to memory starting at location $M$. The flag bit is ignored. If $M$ refers to a main memory location, each word transferred will be taken from that sector of the R-loop whose address agrees with the last four bits of the sector address into which the word is to be stored. For example, when an STR instruction with an operand sector address of 074 is executed, the first word stored will be taken from sector 14 of the R-loop. Successive words will be taken from sectors $15,16,17,0,1 \ldots 13$ of the R-loop and stored in sectors $075,076,077,100,101 \ldots 113$ of the main memory channel specified in M. The D-loop will contain the last two words transferred; that is, D. 0 will hold the contents of the last even numbered location transferred, and D. 1 the contents of the last odd numbered location transferred. In the example given above, D. 0 will hold the new contents of sector 112, and D. 1 will hold the new contents of sector 113 .

If $M$ refers to a main memory channel, and its sector is greater than 160 , the word transmitted from R. 0 will be stored in sector zero of the channel in which M is located.

Example: Given STR 20.166. Resultant storage will be:

$$
\begin{aligned}
& \mathrm{R} .16 \longrightarrow 20.176 \\
& \mathrm{R} .17 \longrightarrow 20.177 \\
& \mathrm{R} .00 \longrightarrow 20.0
\end{aligned}
$$

$$
\mathrm{R} .04 \longrightarrow 20.04
$$

$$
\mathrm{R} .05 \longrightarrow 20.05
$$

If $M$ refers to a location in the $Q$ or R-loop, the contents of $R$ will go into $Q$ or $R$, displaced by two sectors. For example, if a STR instruction is given which has an operand address of $\mathrm{Q}-3$, sectors 3, 4, 5, $6 . . .2$ of $R$ will be transferred to sectors $5,6,7,0 \ldots 4$ of Q. At the end of the execution, D. 0 will contain the new contents Q .4 , and D .1 the new contents of Q .3 . Similarly, if a STR instruction is given which has an operand address of R. 6 , sectors $6,7,10,11 \ldots 5$ of $R$ will be transferred to sectors $10,11,12,13 \ldots 7$ of R. At the end of the execution, D. 0 and D. 1 will hold the new contents of sectors R. 6 and R. 7 respectively.

If $M$ refers to the $D$-loop, the first word transferred appears in the sector of D specified and the last word transferred in the remaining sector of $D$. If the operand address is D.0, the first word with an even numbered address to pass under the read head of the

R-loop after the Instruction Read phase will be the first word transferred. For example, if channel 00, sector 55 contains the instruction STR D.0, the word in R. 16 will appear in D .0 and the word in R .15 will appear in $D .1$ at the end of the execution phase. If the instruction were STR D.1, the contents of R .17 and R. 16 would appear in D. 1 and D. 0 respectively.

## LOAD Q-REGISTER 72 LDQ M Word Time $=20$ <br> LOAD R-REGISTER 72 LDR* M Word Time $=20$

The contents of 20 consecutive words of memory starting with location $M$ are transferred to the specified rapid access loop. The specific loop, i.e., $R$ or $Q$, is determined by the flag bit. A flag bit of 1 indicates the R-loop, while a flag bit of 0 indicates the $Q$-loop.

Each word transferred to $Q$ or $R$ is stored in the sector corresponding to the last four bits of the sector from which the word was transferred. For example, when an LDQ instruction with an operand sector of 074 is executed, the first word transferred will go into sector 14 of the Q-loop. Successive words will go intosectors $15,16,17,0,1,2, \ldots 13$. If M refers to a main memory channel, and the sector is greater than 160 , the word transmitted from sector 177 will be followed by the word from sector 0 of the same channel.

The $A, L$, and $N$-registers and the D-loop are unaffected by the execution of these commands.

If $M=A, L$, or $N$, the contents of the specified register will replace the contents of all 20 sectors of the Q or R -loop. Similarly, if $\mathrm{M}=\mathrm{D} . \mathrm{i}$, the contents of D. 0 will go into all the even numbered sectors of the $Q$ or $R$-loop, while the contents of $D .1$ will replace the contents of all the odd numbered sectors. If the operand of a LDR* command is in the R-loop, or if the operand
of a LDQ command is in the Q-loop, the contents of the referenced loop remain unchanged.

### 2.2.3 TRANSFER

In each of the following instructions, a flag bit of 1 will cause the transfer instruction to be placed in the N -register if and only if the transfer to M is executed. A flag bit of 0 will leave the contents of the N -register unchanged.

## TRANSFER UNCONDITIONALLY

14 TRA M Word Time $=1$

Control is unconditionally transferred to the instruction found in location $M$. The NI is ignored in this instruction, except when the NI contains a 146 channel address. (See Section 3.10.)

TRANSFER ON PLUS 10 TPL M Word Time $=1$

If the contents of the A-register are positive( $\geq$ zero), control is transferred to the instruction in location $M$; if not, control is transferred to the address specified in the NI portion of the transfer command.

## TRANSFER ON ZERO 12 TZE M Word Time $=1$

If the contents of the A-register equal zero, control is transferred to location $M$; if not, control is transferred to the address specified in the NI portion of the transfer command.

## TRANSFER ON OVERFLOW 16 TOV M Word Time $=1$

If overflow occurs, control is transferred to the instruction in location M. If not, control is transferred to the address specified in the NI portion of the transfer command. In either case, the Error light indicator
on the control panel will not flash on execution of this command.

Note that if an instruction which results in overflow is not immediately followed by a TOV command, the Error light will flash, and all computation will halt.

### 2.2.4 SHIFT \& CYCLE

## General.

In the following instructions the operand sector address portion of the instruction, designated by S , does not refer to a location in memory but to the number of binary positions the word in the A-register, or in the A and L-registers, is to be shifted or cycled. All shift and cycle op codes are made up of four octal digits, with the 2 least significant digits in the positions normally occupied by the channel portion of the operand address.

## Overflow.

A change in the value of the sign bit in the Aregister is normally recognized as overflow. However, overflow caused by a right cycle command is always ignored by FADAC, and the next program instruction is exe cuted. Overflow can never occur with a right shift command. The flag bit of all right shift and right cycle instructions MUST be 0 . Overflow caused by a left shift or cycle command will be ignored if the instruction flag bit is 0 . However, a flag bit of 1 with a left shift or cycle command will cause the computer to check for overflow. If overflow occurs, the Error light indicator on the control panel will flash. The next instruction will be exe cuted only if it is a TRANSFER ON OVERFLOW (TOV) command, in which case the Error indicator will stop flashing. If any other instruction follows, the computer will halt, and the

Error indicator light will continue to flash.
CAUTION: The contents of the shifted register(s) are not predictable when overflow occurs since the overflow condition may terminate the command before it is completed.

Execution Time.
The execution time for all shift and cycle instructions is given by the following relationship:

Convert $S$ to decimal $S^{1}$. The number of word times equals the integer portion only of

$$
1 / 2\left(K+S^{\prime}+1\right)
$$

where $K=0$ for odd $S^{\prime}$ and where $K=1$ for zero and even $S^{\prime}$. This equation was derived from the following relationship:

$$
\text { Word Time }=1 / 2\left(1+S^{\prime}+\cos ^{2} \pi S^{\prime} / 2\right)
$$

All shift and cycle commands are summarized in Fig. 2.4.

Accumulator Right Cycle $\quad 7600$ ARC S
The contents of the A-register are shifted right the number of binary positions designated by S. The bits shifted off the right end of the A-register enter the left end of the A-register, through the sign position, in the same sequence.

## Accumulator Right Shift 7602 ARS S

On execution of this command, the contents of the A-register are shifted right the number of binary places designated by S . The sign is spread in the left-hand bits; while bits shifted beyond the right-hand limits of the A-register are lost.


Figure 2.4. Shift and Cycle Command Summary

The contents of the A-register are shifted left through the sign bit position the number of binary places indicated by S . The bits shifted off the left end of the A-register enter the right end, in the same sequence.

Accumulator Left Shift

## 7606 ALS S

The contents of the A-register are shifted left through the sign bit the number of binary places indicated by S. Positions left vacant as the contents of the A-register are shifted are filled with zeros. Any bits shifted beyond the sign position are lost.

## Long Right Cycle

7620 LRC S
The contents of the A and L-registers are shifted right the number of binary places designated by S . The bits shifted off the right end of the A-register are shifted into the L-register through its sign position. The bits shifted off the right end of the L-register are shifted into the left end of the A-register through its sign position.

## Long Right Shift

## 7622 LRS S

The contents of the A and L-registers are shifted right the number of binary places designated by S . The bits shifted off the right end of the A-register are shifted into the L-register through its sign bit position. The bits shifted off the right end of the L-register are lost. The sign of the A-register is spread.

## Long Left Cycle <br> 7624 LLC S

The contents of the A and L-registers are shifted left the number of binary places indicated by S. The bits shifted left through the sign of the A-register are shifted into the right end of the L-register. Bits shifted through the sign position of the L-register are shifted into the right end of the A-register.

The contents of the A and L-registers are shifted left the number of binary places designated by S. The bits shifted through the sign position of the L-register are shifted into the right end of the A-register. Bits shifted beyond the sign position of the A-register are lost. The bits of the L-register vacated in the shifting process are filled with zeros.

### 2.2.5 LOGIC AND SPECIAL INSTRUCTIONS

## General.

With the exception of the EXT, EQS and GES commands, all Logic and Special instructions are specified by a four digit op-code, with the two least significant digits in the positions normally occupied by the channel portion of the operand address. The sector portion of the operand address is ignored in these commands.

EXTRACT $\quad 34$ EXT M Word Time $=1$
The contents of the A-register are replaced by the logical AND of the contents of $A$ and the contents of $M$. In other words, when both the bit in (M) and the corresponding bit in the A-register are 1 , a one will appear in the same bit position of the A-register at the end of the execution. In all other cases, a zero will appear in that bit position of the A-register upon execution of this command. Location $M$ can refer to any addressable location. The flag bit is ignored in this instruction.

## TAKE ABSOLUTE VALUE 3770 ABS Word Time = 1

The contents of the A-register are replaced by their absolute value. If the contents are negative, they will be replaced by their corresponding positive value; while if the contents are positive they will remain unchanged. If the flag bit is zero, no further storage
occurs. If the flag bit is one and the contents of A are negative, this negative value will be transferred to the L-register. If the flag bit is one and the contents of $A$ are positive, $L$ will remain unchanged. The operand sector is ignored in this command.

## REPLACE A ON MINUS FROM L

## 3772 RML Word Time $=1$

If the contents of the A-register are negative and the flag bit of this instruction is zero, the complement of the contents of the L-register is sent to $A$, and $L$ is unchanged. If the contents of the A-register are negative and the instruction flag bit is one, the complement of the contents of $L$ is sent to $A$, and the original contents of $A$ are sent to $L$. If the contents of the A-register are positive, $A$ and $L$ are unchanged regardless of the value of the flag bit. The sign of the original contents of the L-register in no way affects the execution of this command.

## REPLACE A ON MINUS FROM N

## 3774 RMN Word Time $=1$

If the contents of the A-register are negative and the flag bit of this instruction is zero, the complement of the contents of the N -register is sent to $A$. The $L$ and N -registers are unchanged. If the contents of the A-register are negative and the instruction flag bit is one, the complement of the contents of the N -register is sent to $A$; the original contents of $A$ are sent to $L$; and N is unchanged. If the contents of the A-register are positive, $A, L$ and $N$ are unchanged. The sign of the original contents of the N -register in no way affects the execution of this command.

EQUAL SEARCH 64 EQS M Word Time $=\mathrm{i}+2$

$$
(i=0 \text { to } 176)
$$

Starting with $i=0$, the contents of the A-register
are compared with those in memory location ( $M+i$ ), in bit positions indicated by a mask appearing in the L-register. This mask must have been stored previously in the L-register by the program. It directs which bits of the A-register and memory should be compared. Wherever a 1 appears in the L-register, the corresponding bits in the A-register and location ( $M+i$ ) will be checked.

If in all the designated bit positions, the contents of the A-register are found to be equal to the contents of location ( $M+i$ ), the contents of location ( $M+i+1$ ) replace the contents of the A-register, and control is transferred to the next instruction.

If the contents of the A-register in all the designated bit positions are not equal to the contents of location ( $M+i$ ), then $i$ is increased by one and the comparison is repeated. This comparison search continues until agreement has been reached, or until the contents of operand sector 176 have been compared. If no agreement is reached after comparison search through sector 176, then the contents of A remain unchanged, and control is transferred to the next instruction.

If a complete channel is to be searched, the starting operand sector should be 177. Then if agreement were reached, the first and last possible word that would be picked up would be in sectors 000 and 177 respectively. If the flag bit is a 1 , and agreement is reached, the previous contents of the accumulator replace the contents of $L$. If agreement is not reached, the flag bit has no effect. The execution time is the number of sectors reached, i , plus 2.

GREATER THAN 66 GES M Word Time $=\mathbf{i}+2$
OR

$$
(i=0 \text { to } 176)
$$

EQUAL SEARCH
The execution of this instruction is the same as EQS, except that the comparison is satisfied if the value of the bit positions of (M) which are compared is greater than or equal to the value of the corresponding bit positions of (A).

## HALT $\quad 3720$ HLT Word Time $=1$

Program computation is halted. The flag bit, operand sector and next instruction address are ignored in this command.

ZERO-L-REGISTER 3762 ZEL Word Time $=1$
The L-register is cleared to zero. The flag bit and operand sector are ignored in this instruction.

INITIATE DISPLAY MODE 3764 IDM Word Time $=1$
This command will cause the contents of the $D$ loop to be displayed, in BCD format, in the NIXIE tubes located on the control panels and will also activate the Battery display NIXIE. The Battery NIXIE is not part of the D -register, but is a wired-in display of the Battery button (A, B, C, D, or E) depressed on the matrix panel. Note that a battery button must be depressed when an IDM command is sensed, or no display will occur. Once activated, the NIXIE display will continue until a HALT DISPLAY MODE (HDM), or a READ or WRITE command is sensed. Computation can continue while in the display mode. The flag bit and operand sector are ignored in this command.

See Section 4.1 for a full description of the NIXIE display.

## HALT DISPLAY MODE 3766 HDM Word Time $=1$

Visual display mode is halted. The flag bit and the operand sector address are ignored by this command. Successive IDM and HDM commands can be given at any time.

## INITIATE COMPUTE MODE 3726 ICM Word Time $=1$

This command will turn on the Compute light located on the front panel of FADAC. If the Compute light is already on, the command will introduce a one word-time delay. It has no other effect on a program.

HALT COMPUTE MODE 3724 HCM Word Time $=1$
This command will turn off the Compute light located on the front panel of FADAC, but does NOT halt computation. It has no other effect on a program. After an HCM command is sensed, the Compute light can only be relit if:
(1) an ICM command is sensed; or
(2) FADAC is turned off, then turned on, and one of the eight control buttons is depressed.

### 2.2.6 DISCRETE INPUT/OUTPUT

Input.
A single command, DIA, is available for a variety of input configurations to FADAC. These are described below:

DISCRETE INPUT TO
ACCUMULATOR 3640 DIA Word Time $=1$
The DIA command places the following information into the A -register:
a. In bit position 0: A timing signal, with a period of 22 milliseconds, which is generated by an oscillator in FADAC. The signal is a square wave which remains
at zero volts ( 0 to sign bit) for 11 milliseconds, and at -6 volts ( 1 to sign bit) for 11 millise conds. The signal is provided primarily as a synchronization signal for output by bit to a standard two-wire teletypewriter, but may be used as a real-time clock in any application.
b. In bit position 1 : Normally set to zero, this bit is set to one whenever any one of the 8 control buttons (Sample Matrix, Test, Set-up, Recall, Send, Compute, Trig, or Receive) is depressed; and remains at one only as long as pressure remains on the control button.
c. In bit positions 2-18: Up to 17 signals on lines connecting a single, or many, external devices to FADAC through Input Plug J - 17 and Output Plug J - 10. Note that the DIA command places information on specified pins of both plugs into the A-register simultaneously. (Figure 2.5 shows which pins on the $1 / O$ plugs are connected to each bit position.)
d. In bit positions 19-23: A 5 position switch is available to the programmer through the FADAC control panel buttons labelled A through E , which are located on the right hand side of the $8 \times 8$ matrix. With one of these buttons depressed, a DIA command will place its identifying code into bits 19 through 23 of the Aregister. (See figure 2.7.) (Note in the Cannon Program these buttons are referred to as the Battery Buttons. However, they are available for use in any application.)
e. In bit position 24: A 2 position switch is available to the programmer through the control panel buttons labelled (1) and (2) which are located immediately above the 5 position switch, buttons A to E. With Button (1) depressed, a DIA command will place a 1 in bit position 24 of the A-register; with Button (2) depressed, a 0 appears in bit position 24. (In the Cannon Program, these buttons are referred to as the Weapons Buttons.)
f. In bit positions 25-30: The $8 \times 8$ matrix on the FADAC control panel greatly extends the input

| BIT POSITION | F-LINE | PIN DESIGN ATION |
| :---: | :---: | :---: |
| 2 | F-30-I | m |
| 3 | F-29-I | k |
| 4 | F-28-I | J On Input Plug J - 17 |
| 5 | F-27-I | i |
| 6 | F-26-I | h |
| 7 | F-25-I | g |
| 8 | F-24-I | r |
| 9 | F-23-I | q |
| 10 | F-22-I | p On Output Plug J - 10 |
| 11 | F-21-I | n |
| 12 | F-20-I | nt |
| 13 | F-19-I | k |
| 14 | F-18 | MM |
| 15 | F-17 | LL |
| 16 | F-16 | KK On Input Plug J-17 |
| 17 | F-15 | J J |
| 18 | F-14 | HH |

Figure 2.5. DIA Command Pin and Line Connections Bits 2 Thru 18

| COL. $\boldsymbol{y}$ | IDENTIFYING CODES |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| H | 70 | 71 | 72 | 73 | 74 | 75 | 76 | 77 |
| G | 60 | 61 | 62 | 63 | 64 | 65 | 66 | 67 |
| F | 50 | 51 | 52 | 53 | 54 | 55 | 56 | 57 |
| E | 40 | 41 | 42 | 43 | 44 | 45 | 46 | 47 |
| D | 30 | 31 | 32 | 33 | 34 | 35 | 36 | 37 |
| C | 20 | 21 | 22 | 23 | 24 | 25 | 26 | 27 |
| B | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 |
| A | 00 | 01 | 02 | 03 | 04 | 05 | 06 | 07 |
| ROW $\rightarrow$ | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |

Figure 2.6. Identifying Codes for Each FADAC Matrix Position


Figure 2.7. DIA COMMAND - Binary Information
f. In bit positions 25-30 (Continued)
capabilities of FADAC by providing a 64 position switch to the programmer. Each position may be programmed as a pointer to a separate element of a given complex program. With a row (labelled A, B, ...., or H) and column (labelled $1,2, \ldots \ldots$, or 8 ) button depressed, a DIA command will place the identifying code of the combination in bits 25 - 30 of the A-register. Figure 2.6 shows the 2-digit octal identifying code for each combination.
g. Bit 31: If only one button is depressed on the $8 \times 8$ matrix (a row or a column, but not both) a 0 will appear in bit position 31; if both are depressed, a 1 will fill bit position 31.
h. Figure 2-7. summarizes the information sent to the A-register by a DIA command.

The following example illustrates a use of the DIA command:

## Example

The Problem.
Code an input routine for Matrix Position A - 1 through A -5; test if proper row and column button has been depressed; if so, transfer to one of 5 subroutines; if not, cause the No Solution Light to flicker, and halt. Assume that subroutines 1 through 5 begin in locations $10.00,12.00,14.00,16.00$ and 20.00 , respectively, and that a transfer table has been stored which has the following octal format:

| Octal Digit | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 2 | 0 | 0 | 1 | 4 | X | X | X | X | Y | Y |

Explanation of the Transfer Table Format:
Digit 1: Set to 2 to indicate that a match has been found, and a control word picked up from the transfer table (i.e. - the word in A is negative).

Digit 2 and 3: Always 0
Digit 4 and 5: Set to 14, the TRA op code.
Digit 6 through 9(XXXX): Entrance location of each subroutine 1 through 5.

Digit 10 and 11 (YY): Matrix Codes 00 through 04 (A-1 through A-5).

|  | TRANSFER TABLE |  |
| :--- | :---: | :---: |
| Location | $\frac{\text { Contents }}{200000000000}$ |  |
| 0.172 | 20014100001 |  |
| 0.173 | 20014120002 |  |
| 0.174 | 20014140003 |  |
| 0.175 | 20014160004 |  |
| 0.176 | 20014200000 |  |

Assume that when each selected subroutine has been executed, control is to return to the main program which starts at $\mathrm{L}_{0} 2.100(00300)$.

INSTRUCTIONS

| $\mathrm{L}_{0}$ | $\mathrm{~N}_{1}$ | Op Code | Operand | Explanation |
| :---: | :---: | :---: | :---: | :--- |
| 0.0 | 00201 | 36 | 4000 | When sample matrix <br> button is depressed, <br> the program begins <br> by executing DIA <br> command (3640). |
| 2.1 | 00202 | 76 | 0001 | ARC 1: To place bit <br> 31 in sign position. |
| 2.2 | 00203 | 10 | 0403 | TPL: if (+)goto Er- <br> ror Routine; if( -$)$ |
| 2.3 | 00204 | 34 | 0604 | continue. <br> EXT: Matrix code <br> now in bit positions |
|  |  |  |  | 26 to 31 in A- <br> register. L 604 |
| contains the Mask. |  |  |  |  |


| $\mathrm{L}_{0}$ | $\mathrm{N}_{1}$ | Op Code | Operand | Explanation |
| :---: | :---: | :---: | :---: | :---: |
| 2.4 | 20205 | 24 | 0604 | CLA*: Matrix code to L-register, Mask to A-register. |
| 2.5 | 20206 | 25 | 7200 | CLA*: Matrix code to A-register, Mask to L-register. |
| 2.6 | 00207 | 64 | 0172 | EQS: Search table stored in 0.172 through 0.177. |
| 2.7 | 00210 | 10 | 0403 | TPL: If + , wrong combo set in Matrix, go to error routine; if - , A-register now contains transfer instruction. |
| 2.10 | 00213 | 76 | 0206 | ARS 6: Position transfer in struction properly in A. |
| 2.13 | 20300 | 15 | 7000 | TRA* $A$ and return to Main Program. |
| 4.3 | 00404 | 37 | 1600 | NSL-Error Routine Wrong matrix position is set. |
| 4.4 | 00000 | 37 | 2000 | HLT - Halt, correct and start again. |
| 6.4 | 00000 | 00 | 0077 | Mask for EXT and EQS command |

The DIA command may also be used to input any 2 to 16 bit code which is foreign to FADAC. This is accomplished by connecting up to 17 of the J10 and

J17 lines associated with the DIA command to any input device which generates the desired code. A timing signal generated by the external device must be connected to one of the 17 bits positions available. Duration of the timing signal must be long enough to allow input of each new character, and short enough to assure that each new character is sensed only once per input. With the period of the timing signal known, a program can be written to sample for this signal by a DIA, EXT, and TEST. When the timing signal is thus detected, the new code can be input by a second DIA command followed by an EXT (extract the number of bits which define a character), STA, and a delay designed to inhibit reinitiation of a loop back to test for the timing signal until the next character is available for input. An example of a code foreign to FADAC is USASCII, an 8 - bit code which can be input only in the manner described.

## Discrete Output:

The eight discrete output commands set 3 flipflops, OP1, OP2, and OP3, to any one of eight possible combinations. Each combination, in turn, either places a - 6 volt signal on a given output line; or causes the No Solution Light (NSL) on the FADAC control panel to blink; or turns off the NSL or the signal on any output line (DOF). A description of each discrete output command is summarized in Figure 2.8.

Note that each Discrete Output command requires a four-digit operation code. The operand sector address is ignored in the execution of these commands. Also, note that the setting of the Discrete Output control flip-flops remains fixed until a new Discrete Output command is sensed. When FADAC is turned on, the output flip-flops are set to all zeros (effectively, a DOF command is executed).

| Command <br> Mnemonic | Octal Op-Code | OP. Flip-Flop Setting (in Binary) |  |  | OPL Line on which a -6 volt signal appears when the command is executed: | Pin to which given OPL line is connected on specified I/O Plug: | Function of Command |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | Op 3 | Op 2 | Opl |  |  |  |
| DOF | 3700 | 0 | 0 | 0 | None (All | None | Turns off the NSL or Sets all OPL lines to Zero |
| ODI | 3702 | 1 | 1 | 1 | OPL - 1 | X of $\mathrm{J}-10$ | Sets OPL-1 to one |
| OD2 | 3704 | 0 | 1 | 0 | OPL - 2 | Y of J-10 | Sets OPL-2 to one |
| OD3 | 3706 | 0 | 1 | 1 | OPL-3 | Z of J-10 | Sets OPL. 3 to one |
| IDI | 3710 | 1 | 0 | 0 | OPL - 4 | a of J-17 | Sets OPL-4 to one |
| ID2 | 3712 | 1 | 0 | 1 | OPL - 5 | $\overline{\mathrm{b}}$ of J-17 | Sets OPL-5 to one |
| ID3 | 3714 | 1 | 1 | 0 | OPL-6 | तै of J-17 | Sets OPL-6 to one |
| NSL | 3716 | 1 | 1 | 1 | None | None | Causes No Solution Light to blink |

Figure 2.8. Discrete Output Command Summary
The function of the OD and ID commands is to to energize a given pin on the input or output plug, or de-energize all OPL line-connected pins on these plugs. They may be used to control external equipment by the presence or absence of -6 volts on a given line.

### 2.2.7 PARALLEL INPUT/OUTPUT

## Introduction.

The parallel I/O commands cause FADAC to read information input through the keyboard, mechanical reader, magnetic tape, or any external device; or to write information out to another FADAC or to any of
several external devices. The command structure allows the programmer to designate the I/O device; the kind of block (word or character); the block size; the I/O code (octal, BCD, or alpha-numeric); and, in all numeric WRITE commands, the code level (either 5 -level Teletype (TT), or 8 -level Fieldata (FD)). During input, the level is designated by the setting of the SDR (or SDR simulator) TT/FD switch. Parallel I/O commands are so named since each character enters FADAC on 5 (for TT) or 8 (for FD) information lines, (designated " I "' lines) in parallel; and each character is output from FADAC on 5 or 8 data lines (designated "D" lines) in parallel. Input through the Mechanical Reader on the FADAC control panel is specified by the READ TAPE commands. Input through the MLU is specified by the READ EXTERNAL DEVICE commands. The flow of information through the special registers during $\mathrm{I} / \mathrm{O}$ is a function of the particular $\mathrm{I} / \mathrm{O}$ command exe cuted, and is described in the following sections. In general, it is recommended that the contents of all special registers be saved prior to execution of any I/O command if they are of any significance in the balance of a given program. Note that the number of word times required to execute a parallel I/O command cannot be fully predicted.
The I/O Command Format:
Figure 2.9 illustrates the $1 / O$ command format. The flag bit is set to 0 to designate I/O by character, or to 1 to designate $1 / O$ by word. The value of the flag is added to the most significant digit of the NI address to form the first octal digit of the command word. The operation code occupies the 6th, 7th, 8th and 9 th octal digits of the word. The block size is in digits $8,9,10$ and 11. Thus, digits 8 and 9 are the sum of the two least significant digits of the operation code and the two most significant digits of the block size. The block size is designated as one less than the actual

OP-Code and


## Figure 2.9. I/O Command Format

number of words or characters to be input or output. When the block is designated as a word, the maximum block size which may be specified per command is 1023 (1777). When $I / O$ is by character, the maximum number of characters which can be put or output per command is a function of the number of characters which can be packed into a single FADAC word in the specified code.

### 2.2.8 READ COMMANDS

Octal Input:
Execution of a READ Octal command differs from all other parallel I/O in that the block type, block size and the Next Instruction portions of the command are ignored by FADAC; and control reverts to the input device. As each character is received, FADAC hardware inverts the tape code sensed, and searches the sector track (Section 1.4) for a match. The four least significant bits of each sector number represent the binary translation of the tape code for that number. When a match is found, the leading bit of the binary translation is tested. If it is a one, the character is an octal control character or a blank, and FADAC logic is set to perform the indicated operation (Halt, Compute, Fill, Verify, Enter, Location, Clear, or Blank). Both leading and non-leading blanks are ignored. If the leading bit is an $O$, it is stripped and the octal translation of the number is shifted into the right side of the Aregister. If the tape code sensed is not found in the
sector track, the character is ignored. If input is in Fieldata code, parity is checked. The Parity Light on the FADAC control panel will blink when a parity error is detected, halting input. Otherwise, input continues until a Compute or Halt code is received. The Compute code will cause control to be returned to FADAC, and the command contained in the address defined by the 14 least significant bits of the L-register will be executed. (See Appendix C for a list of Octal codes, and their TT and FD tape code equivalents.) The Halt code will stop all FADAC operations. If no Enter code is received before input is halted, the last 11 (or less) octal digits input will remain in the A-register. If more than 11 digits are received with no Enter code, the most significant digit is shifted out of the left side of $A$ and lost, as each new digit enters the right side of the register. The contents of $A, L, N$, and all sectors of the Q-loop are destroyed by the input of no more than one full Octal word, while the contents of the D-loop and R -loop are unchanged. If more than one Octal word is to be input, the 5 -digit octal address of the first location into which the data is to be stored must be entered right adjusted, followed by a location code. This causes the last 14 bits of $A$ to replace the last 14 bits of the L-register. Input of an Enter code will send the operand address in L to the Q -loop and the contents of A to the N-register, a sector of the R-loop, both sectors of the D-loop and thence to the memory location last sent to the Q -loop from the L-register. $L$ is incremented by 1 each time a word is stored. If several words are to be stored in contiguous memory locations, only the initial address must be forced into L, using the location code. To change the contents of L for non-contiguous storage, it is necessary to input the new 5 -digit address, followed by a Location Code.

During octal input, the 5-digit location address may designate any part of FADAC memory. To address

Memory I, the first digit must be 0 ; to address Memory II, the first digit is 1 ; to address Memory III, the first digit is 3 . When a Memory III (location address) is set in L, and an Enter code sensed, the memory control flip-flop (discussed in detail in Section 3.10) is set such that the information in A is stored in the designated Memory III location. Storage into Memory III will continue until a Memory II location address is sent to the Lregister; at which time the memory control flip-flop will be reset to store inputs into Memory II. Designation of a Memory I location address will not change the current setting of the flip-flop. Storage into the designated Memory I location occurs regardless of this setting.

### 2.2.9 BCD AND ALPHA-NUMERIC INPUT

## General.

BCD and alpha-numeric input is entirely under program and FADAC hardware control. When such input commands are executed, the flag bit is checked to determine block type (word or character), and the block size is saved in a temporary counter. The counter is automatically decremented by one each time a block is processed. Control is sent to the Next Instruction address when the counter equals zero, indicating input is completed.

## Input by Character:

With the flag bit set to zero, each character shifted into A causes the block size counter to decrease by one. The characters remain in the A-register only, if the equivalent of one full word or less is input. If more than a full word (8 BCD characters, 5 Alpha-5, or 6 Alpha- 6 characters) is received, the contents of A are loaded into the N -register, and subsequent characters are shifted into A. A maximum of two full FADAC words can be processed by character. No storage to main
memory occurs automatically during character input. If the data received is to be stored, a Store A (STA) and Store N (STN) command must follow the input instruction. The original contents of $A, L$, and $N$ are destroyed during character input, while the $D, R$, and $Q$-loops are unchanged.

Input by Word:
Prior to word input, the L-register must be programmed to contain the address of the first location into which the first word received is to be stored. Each time the A-register is filled with the number of characters which define a full word in the input code specified the address in $L$ is sent to the $Q$-loop, and the contents of A are loaded into $\mathrm{N} . \mathrm{N}$ is loaded into R ; then written into the D-loop, and finally into the memory location specified in $Q$. $L$ is incremented by one each time a word is written into main memory. The input cycle thru the registers is repeated until the indicated block size, plus one, has been input and stored.

## BCD Input:

During the input of BCD codes, FADAC hardware inverts the tape code sensed, and searches the sector track for a match. If a match is found, the four least significant bits of the corresponding sector track number are the BCD translation of the tape code. These bits are shifted into the right hand end of the A-register. If no match is found, the character is ignored. Leading blanks are ignored, while imbedded blanks are processed. (See Appendix $C$ for a list of $B C D$ codes and their TT and FD tape code equivalents.)

## BCD Input By Character:

Example: Since 8 BCD characters constitute one full FADAC word, a maximum of 16 BCD characters may be meaningfully input per command. To illustrate
processing, consider the execution of a command requiring the input of 12 BCD characters. When input is finished, the contents of A and N will be as follows:

N -Register

| Char | Char | Char | Char | Char | Char | Char | Char |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |

## A-Register

Char Char Char Char Char Char Char Char

| 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

A contains characters 5 through 8 of the first word input plus characters 9 through 12---the last 4 characters entered. $N$ contains characters 1 through 8 which constitutes the first full word received. Characters 1 through 4, which were in $A$ when $N$ was loaded, were shifted out of the left side of $A$ when characters 9 through 12 entered the right side of $A$.

## READ BCD Commands:

The following op-codes cause FADAC to set up the necessary hardware logic to accept $B C D$ input from the indicated device:

| Op-Code <br> Mnemonic | Op-Code in <br> FADAC Octal | Interpretation |
| :---: | :---: | :--- |
| RED | 5500 | Rd External Device in BCD |
| RTD | 5520 | Rd Tape (Mechanical <br> Reader) Device in BCD |
| RKD | 5560 | Rd Keyboard in BCD |

Two alpha-numeric codes, designated Alpha-5 and Alpha-6, are processed by FADAC input logic. Alpha-5 code is defined as five 6 -bit characters per word, or 30 information bits. The two additional bits which constitute a FADAC word have no significance. Alpha-6 code is defined as six 6 -bit characters per word, with the 4 leading bits of the first character of each Alpha-6 shifted out of $A$ and lost when the sixth character enters A. All alpha-numeric characters are initially received as an 8 -bit Fieldata tape code. Parity is checked. If a parity error is detected, the Parity light on the FADAC control panel blinks until the RESET button is pressed. If Parity is good, the two leading bits of each tape code are stripped, and the 6 information bits are shifted into the right side of the A-register. No table look-up is required for these codes, since the 6 least significant bits of each FD tape code represents the binary configuration designated for a given character. (See Appendix D.) Therefore, I/O in any alpha-numeric code is much faster than is numeric code I/O. During character input, a maximum of 9 useful characters may be entered in Alpha-5 code, or 11 in Alpha-6 code.

READ Alpha-Numeric Commands:
The following op-codes cause FADAC to set up the necessary hardware logic to accept alpha-numeric input from the indicated devices:

Op Code Op Code in Mnemonic FADAC Octal

| RE5 | 5700 | Rd External Device in <br> Alpha-5 |
| :--- | :--- | :--- |
| RT5 | 5720 | Rd Tape Device in <br> Alpha-5 |

Op Code Op Code in Mnemonic FADAC Octal

| RM5 | 5740 | Rd Magnetic Tape in <br> Alpha-5 |
| :--- | :--- | :--- |
| RK5 | 5760 | Rd Keyboard in Alpha-5 <br> RE6 |
| RT6 | 5600 | Rd External Device in <br> Alpha-6 |
| RM6 | 5640 | Rd Tape Device in <br> Alpha-6 <br> Rd Magnetic Tape in <br> Alpha-6 |
| RK6 | 5660 | Rd Keyboard in <br> Alpha-6 |

### 2.2.10 WRITE COMMANDS

## General.

All parallel output is entirely under program and FADAC hardware control. When such commands are executed, the flag bit is checked to determine block type, and the block size is saved in a temporary counter. The counter is automatically decremented by one each time a block is processed. Control is sent to the Next Instruction address when the block size, plus one, has been output.

## Output by Character:

Prior to execution of a WRITE by character command, the characters to be output must be loaded, left justified, into the A-register. If more than one full word is to be output, the first X characters which define a full word in the specified output code must be sent to $A$, and the remaining characters must be left justified in N . The equivalent of two full FADAC words may be
output by character ( 16 BCD characters, 10 . Alpha-5, or 12 Alpha-6), except when the output code is Alpha-4. When one full word has been output, the contents of N are automatically loaded into A. Each character is shifted out of the left side of $A$, one character at a time until the designated block size, plus one, has been generated. The original contents of $A, L$ and $N$ are destroyed. The contents of all sectors of the $R, Q$ and D-loops are unchanged.

## Output by Word:

Prior to execution of a WRITE by word command, bits 18 to 31 of the L-register must be programmed to contain the address of the first word to be generated. During execution of a WRITE command, the contents of the address in L are written into the R-loop, the N register, and the A-register, from which they are sent to the output lines, one character at a time. The address in L is sent to the Q -loop. L is incremented by one, and the cycle through the registers to the output lines is repeated until the number of words specified in the block size portion of the WRITE command, plus one, has been output. The contents of the D-loop are unchanged during output. When output is completed, the N -register contains the last word output; bits 18 to 31 of the L-register contain the last address from which data was output, plus 22. The Contents of A and all sectors of the R and Q -loop cannot be fully predicted.

## Numeric Output:

During numeric output, in Octal or BCD code, each character shifted out of the left side of $A$ is matched against the four least significant bits of the sector numbers inscribed on the sector track. Note that during look-up of octal characters (3 bits), the missing leading bit is interpreted as zero, since control characters
are not stored in octal. Also, since the most significant octal digit of each word is composed of only 2 bits, both missing leading bits of that character are interpreted as zero. When a match is found, the corresponding tape code designated in the WRITE command (TT or FD) is inverted, sent to the $D$ lines, and output. If output is in Octal code, FADAC hardware automatically generates an Enter code each time a full word is output, whether the block type be word or character. The Enter code must be counted as part of the designated block size when octal output is by character.

WRITE Numeric Commands:
The following op-codes cause FADAC hardware to set up the necessary logic to output data in Octal or BCD code to the indicated devices in either 5 level or 8 level tape code.

Op-Code Op-Code in Mnemonic FADAC Octal

| WEOT | 4400 | Write FADAC to External <br> Device in Octal - TT |
| :--- | :---: | :--- |
| WFOT | 4420 | Write FADAC to FADAC in <br> Octal - TT |
| WEOF | 4440 | Write FADAC to External <br> Device in Octal - FD |
| WFOF | 4460 | Write FADAC to FADAC in <br> Octal - FD |
| WEDT | 4500 | Write FADAC to External <br> Device in Dec. - TT <br> WFDT |
|  | 4520 | Write FADAC to FADAC <br> in Dec. - TT |

Op-Code Op-Code in Mnemonic FADAC Octal

| WEDF | 4540 | Write FADAC to External <br> Device in Dec - FD |
| :--- | :--- | :--- |
| WFDF | 4560 | Write FADAC to FADAC in <br> Dec. - FD |

## Alpha-Numeric Output:

Three alpha-numeric codes are output by FADAC, each of which are discussed below. Since no table lookup is required for alpha-numeric output, these codes are generated at a much faster rate than is possible during numeric output.

Alpha-4:
Alpha-4 is defined for output only. It is composed of 8 bits per character, 4 characters per FADAC word. Any 8 -bit configuration pre-stored in FADAC may be generated by a Write Alpha-4 command. It is used to output a code foreign to FADAC, such as ASCII, or to control external devices by placing a control word on the D-lines. Output in Alpha-4 must be by character only, with a maximum block size of 3 specified (designating 4 characters actually to be generated). Prior to execution of a Write Alpha- 4 command, the A-register must be loaded with the 4 (or less) characters to be generated. As each 8 bit character is shifted out of A, its most significant bit is inverted by FADAC hardware. The inverted first bit and the remaining 7 bits are then placed on the D-lines.

Alpha-5:
Alpha- 5 is defined for both input and output as 5 six-bit characters per FADAC word. Unlike input, Alpha- 5 output must be by word only. Each Alpha-5 word is automatically processed by shifting bits 0 and

1 out of the A-register, after which each six bit character is placed on the data lines, with a parity bit and a timing bit added to the most significant bit positions by FADAC hardware to form the 8-bit Fieldata code for that character.

Alpha-6:
During processing of a Write Alpha-6 command, the 2 leading bits of each Alpha-6 word are recognized as the 2 least significant bits of the first character. On output, these two bits are placed on D-lines 1 and 2 , the next 4 bits are generated by FADAC, and a parity and timing bit added, to form an 8-bit Fieldata code on lines $D-1$ thru $D-8$. The value of the machine generated bits cannot be predicted or controlled. Therefore, the first character of each word output in Alpha-6 code may be meaningless. This characteristic limits the use of Alpha-6 code to communication between two FADAC's or between FADAC and a magnetic tape unit. In both applications, when each Alpha-6 word is read back into FADAC or the magnetic tape unit, the 6 leading bits of the first character are stripped, and only the two original bits are stored. The flow of information through the special registers is as described in the sections on "output by Character" and "Output by Word". Alpha- 6 code may be output by word or character.

## Write Alpha-Numeric Commands:

The following op-codes cause FADAC to set up the necessary hardware logic to allow output in the indicated Alpha-numeric codes:
Op-Code Op-Code in Mnemonic FADAC Octal

| WE4 | 4700 | Write FADAC to External <br> Device in Alpha-4 by <br> Character |
| :---: | :---: | :--- |
| WF4 | 4720 | Write FADAC to FADAC in <br> Alpha-4 by character |
| WE5* | 4700 | Write FADAC to External <br> Device in Alpha-5, by word |
| WF5* | 4720 | Write FADAC to FADAC in <br> Alpha-5, by word |
| WE6 | 4600 | Write FADAC to External <br> Device in Alpha-6 |
| WF6 | 4620 | Write FADAC to FADAC in <br> Alpha-6 |
| WM6 | 4640 | Write FADAC to Mag. Tape <br> Alpha-6 |

* Note that the op-codes for the Write in Alpha-4 commands are identical to those for the Write in Alpha-5 commands. For the former, the flag bit must be 0 , since output can be by character only; while for the latter, the flag bit must be 1 to designate output by word only.


## CHAPTER 3

## CODING FOR FADAC

### 3.1 MINIMUM ACCESS CODING

Multiplexing.
In the processing of all FADAC instructions, the following five phases are automatically completed:
(1) Instruction Search (IS)
(2) Instruction Read (IR)
(3) Operand Search (OS)
(4) Operand Read (OR)
(5) Execution (EX)

As illustrated in Figure 3.1, the phases are multiplexed; that is, performed simultaneously by FADAC hardware. Consider the multiplexing of five instructions, n - 1 through $n+3$, each of which is assumed to be optimally coded, and each of which requires one word time to execute. Assume that $n$ is the current instruction, $\mathrm{n}-1$ the immediately preceding instruction, and $n+1$ through $\mathrm{n}+3$ the subsequent instructions. Let $\mathrm{t}=$ the word time of the execution of the preceding instruction, n - 1. In Figure 3.1, the multiplexing of the five instructions is shown from time $t-2$ through time $t+2$. At $t-2$ (two word times before $t$ ) search is made for instruction $n-1$. At time $t-1$, instruction $n-1$ is read, its operand is searched, and simultaneously, search is made for the current instruction, $n$. At time $t$, search is made for the next instruction, $n+1$, instruction $n$ is read and its operand searched; and at the same time the operand of the preceding instructions, $n-1$, is read and instruction $n-1$ is executed. At time $t+1$,

| Word Time |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| t-2 | t-1 | t | $\cdot \mathrm{t}+1$ | $t+2$ |
| $\mathrm{IS}_{n-1}$ | $\begin{aligned} & \mathrm{IS}_{n} \\ & \mathrm{IR}_{\mathrm{n}}-1 \\ & \mathrm{OS}_{\mathrm{n}-1} \end{aligned}$ | $\begin{aligned} & \mathrm{IS}_{n+1} \\ & \mathrm{IR}_{\mathrm{n}} \\ & \mathrm{OS}_{\mathrm{n}} \\ & O R_{\mathrm{n}-1} \\ & E X_{n-1} \end{aligned}$ | $\begin{aligned} & \mathrm{IS}_{\mathrm{n}+2} \\ & \mathrm{IR}_{\mathrm{n}+1} \\ & \mathrm{OS} \\ & \mathrm{n}+1 \\ & \mathrm{OR} \\ & \mathrm{n} \end{aligned}$ | $\begin{aligned} & \mathrm{IS}_{\mathrm{n}+3} \\ & \mathrm{IR}_{\mathrm{n}}+2 \\ & \mathrm{OS} \\ & \mathrm{n}+2 \\ & \mathrm{OR}_{\mathrm{n}+1} \\ & E X_{\mathrm{n}+1} \end{aligned}$ |

Figure 3.1. Multiplexing Diagram
the operand of the current instruction, $n$, is read, $n$ is executed, and within the same word time, instruction search is made for $n+2, n+1$ is read and its operand searched, etc.

In all of the above it is assumed that each instruction requires only one word time to execute. If time to execute a command is greater than one word time, then execution is initiated at the same time the current operand is read, and execution of the next command is delayed until execution of the current command is completed.

This multiplexing feature of FADAC hardware will decrease program execution time if the programmer places instructions and operands in memory such that:

1. Each instruction comes under a read head no later than the last word time of the execution phase of the previous instruction.
2. Each operand comes under a read head during the first word time following the completion of the previous instruction.

If these conditions are met, the program is optim-ally-coded, and there is no wait to read instructions or operands. The time lapse in fully executing an instruction is its execution time; the total elapsed time to complete a program is equal to the sum of the execution times of all instructions within the program.

Recall that the execution of most FADAC commands requires one word time, with the exception of the following:

Command

Number of Word Times Required for Execution

| MPY | 18 |
| :--- | :---: |
| DIV | 18 |
| LDR* | 16 |
| LDQ | 16 |
| STR | 18 |
| All Shifts \& Cycles | Variable |
| All Parallel I/O | Variable |
| STO | $\left\{\begin{array}{l}3 \text { if M is in Main Memory } \\ \text { STP } \text { if M is in } R \text { or } Q\end{array}\right.$ |

### 3.2 OPTIMIZING ASSIGNMENTS MAIN MEMORY

In the following discussions, let $E=$ number of word times required to execute a given command, and $\mathrm{Se}=$ the current sector at which the read heads are located. In all examples given, the operation code is represented as a 3 -character mnemonic. An asterisk next to the operation code indicates that the flag bit is one. All FADAC locations are shown as one to three octal digits, followed by a period, followed by one to
three octal digits to designate "channel-sector". (Leading zeros are not shown in any address.) Registers are coded as A, L, N, D.0, D.1, R. 0 through R.17, and Q. 0 through Q.17. Throughout the discussion all numbers are in octal, unless otherwise indicated. Recall that at any time, $t$, the read heads of all $64{ }_{10}$ channels of main memory are simultaneously positioned at the same sector, Se , while the write heads are positioned two sectors before the read heads ( $\mathrm{Se}-2$ ). In general, to optimize an operand in main memory, assign the operand to a sector to $(\mathrm{Se}+1)$. To optimize an NI in main memory, assign the NI to any sector from ( $\mathrm{Se}+1$ ) through $(\mathrm{Se}+\mathrm{E})$. If a given NI is less that $(\mathrm{Se}+\mathrm{E})$, then the next optimum operand (OP) must be $\geq(\mathrm{Se}+\mathrm{E})$ +1 .

Example.
Time
Location Op-Code Operand Ni Remarks
$1 \quad S_{1}=100.2 \quad$ CLA $\quad 100.3 \quad 102.3 \quad E$ for $C L A=1$

$$
\begin{aligned}
\therefore & O p_{p}=N I= \\
& \left(S_{e_{1}}+1\right)
\end{aligned}
$$

$2 \quad \mathrm{Se}_{2}=102.3 \quad \mathrm{ADD} \quad 100.4 \quad 102.4 \quad \mathrm{E}$ for $\mathrm{ADD}=1$

$$
\begin{aligned}
\therefore & O p=N I= \\
& \left(S e_{2}+1\right)
\end{aligned}
$$

$3 \quad \mathrm{Se}_{3}=102.4 \quad$ MPY $\quad 100.5 \quad 100.10 \quad E$ for MPY $=22$

4

$$
\mathrm{Se}_{4}=100.10 \quad \text { ADD } \quad 100.27 \quad 120.27 \quad \mathrm{Op}=\left[\begin{array}{r}
\left(\mathrm{Se}_{3}+\right. \\
\left.\left.\mathrm{E}_{3}\right)+1\right]
\end{array}\right.
$$

STORE and LOAD Commands ( $M$ in Main Memory). STA, STL, STN, STD.

These commands require that the contents of the named register be written into the operand, M. Since the write heads in main memory are at $(\mathrm{Se}-2)$ at time t , when the instruction is read, they will be ( $\mathrm{Se}-1$ ) at time $t+1$ when execution is initiated. Therefore, $M$ is optimum at $\mathrm{Se}-1$. The optimum NI is sector $\mathrm{Se}+1$. Note that given an STD command, the contents of D. 0 are stored in $M$ when $M$ is even; the contents of $D .1$ are stored in $M$ when $M$ is odd.

## STO and STP.

When M is in main memory, these commands require 3 word times to complete, with the actual storage (writing) into $M$ occurring during the third word time. The write heads, which are located at $\mathrm{Se}-2$ at the beginning of execution of these commands, have turned through to $\mathrm{Se}+1$ at the time storage occurs. Therefore, M is optimum at $\mathrm{Se}+1$, and the NI is optimum at $\mathrm{Se}+3$ (or $\mathrm{Se}+1 \leq \mathrm{NI} \leq \mathrm{Se}+3$ ) if the next operand is $\geq$ $[(\mathrm{Se}+3)+1]$.

STR.
This command is initiated by FADAC hardware by writing the contents of sector $\mathrm{Se}+\mathrm{i}$ and $\mathrm{Se}+\mathrm{i}+\mathrm{l}$ of the R -loop into the D -loop, and then transferring the information from a sector of D into M . Because of the initial 2 word time delay during which the D-loop is filled, M is optimum at $\mathrm{Se}+1$. Again, the NI may be any sector between $\mathrm{Se}+1$ and $\mathrm{Se}+\mathrm{E}$, if the next operand is $\geq[(\mathrm{Se}+\mathrm{E})+1]$; where $\mathrm{E}=22$ word times. Locations $M$ through $M+17$ are filled with the contents of the R-loop during the execution of this command.

These commands transfer the contents of $M$ through $\mathrm{M}+17$ sequentially into the corresponding sectors of the R or Q -loop. Therefore, M is optimum at $\mathrm{Se}+1$.

## Example.

| Line No. | Location | Op-Code | Operand | NI | Notes |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | $\mathrm{Se}_{1}=014.75$ | STA | $\begin{aligned} & 76.74 \\ & \left(S e_{1}-1\right) \end{aligned}$ | $\begin{aligned} & 60.76 \\ & \left(\mathrm{Se} \mathrm{e}_{1}+1\right) \end{aligned}$ | $(\mathrm{A})-76.74$ |
| 2 | $\mathrm{Se}_{2}=060.76$ | STP | $\begin{aligned} & 70.77 \\ & \left(\mathrm{Se}_{2}+1\right) \end{aligned}$ | $\begin{gathered} 62.101 \\ \left(\mathrm{Se}_{2}+3\right) \end{gathered}$ | NI address in $\mathrm{A} \rightarrow 70.77$; <br> new contents of $70.77 \rightarrow \mathrm{~N}$. <br> E for $\mathrm{STP}=3$ word times $\mathrm{NI}=\mathrm{Se}_{2}+3$ |
| 3 | $\mathrm{Se}_{3}=62.101$ | CLA | 64.102 | 70.102 | $\mathrm{OP}=\mathrm{NI}=\mathrm{Se}_{3}+1$ |
| 4 | $\mathrm{Se}_{4}=70.102$ | STO | 130.103 | 70.105 | OP address in $\mathrm{A} \rightarrow 130.103^{*}$ <br> Modified $\mathrm{M} \rightarrow \mathrm{N}, \mathrm{E}$ for $\mathrm{STO}=$ 3 word times, $\mathrm{NI}=\mathrm{Se}_{4}+3$ |
| 5 | $\mathrm{Se}_{5}=70.105$ | STR | 70.106 | 70.120 | $\begin{aligned} & \mathrm{Se}_{5}+\mathrm{l} \leq \mathrm{NI} \leq \mathrm{Se}_{5}+\mathrm{E} \\ & \mathrm{E} \text { for } \mathrm{STR}=22 \end{aligned}$ |
| 6 | $\mathrm{Se}_{6}=70.120$ | ADD | 76.130 | 24.130 | $\mathrm{M}=\left(\mathrm{Se}_{5}+\mathrm{E}\right)+1$ |
| 7 | $\mathrm{Se}_{7}=24.130$ | LDR* | 100.131 | 24.150 | $\mathrm{M}=\mathrm{Se}_{7}+1 ; \mathrm{NI}=\mathrm{Se}_{7}+\mathrm{E}$ |
| 8 | $\mathrm{Se}_{8}=24.150$ | TRA | R.O | ---- | (See Section 3.6) |

*Assume working storage size is 12. Channel 130 is HOT. (See Section 3.7.)

In all cases, optimization must be based on the current operand, whether it is optimized or not. For example, consider a case where an operand is optimized on first occurrence. Subsequent references to that operand usually cannot be optimized.

Example.

| Line | LO | Op-Code | Op-Ad | NI |
| :---: | :---: | :---: | :---: | :---: |
| 1 | 70.2 | CLA | 110.3 | 120.3 |
| 2 | 120.3 | MPY | 100.4 | 100.25 |
| 3 | 100.25 | STA | 110.24 | 70.26 |
| 4 | 70.26 | ADD | 110.3 | 100.3 |
| 5 | 100.3 | STA | 72.2 | 100.4 |

In line 4 , with $\mathrm{Se}=26$, the program calls for the contents of 110.3 , which were originally used and assigned an optimum location in line 1. The choice of optimum NI for line 4 is based on the location of the read heads at the time the operand is read, in this case, sector 3.

### 3.3 OPTIMIZING ASSIGNMENTS R \& Q-LOOPS

Note that the main read head in the high speed loops, and the only write head are always located at Se, modulo 20, while the auxiliary read head is at ( $\mathrm{Se}+10$ ). Thus at anytime, t , two sectors in R and Q , spaced 10 sectors apart, are simultaneously under a read head. Therefore, for all commands which do not require writing into a sector in R or Q , the programmer has a choice of two optimum sectors in the high speed loops. In contrast, for all STORE (write) commands only one R or Q -loop sector is optimum. In view of this, the following rules apply to optimizing assignments in the high speed loops:

1. In gene ral, to optimize the R or Q operands of STORE instructions, assign the operand to sector ( $\mathrm{Se}+1$ ), modulo 20.
2. To optimize operands for all other commands, assign the operand to one of the two sectors defined by ( $\mathrm{Se}+1$ ), modulo 10 .
3. To optimize an NI in R or Q , assign the NI to any sector between ( $\mathrm{Se}+1$ ) and $(\mathrm{Se}+\mathrm{E})$, modulo 10 . If a given NI is less than ( $\mathrm{Se}+\mathrm{E}$ ), modulo 10 , then the next operand must be $\geq \mathrm{Se}+\mathrm{E}+1$.

Example.

| $\begin{aligned} & \text { Line } \\ & \text { No. } \end{aligned}$ | Location | (Se) <br> Read Head in Main Memory is © Sector | Write and Read Heads In R (\& Q) are @ Sectors: |  | Op-Code | Operand | NI |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | Main Read \& Write | Auxiliary Read |  |  |  |
| 1 | 110.30 | 30 | 10 | 0 | CLA | R. 11 | R. 1 |
| 2 | R. 1 | 31 | 11 | 1 | ADD | R. 12 | R. 2 |
| 3 | R. 2 | 32 | 12 | 2 | MPY | R. 13 | R. 4 |
| 4 | R. 4 | 54 | 14 | 4 | ADD | 102.55 | R. 5 |
| 5 | R. 5 | 55 | 15 | 5 | STA | R. 16 | R. 6 |
| 6 | R. 6 | 56 | 16 | 6 | STN | Q. $1^{7}$ | 72.57 |
| 7 | 72.57 | -- | - | -- | -- | -- | -. |

In line 1 above, the main R -loop read head picks up the operand in R.11, while simultaneously the auxiliary read head reads the contents of R.1. In line 2, precisely the same simultaneous reading of the operand and NI occurs. Because the MPY command of line 3 takes 22 word times to execute, the NI may be any sector between $\mathrm{Se}_{3}+1=33$ and $\mathrm{Se}_{3}+\mathrm{E}=(32+22)$, modulo 10 , while the next operand (line 4) must be $\mathrm{Se}_{3}+\mathrm{E}+1=(32+22+1)=55$. Note that the operand of line 4 could have been optimally assigned to any sector 55 in main memory, or to sectors 5 or 15 in $R$ or $Q$. In line 5 , since the op-code is a STORE command, only sector 16 is the optimum operand in $R$ or $Q$ (i.e., 56 , modulo 20 ).

The user is cautioned that when the operand address of a transfer or test command stored in $R$ or $Q$ is
not located in $R$ or $Q$, execution of the instruction will not be completed until the address of operand is searched and found. Therefore, for example, if one plans to affect the high speed execution of a loop which contains a test, the operand address of the test should be in the R or Q -loop, as follows:

| Line <br> No. | Control <br> Location | First Time <br> Thru Loop, Read Head in Main Memory is © Sector | Write and Read Heads in $\mathrm{R}(\& \mathrm{Q})$ are @ Sectors: |  | Op-Code | Operand | NI |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  | Main Read \& Write | Auxiliary Read |  |  |  |
| 1 | 10.137 | 137 | 17 | 7 | STA | 70.136 | R. 0 |
| 2 | R. 0 | 140 | 0 | 10 | CLA | R. 1 | R. 11 |
| 3 | R. 11 | 141 | 1 | 11 | SUB | R. 2 | R. 12 |
| 4 | R. 12 | 142 | 2 | 12 | STA | R. 1 | R. 13 |
| 5 | R. 13 | 163 | 3 | 13 | TZE | R. 14 | R. 4 |
| 6 | R. 14 | 164 | 4 | 14 | TRA | 20.13 | - |

In the above, if R .13 were coded as "TZE 20.13, R.4', the test would not be executed until location 20.13 was found in main memory. Thus, each time through the loop as much as a whole drum revolution may be lost. In contrast, when coded as in the example, full advantage is taken of the high speed registers. This "costs" one extra instruction to transfer out of the high speed register. Note that the operand of the store instruction of line 4 caused an unavoidable loss of 20 word times since it was optimally assigned in line 2 above.

### 3.4 USE OF REGISTERS FOR TEMPORY STORAGE

Very few FADAC commands use all of the registers during execution. When the commands called in a program do not use a register, that register may be used as temporary storage for operands or instructions.

It is possible to store the contents of $A$ into any register by using the proper command. Note that there is no direct "Store A into $L$ '" command. Instead, it is necessary to use a CLA* or CLS*, which cause the former contents of the A-register to be stored in the Lregister, as illustrated in Line 5 of the following example:

Example.
Line No.

Op-Code Operand
Notes

| 1 | STA | D. 1 | $(\mathrm{A}) \rightarrow$ D. 1 |
| :---: | :---: | :---: | :---: |
| 2 | STA | D. 0 | $(\mathrm{A}) \rightarrow \mathrm{D} .0$ |
| 3 | STA* | N | $(\mathrm{A}) \rightarrow \mathrm{N}$ |
| 4 | STA* | 12.0 | $(A) \rightarrow N$ and 12.0 |
| 5 | CLA* | 12.3 | Former $(A) \rightarrow-L ;(12.3) \rightarrow$ A |
| 6 | LDQ | A | $(A) \rightarrow$ All sectors of $Q$ |
| 7 | LDR* | A | $(A) \rightarrow$ All sectors of $R$ |

It is possible to transfer the contents of D. 0 or D. 1 to $A, L, N, R, Q$, or main memory using the proper command or commands in the proper sequence, as illustrated below:

| Line No. | Location | Op-Code | Operand | Nex+1 | Notes |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 2.17 | CLA | D. 0 | 2.20 | $(\mathrm{D} .0) \rightarrow \mathrm{A}$ |
| 2 | 2.20 | CLA* | A | 2.21 | $(D .0) \rightarrow L$ |
| 3 | 2.21 | STD* | M | 2.22 | (D.i) $\rightarrow M$ and $N$, where the D Sector stored depends on whether $M$ is odd or even. |
| 4 | 2.22 | STD | R. 3 | 2.23 | $(\mathrm{D} .1) \rightarrow \mathrm{R} .3$ |
| 5 | 2.23 | STD | Q. 4 | 2.24 | $(\mathrm{D} .0) \rightarrow$ Q. 4 |
| 6 | 2.24 | STD | 10.23 | 2.25 | $(D .1) \rightarrow 10.23$ |
| 7 | 2.25 | STD | 60.24 | 2.26 | $(\mathrm{D} .0) \rightarrow 60.24$ |
| 8 | 2.26 | LDR* | D. 0 | 2.46 | D.0 $\rightarrow$ all even sectors, of $R$-loop |
|  |  |  |  |  | D. $1 \rightarrow$ all odd sectors of R.loop |

(1) The contents of D. 0 or D. 1 must be in the Aregister in order to transfer them to the L-register. See Lines 1 and 2 of example above.
(2) D. 1 and D. 0 cannot be stored into each other. The instructions STD D. 1 or STD D. 0 will not store D. 0 in D. 1 or D. 1 in D. 0 . The D-loop will remain unaffected.
(3) As illustrated in line 8, when a LDR* D.i (or LDQ D.i) command is executed, all even sectors of the R (or Q ) loop will be filled with the contents of D.0, and all odd sectors will be filled with D.1.

As shown in the next example it is possible to transfer the contents of L or N to $\mathrm{A}, \mathrm{R}, \mathrm{Q}, \mathrm{D} .0$ or D.1, or main memory. L may be stored in N as in Line 5. Note that the STL instruction itself does not store L in the N -register. The flag is required.
Line No.

| 1 | 4.0 | STL | D. 1 | 4.1 | $(L) \rightarrow$ D. 0 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 2 | 4.1 | STN | D. 0 | 4.2 | $(\mathrm{N}) \rightarrow \mathrm{D} .1$ |
| 3 | 4.2 | STL | R. 3 | 4.3 | $(L) \rightarrow$ R. 3 |
| 4 | 4.3 | STN | Q. 4 | 4.4 | $(\mathrm{N}) \rightarrow$ Q. 4 |
| 5 | 4.4 | STL* | 14.3 | 4.5 |  |
| 6 | 4.5 | STL | 2.4 | 4.6 | $(L) \rightarrow 2.4$ |
| 7 | 4.6 | STN | 2.5 | 4.7 | $(\mathrm{N}) \rightarrow 2.5$ |
| 8 | 4.7 | CLA* | L | 4.10 | $(L) \rightarrow A$; <br> previous <br> $(A) \rightarrow L$ |
| 9 | 4.10 | CLA | N | 4.11 | $(\mathrm{N}) \rightarrow \mathrm{A}$ |
| 10 | 4.11 | STL* | N | 4.12 | $(L) \rightarrow N$ |
| 11 | 4.12 | STL | N | 4.13 | * |

* No flag. No storage occurs, only a one word time delay is realized.

N cannot be stored directly into L. To accomplish this, do the following:

Line No. Location Op-Code Operand Next I Note

| 1 | 2.1 | CLA | N | 2.2 | $(\mathrm{~N}) \rightarrow \mathrm{A}$ |
| :---: | :---: | :---: | :---: | :---: | :--- |
| 2 | 2.2 | CLA* $^{*}$ | 6.3 | 2.3 | $(6.3) \rightarrow \mathrm{A}$, <br> former con- <br> tents of $\mathrm{A}=$ |
|  |  |  |  |  |  |
|  |  |  |  |  |  |

Note that the operand of Line 2 above may be any location or register.

Any of the registers may be used to store instructions or data for rapid access during execution. In writing routines using this feature, make certain that the desired contents of the registers used are not destroyed before being read. In the following assume location 40.2 contains some value, X ; and location 40.1 contains the instruction MPY, N, Q.6.

| Location | Op-Code | Operand Next I | Notes |  |
| :---: | :--- | :---: | :---: | :--- |
| 20.0 | CLA | 40.1 | 20.1 | $($ MPY N. Q.6) $\rightarrow$ A |
| 20.1 | CLA* | 40.2 | 44.2 | $($ A $) \rightarrow$ L; X $\rightarrow$ A |
| 44.2 | ADD | A | R.3 | $($ A $)=2 \mathrm{X}$ |
| R.3 | STA $^{*}$ | D.1 | L | $($ A $) \rightarrow$ D.1 and N |
| (L) | MPY | N | Q.6 | $($ A $)=4 \mathrm{X}^{2}$ |
| Q. 6 | ADD | D. 1 | 20.27 | (A) $)=4 \mathrm{X}^{2}+2 \mathrm{X}$ |

Another method of using the registers for storage of instructions is shown in the following example:

Line
No.

| 1 | 2.0 | CLA | 2.1 | 2.2 | $(2.1) \rightarrow$ A |
| :--- | :--- | :--- | :---: | :---: | :--- |
| 2 | 2.1 | ADD | 4.2 | 2.4 | A. |
| 3 | 2.2 | CLA $^{*}$ | A | 2.3 | Former (A) <br> 4 |
| 2.3 | TRA* | L | 6.4 | Transfer <br> Instr to N |  |
| 5 | $(\mathrm{~L})$ | ADD | 4.2 | 2.4 |  |

Note that this example results in executing the instruction which was in line 2 from the L-register; and also causes the transfer instruction in line 4 to be stored in the N -register.

### 3.5 USE OF R AND Q-LOOPS

For rapid access, operands and instructions may be transferred from main memory to loops $R$ and $Q$ by using the instructions LDR* and LDQ. Information may be transferred back to main memory from $R$ with the STR command.

Notice that there is no command to store the contents of Q in main memory. This can be accomplished by:

| Location | Op-Code | Operand | Next I |
| :---: | :---: | :---: | :---: |
| 2.17 | LDR* $^{*}$ | Q.0 | 2.37 |
| 2.37 | STR | 2.42 | 10.62 |

In the above, the contents of Q are first loaded into R ; R is then loaded into main memory.

R and Q are loaded from one channel only. For instance, the instruction LDR* 2.177, go to 4.10, begins loading $R$ from location 2.177. The next location loaded into R is 2.0 not 4.0. All information destined for R or $Q$ is taken from 16 consecutive words of the same channel.

The same restriction is true of the STR command. It causes storage into one channel. The instruction STR 2.177 begins storing at 2.177. The next word stored from R is into 2.0. The loops will always be stored into 16 consecutive words of the same channel of main memory.

It is possible to store R into R with an instruction such as:

## STR R. 16

The contents of R will be displaced by two words; that is, R. 16 will appear in R.0, R. 2 will appear in R.4, etc. The last two transferred words will also appear in D. 0 and D.1.
$R$ and $Q$ can be loaded from any of the registers. The result is that the loop will be loaded with the same number in all of its cells, if loaded from $A, L$ or $N$; or with the two numbers in D. 0 and D.1, if loaded from the D-loop. Location Op-Code Operand Next I Notes

| 6.2 | CLA | 4.3 | 6.4 | $(\mathrm{~A})=\mathrm{X}$ |
| :--- | :--- | :---: | :--- | :--- |
| 6.4 | LDR $^{*}$ | A | 6.24 | $(\mathrm{R} .0-\mathrm{R} .17=\mathrm{X})$ |

### 3.6 INSTRUCTION MODIFICATION

For greater flexibility in a program it is often necessary to modify certain instructions. In FADAC, this can be done using standard techniques involving
standard commands. The process can be greatly simplified, however, by using the commands STORE OPERAND ADDRESS (STO) and STORE PROGRAM ADDRESS (STP), which replace the contents of one portion of an instruction in memory without affecting any other portion.

Care must be taken to introduce a delay between the time a store command which modified an instruction, is called, and the time the modified instruction itself is called. Both STO and STP require 3 word-times for execution, if M is in main memory, or 1 word time if M is in R or Q . Thus, for example, the following will cause execution of the instruction in Location 72.5 before the Operand Address in that location is modified, since the contents of the NI address are sent to the X-register at the same time that execution of the STO command is initiated. (See Figure 3.1-Multiplexing Diagram.)

Example:

| Location | Op-Code | Operand | NI |
| :---: | :---: | :---: | :---: |
| 2.4 | STO | 72.5 | 72.5 |

To properly modify an instruction, and immediately follow the modification with execution of the modified instruction, an extraneous instruction must be coded, such as:

| Line | Location | Op-Code | Operand | NI |
| :---: | :---: | :---: | :---: | :---: |
| No. |  | STO | 72.5 | 2.7 |
| 1 | 2.4 | STA | N | N |

As noted in Chapter 2, given an STO or STP command, the modified contents of $M$ (in this example $M=72.5$ ) are stored in N . The TRA N of line 2 then causes execution of the modified instruction from the N -register.

Line 2 may be coded:

Line
No.
2 (a)
2.7

CLA
A
72.5
or
2 (b)
2.7
TRA
72.5

Both 2(a) and 2(b) will cause execution of the modified instruction in 72.5.

The same precaution must be taken in the use of
oad or store commands, as illustrated in the follow-
The same precaution must be taken in the use of
all load or store commands, as illustrated in the following examples:

Line
No.
Location Op-Code
Op-Address
NI
1
110.1

CLA
110.2
112.2

2
112.2

ADD
110.3
112.3

3
112.3

STA
110.4
110.4

Op-Code
Operand
NI

Etc. ...

Execution of the command in line 3 will cause the unmodified contents of 110.4 to be sent to the X-register. To avoid this, the program should be corrected as follows:

Line
No.
Location
Op-Code
Op-Address
NI
3
112.3

STA
110.4
112.4

4
112.4

CLA
A
110.4

Line 4 is a do-nothing command inserted to allow modification of the contents 110.4 prior to transfer of control
to that location. Consider the effect of coding a LDR* or LDQ command as follows:

| Line <br> No. | Location | Op-Code | Operand | NI |
| :---: | :---: | :---: | :---: | :---: |
| 1 | 70.1 | LDR* | 70.2 | R. 0 |
| 2 | 100.27 | LDQ | 30.130 | Q.7 |

In line 1 , the contents of locations 70.2 thru 70.21 will be loaded into the R-loop. However, at the initiation of execution of this command the present unmodified contents of R. 0 (or any R sector designated) are stored in the X -register, and it is this unmodified command which is executed next. Precisely the same sequence of events occurs when the command in line 2 is executed. This may be desired. However, to execute the modified contents of a loop sector, time must be allowed for completion of the LDR* or LDQ command before control is transferred to the desired loop sector, as:

$$
\begin{array}{lllll}
\text { Line } 1 & 70.1 & \text { LDR* } & 70.2,72.21 \\
\text { Line 1a } & 72.21 & \text { CLA } & \text { A, R. } 0 \\
\text { Line 2 } & 100.127 & \text { LDQ } & 30.130, & 100.47 \\
\text { Line 2a } & 100.47 & \text { TRA } & \text { Q. } 7
\end{array}
$$

In lines 1 and 2, the NI of the load instruction is coded $\geq 17$ word times past the operand sector, to introduce the necessary delay. Transfer of control to the desired loop sector is achieved in line 1 a by a "do-nothing" CLA A with a loop sector NI; and in line $2 a$ with an unconditional transfer to a loop sector.

### 3.7 STORE INSTRUCTION RESTRICTIONS

Because the Write heads of all channels designated as permanent storage (COLD) are de-energized when
the MLU is not connected to FADAC, the programmer must exercise care in the use of all STORE instructions, such that either (1) the Operand of a STORE instruction is in HOT memory or, (2) the information is temporarily stored in a register, then either retrieved from the register and sent to HOT memory for later use, or executed from the register before the contents of that register are changed.

In the following example, assume working storage size is 4 channels (i.e., channels $70-76$ only are HOT).

Example.

| Line <br> No. | Location | Op-Code | Operand | NI |
| :---: | :---: | :---: | :---: | :---: |
| 1 | 2.60 | CLA | 4.61 | 2.61 |
| 2 | 2.61 | ADD | 4.62 | 2.62 |
| 3 | 2.62 | STA | 74.61 | 2.63 |
| 4 | 2.63 | STO | 4.64 | 2.66 |
| 5 | 2.66 | STN | 74.67 | R.0 |

In line 3, the operand is in HOT memory, so that proper storage occurs. In line 4 the instruction in 4.64, modified by the operand in the A-register is sent to N , and in line $5, \mathrm{~N}$ is sent to working storage for later use. The contents of 4.64 are not modified.

See Section 2.2.2 for a detailed description of all store instructions.

### 3.8 LINKAGE TO SUBROUTINES

Entry to a subroutine and subsequent return to the main program is facilitated by using the transfer instructions in conjunction with the STORE PROGRAM

ADDRESS command. In the following example, the entry to the subroutine is made with a transfer instruction in the main program. The NI of this instruction is the return address back to the main program. Upon execution of the transfer instruction, the flag bit puts the entire transfer instruction into N , and control goes to the subroutine. The first instruction of the subroutine will take the return address (NI) out of N and put it into the last instruction of the subroutine; and simultaneously store the argument in A into the L-register. Upon completion of the subroutine, control returns to the main program.

Example.


It is sometimes desirable to store a complete word into the last location of the subroutine, rather than the return address only:


In this example, assume that the MLU is disconnected at execution time. Therefore, Channel 6 is in cold memory, and cannot be written into. Line 1 causes transfer to a subroutine which begins in Channel 6. The first word of the subroutine (line 3) brings the transfer instruction to $A$. The STP command of line 4 sends the contents of 6.12 to the N-register, with the NI in $N$ changed to contain the NI currently in the A-register. Thus, on execution of the command in line 4 , the $N$ register contains

| Op-Code | Operand | NI |
| :---: | :---: | :---: |
| STD | 74.16 | 64.22 |

where the NI address (64.22) is the return to main program. The contents of 6.12 remain unchanged. The command in line 6 stores the modified contents of $N$ into the last word of the subroutine (line 8 ).

### 3.9 TABLE LOOK-UP

The FADAC instruction repertoire includes two table look-up commands, Equal Search (EQS) and Greater

Than or Equal Search (GES), both of which operate in a similar manner. (See Section 2.2.5.) To illustrate a typical application, consider the following equation:

$$
\mathrm{Y}=\mathrm{B}_{\mathrm{x}} \mathrm{X}+\mathrm{X}^{2}
$$

Assume the value of the coefficient $B x$ to be used in evaluating this equation depends on the computed value of X . If a prestored value of X is greater than or equal to the computer value of X , then a particular value of the coefficient Bx is to be submitted in the equation. The programmer must construct a table for the variable coefficients such that each stored value of X is followed by the value to be assigned to the coefficient. Consider the structure of a table for the value of $\mathrm{B}_{\mathrm{x}}$. Values of X are stored in the table in ascending order, each followed by the $\mathrm{B}_{\mathrm{x}}$ to be used with that computed X which is within a given range of values. X is scaled at $0 ; \mathrm{B}_{\mathrm{x}}$ at $(-2)$.

| Location | FADAC Octal Word |  | Remarks |
| :---: | :---: | :--- | :--- |
| 2.170 | 37757000000 |  | $\mathrm{X}=-0.002$ |
| 2.171 | 00015770000 |  | $\mathrm{~B}_{\mathrm{x}}=0.000427$ |
| 2.172 | 37767600000 |  | $\mathrm{X}=-0.001$ |
| 2.173 |  | 00013150000 |  |
| 2.174 |  | $\mathrm{~B}_{\mathrm{x}}=0.000342$ |  |
| 2.175 | 00010140000 |  | $\mathrm{X}=0.001$ |
| 2.176 | 0004061000 |  | $\mathrm{~B}_{\mathrm{x}}=-0.000125$ |
| 2.177 |  | 00001424400 |  |

Assume that in our problem, X has been evaluated as 0.004 . The following program segment will search for the proper value of $\mathrm{B}_{\mathrm{x}}$ :

| Location | Op-Code | Operand | Next I | Notes |
| :---: | :---: | :---: | :---: | :---: |
| 4.20 | CLA | MASK | 4.21 | $(A)=37777777777=$ <br> MASK |
| 4.21 | CLA* | $X$ | 4.22 | $\left\{\begin{array}{l}(A)=x=0.004 \\ (L)=37777777777\end{array}\right.$ |
| 4.22 | GES | 2.170 | 4.00 | Search table for <br> $X \geq 0.004$ |
| 4.00 | STA | 76.1 | 4.01 | $\left\{\begin{array}{l}(A)=B_{x}=0.000047 \\ \text { Store } B_{x} \text { in 76.1 }\end{array}\right.$ |

The GES command searches the table until it finds a word that is greater than or equal to X , at which time the next word in the table is read into $A$. When the proper $\mathrm{B}_{\mathrm{x}}$ has been found, control is transferred to the next instruction. Note that $L$ is filled with the mask before the Search is initiated.

The last sector searched in any channel is 176. Therefore, if the proper X is not found before the search reaches the end of the channel, the A-register will contain the same number as it did at the beginning of the search.

### 3.10 MEMORY II AND MEMORY III ADDRESSING TECHNIQUE

In its original design, FADAC memory contained only 32 channels plus the A, N, L, D, R, and Q-registers. When the design was expanded to the present 64 channel version, it became necessary to devise a special method
to address all channels and registers without increasing the size of the FADAC word. Since a 32 bit FADAC word allows only 7 bits (see Section 2.1) to identify a channel address, an operand or NI address in channels numbered 100 through 136 (Memory II) has the same bit configuration as addresses in channels 300 through 336 (Memory III). To enable the computer to identify the desired channel in these ambiguous areas, FADAC hardware includes a flip-flop which is set to its OFF state to interpret channel address 100-136 as Memory II and Reset to its ON state to interpret the same bit configuration as Memory III.

A 146 in the channel address of the NI of a given instruction is required to change the state of this flipflop. Initially, when FADAC is turned on or when any control button on the front panel of FADAC is depressed, the flip-flop is set for Memory II addressing. The first 146 NI channel address sensed will set the flip-flop to its ON state, either permanently when the 146 NI is used with a TRA operation code, or momentarily when the 146 NI is coded with any other operation code. In all cases, when a 146 NI channel address is coded with a TRA operation code, the change in the state of the memory control flip-flop remains fixed until the next TRA-146 combination is sensed, or until a control button is depressed. In contrast, a 146 NI channel address coded with a non-TRA operation code causes a change in the state of the flip-flop which lasts only until the instruction is executed. See Section 3.11.

No 146 NI is required to call or transfer control to the registers $A, N, L, D, R$, or $Q$.

Within a FADAC word, any reference to a channel in Memory III is coded as 100 through 136. The leading digit, 3 , of a channel address is only used when the location field is specified, as on an input tape or on
keyboard entry. In these cases the five octal digit address must be followed by the octal location code (a period). (See Section 2.2.7,)

Case I through Case IV specify the rules which must be used to properly address Memory II and III. Case I - Location Field in Memory I.

Memory Control flip-flop set to its OFF state to interpret channels 100 through 136 as Memory II.
a. To call an operand in Memory I or II, no 146 is required.

Example.la.
FADAC Octal

| Lo | Op-Code | Operand | NI | Lo | Word |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 72.4 | CLA | 100.5 | 74.5 | 07204. | 07405250005 |
| 74.5 | ADD* | 70.5 | 72.6 | 07405 | 27206007006 |
| 72.6 | STA | 70.5 | 72.7 | 07206. | 07207507005 |
| 72.7 | STA | 100.5 | 70.10 | 07207. | 07010510006 |

b. To call an operand in Memory III, a 146 must appear as the NI channel address, to momentarily set the flip-flop to its ON state.

Example 1b.
FADAC Octal

| Lo | Op-Code | Operand | NI | Lo | Word |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 00.21 | CLA | 300.22 | 146.22 | 00021 | 14622250022 |
| 00.22 | ADD. | 300.23 | 146.23 | 0002 | 14623010023 |
| 00.23 | ADD | 70.24 | 72.24 | 00023. | 07224007024 |
| 72.24 | STA | 70.25 | 000.25 | 07224. | 00025507025 |

c. To transfer control to Memory I or II, no 146 is required with a TRA operation code.

Example lc.
FADAC Octal

| Lo | Op-Code | Operand | NI | Lo | Word |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 00.136 | CLA | 70.137 | 72.137 | 00136. | 07337247137 |
| 72.137 | TRA* | 100.140 | 00.026 | 07337. | 20026150140 |
| 100.140 | - | - | - | 10140. | - |
|  | - | - | - | - |  |
|  | - | - | - | - |  |
| 40.06 | TRA | 74.07 | 00.00 | 04006. | 00000147407 |

d. To transfer control to Memory III, the flip-flop must be reset to its ON state. Therefore, a 146 must appear in the NI field with a TRA operation code.

Example 1d.

| Lo | Op-Code | Operand | NI | FADAC Octal |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | Lo | Word |
| 20.25 | CLA* | 70.26 | 20.26 | 02025. | $\frac{22026247026}{(\text { Note 1) }}$ |
| 20.26 | TRA* | 300.27 | 146.27 | 02026. | $\begin{aligned} & \underline{34627150027} \\ & \text { (Note 1) } \end{aligned}$ |
| 300.27 | ADD | 300.30 | 320.30 | 30027. | 12030010030 |

Note 1:
Recall that the first digit of a FADAC octal word is the sum of the flag bit and the first channel address bit.

Case II - Location Field in Memory I.
Memory Control flip-flop set to its ON state to interpret channels 100 through 136 as Memory III (channels 300 through 336).
a. To call an operand in Memory I or III, no 146 is required.

Example Ila.
FADAC Octal

| Lo | Op-Code | Operand | NI | Lo | Word |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 6.0 | CLA | 300.1 | 72.1 | 00600. | 07201250001 |
| 72.1 | ADD | 4.2 | 300.2 | 07201. | 10002000402 |
| 300.2 | ADD | 300.3 | 40.3 | 30002. | 04003010003 |

b. To call an operand in Memory II, a 146 must appear as the NI channel address to momentarily reset the flip-flop to its OFF state.

Example IIb.
FADAC Octal

| Lo | $\underline{\text { Op-Code }}$ | Operand | NI | Lo | Word |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 70.5 | SUB* | 100.6 | 146.6 | 07005. | 34606030006 |
| 70.6 | ADD | 102.7 | 146.7 | 07006. | 14607010207 |
| 70.7 | MPY | 110.10 | 146.32 | 07007. | 14632211010 |
| 70.32 | --. | ... | ... | 07032. | -.... |

c. To transfer control to Memory I or III, no 146 is required.

Example IIc.

| Lo | $\underline{\text { Op-Code }}$ | Operand | NI | FADAC Octal |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | Lo | W ord |
| 6.107 | CLA | 310.110 | 70.110 | 00707. | 07110251110 |
| 70.110 | TRA* | 306.111 | 70.116 | 07110. | 27116150711 |
| 306.111 | TRA | 04.121 | 0.0 | 30711. | 00000140521 |

d. To transfer control to Memory II, a 146 must appear as the NI channel address with a TRA operation code.

Example IId.

| Lo | Op-Code | Operand | NI | FADAC Octal |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | Lo | Word |
| 22.6 | CLA | 320.7 | 24.7 | 02206. | 02407252007 |
| 24.7 | TRA* | 100.10 | 146.10 | 02407. | 34610150010 |
| 100.10 | CLA | 136.11 | 100.11 | 10010. | 10011253611 |

Case III - Location Field in Memory II.
Memory Control flip-flop set to its OFF state to interpret channels 100 through 136 as Memory II.
a. To call an operand in Memory I or II, no 146 is required. Example Illa.

| Lo | Op-Code | Operand | NI | FADAC Octal |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | Lo | Word |
| 102.57 | TOV | 50.60 | 106.60 | 10257. | 10660165060 |
| 106.60 | ADD | 126.61 | 130.61 | 10660. | 13061012661 |
| 130.61 | STA | 70.60 | 100.62 | 13061. | 10062507060 |

b. To call an operand in Memory III, a 146 must appear in the NI channel address, since the flip-flop must be momentarily reset to its ON state. Example IIIb.

FADAC Octal

| Lo | Op-Code | Operand | NI | Lo | Word |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 134.7 | CLA | 310.10 | 146.10 | 13407. | 14610251010 |
| 134.10 | ADD | 70.11 | 130.11 | 13410. | 13011007011 |
| 130.11 | ADD | 334.12 | 146.12 | 13011. | 14612013412 |
| 130.12 | MPY | 306.13 | 146.34 | 13012. | 14634210613 |
| 130.34 |  |  |  | 13034. |  |

c. To transfer control to Memory III, a 146 must appear in the NI channel address with a TRA operation code.

Example IIIc.

| Lo | Op-Code | Operand | NI | FADAC Octal |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | Lo | Word |
| 126.130 | CLA | 100.131 | 124.131 | 12730. | 12531250131 |
| 124.131 | TRA | 302.132 | 146.00 | 12531. | 14600150332 |
| 302.132 | ADD | 306.133 | 302.133 | 30332. | 10333010733 |
| 302.133 | --- | --- | $\cdots$ | 30333. | - |

d. To transfer control to Memory I or II, no 146 is required with a TRA operation code.
Example IIId.

| Lo | Op-Code | Operand | NI | FADAC Octal |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | Lo | Word |
| 126.0 | CLA | 110.1 | 100.1 | 12600. | 10001251001 |
| 100.1 | TRA | 70.2 | 0.0 | 10001. | 00000147002 |
| 70.2 | TRA* | 136.3 | 0.0 | 07002. | 20000153603 |
| 136.3 | TRA | 106.4 | 0.0 | 13603. | 00000150604 |
| 3.28 |  |  |  |  |  |

Case IV - Location Field in Memory III.
Memory Control flip-flop is set to its ON state to interpret channels 100 through 136 as Memory III.
a. To call an operand in Memory I or III, no 146 is required.
Example IVa.

| Lo | Op-Code | Operand | NI | FADAC Octal |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | Lo | Word |
| 306.21 | CLA | 76.22 | 304.22 | 30621. | 10422247622 |
| 304.22 | ADD | 00.23 | 306.23 | 30422. | 10623000023 |
| 306.23 | STA | 70.22 | 306.24 | 30623. | 10624507022 |

b. To call an operand in Memory II, a 146 must be coded in the NI field.
Example IVb.

| Lo | Op-Code | Operand | NI | FADAC Octal |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | Lo | Word |
| 324.02 | MPY | 100.03 | 146.24 | 32402. | 14624210003 |
| 324.24 | ADD* | 110.25 | 146.25 | 32424. | 34625011025 |
| 324.25 | STA | 70.24 | 306.26 | 32425. | 1062650702 |

c. To transfer control to Memoryl, or II I, no 146 is required with the TRA operation code Example IVc.

| Lo | Op-Code | Operand | Ni | FADACC Octal |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | م | Word |
| 336.130 | CLA | 300.131 | 322.131 | 33730. | 12331250131 |
| 322.131 | tra | 306.132 | 0.0 | 32331. | 00000150732 |
| 306.132 | TRA* | 10.133 | 0.0 | 30732. | 20000141133 |
| 10.133 | TRA | 314.134 | 0.0 | 01133. | 00000151534 |

d. To transfer control to Memory II, a 146 must appear in the NI field with a TRA operation code. Example IVd.

| Lo | Op-Code | Operand | NI | FADAC Octal |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | Lo | Word |
| 320.10 | CLA | R.11 | 320.11 | 32010. | 12011254211 |
| 320.11 | tra | 100.12 | 146.12 | 32011. | 14612150012 |
| 100.12 | ADD | 120.13 | 122.13 | 10012. | 12213012013 |
| 122.13 | ADD | R. 14 | D. 1 | 12213. | 16001014214 |
| (D.1) | TRA | 322.14 | 146.00 | (executed | from D.1) |
| 322.14 | STA | D.0. | 70.22 | 32214. | 07022516000 |
| 70.22 | TRA | Q. 3 | L | 07022. | 17200155203 |

Example IVd also illustrates that no special coding is required to address the registers $\mathrm{A}, \mathrm{L}, \mathrm{N}, \mathrm{R}, \mathrm{Q}$, or D from Memory I, II or III, regardless of the state of the memory control flipoflop.

Note that control may be in Memory I regardless of the memory control flipoflop (see Cases I- and II); while control can only be in Memory II when the flip-flop is set to its OFF state, and in Memory III when the flipflop is set to its ON state.

Channel Return Given a 146-NI Channel Address.
As each instruction is processed, its NI is automatically stored in a special NI address buffer. The contents of the buffer are updated to a new NI address when execution of an instruction is completed and the next instruction read. However, when the new NI contains a 146 channel address, the contents of the channel portion of the NI buffer remain unchanged; that is the last non- 146 NI channel address remains in the NI address buffer until a new non-146 is
sensed. During execution of each instruction, FADAC hardware examines the NI buffer, and causes transfer of control to the channel sector specified when execution is completed.

## Example.

| Line No. | Location Field | Op.Code | Op-Address | NI | Notes |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 0.0 | CLA | 110.1 | 100.1 | Program is started by De. pressing SM (LOCATION 0.0) FF is set OFF (A\|l 100 through 136 references are interpreted as Memory II.) |
| 2 | 100.1 | ADD | 70.2 | 70.3 |  |
| 3 | 70.3 | STA | 70.4 | 72.4 |  |
| 4 | 72.4 | TRA | 110.2 | 000.0 | Transfer from Memory $\mid$ to Memory II, No 146 necessary since FF is set to interpret 100 through 136 as Memory 11 by line 1 . |
| 5 | 110.2 | ADD | 130.3 | 72.3 |  |
| 6 | 72.3 | ADD | 330.4 | 146.4 | Momentarily reset FF to ON to pick up operand in Memory 111. |
| 7 | 72.4 | TRA | 330.5 | 146.0 | FF reset to ON to transfer control to Memory III. |
| 8 | 330.5 | MPY | 300.6 | 70.27 | No 146 necessary to return io Memory 1 . |
| 9 | 070.27 | STA | 70.26 | 300.30 | No 146 necessary to return to Memory 111 , since FF set ON by line 7 . |
| 10 | 300.30 | ADD | 100.31 | 146.31) | Momentarily set $F F$ to OFF to get Operands from Memory |
| 11 | 30031 | ADD | 104.32 | 146.32) | 11. Return to channel 300 as |
| 12 | 300.32 | ADD | 110.33 | 146.33) | each instruction is executed. executed. |
| 13 | 300.33 | ADD | 120.34 | $146.34)$ |  |
| 14 | 300.34 | TRA | 110.36 | 146.00 | FF reset to OFF to transfer control to Memory 11. |
| 15 | 110.36 | STA | 70.36 | 70.37 |  |
| 16 | 70.37 | TRA | 110.50 | 100.50 | No 146 necessary since FF set OFF by line 14. |

In the example above, computation is initiated by depressing the Sample Matrix button (location 0.0 ), and the NI address buffer is filled with 100.01 . When the instruction in line 1 is executed, control is sent to 100.01. As each instruction in lines 2 through 5 is executed, the contents of the NI buffer are updated. However, when the instruction in line 6 is read, only the sector portion of the NI buffer is changed. Thus, during execution of the line 6 instruction, the NI buffer contains 72.04. When execution is completed, control is transferred to location 72.04 and the contents of the NI buffer changed to 72.00 . Since line 7 contains a TRA instruction, the contents of the NI buffer are ignored, and control is transferred to the operand $=330.05$. After execution of the instruction in line $8,300.30$ is stored in the NI address buffer. The channel portion of the buffer remains unchanged throughout the execution of the instructions in lines 10 through 14 , since each contains a 146 NI channel address. Note that the return to the last non- 146 NI channel is automatic, but that the desired sector must be coded.

### 3.11 USE OF 146 NI

When a TRA-146 combination is sensed, transe ferring control to the opposite side of memory, and information is immediately desired from the original side of memory, special care must be taken to keep track of the last non-146 NI sensed. Consider the following example.

Example.

| Lo | Op-Code | Op-Ad | NI-Ad | Line \# | Remarks |
| :--- | :---: | :---: | :---: | :---: | :---: |
| 120.177 | STA | 70.176 | 110.2 | 1 |  |
| 110.0 | TRA | 300.1 | 146.0 | 2 | Reset flip-flop <br> to Memory III |
| 300.1 | CLA | 100.2 | 146.2 | 3 |  |
| 310.2 | ADD | $\ldots$ | $\ldots$ | 4 |  |

Note that the location field of line 4 is in channel 310. This is true since the last non-146 NI channel address, 110 , appearing in line 1 , is interpreted as 310 , because of the TRA-146 coded in line 2 . In contrast, consider:

Lo Op-Code Op-Ad NI-Ad Line \# Remarks

| 120.177 | STA | 70.176 | 72.0 | 1 |
| :---: | :---: | :---: | :---: | :---: |
| 72.0 | SRA | 300.1 | 146.2 | 2 |
|  |  |  |  |  |
| 300.1 | CLA | 100.2 | 146.2 | 3 |
| 72.2 | ADD | 300.3 | 302.3 | 4 |
| 302.3 | $\ldots .$. | $\ldots$ | $\ldots$. | 5 |

Here the location field of line 4 is in channel 72, which is the last non-146 NI channel address sensed. Since no ambiguity exists in Memory I addresses, the channel address in the NI buffer, 72, is not affected by the TRA-146 in line 2.

If the programmer intends a return to channel 300 , in the above examples, then a dummy instruction must be inserted to reset the NI buffer before a call for information from the original side of memory:

| Lo | Op-Code | Op-Ad | Nl-Ad | Line \# | Remarks |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 120.177 | STA | 70.176 | 110.0 | 1 |  |
| 110.0 | TRA | 300.1 | 146.0 | 2 | Reset flip-flop <br> 300.1 |
| STA | A | $\underline{300.2}$ | 3 | Dummy instruc- <br> tion to Reset |  |
| 300.2 | CLA | 100.3 | 146.3 | 4 | Nl buffer <br> Clear and add <br> contents of <br> 100.3 and go <br> to 300.3 <br> Lo field inter. |
| 300.3 | ADD | 302.4 | 300.4 | 5 | preted as 303.3 <br> as desired. |

Whenever a Test for Zero (TZE) or Test for Plus (TPL) (branching) instruction is used, special precautions must be taken:
Example.

| Location of <br> Command | Op-Code | Op-Address | Next Instr' <br> Address | Remarks |
| :---: | :---: | ---: | :---: | :--- |
| 300.15 | CLA | 320.16 | 300.17 |  |
| 300.17 | MPY | 70.20 | 300.20 |  |
| 300.20 | ADD | 320.42 | 300.43 |  |
| 300.43 | TPL | 310.44 | 312.44 | If plus go to |
|  |  |  |  | 310.44 |
|  |  |  |  | If minus go |
|  |  |  |  | to 312.44 |

Note that the command in location 300.43 will cause transfer to channel 310 if the result of the computation is plus, or to channel 312 if the result is minus. In either case, 312.44 is sensed as the Next Instruction address, and is therefore stored in the NI channel buffer. Observe the confusion if the branching instruction is followed by a call for information stored on the opposite side of memory, as shown in the following: Example.

| Location of <br> Command | Op-Code | Op-Address | Next Instr' <br> Address |  |
| :---: | :---: | :---: | :---: | :---: |
| 300.43 | TPL | 310.44 | 312.44 |  |
| 312.44 | CLA | 100.45 | 146.46 |  |
| 312.46 | STA | 70.45 | 312.47 |  |
| 312.47 | HLT | $\ldots$ | 0.0 |  |
| • | P |  | . |  |
| P | P |  | . |  |
| 310.44 | CLA | 102.45 | 146.46 | Control goes |
| $310-46$ | ADD | 300.47 | 310.47 |  |
| 310.47 | STA | 300.48 | 310.48 |  |
|  |  |  | etc. |  |

Should the computation yield a negative result, no problems exist, since the information in 100.45 will be sent to the A-register, after which control will return to channel 312, sector 46 as desired. However, if the result is positive, a transfer to 310.44 will occur, the contents of 102.45 will be sent to the A-register and since the NI channel buffer still contains the 312 sensed as the last non-146 NI address, control will then go to 312.46 , rather than to the 310.46 desired, yielding a program error. Since the programmer cannot usually predict which branch will be taken, he is required to either (1) program the same channel for each branch of a conditional transfer, or (2) avoid a 146 NI channel reference immediately following a conditional transfer instruction by the insertion of a dummy instruction which will change the address stored in the NI channel buffer, as follows:

## Example.

| Location of <br> Command | Op-Code | Op-Address | Next Instr' <br> Address | Remarks |
| :---: | :---: | :---: | :---: | :---: |
| 300.43 | TPL | 310.44 | 312.44 |  |

(Continued)

Example. (Continued)

| Location of Command | Op-Code | Op-Address | Next Instr' <br> Address | Remarks |
| :---: | :---: | :---: | :---: | :---: |
| 310.46 | ADD | 300.47 | 310.47 | Add contents |
| . | - | . | . | of 300.47 to |
| - | - | - | . | A-Reg. Go to |
| . | . | . | . | 310.47 etc . |

### 3.12 WORD INPUT/OUTPUT

The rules specified for designation of Memory II and Memory III are applicable in the execution of all commands except Parallel I/O by word. Recall that these commands require that the L-register be programmed to contain the starting location into which input is to be stored, or from which output is to be fetched. When input is in Octalcode, use of the 146 to differentiate between Memory II and Memory III locations is ineffective. Rather, the value of bit 18 in the L-register when an Enter code is sensed controls FADAC's interpretation of the storage location. When this bit equals 0 , the starting location is interpreted as in Memory I or II. When bit 18 of L equals 1 , the storage location is interpreted as in Memory III.

For all other Parallel I/O, it is recommended that the following procedure be adopted:

1. Program all Parallel I/O Commands into Memory I.
2. Code the NI of all Parallel I/O Commands into Memory I.
3. Program a zero to $L$ before setting up the starting location in L.
4. If the memory control flip-flop is set to interpret all 100 channels as in Memory II and the starting store (for READ) or fetch (for WRITE) location is in Memory III, code the following sequence of instructions.
a. TRA (to a Memory I lo), 146 XX .
b. Set $L$ to the starting Memory III location.
c. Code the desired Read or Write Command with an NI in Memory I.

## Example.

Control is in Memory II. Read External Device, 9 words, in BCD code, storing the input into location 306.10 through 306.2 and return to 104.27 .

| Lo | Op-Code | Operand | NI | Explanation |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 102.21 | ZEL | 0 | , | 102.22 | 0 |

104.27

If the memory control flip-flop is set to interpret all 100 channels as in Memory III, and the starting store or fetch location is in Memory II, precisely the same sequence of instructions should be coded.

## Example.

Control is in Memory III. Write to an External Device, 18 words in Alpha 5 code, starting at location 105.44, and return to 310.27 .

| Lo | Op-Code | Operand | NI | Explanation |
| :---: | :---: | :---: | :---: | :---: |
| 304.21 | ZEL | 0 |  | 304.22 |

If this procedure is used, the danger of loss of control of the setting of the Memory control flip-flop during I/O is eliminated.

## CHAPTER 4

## FADAC OPERATING PROCEDURES

### 4.1 CONTROL PANEL

Functional Description.
The computer control panel is shown pictorially in Figure 4.1. Primarily, the control panel is divided into five basic functional areas as follows:

1. Controls and Indicators
2. Keyboard Input Device
3. Mechanical Tape Reader
4. Matrix Input Device
5. Nixie Display Panel

The following is a summary of the functions of the operating controls and indicators of FADAC.

PWR ON, PWR OFF Switch.
When flicked upward, this switch actuates the power control circuits in the computer, energizing in the proper sequence the power supplies, blowers, and memory. When pulled down, the computer is de-energized. PWR READY (Indicator).

This neon indicator lights after a delay (approximately 20 seconds) suitable for insuring the memory is fully operational. The light remains in its ON condition until the computer power is turned off. When the computer is placed in a Marginal Test Condition, the indicator blinks until normal operation mode is restored.

## $\stackrel{+}{\sim}$



Figure 4.1. Production Engineering FADAC

Normally ON, this neon indicator blinks when a line voltage transient has occurred, or when the input voltage exceeds high or low tolerances. If the blinking cannot be reset with the RESET button, the input voltage is approaching either a low or high voltage kickout condition.

Temp. (Indicator).
Normally ON, this neon indicator blinks when the internal temperature of the machine is approaching marginal operation. The computer may usually be used, however, until temperature kick out actually occurs. Any system application of the computer must insure that the control panel and computer filters are accessable for cleaning, and are provided with unimpeded air flow. Additionally, a cut-out switch is provided on the rear of the computer to de-energize one of the internal blowers for cold temperature operation. This feature is also used as a maintenance tool to determine if proper blower action is present when a temperature warning occurs.

## Parity (Indicator).

Normally ON, this neon indicator blinks when a parity error is detected, either when reading from memory or in reading from an input device in Fieldata or other parity-sensitive code. To clear the blinking, press the RESET button. If this does not correct the condition, turn FADAC OFF, then ON again, and proceed as desired.

## Error (Indicator).

Normally ON, this neon indicator blinks when an overflow error has occurred in computation or when an error is sensed with the computer in the VERIFY mode.

Normally ON, this indicator blinks when a No Solution Light (NSL) Command has been executed. It will continue to blink until one of the Discreet I/O Commands (DOF, OD1, OD2, OD3, ID1, ID2, or ID3) is sensed. The indicator may be programmed to blink to indicate an insoluble problem, or that an incorrect procedure has been employed.

Compute (Indicator).
Normally OFF, this neon indicator lights during the periods when the computer is in a computing state. It may also be turned ON and OFF under program control with execution of an Initiate Compute Mode (ICM) or Halt Compute Mode (HCM) instruction. Thus, the the light can be kept off while the computer is in a compute mode. This function is usually reserved for periods when the computer is either self-testing the memory arithmetic loop registers, or scanning the discrete input/output lines for evidence of an input from (or output to) an external device.

In/out (Indicator).
Normally OFF, this neon indicator lights when information is being transferred to or from an external device. It is automatically controlled by the execution of any computer READ or Write Command.

Keyboard (Indicator).
Normally OFF, this neon indicator lights when the computer calls for information from the keyboard. The light is automatically controlled by the execution of a READ KEYBOARD command.

## Time Meter.

A running time meter is provided on the control panel to show total elapsed operation time of the computer.

Marginal Test (Switch).
While not a physical part of the control panel, a five position switch is provided on the left side of the computer main frame. Normally OF F, this switch provides four different combinations of marginal voltage. When the switch is placed in any of the marginal test positions, the PWR READY indicator will flicker.

Reset (Switch).
This switch, when momentarily depressed, resets the computer to the program or manual Halt mode from the Compute or Input/Output modes. It is used when a parity or overflow error has occurred; or to reset the Temperature Warning or Transient neon indicators when these are blinking.

## Input Matrix.

The matrix consist of 64 windows arranged in an $8 \times 8$ array. The buttons used to select a specific input windows, corresponding to an assigned input parameter, are located on the left-hand column (lettered A through H ) and on the bottom row (numbered 1 through 8). Selection of a given window sets binary information into bit positions 25 through 30 of the A-register, when a DIA command is executed (see Section 2.2.6). In addition to the $8 \times 8$ matrix there are several auxiliary buttons on the Matrix panel as described below:
a. Five buttons on the right hand bank (lettered A through E) which set binary information into bit positions 19 through 23 of the A-register when a DIA command is executed. These buttons are wired directly to the Battery tube of the Nixie display.
b. Two switches on the right hand bank (numbered 1 and 2) which set binary information into bit position 24 of the A-register when a DIA command is executed.

The mechanical tape reader is used for entering data into FADAC. The input code may be either a 5 level teletype code or an 8 level Fieldata code. Starting or stopping of the tape reader may be internal under program control, or external by depressing the proper sequence of buttons. The running speed of the tape reader is approximately 10 characters per second.

## Keyboard.

The computer keyboard is used for manual input of information. It contains keys for the digits 0 through 9 , one key each for + and - , and three control keys (Clear, ., and Enter).

## Control Buttons.

If the computer has been set into the program Halt state via the program, or by activation of the RESET switch, the control buttons may be used to reinitate computation. The buttons serve primarily as a hardware "bootstrap"' to initiate execution of a program.

## Nixie Display.

The nixie display is a programmed visual display. Information to be displayed must first be converted into Binary Code Decimal characters and stored properly in the 2 word D-loop. The display consists of 17 nixie tubes and 16 neon lamps used for decimal point display. Figure 4.2 is a diagram showing the relationship between the Nixie Lights and the D-register. (The Nixie tubes are numbered in the illustration to facilitate discussion). As shown, Nixie tubes 1 through 9 display the contents of the 32 bit positions in D.1; tubes 9 through 15 display the 28 most significant bit positions in D.0. Tubes 16 a and 16 b are both connected to the four least significant bit positions of D. 0 (bits 28-31).


NOTE: The 4 least significant bits in $D$. O will energize either the sign or the charge nixie, but not both.

Figure 4.2. Relationship Between Nixie Lights and D-Register

If these bit positions contain a digit ( BCD 0 to 9 ), tube 16b (the CHARGE nixie) will be activated. If they contain a sign (BCD 12 or 13), tube 16a (the SIGN nixie) will activate. A sign coded in any position other than the four least significant bits of D.O will be interpreted as a blank. Therefore, care must be taken to properly position a sign which is to be displayed. (See Appendix C for a list of BCD codes.) The Nixie tube labelled BATTERY is not part of the D-loop. Rather, it is wired directly to the auxiliary buttons labelled A through E , on the right-hand side of the Matrix panel. Note that no nixie display will occur unless one of the Battery buttons is depressed.

### 4.2 FADAC FIELD TABLE

The Field Table contains the computer power supply, which must be-connected to a 3 -phase, 400 HZ 120/208 volt source. The Field Table has 5 plugs labelled J-1 through J-5. J-5 is connected to the source voltage. J-1 through J-4 are identical, and are used for any peripheral equipment (such as the SDR or Teletypewriter) which requires the same source voltage. A cable, P-11, which is permanently fixed to the bottom of the Field Table behind J-4, is connected to J-11 on FADAC to energize the computer.

### 4.3 SIGNAL DATA REPRODUCER (SDR)

The Signal Data Reproducer (SDR) is the primary input device used with FADAC. It is a high-speed, photoelectric reader which operates at approximately 600 characters per second. Data input to the SDR may be 5 level Teletype, or 8 level Fieldata code. Figure 4.3 is a photograph of the SDR which shows the tape cannister, read head, operation controls, indicators, and SDR output plugs. To connect the SDR to FADAC, SDR-PA is connected to one of the power supply


Figure 4-3. Signal Data Reproducer (SDR)
plugs (J-1 through J-4) on the Field Table.
The following is a description of the SDR controls and indicators:

## Tape Advance On/Off.

Normally OFF, this switch is thrown to ON to cause the tape to mechanically advance to any desired position. No reading occurs while this switch is in the ON position.

Fieldata/Teletype.
The position of this switch determines whether information input to FADAC is interpreted as Fieldata or Teletype Code.

Fill/Verify Switch.
The position of this switch is immaterial, since it has no effect in either position. The switch was designed to control the input and verification of a continuous tape, in which the ends of the tape have been spliced together to form an endless loop.

## Start FADAC.

This button is depressed to initiate Octal input to FADAC under external control. (Note that BCD or Alphanumeric input is only possible under program control.)

Compute Run/Halt.
This switch is thrown into the HALT position when input is under external control; and into the RUN position when FADAC is under program control.

Auxiliary Memory Switch.
With this switch in the OFF position, all FADAC write heads are energized, i.e., all memory is HOT. With the switch in the ON position, only the write heads
associated with HOT memory, as determined by the position of the HOT-Channel Select Switch, are energized. (See Section 4.5.)

## Stop.

This button is depressed to manually stop the SDR tape.

Signal On-Off.
This switch is thrown to the ON position to energize the light source in the SDR Read head.

Power On/Off.
This switch is thrown to the ON position to energize the SDR.

Fill/Verify Indicators.
These neon bulbs are normally off. The appropriate one will light and remain on while information is input in the Fill or Verify mode.

## Open/Close Handle.

This handle is OPEN to allow positioning of a paper tape; and in CLOSE position to lock the tape in place.

Figure 4.4 is a block diagram of a typical FADAC system, which illustrates how an input device (here the SDR), an output device, FADAC, and the Field Table are interconnected.

### 4.4 INPUT

General.
In most applications, paper tape is used to input information to FADAC. Either of two paper tape codes


Figure 4-4. Block Diagram of a Typical FADAC Sysiem
may be used, 8 -level Fieldata or 5 -level Teletype code. When FADAC is under program control these input codes are interpreted and stored as Octal, BCD, Alpha-5 or Alphao6, depending on the Read command sensed. Any input device which will generate Teletype or Fieldata paper tape codes may be used, whether or not paper tape is in fact the input vehicle. However, when FADAC is under external control, information can enter the computer in Octal format only.

Octal Input.
Recall that during octal input, control is with the input device, not with the computer itself. Therefore, in addition to the octal digits 0 through 7, FADAC will recognize seven special control characters. These characters are never stored in computer memory; rather each has a specific function. The blank, which can only be generated by a sprocket hole with no punches on paper tape, is ignored. (Paper tape codes for Octal, BCD digits and octal control characters are listed in Appendix C.)

### 4.4.1 OCTAL CONTROL CHARACTERS

Halt Code (8).
When this code is sensed, FADAC and the input device will halt. The I/O indicator light on the computer panel will turn off. If input is via the keyboard, the Keyboard Indicator light will also turn off.

Compute Code (9)
When this code is sensed, the Compute Indicator on the control panel will light, and computation will begin at the location specified by the 14 least significant bits in the L-register. (Note the L-register must be filled with the address of the first instruction to be executed prior to transmission of the Compute code.)

If the SDR is connected, computation will not start until the RUN/HALT switch is thrown into the RUN position.

## Fill Code (RIGHT/UP/ADD/+).

This code may be used to put FADAC back into the FILL mode after it has been placed in the VERIFY mode.

## Verify Code (LEFT/DROP/DOWN/-).

The VERIFY code causes the computer to compare the information following the VERIFY code with the information already stored in memory. If the data being ! not the same as that in memory, FADAC vil. welf and the input device and the ERROR indicat. .. flash.

Location Code (.).
This code causes the 14 least significant bits in the A-register to be transferred to the L-register, thus setting up the initial location where the information following is to be stored, or where computation or verification is to start. Following a LOCATION code, each FADAC-Octal word entered into or verified by the computer causes the location in the L-register to be increased by one. This code can be given anywhere in a sequence of input data. The contents of the Aregister are not changed by this control code.

## Clear Code (CLEAR).

This code is ignored by FADAC during octal input.

## Enter Code (ENTER).

This code will cause the contents of the A-register to be stored in the location specified by the 14 least
significant bits in the L-register, after which the contents of $L$ will be incremented by one. The contents of A remain unchanged.

## Blank Code.

This code is ignored on octal input.

### 4.4.2 PREPARATION OF PAPER TAPES

FADAC paper tape is black, unoiled photo-electric tape, fan-folded into $8-1 / 2$ inch folds. For teletype code, the tape is $11 / 16$ inch wide; for Fieldata, one inch wide. Characters are punched 10 per inch. Any paper tape punch device may be used which will generate the paper tape codes listed in Appendix C and Appendix D. Note that the DELETE code (all ones) is used to cover a tape code which is punched in error on a program tape.

## Octal Format.

All octal tapes should begin with a leader of approximately 2 folds of sprocket holes (17''), followed by a VERIFY code (-), followed by one fold of sprocket holes. The first data punched should be a 5-digit FADACaddress, followed by a Location code (.), followed by an 11 digit FADAC-octal word, followed by an Enter Code. (Note: On the Friden and Soroban paper tape punch units the Enter Code is represented by the Carriage Return.) To shorten the length of a program tape, all information to be stored sequentially should be punched sequentially. This will eliminate the need to punch a 5 -digit address for each entry. Each non-sequential entry MUST be preceded by the desired address followed by a Location code. The tape should end with a Halt code (8). If desired, however, the Halt code may be omitted and a 5 -digit address, followed by the Location code, followed by the Compute code (9), may be the last punches on a tape. (If this latter method is used, computation will
begin at the last location address specified, as soon as the SDR RUN/HALT switch is thrown into the RUN position). A leader of approximately two folds should follow the last character punched.

Non-Octal Tapes.
As stated, codes other than Octal are input under program control. The Read Command given will cause FADAC hardware to interpret and store the specified blocks of data in the desired format, beginning at the address specified in the L-register. The block size and type must be specified in all non-Octal Read Commands.

Tape Preparation - Non-Octal.
Non-octal tapes to be read into FADAC under program control are punched as an endless string of characters with no control characters, as such. Codes for alphanumeric input (Read Alpha-5 or Read Alpha-6) are listed in Appendix D, BCD in Appendix C.

## Caution.

The SDR does not stop on a character. Rather approximately three spaces are lost each time an SDR stop is executed either manually or under program control. Therefore, if a tape is to be input via this device, it is essential that a space of at least 5 sprocket holes be left between each segment of information to be input by each Read External Device command.

### 4.5 PREPARATION OF THE FADAC SYSTEM

## HOT Channel Select Switch.

This switch, located behind the Keyboard panel, is normally set for Hot 12 Channels. To change this setting, remove the 8 screws which fit the keyboard to
the control panel. The Hot Channel Select switch is marked for 16 channels (left hand position), 4 channels (middle position), and 12 channels (right hand position), and may be manually set to the desired position.

SDR Settings.
With the SDR connected to FADAC, SDR controls are set as follows:

POWER - ON
SIGNAL - ON
TAPE ADVANCE - OFF
FD/TT - Set for input code desired.
FILL/VERIFY - Setting immaterial.
AUXILIARY MEMORY - OFF
COMPUTE RUN/HALT
Halt position for input under external control. In RUN position for all FADAC operations under program control.

### 4.5.1 STARTING THE SYSTEM

Turn on all peripheral equipment first.
Throw FADAC power switch to ON.
Do not perform any FADAC operations until the Power Ready indicator is on. This will occur approximately 20 seconds after power is set ON.

### 4.6. INPUT THROUGH EXTERNAL CONTROL (INPUT IN OCTAL)

### 4.6.1 PAPER TAPE (SDR)

To Insert Tape.
Check settings of the FD/TT switch and the tape guides at the SDR read head. To adjust the tape guides, turn screw head at each tape guide clockwise as far as possible to increase size for 8 level tape, and counterclockwise to dectease size for 5 level tape. Pusition the OPEN/CLOSE handle to OPEN. Piace tape through the tape guide slit located on the left side of the tape cartridge holder, and through the tape guides at the read head. The sprocket holes of the tape being between the 3rd and 4th levels of same, the tape should be positioned such that space for the 1 st three levels is toward the inside wall of the SDR, Set the OPEN/CLOSE handle to CLOSE. (Check the tape guide settings by running a dummy tape through the SDR, as an improper setting may destroy a tape.)

## Operation.

Press RESET.
Check setting of SDR controls. (Section 4.5.) Throw RUN/HALT switch to HALT.

Press RESET button on FADAC control panel, and START Button on SDR. The FILL indicator on the SDR and I/O indicator on FADAC will light, and remain lit throughout the tape input operation.

## CAUTION.

Be sure the tape is positioned such that the Verify code is to the right of the read head.

## To Verify.

Reposition the tape such that the VERIFY code is to the left of the read head. Press RESET on FADAC, and START on the SDR. The SDR VERIFY indicator and I/O indicator on FADAC will light. If an error is
sensed, the ERROR indicator on FADAC will blink, and input will stop. Repeat tape input procedure described above and re-verify.

### 4.6.2 PAPER TAPE (MECHANICAL READER)

## To Insert Tape.

Raise the clamp armature on the right hand side of the reader (Figure 4.5). The sprocket holes of the tape being between the 3rd and 4th levels of same, position the tape such that the space for the first three levels is toward the computer. Insure that all information to be input is to the left of the read head. Place the tape under the read head, engage the tape sprocket holes with the reader sprockets, and close the armature clamp. Turn the sprocket knob on the upper right side of the reader a few times to insure correct engagement. If the tape does not move freely, verify that the sprocket holes have made proper contact with the reader sprocket and that the tape is properly threaded between the read and the sprocket (Figure 4.5).

## Operation.

Press RESET. Check setting of SDR controls. (See Section 4.5.) Throw RUN/HALT switch to HALT.

Press RESET and RECALL on FADAC, and START on the SDR. The FILL indicator on the SDR and the I/O indicator on FADAC will light, and remain lit throughout the tape input operation.

## CAUTION.

Be sure the tape is positioned such that the VERIFY code is to the right of the read head.

To Verify.
Press RESET. Reposition the tape such that the


Figure 4.5. Mechanical Reader

VERIFY code is to the left of the read head. Press RESET and RECALL on FADAC and START on the SDR. The VERIFY indicator on the MLU and the I/O indicator on FADAC will light. If an error is sensed, the ERROR indicator on FADAC will blink, and input will stop. Repeat the tape input procedure described above and re-verify.

Keyboard Input.
Check setting of SDR controls. (See Section 4.5.) Throw RUN/HALT switch to HALT.

Press RESET and RECEIVE (or RESET and TRIG) on FADAC, and START on the SDR. The KEYBOARD and I/O indicators on FADAC and the FILL indicator on the SDR will light and remain on as long as FADAC remains in the Keyboard Input mode.

Type the 5-octal digit address to which the first entry is to be sent, followed by a LOCATION code (.), followed by the 11 octal digit word to be stored, followed by the ENTER code (ENTER key). If entries are to be sent to sequential locations, only the first assigned location followed by the LOCATION code must be typed. All subsequent entries may be 11 octal digits followed by the ENTER key. Each non-sequential entry MUST be preceded by the 5 -digit assigned location followed by the LOCATION CODE.

Example:

| Line 1 | 10125.10126250326 E |
| :--- | :--- |
| Line 2 | 10126.30127010327 E |
| Line 3 | 10127.10130511326 E |
| Line 4 | 10130.14600150126 E |
| Line 5 | 30126.00000372000 E |

Since the entries in lines 1 through 4 are sequential, only the first location field, 10125, must be typed. The
contents of the L-register are automatically incremented by one each time an enter code (E) is sensed. To break the sequence, a new location field followed by a location code is entered, as in line 5 (30126). Note that this location in Memory III is identified by a 1 in bit position 18 of the L-register as follows:


The format of the typed entries shown in the Example is precisely the same as that used to punch an octal paper tape.

## To Verify Keyboard Entries.

Press LEFT DROP key on the keyboard and retype all entries to be verified. If an error is detected, the ERROR indicator will blink. Re-enter the data, as described above and re-verify.

### 4.7 INPUT THROUGH PROGRAM CONTROL

(Note: Input may be in BCD, Alpha-5 or Alpha-6 format) With the SDR Connected.

Check settings on the SDR. (See Section 4.5.) RUN/HALT switch must be in the RUN position. Check all peripheral equipment connections.

## If Via SDR or Mechanical Reader.

Position tape properly in the input device and initiate computation. Tape will be input automatically when any Read command is sensed. The FILL indicator
on the SDR and I/O indicator on FADAC will light, and remain on until the specified block size is entered.

## If $\mathrm{Via}_{\text {K }}$ Keyboard.

Initiate computation. The Keyboard and I/O indicators on FADAC will light, and FADAC will halt to await a keyboard entry when any Read Keyboard command is sensed.

To Initiate Computation.
With the SDR (or SDR Simulator) Settings as specified in Section 4.5, the SDR RUN/HALT switch in the run position, and all peripheral equipment connected and operational, press the FADAC control button which is programmed to initiate computation.

## APPENDIX A

Powers of Two Table

```
            2N}\quadN\quad2-
            1 0 1.0
            2 1 0.5
            4 2 0.25
            8 0.125
            16 4 0.062 5
            32 5 0.031 25
            64 6 0.015 625
            128 7 0.007 812 5
            256 8 0.003 905 25
            512 ' 9 0.001 953 125
            1 024 10 0.000 976 562 5
            2}0481110.000488\quad28112
            4 096}12120.000 244 140 625
            8 192 13 0.000 122 070 312 5
            16}38414\quad0.000061 035 156 25
            32768}1500.000 030 517 578 125
            65 536 16 0.000 015 258 789 062 5
            131 072 17 0.000 007 629 394 531 25
            262 144 18 0.000 003 814 697 265 625
            524 288 19 0.000 001 907 348 632 812 5
            1 048 576 20 0.000 000 953 674 316 406 25
            2 097 15, 21 21 0.000 000 476 837 158 203 125
            4}19
            8 388 608 23 0.000 000 119 209 289 550 781 25
                    16}777216 24 0.000 000 059 604 644 775 390 625,
                    33 554 432 25 0.000 000 029 802 322 387 695 312 5
                            67 108 864 26 0.000 000 014 901 161 193 847 656 25
                            134 217 728 27 0.000 000 007 450 580 596 923 828 125
                    268435456 28 0.000 000 003 725 290 298 461 914 062 5
                    536 870 912 29 0.000 000 001 862 645 149 230 957 031 25
1 073 741 824 30 0.000 000 000 931 322 574 615 478 515 625
2 147483 648 31 0.000 000 000 465 661 287 307 739 257 812 5
```


## APPENDIX B

OCTAL-DECIMAL INTEGER CONVERSION TABLE

| 0000 | 0000 |
| :---: | :---: |
| 10 | 10 |
| 0777 | 0511 |
| 10 (1al) | (Decimal) |


| 0000 | 0000 | 0001 | 0002 | 0003 | 0004 | 0005 | 0006 | 0007 |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0010 | 0008 | 0009 | 0010 | 0011 | 0012 | 0013 | 0014 | 0015 |
| 0020 | 0016 | 0017 | 0018 | 0019 | 0020 | 0021 | 0022 | 0023 |
| 0030 | 0024 | 0025 | 0026 | 0027 | 0028 | 0029 | 0030 | 0031 |
| 0040 | 0032 | 0033 | 0034 | 0035 | 0036 | 0037 | 0038 | 0039 |
| 0050 | 0040 | 0041 | 0642 | 0043 | 0044 | 0045 | 0046 | 0047 |
| 0060 | 0048 | 0049 | 0050 | 0051 | 0052 | 0053 | 0054 | 0055 |
| 0070 | 0056 | 0057 | 0058 | 0059 | 0060 | 0061 | 0062 | 0063 |
|  |  |  |  |  |  |  |  |  |
| 0100 | 0064 | 0065 | 0066 | 0067 | 0068 | 0069 | 0070 | 0071 |
| 0110 | 0072 | 0073 | 0074 | 0075 | 0078 | 0077 | 0078 | 0079 |
| 0120 | 0080 | 0081 | 0082 | 0083 | 0084 | 0085 | 0086 | 0087 |
| 0130 | 0088 | 0089 | 0090 | 0091 | 0092 | 0093 | 0094 | 0095 |
| 0140 | 0096 | 0097 | 0098 | 0099 | 0100 | 0101 | 0102 | 0103 |
| 0150 | 0104 | 0105 | 0106 | 0107 | 0108 | 0109 | 0110 | 0111 |
| 0160 | 0112 | 0113 | 0114 | 0115 | 0116 | 0117 | 0118 | 0119 |
| 0170 | 0120 | 0121 | 0122 | 0123 | 0124 | 0125 | 0126 | 0127 |
| 0200 | 0128 | 0129 | 0130 | 0131 | 0132 | 0133 | 0134 | 0135 |
| 0210 | 0136 | 0137 | 0138 | 0139 | 0140 | 0141 | 0142 | 0143 |
| 0220 | 0144 | 0145 | 0146 | 0147 | 0148 | 0149 | 0150 | 0151 |
| 0230 | 0152 | 0153 | 0154 | 0155 | 0156 | 0157 | 0158 | 0159 |
| 0240 | 0160 | 0161 | 0162 | 0163 | 0164 | 0165 | 0166 | 0167 |
| 0250 | 0168 | 0169 | 0170 | 0171 | 0172 | 0173 | 0174 | 0175 |
| 0260 | 0176 | 0177 | 0178 | 0179 | 0180 | 0181 | 0182 | 0183 |
| 0270 | 0184 | 0185 | 0186 | 0187 | 0188 | 0189 | 0190 | 0191 |
| 0300 | 0192 | 0193 | 0194 | 0195 | 0196 | 0197 | 0198 | 0199 |
| 0310 | 0200 | 0201 | 0202 | 0203 | 0204 | 0205 | 0206 | 0207 |
| 0320 | 0208 | 0209 | 0210 | 0211 | 0212 | 0213 | 0214 | 0215 |
| 0330 | 0216 | 0217 | 0218 | 0219 | 0220 | 0221 | 0222 | 0223 |
| 0340 | 0224 | 0225 | 0226 | 0227 | 0228 | 0229 | 0230 | 0231 |
| 0350 | 0232 | 0233 | 0234 | 0235 | 0236 | 0237 | 0238 | 0239 |
| 0360 | 0240 | 0241 | 0242 | 0243 | 0244 | 0245 | 0246 | 0247 |
| 0370 | 0248 | 0249 | 0250 | 0251 | 0252 | 0253 | 0254 | 0255 |


|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | $\mathbf{7}$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0400 | 0256 | 0257 | 0258 | 0259 | 0260 | 0261 | 0262 | 0263 |
| 0410 | 0264 | 0265 | 0268 | 0267 | 0268 | 0269 | 0270 | 0271 |
| 0420 | 0272 | 0273 | 0274 | 0275 | 0276 | 0277 | 0278 | 0279 |
| 0430 | 0280 | 0281 | 0282 | 0283 | 0284 | 0285 | 0286 | 0287 |
| 0440 | 0288 | 0289 | 0290 | 0291 | 0292 | 0293 | 0294 | 0295 |
| 0450 | 0296 | 0297 | 0298 | 0290 | 0300 | 0301 | 0302 | 0303 |
| 0460 | 0304 | 0305 | 0306 | 0307 | 0308 | 0309 | 0310 | 0311 |
| 0470 | 0312 | 0313 | 0314 | 0315 | 0316 | 0317 | 0318 | 0319 |
|  |  |  |  |  |  |  |  |  |
| 0500 | 0320 | 0321 | 0322 | 0323 | 0324 | 0325 | 0326 | 0327 |
| 0510 | 0328 | 0329 | 0330 | 0331 | 0332 | 0333 | 0334 | 0335 |
| 0520 | 0336 | 0337 | 0338 | 0339 | 0340 | 0341 | 0342 | 0343 |
| 0530 | 0344 | 0345 | 0346 | 0347 | 0348 | 0349 | 0350 | 0351 |
| 0540 | 0352 | 0353 | 0354 | 0355 | 0356 | 0357 | 0358 | 0359 |
| 0550 | 0360 | 0361 | 0362 | 0363 | 0364 | 0365 | 0366 | 0367 |
| 0560 | 0368 | 0369 | 0370 | 0371 | 0372 | 0373 | 0374 | 0375 |
| 0570 | 0376 | 0377 | 0378 | 0379 | 0380 | 0381 | 0382 | 0383 |
| 0600 | 0384 | 0385 | 0386 | 0387 | 0388 | 6 | 19 | 0390 |
| 039 |  |  |  |  |  |  |  |  |$|$


| 1000 | 0512 |
| :---: | :---: |
| 10 | 10 |
| 1777 | 1023 |
| (Octal) | (Deciseal) |


|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1000 | 0512 | 0513 | 0514 | 0515 | 0516 | 0517 | 0518 | 0519 |
| 1010 | 0520 | 0521 | 0522 | 0523 | 0524 | 0525 | 0526 | 0527 |
| 1020 | 0528 | 0529 | 0530 | 0531 | 0532 | 0533 | 0534 | 0535 |
| 1030 | 0536 | 0537 | 0538 | 0539 | 0540 | 0541 | 0542 | 0543 |
| 1040 | 0544 | 0545 | 0546 | 0547 | 0548 | 0540 | 0550 | 0551 |
| 1050 | 0552 | 0553 | 0554 | 0555 | 0556 | 0557 | 0558 | 0559 |
| 1060 | 0560 | 0561 | 0562 | 0563 | 0564 | 0565 | 0566 | 0567 |
| 1070 | 0568 | 0569 | 0570 | 0571 | 0572 | 0573 | 0574 | 0575 |
| 1100 | 0576 | 0577 | 0578 | 0579 | 0580 | 0581 | 0582 | 0583 |
| 1110 | 0584 | 0585 | 0586 | 0587 | 0588 | 0588 | 0590 | 0591 |
| 1120 | 0592 | 0593 | 0594 | 0595 | 0596 | 0597 | 0598 | 0599 |
| 1130 | 0600 | 0601 | 0602 | 0603 | 0604 | 0605 | 0606 | 0607 |
| 1140 | 0608 | 0609 | 0610 | 0611 | 0617 | 0813 | 0614 | 0615 |
| 1150 | 0516 | 0617 | 0618 | 0619 | 0620 | 0621 | 0622 | 0623 |
| 1160 | 0624 | 0625 | 0626 | 0527 | 0628 | 0629 | 0630 | 0631 |
| 1170 | 0632 | 0633 | 0634 | 0635 | 0636 | 0637 | 0638 | 0639 |
| 1200 | 0640 | 0641 | 0642 | 0543 | 0644 | 0645 | 0646 | 0647 |
| 1210 | 0648 | 0649 | 0650 | 0651 | 0652 | 0653 | 0654 | 0655 |
| 1220 | 0656 | 0657 | 0658 | 0659 | 0660 | 0661 | 0662 | 0663 |
| 1230 | 0664 | 0665 | 0668 | 0667 | 0568 | 0669 | 0670 | $0671 \mid$ |
| 1240 | 0672 | 0673 | 0674 | 0675 | 0676 | 0677 | 0678 | 0679 |
| 1250 | 0680 | 0681 | 0882 | 0683 | 0584 | 0685 | 0688 | 0687 |
| 1260 | 0688 | 0689 | 0690 | 0691 | 0692 | 0693 | 0694 | 0695 |
| 1270 | 0696 | 0697 | 0698 | 0699 | 0700 | 0701 | 0702 | 0703 |
| 1300 |  | 0704 | 0705 | 0706 | 0707 | 0709 | 0708 | 0710 |
| 071 |  |  |  |  |  |  |  |  |
| 1310 | 0712 | 0713 | 0714 | 0715 | 0716 | 0717 | 0718 | 0719 |
| 1320 | 0720 | 0721 | 0722 | 0723 | 0724 | 0725 | 0726 | 0727 |
| 1330 | 0728 | 0729 | 0730 | 0731 | 0732 | 0733 | 0734 | 0735 |
| 1340 | 0736 | 0737 | 0738 | 0739 | 0740 | 0741 | 0742 | 0743 |
| 1350 | 0744 | 0745 | 0740 | 0747 | 0748 | 0749 | 0750 | 0751 |
| 1360 | 0752 | 0753 | 0754 | 0755 | 0758 | 0757 | 0758 | 0759 |
| 1370 | 0760 | 0761 | 0762 | 0763 | 0764 | 0765 | 0766 | 0767 |


|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1400 | 0768 | 0769 | 0770 | 0771 | 0772 | 0773 | 0774 | 0775 |
| 1410 | 0776 | 0777 | 0778 | 0779 | 0780 | 0781 | 0782 | 0783 |
| 1420 | 0784 | 0785 | 0786 | 0787 | 0788 | 0789 | 0790 | 0791 |
| 1430 | 0792 | 0793 | 107:1 | c795 | 0796 | 0797 | 0798 | 0799 |
| 1443 | 0809 | 0801 | 6802 | 0803 | 0804 | 0805 | 0806 | 0807 |
| 1450 | 0808 | 0804 | 0810 | 9811 | 0812 | 0813 | 0814 | 0815 |
| 1460 | 0816 | 0817 | 0818 | 0819 | 0820 | 0821 | 0822 | 0823 |
| 1470 | 0824 | 0825 | 0826 | 0827 | 0828 | 0829 | 0830 | 0831 |
| 1500 | 0832 | 0833 | 0834 | 0835 | 0836 | 0837 | 0838 | 0839 |
| 1510 | 0840 | 0841 | 0842 | 0843 | 0844 | 0845 | 0846 | 0847 |
| 1529 | 0848 | 0849 | 0850 | 0851 | 0852 | 0853 | 0854 | 085s |
| 1530 | 0856 | 0857 | 0858 | 0859 | 0860 | 0861 | 0862 | 0863 |
| 1540 | 0864 | 0865 | 0866 | 0867 | 0868 | 0869 | 0870 | 0871 |
| 1550 | 0872 | 0873 | OB74 | 0875 | 0876 | 0877 | 0878 | 0879 |
| 1560 | 0880 | 0881 | 0882 | 0883 | 0884 | 0885 | 0886 | 0887 |
| 1570 | 0888 | 0889 | 0890 | 0891 | 0892 | 0893 | OB94 | 0895 |
| 1600 | 0896 | 0897 | 0898 | 0899 | 0900 | 0901 | 0902 | 0903 |
| 1610 | 0904 | 0905 | 0906 | 0907 | 0908 | 0909 | 0910 | 0911 |
| 1620 | 0912 | 0913 | 0914 | 0915 | 0916 | 0917 | 0918 | 0919 |
| 1630 | 0920 | 0921 | 0922 | 0923 | 0924 | 0925 | 0926 | 0927 |
| 1640 | 0925 | 0929 | 0530 | 0931 | 0932 | 0933 | 0934 | 0935 |
| 1650 | 0936 | 0937 | 0938 | 0939 | 0940 | 0941 | 0942 | 0943 |
| 1660 | 0944 | 0945 | 0946 | 0947 | 0948 | 0949 | 0950 | 0951 |
| 1870 | 693 | 9953 | 0954 | 0955 | 0956 | 0957 | 0958 | 0959 |
| : 170 | O2.30 | $0 \leq 61$ | 0362 | 0963 | 0564 | 0965 | 0962 | 0967 |
| 1710 | 0968 | 0969 | 0970 | 0971 | 0972 | 0973 | 0974 | 0975 |
| 1720 | 0976 | 0977 | 0978 | 0978 | 0980 | 0981 | 0982 | 0983 |
| 1730 | 0984 | 0985 | 0988 | 0987 | 0988 | 0989 | 0990 | 0991 |
| 1740 | 0992 | 0993 | 0994 | 0995 | 0998 | 0997 | 0998 | 0999 |
| 1750 | 1000 | 1001 | 1002 | 1003 | 1004 | 1005 | 1006 | 1007 |
| 1760 | 1008 | 1009 | 1010 | 1011 | 1012 | 1013 | 1014 | 1015 |
| 1770 | 1016 | 1217 | 1018 | 1019 | 1020 | 1021 | 1022 | 1023 |

OCTAL-DECIMALINTEGER CONVERSION IABLE (Cont'd)

| 2000 | 1024 |
| :---: | :---: |
| 10 | 10 |
| 2777 | 1539 |
| (Octal) | (Decimal) |


|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 2000 | 1024 | 1025 | 1026 | 1027 | 1028 | 1029 | 1030 | 1031 |
| 2010 | 1032 | 1033 | 1034 | 1035 | 1036 | 1037 | 1038 | 1039 |
| 2020 | 1040 | 1041 | 1042 | 1043 | 1044 | 1045 | 1046 | 1047 |
| 2030 | 1048 | 1049 | 1050 | 1051 | 1052 | 1053 | 1054 | 1055 |
| 2040 | 1056 | 1057 | 1058 | 1059 | 1060 | 1061 | 1062 | 1063 |
| 2050 | 1064 | 1065 | 1066 | 1067 | 1068 | 1069 | 1070 | 1071 |
| 2060 | 1072 | 1073 | 1074 | 1075 | 1076 | 1077 | 1078 | 1079 |
| 2070 | 1080 | 1081 | 1082 | 1083 | 1084 | 1085 | 1086 | 1087 |
|  |  |  |  |  |  |  |  |  |
| $2: 50$ | 1088 | 1089 | 1090 | 1091 | 1092 | 1093 | 1094 | 1095 |
| 2110 | 1096 | 1097 | 1098 | 1099 | 1100 | 1101 | 1102 | 1103 |
| 2120 | 1104 | 1105 | 1106 | 1107 | 1108 | 1109 | 1110 | 1111 |
| 2130 | 1112 | 1113 | 1114 | 1115 | 1116 | 1117 | 1118 | 1119 |
| 2140 | 1120 | 1121 | 1122 | 1123 | 1124 | 1125 | 1126 | 1127 |
| 2150 | 1128 | 1129 | 1130 | 1131 | 1132 | 1133 | 1134 | 1135 |
| 2160 | 1136 | 1137 | 1138 | 1139 | 1140 | 1141 | 1142 | 1143 |
| 2170 | 1144 | 1145 | 1146 | 1147 | 1148 | 1149 | 1150 | 1151 |
| 2200 | 1152 | 1153 | 1154 | 1155 | 1158 | 1157 | 1158 | 1159 |
| 2210 | 1160 | 1161 | 1162 | 1163 | 1164 | 1165 | 1166 | 1167 |
| 2220 | 1168 | 1169 | 1170 | 1171 | 1172 | 1173 | 1174 | 1175 |
| 2239 | 1176 | 1177 | 1178 | 1179 | 1180 | 1181 | 1182 | 1183 |
| 2240 | 1184 | 1185 | 1186 | 1187 | 1188 | 1189 | 1190 | 1191 |
| 2250 | 1192 | 1193 | 1194 | 1195 | 1196 | 1197 | 1198 | 1199 |
| 2260 | 1200 | 1201 | 1202 | 1203 | 1204 | 1205 | 1208 | 1207 |
| 2270 | 1208 | 1209 | 1210 | 1211 | 1212 | 1213 | 1214 | 1215 |
| 2300 | 1216 | 1217 | 1218 | 1219 | 1220 | 1221 | 1222 | 1223 |
| 2310 | 1224 | 1225 | 1226 | 1227 | 1228 | 1229 | 1230 | 1231 |
| 2320 | 1232 | 1233 | 1234 | 1235 | 1238 | 1237 | 1238 | 1239 |
| 2330 | 1240 | 1241 | 1242 | 1243 | 1244 | 1245 | 1246 | 1247 |
| 2340 | 1248 | 1249 | 1250 | 1251 | 1252 | 1253 | 1254 | 1255 |
| 2350 | 1256 | 1257 | 1258 | 1259 | 1260 | 1261 | 1262 | 1263 |
| 2360 | 1264 | 1265 | 1268 | 1267 | 1268 | 1269 | 1270 | 1271 |
| 2370 | 1272 | 1273 | 1274 | 1275 | 1278 | 1277 | 1278 | 1279 |


|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 2400 | 1280 | 1281 | 1282 | 1283 | 1284 | 1285 | 1286 | 1287 |
| 2410 | 1288 | 1289 | 1290 | 1291 | 1292 | 1293 | 1294 | 1295 |
| 2420 | 1296 | 1297 | 1298 | 1299 | 1300 | 1301 | 1302 | 1303 |
| 2430 | 1304 | 1305 | 1306 | 1307 | 1308 | 1309 | 1310 | 1311 |
| 2440 | 1312 | 1313 | 1314 | 1315 | 1316 | 1317 | 1318 | 1319 |
| 2450 | 1320 | 1321 | 1322 | 1323 | 1324 | 1325 | 1326 | 1327 |
| 2460 | 1328 | 1329 | 1330 | 1331 | 1332 | 1333 | 1334 | 1335 |
| 2470 | 1336 | 1337 | 1338 | 1339 | 1340 | 1341 | 1342 | 1343 |
| 2500 | 1344 | 1345 | 1348 | 1347 | 1348 | 1349 | 1350 | 1351 |
| 2510 | 1352 | 1353 | 1354 | 1355 | 1355 | 1357 | 1358 | 1359 |
| 2520 | 1360 | 1361 | 1362 | 1363 | 1364 | 1365 | 1366 | 1367 |
| 2530 | 1368 | 1369 | 1370 | 1371 | 1372 | 1373 | 1374 | 1375 |
| 2540 | 1376 | 1377 | 1378 | 1379 | 1380 | 1381 | 1382 | 1383 |
| 2550 | 1344 | 1385 | 1366 | 1387 | 1388 | 1389 | 1390 | 1391 |
| 2560 | 1392 | 1393 | 1394 | 1395 | 1396 | 1397 | 1398 | 1399 |
| 2570 | 1400 | 1401 | 1402 | 1403 | 1404 | 1405 | 1406 | 1407 |
| 2600 | 1408 | 1409 | 1410 | 1411 | 1412 | 1413 | 1414 | 1415 |
| 2610 | 1416 | 1417 | 1418 | 1419 | 1420 | 1421 | 1422 | 1423 |
| 2620 | 1424 | 1425 | 1426 | 1427 | 1428 | 1429 | 1430 | 1431 |
| 2630 | 1432 | 1433 | 1434 | 1435 | 1436 | 1437 | 1438 | 1439 |
| 2640 | 1440 | 1441 | 1442 | 1443 | 1444 | 1445 | 1446 | 1447 |
| 2650 | 1448 | 1449 | 1450 | 1451 | 1452 | 1453 | 1454 | 1455 |
| 2660 | 1456 | 1457 | 1458 | 1459 | 1460 | 1461 | 1462 | 1463 |
| 2670 | 1464 | 1465 | 1466 | 1467 | 1468 | 1469 | 1470 | 1471 |
| 2700 | 1472 | 1473 | 1474 | 1475 | 1476 | 1477 | 1478 | 1479 |
| 2710 | 1480 | 1481 | 1482 | 1483 | 1484 | 1485 | 1436 | 1487 |
| 2720 | 1488 | 1489 | 1490 | 1491 | 1492 | 1493 | 1494 | 1495 |
| 2730 | 1496 | 1497 | 1498 | 1499 | 1500 | 1501 | 1502 | 1503 |
| 2740 | 1504 | 1505 | 1506 | 1507 | 1508 | 1509 | 1510 | 1511 |
| 2750 | 1512 | 1513 | 1514 | 1515 | 1518 | 1517 | 1518 | 1519 |
| 2760 | 1520 | 1521 | 1522 | 1523 | 1524 | 1525 | 1526 | 1527 |
| 2770 | 1528 | 1529 | 1530 | 1531 | 1532 | 1533 | 1534 | 1535 |


| 3000 | 1536 |
| :---: | :---: |
| 10 | 10 |
| 3777 | 2047 |
| Octal) | (Decimal) |


|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 3000 | 1536 | 1537 | 1538 | 1539 | 1540 | 1541 | 1542 | 1543 |
| 3010 | 1544 | 1545 | 1548 | 1547 | 1548 | 1549 | 1550 | 1551 |
| 3020 | 1532 | 1553 | 1554 | 1555 | 1556 | 1557 | 1558 | 1559 |
| 3030 | 1580 | 1561 | 1582 | 1583 | 1564 | 1565 | 1568 | 1587 |
| 3040 | 1568 | 1569 | 1570 | 1571 | 1572 | 1573 | 1574 | 1575 |
| 3050 | 1578 | 1577 | 1578 | 1579 | 1580 | 1581 | 1582 | 1583 |
| 3060 | 1584 | 1585 | 1586 | 1587 | 1588 | 1589 | 1590 | 1591 |
| 3070 | 1592 | 1593 | 1594 | 1595 | 1596 | 1597 | 1598 | 1589 |
| 3100 | 1600 | 1801 | 1602 | 1803 | 1804 | 1805 | 1608 | 1607 |
| 3110 | 1608 | 1609 | 1610 | 1611 | 1612 | 1813 | 1614 | 1615 |
| 3120 | 1618 | 1617 | 1818 | 1619 | 1620 | 1621 | 1822 | 1623 |
| 3130 | 1624 | 1825 | 1625 | 1627 | 1628 | 1629 | 1630 | 31 |
| 3140 | 1632 | 1633 | 1634 | 1635 | 1636 | 1637 | 1838 | 39 |
| 3150 | 1640 | 1641 | 1842 | 1643 | 1644 | 1645 | 1646 | 647 |
| 3160 | 1648 | 1649 | 1650 | 1651 | 1652 | 1653 | 16 | 655 |
| 3170 | 1656 | 1657 | 1658 | 1859 | 1680 | 1681 | 1662 | 1663 |
| 3200 | 1664 | 1665 | 1666 | 1667 | 1668 | 1669 | 1670 | 1671 |
| 3210 | 1672 | 1673 | 1674 | 1675 | 1676 | 1677 | 1678 | 1679 |
| 3220 | 1680 | 1681 | 1682 | 1683 | 1684 | 1685 | 1686 | 1687 |
| 3230 | 1688 | 1689 | 1690 | 1691 | 1692 | 1693 | 1694 | 1695 |
| 3240 | 1696 | 1697 | 1698 | 1699 | 1700 | 1701 | 1702 | 1703 |
| 3250 | 1704 | 1705 | 1706 | 1707 | 1708 | 1709 | 1710 | 1711 |
| 3260 | 1712 | 1713 | 1714 | 1715 | 1716 | 1717 | 1718 | 1719 |
| 3270 | 1720 | 1721 | 1722 | 1723 | 1724 | 1725 | 1726 |  |
| 3300 | 1728 | 1729 | 1730 | 1731 | 1732 | 1733 | 1734 | 1735 |
| 3310 | 1736 | 1737 | 1738 | 1739 | 1740 | 1741 | 1742 | 1743 |
| 3320 | 1744 | 1745 | 1746 | 1767 | 1748 | 1749 | 1750 | 1751 |
| 3330 | 1752 | 1753 | 1754 | 1755 | 1736 | 1759 | 1758 | 1759 |
| 13340 | 1760 | 1761 | 1762 | 1753 | 1764 | 1765 | 1765 | 1767 |
| 13300 | 178 R | 1769 | 1770 | 1771 | 1772 | 1773 | 1774 | 1:79 |
| 3360 | 1775 | 1777 | 1778 | 1779 | 1780 | 1781 | 1782 | 1783 |
| 3370 | 1784 | 1785 | 1786 | 178 | 1788 | 1789 | 1790 | 1791 |


|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 3400 | 1792 | 1793 | 1794 | 1795 | 1796 | 1797 | 1798 |  |
| 3410 | 1800 | 1801 | 1802 | 1803 | 1804 | 1805 | 1808 | 1807 |
| 3420 | 1808 | 1809 | 1810 | 1811 | 1812 | 1813 | 1814 | 1815 |
| 3430 | 1816 | 1817 | 1818 | 1819 | 1820 | 1821 | 1822 | 1823 |
| 3440 | 1824 | 1825 | 1826 | 1827 | 1828 | 1829 | 1830 | 1831 |
| 3450 | 1832 | 1833 | 1834 | 1835 | 1836 | 1837 | 1838 | 1839 |
| 3460 | 1840 | 1841 | 1842 | 1843 | 1844 | 1845 | 1848 | 1847 |
| 3470 | 1848 | 1819 | 1850 | 1851 | 1852 | 1853 | 1854 | 1855 |
| 3500 | 1858 | 1857 | 1858 | 1859 | 1860 | 1861 | 1862 | 1883 |
| 3510 | 1864 | 1855 | 1886 | 1867 | 1868 | 1869 | 1870 | 1871 |
| 3520 | 1872 | 1873 | 1874 | 1875 | 1878 | 1877 | 1878 | 1879 |
| 3530 | 1880 | 1881 | 1882 | 1883 | 1884 | 1885 | 1886 | 1887 |
| 3540 | 1888 | 1889 | 1890 | 1891 | 1892 | 1893 | 1894 | 1895 |
| 3550 | 1896 | 1897 | 1898 | 1899 | 1900 | 1901 | 1902 | 1903 |
| 3560 | 1904 | 1905 | 1906 | 1907 | 1908 | 1909 | 1910 | 1911 |
| 3570 | 1912 | 1913 | 1914 | 1915 | 1916 | 1917 | 1918 | 1919 |
| 3600 | 1920 | 1921 | 1922 | 1923 | 1924 | 1925 | 1926 | 1927 |
| 3610 | 1928 | 1929 | 1930 | 1931 | 1932 | 1933 | 1934 | 1835 |
| 3620 | 1936 | 1937 | 1938 | 1939 | 1940 | 1941 | 1842 | 1943 |
| 3630 | 1944 | 1945 | 1946 | 1947 | 1948 | 1949 | 1950 | 1951 |
| 3640 | 1952 | 1953 | 1954 | 1955 | 1956 | 1957 | 1958 | 1959 |
| 3650 | 1960 | 1961 | 1962 | 1963 | 1964 | 1965 | 1966 | 1967 |
| 3660 | 1968 | 1969 | 1970 | 1971 | 1972 | 1973 | 1974 | 1975 |
| 3670 | 1976 | 1977 | 1978 | 1979 | 1980 | 1981 | 1982 | 1983 |
| 3700 | 1984 | 1985 | 1986 | 1987 | 1988 | 1989 | 1990 | 1991 |
| 3710 | 1992 | 1993 | 1994 | 1995 | 1996 | 1997 | 1998 | 1999 |
| 3720 | 2000 | 2001 | 2002 | 2003 | 2004 | 2005 | 2006 | 2007 |
| 3730 | 2008 | 2009 | 2010 | 201: | 2012 | 2013 | 2014 | 2015 |
| 3742 | 2016 | 2017 | 2018 | 2019 | 2020 | 2021 | 2022 | $2 \mathrm{az3}$ |
| 3"50 | 2024 | 2025 | 2028 | 2027 | 2028 | 2029 | 2030 | 2031 |
| 3760 | 2032 | 2033 | 2034 | 2035 | 2038 | 2037 | 2038 | 2039 |
| 3770 | 2040 | 2041 | 2042 | 2043 | 2044 | 2045 | 2046 | 2047 |

## OCTAL-DECIMAL INTEGER CONVERSION TABI.E (Cont'd)

| 4000 | 2048 |
| :---: | :---: |
| 90 | 10 |
| 479 | 2559 |
| Ocial | Decimol |


|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 4000 | 2048 | 2049 | 2050 | 2051 | 2052 | 2053 | 2054 | 2055 |
| 4010 | 2056 | 2057 | 2058 | 2059 | 2060 | 2061 | 2062 | 2063 |
| 4020 | 2064 | 2065 | 2066 | 2067 | 2068 | 2069 | 2070 | 2071 |
| 4030 | 2072 | 2073 | $20: 4$ | 2075 | 2076 | 2077 | 207A | 2079 |
| 4040 | 2080 | 2081 | 2082 | 2083 | 2084 | 2085 | 2086 | 2087 |
| 4050 | 208 B | 2089 | 2090 | 2091 | 2092 | 2093 | 2094 | 2095 |
| 4060 | 2996 | 2097 | 2098 | 2099 | 2100 | 2101 | 2102 | 2103 |
| 4070 | 2104 | 2105 | 2106 | 2107 | 2108 | 2109 | 2110 | 2111 |
| 4100 | 2112 | 2113 | 2114 | 2115 | 2116 | 2117 | 2118 | 2119 |
| 4110 | 2120 | 2121 | 2122 | 2123 | 2124 | 2125 | 2126 | 2127 |
| 4120 | 2128 | 2129 | 2130 | 2131 | 2132 | 2133 | 2134 | $213 s$ |
| 4130 | 2136 | 2137 | 2138 | 2139 | 2140 | 2141 | 2142 | 2143 |
| 4140 | 2144 | 2145 | 2146 | 2147 | 2148 | 2149 | 2150 | 2151 |
| 4150 | 2152 | 2153 | 2154 | 2155 | 2156 | 2157 | 2158 | 2159 |
| 4160 | 2160 | 2161 | 2162 | 2163 | 2164 | 2165 | 2166 | 2167 |
| 4170 | 2168 | 2169 | 2170 | 2171 | 2172 | 2173 | 2174 | 2175 |
| 4200 | 2176 | 2177 | 2178 | 2179 | 2180 | 2181 | 2182 | 2183 |
| 4210 | 2184 | 2185 | 2186 | 2187 | 2188 | 2189 | 2190 | 2191 |
| 42201 | 2192 | 2193 | 2194 | 2195 | 2196 | $219 \%$ | 2198 | 2199 |
| 4230 | 2200 | 2201 | 2202 | 2203 | 2204 | 2205 | 2206 | 2207 |
| 4240 | 2208 | 2209 | 2210 | 2211 | 2212 | 2213 | 2214 | 2215 |
| 4250 | 2216 | 2217 | 2218 | 2219 | 2220 | 2221 | 2222 | 2223 |
| 4250 | 2224 | 2225 | 2226 | 2227 | 2228 | 2229 | 2230 | 2231 |
| 4.? 0 | 2232 | 2233 | 2234 | 2235 | 2236 | 2237 | 2238 | 2239 |
| 4300 | 2240 | 2241 | 2242 | 2243 | 2244 | 2245 | 2246 | 2247 |
| 4310 | 224 H | 2249 | 2250 | 2251 | 2252 | 225] | 2254 | 2255 |
| 4320 | 2256 | 2257 | 2258 | 2259 | 2260 | 2251 | 2262 | 2263 |
| 4330 | 2264 | 2265 | 2266 | 2267 | 2258 | 2269 | 2270 | 2271 |
| 4340 | 2272 | 2273 | 2274 | 2274 | 2276 | 2277 | 2274 | 2279 |
| 4350 | 2280 | 22R1 | 2282 | 2283 | 2284 | 224, | 22H6 | 228? |
| 4360 | 228 A | 22 Hy | 2291) | $22 \pm 1$ | 2242 | 2293 | 22.14 | 229\% |
| 4370 | 2296 | 2297 | 224H | 2299 | 2300 | 2301 | 2302 | 23013 |


|  | 0 | 1 | 2 | J | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 4409 | 2304 | 2305 | 2306 | 2307 | 2308 | 2309 | 2310 | 2311 |
| 4410 | 2312 | 2313 | 2314 | 2315 | 2316 | $2317^{\circ}$ | 2318 | 2319 |
| 4420 | 2320 | 2321 | 2322 | 2323 | 2324 | 2325 | 2326 | 2327 |
| 4430 | 2328 | 2329 | 2330 | 2331 | 2332 | 2333 | 2334 | 2335 |
| 4440 | 2336 | 2337 | 2338 | 2339 | 2340 | 2341 | 2342 | 2343 |
| 4450 | 2344 | 2345 | 2346 | 2347 | 2348 | 2349 | 2350 | 2351 |
| 4460 | 2352 | 2353 | 2354 | 2355 | 2356 | 2357 | 2358 | 2359 |
| 4470 | 2360 | 2361 | 2362 | 2363 | 2364 | 2365 | 2366 | 2367 |
| 4500 | 2368 | 2369 | 2370 | 2371 | 2372 | 2373 | 2374 | 2375 |
| 4510 | 2376 | 2377 | 2378 | 2379 | 2380 | 2381 | 2382 | 2383 |
| 4520 | 2384 | 2385 | 2386 | 2387 | 2388 | 2389 | 2390 | 2391 |
| 4530 | 2392 | 2393 | 2394 | 2395 | 2396 | 2397 | 2398 | 2399 |
| 4540 | 2400 | 2401 | 2402 | 2403 | 2404 | 2105 | 2406 | 2407 |
| 4550 | 2408 | 2409 | 2410 | 2411 | 2412 | 2413 | 2414 | 2415 |
| 4560 | 2416 | 2417 | 2418 | 2419 | 2420 | 2421 | 2422 | 2423 |
| 4570 | 2424 | 2425 | 2426 | 2427 | 2428 | 2429 | 2430 | 2431 |
| 4600 | 2432 | 2433 | 2434 | 2435 | 2436 | 2437 | 2438 | 2439 |
| 4610 | 2440 | 2441 | 2442 | 2443 | 2444 | 2445 | 2446 | 2447 |
| 4620 | 2448 | 2449 | 2450 | 2451 | 2452 | 2453 | 2454 | 2455 |
| 4630 | 2456 | 2457 | 2458 | 2459 | 2460 | 2461 | 2462 | 2463 |
| 4640 | 2464 | 2465 | 2466 | 2467 | 2468 | 2469 | 2470 | 2471 |
| 4650 | 2472 | 2473 | 2474 | 2475 | 2476 | 2477 | 2478 | 2479 |
| 4660 | 2480 | 2481 | 2482 | 2483 | 2484 | 2485 | 2486 | 2487 |
| 4670 | 2488 | 2489 | 2490 | 2491 | 2492 | 2493 | 2494 | 2495 |
| 4700 | 2496 | 2497 | 2498 | 2499 | 2500 | 2501 | 2502 | 2503 |
| 4710 | 2504 | 2505 | 2506 | 2507 | 2508 | 2509 | 2510 | 2511 |
| 4720 | 2512 | 2513 | 2514 | 2515 | 2516 | 2517 | 2518 | 2519 |
| 4730 | 2520 | 2521 | 2522 | 2523 | 2524 | 2525 | 2526 | 2527 |
| 4340 | 2528 | 2ヶ29 | 2530 | 2531 | 2532 | 2533 | 2534 | 2535 |
| 4750 | 2535 | 2537 | 2538 | 2539 | 2540 | 2541 | 2542 | 2543 |
| 4760 | 24,44 | 2545 | 2546 | 2547 | 2548 | 2549 | 2550 | 2551 |
| 4770 | 2552 | 2553 | 2554 | 2555 | 2556 | 2557 | 2558 | 2559 |


| 5000 | 2500 |
| :---: | :---: |
| 10 | 10 |
| 5777 | 3071 |
| Ociol | Decimal) |


|  | 7 | 1 | 2 | 3 | 4 | 5 | 6 | $\stackrel{\rightharpoonup}{4}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 5000 | 2560 | 2561 | 2562 | 2563 | 25,4 | 2,54 | 2566 | 2567 |
| 5010 | 2568 | 2,599 | 2570 | 2371 | $25+2$ | 2303 | $25,-4$ | 2575 |
| , 5020 | 2576 | 2577 | 2578 | 2599 | 2580 | 2581 | < 682 | 2583 |
| 5030 | 2584 | 2585 | 2586 | 2587 | 2588 | 25.84 | 25.40 | 2591 |
| 5046 | 2592 | 2:93 | 2554 | 2595 | 2596 | 259\% | 2396 | 2599 |
| 5050 | 2600 | 2601 | 2602 | 2603 | 2504 | 2605 | 2506 | 2507 |
| 5060 | 2608 | 2609 | 2510 | 2611 | 2612 | 2613 | 2614 | 2615 |
| 5070 | 2616 | 2817 | 2618 | 2619 | 2620 | 2621 | 2622 | 2623 |
| 3100 | 2624 | 2625 | 2626 | 2627 | 2628 | 2629 | 2630 | 2631 |
| 5110 | 26.32 | 2633 | 2634 | 2635 | 2636 | 2637 | 2638 | 2639 |
| 5120 | 2640 | 2811 | 2642 | 2643 | 2644 | 2645 | 2646 | 2647 |
| 5130 | 2648 | 2649 | 2650 | 2651 | 2652 | 2853 | 2654 | 2655 |
| 5140 | 2556 | 2857 | 2658 | 2659 | 2650 | 2661 | 2562 | 2663 |
| 5150 | 2664 | 2665 | 2666 | 2667 | 2668 | 2669 | 2670 | 2671 |
| 5160 | 2872 | $26^{73}$ | 2674 | 2675 | 2676 | 2677 | 267B | 2699 |
| 5170 | 2680 | 2681 | 2632 | 2693 | 2584 | 2685 | 2686 | 2687 |
| - 5200 | 2688 | 2689 | 2690 | 2691 | 2692 | 2693 | 2694 | 2695 |
| 5210 | 2696 | 2697 | 26.98 | 2599 | 2700 | 2701 | 2:02 | 2703 |
| 5220 | 2704 | 2705 | 2706 | 2707 | 2708 | 2709 | 2710 | 2711 |
| 5230 | 2712 | 2713 | 2714 | 2715 | 2716 | 2717 | 2718 | 2719 |
| 5240 | 2720 | 2721 | 2722 | 27.23 | 2724 | 2725 | 27.26 | 2727 |
| 5250 | 2728 | 2:29 | 2:30 | 2:31 | 2732 | 2733 | 2734 | 27.75 |
| 5260 | 2736 | 2:37 | 2738 | 2734 | 2740 | 2741 | 2742 | 2743 |
| 5270 | 2744 | 2745 | 2746 | 2747 | 2748 | 2749 | 2750 | 2751 |
| 5300 | 2752 | 2753 | 2754 | 2755 | 2756 | 2757 | 2758 | 2759 |
| 5310 | 2760 | 2761 | 2762 | 2763 | 2764 | 2765 | 2766 | 2767 |
| 5320 | 2768 | 2769 | 2770 | 2771 | 2772 | 2773 | 2774 | 2775 |
| 5339 | 2778 | 2777 | 2778 | 2779 | 2780 | 2781 | 2782 | 2783 |
| 5310 | 2784 | 2785 | 2786 | 2787 | 2788 | 2789 | 2790 | 2791 |
| 5350 | $2: 92$ | 2993 | c $7!4$ | 2795 | 2796 | 2797 | 2798 | 2799 |
| 53 EO | 2600 | 2501 | 2802 | 2803 | 2804 | 2805 | 2806 | 2807 |
| $53: 0$ | 2008 | 2899 | 2810 | 2811 | 2812 | 281? | 2814 | 2815 |


|  | 0 | 1 | 2 | 3 | 4 | j | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 5400 | 2816 | 2817 | 2818 | 2819 | 2820 | 2831 | 2822 | 2823 |
| 5410 | 2824 | 2825 | 2826 | 2827 | 2828 | 2829 | 2830 | 2831 |
| 5420 | 2832 | 2837 | 2834 | 2835 | 2836 | 2837 | 2818 | 2839 |
| 4430 | 2810 | 28:1 | 2842 | 2843 | 2844 | 28:5 | 2846 | 2847 |
| 5440 | 2846 | 2849 | 2850 | 2851 | 2852 | 28,3 | 2854 | 2855 |
| 5450 | 2856 | 2857 | 2858 | 2859 | 2860 | 2861 | 2862 | 2863 |
| 5460 | 2864 | 2865 | 2866 | 2867 | 2868 | 2869 | 2870 | 2871 |
| 5470 | 2872 | 2873 | 2874 | 2875 | 2876 | 2877 | 2878 | 2879 |
| 5500 | 2880 | 2881 | 2882 | 2883 | 2884 | 2885 | 2886 | 2887 |
| 5510 | 2888 | 2889 | 2890 | 2891 | 2892 | 2893 | 2894 | 2895 |
| 5520 | 2896 | 2897 | 2898 | 2899 | 2900 | 2901 | 2902 | 2903 |
| 5530 | 2904 | 2905 | 2906 | 2907 | 2908 | 2909 | 2910 | 2011 |
| 5540 | 2912 | 2913 | 2914 | 2915 | 2916 | 2917 | 2918 | 2919 |
| 5550 | 2920 | 2921 | 2922 | 2923 | 2924 | 2925 | 2926 | 2927 |
| 5560 | 2928 | 2929 | 2930 | 2931 | 2932 | 2933 | 2934 | 2935 |
| 5570 | 2936 | 2937 | 2938 | 2939 | 2940 | 2941 | 2942 | 2943 |
| 5600 | 2944 | 2945 | 2946 | 2947 | 2948 | 2949 | 2950 | 2951 |
| 5610 | 2952 | 2953 | 2954 | 2955 | 2956 | 2957 | 2958 | 2959 |
| 4520 | 2960 | 2961 | 2762 | 2963 | 2964 | 2965 | 2966 | 2967 |
| 5630 | 2958 | 2969 | 2970 | 2971 | 2972 | 2973 | 2974 | 2975 |
| 5640 | 2976 | 2977 | 2978 | 2979 | 2980 | 2981 | 2962 | 2983 |
| 5650 | 2984 | 2985 | 2986 | 2987 | 2988 | 2989 | 2990 | 2991 |
| 5660 | 2992 | 2993 | 2994 | 2995 | 2996 | 2997 | 2998 | 2999 |
| :670 | 3000 | 3001 | 3002 | 3003 | 3004 | 3005 | 3008 | 3007 |
| , 700 | 3008 | 3009 | 3010 | 3011 | 3012 | 3013 | 3014 | 3015 |
| 5710 | 3016 | 3017 | 3018 | 3019 | 3020 | 3021 | 3022 | 3023 |
| 5720 | 3024 | 3025 | 3026 | 3027 | 3028 | 3029 | 3030 | 3031 |
| 5730 | 3032 | 3033 | 3034 | 3035 | 3036 | 3037 | 3038 | 3039 |
| 5740 | 3040 | 3041 | 3042 | 3043 | 3044 | 3045 | 3046 | 3047 |
| 5750 | 3048 | 3049 | 3050 | 3051 | 3052 | 3053 | 3054 | 3055 |
| 57.60 | 3056 | 3057 | 3058 | 3059 | 3060 | 3061 | 3062 | 3063 |
| 5770 | 3064 | 3065 | 3066 | 3067 | 3068 | 3069 | 3070 | 3071 |


| 6000 | 3072 |
| :---: | :---: |
| 10 | 10 |
| 6777 | 3583 |
| (Octal) | (Decimal) |


|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 6000 | 3072 | 3073 | 3074 | 3075 | 3078 | 3077 | 3078 | 3079 |
| 6010 | 3080 | 3081 | 3082 | 3083 | 3084 | 3085 | 3086 | 3087 |
| 6020 | 3088 | 3089 | 3090 | 3091 | 3092 | 3093 | 3094 | 3095 |
| 6030 | 3096 | 3097 | 3088 | 3099 | 3100 | 3101 | 3102 | 3103 |
| 6040 | 3104 | 3105 | 3106 | 3107 | 3108 | 3109 | 3110 | 3111 |
| 6050 | 3112 | 3113 | 3114 | 3115 | 3116 | 3117 | 3118 | 3119 |
| 6080 | 3120 | 3121 | 3122 | 3123 | 3124 | 3125 | 3126 | 3127 |
| 6070 | 3128 | 3129 | 3130 | 3131 | 3132 | 3133 | 3134 | 3135 |
| 6100 | 3136 | 3137 | 3138 | 3139 | 3140 | 3141 | 3142 | 3143 |
| 6110 | 3144 | 3145 | 3146 | 3147 | 3148 | 3149 | 3150 | 3151 |
| 6120 | 3152 | 3153 | 3154 | 3155 | 3156 | 3157 | 3158 | 3159 |
| 6130 | 3160 | 3161 | 3162 | 3163 | 3164 | 3165 | 3168 | 3187 i |
| 6140 | 3168 | 3169 | 3170 | 3171 | 3172 | 3179 | 3174 | 3175 |
| 6150 | 3176 | 317: | 3178 | 3178 | 3180 | 3181 | 3182 | 3183 |
| 6180 | 3184 | 5185 | 3188 | 3187 | 3188 | 3189 | 3190 | 3191 |
| 6170 | 3192 | 3193 | 3194 | 3195 | 3196 | 3197 | 3198 | 3199 |
| 0200 | 3200 | 3201 | 3202 | 3203 | 3204 | 3205 | 3206 | 3207 |
| 6210 | 3208 | 3209 | 3210 | 3211 | 3212 | 3213 | 3214 | 321 |
| 6220 | 3216 | 3217 | 3218 | 3219 | 3220 | 3221 | 3222 | 322 |
| 8230 | 3224 | 3225 | 3226 | 3227 | 3228 | 3229 | 3230 | 3231 |
| 6240 | 3232 | 3233 | 3234 | 3235 | 3238 | 3237 | 3238 | 3239 |
| 6250 | 3240 | 3241 | 3242 | 3243 | 3244 | 3245 | 3248 | 3247 |
| 6280 | 3248 | 3249 | 3250 | 3251 | 3252 | 3253 | 3254 | 3255 |
| 6270 | 3250 | 3257 | 3258 | 3250 | 3280 | 3261 | 5282 | 3263 |
| 6300 | 3264 | 3265 | 3206 | 3267 | 3268 | 3289 | 3270 | 3271 |
| 6310 | 3272 | 3273 | 3274 | 3275 | 3276 | 3277 | 3278 | 3279 |
| 6320 | 3280 | 3281 | 3282 | 3283 | 3284 | 3205 | 3206 | 3287 |
| 6330 | 3288 | 3289 | 3290 | 3201 | 3202 | 3293 | 3204 | 3295 |
| 6340 | 3296 | 3297 | 3298 | 3299 | 3300 | 3301 | \$302 | 3303 |
| 6350 | 3304 | 3305 | 3306 | 3307 | 3303 | 3309 | 3310 | 3311 |
| 6360 | 3312 | 3313 | 3314 | 3315 | 3316 | 3317 | 3318 | 3319 |
| 8370 | 3320 | 3321 | 3322 | 3323 | 3324 | 3325 | 3326 | 3327 |


|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 6400 | 3328 | 3329 | 3330 | 3331 | 3332 | 3935 | 3334 | 3335 |
| 6410 | 3336 | 3337 | 3338 | 3339 | 3340 | 3341 | 3342 | 1343 |
| 6420 | 3344 | 3343 | 3346 | 3347 | 3348 | 3349 | 3350 | 3351 |
| 6430 | 3352 | 3353 | 3354 | 3355 | 3356 | 3357 | 3358 | 3359 |
| 6440 | 3360 | 3361 | 3362 | 3383 | 3364 | 3365 | 3366 | 3187 |
| 6450 | 3368 | 3369 | 3370 | 3371 | 3372 | 3373 | 3374 | 3375 |
| 6460 | 3376 | 3377 | 3378 | 3379 | 3380 | 3381 | 3382 | 3383 |
| 6470 | 3384 | 3385 | 3386 | 3387 | 3388 | 3389 | 3390 | 3391 |
| 8500 | 3392 | 3393 | 3394 | 3395 | 3396 | 3397 | 3398 | 3399 |
| 6510 | 3400 | 3401 | 3402 | 3403 | 3404 | 3405 | 3408 | 3407 |
| 0520 | 3408 | 3409 | 3410 | 3411 | 3412 | 3413 | 3414 | 3415 |
| 6530 | 34:6 | 3417 | 3418 | 3419 | 3420 | 3421 | 3422 | 3423 |
| 6540 \| | 3424 | 3425 | 3426 | 3427 | 3428 | 3429 | 3430 | 3431 |
| 0550 | 3432 | 3433 | 3434 | 3435 | 3436 | 3437 | 3438 | 3439 |
| 6560 | 3440 | 3441 | 3442 | 3443 | 3444 | 3445 | 3446 | 3447 |
| 6570 | 3448 | 3449 | 3450 | 3451 | 3452 | 3453 | 3454 | 3455 |
| 6600 | 3456 | 3457 | 3458 | 3459 | 3460 | 3461 | 3462 | 3463 |
| 6810 | 3464 | 3465 | 3466 | 3467 | 3488 | 3469 | 3470 | 3471 |
| 6620 | 3472 | 3473 | 3474 | 3475 | 3476 | 3477 | 3478 | 3479 |
| 6630 | 3480 | 3481 | 3482 | 3483 | 3484 | 3485 | 3486 | 3487 |
| e640 | 3488 | 3489 | 3490 | 3491 | 3492 | 3493 | 3494 | 3495 |
| 6650 | 3496 | 3497 | 3498 | 3499 | 3500 | 3501 | 3502 | 3503 |
| 6680 | 3504 | 3505 | 3506 | 3507 | 3508 | 3509 | 3510 | 3511 |
| 6670 | 3512 | 3513 | 3514 | 3515 | 3516 | 3517 | 3518 | 3519 |
| 6700 | 3520 | 3521 | 3522 | 3523 | 3524 | 3525 | 3526 | 3527 |
| 6710 | 3528 | 3529 | 3530 | 3531 | 3532 | 3533 | 3534 | 3535 |
| 6720 | 3536 | 3537 | 3538 | 3539 | 3540 | 3541 | 3542 | 35431 |
| 6730 | 3544 | 3545 | 3540 | 3547 | 3548 | 3549 | 3550 | 3551 |
| 6740 | 3552 | 3553 | 3554 | 3555 | 3556 | 3557 | 3558 | 3558 |
| 6750 | 3560 | 3561 | 3562 | 3563 | 3564 | 3565 | 3568 | 3587 |
| 6780 | 3568 | 3569 | 3570 | 3571 | 3572 | 3573 | 3574 | 3575 |
| 6770 | 3576 | 3577 | 3578 | 3578 | 3580 | 3581 | 3582 | 3583 |

# OCTAL-DECIMAL INTEGER CONVERSION TABLE (Cont'd) 

| 7000 | 3984 |
| :---: | :---: |
| 10 | 10 |
| 7777 | 4095 |
| (Octol) | (Decimal) |


|  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 7000 | 3584 | 3585 | 3586 | 3587 | 3588 | 3589 | 3590 | 3591 | 7400 | 3840 | 3841 | 3842 | 3843 | 3844 | 3845 | 3846 | 3847 |
| 7010 | 3592 | 3593 | 3594 | 3595 | 3596 | 3597 | 3508 | 3509 | 7410 | 3848 | 3849 | 3850 | 3851 | 3852 | 3853 | 3854 | 3855 |
| 7020 | 3000 | 3601 | 3602 | 3603 | 3604 | 3005 | 3006 | 3607 | 7420 | 3856 | 3857 | 3858 | 3859 | 3880 | 3861 | 3862 | 3883 |
| 7030 | 3808 | 3609 | 3010 | 3611 | 3612 | 3013 | 3614 | 3615 | 7430 | 3864 | 3865 | 3866 | 3867 | 3868 | 3869 | 38:0 | 3871 |
| 1040 | 3616 | 3617 | 3018 | 3319 | 3620 | 3021 | 3022 | 3023 | 7440 | 3872 | 3873 | 3874 | 3875 | 3878 | 3877 | 3878 | 3879 |
| 7050 | 3624 | 3025 | 3828 | 3827 | 3028 | 3029 | 3030 | 3631 | 7450 | 3000 | 3881 | 3882 | 3883 | 3884 | 3885 | 3886 | 3887 |
| 7000 | 3632 | 3033 | 3834 | 3035 | 3636 | 3037 | 3638 | 3639 | 7460 | 5888 | 3889 | 3890 | 3891 | 3892 | 3893 | 3894 | 3895 |
| 7070 | 3640 | 3641 | 3842 | 3643 | 3644 | 3645 | 3846 | 3647 | 7470 | 3898 | 3897 | 3898 | 3899 | 3900 | 3901 | 3902 | 3903 |
| 7100 | 3648 | 3649 | 3850 | 3651 | 3652 | 3853 | 3854 | 3655 | 7500 | 3904 | 3905 | 3906 | 3907 | 3908 | 3900 | 3910 | 3911 |
| 7110 | 3656 | 3057 | 3058 | 3859 | 3060 | 3661 | 3662 | 3663 | 7510 | 3912 | 3913 | 3914 | 3915 | 3916 | 3917 | 3918 | 3919 |
| 7120 | 3664 | $3 \mathrm{CO5}$ | $3 \mathrm{C30}$ | 3037 | 3068 | 3009 | 3870 | 3671 | 7520 | 3020 | 3921 | 3922 | 3923 | 3924 | 3925 | 3926 | 3927 |
| 7130 | 3072 | 3073 | 3074 | 3675 | 3870 | 3877 | 3678 | 3679 | 7530 | 3928 | 3929 | 3930 | 3931 | 3932 | 3933 | 3934 | 3935 |
| 7140 | 3600 | 3681 | 3 C 82 | 3683 | 3004 | 3035 | $3 \mathrm{CB6}$ | 3687 | 7540 | 3936 | 3937 | 3938 | 3939 | 3940 | 3941 | 3942 | 3943 |
| 7150 | 3688 | 3039 | 3000 | 3001 | 3002 | 3093 | 3694 | 3095 | 7530 | 3944 | 3945 | 3946 | 3947 | 3948 | 3949 | 3950 | 3951 |
| 7160 | 3696 | 3097 | 3698 | 3099 | 3700 | 3701 | 3702 | 3703 | 7560 | 3952 | 3953 | 3954 | 3955 | 3956 | 3957 | 3958 | 3959 |
| 7170 | 3704 | 3705 | 3708 | 3707 | 3708 | 3709 | 3710 | 3711 | 7570 | 3960 | 3961 | 3982 | 3963 | 3964 | 3985 | 3966 | 3967 |
| 7200 | 3712 | 3713 | 3714 | 3715 | 3718 | 3717 | 3718 | 3719 | 7800 | 3988 | 3969 | 3970 | 3971 | 3972 | 3973 | 3974 | 3975 |
| 7210 | 3720 | 3721 | 3722 | 3723 | 3724 | 3725 | 3728 | 3727 | 7610 | 3976 | 3977 | 3978 | 3979 | 3980 | 3981 | 3982 | 3983 |
| 7220 | 3728 | 3729 | 3730 | 3731 | 3732 | 3733 | 3734 | 3735 | 7620 | 3984 | 3885 | 3986 | 3987 | 3988 | 3989 | 3980 | 3991 |
| 7230 | 3736 | 3737 | 3738 | 3738 | 3740 | 3741 | 3742 | 3743 | 7630 | 3992 | 3993 | 3994 | 3995 | 3998 | 3997 | 3998 | 3999 |
| 7240 | 3744 | 3745 | 3746 | 3747 | 3748 | 3749 | 3750 | 3751 | 7640 | 4000 | 4001 | 4002 | 4003 | 4004 | 4005 | 4008 | 4007 |
| 7250 | 3752 | 3753 | 3754 | 3755 | 3756 | 3757 | 3758 | 3759 | 7850 | 4008 | 4009 | 4010 | 4011 | 4012 | 4013 | 4014 | 4015 |
| ¢260 | 3780 | 3761 | 3762 | 3763 | 3764 | 3765 | 3766 | 3767 | 7660 | 4016 | 4017 | 4018 | 4019 | 4020 | 4021 | 4022 | 4023 |
| 7270 | 3768 | 3769 | 3770 | 3771 | 3772 | 3773 | 3774 | 3775 | 7870 | 4024 | 4025 | 4028 | 4027 | 4028 | 4029 | 4030 | 4031 |
| 7300 | 3778 | 3777 | 3778 | 3779 | 3780 | 3781 | 3782 | 3783 | 7700 | 4032 | 4033 | 4034 | 4035 | 4036 | 4037 | 4038 | 4039 |
| 7310 | 3784 | 3785 | 3786 | 3787 | 3788 | 3789 | 3790 | 3791 | 7710 | 4040 | 4041 | 4042 | 4043 | 4044 | 4045 | 4046 | 4047 |
| 7320 | 3792 | 3793 | 3794 | 3795 | 3796 | 3797 | 3798 | 3799 | 7720 | 4048 | 4049 | 4050 | 4051 | 4052 | 4053 | 4054 | 4055 |
| 7330 | 2800 | 3801 | 3802 | 3803 | 3804 | 3805 | 3806 | 3807 | 7730 | 4056 | 4057 | 4058 | 4059 | 4060 | 4081 | 4062 | 4063 |
| 7340 | 3808 | 3809 | 3810 | 3811 | 3812 | 3813 | 3814 | 3815 | 7740 | 4064 | 4065 | 4068 | 4067 | 4088 | 4089 | 4070 | 4071 |
| 7353 | 3816 | 3817 | 1818 | 3819 | 3820 | 3821 | 3822 | 3823 | 7750 | 4072 | 4073 | 4074 | 4075 | 4076 | 4077 | 4078 | 4079 |
| 7360 | 3824 | 3825 | 3826 | 3827 | 3828 | 3829 | 3830 | 3831 | 7760 | 4080 | 4081 | 4082 | 4083 | 4084 | 4085 | 4086 | 4087 |
| 7370 | 3832 | 3833 | 3834 | 3835 | 3836 | 3837 | 3838 | 3839 | 7770 | 4088 | 4089 | 4090 | 4091 | 4092 | 4093 | 4094 | 4095 |

## APPENDIX C

## NHMERIC: INPUT-OUTPUT CODES

a. The following codes are employed in numeric input-output in the BCD and Oetal modes:

| Teletype Meaning |  | FADAC Meaning |  |  | Fieldata Meaning |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | Octal <br> Mode | $\begin{gathered} \mathrm{BCD} \\ \text { Mode } \end{gathered}$ | Binary Conversion |  |  |
| Character | Tape code |  |  |  | Tape Code | Character |
| O, P | (0)101 | () | 0 | 0000 | 00110000 | 0 |
| 1, Q | 11101 | 1 | 1 | 0001 | 10110001 | 1 |
| 2, w | 11001 | 2 | 2 | 0010 | 10110010 | 2 |
| 3, E | 10000 | 3 | 3 | 0011 | 00110011 | 3 |
| 4, R | 01610 | 4 | 4 | 0100 | 10110100 | 4 |
| 5, T | (0)001 | 5 | 5 | 0101 | 00110101 | 5 |
| 6, Y | 10101 | 6 | 6 | 0110 | 00110110 | 6 |
| 7, 1 | 11100 | 7 | 7 | 0111 | 10110111 | 7 |
| 8, I | 01.100 | Halt | 8 | 1000 | 101111000 | 8 |
| 9,0 | (0)(1)11 | Compute | 9 | 1001 | 00111001 | 9 |
| (+) ' ${ }^{\prime}$, Z | 10001 | Fill | + | 1010 | 00100010 | + |
| -, A | 11000 | Verify | - | 1011 | 00100001 | - |
| ., M | OM111 | Location |  | 1100 | 10111101 |  |
| /, X | 10111 | Clear | Clear | 1101 | 1101L101 | X |
| Car. Ret | Oukiol | Enter | Enter | 1110 | 01000100 | Car. Ret |
| Blank | 0 monos | Blank | Blank | 1111 | 00000000 | Blank |

b. The line through the tape code field indicates the position of the sprocket holes. (ode is not applicabl to nametic tape.

## APPENDIX D

## ALPHANUMERIC INPUT•OUTPUT CODES

a. The following paper tape codes are used in alphanumeric inputooutput:
b. The line through the tape code field indicates the position of the sprocket holes.
c. The bit positions of the above codes are numbered 1 through 8 from right to left. Bit number 8 is parity bit; bit number 7 is control bit; bits number 6 through 1 are information bits.

| FIELDAIA |  | $\left[\begin{array}{c} 3 \\ y \\ z \end{array}\right.$ | $\left\lvert\, \begin{array}{\|l\|} \text { FADAC } \\ \text { Binary Code } \end{array}\right.$ | FIELDATA |  | - | $\frac{\text { FADAC }}{} / \text { Binary Code }$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Tape Code | lipper ciase |  |  | Tape Code | Upper Case |  |  |
| $11000 \text { 000 }$ | Master Sp | U1 | 0000000 | 10100kote | ) |  | 100000 |
| 010000001 | l'ppercase | 2 | noworl | 00100001 | - |  | 100001 |
| (1)10000310 | Lower case | $\%$ | (6)H010 | 001009010 | + |  | 100010 |
| $110 \mathrm{coshl1}$ | Tab | $\bigcirc$ | 000011 | 10100 (1) 1 | $<$ |  | 100011 |
| 010nction | Car. ret. | $<$ | (100)100 | 00100.100 | $=$ |  | 100100 |
| 11000101 | Space | $\triangle$ | 000101 | 1010010 ! | $>$ |  | 100101 |
| 11006110 | A | A | 00 n 110 | 10100110 |  |  | 100110 |
| 01000111 | B | B | 000111 | 0010nt11 | \$ |  | 100111 |
| O1001600 | C | C. | 001000 | (i) 101000 | * |  | 101000 |
| 11001061 | D | D | 001001 | 16atom | ( |  | 101001 |
| 11001010 | E | E | 001010 | 10101010 | " |  | 101010 |
| 01001011 | F | F | 001011 | 00101011 | : |  | 101011 |
| 11001100 | G | G | 001100 | 10101100 | ? |  | 101100 |
| 01001101 | H | H | 001101 | 00101101 | ! |  | 101101 |
| 010011110 | I | I | 001110 | 00101110 |  |  | 101110 |
| 11001111 | 1 | J | 001111 | 10101111 | Stop $\oplus$ |  | 101111 |
| 01010 ¢0\% | K | K | 010000 | 00110000 |  | 0 | 110000 |
| 1101000 : | L | L | 010001 | 10110001 |  | 1 | 110001 |
| 11010010 | M | 11 | 010010 | 10110010 |  | 2 | 110010 |
| 01010011 | N | $\cdots$ | 010011 | 00110011 |  | 3 | 110011 |
| 11010100 | O | 0 | 010100 | 101101100 |  | 4 | 110100 |
| 01010161 | P | P | 010101 | 00110101 |  |  | 110101 |
| 01010,110 | Q | Q | 010110 | 00110110 |  | 6 | 110110 |
| 11010111 | R | R | 010111 | 10410\|111 |  | 7 | 110111 |
| 11011000 | S | S | 011000 | 10111)000 |  | 8 | 111000 |
| 01011001 | T | T | 011001 | 00111001 |  | - | 111001 |
| 0101 mbl | ! | ${ }^{\prime}$ | 011010 | 00111010 |  | - | 111010 |
| 11016111 | $V$ | $\checkmark$ | 011011 | 10111011 |  | ; | 111011 |
| 01011169 | W | W | 011100 | 00113100 |  | , | 111100 |
| 11011\|111 | X | X | 011101 | 10111101 |  |  | 111101 |
| 11011110 | Y | Y | 011110 | 10111110 | Special $\square$ |  | 111110 |
| 010111111 | 7. | Z | 0111111 | 1001111111 | Back space |  | 111111 |

APPENDIX E

OPTIMIZATION TABLE

| Operation Code | $\underset{\text { Time }}{\text { Execution }}$ <br> Time | Optimum Operand Sector* | Optimum <br> Next I Sector |
| :---: | :---: | :---: | :---: |
| CLA | 1 | $t+1$ | $\mathrm{t}+1$ |
| CLS | 1 | t+1 | t+1 |
| ADD | 1 | t+1 | t+1 |
| SUB | 1 | $\mathrm{t}+1$ | $t+1$ |
| MPL | 18 | $\mathrm{t}+1$ | $t+1 \quad t+18$ |
| DIV | 18 | $t+1$ | $t+1 \quad t+18$ |
| STA ${ }^{1}$ | 1 | $t-17$ | $t+1$ |
| STN ${ }^{1}$ | 1 | $t-1\}$ | $\mathrm{t}+1$ |
| STD ${ }^{1}$ | 1 | $t-1\}^{*}$ | $t+1$ |
| STL ${ }^{1}$ | 1 | $t-1$ | $t+1$ |
| STO | ( $1 \cdot R \& Q$ |  | $t+1 \quad R \& Q$ |
| STP | $\int \begin{gathered}3-\text { anything } \\ \text { else }\end{gathered}$ | t +1 | t+3 ** |
| LDR | 16 | $\mathrm{t}+1$ | $t+1 \quad t+16$ |
| LDQ | 16 | t+1 | $t+1 \quad \mathrm{t}+16$ |
| STR | 18 | $\mathrm{t}+1$ | $\mathrm{t}+1 \quad \mathrm{t}+18$ |
| TZE | 1 | $t+1$ | $t+1$ |
| TPL | 1 | t +1 | $\mathrm{t}+1$ |
| TRA | 1 | $\mathrm{t}+1$ | $t+1$ |
| TOV | 1 | $t+1$ | $t+1$ |
| ARC LLS | Variable | Variable | Variable |
| EXT | 1 | $\mathrm{t}+1$ | $t+1$ |
| ABS | 1 | N/A | $\mathrm{t}+1$ |
| RML | 1 | N/A | $t+1$ |

* $\mathrm{t}-1$ in main memory; $\mathrm{t}+1$ in R or Q -Loops.
** $\mathrm{t}+3$ in main memory.


## OP TIMIZATION TABLE (Cont ${ }^{\text {d }}$ d

| Operation <br> Code | Execution <br> Time | Optimum <br> Operand <br> Sector | Optimum <br> Next I <br> Sector |
| :---: | :---: | :---: | :---: |
| RMN | 1 | N/A | $\mathrm{t}+1$ |
| ZEL | 1 | N/A | $\mathrm{t}+1$ |
| EQS | Variable | Variable | Variable |
| GES | 1 | Variable | Variable |
| HLT | 1 | N/A | $\mathrm{t}+1$ |
| HDM | 1 | N/A | $\mathrm{t}+1$ |
| IDM | 1 | N/A | $\mathrm{t}+1$ |
| DIA | 1 | N/A | $\mathrm{t}+1$ |
| DOF NSL | 1 | N/A | $\mathrm{t}+1$ |

## APPENDIX F

## SUMMARY OF FADAC INSTRUCTION REPERTOIRE

| Operation |  | Mnernonic | Code | $M$ is the Operand Address $[\mathrm{M}$ ] indicates Contenta of M ) |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | Effect of Flag \& Notes |  | No. of Word Times for Execution |
| I. ARITHMETIC |  |  |  |  |  |
|  | Add |  | ADD | 00 | ADD* $=[\mathrm{M}]-\mathrm{N}$ | 1 |
|  | Subtract | SUB | 02 | SUB* $=[\mathrm{M}]-\mathrm{N}$ | 1 |
|  | Multiply | MPY | 20 | No Flag. $[\mathrm{M}] \rightarrow \mathrm{N}$ | 18 |
|  | Clear * Add | CLA | 24 | $\begin{aligned} C L A *= & \text { previous contents of } \\ & A-L \end{aligned}$ | 1 |
|  | Clear \& Subtract | CLS | 26 | $\begin{aligned} \text { CLS* }= & \text { previous contents of } \\ & A-L \end{aligned}$ | 1 |
|  | Divide | DIV | 30 | $\left\{\begin{aligned} & \text { DIV } *= \text { unrounded Quotient in } A . \\ & \text { Remainder in } L \\ & \text { DIV }= \text { rounded Quotient in } A \\ & \text { Contents of L meaningless } \\ & \text { For both DIV \& DIV }[M] \rightarrow N \end{aligned}\right.$ | 18 |
| II. STORE * LOAD |  |  |  |  |  |
|  | Store N | STN | 40 | No flag | 1 |
|  | Store D | STD | 42 | $\left\{\begin{array}{l} S T D *=D . i-N \\ D . i=D .0 \text { if } M \text { is even; } \\ D . \mathrm{i}=\mathrm{D} .1 \text { if } \mathrm{M} \text { is odd } \end{array}\right.$ | 1 |
|  | Store A | STA | 50 | STA* $=\mathbf{A} \rightarrow \mathrm{N}$ | 1 |
|  | Store L | STL | 52 | STL* $=\mathbf{L}-\mathrm{N}$ | 1 |
|  | Store Op Ad | STO | 70 | Noflag | $\int 1$ if $M$ in $R$ or $Q$ |
|  | Store (Ni) Prog Ad | STP | 60 | Noflag | , 3 if M in Main Memory |
|  | Store R loop | STR | 62 | Noflag (Note: No store $Q$ instruction exists) | 18 |
|  | Load R loop Load O loop |  | 72 72 | $\left\{\begin{array}{l} \text { LDR } \# \text { = } R \text { loop filled } \\ \text { LDQ (no flag) }=Q \text { loop filled } \\ \text { Note: flag deter mines whether } \\ R \text { or Q loop is filled. } \\ \text { Code } 72 \text { is used for both. } \end{array}\right.$ | $\{16$ |
| III. | TRANSFER |  |  |  |  |
|  | Transfer on Plus | TPL | 10 | TPL* $=$ Transfer instruction -N | 1 |
|  | Transfer on Zero | TZE | 12 | TZE* $=$ Transfer instruction -N | 1 |
|  | Unconditional Transfer | TRA | 14 | TRA* $=$ Transfer instruction $\rightarrow \mathrm{N}$ | 1 |
|  | Transfer on Overflow | TOV | 16 | TOV* $=$ Transfer instruction $\rightarrow \mathrm{N}$ | 1 |



Note: Bits which normally designate the channel of the Op Ad are part of the Op Code of most Commands in Sections IV, V, and VI.

| Operation | Mnemonic | Code | Effect of Flag \& Notes | No. of Word Times for Execution |
| :---: | :---: | :---: | :---: | :---: |
| VI. SERIAL $1 / 0$ |  |  |  |  |
| Discrete Input to A | DIA | 3640 | No flag | 1 |
| Discrete Outp Off | DOF | 3700 | No flag | 1 |
| Outp Dev Stepping - 1 | ODI | 3702 | No flag | 1 |
| Outp Dev Stepping - 2 | OD 2 | 3704 | No flag | 1 |
| Outp Dev Stepping - 3 | OD3 | 3706 | Noflag | 1 |
| Inp Dev Stepping - 1 | ID1 | 3710 | No flag | 1 |
| Inp Dev Stepping -2 | ID2 | 3712 | No flag | t |
| Inp Dev Stepping -3 | LD3 | 3714 | No flag | 1 |
| No Solution Light | NSL | 3716 | Noflag | 1 |

VII. PARALLEL INPUT/OUTPUT COMMAND SUMMARY

In all Parallel I/O commands, a flag indicates input (output) by WORD, no flag indicates input (output) by CHARACTER. No Op Ad is allowed in I/O commands. The bits which normally designate the channel of the Op Ad are part of the $I / O$ Op code. Bits which normally designate the sector of the Op Ad are used to designate the num. ber of characters (words) to be input (output), less one. Execution time for all $1 / 0$ is variable.

| Mnemonic | Op Code | Interpretation | Comments |
| :---: | :---: | :---: | :---: |
| WEOT | 4400 | Write FADAC to Ext Dev in Oct (TT) |  |
| WFOT | 4420 | Write FADAC to FADAC in Oct (TT) |  |
| WEOF | 4440 | Write FADAC to Ext Dev in Oct (FD) |  |
| WFOF | 4460 | Write FADAC to FADAC in Oct (FD) |  |
| WEDT | 4500 | Write FADAC to Ext Dev Dec (TT) |  |
| WFDT | 4520 | Write FADAC to FADAC Dec (TT) |  |
| WEDF | 4540 | Write FADAC to Ext Dev Dec (FD) |  |
| WFDF | 4560 | Write FADAC to FADAC Dec (FD) |  |
| WE6 | 4600 | Write FADAC to Ext Dev 9.6 |  |
| WF6 | 4620 | Write FADAC to FADAC $\boldsymbol{\alpha}$-6 |  |
| WES* | 4700 | Write FADAC to Ext Dev $\alpha-5$ | Output is by word only. Flag required. |
| WF5* | 4720 | Write FADAC to FADAC $\alpha-5$ | Output is by word only. Flag required. |
| WE4 | 4700 | Write FADAC to Ext Dev $\alpha-4$ | Output is by character only. Flag mugt be omitted. |
| WF4 | 4720 | Write FADAC to FADAC $\alpha-4$ | Output is by character only. Flag must be omitted. |
| REO | 5400 | Rd Ext Dev in Oct |  |
| R TO | 5420 | Rd Tape Dev in Oct |  |
| RKO | 5460 | Rd Keyboard in Oct |  |
| RED | 5500 | Rd Ext Dev in Dec |  |
| RTD | 5520 | Rd Tape Dev in Dec |  |
| RKD | 5560 | Rd Keyboard in Dec |  |
| RE6 | 5600 | Rd Ext Dev in $\alpha-6$ |  |
| R T6 | 5620 | Rd Tape Dev in $\alpha-6$ |  |
| RM6 | 5640 | Rd Mag Tape in $\%$-6 |  |
| RK6 | 5660 | Rd Keyboard in 0.6 |  |
| RES | 5700 | Rd Ext Dev in 0.5 |  |
| RT5 | 5720 | Rd Tape Dev in 0.5 |  |
| RM5 | 5740 | Rd Mag Tape in $\alpha-5$ |  |
| RK5 | 5760 | Rd Keyboard in $\alpha$ - 5 |  |

SUMMARY OF CONTROL BUTTON ADDRESSES

| Controi Button | Address |
| :--- | :--- |
| Sample Matrix | 000.00 |
| Test | 000.01 |
| Set Up | 000.02 |
| Recall | 000.03 |
| Send | 000.04 |
| Compute | 000.05 |
| Trig | 000.06 |
| Receive | 000.07 |

SUMMARY OF REGISTER ADDRESSES

## Register

R
Q
D. 0
D. 1

A

L
N

Address
142. (00 to 17)
152. (00 to 17)
160.00
160.01
170.00
172.00
174. 00

TIMING: $\left\{\begin{array}{l}1 \text { word time }=78 \text { micro-sec } \\ 1 \text { drum revolution }=10 \mathrm{milli}-\mathrm{sec}\end{array}\right.$

MAXIMUM I/O RATES: $\left\{\begin{array}{l}4000 \text { characters/sec in Alpha } 6 \\ 3550 \text { characters/sec in Alpha } 5 \\ 650 \text { characters/sec in Oct or Dec }\end{array}\right.$

WORKING STORAGE DESIGNATIONS

## SUMMARY OF MEMORY DESIGNATIONS

Memory I - Channels 00.76
Memory II - Channels 100-136
Memory III - Channels 300-336

Size of WS
4
$12\left\{\begin{array}{c}70-76 \\ 110-116 \\ 130-136\end{array}\right.$
16

WS-Channels

70-76

40-76

