

digital . INTEROFFICE MEMORANDUM

SUBJECT: techniques, cute ideas, and programming tricks for the PDP-11.
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 The Buffer

DEPARTMENT: Programming

The following is a collection of items concerning all sorts of odds and ends about programming the PDP-11.

- One of the features of the PDP-11 is the ability to trap on various conditions such as illegal instructions, reserved instructions, power failure, etc. However, if ~~any of~~ the trap vectors are not loaded with meaningful information, the occurrence of any of these traps will cause unpredictable results. By ~~using~~^{using} the following, it is possible to avoid these problems as well as gain meaningful information about any unexpected traps ~~which~~^{that} ~~may (and will)~~ occur. This technique, which makes it easy to identify the source of a trap, is to load each unused trap vector with:

```
. = trapaddress
.WORD .+2,0
```

This will load the first word of the vector with the address of the second word of the vector (which contains a halt). Thus, for ~~instance~~^{example}, a halt at location 6 ~~implies~~^{means that} a trap through the vector at location 4 has occurred. The old PC and status may be examined by looking at the stack pointed to by register 6.

The trap vectors of interest are:

| <u>Vector location</u> | <u>halt at</u> | <u>vector meaning</u> |
|------------------------|----------------|---|
| 4 | 6 | ✓ bus error, illegal instruction, ✓ stack overflow, non-existent memory, non-existent device, word at odd address, etc. |
| 10 | 12 | executed reserved instruction |
| 14 | 16 | trace trap instruction (000003) executed or T-bit set in status word (used by ODT) |

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| (Cont'd) <u>Vector location</u> | <u>halt at</u> | <u>vector meaning</u> |
|------------------------------------|----------------|-------------------------------|
| 20 | 22 | IOT executed (used by IOX) |
| 24 | 26 | Power failure or restoration |
| 30 | 32 | EMT executed (used by FPP-11) |
| 34 | 36 | TRAP executed |

2. Cute instructions and tricks .

Note: REG refers to a register in general.

a. TST (REG)+

Add two to a register. Use with care since condition codes are clobbered, the register's contents must be even, and the value cannot be outside addressable memory.

b. TST -(REG)

Subtract two from a register. Same cautions as (a).

c. CMP (REG)+,(REG)+

Add four to a register. Same cautions as (a).

d. CMP -(REG),-(REG)

Subtract four from a register. Same cautions as (a).

e. JSR REG,XXX
BNE ABC

XXX: SEZ
RTS PC

This is a very useful tool, in other words, use the condition codes to pass two valued parameters or flags upon subroutine returns.

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4. An example of recursion.

This problem is to scan through a string of characters containing parenthesis nesting to find matching parentheses.

E.G. - find the right paren in the following string which matches the first occurrence of a left paren.

B=A(EXP(A+B(15,12)),)

Recursive form:

```

A:   JSR   R7,@R7           ;CALL B RECURSIVELY
B:   MOVB  (REG)+,R0        ;PICK UP NEXT CHARACTER
     CMPB  #'(',R0         ;IS IT A "("?
     BEQ   A                ;YES, TRY AGAIN
     CMPB  #'),R0         ;IS IT A ")"?
     BNE   B                ;NO, RELOOP
     RTS   R7              ;DO A POP JUMP

```

To use the above, B is entered with REG pointing to the character following the first open parenthesis. B is called with a JSR R7,B. The recursive form is 9 words long.

A non-recursive equivalent is:

```

B:   CLR   R1
B1:  MOVB  (REG)+,R0
     CMPB  #'(',R0
     BNE   B2
     INC   R1
B2:  CMPB  #'),R0
     BNE   B1
     DEC   R1
     TST   R1
     BGE   B1
     RTS   R7

```

The non-recursive example is 13 words long.

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5. Another cute trick

```

    INCB  @(R1)      (and)    INCB  @(R1)+
    TSTB  @(R1)      TSTB  @-(R1)

```

are equivalent except that the first is four words long
and the second is two words long.

6. Pseudo-assembly options

The following example shows how to make assembly options
available or an assembler not having explicit optional
assembly pseudo-ops. This neat nasty depends on the fact
that when two pieces of code are loaded in the same place,
the last one (usually!) takes precedence.

Assume the existence of two short subroutines A and B.

```

A:  INCB  @R1          B:  INC    R5
    TSTB  @R1          MOVB  SP, R0
    BPL   -4           RTS   R5
    MOVB  @R0,R2
    RTS   R5

```

Assume that subroutine B is to be made optional. The program
may be arranged thus:

```

OPT  =  -1
TMP  =
B:   INC    R5
     MOVB   SP, R0
     RTS   R5
LG   =     ·-TMP
     =     LG & OPT + TMP
A:   INCB   @R1
     TSTB   @R1
     BPL   -4
     MOVB   @R0,R2
     RTS   R5

```

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6. Pseudo-assembly options (Cont'd)

If OPT is set to -1, subprogram B will be assembled normally,
If OPT is set to 0, B will be overlaid by A.

Cautions:

- 1) This is not a true conditional assembly thus symbol redefinition may not be done (i.e. two subprograms named B) without some effort.

For example:

```

OPT = -1
B:           ; label definition
TMP = .
. = xx & -OPT +TMP2
  INC R5
  MOVB SP, R0.
  RTS
LG = .-TMP
  = LG & OPT+TMP
  INCB R1
.
.
.

```

- 2) The actual amount of code to be loaded will not change, only location where it will be loaded.

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7. I/O buffering with IOX.

This is a relatively obscure method to double buffer input and output. There are two input buffers (I1, I2) and two output buffers (O1, O2). This is a process which processes an input buffer and in the process loads an output buffer. The process uses pointers to the buffers and buffer headers which are set by the I/O routines. Thus the process does not know (or care) which buffer is being worked upon.

```

PC=%7
BEGIN:  (do I/O resets, inits, etc.)
        .
        .
        .
        IOT          ;read into I1
        .WORD I1
        .BYTE READ,INSLOT
        MOV #A,-(6)  ;initialize stack
B:      JSR PC,@(6)+ ;do I/O
        .
        .   Perform processing.
        .
        BR B        ;do it again, etc.
;
;END OF MAIN LOOP
;
;I/O CO-ROUTINES
;
A:      IOT          ;read into I2
        .WORD I2
        .BYTE READ,INSLOT
        .
        .   set parameters to process I2,O1.
        .
        JSR PC,@(6)+ ;return to process
        IOT          ;write from O1
        .WORD O1
        .BYTE WRITE,OUTSLOT
        IOT          ;read into I1
        .WORD I1
        .BYTE READ,INSLOT
        .
        .   set parameters to process I2, O2.
        .

```

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```

JSR PC,@(6)+      ;return to process
IOT               ;output from 02
.WORD 02
.BYTE WRITE,OUTSLOT
BR. A             ;go read into I2

```

The above is a good example of the usage of the co-routine call form of JSR:

```

JSR PC,@(6)+

```

which does a jump to the address specified on top of the stack and replaces that address with the new return address.

8. More cute tricks

a. CMPB (R6)+, (REG)+

will increment R6 by two and increment REG by one. This can be extended to:

| | | |
|--------------------|------|-----|
| | 1000 | 400 |
| CMPB (R6)+, -(REG) | 1002 | 377 |
| CMPB -(R6), (REG)+ | 1000 | 400 |
| CMPB -(R6), -(REG) | 777 | 377 |

This trick depends on the fact that the stack register is always incremented or decremented by two in auto-increment/decrement mode. The same cautions as in section 2a apply. Also, REG is in the range 0-5 only.

b. The above can be generalized slightly.

```

CMPB (REG1)+, (REG2)+
CMP (REG1)+, (REG2)+ , ad nauseum

```

This allows two separate registers to be incremented (or decremented, e.g. CMP (REG1)+, -(REG2)) at the same time by the same value in one instruction. This is sometimes useful in table scanning algorithms, etc. The same cautions apply as in section 2a.