

# RSX-11D System Manager's Guide

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# PREFACE

# **0.1** MANUAL OBJECTIVES AND READER ASSUMPTIONS

The RSX-llD System Manager's Guide is intended for the individual who has overall responsibility for the operation of an RSX-llD system.

The system manager must have knowledge of RSX-llD in order to benefit from this guide. In particular, the system manager should consider the following documents prerequisites:

RSX-11D User's Guide, RSX-11D Task Builder Reference Manual, RSX-11D System Generation Reference Manual.

# 0.2 DOCUMENT STRUCTURE

This manual comprises the following chapters:

Chapter 1 Introduction to the manual and the significance of UICs and passwords,  Chapter 2 Operation of and output from task accounting,  Chapter 3 Use of spooling, despooling, and input spooling,  Chapter 4 Checkpointing,  Chapter 5 Global libraries and common areas,  Chapter 6 Core dump analysis,  Chapter 7 Time-scheduled partitions,  Chapter 8 Error logging,  Chapter 9 System management notes,  Chapter 10 Executive debugging aids,  Chapter 11 Parity support,  Appendix A Sample core dump analyzer output,  Appendix B TKB15.CMD file,		
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Appendix C Error logging sample output and data record formats.

Appendix D 11/70 parity error response

# 0.3 ASSOCIATED DOCUMENTS

Refer to the RSX-llD Documentation Directory, order number DEC-ll-OXUGA-D-D for a description of associated RSX-llD documents and their readerships.

#### CHAPTER 1

#### INTRODUCTION

The RSX-llD system manager is responsible for system features that fall into the following three categories:

- 1. Coordination and accounting of the multiuser environment,
- 2. Control of system operation,
- 3. System reliability monitoring.

The coordination of the multiuser environment consists of controlling the assignment and use of user identification codes (UICs) and their associated passwords as described in Section 1.1. UICs and passwords uniquely identify system users and establish their access rights. UICs also provide the basis for the accumulation of accounting statistics on executing tasks. Accounting details system use on a task-by-task basis, as described in Chapter 1.

Control of system operation consists of such features as output spooling and despooling of print files, input spooling of card files, the operation of checkpointing, and the implementation and use of global libraries and common areas.

Output spooling and despooling improves system operation by eliminating individual task contention for printers, teleprinters, and other serial output devices. Input spooling allows tasks to process card files from the disk rather than requiring them to wait for card reader availability. Output spooling and input spooling are detailed in Chapter 3.

Checkpointing is the mechanism used to aid high priority tasks to obtain the memory they require to execute as soon as possible after they are requested. Lower priority tasks are checkpointed to disk to allow higher priority tasks to run. Chapter 4 describes this process.

Global libraries are collections of routines that can be shared by many simultaneously active tasks. Common areas are data areas in memory that can be shared by many simultaneously active tasks. Both global libraries and common areas are loaded into memory dynamically when needed and their memory is released when they are no longer required as described in Chapter 5.

System reliability monitoring includes the core dump analyzer and the hardware error logging facility.

The core dump analyzer provides a means of interpreting the state of the software when a system crash occurred. It is described in Chapter 6.

Chapter 7 contains a description of the algorithm used in time-scheduled partitions. This information should be taken into consideration when configuring a system that is to contain time-scheduled partitions.

The error logging facility allows the system manager to obtain a detailed and/or summary report of peripheral device errors. Error logging reports can be produced for specified time frames and for specified devices or volumes, as described in Chapter 8.

The system manager also is responsible for optimizing system operation. Chapter 9 contains system management notes for that purpose.

Executive trace and ODT are provided in RSX-11D for the debugging of user-initiated modifications to the Executive. These features are described in Chapter 10.

Chapter 11 details the software support for handling hardware parity errors. This information should be used in conjunction with the error logging description in Chapter 8.

#### 1.1 UICS AND PASSWORD PROTECTION

Under RSX-11D, the key to controlling the multiuser environment is UIC and password assignment. Each system user should have a unique UIC. Each UIC is associated with a UFD on the system volume (SY). The UFD command, described in Chapter 6 of the RSX-11D User's Guide, is used to place UFDs on directoried devices. Any UFD placed on the SY volume also serves as a UIC. A UIC is not valid unless the same UFD exists on SY.

The group code portion of the UIC is used to determine the privileges extended to the user once the user has logged onto the system. Any group code of 10 (octal) or less indicates a privileged UIC. Its owner is granted access to system functions as a privileged user. Group codes greater than 10 (octal) indicate a general user with restricted access to system functions. The terms general user and privileged user are defined in the RSX-11D User's Guide. The following are examples of privileged UICs: [1,12], [5,20] and [3,3].

UICs are specified in the HELLO command that allows a user to log onto the system. The HELLO command is detailed in Chapter 5 of the RSX-llD User's Guide. The system validates UICs typed in the HELLO command to determine that the UIC is valid for the system.

Each UIC can have a password associated with it. When passwords are used, once the user types the HELLO command, a request is printed on the terminal as follows.

#### PASSWORD>

The user types the password at this point. The password is not printed on the terminal as it is typed. If the password corresponds to the UIC specified in HELLO, the user gains access to the system.

A UIC created by the UFD command does not have a password associated with it. The PWD command to MCR is used to assign a password if one is desired. The PWD command is detailed in Chapter 6 of the  $\frac{RSX-lld}{USER}$  User's Guide.

When assigning UICs to system users, the system manager. should take into consideration the ways that UICs might be used. For example, if accounting is to be performed on tasks executing in the system, the UIC serves as the account number. The system manager should ensure that a meaningful system of assigning UICs is developed.

# INTRODUCTION

The RSX-11D system as distributed by DIGITAL contains privileged and nonprivileged UICs. These UICs are not password protected. The system manager can log on under any one of the privileged UICs to perform a system generation. The system manager can use the PWD command to assign passwords to any of the distribution system UICs as desired.

# NOTE

All functions described in this manual must be performed by a privileged user.

#### CHAPTER 2

#### ACCOUNTING

RSX-11D accounting provides the system manager with a means of knowing how the system resources are being used by executing tasks. Accounting provides statistics on a task-by-task basis and also accumulates totals for each UIC (account). The user specifies at task build time whether accounting is to be performed for the task. Therefore, accounting information is accumulated only when the information is desired.

#### 2.1 INTRODUCTION

Accounting statistics provide the following information for a task:

- The amount of computer memory used by a task in terms of 32-word blocks,
- 2. The elapsed execution time; i.e., time in memory,
- Central processor time used by the task indicated in clock ticks,
- 4. The number and types of accesses to each peripheral device used by the task.

Summary statistics contain totals for each UIC under which tasks specified for accounting are run. Figure 2-1 provides an example of task statistics and associated statistics.

Elapsed time is total time starting when the task is first made active and ending when the task exits. The time between a suspend request and a resume request is included in the elapsed time. Interrupt and Executive overhead time for the task are included in both central processor time and elapsed time.

Tasks that are requested more than once have a set of statistics for each time they execute. Each version of a multiuser task has a separate set of statistics.

ACCOUNTING REPORT FOR 19-FEB-75 TIME: 17:43:58

ACCOUNTS FOR UIC [5,73]

TASK: ...FLX STARTED AT: 19-FEB-75 17:43:49:22

CORE USED: CPU TIME: +415 +2

ELAPSED TIME: 0-0-0 0:0:1:12

DEVICE ACCESS SUMMARY:

DEVICE: UNIT: READS: WRITES: OTHER:

DP 0 0 0 0 1 1 1 1 1

SUMMARY FOR UIC [5,73]

CORE USED: CPU TIME: +415 +2

ELAPSED TIME:  $\emptyset-\emptyset-\emptyset$   $\emptyset:0:1:12$ 

DEVICE ACCESS SUMMARY:

+1 READS +1 WRITES +4 OTHERS

Figure 2-1 Accounting Output

Accounting provides data that can be used for billing based on computer usage. Because all accounting statistics are accumulated using UICs as account numbers, each account to be billed should be given a specific UIC or a related set of UICs so that statistics related to that account can be identified readily.

In addition, the system manager should be aware of how the system uses UICs for various functions. For example, a task run by a general (nonprivileged) user executes under the UIC that was specified in the HELLO command. On the other hand, if the same task is run by a privileged user, it executes under the UIC for which it was built or the UIC specified using the /UIC= switch with the RUN command. The UIC for which the task was built is the one given in the UIC= command to the task builder or the value to which it defaulted. When using batch, the /ACC= option on the \$JOB command dictates the UIC under which tasks within that job execute, as described in the RSX-11D Batch Reference Manual.

Tasks which run at a priority above the priority of ACCLOG should NOT be built accountable.

Privileged tasks should not be made accountable if there is any possibility that they might be terminated for time limit by ACCLOG while modifying system data. It is NOT recommended that any distributed task be built accountable, except for the non-privileged system utilities. A batch file is provided for this purpose (see Section 2.3.4).

#### 2.2 FUNCTIONAL DESCRIPTION

The actual statistics for each task are accumulated by a set of system subroutines and placed in the task's header area. Storing information in the task header means that statistics for a task usually need to be written to a disk file once only; i.e., when the task exits.

Two accounting tasks are required to 1) log the accounting information into a permanent disk file, and 2) analyze the accounting and produce the final output files.

# 2.2.1 Logging Task -- ACCLOG

The logging task, ACCLOG, takes statistics from the task header when the task exits and writes that information in a disk-resident file.

# NOTE

ACCLOG is a privileged, high-priority task that must be running at a priority of 235 or greater for accounting to take place.

ACCLOG creates two files on SY under UFD [1,5] in which statistics are stored. First, it opens and uses a file named ACCTS1.DAT and writes statistical data into it. When the analysis task runs, as described below, it sets global event flag 61 to indicate that it wants to process ACCTS1.DAT. ACCLOG tests the same flag. When it is set, ACCLOG stops updating ACCTS1.DAT and opens another file named ACCTS2.DAT. ACCLOG also sets global event flag 62 to indicate that analysis can proceed.

When the analysis task runs again to process the second file, it sets a flag to indicate this fact. At that point, ACCLOG creates a file named ACCTS1.DAT. ACCLOG maintains two files: one to store statistics in and one to be processed by the analysis task. ACCLOG alternates between creation of ACCTS1.DAT and ACCTS2.DAT.

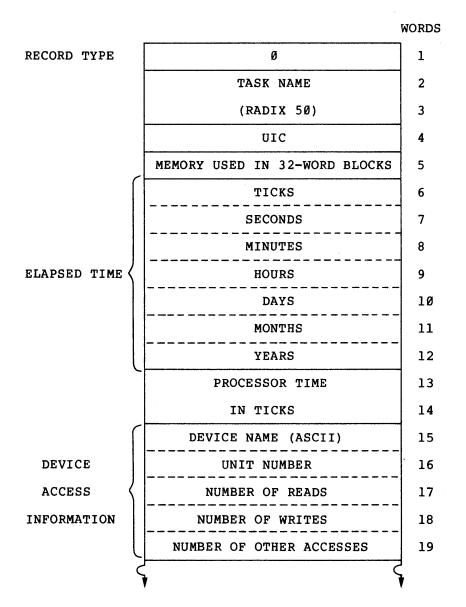
# 2.2.2 Analysis Task -- ACCRPT

The analysis task, ACCRPT, analyzes the data file produced by ACCLOG. ACCRPT can be installed and run or scheduled at the user's discretion; unlike ACCLOG, it need not be resident constantly. When ACCRPT executes, it analyzes all information in the data file.

ACCRPT produces two files: a report file (ACCTS.RPT) that can be listed on a line printer and a data file (ACCTS.TOT) that can be processed by a user-written task, if desired. All files are under UFD [1,5]. Additionally, ACCRPT deletes the raw data file (ACCTS1.DAT or ACCTS2.DAT).

The records in ACCTS.TOT are variable in length and contain octal information providing task and UIC statistics. Figures 2-2 and 2-3 illustrate the fields and their lengths for task-specific and summary records, respectively.

Each time ACCRPT executes, it resumes ACCLOG and sets an event flag to tell ACCLOG to open a new file. When the old file is closed, ACCRPT reopens it and begins processing the data on a UIC and task basis. When finished, ACCRPT deletes the input file (ACCTS1 or ACCTS2).



# NOTE

- The minimum record size is 14 words.
   The maximum is 110 words.
- 2. The device access information is repeated once for each device used by the task.
- Except where noted, all information is in binary.

Figure 2-2 Format of Task Statistics Record (ACCTS.TOT)

		WORDS
RECORD TYPE	1	] 1
	UIC	2
	MEMORY USED IN	3
_	32-WORD BLOCKS	4
	TICKS	5
	SECONDS	6
	MINUTES	7
ELAPSED TIME	HOURS	8
	DAYS	9
	MONTHS	10
	YEARS	11
	PROCESSOR	12
	TIME	13
	IN TICKS	14
	TOTAL NUMBER	15
DEVICE .	OF READS	16
ACCESS	TOTAL NUMBER	17
INFORMATION	OF WRITES	18
	TOTAL NUMBER OF	19
L	OTHER ACCESSES	20

# NOTES

The record containing UIC statistics is 20 words in length.

All information is in binary.

Figure 2-3 Format of UIC Summary Records (ACCTS.TOT)

Each time ACCRPT runs, it creates new output files rather than updating the existing ones, thereby making it easier for the user to keep track of accounting information. It is the system manager's responsibility to delete these files when they are no longer needed.

If a system crash occurs, ACCRPT is not able to process the raw data file (either ACCTS1.DAT or ACCTS2.DAT) that was being used by ACCLOG when the failure occurred. However, the file remains on the disk and can be processed by the user. Figure 2-4 illustrates the record lay-out for the raw data file.

#### 2.3 OPERATIONAL INFORMATION

Operational information describes 1) the procedures for initiating and terminating accounting and 2) the task builder and task installation requirements of the task to be accounted. Optionally, accounting can be performed for utility programs as described in this section.

#### NOTE

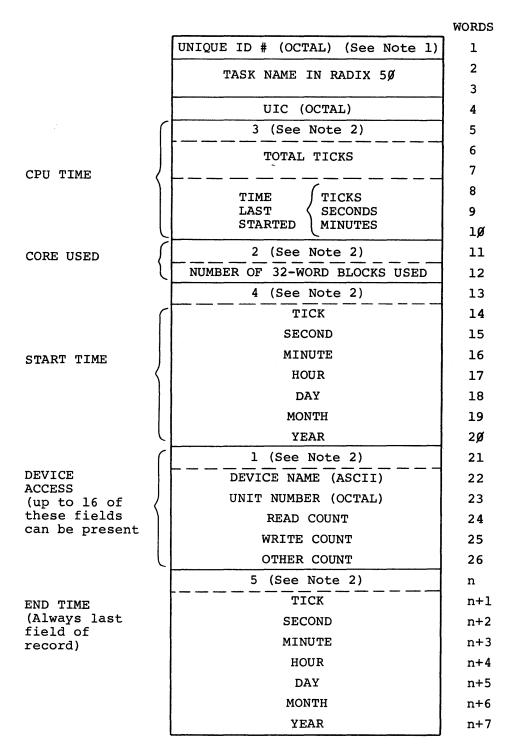
If the system clock is not running at 60 (decimal) ticks per second, the analysis task (ACCRPT) must be rebuilt before it is used.

To rebuild ACCRPT, first edit the file [11,26] RPTBLD.CMD by locating the following line.

#### GBLDEF=TICKS:74

Change the 74 (octal) to the correct number of ticks per second in octal. Then issue the following string to TKB.

TKB>@[11,26]RPTBLD.CMD



#### NOTES

- Uniquely identifies two tasks with the same UIC and name.
- 2. Field type code.

Figure 2-4 Record Format for Raw Data File

# 2.3.1 Initiating Accounting

Use the following procedures to initiate accounting.

- If it has not been done already, mount the system disk using the MOU command and update the time and date using the TIM command.
- Run ACCLOG from a privileged UIC using the following commands.

MCR>RUN ACCLOG

At this point, the system is ready to start accumulating statistics. If tasks designated for accounting execute before ACCLOG runs, no statistics accumulate for them.

3. When accounting statistics are desired, install and run ACCRPT using the following commands.

MCR>INS[11,1]ACCRPT
MCR>RUN ACCRPT

4. When ACCRPT terminates, use PIP to print ACCTS.RPT. Additionally ACCTS.TOT can be printed using the DMP utility program. All files are on UFD [1,5].

# 2.3.2 Terminating Accounting

Two separate tasks are provided to terminate accounting. The first, ACCABT, allows the operator to abort accounting for the entire system. The second, ACCOFF, allows the operator to terminate accounting in an orderly fashion.

When ACCABT runs, it aborts ACCLOG and sets an internal switch to indicate that no further accounting is to be performed. It makes no attempt to finish the logging of information for tasks that have not terminated nor does it attempt to analyze information in the raw data file.

On the other hand, ACCOFF first runs ACCRPT to preserve all accumulated data and then aborts ACCLOG. Normally, ACCOFF is used to perform an orderly shut down of accounting. ACCABT is used only in emergency situations to abort accounting.

Regardless of which method is used to terminate accounting, it can be restarted at any time by following the procedures in Section 2.3.1.

Use one of the following procedures to terminate accounting.

 To perform an orderly shut down, install ACCRPT and ACCOFF and run ACCOFF using the following commands.

> MCR>INS [11,1]ACCRPT MCR>INS [11,1]ACCOFF MCR>RUN ACCOFF

2. To abort accounting, type the following commands.

MCR>INS [11,1]ACCRPT MCR>RUN ACCABT

Both procedures cause a normal termination message to be printed on the terminal. Do not under any circumstances use ABO to abort ACCLOG, as it causes misinformation to be left in the system data base and could cause serious problems. Always use one of the above first.

# 2.3.3 Task Requirements (/TA and /TIM)

All tasks for which accounting information is to be accumulated must have accounting specified for them during task building. The /TA switch is used in the command string to specify accounting as described in the RSX-11D Task Builder Reference Manual.

#### NOTE

The default for the Task Builder can be changed so that /TA becomes the default. Refer to Chapter 9.

Another option for controlling accounted tasks is the /TIM switch that can be used on the INS command to MCR. The /TIM switch allows the user to specify the amount of central processor time that a task can use. If the specified amount is exceeded, the task is aborted and its accounting information is recorded. The /TIM switch has the following format.

#### /TIM=nx

- n is a decimal number indicating the number of time units (x) of central processor time that the task is allowed.
- x indicates the type of time unit and can be one of the following.

H = hour

M = minutes

S = seconds

T = ticks

The amount of time specified must be no greater than 24 hours. For example, 24H, 1440M, and 86400S are the maximum amounts for hours, minutes, and seconds, respectively.

If /TIM is not included and accounting is specified for the task, the maximum amount of processor time allowed is five minutes. If the task is still running after five minutes of processor time has been used and /TIM was not specified, the system aborts the task.

# 2.3.4 Rebuilding Utilities to be Accounted

The system manager can rebuild all the RSX-11D utility programs, MACRO-11, and the FORTRAN compiler to have their use accounted. To simplify the rebuilding process, a batch command file is included in the system to rebuild all utilities, MACRO-11, and FORTRAN. The file is located under UFD [11,26] and is called UTLACCBLD.BIS. After UTLACCBLD.BIS is processed by batch, the new task images are located under UFD [11,26].

Use the following procedures to rebuild the utilities, MACRO-11, and FORTRAN to be accounted.

- 1. Redirect SY to the device containing the system object files.
- Type the following command to request batch and process the file.

# MCR>BAT [11,26]UTLACCBLD

If some, but not all, are to be accounted, the utility tasks, MACRO-11, and/or the FORTRAN compiler can be rebuilt individually. Table 2-1 lists the individual task build files to make a utility, MACRO-11, or FORTRAN accounted. All of the files are under UFD [11,26].

Table 2-1
Task Build Files for Accounted System Tasks

ASSOCIATED MODULE
FLX
PIP
SLIPR
Librarian
MACRO
Task Builder
Editor
FORTRAN

# 2.3.5 Disk Storage Statistics

The /TB switch for PIP can be used to obtain statistics on the amount of disk storage that is used for a UFD. To determine the amount of storage used by a UFD on a specific device, type the following command to PIP.

PIP>dev:[ufd]/TB

dev is the device for which statistics are desired.

[ufd] is the UFD whose usage is to be determined.

To determine the amount of disk storage used for all UFDs on a device, type the following command to PIP.

PIP>dev:[\*,\*]/TB

# 2.3.6 Suggestions

- 1. Run ACCRPT only when necessary. Random running of ACCRPT results in many report files (ACCTS.RPT and ACCTS.TOT). Normally, ACCRPT needs to be run only when the system is to be shut down.
- When the system is to be shut down, run ACCOFF rather than running ACCRPT and ACCOFF individually. ACCOFF automatically calls ACCRPT.

- Use ACCTS.TOT as a basis for billing or further analysis. ACCTS.TOT is an octal file that can be processed by user-written tasks using FCS.
- 4. If accounting is to be used for billing purposes, ensure that UICs are assigned to individual accounts in a meaningful and easily processed fashion.
- 5. If many .TOT files are to be processed to produce one set of output, combine the .TOT files into one file using PIP. See the RSX-ll Utilities Procedures Manual.

# 2.4 ERROR MESSAGES

ACCLOG, ACCRPT, ACCOFF, and ACCABT issue error messages to the user. All of the accounting tasks assume that the SY handler is loaded, the device is mounted with correct access, and the device is not full. If this is not the case, error messages are printed on the terminal and the user must correct the situation.

# 2.4.1 ACCLOG Error Messages

ACCLOG -- FAILED TO OPEN NEW FILE, RUN ACCABT.

Use the following procedures to correct the situation:

- 1. Ensure that UFD [1,5] is present on the system disk,
- 2. Ensure that ACCLOG was running under a privileged UIC so that it has correct access rights,
- If neither 1 or 2 above was the cause of the problem, check the device,
- 4. Run ACCABT,
- 5. Rerun ACCLOG.

ACCLOG -- IO ERROR; SOME DATA LOST

This message indicates that write errors have occurred and data has been lost for a task's account. Use the following procedures:

- 1. Check the file and device to see if they are usable,
- If the errors persist, run ACCOFF to preserve accumulated statistics and turn accounting off until the cause of the problem is fixed.

# 2.4.2 ACCRPT Error Messages

ACCRPT -- FAILED TO OPEN NEEDED FILE, CANNOT PROCEED

#### filename

If the filename specified in the message is ACCTS1.DAT or ACCTS2.DAT, use the following procedures:

- 1. Ensure that ACCTS1 or ACCTS2 exists under UFD [1,5],
- Ensure that ACCRPT was run under a privileged UIC so that it has the correct access rights,
- 3. Ensure that the file is closed properly and not locked,
- If none of the above is the cause of the problem, check the device,
- 5. Rerun ACCRPT.

If the filename specified in the message is ACCTS.TOT, ACCTS.RPT, or ACCTS.TMP, use the following procedures:

- Ensure that ACCRPT is running under a privileged UIC so that it has correct access rights,
- 2. Check the device,
- 3. Rerun ACCRPT.

If rerunning ACCRPT does not produce the desired results, use the DMP utility program to retrieve all data files except ACCTS.RPT. Use PIP to retrieve ACCTS.RPT or try to reanalyze the file.

ACCRPT -- IO ERROR, SOME DATA LOST

#### filename

If the filename specified in the message is ACCTS1.DAT, ACCTS2.DAT, or ACCTS.TMP, some input data has been lost. One task's account will be missing from the final analysis. Ensure that the file is not corrupted or the device is not malfunctioning.

If the error persists, stop analysis and print or dump all files under [1,5] to retrieve data and locate the problem.

If the filename specified in the message is ACCTS.TOT or ACCTS.RPT, a write error has occurred. Some data will be missing from the file. Data can be retrieved by comparing the ACCTS.TOT file with ACCTS.RPT unless both files are reporting problems.

If errors persist, print or dump all files under [1,5] to retrieve data and locate the problem.

# 2.4.3 ACCOFF Error Messages

ACCOFF -- PLEASE RUN ACCRPT TO PRESERVE DATA

The request from ACCOFF to run ACCRPT failed; probably ACCRPT was not installed. ACCOFF does not terminate accounting until all data is preserved. Type the following commands.

INS [11,1]ACCRPT RUN ACCRPT

ACCOFF -- ABORT FAILED; ABORT ACCLOG FROM CONSOLE

For some reason ACCOFF was unable to abort ACCLOG. Use the ABO command to terminate ACCLOG, as follows.

ABO ACCLOG

# 2.4.4 ACCABT Error Message

ACCABT -- ABORT FAILED; ABORT ACCLOG FROM CONSOLE

For some reason ACCABT was unable to abort ACCLOG. Use the ABO command to terminate ACCLOG, as follows.

ABO ACCLOG

#### 2.5 ACCOUNTABLE TASKS AND SAVING THE SYSTEM

The SAVE command (see the <u>RSX-11D User's Guide</u>, Section 6.23) preserves on disk a copy of the memory part of an RSX-11D system. The disk can thereafter be bootstrapped to restart the system.

A SAVE may be performed only when the system is quiescent, as defined in the section referred to. This includes, for example, no loading of tasks or peripheral transfers in progress. However, an accountable task can be resident and if a SAVE is performed, the time limit will be recorded and, on the disk being bootstrapped, will continue from where it was before the save. This can result in terminations for time limit for tasks that seem not to have run for very long.

#### CHAPTER 3

# SPOOLING, DESPOOLING AND INPUT SPOOLING

RSX-11D provides three types of file spooling:

- 1. Automatic output spooling to disk,
- 2. Automatic despooling of disk files queued for printing,
- 3. Input spooling from a card reader to disk.

Automatic output spooling eliminates program contention for line printers, teleprinters, and other serial output devices. Because programs do not have to wait to use a serial device, higher system throughput is achieved.

The need to wait for a particular device is removed by temporarily redirecting the output intended for the serial device to a disk file. This temporary redirection is defined as automatic output spooling. The files redirected to the disk are later produced on the original device. This process is defined as despooling.

RSX-llD provides several levels of automatic output spooling/despooling. On the simplest level, the automatic output spooler can queue files for the single device despooler which prints them one at a time on a first-in/first-out basis. Only one output device is used at a time.

On a more sophisticated level, the automatic output spooler and the console user both can queue files for printing. In this case, the multiple device despooler is used. It sends queued files for printing on a priority basis. Up to 26 devices can be used for despooled output simultaneously.

The input spooler eliminates program contention for the card reader and serves as a batch job input device. Card files are read onto a disk and placed under a user-specified UFD.

# 3.1 AUTOMATIC OUTPUT SPOOLING

Automatic output spooling is implemented for use with programs using RSX-llD FCS. It is not operational with programs using the QIO directive. However, both automatic output spooling and programs issuing QIO directives can coexist when using the same output device. The program using QIO directives should adhere to the following procedure:

- 1. Attach the output device,
- 2. Space the form to headmof-form if using a line printer,

- 3. Issue the QIO directives,
- 4. Detach the output device.

If the task issuing the QIO directives is executing under a privileged UIC of [1,n], it must interrogate the status returned from the attach QIO tio ascertain whether the device already is attached. If the device already is attached, the calling task should wait for a significant event and then retry the attach.

# 3.1.1 System Requirements for Automatic Output Spooling

The following three system requirements must be met for automatic output spooling.

 The SY device must have a UFD of [1,4] with read, write, extend and delete privileges for owner, group, system, and world. The following command is used to put the UIC on SY.

MCR>UFD SY: [1,4]/PRO=[RWED,RWED,RWED]

2. A PUD for SP (spooler) must be created during system generation. A DEV directive to create a PUD for SP is included in the system generation command files on the distribution medium. The command has the following format.

DEV=SP

3. The SP PUD is normally assigned to SY during Phase 2 of system generation by the redirect command:

MCR>RED SY=SP

Once the above command has been issued, the system manager has the option of redirecting the SP PUD to any disk device. Redirection can be performed at any time. The device that SP is assigned or redirected to must have a UFD of [1,4] with the same protection as defined above.

4. Both the devices assigned to SY and SP must be write enabled.

# 3.1.2 Program Conventions For Automatic Output Spooling

The following requirements must be met by the MACRO-11 program producing the file to be spooled.

- 1. The file storage region must contain at least a 512-byte output area to contain the output for the file that is targeted for the serial device. It is allocated using the FSRSZ\$ macro. An open failure results on the output file if less than 512 bytes are supplied.
- The program must open and close the output file.
- 3. An open by file identification or an open by filename block must be preceded by a call to .PARSE or by obtaining the directory identification.
- 4. The program must not attempt to regain access to the file once it is closed.

5. The maximum output record size must not exceed 256 bytes.

The FORTRAN compiler automatically adheres to the above procedures within the FORTRAN OTS. However, FORTRAN programs using the DEFINE FILE statement, must ensure that the maximum output record size of 256 bytes is not exceeded.

# 3.1.3 Enabling or Disabling Automatic Output Spooling

Automatic output spooling is enabled or disabled on an output device by typing the SET command to MCR. The syntax of the SET command to enable spooling follows with brackets indicating optionality.

MCR>SET /SP=dev:[,dev:,...,dev:]

/SP= is the command switch for automatic output spooling.

dev: indicates which device(s) are to have files spooled.
dev is a 2-character device name and a 1- or 2-digit
unit number. The device must be both a
carriage-controlled and a record-oriented device, e.g.,
line printer or teleprinter.

Any attempt to set spooling of files for a device that is not a teleprinter or a line printer results in the printing of the following error message on the console.

SET -- SPOOLING NOT SUPPORTED IN THIS RELEASE FOR THIS DEVICE

The syntax of the SET command to disable spooling follows.

MCR>SET /-SP=dev:[,dev:,...,dev:]

/-SP= is the command switch to disable spooling.

The SET command can be used at any time to enable or disable spooling. The SYS /DEV command can be typed to MCR to determine whether spooling is enabled or disabled for a particular device. Any device for which spooling is enabled is designated as follows.

dev: SPOOLED x

x is the designated form type for the device. Refer to Section 3.2.1.5.

# 3.1.4 Automatic Output Spooling Technique

The RSX-llD system performs special functions for programs writing spooled files when the file is opened, as individual records are written, and when the file is closed.

During open, when the output LUN is assigned to the file, the system determines whether the output device is spooled. If it is, the output device is redirected to SP. Normally, a file is created in UFD [1,4] with a 2-character filename indicating the final output device and a file type of .SPR. The file specification has the following format.

SP: [1,4] dv.SPR;n

- dv is the device on which the file is to be printed, e.g.,
  LP.
- n is the file version. It is one greater than any current file having the same filename and file type in UFD [1,4].

FORTRAN programs differ in that they have the following file specification.

SP: [uic] FORnnn.DAT; n

uic is the terminal user's UIC.

nnn corresponds to the logical unit number used to create the file.

During program execution, each successive PUT\$ results in placing the contents of the output buffer into the 512-byte block. When the file storage region block is filled, its contents are written into the file space on the disk.

When the file is closed, the file identification information is taken from the file description block (FDB) and sent to the despooler task using a SEND/REQUEST DIRECTIVE.

# 3.2 DESPOOLING

RSX-11D provides two despooling functions:

- 1. Single device despooler,
- 2. Multiple device despooler.

NOTE

The Multiple device despooler must be used if BATCH is to be run.

The single device despooler services SEND/RECEIVE requests on a first-in/first-out (FIFO) basis. The single device despooler is intended for smaller systems not requiring a system manager's operator interface by means of the OPR command to MCR or a terminal interface command by means of the QUE command to MCR.

The single device despooler has a file name of [11,1] PRT and a task name of SPR....

The multiple device despooler can service up to 26 devices simultaneously. It consists of four tasks as listed in Table 3-1.

#### SPOOLING, DESPOOLING AND INPUT SPOOLING

Table 3-1 Multiple Device Despooler Tasks

FILE NAME	FUNCTIONS	TASK NAME
[11,1]OPR	System manager's operator interface to despooling	OPR
[11,1]QUE	User´s console interface	QUE
[11,1]SPR	Despooler queue manager	SPR
[11,11]SPR2	Multiple device despooler	SPR2

# 3.2.1 Multiple Device Despooler Options

The multiple device despooler has the following options that are not available with the single device despooler:

- QUE and OPR commands to provide user and system manager interface with despooling,
- 2. Up to 26 output devices,
- 3. Despooling of any file on a Files-11 directory device.
- 4. Output by priority,
- 5. Multiple form types,
- 6. Capability to obtain multiple copies of a file,
- 7. Capability to delete or preserve the file after printing,
- 8. Batch output support.

# 3.2.1.1 $\underline{\text{QUE}}$ and $\underline{\text{OPR}}$ Commands to $\underline{\text{MCR}}$ - The $\underline{\text{QUE}}$ command provides the user with a terminal interface with despooling. The $\underline{\text{QUE}}$ command allows the user to perform the following:

- Queue files to be despooled for printing on any teleprinter or printer,
- 2. List files that are queued for despooling,
- Delete (kill) entries from the queue,
- 4. Modify entries in the queue. The following options can be modified for an entry:
  - a. The output device to be used for printing,
  - b. Priority of the entry,
  - c. Form type to be used,
  - d. Number of copies to be printed,
  - e. Whether the file is to be deleted or preserved after printing.

# SPOOLING, DESPOOLING AND INPUT SPOOLING

The despooler handles all files identically regardless of whether they were queued by the automatic output spooler or the QUE command.

The OPR command is the system manager's interface with despooling. It allows the system manager to perform the following:

- 1. Stop output,
- 2. Resume output,
- 3. Resume output from the beginning of the input file,
- Resume output from the last top-of-page encountered in the input file,
- 5. Abort output,
- 6. Begin despooling output,
- 7. Change the form type,
- 8. Indicate that a form type change is to occur and set up the printing of files to test form type alignment,
- Recycle all devices that were being used for despooling prior to a system crash. This option provides a recovery of despooling operations.

The QUE and OPR commands are detailed in the RSX-llD User's Guide.

- 3.2.1.2 <u>Multiple Output Devices</u> Files can be despooled to up to 28 output devices simultaneously. Any printer or teleprinter can be specified as the output device. The output device does not have to be specified for spooling in an MCR SET command. Setting a device for spooling only affects automatic output spooling; it has no effect on the despooling operation. That is, a device does not have to be spooled to have files gueued to it for despooling.
- 3.2.1.3 <u>Despooling of any File</u> Any file on a Files-11 directory device can be designated in a QUE command for despooling. The file to be printed should be other than a binary file, such as a data, source, list, or map file.

An attempt to print a binary file such as an object or task-image file results in a despooler error and an automatic abort of despooling operations for the file.

- 3.2.1.4 <u>Despooling Priority</u> Priority determines the order in which files are despooled. The priority can be specified in the QUE command or the PRINT\$ macro or it can default to the task's running priority. Priority is a decimal number in the range of 1 through 250.
- 3.2.1.5 <u>Multiple Form Types</u> The type of form on which the file is to be printed can be specified by a number in the range of one through seven. The form type value itself is contained in the PUD for the device and can be displayed using the SYS /DEV command. The system is generated with a form type of Ø for every printer and teleprinter.

#### SPOOLING, DESPOOLING AND INPUT SPOOLING

The PUD form type value can be changed dynamically by the system manager using the OPR command. Two switches to the OPR command are provided for this purpose.

- 1. /CHG alters the form type of the specified device to form type 7. Form type 7 is reserved for the despooling of test files to verify the form alignment. Test files can be gueued by using the /TE option of the QUE command.
- /FO:n alters the PUD form type to the value n; n is in the range from Ø through 6.

The system manager determines which form types correspond to the numbers zero through six. The OPR command is used to specify the value for the current form type for a particular device by placing that value in the PUD.

The system manager can assign separate values for different form types or have several form type values correspond to the same type of printer paper. However, the following facts should be considered.

- Files queued to be printed on a particular device with a specified form type are printed only when the PUD for that device has a matching form type value.
- Form type Ø is the form type initially contained in the PUD for each printer or teleprinter; therefore, it is logical to define the most commonly-used printer form as form type Ø.
- 3. Despooling of queued files onto a device with a matching form type is automatic. Differing form types require the system manager to alter the PUD form type using the OPR command.
- 3.2.1.6 <u>Multiple Copies of the Printed File</u> The user can specify the number of copies of a file that are to be printed. The number of copies is a decimal number in the range of 1 through 31.
- 3.2.1.7 <u>Delete/Preserve Indicator</u> The user can specify whether the file is to be deleted after printing or to be preserved for subsequent use.
- 3.2.1.8 Account Number The account number is the UIC of the user who gueued the despool request. The account number determines access rights of a nonprivileged user to modify or delete a gueued request. It also determines which entries are to be listed in response to a Queue List MCR command.

#### 3.2.2 Functions of the Despooler Tasks

QUE, OPR, SPR, and SPR2 are the four tasks used to achieve multiple device despooling. The functions of QUE and OPR are presented in Section 3.2.1.1. SPR and SPR2 are described below.

3.2.2.1 <u>SPR (Queue Manager)</u> - SPR is the queue manager. It receives SEND/REQUEST directives that are directed to the multiple device

despooler and creates an entry in the gueue. The despooler gueue is a file that is maintained on the system disk under UFD [1,4]. The file is created by the gueue manager should no file exist. The file specification for the gueue follows.

SY: [1,4] SPRQUEUE.SYS

SPRQUEUE.SYS is a dynamically expanding file capable of holding up to 2047 (decimal) queued items. The queue manager selects files (queue entries) to be despooled based on the number of devices the multiple device despooler can service simultaneously. It issues a SEND/REQUEST directive to the multiple device despooler. The queue manager is active only when it is queuing an entry or engaged in sending a new file specification to the multiple device despooler.

The selection algorithm for determining which entry to send to the multiple device despooler is illustrated in Figure 3-1.

3.2.2.2 <u>SPR2 (Multiple Device Despooler)</u> - SPR2 is the multiple device despooler task. It writes the files indicated to it by SPR on the designated printer or teleprinter. SPR2 is AST-driven and handles up to 26 devices simultaneously depending on the size of the task when installed.

SPR2 supports one device automatically. For each additional device to be supported, another 460 words must be added to the size of SPR2 when it is installed. The /INC switch is used in the INS command to MCR to increase the task size. Multiply the number of devices to be supported by 460 to determine the appropriate value to use in the /INC switch. For example, if three devices are to be supported, type the following command when SPR2 is installed.

MCR > INS [11,1] SPR2 / INC = 1380 (3 \* 460 = 1380)

### 3.2.3 Interfaces to the Despooling Tasks

RSX-11D provides four interfaces with the despooling tasks:

- Automatic output spooling triggered by the closing of a file that is assigned to a spooled device,
- 2. The PRINT\$ macro call,
- 3. The QUE command to MCR,
- 4. The OPR command to MCR.

Automatic output spooling is operational for both the single and multiple device despoolers.

3.2.3.1 Print\$ Macro Call - The PRINT\$ macro call is the program interface with the despooler. It is operational with both the single and multiple device despooler. The syntax of the PRINT\$ macro is described in the  $\overline{\mbox{IAS/RSX-ll I/O}}$  Operations Reference Manual.

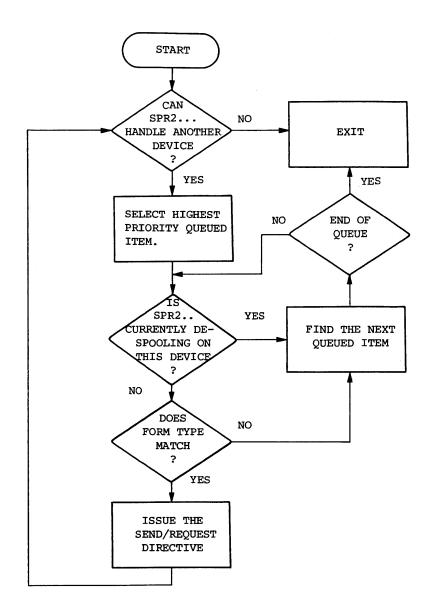


Figure 3-1 Queue Entry Selection Algorithm

The multiple device despooler is capable of using more of the parameters passed to it by PRINT\$ than the single device despooler can. Table 3-2 summarizes the differences.

Table 3-2 PRINT\$ Capabilities of the Single and Multiple Device Despoolers

PARAMETER	SINGLE DEVICE	MULTIPLE DEVICE
Output device Input filename Priority Form type Number of copies Delete/preserve the file	yes yes no no l only delete	yes yes yes yes l through 3l either

If parameters that do not apply for the single device despooler are included in a PRINT\$ macro call to that despooler, they are ignored.

3.2.3.2 QUE and OPR Commands - The QUE and OPR commands to MCR are the user's and system manager's terminal interface to the despooling tasks. Both commands are summarized in Section 3.2.1.1.

This interface is operational only for the multiple device despooler. The system manager should install these tasks only if the multiple device despooler is installed.

### 3.2.4 Priorities of the Despooling Task

All of the despooler tasks run in the GEN partition with a priority of 100. If the system manager wishes to change their priorities, he should consider the following paragraphs.

- 3.2.4.1  $\underline{OPR}$  The priority of OPR should be maintained at a high value. The task must be able to communicate instantaneously with SPR2 (multiple device despooler).
- 3.2.4.2 QUE The QUE task is a multiuser task. Its running priority is not critical and can be altered to any value. However, ideally, its priority should not exceed that of the gueue manager (SPR). Multiple SEND/REQUEST directives to SPR, such as a number of modify or kill QUE commands could cause the pool limit of 100 to be exceeded.
- 3.2.4.3 SPR (Queue Manager) The queue manager should operate at a priority of 100. It should process SEND/REQUEST orders in a reasonably short time. SEND/REQUEST orders originate from any task desiring a despooling operation, e.g., MACRO-11, FORTRAN, and the task builder.

It is not recommended that SPR's priority be increased above 100 because the efficiency increases proportionally to the number of

SEND/REQUEST orders queued. Under these conditions, the gueue manager performs a minimum number of read operations or write operations of the gueued file.

- 3.2.4.4 SPR2 (Multiple Device Despooler) The multiple device despooler can operate at any priority from 1 through 240. The system manager should determine the optimum priority for his system. The following points should be considered.
  - The number of queued input files increases if the despooler task is not able to despool files at the same rate that they are queued. An increase in the number of files queued slightly reduces the performance of the queue manager.
  - If the system manager desires more work performed by other processes than the despooler, the despooler can be lowered to a priority that is less than the priority of the other processes.
  - 3. If it is desirable to expedite the output of certain files, the priority of the despooler can be altered dynamically using the MCR Alter Priority command. The change in priority should be lowered after the file has been despooled.
  - 4. The multiple device despooler should not have a priority greater than 210. A priority greater than 210 may cause the despooler to interfere with device handler or FlIACP task execution.

### 3.2.5 Error Recovery Procedures

Error recovery procedures for device or file errors, internal despooler errors, and queue manager errors are described in the following paragraphs.

- 3.2.5.1 <u>Device or File Errors</u> The following procedures should be used if persistent errors occur on a given device.
  - Execute the MCR OPR command to change the form type for the device.

MCR>OPR dev:/CHG

2. Abort the file being processed on the device receiving errors.

MCR>OPR dev:/AB

- Perform a redirect of the device to an alternate device, if appropriate.
- 4. Once the redirect has been performed or the device has been repaired, reset the form type to its original value.

MCR>OPR dev:/FO:n

3.2.5.2 <u>Internal Despooler Errors (Task SPR2..)</u> - The following procedures should be used if an internal error occurs in the despooler task.

1. Obtain a list of all gueued files.

MCR>QUE/AL

Repeat the MCR OPR change of form type command for all output devices in the queue.

MCR>OPR dev:/CHG

3. Abort the despooler task.

MCR>ABO SPR2.../TI=TTØ:

- 4. Perform an MCR OPR recycle command to make all active orders inactive. The /RE switch is used to recycle.
- 5. Reset all changed forms to their original value.
- 3.2.5.3 Queue Manager Errors (Task SPR...) If persistent queue manager errors such as read/write errors on the queue file occur or if the queue manager enters a loop, the following error recovery procedures should be used.
  - 1. Abort the queue manager if it is active.

MCR>ABO SPR.../TI=TTØ:

2. If possible, dump the queue file. Use the following commands to disable spooling for the line printer and to dump the file.

MCR>SET /-SP=LP: MCR>DMP LP:=[1,4]SPRQUEUE.SYS

- Delete the queue file ([1,4]SPRQUEUE).
- 4. Obtain a directory of [1,4] and if it contains any files of type .SPR requeue these files using the following command.

MCR>QUE [1,4]\*.SPR;\*/DE

5. Enable spooling for the line printer, if it was disabled.

MCR>SET /SP=LP:

### 3.3 INPUT SPOOLING

RSX-11D provides an input spooling task called [11,1]SPL. Once installed, the task can be called by any task issuing a SEND DATA AND RESUME OR REQUEST directive. Th input spooler is a multiuser task that can be called simultaneously by more than one task. The task reads variable length data from an input device.

To call the input spooler, data sent to the input spooler should have the following format: the first word contains the 2-character ASCII device name and the second word contains the unit number in binary of the input device.

Once the input spooler task is called by the send directive, it receives the same TI as the TI assignment of the task issuing the directive. When invoking several simultaneous versions of the input spooler task, the user should ensure that the tasks issuing the send/request directives have different TI assignments.

The input spooler reads double-buffered, variable-length input data. The data can have a maximum length of eighty characters.

All RSX-llD card reader handlers (CRBR and CRNP) automatically call the input spooler if input spooling is enabled on the card reader device.

During the input spooling of card decks, files are spooled directly into the user's UFD which must be specified on the \$CREATE card. For input spooling of BATCH job files the same system requirements must be met as for output spooling (see Section 3.1.1).

## 3.3.1 Enabling and Disabling Input Spooling

Input spooling is enabled and disabled using the SET command to MCR. Use the following command to enable spooling.

MCR>SET /SP=CRn:

Use the following command to disable spooling.

MCR>SET /-SP=CRn:

The DEV command can be issued to MCR to determine whether a card reader is spooled or not. If it is spooled, it has the word SPOOLED printed with the card reader status.

# 3.3.2 Basic Operations

When files are to be spooled from the card reader, set the device as being spooled, place a deck in the card reader input hopper and depress the READY switch. This action causes the card reader handler to request the input spooler task (SPL...). If the input spooler is not installed, the following message is printed on the terminal.

#### \*\*\*CRn SPL TASK NOT INSTALLED

The console operator should install the input spooler; operation resumes automatically with a maximum delay of one minute.

Once the input spooler is loaded, it reads cards until no more remain in the card reader input hopper.

#### 3.3.3 Format of the Input Deck

Input decks consist of one or more data decks. Each data deck must be peceded by a header card and followed by an EOF card. The CREATE card is the data deck header. It is followed by the actual data cards and an EOF card as illustrated in Figure 3-2. The formats of the header and EOF cards are described below.

CREATE filespecification/switches

data cards

.
.
.
.
.
.
.
.
.
.
.
.
EOF card

Figure 3-2 Format of Input Data Deck

3.3.3.1 Format of the CREATE Card (Data Deck Header) - The CREATE card has the following format.

col 1

CREATE dev:[ufd]filename.type;version/switches

CREATE must appear in columns 1 through 6. The file specification is the standard RSX-11D file specifier. Table 3-3 lists default fields in the file specifier. Table 3-4 lists and defines the switches that can be used on the CREATE card.

Table 3-3
Defaults in CREATE File Specifier

FIELD	DEFINITIONS	DEFAULT
dev*	The device onto which the file is to be spooled. It must be a directoried disk device.	SP:
[ufd]	A valid UFD that is defined on dev:.	No default. Must be specified.
filename	Name to be used for the file spooled onto the disk.	No default. Must be specified.
type	File type	DAT
ver	File version can be specified or omitted.	If /-NV, ver is l. If /NV, ver is one greater than the exist-ing version number.

Table 3-4
CREATE Card Switches

SWITCH	DEFINITION	DEFAULT
/BL /-BL	Suppress all trailing blanks Do not suppress trailing blanks	/BL
/SE /-SE	Delete card columns 73 through 80 (seguence number) Do not delete columns 73 through 80	/-SE
/NV	Create a version number that is one greater than the existing version of the file.	/NV
/-NV	Do not create a version number that is one greater than the existing version of the file.	

The input spooler first searches the UFD on the indicated device to determine whether the specified file already exists. If the file exists, a new file is not created unless the /NV switch has been specified. If /NV is not specified, the following message is printed on the terminal.

#### SPL -- FILE ALREADY EXISTS

This message is followed by a printout of the header card.

If there are no errors in the file specification, the data cards up to but not including, the EOF card are written onto the disk.

#### NOTE

Neither the CREATE card nor the EOF card is written in the file.

3.3.3.2 Format of the EOF Card - The EOF card contains a 12, 11,  $\emptyset$ , 1, 6, 7, 8, 9 multipunch in card column 1.

# 3.3.4 Input Spooling Support of BATCH

In addition to processing CREATE cards, the input spooler processes standard RSX-11D batch job control cards. A \$JOB card either as the first card in the card reader input hopper or following an EOF card is treated in the same manner as the CREATE card.

3.3.4.1 \$JOB card - The format of the \$JOB card is identical to the \$JOB card described in the RSX-11D Batch Reference Manual. The user should use the /ACCOUNT switch to establish the proper UIC for accounting. The batch file is spooled onto the physical device to which SP is directed, under UFD [1,4]. If the /NAM switch is included, the specified job name becomes the name of the file on disk. In the following example, the file is named BJOB12. The file type defaults to BIS.

#### \$JOB/NAM=BJOB12/LIM=200

If the /NAM switch is not included, the default filename is RSX11DBAT, which also is the default batch job name.

3.3.4.2 Processing of Spooled Batch Input - The \$JOB card causes a file to be created on disk. The disk file contains the \$JOB card followed by data cards, followed by the \$EOJ card.

The input spooler closes the output file once it reads the \$EOJ card. To combine several jobs in one disk file, the user should omit the \$EOJ card from each separate job and terminate only the last job with a \$EOJ card.

3.3.4.3 Automatic Calling of Batch - Provided the \$EOJ card is followed by two successive EOF cards (3.3.3.2), the input spooler closes the spooled file on disk, then issues a SEND/REQUEST directive to batch (...BAT). If ...BAT is installed, it automatically starts processing the file. The input file created by the input spooler is deleted automatically by batch after it processes the input data. If it is not installed, the operator receives no notification. This approach permits batch job files to be created with or without automatic calling of batch.

#### 3.3.5 End of Input Spooling

When the input spooler reads two successive EOF cards, it prints the following message on the operator's terminal.

### SPL -- END OF SPOOL INPUT

If the two successive end cards are followed by an input hopper empty condition, the card reader handler enters a waiting state and the input spooler exits. If the input hopper is not empty and spooling is still requested for files, the card reader begins reading cards. Input decks can be placed one behind the other in the input hopper or the hopper can be allowed to empty before a new deck is placed in the hopper. Pressing the RESET button again resumes input spooling.

# 3.3.6 Disk Files Created by Input Spooler

With the exception of batch files, the disk files created by the input spooler are not deleted by it. The files must be deleted under user control. The files cannot be accessed using the CR: input designator; they must be accessed using the full input file specification. This specification is identical to the file specification supplied on the CREATE card or to the \$JOB default specification, as described in Section 3.3.4.1.

## 3.3.7 Aborting a Batch Job

Ordinarily batch jobs should not be aborted. However, if it becomes necessary to abort a batch job, use the following procedure.

1. Abort the task BPR.

ABØ BPR.../TI=BPØ

2. Obtain an active task list.

ACT /ALL

- 3. Abort any task that has a TI designation of BP0.
- 4. Delete all batch files created for the execution of the aborted batch job. If only one batch job is queued or executing under the UFD that contained the batch job, the files should be deleted using the command:

PIP RSX11DBAT.\*,\*/DE

If more than one batch job is queued for the particular UFD, wait until the queued batch jobs complete and then deleted the files using the PIP command described above.

5. Set the invoking terminal to a nonspooled state.

SET /-SP=TTn:

6. Rerun the task BPR... to resume processing of other batch jobs within the queue.

RUN BPR...

7. The partial listing output from the aborted batch job is printed at the completion of the next batch job.

#### CHAPTER 4

#### CHECKPOINTING

Checkpointing is a mechanism to aid high priority tasks obtain the memory required to execute as soon as possible after they are requested. The process provides for temporarily halting the execution of one or more tasks in an orderly fashion to free memory for execution of a higher priority task. A task image of each halted task is rolled onto disk to free memory for the higher priority task. Checkpointing also is known as rollout/rollin.

Individual tasks can be built at task build time as checkpointable or non-checkpointable (see the RSX-11D Task Builder Reference Manual). The checkpointability of a task can be disabled/enabled using the directives described in the RSX-11D Executive Reference Manual.

Two factors are critical to the operation of checkpointing:

- Whether any executing tasks have been designated as checkpointable during task building.
- 2. The priority of tasks that require memory for execution.

To control these factors, the Executive maintains information about which tasks can be checkpointed and maintains the memory required list (MRL) of tasks waiting for execution. Checkpointable task list (CTL) information is threaded through two system tables: the active task list (ATL) and the memory required list (MRL). The layout of the ATL and MRL is provided in Appendix A of the RSX-llD Executive Reference Manual.

# 4.1 CHECKPOINTABLE TASK LIST (CTL)

Every system-controlled partition has an associated checkpointable task list (CTL). Each CTL is a priority-ordered list of all checkpointable tasks that are active in or queued to run in the partition. All tasks are checkpointable unless specified as not checkpointable (/-CP) during task building. CTL is scanned in reverse order to ensure that the lowest priority tasks are the first candidates for checkpointing.

# 4.2 MEMORY REQUIRED LIST (MRL)

The memory required list (MRL) is a priority-ordered list of tasks that are waiting for memory to become available in a partition. There is one MRL for each partition. If more than one task is in the MRL, lower priority tasks are not considered for execution until the highest priority task starts execution.

#### CHECKPOINTING

### 4.3 FUNCTIONAL DESCRIPTION

When a task is listed in the MRL, the following series of events occurs.

- A routine scans the CTL to determine whether a checkpointable task with a priority lower than the task awaiting memory is executing. The scan occurs wjenever a task is inserted or reinserted in the MRL.
- If a checkpointable task is running, a routine allocates a checkpoint image for the task on disk and, if successful, performs either of the following:
  - a. If the task has no I/O in progress, calls the rollout routine.
  - b. If the task has I/O in progress, marks the task for checkpointing. When a task is marked for checkpointing, no I/O requests are degueued for it. When the I/O in progress is completed, the I/O done function calls the rollout routine.
- The rollout routine initiates the actual rolling out of the task image and sets its status to suspended for checkpointing.
- 4. The task remains in the ATL until it is recorded on the disk.
- Another routine relinks the task's ATL node to the MRL at the task's original priority.
- 6. after the rolled out task's memory is freed, an attempt is made to activate the task that forced checkpointing. If this routine fails to find sufficient memory, it calls the routine in step 1 to try to roll out another checkpointable task.
- when a checkpointed task returns to memory, the appropriate handlers are restarted, ASTs are restarted, and the disk area used for the task image is released.

### NOTE

When the checkpointed task resumes execution, it does not necessarily occupy the same memory area in which it was executing; i.e., any available memory area of sufficient size in the same partition can be used.

#### CHAPTER 5

#### GLOBAL LIBRARIES AND COMMON AREAS

RSX-llD provides a means of using and maintaining global libraries and common areas. Global libraries are collections of routines that can be shared by many simultaneously active tasks. Global common areas are data areas in memory that can be shared by many simultaneously active tasks.

In fact, global libraries and common areas are the same type of entity. Both are treated similarly by the system. The only real difference is in terminology to aid the user in understanding their functions. For this reason, global libraries and common areas are referred to as shared global areas (SGAs). SGAs are completely dynamic; without requiring a system generation. Each task can use a maximum of seven SGAs.

# 5.1 PURPOSE AND USE OF SGAS

The purpose of a shared global area (SGA) is to allow the system designer to share code and data among simultaneously active tasks. For example, the FORTRAN compiler and the editor share the SGA (library) named SYSRES when they are both active at the same time. Similarly, two or more FORTRAN programs can exchange data through an SGA (common area).

An SGA is loaded into real memory whenever a task that links to it is requested for execution. However, if the SGA already has been loaded into memory by a previous task execution, the in-memory version is used; duplicate copies of SGAs never occur.

An access count is maintained for each SGA. An SGA remains in memory as long as the count is greater than zero. When a task exits or is checkpointed to the disk, the access count for any SGAs that the task refers to is decreased by one. If the count becomes zero, the memory allocated to the SGA is released. In addition, if the SGA was installed in the system as a read/write, as often is the case with common areas, the SGA is written to the disk when the access count becomes zero.

# NOTE

SGAs are not checkpointable and never are checkpointed. However, read/write SGAs are written back onto the disk over

the original file when their access count becomes zero.

Read-only SGAs simply have their memory released when the access count becomes zero.

The access count becomes zero if, for example, all tasks linked to the SGA are checkpointed.

The read-only (pure) area of multiuser tasks is treated in exactly the same manner as the read-only SGAs. The only difference is that the SGA is global to many tasks (i.e., referred to by name) and the pure area of a task is local to the task with which it was built. Multiuser task read-only areas are created, installed, loaded, and removed, automatically by the system.

### 5.2 IMPLEMENTATION OF SGAS

An SGA is contructed and introduced to the system in three distinct phases, namely:

writing the SGA,

Task building,

Installation and removal.

### 5.2.1 Writing the SGA

Because SGAs are used to perform a variety of functions, the following generalizations should be considered when writing an SGA.

- Read-only libraries of routines must be written in MACRO-11 assembly language. The user is responsible for all re-entrancy problems. The user must write his own re-entrant code.
- Read/write areas can be written in MACRO-11 language or as FORTRAN block data subprograms. Re-entry problems do not exist for read/write SGAs.
  - If the SGA is written as a FORTRAN block data subprogram, it is recommended that the data be initialized by the DATA statement.
- The user is responsible for synchronization of accesses to the data base in read/write areas.

### 5.2.2 Task Building the SGA

The RSX-11D Task Builder Reference Manual contains a complete description of the task builder and its options. The following descriptions are intended to provide an outline of the procedure for and the main points of task building an SGA.

First, the user is responsible for determining whether an SGA is position independent. Refer to the <a href="IAS/RSX-11">IAS/RSX-11</a> MACRO-11 Reference Manual for a discussion of position independent code. Position independency indicates the ability of a bound module (either code or data area) to exist anywhere in the virtual address space of a task. All FORTRAN block data subprograms must be built as being position independent; they are so by definition.

Secondly, the SGA must be task built. The user must be operating under UIC [1,1] when TKB is requested. The output from the task build operation is the following:

- 1. A task image of the SGA under UFD [1,1],
- 2. A symbol table file under UFD [1,1],
- 3. Optionally, a map file.

The following task builder input is mandatory:

- The no header switch (/-HD) must be used on the task image output file,
- 2. The input options STACK=0 and UNITS=0 must be specified.

If the SGA is position independent, the /PI switch can be specified. The /PI switch must be included on FORTRAN block data SGAs.

# 5.3 INSTALLATION AND REMOVAL OF AN SGA

SGAs are installed into the system and removed from it in a manner similar to task installation. The system MCR functions INS and REM are used. See the RSX-11D User's Guide for a complete description of these functions and their options. Below is a description of the main points of the installation and removal process.

- 1. At installation, the user specifies that the file being installed is an SGA by using either the /LI switch to indicate that the file is a global library or the /CM switch to indicate that it is a common area. INSTALL determines that the entity is an SGA and assumes that the SGA is a common area unless /LI is specified.
- The owner UIC is specified at installation by means of the /UIC=[n,n] switch. The owner UIC always should be specified when the SGA is installed.
- 3. The nonowner protection that is applied to the SGA is specified at installation by means of the /ACC= switch. The switch can be specified in the form /ACC=RO for read-only or

/ACC=RW for read/write. The default is no access; be sure to include the switch if tasks running under a nonowner UIC are to have access to the SGA.

# 5.3.1 Effect of Installation

Within the system, the main effect of installing an SGA is to create a global common directory (GCD) entry. The GCD entry is a control block that serves to identify the SGA to the system. Refer to the RSX-11D Executive Reference Manual for the format of a GCD entry. One important field in each GCD entry is the installed reference count.

5.3.1.1 <u>Installed Reference Count</u> - The system uses the installed reference count to keep track of the number of tasks that currently are installed and that refer to the SGA.

#### NOTE

Do not confuse the installed reference count with the active task access count.

An SGA must be installed prior to installation of the tasks that link to it.

The maximum number of tasks that can refer to one SGA is 255 (decimal). This count is set to zero when the SGA is installed. The count is increased by one whenever a task that links to that SGA is installed.

Similarly, whenever a task that links to the SGA is removed, the count is decreased. Therefore, if an attempt is made to remove an SGA using the REM command to MCR, the installed reference count is checked. The removal is not allowed if the count is nonzero, i.e., if installed tasks refer to that SGA. Before an SGA can be removed, all tasks that link to it must be removed.

### 5.4 SGA OPERATION AT RUN TIME

Once the SGA is installed in the system and a GCD entry exists, it can be loaded into memory for use by tasks that refer to it. A brief description of the basic mechanism for loading and unloading SGAs follows.

#### 5.4.1 Access Count in the GCD

At run time, the most important field of the GCD entry is the access count. This count is increased whenever a task that links to the SGA is loaded successfully into memory. It is decreased whenever the task exits or is checkpointed to the disk. The access count determines when as SGA is to loaded or unloaded.

## 5.4.2 Loading Principle for an SGA

When a task that links to an SGA is requested, the Executive prepares to load the task and at the same time checks the GCD entries for the SGA to determine their status. If the SGA is not loaded, the Executive prepares to load it as well.

The actual loading of the task and any of its related SGA is attempted only if memory can be found for all of the components (task and SGAs) of the load. If an SGA already is loaded, no attempt is made to load it again. The access count provides the Executive with this information. When the task and its SGAs have been loaded successfully, the access count for each SGA is increased by one.

## 5.4.3 Task Exit/Checkpointing and Unloading SGAs

When a task relinquishes its memory by exiting or being checkpointed, the access count for each SGA that it refers to is decreased by one. If the Executive determines that the access count is now zero for an SGA, it unloads the SGA by releasing the memory allocated to it. If the SGA is a common area, i.e., if it is read/write, it is written to disk also, and then memory is released. Read/write SGAs are written back to the disk at the physical block where the file was positioned at install time. The writing back will occur even if the file is deleted and may corrupt any subsequent file on the disk. Therefore, do not delete the task image file until the SGA has been removed.

The result of this approach is that a single SGA remains in memory as long as its access count is greater than zero regardless of how many tasks refer to it, exit, or are checkpointed.

#### 5.5 CHARACTERISTICS OF SGAS

The characteristics of SGAs described in the following paragraphs should be kept in mind when designing the use of SGAs.

 Every load of an SGA requires another disk access if the SGA is not fixed or already loaded. The result is that a task that refers to two SGAs, neither of which is loaded, requires three disk accesses to activate the task, i.e., one for the task and one for each SGA.

This process applies to multiuser tasks with pure areas as well. These tasks require two disk accesses to load them: one for the read/write area and one for the read-only (pure) area.

This does not apply to nonmultiuser tasks or to multiuser tasks with no pure area.

Every request of a task and its related SGAs requires one I/O request node to be picked for every component of the load. For example, if a task has a read-only area and links to an SGA, three I/O request nodes are picked at task request time. If the nodes cannot be picked, the request is rejected with the standard error return.

2. SGAs can map over the external page by use of the pseudo-partition name \$\$\$EXT at INSTALL time; however, an actual load of such an SGA does not occur. The simplest procedure for such an SGA is shown in the following example.

Assembly code example:

```
.TITLE
                 EXTPAG
START:
        .BLKW
                 5Ø
                      ; IGNORE FIRST 50 WORDS
                           OF EXTERNAL PAGE
XX::
        . WORD
                  Ø
                      ; AN INTERESTING REGISTER
                      ; ANOTHER INTERESTING REGISTER
YY::
        .WORD
                  Ø
                       ; BOTH GLOBALLY DEFINED
        .END
                 START
```

Task build procedure example:

```
TKB [1,1] EXTPAG/-HD, LP:/CR, SY: [1,1] EXTPAG=[200,200] EXTPAG

STACK = 0
UNITS = 0
BASE = 12 ; MAP IT UNDER ASR 3

TKB TASK, MAP=TASK

COMMON=EXTPAG:RN
```

3. The original copy of an SGA installed as a common area is updated when the SGA is unloaded. This characteristic should be considered when designing a system restart procedure.

To retain the original task-built version of a common area, use PIP to make a second version of the file. When the SGA is installed, the file with the higher version number is selected for installation. Because the original version is never installed, it is not modified when the common area is written onto the disk. The modified version is the one with the higher version number.

To reinitialize the common area with its original data, remove the common area from the system, delete the copy of the common area with the higher version number, make another copy of the original data file, and reinstall the common area.

### 5.5.1 Fragmentation and Lockouts

The system manager must avoid lockout conditions due to the unavailability of memory. The following points must be understood.

- No attempt is made to load a task until memory has been found for all components (task and SGAs) of the load.
- 2. The necessary memory need not be contiguous for each component; i.e., the memory for an SGA can be discontiguous from the memory required for the task. Memory is allocated contiguously on a component basis, i.e., for the task or SGA level.
- Attempts to allocate memory for a task are always made for the highest priority task in the MRL.
- 4. If a lower priority task fits in an area of memory that is too small for the higher priority task, the lower priority task is not loaded.
- 5. A task is not considered checkpointable unless it is being loaded or is actually running.
- 6. No moving of currently running tasks is performed to make room for a task desiring memory.

The following rules should be applied when designing the use of SGAs.

- 1. SGAs should be installed into their own partition or into the same partition as the tasks that use them.
- 2. Multiuser tasks with read-only areas should be requested into the partition into which they were installed.

# 5.6 FIXING AN SGA IN MEMORY

An SGA can be fixed in memory by fixing a dummy task that uses the SGA using the FIX command to MCR.

The dummy task should be built with the following options, which are described in the RSX-11D Task Builder Reference Manual:

```
/-CP /-FP /FX STACK = 10 UNITS = 0 LIBR = sga name or COMMON = sga name
```

The dummy task uses 96 words; this includes the task header and minimum stack size, plus one FTL entry. The above approach is the optimal solution to fixing an SGA.

Fixing the dummy task results in the loading and fixing of the SGA that it uses.

The following is an example of a dummy task.

```
.MCALL EXIT$C
START: EXIT$C ; EXIT$C is included in ; case the task is ; ever executed
```

#### CHAPTER 6

#### CORE DUMP ANALYZER

The Core Dump Analyzer (CDA) program allows the user to analyze the state of the software at the time of a system crash. The term 'crash' refers throughout this chapter either to an actual system crash or to a crash forced by the operator.

When a crash occurs, and provided the Executive has been built with a crash module appropriate to some dumping device, an area of memory is automatically dumped on that device. When the system is again running, the operator can request CDA to analyze the core dump. The same system should be used as the one that crashed. CDA can be run at the same time as other tasks.

#### NOTE

In order to benefit fully from CDA output, the user must be familiar with the RSX-llD Executive and with the PDP-ll Processor Handbook.

The output from CDA is a listing. The information listed falls into two categories, standard and optional.

The standard information is:

- An analysis of the system status at the time the core dump was taken and an indication of the cause of the system crash,
- The content of all the central processing unit registers and memory management (KTll) registers,
- 3. The content of the SCOM data area and checkpoint bit map.

The optional information is:

- An analysis of all tasks active at the time of the crash including the interpretation of their headers,
- An analysis of all system lists,
- 3. A verification of the system node pool including a comparison of the nodes allocated with those nodes in use in lists,
- 4. A complete dump of the system node pool in three output formats (octal, ASCII, and Radix 50),
- 5. A dump of the address space of any task that has memory allocated at the time of the crash.

6. A dump of a named Shareable Global Area (SGA).

Standard information is always listed. Optional information is listed only if the appropriate option switches are included in the command sequence that calls CDA. Standard information is described in Section 6.1. Optional information is described in Section 6.2.

CDA may be called more than once, if desired to analyze the same core dump.

Appendix A contains a sample CDA listing that is referred to for all examples throughout this chapter.

### 6.1 STANDARD INFORMATION PRODUCED BY CDA

The standard CDA information is system data and is printed at the beginning of the core dump analysis. See pages 1 through 3 of the sample analysis in Appendix A.

### 6.1.1 Crash Stack

The first page of output from CDA displays the data recorded on the crash stack at the time the dump was taken. The crash stack refers to the data area that is internal to the crash module. It contains the following information.

 Cause of the crash -- The specific reason for the system crash, e.g.,

# SST EXECUTED IN KERNEL MODE

If the crash dump was obtained by some means other than an Executive crash module, i.e., by a stand alone program as described in Section 6.5, the following message is printed.

#### MANUAL DUMP BY OPERATOR

- 2. The processor status register (PS), the kernel stack pointer (SP) and the user SP after the crash -- If the crash was forced by the Executive, these registers reflect the content of the respective registers after the IOT (I/O trap) instruction has been executed.
- 3. The program counter (PC) and PS before the crash -- If the crash was forced by the Executive, the content of these two registers is pushed onto the kernel stack by the IOT instruction. The PC value is equated to a specific global crash code whenever the .CRASH macro appears in the Executive code.
- 4. RØ through R5.
- 5. Segmentation status registers 0 through 3. On a PDP-11/40, not all of the registers are present. Refer to the  $\frac{\text{PDP-11}}{\text{Processor Handbook}}$ .
- 6. All of the segmentation registers for user, supervisor, and kernel modes for both I and D space. On a PDP-11/40, not all of these registers are present. Refer to the PDP-11 Processor Handbook.

The information contained in the crash stack is the key to the operation of CDA. In particular, the contents of the kernel PARs enable CDA to convert virtual memory addresses into the physical block and word addresses of the dump medium.

# 6.1.2 Kernel Stack

The kernel stack of 128 words is dumped on the second page of the CDA output. The location of the kernel stack is determined from the symbol .SG.TS in the Executive, which is the top of the Executive stack.

# 6.1.3 System Communication Area (SCOM)

The system communication area (SCOM) data area is dumped on the third page of the CDA output. The contents of the following SCOM items are interpreted.

- 1. Time of the crash
- 2. Current task name and ATL address (.CRTSK)
- PDP-11/45 floating point hardware indicator (.FP45X)
- 4. Processor type (.PDP11)
- 5. Size of memory in 32-word blocks (.MSIZE)
- 6. Size of the checkpoint file (.CKFSZ)
- 7. Significant event flags (.SERFG)
- System event flags (.COMEF)
- 9. The entire SCOM data area from .SG.BC to .SG.EC

# 6.2 OPTIONAL OUTPUT FROM CDA

Optional output from CDA is obtained by including the appropriate switches in the command string that calls CDA, as described in Section 6.3.

All switches are presented in alphabetic order by switch name for ease of reference. However, anyone unfamiliar with CDA should read all descriptions before using the following information for reference purposes. Also, refer to the tables and lists described in the RSX-11D Executive Reference Manual.

# 6.2.1 All Lists (/ALL)

When the /ALL switch is included in the command string, CDA includes the following information in the printed output.

- 1. AST queues
- 2. Active task list
- 3. Batch command buffers
- 4. Clock gueue
- 5. Fixed task list
- 6. Global common directory
- 7. Task headers
- 8. MCR command buffers
- 9. System node pool
- 10. Physical unit directories (PUDs)
- 11. SEND/RECEIVE gueues
- 12. STD entries
- 13. Task partition directories
- 14. Verification of the system node pool

Using the /ALL switch is equivalent to including all CDA command switches except /TAS, /DMP and /SGA. These last three are used to dump specified portions of memory.

## 6.2.2 AST Queues (/ASQ)

A dump of the AST queues is included if the /ASQ switch is used. See page 46 of the sample listing for an example of AST queue output.

The /ASQ switch causes all tasks which have any entries in their AST queue to have such output. Note that the AST queue also includes I/O requests which have been completed but for which the task has not yet been notified.

For each task with ASTs pending, the task name, ATL address (real and virtual) and STD address are output as a heading.

I/O completion entries are output as:

- 1. AST entry point, if any
- 2. Event flag number specified in the QIO
- 3. I/O status block address
- 4. I/O status block contents (two words)
- Intermediate buffer address (real and virtual) (RSX-11D intermediate buffered transfer only)
- User virtual buffer address (RSX-11D intermediate buffered transfers only).

For true AST entries the output is:

- 1. AST type
- 2. AST entry point
- Additional information, if any, which will be pushed on to the user's stack.

# 6.2.3 Active Task List (/ATL)

A dump of the active task list is included in CDA output if the  $/ \mathrm{ATL}$  switch is used. See page 4 of the sample listing for an example of ATL output.

The /ATL switch causes all entries in the ATL to be listed and provides interpretation of the following items in the order described.

- 1. ATL address -- virtual and real address nodes
- 2. STD address -- virtual and real address nodes
- 3. REQUESTER -- name of the task that requested this task
- 4. Load address -- real memory address of the task's header
- 5. Name of the partition in which the task was running
- 6. The task's run priority
- 7. I/O in progress count
- I/O pending count
- 9. Status -- the same 2-character codes as those used by ...ACT as described in the RSX-11D User's Guide
- 10. TI assignment -- the terminal from which the task is being run
- 11. Flag word -- A.TF
- 12. Flags -- task's 32 event flags (two words)
- 13. Masks -- task's 64 event flag masks (four words)
- 14. Entire ATL node dumped with virtual addresses

# 6.2.4 Batch Command Buffers (/BAT)

A dump of the batch command buffers is included in CDA output if the /BAT switch is used in the command string. On page 44 of the example, the sample listing indicates that no batch command buffers were in use at the time of the crash.

The format of CDA output for the batch command buffer is the same as the format for the MCR command buffers shown on page 43 of the sample listing.

The /BAT switch causes all the batch command buffers to be dumped and provides interpretation of the following items for each buffer.

- 1. Node address -- virtual address of the node
- 2. Name of the MCR task requested
- 3. TI assignment -- terminal initiating the batch job
- 4. Length of the buffer
- 5. The first 40 characters in the buffer converted to ASCII

# 6.2.5 Clock Queue (/CKQ)

A dump of the clock gueue is included in CDA output if the /CKQ switch is used. See page 48 of the sample listing for an example of CKQ output.

The /CKQ switch causes all entries in the clock queue to be dumped. The format of the output depends on the type of the clock queue entry. For all types, the first line indicates the type and the time in ticks which must elapse after the previous entry before this entry will be dequeued (the 'Delta' time).

#### The types are:

- a. Time-slice, for the RSX-llD time-slicing scheduler. No further information is output.
- b. Mark time. The following is supplied:
  - 1. Task name
  - 2. ATL address virtual and real node address.
  - 3. Event flag number
  - 4. AST entry point address
- c. Schedule, that is, a request for a task to run at a specified tiime. The additional information is:
  - 1. Name of requested task
  - STD address of requested task real and virtual addresses.
  - 3. Task name of requestor.
  - 4. STD address of requestor real and virtual addresses.
  - 5. Partition specified for task to run in, or blank if no task specified.
  - Priority specified for task, or zero if no priority specified.
  - 7. UIC specified for task to run under, or  $[\emptyset,\emptyset]$  if no UIC specified.
  - 8. TI for which the task will run.
- d. Reschedule, that is, a schedule which is to be repeated at a regular interval. The additional information is as for 'schedule' in c. above, with the addition of the interval between periodic reschedules, in ticks.

# 6.2.6 Dump of Real Memory (/DMP)

The /DMP switch causes a portion of real memory to be dumped in three notations: octal, ASCII, and radix 50. The /DMP switch has the following syntax.

#### /DMP:start:end

start is the starting real memory address. An octal number representing up to 22 bits is used.

end is the ending real memory address. An octal number representing up to 22 bits is used.

Both the starting and end addresses must be specified. Only one /DMP switch can be specified.

# 6.2.7 Fixed Task List (/FTL)

A dump of the fixed task list is included in CDA output if the /FTL switch is used. See Page 45 of the sample listing for an example of FTL output.

The /FTL switch causes all entries in the FTL to be listed and provides interpretation of the following items in the order described.

- 1. FTL address -- virtual and real address nodes
- 2. STD address -- virtual and real address nodes
- 3. Load address -- real memory address of the task's header
- 4. Name of the partition in which the task was running
- 5. The task's run priority
- 6. I/O in progress count
- 7. I/O pending count
- 8. Status -- the same 2-character codes as those used by ...ACT as described in the RSX-11D User's Guide
- TI assignment -- the terminal from which the task is being run
- 10. Flag word -- A.TF
- 11. Flags -- task's 32 event flags (two words)
- 12. Masks -- task's 64 event flag masks (four words)
- 13. Entire FTL node dumped with virtual addresses

# 6.2.8 Global Common Directory (/GCD)

A dump of the global common directory is included in CDA output if the /GCD switch is included in the command string. See pages 49 and 50 of the sample listing for an example of GCD output.

In addition to the entries in the GCD itself, entries linked to STDs for the pure area of a multiuser task also are dumped. The formats of both types of nodes are identical except that pure area GCD nodes do not contain the name of the task. Instead, an octal 35 (an illegal radix 50 character) is in the first word of the name. The following items are interpreted for each GCD when the /GCD switch is specified.

- Name of the common block or library, or octal 35 if the node is for the pure area of a task
- 2. Virtual address of the GCD node
- 3. Size -- 1/64th of the size of the common block, library, or pure area
- 4. Name of the partition in which the shared area is installed
- 5. Flags byte
- 6. Status of the global area using the same 2-character codes as ...ACT with the addition of NL to indicate not in use or not loaded
- Active reference count and installed reference count separated by a slash (/)
- 8. Creation time
- 9. UIC

## 6.2.9 Task Headers (/HDR)

A dump of the task headers is included in CDA output if the /HDR switch is included in the command string. See pages 8 and 9 for an example of task header output.

Interpretation of the following items is provided in addition to the dump of the headers.

- 1. Header address -- real address of the task's header
- 2. ATL address for the task
- 3. H.CRl -- context buffer reference point 1
- 4. User PDRs
- 5. User PARs
- 6. PS, PC, RØ through R5, and SP
- 7. Initial PS, PC and SP
- 8. Size of the header in 32-word blocks
- 9. Directive status word (DSW) -- task's virtual zero
- 10. Number of LUNs assigned to the task
- 11. Logical unit table (LUT) -- the device and unit address for each entry in the LUT. If there is a window open for the entry, the following additional items are dumped.

From the window: W.CTL, W.VBN, W.FCB

From the FDB: F.FNUM, F.FSEQ, F.STAT, F.NACS, F.NLCK

12. The entire header is dumped.

### 6.2.10 MCR Command Buffers (/MCR)

A dump of the MCR command buffer is included in CDA output if the /MCR switch is used in the command string. See page 43 of the sample listing.

The /MCR switch causes all MCR command buffers to be dumped and provides interpretation of the following items for each buffer.

- 1. Node address -- virtual address of the node
- 2. Name of the MCR task requested
- 3. TI assignment -- terminal initiating the batch job
- 4. Length of the buffer
- 5. The first 40 characters in the buffer converted to ASCII

# 6.2.11 Dump of System Node Pool (/POL)

Inclusion of the /POL switch causes the entire system node pool to be dumped with virtual and real addresses (see pages 54 and 55 of the sample listing). Each node (8-word block) is dumped in three formats: octal, ASCII, and radix 50. In addition if the node is in use (allocated), an asterisk (\*) is printed at the lefthand margin.

# 6.2.12 PUDs and I/O Request Queues (/PUD)

Inclusion of the /PUD switch causes all peripheral unit directories (PUDs) to be dumped. If any I/O requests are queued for the device type, the I/O request nodes are dumped after each associated PUD. See pages 31 and 32 of the sample listing.

The following items are interpreted for each PUD entry.

- 1. Device name and unit number
- 2. Flag byte
- 3. Name of the task attaching the device, if any
- 4. Device to which this one has been redirected
- 5. Handler task name
- Count of express requests in the queue (EXP CNT)
- 7. U.CH -- characteristics word in PUD (CHAR)
- 8. First three characters of the default ACP task name in radix 50 (DACP)
- 8. Name of the ACP task
- 9. Terminal flag bytes (TERM)
- ll. A dump of the entire PUD The following items are interpreted for the I/O request node.
  - 1. Name of the task generating the I/O request
  - 2. ATL node address of the requester
  - 3. Size of the DPB (DPB SZ)
  - 4. Priority of the request
  - 5. LUN
  - 6. Event flag number (EFN)
  - 7. Function code -- if the function code is one of the more common function codes, it is printed symbolically, e.g., IO.WLB. If it is not a common code, IO.XXX is printed.
  - Request type -- either normal or express
  - 9. A dump of the entire I/O request node

## 6.2.13 Dump of an SGA (/SGA)

The /SGA switch is used to dump a shareable global area (library or common area). The SGA must have been resident in memory.

The /SGA switch has the following syntax, with brackets indicating the optional portion:

/SGA=sqanamel[:start:end]

#### where:

sganame is the name of the SGA to be dumped

start is the starting virtual address in octal of the portion of the SGA to be dumped.

end is the ending virtual address in octal of the portion of the SGA to be dumped.

Default: if the start and end addresses are omitted the entire SGA is dumped.

For a position-independent SGA, addresses are assumed to start at 0, both in the dump and in reckoning the 'start' and 'end' values. For a non-position-independent SGA, addresses correspond to those used when the task was built, e.g. an SGA mapped through APR3 would start at 60000.

This switch may only appear once in a CDA command string. See pages 102 and 103 of the sample listing for an example of the output obtained by using the /SGA switch.

## 6.2.14 SEND/RECEIVE Queues (/SRQ)

Inclusion of the /SRQ switch (see page 47 of the sample listing) causes all the SEND/RECEIVE queues to be listed and provides interpretation of the following items for each queue.

- 1. Task name of the receiver
- 2. Task name of the sender
- 3. TI -- terminal of receiver or sender task
- 4. Priority of the send
- 5. Data block contents

## 6.2.15 System Task Directory Entries (/STD)

Inclusion of the /STD switch causes all the entries in the system task directory (STD) to be dumped. See pages 22 and 23 of the sample listing. When /STD is used, CDA provides interpretation of the following items for each STD entry.

- 1. Name of the task
- 2. Name of the partition in which the task is installed (PAR)
- 3. STD flag word
- 4. Priority at which the task is installed (PRI)
- 5. Disk from which the task is installed (DISK)
- 6. Size of the task (SIZE) -- S.TZ
- 7. Number of active versions of the task (AV)
- 8. Pool limit specified in 8-word blocks (LIM)
- 9. Pool utilization specified in 8-word blocks (UTIL)
- 10. SRQ listhead

## 6.2.16 Dump of Task's Address Space (/TAS)

Inclusion of the /TAS switch causes the address space of a task that has memory allocated to it at the time of the crash to be dumped (see pages 100 and 101 of the sample listing). The task must have been either active or fixed in memory. Each task area is dumped in three notations: octal, ASCII, and radix 50.

The /TAS switch has the following syntax with brackets indicating the optional portion.

# /TAS=taskname[:start:end]

taskname is the name of the task to be dumped

start is the starting virtual address in octal for the portion of the task to be dumped

end is the ending virtual address in octal of the portion of the task to be dumped.

If the starting and ending addresses are omitted, the entire task is dumped. When this switch is used to dump a multiuser task, only its impure area is dumped.

#### NOTE

Only tasks with entries in the STD can be dumped using the /TAS switch. Pure areas of multi-user tasks cannot be dumped.

# 6.2.17 Task Partition Directory (/TPD)

Inclusion of the /TPD switch causes all entries in the task partition directory to be dumped. Also dumped are any entries in the memory required list (MRL) or the checkpointable task list (CTL) for each partition. An allocation map is printed for each partition showing the unused area, active and fixed tasks, and shared areas in the partition. See pages 19 and 20 of the sample listing.

The following items are interpreted for each partition.

- 1. Partition name
- 2. TPD entry address -- virtual and real address of TPD entry
- 3. Base address -- real starting address of the partition
- 4. Size of the partition in bytes
- 5. TPD flag word
- Unused areas (HOLES) -- if the partition is system-controlled, the address and size of each unused area is printed.
- 7. Allocation map -- all active or fixed tasks and shared areas in the partition are listed with the real address of the space allocated. One of the following letters is used to indicate the type of task or area.
  - A -- active task
  - F -- fixed task
  - P -- pure area of multiuser task
  - L -- library
  - C -- common block

If the partition has any entries in the memory required list (MRL), the following items are interpreted.

- 1. Virtual address of the MRL node (NODE ADR)
- 2. Name of the task
- 3. Size of the task (SIZE) -- S.TZ
- 4. Run priority of the task
- 5. Checkpoint priority (CP PRI)
- 6. Status (STS)
- 7. Checkpoint status (CP STS)
- 8. I/O node address
- 9. TI

If the partition has any entries in the CTL, the following items are interpreted.

- 1. CTL node address
- 2. Name of the task
- 3. Run priority
- 4. Status (STS)
- 5. TI

# 6.2.18 Verification of System Node Pool (/VFY)

Inclusion of the /VFY switch causes the system node pool to be verified (see pages A-16 and A-17 of the sample listing). The following procedures are used to verify the node pool.

- 1. The number of nodes (8-word blocks) in use is determined.
- The total number of nodes in the pool also is calculated using the following formula.

Total number of nodes = (Mod32(.STDTA+(.STDTC\*2))-.PUDBA/8

.STDTA is the starting address of the STD alpha table. .STDTC is the capacity of the STD. .PUDBA is the beginning address of the PUD.

This formula provides the total number of nodes available in the pool, i.e., those that are not allocated permanently to fixed lists such as the PUD, TPD, and STD alpha table.

 All the dynamic system lists are scanned to find the nodes that have been picked from the pool. The lists that are searched include the following.

> Fixed task list Batch command buffer I/O request nodes Volume control block Vector connection (..CINT) nodes AST nodes Window nodes AST nodes linked to task headers Active task list CKQ (clock gueue) GCD nodes I/O nodes linked to GCD nodes MCR command buffer Memory required list I/O nodes linked to MRL nodes AST nodes linked to MRL nodes System task directory SEND/RECEIVE Queue Pure area GCD nodes

As a result of this search, a bit map is constructed to show which nodes are in use in system lists. Although every list is searched, it is still possible to miss some nodes because any privileged task can pick nodes. For example, the teletype handler uses nodes for read-ahead buffers. Nodes in

use for privileged tasks are not found when the system lists are searched.

- 4. The bit map created in step 3 above is scanned and the number of nodes marked as being in use is printed.
- 5. The pool is tested for discrepancies. First, the nodes actually allocated, but not found in any list, are dumped. There always are some of these.

Next, the nodes found in the lists, but not allocated, are dumped. If there are any of these, they should be examined carefully. They indicate that the current pool allocation does not reflect the status of the nodes in the pool accurately.

Finally, the bitmaps representing the nodes allocated and the nodes found in the lists are dumped.

## 6.3 OPERATIONAL INFORMATION

The operational information required for CDA includes the process of linking a crash module into the Executive (Section 6.3.1) and the command string used to call CDA (Sections 6.3.2 and 6.3.3).

# 6.3.1 Linking a Crash Module into the Executive

In order for a crash module to write the dump that provides CDA with input, a crash module must be linked into the Executive.

The command file [11,15] TKB15.CMD is used as input to the task builder for this purpose. Appendix B provides a listing of the command file.

To link the crash module into the Executive, replace NCRASH.OBJ with one of the following crash modules:

TC11CM or TC11CR to dump onto DECtape,

TUl0CM or TUl0CR to dump onto TUl0 magnetic tape,

TU16CM or TU16CR to dump onto TU16 magnetic tape,

RKØ5CM or RKØ5CR to dump onto RKØ5 disk.

The modules with names ending in the letters CM are the same as those with names ending in CR except that the CM versions cause more detailed messages to be printed on the console at the time of a crash.

If none of the crash modules is needed, but the more detailed console messages are desired, replace NCRASH.OBJ with the module MCRASH.OBJ.

Refer to the listing in Appendix B to determine default unit numbers for the crash modules and how to modify them.

Rebuild the Executive using the TKB15.CMD file and rebuild the bootstraps using [11,17]BOOTSBLD.CMD. Then rebuild CDA using [11,13]CDABLD.CMD.

# 6.3.2 Calling CDA

Core dump analyzer is a task that can execute online with other tasks. Type the following command to MCR to install CDA.

MCR>INS [11,1]CDA

The following command syntax is typed to MCR to request CDA.

MCR>CDA [option switches]

option switches

are those switches described previously in Section 6.2 and summarized in Table 6-1 below. Any combination of switches can be used, or all switches can be omitted.

Table 6-1
Summary of CDA Option Switches

SWITCH	FUNCTION
/ALL	Lists the same information as produced by using all CDA option switches except /TAS, /SGA and /DMP.
/ASQ	Lists and interprets ASTs and I/O requests completed but not yet notified to the task.
/ATL	Interprets and dumps ATL entries.
/BAT	Interprets and dumps batch command buffers.
/CKQ	Interprets and dumps clock queue entries.
/DMP:start:end	Causes a specified portion of memory to be dumped.
/FTL	Interprets and dumps the fixed task list.
/GCD	Interprets and dumps GCD entries for libraries, global common areas, and multi-user task pure areas.
/HDR	Interprets and dumps individual task headers.
/MCR	Interprets and dumps MCR command buffers.
/POL	Dumps the node pool in octal, ASCII, and Radix 50.
/PUD	Interprets and dumps PUD information.

(Continued on next page)

# Table 6-1 (Cont.) Summary of CDA Option Switches

SWITCH	FUNCTION
/SGA=sganame[:start:end]	Dumps all or part of a named SGA
/SRQ	Interprets and dumps SEND/RECEIVE queues.
/STD	Interprets and dumps STD entries.
/TAS=taskname[:start:end]	Dumps all or part of the task's address space for a task that has memory allocated to it.
/TPD	Interprets and dumps all entries in the task partition directory.
/VFY	Causes verification of the system node pool.

## NOTE

The CDA input LUN is assigned to DTØ. To change this assignment, use the REA command as in the following example before calling CDA.

MCR>REA ...CDA 3 DK1:

## 6.3.3 Examples of CDA Command Strings

The following are examples of CDA command strings:

Example 1:

MCR>CDA

Result: Only the system data is printed.

Example 2:

MCR>CDA /ALL

Result: All the optional output is printed.

Example 3:

MCR>CDA /POL

Result: The system node pool is dumped.

Example 4:

MCR>CDA /ATL/HDR

Result: The ATL nodes and task headers are dumped for those tasks that were active when the system crashed.

## Example 5:

MCR>CDA /PUD/TAS=TT....

Result: All the PUD's are dumped followed by a dump of the entire task "TT....".

# Example 6:

MCR>CDA /TASK=XYZ:0:2000

Result: If the task XYZ had memory allocated at the time of the crash, its virtual address space from 0 to 2000 is dumped. If no memory was allocated for XYZ a message is printed instead.

# Example 7:

MCR>CDA /DMP:4400000:4430000

Result: Real memory from 4400000 to 4430000 is dumped.

# 6.4 OPERATOR ERROR MESSAGES

The messages listed in Table 6-2 are reported on the console requesting CDA when an error is detected that prevents CDA performing the analysis. 'Input device' is the device on which the memory dump is held and which provides the input to CDA.

Table 6-2 CDA Error Messages

Message	Meaning	Suggested Action
CDA DEVICE READ ERROR	Input device is not ready or is other-wise unavailable.	Check status of device. Call field service if problem persists.
CDA PLEASE MOUNT INPUT DEVICE AS FOREIGN	Input device not mounted as a foreign volume.	Mount the input de- vice as foreign.
CDA DEVICE HANDLER MISSING	The device handler for the input device is not loaded.	Load the handler, then retype the command string.
CDA FAILED TO READ MCR COMMAND BUFFER	CDA was not invoked as an MCR function.	Execute CDA as an MCR function, not with RUN.
CDA SYNTAX ERROR	There is a mistake in the command string.	Correct the error and retype the command.
CDA ILLEGAL SWITCH	An unrecognized switch was found in the command string.	Type in the command string again using the correct switch.
CDA DUMP ABORTED - KERNEL PAR'S CLOBBERED	Either the crash clobbered the PAR's or the format of the crash stack on the dump medium is incorrect.	Try to obtain another crash.

## 6.5 INPUT TO CDA

The usual source for input to CDA is the dump produced by one of the RSX-11D Executive crash modules that can be linked in when the Executive is built (see Section 6.3). If the dump has been made by one of these modules, it is possible for CDA to report the cause of the crash.

# 6.5.1 Forcing a Crash

It is also possible that the user would like to crash the system in order to study the dump produced.

The executive can crash for any of the following reasons:

- 1. Odd Address or TRAP 4
- 2. Undefined System Event
- 3. AST Node Pick Failure
- 4. Floating Point Error
- 5. Memory Parity Error
- 6. Non-RSX-EMT execution
- 7. Reserved Instruction
- 8. SST in Kernel Mode
- 9. Red or Yellow Stack
- 10. EMT 377 in Kernel Mode
- 11. Undefined System Error.

A crash is forced by generating one of the above conditions. The two easiest methods are 1 and 9.

An 'odd address or trap 4' can be generated by halting the processor, setting the Program Counter (PC) value odd, and continuing.

A 'Red or Yellow Stack' can be generated by halting the processor, depositing zero into the Kernel Stack Pointer, and continuing.

The Console Switches are used to halt, set an appropriate value and continue. See the section on the Operator's Console in the Processor Handbook for the model involved.

#### NOTE

Causing an active, running system to crash may result in serious consequences. It always causes a complete loss of all data in memory. It can cause corruption of disk file structures. Use extreme caution, e.g., write protect disks first.

## 6.5.2 Dumps from Stand-alone Programs

Occasionally, none of these methods causes a crash; e.g., the system is dead or in a very tight loop. For this reason, various techniques have been developed for dumping a system without using the crash modules.

All of these techniques involve bootstrapping a stand-alone program into memory. The most common technique is to bootstrap a program from DECtape. This program proceeds to write the dump on the DECtape.

CDA can interpret the dumps from stand-alone programs if the following conditions are met.

- The memory dump begins at block 101 (octal) for disks or DECtape or at physical record 2 for magnetic tape. The memory dump must begin at real memory location 0.
- Block 100 (octal) or record 1 must be formatted in a specific way and must include at least the kernel PARs in particular locations. For the precise format of this information, refer to Figure 6-1.

## NOTE

If the dump was not made by an Executive crash module, the first 2000 (octal) bytes of real memory may be lost because the hardware bootstraps for some devices, notably DECtape, use the first 2000 bytes. In certain cases, losing these 2000 bytes can be critical, especially if the kernel stack extends into that area.

THE INTERNAL STACK (AT 'ENDST') CONTAINS THE FOLLOWING INFORMATION:

(SP+34P IS AT LOCATION 'ENDST') SP+334 -= PROGRAM STATUS WORD AFTER BOOT
SP+332 -= KERNEL SP AFTER BOOT
SP+338 -= PC BEFORE CRASH IOT (FROM KERNEL STACK)
SP+326 -= PS BEFORE CRASH IOT (FROM KERNEL STACK)
SP+324 -= USER SP
SP+320 -= R1
SP+326 -= R1 SP+316 -- R2 SP+314 -- R3 3P+314 -- R3
SP+310 -- R4
SP+310 -- R5
SP+306 -- SEGMENTATION STATUS REGISTER 8
SP+394 -- SR1
SP+392 -- SR2
SP+396 -- SR3
SP+276 -- USER I SPACE DESCRIPTOR REGISTER 8 SP+260 -- USER I SPACE DESCRIPTOR REGISTER 7 SP+256 -- USER D SPACE DESCRIPTOR REGISTER 0 SP+240 -- USER D SPACE DESCRIPTOR REGISTER 7 SP+236 -- USER I SPACE ADDRESS REGISTER 0 SP+220 -- USER I SPACE ADDRESS REGISTER 7 SP+162 -- USER D SPACE ADDRESS REGISTER 6 SP+200 -- USER D SPACE ADDRESS REGISTER 7 SP+176 -- SUPERVISOR I SPACE DESCRIPTOR REGISTER 0 SP+160 -- SUPERVISOR I SPACE DESCRIPTOR REGISTER 7
SP+156 -- SUPERVISOR D SPACE DESCRIPTOR REGISTER 0 SP+140 -- SUPERVISOR D SPACE DESCRIPTOR REGISTER 7 SP+136 -- SUPERVISOR I SPACE ADDRESS REGISTER Ø SP+120 -- SUPERVISOR I SPACE ADDRESS REGISTER 7
SP+116 -- SUPERVISOR D SPACE ADDRESS REGISTER 0 SP+100 -- SUPERVISOR D SPACE ADDRESS REGISTER 7
SP+076 -- KERNEL I SPACE DESCRIPTOR REGISTER 0 \*\*
\$P+060 -- KERNEL I SPACE DESCRIPTOR REGISTER 7
\$P+056 -- KERNEL D SPACE DESCRIPTOR REGISTER 0 SP+040 == KERNEL D SPACE DESCRIPTOR REGISTER 7 SP+036 -- KERNEL I SPACE ADDRESS REGISTER Ø \*\* SP+020 \*\* KERNEL I SPACE ADDRESS REGISTER 7
SP+016 \*\* KERNEL D SPACE ADDRESS REGISTER 0 SP+000 -- KERNEL D SPACE ADDRESS REGISTER 7

\*\*\*\*\* NOTE: THE D=SPACE AND SUPERVISOR REGISTER LOCATIONS WILL CONTAIN RANDOM GARBAGE WHEN THE DUMP IS MADE ON AN 11/40.

Figure 6-1 Format of CDA Input (Record 1 or Block 100)

## CHAPTER 7

#### TIME-SCHEDULED PARTITIONS

RSX-11D supports a time-based scheduling algorithm as an aid to program development. The algorithm performs round-robin scheduling on all tasks of which the following are true:

- The task runs in a time-slice controlled partition, as defined at System Generation.
- The task, when installed, has a priority lying between limits defined in the system.

The running of time slicing tasks takes place within the priority structure of RSX-11D. Tasks satisfying the above conditions will come under time slicing whether or not they are checkpointable (liable to rollout/rollin).

## 7.1 TIME-SCHEDULING OF TASKS

Round-robin scheduling has two time parameters:

- The CPU Burst, which is the length of time (number of clock ticks) that must pass before the current user of the CPU is changed.
- 2. The Memory Swap Slice, which is the length of time (number of CPU bursts) that must pass before a resident checkpointable task is rolled out of memory, if the MRL is not empty.

Other parameters are those associated with task priority.

The mnemonics for these parameters and the values supplied in the released system are described in Table 7-1.

A time sliceable task is a task which has been installed into a time-slice controlled partition at a priority of between 3 and PR.MAX.

When such a task is loaded the task's priority is set to the default value PR.DFL associated with time slicing. Normal RSX-11D processing then resumes. No time sliceable task will execute unless all tasks, whose installed priority is greater than PR.MAX, are idle. When the last time sliceable task exits, time slicing is turned off.

When a CPU burst has been completed, the time scheduler is again activated. The Active Task List (ATL) is searched for the first task whose priority is equal to the default. This task is deleted from the ATL, and from the Checkpointable Task List (CTL) if it is checkpointable, and then reinserted into these lists as appropriate. This has the effect of 'lowering' the task's priority in the sense that all other equal priority tasks now precede it in the ATL. Another CPU burst is initiated and the scheduler becomes inactive.

#### TIME-SCHEDULED PARTITIONS

By contrast, when a memory swap slice has been completed the priority of each task, whose current MRL priority is between 3 and PR.MAX, is incremented by 1. In addition, the status of any task on the MRL whose priority is 2 is separately checked. The status of such a task can imply that it is now runnable, for example because of one of the following:

- 1. an awaited event flag has been set
- 2. the task has been resumed from a suspended state
- 3. its status has been changed to 'RUN'
- 4. an AST has been declared for it
- 5. the task has been aborted
- 6. an intermediately buffered I/O request has been terminated.

If such a task with priority 2 is now runnable, its priority is raised either to the maximum time-slice value of PR.MAX or to its installed priority, whichever is the greater.

At this point an attempt is made to activate the highest priority task on the MRL, firstly by a normal MRL scan and secondly by trying to force checkpointing for lower priority tasks. The first candidate for swapping out is the task now at the bottom of the list of those sharing the default priority.

These attempts to activate the highest priority task on the MRL can still fail to provide the necessary memory. Normally the next step is to begin another memory swap slice. But if a task with priority PR.MAX remains at the top of the MRL for 3 memory swap times, then this task is 'demoted'. Demotion is accomplished by subtracting PR.DEC from the task priority and reinserting the task, by priority, in the MRL. Demotion can be inhibited by setting PR.DEC=0.

# 7.2 PRIORITY OF TASKS ON THE MRL

When a task is put on the MRL, either as a new request or as a roll-out, it has one of two priorities. These are:

- 1. Its installed priority
- 2. A priority value of 2. A task is given this priority whenever it is rolled out awaiting some external event which blocks additional execution. In this case the task is inhibited from contending for memory until the external event becomes due (see previous section).

A task with a installed priority greater than the maximum time-slice priority will always be rolled in before any task whose priority is less than or equal to this maximum. This is because, during time slicing, MRL priority will never be incremented beyond this maximum.

If, for example, PIP is installed into a time sliced partition with a priority of 55 and TKB is installed into a time sliced partition with a priority of 50, then as long as both PIP and TKB are on the MRL, PIP will be rolled in first. However, once PIP is rolled in, the MRL priority of TKB will continue to be incremented until, at some subsequent rollout of PIP, PIP will appear in the MRL after TKB, whose priority by now has been incremented to 55 or more.

#### TIME-SCHEDULED PARTITIONS

Rollout is not initiated until there is a task on the MRL whose priority is greater than or equal to the default time-slice priority. This ensures that low-priority 'back-ground tasks' do not get rolled in as often as others. As an example, if task X has been installed at priority 30, then since the default priority is 50, 20 memory swap slices occur before task X is even considered as a candidate for memory.

A task will not be rolled out if its executing priority is greater than the MRL priority of the first (highest priority) task on the MRL. If there is not room in memory for the highest priority task on the MRL, no task of lower priority will be rolled in until the demotion algorithm, described above, takes effect.

## 7.3 TIME SLICING PARAMETERS

The following table summarizes the parameters involved in 7.1 and 7.2 above. The period (.) indicates values in decimal. B indicates a byte value,  $\bar{W}$  a word.

Table 7-1
Time Slicing Parameters

Parameter	V6.2 System Value	Meaning
PR.DFL	50.	The execution priority of any task whose installed priority is between 3 and PR.MAX inclusive. (B)
PR.MAX	75.	The maximum priority for CPU time slicing. If a task is installed with a priority greater than PR.MAX it will always execute at its installed priority. Also, if the task is on the MRL, it will always take precedence over all time sliceable tasks. (B)
SL.BST	2.	The CPU burst which is the number of clock ticks before putting another equal priority task at the top of the 'equal priority ATL subset'. (W)
SL.SWP	15.	The memory swap slice which is the number of SL.BSTs before swapping is initiated. (W)
SL.BLK	3.	The number of successive SL.SWPs before a task is demoted from the top of the MRL. (W)
PR.DEC	Ø.	The amount by which a task's priority is decremented when the time slicer decides to 'demote' the task. (W)

## TIME-SCHEDULED PARTITIONS

# 7.4 LIMITATIONS OF TIME SLICING

RSX-llD Time slicing is designed to cover program development - editing, assembling, linking - by one or more interactive users. In this way better service can be provided than by the standard priority scheduling of RSX-llD. It is not designed as a full timesharing system.

#### CHAPTER 8

#### ERROR LOGGING

In RSX-llD the logging of hardware is performed for both memory errors and device errors. The Executive log memory errors, which can be either main memory parity errors or cache memory errors on the PDP-l1/70. Device handler tasks developed for RSX-llD log device errors.

A single error log contains items logged by both the Executive and the device handler tasks. The system manager can use the output from the error logging analysis task to determine the reliability of memory and devices used in the RSX-11D system.

Error statistics are accumulated by handlers for the following devices:

- · Disks,
- DECtape,
- Magnetic tape.

For further information on the handling of parity errors, refer to Chapter 11 of this manual.

The report produced by error logging can contain itemized error statistics with optional summary information, or it can contain only the summary information. Additionally, the system manager can select the time frame that the report is to encompass and can indicate that the report is to include only memory parity errors or device errors. Further, if a report on device errors is desired, the device type, unit, or volume for which the report is desired can be specified.

Figure 8-1 provides a sample report for a device error. Figure 8-2 provides a sample of summary information. All numbers followed by a period (.) are decimal values. All others are octal. Appendix C contains sample reports for various types of errors.

PAGE 1.

RSX 11 SYSTEM ERROR REPORT COMPILED AT 20-AUG-75 13:46:42

```
ERROR LOGGING SYSTEM STARTED
AT 12-AUG-75 10:13:45
***********
TAPE HARDWARE ERRUR
LOGGED AT 12-AUG-75 10:14:54
                               ERROR NUMBER 1.
TAPE PARAMETERS
                        DT1
      UNIT NAME
      VOLUME LABEL
      VOLUME OWNER UIC DEVICE TYPE
                         1,1
                        TU56 UNIT-1 CONTRULLER-0
TAPE REGISTERS AT ERROR TIME
                  004037 SELECTION ERROR
      TEST
                  100702
      TCCM
      TCWC
                  177400
      TCBA
                  113534
      TCDT
                  100770
ERROR DIAGNOSIS
RECOVERED
SELECTION ERROR
RETRIES PERFORMED
                 0.
USER TASK PARAMETERS
      TASK NAME
                        F11ACP
      TASK UIC
                        2,2
      PHYSICAL START ADDRESS 102200
      USER FUNCTION REQUESTED
                               READ
                                    (1000)
      FUNCTION INTERPRETED FROM REGISTERS
                                     READ BLOCK NUMBER FORWARD
            PHYSICAL BUFFER ADDRESS START
                                     113534
            TRANSFER SIZE IN BYTES
                                     1000
            LOGICAL BLOCK NUMBER AT TO GO
                                     13
            COUNT OF IO IN PROGRESS
                                     1.
            COUNT OF IO PENDING
                                     1.
```

Figure 8-1 Device-Specific Error Log

7

NUMBER OF FUNCTIONS ISSUED TO THIS UNIT

VECTORS WITH ACTIVE IO

```
SYSTEM ERROR REPORT SUMMARY (SYSTEM)
COMMAND LINE USED
LP:=/bREAKDOWN:ALLALL/DETAIL
REPORT FILE ENVIRONMENT
         SOURCE FILE
                                    SYO: [1,6] ERROR.SYS
         OUTPUT FILE
                                    LP: [200, 200] ALLALL.LST
         DATE OF FIRST ENTRY
                                    12-AUG-75 10:13:45
         DATE OF LAST ENTRY
                                    12-AUG-75 10:21:01
ENTRIES PROCESSED
                                               19.
                                                4.
ENTRIES MISSING
                                                Ø.
UNKNOWN ENTRIES ENCOUNTERED
                                                ø,
UNKNOWN DEVICES ENCOUNTEED
FIELD FORMAT ERRORS ENCOUNTERED DEVICE ERRORS PROCESSED
                                                0.
                                               14.
                                                0.
UNDEFINED INTERRUPTS PROCESSED
UNDF. INTR. MISSED DUE TO UNDF. INTR. PROCESSING TRAPS THROUGH LOCATION @ PROCESSED 0.
                                                          0.
DEVICE TIMEOUTS PROCESSED
                                                Ø.
MEMORY MANAGEMENT ERRORS ENCOUNTERED
SYSTEM POWER FAILS ENCOUNTERED
SYSTEM PARITY ERRORS ENCOUNTERED
                                                1.
         REPRODUCIBLE
         NON REPRODUCIBLE
SYSTEM LUADS
                                                2.
SYSTEM ERROR REPORT SUMMARY
                                 (TAPE)
TU56 UNIT-1 CONTROLLER-0
```

Figure 8-2 Summary Error Log

HARD

SOFT

0.

14.

## 8.1 FUNCTIONAL DESCRIPTION

Error logging consists of two distinct functions. The first function is the gathering of information pertinent to the errors that occur and the second is error analysis and the creating of a list file. These functions are performed by three tasks: ERRLOG, PSE, and SYE.

ERRLOG gathers volatile information when a device error occurs. It places this information in a temporary file named ERR.TMP under UFD [1,6] on the system device or a user-specified device. To specify a device other than SY, type the following command.

MCR>REA ERRLOG 4 dev: dev: is the device on which ERR.TMP is to be placed.

When a report of errors is desired, the preanalyzer (PSE) and the analyzer (SYE) tasks are run. When PSE starts, it attempts to open the ERROR.TMP file that has the highest version number. If PSE cannot locate an ERROR.TMP file, it requests that ERRLOG close its current log file (ERR.TMP) and rename that file to ERRLOG.TMP.

When the next error occurs, ERRLOG creates a new ERR.TMP file and continues logging errors. The preanalyzer uses the information in ERROR.TMP to produce a formatted file. The default filename and UFD are ERROR.SYS and [1,6], respectively. However, any operator-specified name and UFD can be used. When the analyzer is run, it uses that formatted file to produce a list file as illustrated in Figures 8-1 and 8-2.

ERROR.SYS remains on disk until the system manager deletes it because it contains information that can be reprocessed by the analyzer or processed by user-written tasks to provide a report with different content. Appendix C provides the layout of records in ERROR.SYS.

## 8.2 OPERATIONAL INFORMATION

This section provides operating procedures for the three error logging and analysis tasks (ERRLOG, PSE, and SYE) and for the task that terminates error logging (ERROFF).

# 8.2.1 Running ERRLOG

ERRLOG must be running in order for error statistics to be accumulated and for the raw data file to be passed to the preanalyzer (PSE), if an ERROR.TMP file is not ready for PSE to process. Normally, ERRLOG is installed during system generation. To run ERRLOG, type the following command to MCR and press ALTMODE.

MCR>RUN ERRLOG

The task responds with the following message.

INPUT MINIMUM NUMBER OF ERRORS CAPABLE OF BEING LOGGED IN A 5 SECOND PERIOD "CARRIAGE RETURN." THIS VALUE SHOULD NOT EXCEED 5.
IF ERROR LOGGING NOT WANTED INPUT "CONTROL Z." 5 SECOND ERROR RATE =

At this point, type a value in the range from 1 through 5 and press RETURN. The value indicates the number of 72-word nodes to be assigned permanently to the error log task. The number of nodes

allocated determines the minimum number of errors that can be logged by ERRLOG without a disk write.

Errors can occur more rapidly than they can be handled by ERRLOG. In this case, the summary report contains an entry under the heading NUMBER OF ERRORS MISSED. This entry indicates the number of errors that were not logged due to insufficient node space. Because the device-specific reports provide sequential numbers for errors, the user can determine at which point errors occurred but were not logged.

If a large amount of node space is allocated, it may adversely affect the ability of other tasks to acquire enough dynamic memory to run. The number of nodes required is determined by the system manager and varies from installation to installation.

- 8.2.1.1 Abnormal Termination of ERRLOG The ERRLOG task terminates automatically in three cases:
  - 1. When the desired number of nodes cannot be obtained,
  - 2. If the error logging device used by ERRLOG becomes full,
  - 3. If an error occurs when writing to the logging device.

#### NOTE

For procedures to terminate error logging normally, see Section 8.3.

The ERRLOG task terminates at task startup when the task cannot obtain the number of nodes specified by the user. When this situation occurs, the following message is printed on the console.

"ERRLOG" TASK FAILED TO PICK LARGE ENOUGH ERROR LOG NODE BUFFER. "ERRLOG" TASK TERMINATED. IF YOU WANT TO TRY AGAIN EXECUTE THE FOLLOWING SEQUENCE.

#### RUN ERRLOG

To attempt to run ERRLOG again, type the request to run ERRLOG again. When the request for the number of nodes is printed respond with a smaller number to 5 SECOND ERROR RATE = .

The second case that causes ERRLOG to terminate is when the error logging device becomes full. The following message is printed on the console.

ERROR LOGGING DEVICE device and unit number FULL. "ERRLOG' TASK TERMINATED. IF YOU WISH TO CONTINUE LOGGING ERRORS EXECUTE THE FOLLOWING SEQUENCE.

REA ERRLOG 4 device and unit number RUN ERRLOG

The REA command is detailed in the RSX-11D User's Guide.

Prior to reassigning the logging device, the new device to which it is to be reassigned must be given a UFD of [1,6] if it does not have one already. The UFD must have the following access rights [RWED,RWED,RWED].

Type the following command to place the UFD on the disk.

UFD dev:[1,6]/PRO=[RWED,RWED,RWED]

The UFD command is detailed in the RSX-11D User's Guide.

After creating the UFD, type the sequence provided in the console printout.

It is suggested that the entire volume be cleared of old files using PIP and VFY (verify). PIP and VFY are described in the  $\frac{RSX-l1}{L}$  Utilities Procedures Manual.

The third case in which ERRLOG terminates occurs if an error occurs while trying to write to the logging device. The following information is printed on the console.

ERROR - xx ON ERROR LOGGING DEVICE device name "ERRLOG" TASK EXITING.
TASK "ERRLOG" TERMINATED
VIA 'EXIT' WITH PENDING I/O

xx is the standard system code as defined in the <a href="IAS/RSX-11D Device Handlers Reference Manual">IAS/RSX-11D Device Handlers Reference Manual</a>. If MO is loaded additional error information is printed on the terminal. If it is desirable to continue error logging, reassign the logging device.

## 8.2.2 Running PSE

The function of the preanalyzer, PSE, is to format the raw data collected by ERRLOG into a file to be processed by SYE. In order to run PSE, PSE must be installed, and the user must be operating under a privileged UIC. In addition, if no ERROR.TMP file is available for PSE to process, the ERRLOG task must be running. ERRLOG can rename the ERR.TMP file and pass it to PSE.

Use the following command to install PSE.

MCR>INS [11,1]PSE

To initiate PSE after it has been installed, type the following command to MCR.

MCR>PSE

The preanalyzer responds with the prompt PSE> and waits for the user to type a command line. The format of the PSE command line follows.

OUTDEV: [ufd] file.ext=indev:

The output file specification is a standard RSX-11D specification.

The input file specification consists only of the input device specification. The file name is always ERROR.TMP and it is under [1,6]. The name is assigned by ERRLOG.

The following defaults are used for omitted portions of the file specifications.

outdev defaults to SY:
ufd defaults to [1,6]
file.ext defaults to ERROR.SYS
indev defaults to SY:

If the default values are to be used, press RETURN in response to the PSE prompt. When PSE prompts again, type CTRL Z to return to MCR.

# 8.2.3 Running SYE

The analyzer produces an error report in the form of a printed listing or a listing file.

Before SYE can run, the user must be operating at a privileged terminal and SYE must be installed. Type the following command to install SYE.

MCR>INS [11,1]SYE

Once SYE is installed, type the following command to run it.

MCR>SYE

The analyzer responds with the following prompt SYE> and waits for the user to type a command line. The format of the SYE command line follows.

outdev:[ufd]=index:[ufd]file.ext/switchl.../switchn

The output filename corresponds to the values specified for xxx and yyy for the /BR: switch described below. The file type is LST. The output filename and type should not be included in the command line.

The input file specification is a standard RSX-llD file specification. The input device, UFD, filename, and extension must be identical with the output file specification used when running PSE. It is the output of PSE that SYE analyzes.

The following switches can be used as part of the input file specification.

/BR:xxxyyy

is the breakout switch that determines what information is to be included in the report.

xxx can have one of the following values:

- ALL indicates that error statistics for all disk, magnetic tape, and DECtape units are to be included.
- DEV indicates that only device errors are to be itemized in the report; that is, no memory parity errors are to be included.
- SYS indicates that only memory parity errors are to be included in the report.
- DSK indicates that error statistics for all disk units are to be included.
- MAG indicates that error statistics for all tape devices, both magnetic tape and DECtape, are to be included.

ALL is the default value for xxx.

yyy can have one of the following values:

ALL indicates that both the device-specific and the summary information is to be included in the report.

SUM indicates that only the summary information is to be included in the report.

DSK indicates that only disk errors are to be included in the summary report.

MAG indicates that only magnetic tape errors are to be included in the summary report.

SUM is the default for yyy.

The output filename is formed from the xxxyyy portion of the /BR switch. The output file type is .LST.

/ID:name

indicates that the report of errors is to contain only those errors that occurred while a specified volume is mounted. The value name provides the volume identification.

The volume can be mounted using the /OVE switch if the volume identification is not known. The volume information printed on the console gives the volume identification once the device is mounted.

/DV:devn

indicates that the report is to contain only those errors that occur on a specified device type or on a specified unit. For example, if devn is specified as DK, error statistics for all RK03 or RK05 units are provided. If devn is specified as DK1, error statistics for RK03 or RK05 unit 1 are provided.

/BG:time:date

indicates that only those errors that occur after the specified time and date are to be included in the report. The format of the time and date specification follows.

hh:mm:ss:dd:mmm:yy

All numbers are decimal and all six fields must be specified.

/ED:time:date

indicates that only those errors that occurred on or before the specified time and date are to be included in the report. Time and date have the same format as in the /BG switch.

/DE

indicates that a more detailed report than that produced using the defaults is desired. When /DE is specified additional information is included if it is available:

- Device error register bit description
- Task and function that caused the entry to be logged
- Statistical information about the device
- Information relating to concurrent activity on the UNIBUS

The default is NODE or -DE.

See Appendix A for examples of the report when /DE is specified.

/-SU

indicates that the summary report is not to be produced. The default is  $/\mathrm{SU}$ .

The following are the default values for the SYE command string.

SY0: [user uic]ALLSUM.LST=SY0: [1,6]ERROR.SYS/BR:ALLSUM/NODE/SU

8.2.4 Terminating Error Logging - Use the following procedures to terminate the error logging task (ERRLOG) in an orderly fashion.

MCR>INS [11,1]ERROFF

MCR>RUN ERROFF

One of two messages is printed on the terminal when ERROFF has terminated error logging. The following message is printed to indicate that ERROFF executed successfully.

ERRLOG TASK TERMINATED

The following message is printed to indicate that ERRLOG was not active, and therefore, was not terminated.

ERRLOG TASK NOT ACTIVE

## 8.3 ERROR MESSAGES

Both the preanalyzer (PSE) and the analyzer (SYE) issue error messages to inform the user of operational difficulties.

## 8.3.1 PSE Error Messages

After each error message, PSE prompts again for a new command line.

#### COMMAND STRING PARSE ERROR

A syntax or semantic error was encountered while examining the input command string to PSE. PSE prompts again for a new command line.

Type a corrected version of the command line.

#### DELETE ERROR

After the preanalyzer processed the input file ERROR.TMP, it was unable to delete it.

Attempt to delete the file using PIP.

#### INPUT FILE ERROR

An error was encountered while trying to open or obtain data from the input file ERROR.TMP. ERROR.TMP is closed, processing is terminated, and the input file is not deleted.

PSE prompts again for a new command string. If this fails, delete the file.

## NO ERROR FILES FROM SYSTEM

The preanalyzer is unable to locate a file named ERROR.TMP. This message can be caused by one of the following situations.

- No errors have occurred. Therefore ERRLOG has no raw data file to pass to PSE.
- ERRLOG is not running and, therefore, cannot rename the ERR.TMP file to ERROR.TMP and pass it to PSE.
- ERRLOG is not writing to the specified device. Check the LUN assignment of ERRLOG. The logging device is assigned to LUN 4.

If the cause of the message is that ERRLOG is not running, follow the procedures in Section 2.2.1 to run the task.

If the cause of the message is that the ERROR.TMP file resides on a device other than the one specified, use the device assigned to LUN 4 of ERRLOG as the index for the PSE command line.

#### OUTPUT FILE ERROR

An error was encountered while working with the PSE output file. Both the input and output files are closed. ERROR.TMP is not deleted.

Try to rerun PSE.

## PREANALYZER OUTPUT DEVICE FULL

The output device became full while PSE was writing data to the output file. Both the input and output files are closed. ERROR.TMP is not deleted.

Rerun PSE using a different output volume.

## UNABLE TO CLOSE INPUT FILE

PSE is unable to close the file ERROR.TMP. The file is not deleted. File processing is terminated.

Use PIP to delete the file.

## UNABLE TO CLOSE OUTPUT FILE

PSE is unable to close the output file.

Rerun PSE or use PIP to delete the file.

## 8.3.2 SYE Error Messages

SYE--COMMAND STRING ERROR portion of the string in error

The format convention within a particular portion of the command string is violated. No files remain open. SYE issues a prompt.

Correct the error and type the command.

# SYE--COMMAND STRING ERROR ERROR NUMBER n

The command string interpreter detects an error while attempting to get a command line. n is a CSI error code. Refer to the  $\underline{IAS/RSX-l1}$   $\underline{I/O}$  Operations Reference Manual to determine the meaning of n. No files remain open. SYE issues a prompt.

Correct the error and type the command.

# SYE--COMMAND STRING SYNTAX ERROR command string typed

The proper format was not used in the command string. No files are open. SYE issues a prompt.

Type the corrected command.

SYE--DEVICE ERROR INPUT FILE FATAL ERROR - n

SYE was attempting to obtain further information from the input file but could not get the next record. n is an FCS error code. Refer to the IAS/RSX-11 I/O Operations Reference Manual to determine the meaning of n.

Both the input and output files are closed. SYE issues a prompt for the next command.

SYE--DEVICE ERROR OUTPUT FILE FATAL ERROR - n

SYE was unable to write information in the output file. n is an FCS error code. Refer to the <u>IAS/RSX-11 I/O Operations Reference Manual</u> to determine the meaning of n.

Both the input and output files are closed. SYE issues a prompt for the next command.

SYE--ILLEGAL BREAKOUT SWITCH /BR:xxxyyy

SYE issues this message when the operator attempts to request a breakout of the input file that is not legal. No files remain open. SYE issus a prompt for another command.

Retype the command with a correct use of /BR:.

SYE-- OPEN FAILURE ON INPUT FILE FATAL ERROR - n

SYE was unable to open the input file. No files remain open. n is an FCS error code. Refer to the  $\underline{IAS/RSX-11}$  I/O Operations Reference Manual to determine the meaning of n.

SYE--OPEN FAILURE ON OUTPUT FILE FATAL ERROR - n

SYE was unable to open the output file. No files remain open. n is an FCS error code. Refer to the <a href="IAS/RSX-11">IAS/RSX-11</a> I/O Operations Reference Manual to determine the meaning of n.

SYE--n. PAGES text

This message does not indicate an error; rather, it indicates the number of pages (n) in the finished report. The value "text" indicates the filename for future reference.

SYE--SUMMARY TABLE OVERFLOW REPORT CONTINUES WITHOUT SUMMARIES

Summary table overflow occurred because more devices received errors than the analyzer is built to handle.

Use the breakout switches to reduce the number of devices included in the summary.

#### CHAPTER 9

#### SYSTEM MANAGEMENT NOTES

This chapter contains information that can be used to optimize system performance and to alter certain system defaults.

## 9.1 ANCILLARY CONTROL PROCESSORS AND THEIR USE

An ancillary control processor (ACP) is a file processing task that interfaces between a user task and FCS on the one hand and a device handler on the other. The device handler performs basic operations, such as read and write, without aid of an ACP. However, file processing functions, such as create, delete, and access, are performed by an ACP task associated with the device. Not all devices have ancillary control processors associated with them.

The ancillary control processor to be associated with a device is specified either explicitly or by default during system generation as a parameter in the DEV directive. RSX-11D supplies three ACPs:

F11ACP for Files-11 devices (disk and DECtape),

MTAACP for magnetic tape handling,

DTAACP for installations that require significant amounts of DECtape processing.

F11ACP and MTAACP are the default ACPs for Files-11 devices and magnetic tape, respectively. DTAACP is used for DECtape instead of F11ACP if so specified during system generation. If desired, a user-written ACP can be associated with any device during system generation. A user-written ACP can be specified instead of a standard RSX-11D ACP or to perform special functions on a device that does not have a standard ACP associated with it. Only one ACP can be associated with each device.

System generation establishes the default ACP to be associated with a device when it is mounted. However, use of the /ACP switch with the MOUNT command can override the default ACP for a device.

All ACPs must have a 6-character task name that ends with the letters ACP. ACPs must be installed either during system generation or by means of the INSTALL command to MCR.

### 9.1.1 Files-11 ACPs

Two versions of FllACP are supplied with RSX-11D: FCP.TSK and BIGFCP.TSK. They appear under different filenames in UFD [11,1].

The file named FCP.TSK is a minimum size, heavily overlaid version occupying approximately 2.4K words of memory. It is recommended for small systems with limited space and where high file processing throughput is not essential.

The file named BIGFCP.TSK is minimally overlaid and has additional buffer space. It incurs almost no overlay switching during normal operations and provides some optimization of disk accesses due to its additional buffer. BIGFCP.TSK occupies approximately 4.5K words of memory.

Both FCP.TSK and BIGFCP.TSK are built with a task name of FllACP.

The file DTAACP.TSK is a version of FllACP that is optimized for DECtape operation. It is heavily overlaid, but has the additional buffer. It occupies approximately 2.6K words of memory. DTAACP.TSK is built with a task name of DTAACP.

In large systems with a heavy file processing requirement (for example, a multiuser program development system), multiple versions of an ACP can be installed to achieve some parallelism in file processing. Each ACP task processes all functions in a strictly sequential manner. Having two versions of FllACP installed in the system allows each version to service requests for a different set of devices.

For example, in a system with two RP04 drives and two RK05 drives, the system manager could specify F11ACP as the ACP for the RP04 drives and RKAACP as the ACP for the RK05 drives during system generation. By installing BIGFCP twice, the system manager could provide each device type with its own ACP. File processing then occurs concurrently on RP04s and RK05s. In addition, the extra buffer space available in the two copies of the file processor, whether it be FCP or BIGFCP, saves a certain percentage of disk accesses.

An installation that does a significant amount of Files-11 DECtape I/O in a multiuser environment should consider using DTAACP for DECtape and FlIACP for disks. Use of DTAACP ensures that the processing of DECtapes does not interfere with disk file processing.

## 9.1.2 Magnetic Tape ACP

The magnetic tape ACP is named MTAACP and provides support of ANSI magnetic tape file structures. Unlike FllACP, it is both nonfixable and checkpointable.

MTAACP does not need to be fixed in memory for performance reasons. When a magnetic tape is mounted, MTAACP is loaded into memory if it is not already there. It does not exit from the system after it services a request; the only time MTAACP exits is when no magnetic tape volumes are mounted and no outstanding I/O requests remain to be serviced. The result is that once a magnetic tape volume is mounted, which requires use of MTAACP, the ACP task is effectively locked in memory.

#### SYSTEM MANAGEMENT NOTES

MTAACP is built with a priority of 200. Because it is checkpointable, a deadlock can occur on either a normal I/O request or I/O rundown in the following situation:

- A task with a higher priority than MTAACP causes it to be checkpointed,
- That task then needs the services of MTAACP but insufficient memory remains to load MTAACP.

## 9.1.2.1 File and Record Attribute Support under MTAACP

The ANSI magnetic tape file system (MTAACP) supports the FCS record types and attributes with the exception of symbolic offset FD.BLK in the FDB. FD.BLK is set unconditionally to indicate that records cannot cross block boundaries. The result is that, for files containing variable-length records, the block length (that is, buffer size) must be equal to or greater than the maximum record size plus four. Therefore, the default block size of 512 bytes yields a maximum variable-length record size of 508 bytes.

MTAACP does not support the FCS file characteristics (attributes) in that it treats all files as noncontiguous. If a file on magnetic tape is to be contiguous on disk or DECtape, the following steps must be performed using PIP:

- Transfer the file from magnetic tape to either a Files-11 disk or DECtape to cause FCS to determine the real size of the file in blocks.
- Transfer the file from the disk or DECtape to the desired device using the /CO switch in the output file specification to PIP.

# 9.2 /LRU SWITCH ON INITVOL AND MOUNT

Another optimization feature is the /LRU switch that can be specified on the INITVOL and MOUNT commands to MCR. All versions of F11ACP maintain in memory a list of recently referred to directories. Each mounted volume has an associated list. If a directory is present in the list, the overhead involved in gaining access to the directory is almost completely eliminated. Entries in the list are replaced on a least-frequently-used basis.

The /LRU switch specifies the number of entries that are to be retained in the list for a particular volume while it is mounted. The INITVOL and MOUNT default value is three, which is an appropriate value for a single-user volume. In general, the /LRU value should be a few units greater than the expected number of concurrent users of the volume.

# 9.3 TAILORING FllaCP

The file primitives (F11ACP) in RSX-11D Version 6B can be tailored to suit the needs of a particular installation in several ways. All involve balancing certain size/performance tradeoffs.

## 9.3.1 Overlay Structure

The file primitives are almost always run overlayed. Two task builder command files and overlay description files are provided for two different overlay structures. The smallest FCP (FCP.CMD and FCP.ODL) occupies roughly 2.4K words of memory. It is structured with one overlay for each FCP function and is recommended for small systems where space is limited and high file processing throughput is not important. A larger overlay structure is also provided (BIGFCP.CMD and BIGFCP.ODL). It incurs almost no overlay switching during normal file processing; overlays are used only for errors and exceptions.

## 9.3.2 Additional Directory Buffer

A separate buffer for directories can be allocated in FCP with either overlay structure. The presence of this buffer speeds up operation when repeated references are made to the same directory; for example, such as file purges and wild card operations. The directory buffer is allocated by including the following statement in the task builder command file to build FCP:

EXTSCT=\$\$BUF1:1006

## 9.3.3 Extended Directory Buffer

If a significant percentage of the directories used in a system are large (greater than 30 entries), an additional performance improvement can be realized by increasing the size of the directory buffer. Then fewer physical disk accesses are necessary to search a large directory. The directory buffer should be extended in units of one block, expressed in multiples of 1000 bytes (octal). For example, the following task builder statement:

EXTSCT=\$\$BUF3:2000

adds two blocks to the directory buffer, allowing the file system to handle directories in three block units. Extending the directory buffer can be done only if a separate directory buffer has also been allocated as described above in Section 9.3.2.

## 9.3.4 Additional Storage for FCB's

The file primitives use one file control block (FCB) for each file that is currently open in the system. In addition, each directory LRU list entry occupies one FCB, and one FCB is used for each mounted volume for the index file. Therefore, the total number of FCBs allocated in the system is the sum of the number of mounted volumes plus the sum of all mounted volumes´ LRU limits plus the total number of open files. Each FCB is 21 words in length. FCBs are normally allocated in the system node pool, where they occupy 24 words apiece. Pool space is allocated in 8-word units. Clearly, in a large multiuser system FCBs can take up a substantial amount of pool space (several K, in fact). The system manager can allocate space for FCBs within the task body of the file primitives with a task builder command of the following form:

EXTSCT=\$\$FRE1:2000

#### SYSTEM MANAGEMENT NOTES

The space is expressed in octal bytes. Each unit of 1000 has space for approximately 21 FCBs. FCP allocates FCBs out of this area until it is full; it then uses the system node pool.

NOTE

In the system as distributed, FCP has no internal FCB space and BIGFCP has 2000 bytes of internal space.

## 9.4 REASSIGNMENT OF TASK BUILDING AND CROSS REFERENCE LUNS

Both the Task Builder and the cross reference task (CRF) use dynamic memory for symbol processing. System performance when TKB or CRF execute can be enhanced given appropriate resources.

TKB uses LUN 8 for its dynamic workfile, while CRF uses LUN 7. In the released system, these logical unit numbers are assigned to the device on which the TKB and CRF tasks reside. Both tasks are overlaid so that, in any configuration, it may be advantageous to reassign the logical unit numbers to another device.

Example:

MCR>REA CRF... 7 DPl:

## 9.5 COLUMN WIDTH FOR CRF OUTPUT

The released CRF is built with 132-column line output.

The column width for printed output from CRF can be altered by rebuilding the cross reference task. The task build command file for CRF resides under UFD [11,32] on the auxiliary distribution medium and is called CRFBLD.CMD. It contains the following line:

EXTSCT=\$\$RCB0:204

This line establishes a 132-column line width. Omission of this option by inserting a semicolon (;) before the line causes the building of a CRF with an 80-column line output.

## 9.6 CHANGING TASK BUILDER DEFAULTS

The RSX-llD Task Builder contains default values that are used when switches and options are omitted from a command to the Task Builder. These switches and options and their defaults are detailed in the RSX-llD Task Builder Reference Manual.

Some installations may prefer a different set of defaults in TKB: in particular, a default of "task unaccountable" may be desirable. Defaults can be changed by rebuilding the Task Builder with a global patch applied.

#### SYSTEM MANAGEMENT NOTES

To apply the patch, three files must be obtained from UFD [11,11] of the auxiliary distribution medium: TKBBLD.CMD, TKBODL.ODL, and TKB.OLB. These files are the build command file, the overlay-descriptor file, and the relocatable object library of all TKB modules, respectively.

The list below provides the available task image attributes and the bit pattern required to set them. The default in the released version of the Task Builder was achieved by initializing the task image with SW\$FX!SW\$FP.

```
; TASK CHECKPOINTABLE (Ø=YES)
SW$CP==100000
SWSFP==040000
                      ; TASK USES FLOATING POINT (1=YES)
SWSDA==020000
                      ; TASK DEBUGGING AID SPECIFIED (1=YES)
SWSPI==001000
                      ; PIC OUTPUT (1=YES) (ALSO APPLIES TO STB)
                      ; TASK IS PRIVILEGED (1=YES)
; SET TRACE BIT IN PS (1=YES)
SWSPR==000400
SW$TR==000200
                       ; BUILD TASK IMAGE WITHOUT HEADER (1=YES)
SW$NH==010000
SWSFX==004000
                       ; TASK IS FIXABLE (0=YES)
                       ; TASK ABORTABLE (Ø=YES)
SW$AB==000020
                       ; TASK DISABLABLE (Ø=YES)
SWSDS==000010
                       : +++016 MULTI-USER SWITCH (1=YES)
SW$MU==000002
                       ; +++016 ACCOUNTING SWITCH (1=YES)
SW$TA==000001
```

Referring to the list above, the defaults produce a pattern of 44000.

To build a TKB that has these defaults (not fixable, floating point, etc.) plus the task accountable default, the following pattern is required:

#### SW\$FX!SW\$FP!SW\$TA = 44001

Having calculated the required default bit pattern from the list above, edit the file [11,11]TKBBLD.CMD to insert the option in the following form:

# GBLPAT=P11OVR:\$MRFLG+134:value

#### CHAPTER 10

#### EXECUTIVE DEBUGGING AIDS

RSX-11D provides two debugging aids that can be incorporated into the Executive:

- 1. Executive ODT
- 2. Executive trace.

## 10.1 EXECUTIVE ODT

A simplified version of ODT can be task built into the Executive if needed for the debugging of user-initiated modifications to the Executive. This version of ODT can perform the following functions:

- 1. Set and remove breakboints,
- 2. Examine the content of a location and modify the content,
- 3. Search for bit patterns.

Executive ODT operates only in Kernel space. For example, breakpoints can be set only in Kernel space and only Kernel locations can be examined.

## 10.1.1 Building ODT into the Executive

Use the following procedures to build ODT into the Executive and generate the new system.

- 1. Log onto the system under [1,1].
- In the file [11,15]TKB15.CMD, locate the line [11,15]NODT and change it to [11,15]ODT.
- 3. Rebuild the Executive using [11,15] TKD15.CMD as in the following command.

#### TKB>[11,15]TKB15

- 4. Rebuild the bootstraps using the command file [11,17]BOOTSBLD.CMD.
- 5. Perform a system generation using an appropriate command file, examples of which are distributed under UFD[11,17].

After phase 1 of system generation when the disk is booted, the following message is printed at the terminal.

#### EXECUTIVE DEBUGGING AIDS

S/T RSX ODT V7.00

- 6. Press CTRL/C to initiate phase 2 of system generation.
- 7. Follow the formal procedures for saving the system at the end of phase 2 as described in the RSX-11D System Generation Reference Manual.

# 10.1.2 Using Executive ODT

NOTE

It is not possible to use both Executive ODT and user task ODT at the same time.

Once the system generation is completed, the ODT message (S/T RSX EXEC ODT V.07) is printed at the terminal each time the system is bootstrapped or a power failure recovery occurs. When the message is printed, it is followed by the ODT prompt, which is an asterisk (\*). The user can perform either of the following in response to the ODT prompt:

- 1. Interact with ODT,
- 2. Press CTRL/C to allow normal system operation.

Before CTRL C is depressed, the user can issue the commands to set or remove breakpoints, examine or modify the contents of locations, or search for bit patterns. These commands are detailed in the IAS/RSX-11 ODT Manual. Examples are provided below.

Once the necessary initial set-up for Executive ODT is complete, press CTRL C to obtain MCR. Normal RSX-11D operation begins and continues until an instruction in a breakpoint location is executed. Then ODT prompts with an asterisk. At that point, any of the commands to ODT named in Section 10.1 can be used. When ready to continue normal RSX-11D operation, use the P command as described in the  $\frac{IAS/RSX-11}{IAS/RSX-11}$  ODT Manual.

- 10.1.2.1 <u>Setting a Breakpoint</u> In order to set a breakpoint, the user needs a listing of [111,15] EXEC.MAP and a listing of each module in which a breakpoint is to be set. This is similar to the process used with task ODT. The following is an example of setting a breakpoint at location 3612.
  - \*3612;OB The asterisk (\*) is the ODT prompt.
- 10.1.2.2 Examining and Modifying a Location The following command causes ODT to print the contents of location 3716 on the terminal and allows the user to modify the content. Location 3716 contains 000001; the user types 4 to indicate that ODT is to change the content at the location to 000004.
  - \*3716/0000001 4

#### EXECUTIVE DEBUGGING AIDS

10.1.2.3 <u>Sample Use of ODT</u> - The following example illustrates the use of Executive ODT.

S/T RSX ODT V7.00 ;SYSTEM IS BOOTED \*3612;OB ;SET INITIAL BREAKPOINT \* ^C ; PRESS CTRL/C ; NORMAL SYSTEM OPERATION COMMENCES ; MOUNT SYSTEM DISK RSX-11D V6.2 MCR>MOU DKØ:/OVR ;LOG ONTO SYSTEM MCR>HEL [1,1] MCR>INS TEST ; INSTALL TASK NAMED TEST MCR>RUN TEST(\$) ; RUN TEST AND DEPRESS ALTMODE BØ:3612 ;TEST CAUSED EXECUTIVE TO EXECUTE ;INSTRUCTION AT LOCATION 3612 ;SET ANOTHER BREAKPOINT \*3646;1B ; EXAMINE CONTENTS OF REGISTER 1 \*\$1/000542 \*P ; CONTINUE RSX-11D OPERATION MCR>

# 10.2 EXECUTIVE TRACE

The Executive trace facility is part of the resident Executive and is useful for tracing ISRs. The Executive trace allows the user to trace ranges within an ISR and to halt at specified locations.

## 10.2.1 Building Trace into the Executive

The resident Executive in the distributed system does not contain the trace facility. To include trace, the Executive must be rebuilt according to the steps below.

- 1. Put UFD [11,15] on disk and set it as the current UIC.
- 2. Use FLX to transfer the [11,15] files from the distribution binary medium.
- 3. Repeat steps 1 and 2 for UFD [11,17].
- 4. Modify [11,15]TKB15.CMD by changing the line [11,15]NTRACE to [11,15]TRACE.
- Task Build the Executive using TKB15.CMD and task build the bootstraps using [11,17]BOOTSBLD.CMD.
- 6. Change the system generation phase 1 command file to reflect the increased size of the Executive.
- 7. Generate the new system.

## 10.2.2 Using Executive Trace

Once the system generation is completed, use the steps listed below to cause trace to execute.

- 1. Find the address of the global symbol .DBTRA in the EXEC.STB listing which can be obtained from [11,15] EXEC.MAP.
- 2. Using the OPE command to MCR, set up to four trace ranges in the four word-pairs starting at location .DBTRA-30. Addresses are in Kernel virtual space. The following is an example of using the OPE command to specify one range.

MCR>OPE .dbtra-30/KNL 001160 000000/4140 001162 000000/4147 001164 000000/\$ Type in actual address of .DBTRA-30

Depress ALTMODE to return to MCR.

Alternatively, set console switch  $\emptyset$  up to trace the entire ISR.

- 3. If desired, set up to five halt addresses in the four word-pairs starting at .DBTRA-40 or set one halt address in console switches 15 through 1. Use the OPE command as in step 2 to set the halt addresses.
- 4. Change Kernel address 14 to contain the address of .DBTRA as follows:

MCR>OPE 14/KNL 000014 nnnnnn/xxxxxx

nnnnnn = present contents of 14.

5. Set the T-bit (bit 4) in the second word of the interrupt vector for the interrupt service routine to be traced; for example in the DP ISR, set bit 4 in vector 254.

The trace printed on CL is essentially in the same format as that for the task trace package; only the relative PC is omitted. Refer to the task trace chapter of the <a href="IAS/RSX-11">IAS/RSX-11</a> ODT Manual.

If one of the requested halt addresses is encountered, the computer halts with the requested address in the data lights. Execution can be resumed by pressing CONT.

#### CHAPTER 11

#### PARITY SUPPORT

Under RSX-llD, the handling of errors is divided into three categories:

- 1. Main memory parity support,
- 2. Cache memory parity support on a PDP-11/70,
- 3. Additional PDP-11/70 hardware error support.

For a description of cache memory, refer to the  $\underline{\texttt{PDP-11/70 Processor}}$  Handbook.

## 11.1 MAIN MEMORY PARITY SUPPORT

Main memory parity errors are an indication of a memory malfunction. When a parity error occurs, it causes a trap to the RSX-llD Executive. The Executive determines whether the error was fatal (nonrecoverable) or not and then takes appropriate action. In the PDP-ll/40, 11/45, and 11/50, all parity errors are fatal and are treated as main memory errors.

The PDP-11/70 hardware recognizes two types of main memory errors: address and data. The Executive handles both types of errors identically, although errors are logged as being one of the two types; refer to Chapter 8 for information on the logging of errors.

When a fatal main memory parity error occurs, the Executive terminates the task executing where the error occurred unless that task has a hardware priority greater than three. The aborted task remains in memory and no further access to or execution of the task is allowed. Locking the task in memory also makes the entire area of memory allocated to the task inaccessible to the system or other tasks. When the Executive terminates a task, an error message is printed on CO and on the terminal that initiated the task.

#### PARITY SUPPORT

If the Executive encounters any of the following conditions, it halts system operation; that is, a system crash occurs:

- 1. The hardware priority of the task executing where the error occurred is greater than three,
- 2. No task is executing currently,
- 3. The error occurred in Kernel mode.

The only difference between parity handling on the PDP-11/70 and other PDP-11's is that the 11/70 provides the Executive with an indication of whether the error was fatal. Fatal errors are handled as described above. Nonfatal errors are transparent to the task executing where the error occurred, but are logged by the Executive as described in Chapter 8.

#### 11.2 CACHE PARITY SUPPORT (PDP-11/70 ONLY)

Because cache contains a copy of information also maintained in main memory, it can recover from cache parity errors by obtaining the information from main memory. Fatal parity errors are never returned for cache. Rather, the Executive performs the following sequence:

- 1. Determines the cache group in which the error occurred,
- 2. Determines whether the group has exceeded the number of errors allowed in a 1-minute period,
- 3. Turns the group off if the limit has been exceeded,
- 4. Logs the error.

This process is transparent to the task using the cache group in which the error occurs. However, program execution speed may be affected.

The default number of errors allowed in a 1-minute period is 50 (decimal). This number can be changed by the system manager as described in Section 11.4.

The PDP-11/70 hardware recognizes two types of cache errors: address and data. The Executive handles both types of errors identically, although errors are logged as being of the two types; refer to Chapter 8 for information on error logging.

#### PARITY SUPPORT

#### 11.3 ADDITIONAL PDP-11/70 HARDWARE ERROR SUPPORT

The RSX-llD Executive handles three other types of hardware errors that the PDP-11/70 reports:

Nonexistent memory,

UNIBUS error,

CPU/UNIBUS abort.

## 11.3.1 Nonexistent Memory

A nonexistent memory error implies that faulty mapping or a bad address has caused the central processing unit to refer to a nonexistent memory location. This type of error is treated in the same manner as a main memory parity error.

## 11.3.2 UNIBUS Error

A UNIBUS error occurs when bad memory is referred to during a direct memory access (DMA) transfer. Since the task currently executing is probably not the task for which the I/O is being done, the Executive makes no attempt to determine which task issued the I/O request or to locate the bad memory. The error is only logged. Some UNIBUS device controllers detect the memory error and report it at the end of the transfer.

## 11.3.3 CPU/UNIBUS Abort

The CPU/UNIBUS abort occurs when bad data is transferred during a non-DMA I/O transfer or a read of UNIBUS I/O page registers. The Executive terminates the task currently executing because it is the task that issued the I/O request. The error is logged.

#### 11.4 SYSTEM MANAGER'S CONTROL OF PARITY HANDLING

Two MCR functions, SET and MEM, provide the system manager with a means of controlling the effects of parity error handling. Both of these functions are detailed in the RSX-llD User's Guide.

## 11.4.1 SET Functions for Controlling Parity

The SET command is used in conjunction with cache parity. It allows the system manager to turn cache groups on and off and to change the default number of errors that are allowed in a 1-minute period.

#### Examples:

MCR>SET /CAC=75 This command indicates that 75 cache parity errors are to be allowed in a 1-minute period before the Executive turns the group off. The system default is 50.

MCR>SET /CAC=ON1 This command turns on cache group 1.

MCR>SET /CAC=OFØ This command turns off cache group Ø.

The system manager can determine which groups are receiving errors by consulting the error log as described in Chapter 8.

# 11.4.2 MEM Function for Freeing Memory

The MEM command to MCR allows the system manager to free a task locked in memory due to main memory parity errors and to free the memory allocated to it. That task then can be run again; however, the memory can also be used for subsequent task execution and can result in further task terminations.

#### Examples:

MCR>MEM PROG1 This command frees the first task encountered that has the name PROG1.

MCR>MEM PROG1/TI=TT3 This command frees the task PROG1 that was initiated from terminal 3.

MCR>MEM PROG1/AL This command frees all locked versions of PROG1.

#### 11.5 TASK CONTROL OF PARITY HANDLING

The user can avoid having a task terminated by the Executive when a parity error occurs by declaring an SST (synchronous system trap) in the task. SSTs are described in the RSX-llD Executive Reference Manual.

## APPENDIX A

## CDA SAMPLE LISTING

RSX-11D CORE-DUMP ANALYZER V06.2 5-OCT-76 14:25 PAGE 1
CRASH STACK

CAUSE OF CRASH: NON-TASK LEVEL SEGMENT FAULT OR TRAP FOUR

AFTER CRASH: PS=004344 SP(K)=001454 SP(U)=170542

BEFORE CRASH: PC=004456 PS=004301

R0=013173 R1=000001 R2=000001 R3=006000 R4=000774 R5=012667

SR0=000003 SR1=000000 SR2=035576 SR3=000000

#### USER

I S	PACE	DSP	ACE
PDR	PAR	PDR	PAR
000406	000021	000000	000377
600000	000000	000000	007717
000000	000000	000000	002437
600000	000000	000000	047733
000000	000000	000000	007660
000000	000000	000000	127400
000000	000000	000000	007624
000000	000000	000000	006000

## SUPERVISOR

I	S	P	A	C	٤		D	s	P	A	С	Ε	
PDF	₹			P	AR		PD	R			P	AR	
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0000	000		(	30	000	90	000	999		(	000	001	7
000	110		- (	00	006	00	000	010		(	001	021	7
0000	114		(	30	000	90	000	000		•	100	100	7
0000	900		1	00	000	80	000	000		(	00	800	Ø
0000	116		4	00	006	00	000	007		- (	00	100	7
0000	010		- (	20	000	<b>3</b> 0	000	900		(	000	000	0
0000	317		í	30	000	30	000	000		- (	00	000	0

#### KERNEL

I SF	ACE	D 8 P	ACE
PDR	PAR	PDR	PAR
077506	000004	Ø60017	125252
077106	000204	044017	125252
00000	000000	010015	125252
077406	000021	077415	125252
W00116	006402	030013	125252
077406	000602	047010	165542
077506	001002	010011	000000
077506	177600	000000	000000

RSX-11D CORE- KERNEL STACK	DUMP ANALY	ZER V	Ø6 <b>.</b> 2	5-OCT-76	141	25	PAGE 2
KERNEL STAC	CK:						
001000 001020 001040 001060 001100 001120	000000 0	824348 9999 829999 9299 829999 9299 829999 9299 829999 9299	00 000000 00 000000 00 000000	999999 999999 999999	000000 000000 000000 000000 000000	000000 000000 000000 000000	969999 969999 969999 969999
001140 001160 001200 001220 001240	000000 00000 0000000000000000000000000	00000 00000 00000 00000 00000 00000 00000 00000	00 000000 00 000000 00 000000 00 000000 00 000000	000000 000000 000000 000000 000000	999999 999999 999999 999999	999999 999999 999999 999999	808888 808888 808888 808888 808888 808888 808888 808888 808888
001260 001300 001320 001340 001360 001400	000000 0 000000 0 000000 0 060166 0 062574 0	000000 0000 000000 0000 00000 0000 00000 0000 000012 0401 004250 0612	00 00000 00 000000 00 002352 42 071102 30 000400	000000 000000 034110 004240 051312	000000 000000 000000 142040 040140	000000 000000 000000 062352 000001 103360	000000 000000 000000 004240 000400 141740
001420 001440 001460	110254 0	034040 0000 034111 0000 004320 0043	00 015474	035266	034111 035122 034300	145170 000006 000074	147554 004456 174011
RSX=11D CORE- SYSTEM COMMUN		<del>-</del>	<b>06.</b> 2	5-001-76	14:	25	PAGE 3
CRASH OCCURRE	ED AT 14124	1:03 5-OCT	<b>-</b> 76				
.CRTSK = NULL	. TASK						
.FP45X*1 .SERFG*00000	,PDP11*70.	.MSIZE=8	192.=256K 000	.CKFSZ#	515.		
CHECKPOINT BI	LTMAP :						
037500	177777 1	77777 1777	77 177777	177777	177777	177777	177777
SCOM DATA ARE	EAI						
146100 146120 146140	070100 1	00001 0001 77777 1700 00071 0000	01 177777	000002	000000 105721 000012	000400 041104 000114	000062 000062 177546
146160 146200 146220	000100 0 000454 0	000001 0000 026226 0000 00100 0000	00 900000 00 904401	040000 060146	051522 060146	030530 000000 104760	042061 000000 103220
146240	144724 0	0000 00000 0000 00000	00 000000 00 000000	144724 000000	000000 000000 005111	000000 025167	000000 003310
146300 146320 146340	144724 1	100000 0000 17777 0004 177406 1446	02 000602		000000 177600 145100	000003 000016 145174	104760 077406 142240
146360 146400 146420	106060 1	42300 1450 36320 1074 100000 1000	40 107440	111060	000200 111060 037520	145600 111720 001003	145740 111720 000002
146440	000002 0	000017 0454 00001 1451	62 900000	000003	000050 014554	000001 000172	000500 000005

RSX-11D CORE-DUMP ANALYZER 5-0CT-76 14:25 PAGE 4 VØ6.2 ATL SCAN DB... AT: ADDRESS = 145600V/106000R STD ADDRESS = 142100V/102300R REQUESTOR = DB.... LOAD ADDRESS = 00120200

PAR\*SYDISK PRI\*250, I/O IN PROG\*0. I/O PEND\*0. STATUS\*
EVENT FLAGS\*\*000200 000000 MASKS\*\*000007 000000 000000 000000 STATUS=WØ TI=TTØ TASK FLAGS: 000000 100260 146374 142100 142300 000372 000000 000000 001202 000016 142100 000200 000000 000007 000000 000000 000000 145100 145642 145642 000000 000000 000000 145654 145654 145600 145620 145640 TT.... ATL ADDRESS = 100260V/040460R STD ADDRESS = 141740V/102140R REQUESTOR = TT.... LOAD ADDRESS = 06150100 PAR=GEN PRI=248. I/O IN PROG\*0. I/O PEND=0. STATUS=VEVENT FLAGS=000000 040000 MASKS=007377 000000 000000 000000 STATUS#WØ TI#TT2 TASK FLAGS: 000000 101040 145600 141740 144332 000370 000000 000001 001501 000016 141740 000000 040000 007377 000000 000000 000000 145150 100322 100322 000002 027733 000000 100334 100334 100300 100320 LP. . . . -----ATL ADDRESS = 101040V/041240R STD ADDRESS = 104360V/044560R REGUESTOR = ...MFT LOAD ADDRESS = 10225600 PAR=GEN PRI=248. I/O IN PROG=0. I/O PEND=0. STATUS=0 EVENT FLAGS=000000 0000000 MASKS=000003 006401 000000 0000000 TASK FLAGS: 000000 
 101040
 106460
 100260
 102360
 142300
 000370
 000000
 000000
 002256

 101060
 000016
 104360
 000000
 000000
 000000
 000003
 006401
 000000
 000000

 101100
 145150
 101102
 101102
 000001
 141022
 000000
 101114
 101114
 DK.... ATL ADDRESS = 106460V/046660R STD ADDRESS = 103420V/043620R REQUESTOR = ...MFT LOAD ADDRESS = 00251000 PAR=GEN PRI=248. I/O IN PROG=0. I/O PEND=0. STATUS= STATUS#W0 TI#TT0 EVENT FLAGS=000000 0000000 MASKS=000003 006401 000000 000000 TASK FLAGS: 000000 
 196720
 101040
 102360
 142300
 000370
 000000
 000000
 002510

 000016
 103420
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 000003
 006401
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 145150
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 106534
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 106460 106500 106520

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RSX-11D CORE-DUMP ANALYZER VØ6.2 5-OCT-76 14:25 PAGE 8 TASK HEADERS

# DB...

LEADER ADDRESS \* 00120200 ATL ADDRESS \* 145600V/106000R
H.CRI\*000000
USER PDR'S: 035606 000000 000000 000000 000016 077406 077406 077406
USER PAR'S: 001205 000000 000000 000000 000402 000602 0V1002 177600
PS\*174000 PC\*000412
R0\*007034 R1\*112060 R2\*000000 R3\*000001 R4\*100200 R5\*145600 SP\*000356
INITIAL PS\*174017 INITIAL PC\*000100 INITIAL SP\*000100
HEADER SIZE\*000003 DSW\*000001 NUMBER OF LUNS\*6,

## LOGICAL UNIT TABLES

- # DEV WINDOW W.CTL W.VBN W.FCB F.FNUM F.FSEQ F.STAT NAC NLCK
- 3Y0 000000
- 2 SY0 000000
- 3 SY0 000000 4 SY0 000000
- 4 3Y0 000000 5 TIO 000000
- 6 CL0 000000

#### HEADER:

000000	000000	035606	000000	900000	999999	000016	977496	077406
000020	977486	001205	999999	000000	000000	000402	000602	001002
000040	177600	000003	600000	000000	000000	000402	000602	001002
000060	177600	174000	000412	007034	112060	000000	000001	100200
000100	145600	000356	174017	000100	000100	000000	000000	000000
000120	100000	000000	000401	000401	000003	000010	000000	000000
000140	137665	000000	000000	060146	060146	000000	000000	000006
000160	145006	000000	145006	000000	145006	000000	145006	000000
000200	000001	000000	142526	000000	000000	000000	000000	000000
000220	000000	000000	000000	000000	000000	000000	000000	000000
000240	000000	000000	000000	000000	000000	000000	909090	000000
000260	000000	000000	000000	000000	000000	000000	000000	000000

#### TT....

HEADER ADDRESS = 00150100 ATL ADDRESS = 100260V/040460R

H.CRI = 0000000

USER PDR'S: 061206 077406 023106 065106 000016 077406 077406 077406

USER PAR'S: 001503 004227 002034 002103 000402 000602 001002 177600

PS=174000 PC=004114

R0=004010 R1=000044 R2=000000 R3=000001 R4=177777 R5=040040 SP=000172

INITIAL PS=174017 INITIAL PC=000200 INITIAL SP=000200

HEADER SIZE=000002 DSW=000001 NUMBER OF LUNS=3.

LOGICAL UNIT TABLES

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RSX=11D CORE-DUMP ANALYZER VØ6.2 5-0CT-76 14:25 PAGE 19 TPD SCAN

SYDISK

TPD ENTRY ADDRESS = 145100V/105300R

BASE ADDRESS = 00120200 SIZE = 00007700 FLAGS = 000003

ALLOCATION MAP:

SPACE ALLOCATED OCCUPIED BY

00120200-00130100 DB.... (A)

\*\*\*\*\* NO ENTRIES IN MRL \*\*\*\*

\*\*\*\*\* NO ENTRIES IN CTL \*\*\*\*\*

MCR

TPD ENTRY ADDRESS = 145124V/105324R

BASE ADDRESS = 00130100 SIZE = 00020000 FLAGS = 000000

HOLES:

ADDRESS SIZE

00132300 00015600

ALLOCATION MAP:

SPACE ALLOCATED OCCUPIED BY

00130100-00131100 ...MCR (A) 00131100-00132300 ...MCR (P)

00132300-00150100

\*\*\*\* NO ENTRIES IN MRL \*\*\*\*

\*\*\*\*\* NO ENTRIES IN CTL \*\*\*\*\*

GEN

RSX-11D CORE-DUMP ANALYZER VØ6.2 5-0CT-76 14:25 PAGE 20 TPD SCAN TPD ENTRY ADDRESS = 145150V/105350R BASE ADDRESS # 00150100 SIZE \* 01627700 FLAGS = 031014 HOLES: ADDRESS SIZE 00236700 00003400 00244400 00004400 00323500 00011700 00436500 01341300 ALLUCATION MAP: SPACE ALLOCATED OCCUPIED BY 00150100-00164600 TT.... (A) SYSRES (L) 00164600-00203400 00203400-00210300 TTLIBA (L) 00210300-00225600 TTLIBB (L) 00225600-00231400 LP.... (A) MO... (A) 00231400-00235300 00235300-00236700 00236700-00242300 00242300-00244400 USRTSK (A) 00244400-00251000 00251000-00254700 DK... (A) 00254700-00261500 DT.... (A) 00261500-00270600 MM... (A) 00270600-00274300 DS... (A) DTAACP (F) 00274300-00310400 BPR ... (A) 00310400-00323500 00323500-00335400 F11ACP (A)
...LBR (A) 00335400-00365200 00365200-00421400 ...DEM (A) 00421400-00436500 00436500-02000000 \*\*\*\*\* NO ENTRIES IN MRL \*\*\*\* CTLI NODE ADR TASK PRI STS TI 50. 105034 ...LBR W1 TTIB 194769 106540 111440 102420 143656 000462 000001 000000 003652

•

RSX-11D CORE-DUMP ANALYZER VØ6.2 5-0CT-76 14:25 PAGE 22 STD SCAN

ALPHA TABLE STARTING ADDRESS = 145200V/105400R NO. OF ENTRIES = 69.

RSX-11D CORE- STD SCAN	DUMP ANALYZER	VØ6.2	5=0CT=76 141	125 PAGE 23
TASK PAR		DISK SIZE	AV LIM UTIL	
DS GEN	064200 248	. DB0 000035	1. 40. 4.	107242 107242
107220 107240	•	145150 064200 107242 000001	014770 000035 137716 000000	000035 024001 000000 000000
DTAACP GEN	040202 210	. DB0 000141	0, 255, 0,	103042 103042
103020 103040	016041 003310 000000 103042	145150 040202 103042 000001	014722 000072 147772 000000	000141 177490 000000 000000
DT GEN	064200 248	. DB0 000046	1, 40, 4,	103702 103702
103660 103700		145150 064200 103702 000001	014770	000046 024001 000000 000000
ERRLOG GEN	044200 200	. 080 000104	0. 40. 0.	104042 104042
104020 104040		145150 044200 104042 000001	014710 000104 123515 142140	000104 024000 000000 000000
F11ACP GEN	040202 220	. DB0 000276	1, 255, 0,	101142 101142
101120		145150 040202 101142 000002	014734 000210 052604 000000	000276 177401 000000 000000
F11MSG GEN	050000 500	. 080 000053	0. 40. 0.	104142 104142
194129 194149		145150 920000 104142 900001	014710 000045 045466 000000	000053 024000 000000 000000
LP GEN	064200 248	. DB0 000036	1. 40. 4.	104402 104402
104360 104400		145150 064200 104402 000001	014770 000036 141022 000000	000036 024001 000000 000000
MCRERR GEN	060300 75.	DB0 000006	0, 20, 0,	102342 102342
102320 102340		145150 060300 102342 000001	014513 000006 156004 102260	000000 000000 000000

RSX-11D CORE-DUMP ANALYZER VØ6.2 5-0CT-76 14:25 PAGE 31 PUD SCAN TTO ATTACH HANDLER EXP CNT DACP ACP TERM FLAGS RED CHAR 000210 TTO TT.... 0. 000200 000005 142300 052124 104000 000007 000000 000000 001000 ANAMAN 142300 142320 100260 000000 002756 142324 000060 000200 043140 000000 142346 000401 000200 000000 000000 000000 000005 000000 000000 888888 142360 I/O REQUEST NODES: PRI EFN FUNC CODE REG TYPE REQUESTOR ATL NODE DPB SZ LUN 8, 50. 0, USRTSK 110200 1. IO.RLB NORMAL 110720 105520 110200 004062 000001 001000 000000 060146 060146 000120 000000 000401 000400 000000 000000 000001 000001 000000 000000 000000 000401 143740 000000 110740 001336 110760 000000 CIØ FLAGS ATTACH RED HANDLER EXP CNT CHAR DACP ACP TERM 000000 0. 000000 699999 TTØ 142362 044503 000000 000000 000000 000000 000000 000000 142402 999999 000000 000000 000000 000000 000000 000000 000000 000000 000000 000000 142422 000000 000000 000000 000000 000000 142442 000000 \*\*\*\*\* I/O REQUEST QUEUE EMPTY \*\*\*\*\* COØ FLAGS ATTACH RED HANDLER EXP CNT CHAR DACP ACP TERM 000000 0. 000000 000000 TTO

000000 000000

000000 000000

000000

000000

142300

999999

047503 000000 000000 000000 000000

000000 000000 000000

142444

142464

RSX-11D CORE-DUMP ANALYZER V06.2 5-0CT-76 14:25 PAGE 32 PUD SCAN

\*\*\*\* I/O REQUEST QUEUE EMPTY \*\*\*\*

# CLØ

FLAGS	ATTACH	RED	HAN	DLER	EXP CNT	CHAR	DACP	AÇP	TERM
000000		LP0			0.	090000			000000
142526 142546 14256	0000	00 00	90000 90000 90000	000000 000000 000000		000000 000000 000000	000000 000000 000000	909990 909990 909000	143264 000000 000000
14260				000000	000000	0,000,00	200000		

\*\*\*\* I/O REQUEST QUEUE EMPTY \*\*\*\*

# BP0

FLAGS	ATTACH R	ED HAI	NDLER E	EXP CNT	CHAR	DACP	AÇP	TERM
990000	1	T4		0.	000200			000001
142616 142636 142656 142676	000000	000200 000000	004003 142634 000000	000000 142634 000000	000000 000000 000000	000001 000000	000000 000000 000000	144166 000000 000000

\*\*\*\* I/O REQUEST QUEUE EMPTY \*\*\*\*

# SPØ

FLAGS ATTACH RED HANDLER EXP CNT CHAR DACP ACP TERM 000000 000000 SYØ 0. 000200 142672 142712

•

RSX=11D CORE-DUMP ANALYZER PAGE 43 V06.2 5-0CT-76 14:25 MCR COMMAND BUFFERS NODE ADDRESS TASK ŢΙ LENGTH FIRST 40 CHARACTERS IN BUFFER ----111060 TT10 LBR CDAOLB, LP:/LI.CDACKQ ...LBR 24. RSX-11D CORE-DUMP ANALYZER VØ6.2 5-0cT-76 14:25 PAGE 44 BATCH COMMAND BUFFERS TASK LENGTH NODE ADDRESS TI FIRST 40 CHARACTERS IN BUFFER 10. VFY SYI/LO 111720 ...VFY BPO RSX-11D CORE-DUMP ANALYZER VØ6.2 PAGE 45 5-0CT-76 14:25 FTL SCAN DTAACP FTL ADDRESS = 107440V/047640R STD ADDRESS = 103020V/043220R LOAD ADDRESS = 00274300

PAR=GEN PRI=210, I/O IN PROG=0, I/O PEND=0, STATUS=LS TI=TT7

FLAGS WORD=002000 FLAGS=000000 000000 MASKS=144642 000401 000000 000000 
 146404
 146404
 102360
 143740
 000322
 000000
 000000
 000000
 000000
 000000

 145150
 107502
 107502
 000001
 147772
 002000
 107514
 107514
 107440 107460 107500 VØ6.2 RSX-11D CORE-DUMP ANALYZER 5-0CT-76 14125 PAGE 46 AST QUEUE SCAN ATL ADDRESS . 110200V/050400R TASK = USRTSK STD ADDRESS = 105520V/050422R TYPE . MARK TIME ENTRY POINT # 001000 ADDITIONAL INFORMATION: 000000 110448 110460 TYPE # I/O DONE ENTRY POINT # 001330 1038 ADR # 000000 EVENT FLAG . D. IOSB CONTENTS . 006401 000003 110340 110440 105520 000401 004062 000001 001000 000000 001330 000401 020041 000120 006401 000003 000000 000000 143740 000000 002000 002441 000120 001000 000000 000001 110500 110520 110540 TYPE . TERMINAL ENTRY POINT # 001330 NO ADDITIONAL INFORMATION 110340 110360

RSX-11D CORE-DUMP ANALYZER VØ6.2 5-OCT-76 14:25 PAGE 47 SRO SCAN RECEIVER SENDER TI PRI DATA BLOCK 050114 031000 040001 041104 000000 000172 000001 000000 004477 052140 000000 047014 000003 SPR... ... QUE TT2 100. RSX-11D CORE-DUMP ANALYZER VØ6.2 5-0CT-76 14:25 PAGE 48 CLOCK QUEUE SCAN TYPE \* TIME=SLICE DELTA \* 1. TICKS 104320 146400 142100 145600 000000 000001 000400 145150 000000 106102 106102 000000 153252 142140 000000 000000 106060 106100 TYPE = MARK TIME DELTA = 22. TICKS
TASK = ...DEM STD ADDRESS = 103360V/044524R ATL ADDRESS = 111360V/044526R EVENT FLAG = 2. AST ADDRESS = 000000 TYPE = MARK TIME DELTA = 8, TICKS
TASK = TT.... STD ADDRESS = 141740V/041044R
ATL ADDRESS = 100260V/041046R EVENT FLAG = 6. AST ADDRESS . 000000 100640 100660 TYPE = MARK TIME DELTA = 22595. TICKS
TASK = BPR... STD ADDRESS = 102100V/050604R TASK = BPR... STD ADDRE ATL ADDRESS = 111440V/050606R EVENT FLAG . D. AST ADDRESS = 002124 105320 100640 102100 111440 000000 054103 000000 000000 000000 000000 002124 000000 000000 000000 000000 000000 110420 TYPE = MARK TIME DELTA = 509673, TICKS
TASK = USRTSK STD ADDRESS = 105520V/046524R

ATL ADDRESS = 110200V/046526R EVENT FLAG = 1. AST ADDRESS = 001000

101460

RSX-11D CORE-	DUMP ANALYZER	v06•2	5-001-76 14:25	PAGE 5
MAC 00000	0 000130 GE	N 000002 NL	0./1. 5-0CT-76 [0,0]	
101340 101360	000000 00000 000400 00000			2412 3273
MCR 00131	1 000012 MC	R 000002 LS	1./1. 30-SEP-76 [1,1]	
102160 102200	600000 00000 000404 00040			7011 6053
MFT 00000	0 000023 MC	R 000002 NL	0./1. 30-SEP-76 [2,2]	
102220 102240	000000 00000 000400 00100			7011 6105
PIP 00000	0 00006 <b>3</b> GE	N 000002 NL	0./1. 28-SEP-76 [0,0]	
104620 104640	<ul><li>Ø0000</li><li>Ø0000</li><li>Ø0000</li></ul>			6011 6107
PRE 00000	000215 GE	N 000002 NL	0./1. 29-SEP-76 [1,1]	
104720 104740	000000 00000 002000 00040			6411 1676
QUE 00000	000050 GE	N 000005 NF	0./1. 28-SEP-76 (200,	200]
105160 105200	000000 00000 000400 10020			6011 7001

RSX-11D CORE-DUMP ANALYZER POOL VERIFICATION

V06.2

5-OCT-76 14:25

PAGE 51

CONSISTENCY CHECK OF POOL BITHAP AND SYSTEM LISTS: BITMAP INDICATES 347. NODES IN USE OUT OF 1112. LISTS INDICATE 307, NODES IN USE

> \*\*\*\*\*\*\*\*\* IMPORTANT NOTE \*\*\*\*\*\*\*\*\* THERE WILL ALWAYS BE SOME NODES
>  THAT ARE ALLOCATED IN THE BITMAP BUT
>  NOT FOUND IN THE LISTS, SINCE ANY
>  PRIVILEGED TASK CAN PICK NODES

NODES ALLOCATED IN BITMAP BUT NOT FOUND IN LISTS:

100140	100140	100140	100144	100144	100150	100150	100154	190154
100160	100160	100160	100164	100164	100170	100170	100174	100174
100200	100200	100200	000000	901270	000000	000000	000000	000000
100220	000000	000000	000000	000000	000000	000000	000000	000000
100240	900000	000000	000000	000000	000000	000000	000000	000000
107100	000000	000647	999999	000000	000000	000000	000000	000000
107120	000000	000000	000000	000000	000000	000000	000000	000000
107140	000000	000000	000000	000000	000000	000000	000000	000000
107320	000000	000000	000000	000000	000000	000000	999999	000000
107340	000000	000000	000000	000000	000000	000000	000000	000000
107560	000000	000000	000000	000000	000000	000000	909090	000000
107600	000000	000000	000000	000000	000000	000000	000000	000000
110040	110040	110040	110044	110044	000000	011022	000000	000000
110060	142000	031560	000000	000000	000000	000000	000000	000000
110100	000000	000000	000000	000000	000000	000000	000000	000000
110260	869146	110720	110120	104760	006062	016021	001030	026542

	RSX-11D CORE- POOL VERIFICA		LYZER	V06.	2	5=0CT=7	6 14:	25	PAGE	52
	110300	000000	000401	020046	000001	000000	000000	000000	000137	
	110320	143656	000000	002000	004142	000001	011022	000000	110334	
	110560	060146	060146	105520	110200	004062	000001	001000	000000	
	110600	000000	000401	020036	900120	600000	000000	000000	000000	
	110620	143740	888888	002000	002441	000120	000000	000000	999998	
	119640	110720	060146	105520	110200	004062	000001	001000	000000	
	110660	000000	000401	001336	000120	888888	000000	000000	000000	
	110700	143740	000000	000000	122164	000001	901090	000000	000001	
	111300	110720	060146	102420	111000	005746	000401	004600	000370	
	111320	000000	000401	020074	000120	000037	021112	000006	000000	
	111340	144332	000000	902090	000522	000120	011022	000000	145170	
	111560	000000	000000	000000	900000	000000	000000	000000	000000	
	111600	000000	000000	000000	P00000	000000	000000	909099	000000	
	111660	000000	000000	000000	000000	000000	000000	000000	000000	
	111700	000000	000000	000000	000000	000000	000000	000000	000000	
	112060	112220	027420	102420	000000	144332	000000	001000	000000	
	112100	111000	000000	000171	130100	177400	005422	000443	000000	
	112120	144642	000000	004000	000016	000000	011022	000000	000000	
	112140	142000	027500	103360	000401	000000	000401	000410	000000	
	112160	000000	000401	020221	000000	000001	000131	000000	969999	
	112200	143574	000000	002400	004227	000131	000044	000000	112214	
	146020	000000	000000	004767	166244	012600	000207	004767	165316	
	146940	005002	012700	007552	004767	165524	016702	156022	004767	
	146969	177666	005046	016705	156006	072527	000006	012701	001340	

NODES FOUND IN LISTS BUT NOT ALLOCATED IN BITMAP:



RSX-11D CORE-DUMP ANALYZER NODE POOL DUMP

VØ6.2

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PAGE 54

SYSTEM NODE POOL

. . NODE ALLOCATED

	VIRT	REAL		CONT	ENTS		ASCII		RA	50	
			_	_							
•	100000	040200	145642	145642	142100	000003	"K"KPD	2WJ	2WJ	1 P	С
*	100010	040210	003212	000000	000000	000000		AA4			
	100020	040220	000000	000000	000000	000000					
*	100030	040230	000000	000000	000000	000000					
	100040	040248	100100	041540	142100	013746	• C•D	TT2	JIH	1 P	C28
*	100050	040250	172346	042737	004000	100104	_E D	9GN	KGG	AKH	TT6
•	100060	040260	012737	001205	172346	000137	<u>.</u> +	C S 9	PE	9GN	80
*	100070	040270	064102	000000	000000	000000	В	PSJ			
*	100100	040300	000000	000000	000000	000000					
•	100110	040310	000000	000000	000000	000000					
٠	100120	040320	200000	000000	000000	000000					
*	100130	040330	000000	000000	000000	000000					
٠	100140	040340	100140	100140	100144	100144				TU.	TU.
*	100150	040350	100150	100150	100154	100154		TU2	TÜ2		TU6
	100160	040360	100160	100160	100164	100164		TV	TV	TVD	TVD
*	100170	040370	100170	100170	100174	100174			TVH	TVL	TVL
*	100200	040400	100200	100200	000000	001270	. 8	TVP	TVP		QP
	100210	040410	000000	000000	000000	000000					
*	100220	040420	000000	000000	000000	000000					
٠	100230	040430	000000	000000	000000	000000					
	100240	040440	000000	000000	000000	000000					
*	100250	040450	000000	000000	000000	000000	V 6711	*00	040	4 14 4	050
٠	100260	040460	101040	145600	141740	144332	K CZH	122	2VP	1 MX	2ER
	100270	040470	000370	000000	000001	001501	. A	FH	4 44 54	A	Ţ3
	100300	040500	000016	141740	000000	040000	C •	N	1 M X		JIX
•	100310	040510	007377 145150	000000	000000	000000	10 D	809 20P	***	TXR	В
*	100320	040520 040530	027733	100322 000000	100322 100334	000002	JR R 【/ \ \	GZK	TXR	TX.	TX.
	100340	040540	100400	000240	141740	100334 013746	' 'c'	TYX	D	1 MX	C28
-	100350	040550	172346	042737	010000	100100	E •	9GN	KGG	BVP	112
	100360	040560	012737	002103	172346	000137	ž	CS9	SK		B0
	100370	040570	060174	004040	000526	000000	* L ^ *	Out	AL	HV	bu
	100400	040600	100440	000200	141740	013746	Č	TZP		1 M X	C28
÷	100410	040610	172346	042737	020000	100100	ھِ ق	9GN	KĞĞ	ED2	TT2
-	100420	040620	012737	002103	172346	000137		CS9		9GN	BO
	100430	040630	060446	000000	001000	000000	į ·	OÜØ	•	L2	
÷	100440	040640	100500	000140	141740	013746	ē c	TSH	BP	1 MX	C28
-	100450	040650	172346	042737	000400	100106	_E F	9GN	KGG	FP	118
-	100460	040660	012737	002103	172346	000261	- 1	CS9		9 G N	DQ
	100470	040670	000137	060174	000000	000000	•	ВО	OUT	• • • •	-
	100500	040700	100540	000100	141740	013746	* • c	T.	ÁX	1 MX	C28
*	100510	040710	172346	042737	001000	100106	EF	9GN	KGG	ΓŜ	TTB
	100520	040720	012737	002103	172346	000261	ē 1	C89	SK	9GN	DQ
•	100530	040730	000137	066446	000526	000001	1 8 V	80	000	Н۷	Ā
	100540	040740	100600	000040	141740	013746	Ť C	1.2	5	1 MX	c 28
	100550	040750	172346	042737	000001	100106	LE F	9GN	KGG	Â	118
	100560	040760	012737	002103	172346	000137	~	CS9	ŠĶ		80
-	100000	546,05	20 m/ U/	~ 5 5 5 5 5	., 20-0	~0010/		<b></b>	₩:1	~ · · · ·	

NODE POOL D		ANALYZER	VØ	6.2	5=0CT	<b>-</b> 76⁻	141	25		P	AGE: 55
+ 100570	40770	060302	000000	001000	000000	В		OSJ		L2	
* 100600 6	41000	100700	000040	141740	013746	•	C	TOP	2	1MX	C28
	41010	172346	042737	000002	100106	_E	F	9GN	KGG	В	<b>718</b>
	41020	012737	002103	172346	000137	∡ Ċ	•	CS9	SK	9GN	BO
	41030	060502	011022	000000	000000	В		OVR	B5Z		
	41040	110400	104320	141740	100260	P	CØ	.MH.		1MX	
	41050	000000	000010	000000	000040	_			H		2
	141060	100304	999996	000000	000000	D		TXD	F		
	141070	000000	000000	000000	000000		_	Ŧ14	• • • •	. 4	608
	41100	100740	040700 042737	141740	013746	PA	C F	T1H	JT2 KĞG	1MX	
	941110 941120	172346 012737	002103	000004 172346	100106	ž.E.	1	CS9	ŠK	D CN	TTB
	41130	000137	000302	000000	000261 000000	_ C	1	BO	oSJ	9GN	) W
	341140	101000	040640	141740	013746	* B	С	12	JŤ	1MX	C28
	41150	172346	042737	000010	100106	μÊ	F	9GN	KĜG	H	118
	41160	012737	002103	172346	000261	•	ì	CS9	SK	9GN	DG
	41170	000137	060502	000000	000000	ΪĎ	•	80	DVR	• • • • • • • • • • • • • • • • • • • •	•
	41200	100322	100322	141740	000003	RR	C	TXR	_	1MX	C
+ 101010 6	41210	002616	000000	000000	000000		•	bν		•	-
	41220	000000	000000	000000	000000						
· 101030 0	41230	000000	000000	000000	000000						
	41240	106460	100260	102360	142300	00	●D	VWX	TWX	UJ2	1 SH
	41250	000370	000000	000000	002256			FH			/8
	41260	000016	104360	969999	000000			N	U6P		
	41270	000003	006401	000000	000000			Ç	BČI		
	941300	145150	101102	101102	000001	JB 8	3	-	T3Z	T3Z	· 🛕
	41310	141022	000000	101114	101114	BL	_	188			T36
	141320	025167	003310	145150	040202	*H	J ·	F11	ACP	20P	JL4
	141330	014734	000210	000276	177401	, ,	•	DET	CP	00	- 5 ¥
	141340	000000	101142	101142	000002				TAR	T4R	В
	41350	052604	000000	000000	000000	U		MSL		4.5	6 O LI
	41360	101500 001772	000020 111220	142100	142300		D.D	YR	. WR	IP 2JR	1SH BVP
	941370 941400	131053	077514	144642	010000 040200	+2L	_	-TK		20P	JL2
	41410	014771	000031	000031	024000	425	J •	DFI	' 'Y	2 UP	FP
	41420	000000	101222	101222	000001		•	DI I	-	T52	<b>A</b>
	041430	156260	000000	000000	000000	0\		SLP		,	_
	41440	131574	073376	145150	040200	3	J •	• • •	SAV	20P	JL2
	41450	014735	000151	000151	024000	1	•	DEU	БУ	BY	FP
	41460	000000	101262	101262	000001	2 2				TOR	A
	41470	136775	000000	000000	000000			907	•		
	41500	113000	000000	145150	040200		J ·	X		20P	JL2
· 101310 6	41510	014734	000210	000276	177400	\	• •	DET	ÇP	DØ	2
+ 101320	141520	000000	101322	101322	000002	RR	₹ .		・ナブゴ	<b>T7J</b>	8
+ 101330 R	41530	052604	000000	000000	000000	U		MSL			
+ 101340 0	141540	000000	000000	000035	000000					1	
	41550	000000	000130	000114	002412	ΧL			ВН	A6	_2J
	41560	000400	000000	145150	014402		J	FP		20P	D B
	141570	000400	000000	000000	103273		,	FP	La A =		UVK
	41600	131574	050553	145150	020100	3 0	-	DAJ	MAC	_	
	41610	014462	000131	001166	024000	2 Y	(	DAJ	81	_	FP
	141620	000000	101422	101422	000000			1111.0		T84	
• • • • • •	41630	103257	101340	142140	000000	/	D	009	TŽX	1P2	
+ 101440 0	141640	000000	000000	000035	000000					- /	

RSX-11D CORE-DUMP ANALYZER V06.2 5-0CT-76 14:25 PAGE 100 TASK DUMP

TASK ...VFY

•	969999	000002	000000	000000	000000	8
6	99919	000000	000000	000000	999999	
ė	300020	000000	000000	000000	000000	
6	00030	000000	000000	000000	000000	
6	00040	000000	999999	000000	000000	
•	300050	000000	000000	000000	000000	
6	999999	000000	000000	000000	000000	
6	90076	000000	000000	000000	000000	
6	90100	000000	000000	000000	000000	
	000110	000000	000000	000000	888888	
(	00120	000000	000000	000000	900000	
	100130	000000	000000	000000	000000	
•	00140	000000	000000	900000	000000	
	00150	000000	000000	000000	000000	
6	00160	000000	000000	000000	000000	
•	00170	000000	000000	000000	900000	
	90200	000000	000000	000000	000000	
(	300210	000000	000000	000000	000000	
	800220	000000	000000	000000	000000	
	966536	000000	000000	000000	000000	
	000240	000000	000000	808888	000000	
6	100250	000000	000000	000000	000000	
	860260	000000	000000	000000	000000	
_	00270	000000	000000	000000	000000	
	000300	000000	000000	000000	000000	
	300310	000000	900000	000000	88888	
	00320	000000	000000	000000	808088	
	100330	000000	000000	000000	000000	
_	300340	999999	000000	000000	000000	
	00350	000000	000000	000000	000000	
	000360	000000	000000	000000	000000	
-	000370	000000	000000	000000	000000	
-	000400	000000	000000	000000	000000	
	000410	000000	000000	999999	900000	
_	808428	99 <b>99</b> 99	000000	000000	000000	
	00430	000000	000000	000000	00000	
-	800440	900000	000000	000000	909999	
	00450	000000	000000	000000	900000	
	800460	000000	000000	000000	000000	
	00470	000000	000000	000000	909099	
	00500	000000	000000	000000	900000	
	000510	000000	000000	000000	00000	
	800520	000000	000000	999999	00000	
-	000530	000000	000000	000000	00000	
	000540	000000	000000	000000	00000	
	00550	000000	000000	000000	000000	
	000560	000000	000000	000000	000000	
	000570	000000	000000	000000	00000	
	00600	000000	000000	000000	000000	
(	000610	888888	000000	000000	00000	

SX-1 ASK	DUMP	CORE-DUMP P	ANAL	YZER	VØ6.	2	5-00	1-7	6	14:25	5		PAGE	101
_		20 000		22222										
	0062		1000	000000	000000	900000								
	10063		000	000000	000000	000000								
	0064		000	000000	000000	000000								
	0066		1000	000000	000000	000000								
	0067		1000 1000	000000 000000	000000	000000								
	1007			000000	000000 000000	999999 999999								
	007		000	000000	969999	000000								
	0072	•	000	000000	000000	000000								
	0073		000	000000	000000	000000								
	007		000	000000	900000	000000								
	007		1000	000000	000000	000000								
-	007		000	000000	000000	000000								
-	007		800	000000	000000	000455			_				GU	
	0100		746	000455	104377	000000			_	CTF	GU	U61	•	
	0101		1000	000000	000000	999999		-		<b>U</b> 11	30	00.		
-	010		000	000000	000000	000000								
	010		000	000000	000000	000000								
	010		000	000000	000000	000000								
	010		1000	000000	999999	000000								
-	10100		000	000000	000000	999999								
-	010		000	600000	000000	000000								
	0110		444	000034	000000	000000	\$			26				
Q	011	10 077	406	000016	077406	077406	_			TL8	Ñ	TL8	TL8	
Q	011		406	002372	000000	000000				TL8	14			
6	10113	30 004	365	000402	000602	001002				AQM	FŘ	ΙZ	L4	
	1011		600	000003	000000	000000				5H	C			
Q	10115	50 000	000	000402	000602	001002					FŘ	12	L4	
6	0110	60 177	600	174000	000656	111000				5H	95A	JØ	WNP	
Q	10117		000	000064	000000	001277		4	?		AL		G M	
Q	10120	00 060	104	000100	174017	000100	D	•	•	OPD	AX	9\$W	Ax -	
	10121		106	000000	000000	000000	•			AX				
	0122		000	000000	001002	000401						L4	FQ	
	10123		003	000004	000000	000026				C	D		٧	
	10124		268	000000	000000	060146	0 \	,		5LP			OP8	
	012		146	000000	000000	000006				OP8			F	
	012		444	66666	000001	000000	\$ E			10.		A		
	0127		006	900000	145006	000000	J	٠ '	J	2L8		2L8		
	0130		001	000000	142526	900000		٧	£	A		148		
	013		000	000000	000000	000000								
	013		999	000000	000000	999999								
	013		000	000000	000000	000000								
	0134		1000	000000	000000	000000								
	013		000	000000	000000	000000								
	013		1000	000000	000000	000000								
	013		000	000000	000000	000000								
6	0140	00 000	1001	000000	000000	000000				A				

RSX-11D CORE-DUMP ANALYZER V06.2 5-OCT-76 14:25 PAGE 102 8GA DUMP

SGA = SYSRES

DUMP	DUMP ANA	LYZER	VØ6.	2	5-0CT-76	14:25		PAGE 1
000620	000017	001416	122760	000001	PX	0 \$9	ZV	A
000630	000000	001407	004767	005656	W	80		A40
000640	103007	012770	177777	900030	· ·	UQ1 CTX		X
000650	000403	016002	000030	160102	^ <del>**</del> B#	FS DSJ		55Ĵ
999669	004767	015012	016001	000070	W 8	AW1 DF2		AP
889670	004767	011532	116002	000052	₩Z +	ANI CC4		ÂÂ
000700	005005	156005	000055	000412	\•	AXE SHE		FZ
000710	004767	011512	103007	005702	W J B	AW1 CCR	_	A5J
200720	002405	132760	000002	000056	P5 .	2E /0F		AF
000730	001401	111102	016101	000012	B A	SI WPE		j
000740	077515	110260	000052	116046	M_0 + 8	TN/ WF	AB	X9Ñ
000750	000052	004767	007252	005025	* W *	AB AWI	-	AXU
000760	010604	005046	132760	000010	8 P5	BZD AXÊ		Н
000770	000043	001426	132760	000041	# P51	5 30		3
001000	000017	001022	132760	000020	P5	O MJ		مَ
001010	000051	001405	010046	962716	) & NE	AA SM		PKV
001020	000120	012746	001007	013702	PFB	B CTF		C2B
001030	000002	062702	000010	005712	BE J	B PKJ		Ă5R
001040	001403	010246	012746	001001	& F	SK BZV		L3
001050	132760	000001	000043	001003	P5 #	/DP A	_	ĹŠ
001060	010046	012746	007004	010625	8 F	BWN CTF		82U
001070	010446	012784	010000	005760	8 D P	B/Ø CSL		AGP
001100	000076	100006	062705	000004	> EE	AV TSN		D
001110	016015	000076	042715	100000	> ME			TSH
001120	004767	011262	012606	112601	W 2	AN1 B9Z		W63
001130	003002	110160	.000052	132760	P + P5	BR WDP	**	/DP
001140	000200	000051	001516	162706	) N FE	CH AA		60 V
001150	000040	010602	010246	012746	8 F		BZV	
001160	001077	104377	016203	000014	?	NO U61		L
001170	062706	000040	012746	131574	FE F \3		CTF	
001200	012746	074522	010601	010002	FRY	CTF SPR		
001210	062702	000102	016246	000016	BEB &	PKJ AZ		N
001220	016246	000014	016246	000012	8 8	DWN L		j
001230	016246	000010	016246	000006	8 6	DWN H	-	F
001240	016246	000030	016246	000026	8 8	DWN X		Ÿ
001250	016246	000024	016246	000034	ii	DWN T		·
001260	016246	000032	012746	000001	8 F	DWN Z		· Å
001270	116046	000074	110366	000001	8 < V	X9N AT		Δ
001300	016046	000072	010602	005046	8 1 8	DIF AR		AX8
001310	010246	005046	005046	005046	8 8 8 8	BZV AXE		BXA
001320	005046	016146	000002	R11146	8 F F	AX8 DUE		BZÖ
001330	012746	004511	104377	9627 <b>9</b> 6	E 1 EE	CTF ASG		PKN
001340	000036	116001	000042	016002	, ,	Ø X80		DSJ
001350	000072	116003	000074	P10346	: < F	AR X85		B.F
001360	010246	010146	012746	002007	8 F F	BZV BX8		ΥĪ
001370	104377	005060	000072	105060	. 0 : 0	U61 AYH	= .	VDP
001400	000074	004767	011142	000167	< W B W	AT AWI		89
001410	006226	004567	004376	004767	WAW	B V ATH		AW1
001420	007452	103443	132760	000001	* * P5	BOB UXS		Ä
001430	000017	001441	132760	000010	1 P5	O TA		H
001440	000051	001404	112760	177766	) P.V.	AA SL		8 F
001450	000052	000427	004767	011360	WP	AB F9		CAH
001460	016001	000070	016103	000002	" a ĉ ĺ		DTS	8

#### APPENDIX B

#### TKB15.CMD FILE

# See Section 6.3.1

[11,15]EXEC/LB:DMIAR

```
[11,15] T K B 15. C M D
   INDIRECT TKB COMMAND FILE FOR THE RSX-11D EXEC
[11,17]EXEC/-HD/-FP,[111,15]EXEC/CR/-SP,[1,1]EXEC.STB=
; IF A MAP IS NOT REQUIRED PUT A ';' AT THE BEGINNING OF THE ABOVE LINE
; AND REMOVE THE ";" FROM THE BEGINNING OF THE FOLLOWING LINE
;[11,17]EXEC/-HD/-FP,,[1,1]EXEC.STB=
[11,15]EXEC/LB:EM00
[11,15]EXEC/LB:EM01
[11,15]EXEC/LB:EM02
[11,15]EXEC/LB:EM03
[11,15]EXEC/LB:EM04
[11,15] EXEC/LB: EM05
[11,15]EXEC/LB:EM06
[11,15]EXEC/LB:EM07
[11,15] EXEC/LB: EM08
[11,15]EXEC/LB:EM09
[11,15]EXFC/LB:EM10
[11,15]EXEC/LB:ACCTET
[11,15]EXEC/LB:DMGCL
[11,15]EXEC/LB:DM010
[11,15]EXEC/LB:DMGLI
[11,15] EXEC/LB: DMASS
[11,15]EXEC/LB:DMALP
[11,15]EXEC/LB:DMREQ
[11,15]EXEC/LB:DMSCH
[11,15] EXEC/LB: DMMKT
[11,15]EXEC/LB:DMCSR
[11,15]EXEC/LB:DMCMT
[11,15]EXEC/LB:DMSED
[11,15]EXEC/LB:DMSUS
[11,15]EXEC/LB:DMEXT
[11,15]EXEC/LB:DMGTP
[11,15]EXEC/LB:DMGPP
[11,15]EXEC/LB:DMGMP
[11,15]EXEC/LB:DMGCP
[11,15]EXEC/LB:DMSAR
[11,15]EXEC/LB:DMABO
[11,15]EXEC/LB:DMFIX
[11,15]EXEC/LB:DMDST
[11,15]EXEC/LB:DMDCP
```

# TKB15.CMD FILE

```
[11,15]EXEC/LB:DMSDV
[11,15]EXEC/LB:DMAST
[11,15]EXEC/LB:DMATX
[11,15]EXEC/LB:DMGSS
[11,15]EXEC/LB:TSCH
[11.15]EXEC/LB:NTRACE
; THE EXEC IS NORMALLY BUILT WITHOUT ODT. TO INCLUDE ODT IN THE EXEC
; CHANGE THE NEXT SPECIFICATION TO BE: [11,15]EXEC/LB;ODT
; THE SYSTEM IS RESUMED FROM ODT BY KEYING "CONTROL C"
; THE SYSTEM IS RESUMED FROM A BREAKPOINT BY KEYING "P"
[11,15]EXEC/LB:NODT
; THE EXEC IS BUILT INCLUDING A CRASH MODULE WHICH CAN OPTIONALLY WRITE
; MEMORY ONTO A DUMP MEDIUM AND/OR DISPLAY A CRASH MESSAGE. IF MESSAGES
; ARE SELECTED ONE OF THE FOLLOWING WILL BE PRINTED ON THE SYSTEM
; CONSOLE IF THE SYSTEM CRASHES.
         ODD ADDRESS OR TRAP 4
         UNDEFINED SYSTEM EVENT
         AST NODE PICK FAILURE
         FLOATING POINT ERROR
        MEMORY PARITY ERROR
         NON-RSX EMT EXECUTION
        RESERVED INSTRUCTION
         SST IN KERNEL MODE
        RED OR YELLOW STACK
         EMT 377 IN KERNEL MODE
         UNDEFINED SYSTEM ERROR
 THE OPTIONS AVAILABLE WITH EACH CRASH MODULE ARE SHOWN BELOW
        MODULE
                          MESSAGES
                                            MEMORY DUMPING
:
                                            NO
        NCRASH
                          NO
                                            NΩ
         MCRASH
                           YES
                                            YES -- TC11 DECTAPE **
YES -- TC11 DECTAPE **
         TC11CR
                           NO
                          YES
         TC11CM
                                            YES -- TU10 MAGTAPE
YES -- TU10 MAGTAPE
         TU10CR
                           NO
         TU10CM
                           YES
;
                                            YES -- TU16 MAGTAPE
         TU16CR
                           NΩ
                                            YES -- TU16 MAGTAPE
YES -- RK05 DISK ++
YES -- RK05 DISK ++
         TU16CM
                           YES
         RK05CR
                           NO
         RK05CM
                           YES
         TC11 DECTAPE CAN NOT DUMP MORE THAN 124K WORDS
         RKO5 DISK CAN NOT DUMP MORE THAN 1152K WORDS
; THE DEFAULT UNIT NUMBERS FOR THE RESPECTIVE DEVICES ARE:
         TC11 -- #7
         TU10 -- #7
         TU16 -- #0
; THE DEFAULT MAY BE CHANGED BY USING THE TKB GLOBAL PATCH OPTION
; TO MODIFY THE CRASH UNIT NUMBER (.CRSUN) AS DESIRED. FOR EXAMPLE, ; TO CHANGE THE UNIT NUMBER TO 2, INCLUDE THE FOLLOWING LINE:
         GBLPAT=EXEC:.CRSUN:2
[11,15]EXEC/LB:NCRASH
[11,15]EXEC/LB:LVEND
```

[11,15] EXEC/LB:SCOMM

## TKB15.CMD FILE

```
STACK=0
UNITS=0
EXTSCT=:130000
;GBLPAT=EXEC:.CRSUN:0

** NOTE **

** WHENEVER THE EXECUTIVE IS REBUILT USING THE COMMAND FILE [11,15]TKB15.CMD
; THE BOOTSTRAPS MUST ALSO BE REBUILT USING [11,17]B00TSBLD.CMD BEFORE
; PERFORMING A SYSTEM GENERATION TO INCLUDE THE NEW EXECUTIVE IN THE STSTEM.
```

# APPENDIX C

# ERROR LOGGING SAMPLE OUTPUT AND DATA RECORD FORMATS

This appendix contains sample output from error logging and the format of records in the ERROR.SYS file produced by the preanalyzer (PRE).

# C.1 SAMPLE ERROR LOGGING OUTPUT

PAGE 15

RSX 11 SYSTEM ERROR REPORT COMPILED AT 20-AUG-75 13:48:53

**************************************	NUMBER 6.
TAPE PARAMETERS  UNIT NAME  VOLUME LABEL  VOLUME OWNER UIC  DEVICE TYPE  DT1  1,1  TU56 UNIT-1 CON	TROLLER-0
TAPE REGISTERS AT ERROR TIME  TCST 04034 SELECTION ERROR¹ TCCM 104702 TCWC 177400 TCBA 113534 TCDT 000007	
ERROR DIAGNOSIS RECOVERED SELECTION ERROR <sup>1</sup>	
RETRIES PERFORMED 0.	
USER TASK PARAMETERS  TASK NAME TASK UIC TASK UIC PHYSICAL START ADDRESS 102200 USER FUNCTION REQUESTED WRITE FUNCTION INTERPRETED FROM REGISTERS PHYSICAL BUFFER ADDRESS START TRANSFER SIZE IN BYTES LOGICAL LOCK NUMBER AT IO GO COUNT OF IO IN PROGRESS COUNT OF IO PENDING	(400) READ BLOCK NUMBER REVERSE 113534 1000 2 1.
NUMBER OF FUNCTIONS ISSUED TO THIS UNIT	66
VECTORS WITH ACTIVE IO	

PRE ANALYZER RUN

Figure C-1 Tape Error Report

\*

AT 12-AUG-75 10:21:01

<sup>&</sup>lt;sup>1</sup>Included only if /DE switch is specified. <sup>2</sup>System information.

# ERROR LOGGING SAMPLE OUTPUT AND DATA RECORD FORMATS

RSX 11 SYSTEM ERROR REPORT COMPILED AT	20-AUG-75 13:48:04 PAGE 9.
**************************************	ERROR NUMBER 13.
MEMORY REGISTERS AT ERROR TIME PARITY CSR ADDRESS PARITY CSR CONTENTS PROCESSOR STATUS WORD PRIOR TO THIS ER PHYSICAL START ADDRESS OF ERRING TASK PROGRAM COUNTER PRIOR TO THIS ERROR	172124 000000 ROR 174000 171500 254
ERROR DIAGNOSIS 11/45 MEMORY PARITY ERROR AT LOCATION 0XX WITH MEMORY MANAGEMENT REGISTER 0 ACTION USER MODE	VE
**************************************	
**************************************	
1System information.	

System information.

Figure C-2 Memory Parity Error (No /DE Switch)

RSX 11 SYSTEM ERROR REPORT COMPILED AT 20-AUG-75 13:48:04

PAGE 9.

PARITY CSR ADDRESS 172124
PARITY CSR CONTENTS 000000

PROCESSOR STATUS WORD PRIOR TO THIS ERROR 174000
PHYSICAL START ADDRESS OF ERRING TASK 171500
PROGRAM COUNTER PRIOR TO THIS ERROR 254

ERROR DIAGNOSIS
11/45 MEMORY PARITY ERROR
AT LOCATION ØXX
WITH MEMORY MANAGEMENT REGISTER Ø ACTIVE
IN USER MODE

Figure C-3 Memory Parity Error (/DE Switch Specified)

<sup>&</sup>lt;sup>1</sup>Included only if /DE switch is specified.

# C.2 FORMAT OF RECORDS IN ERROR.SYS FILE

HEADER FIELD (HF)
ALL RECORDS CONTAIN A HEADER FIELD. WORDS Ø-7 OF THE HEADER
FIELD ARE STANDARD FOR ALL HEADER FIELDS.

ERROR HEADER FIELD

WD	HIGH BYTE LOW BYTE						
Ø	SIZE OF HEADER FIELD						
1	FORMAT = Ø ENTRY TYPE (See 1)						
2	YEAR						
3	MONTH						
4	DAY						
5	HOUR						
6	MIN						
7	SEC						
8-N	VARIABLE ACCORDING TO ENTRY TYPE						

THE FOLLOWING IS A DESCRIPTION OF WORD 8-N OF THE HEADER FIELDS WHICH ARE PRESENTLY USED.

ENTRY TYPE \$4\$ ERRLOG STARTUP, TYPE 41 ERRLOG SHUTDOWN, TYPE 43 SOURCE FILE RE-INITIALIZE

8	NEXT UNIVERSAL ERROR SEQUENCE NUMBER
9	POINTER TO INITIALIZE DESCRIPTOR FIELD (IDF)
	IF ZERO FIELD NOT PRESENT. CURRENTLY = Ø
10	POINTER TO INITIALIZE COMMENTS FIELD (ICF)
	IF ZERO FIELD NOT PRESENT. CURRENTLY = Ø
11	POINTER TO SYSTEM HARDWARE HISTORY FIELD (SHF)
<u>l</u>	IF ZERO FIELD NOT PRESENT. CURRENTLY = Ø

ENTRY TYPE 1, DEVICE ERROR BIT SET

8	UNIVERSAL ERROR SEQUENCE NUMBER
9	DEVICE TYPE (see 3) DEVICE CLASS (see 2)
10	PHYSICAL UNIT NUMBER (THIS BYTE NOT USED)
11	TWO BYTE ASCII OF THIS DEVICE MNEMONIC
	WITHIN THE SYSTEM
12	UNIT NUMBER AS SEEN BY THE SYSTEM
13	POINTER TO DEVICE REGISTER FIELD
14	POINTER TO PROGRAM FIELD
15	POINTER TO VOLUME FIELD
16	POINTER TO STATISTICS FIELD
17	POINTER TO BUS ACTIVITY FIELD

ENTRY TYPE 2, MEMORY PARITY ERROR

8	UNIVERSAL ERROR SEQUENCE NUMBER
9	DEVICE TYPE (see 3) DEVICE CLASS (see 2)
10	THIS WORD NOT USED
11	PROCESSOR STATUS WORD BEFORE PARITY ERROR
12	PROGRAM COUNTER AT TIME OF PARITY ERROR
13	1/64 PHYSICAL BASE ADDRESS OF TASK WHICH HAD
	PARITY ERROR
14	POINTER TO DEVICE REGISTER FIELD

#### ERROR LOGGING SAMPLE OUTPUT AND DATA RECORD FORMATS

DEVICE REGISTER FIELD (DRF)
THIS FIELD IS A COPY OF THE DEVICE REGISTERS WHICH REPRESENT
THE STATE OF THE DEVICE AND CONTROLLER AT INTERRUPT TIME. THIS
FIELD IS PRESENT WITH MOST DEVICE RELATED ERRORS.

DEVICE REGISTER FIELD

_ <b>W</b> D	HIGH BYTE	LOW BYTE
Ø		SIZE OF DEVICE REGISTER FIELD
1		$FORMAT = \emptyset$
2-N	DEVICE AND CON'	FROLLER REGISTERS AS SEEN ON
	THE UNIBUS	

PROGRAM FIELD (PGF)

THIS PROGRAM FIELD IS USED FOR DEVICE ERROR ENTRIES.

THIS FIELD IS PRESENT WITH ERRORS WHICH CAN BE TRACED TO A PROGRAM

PROGRAM FIELD

WD	HIGH BYTE LOW BYTE
Ø	SIZE OF PROGRAM FIELD
1	2 = 18 BIT ADDRESS SY FORMAT = 1
	STEM
1	4 = 22 BIT ADDRESS SY
	STEM
2-3	TASK NAME (6 CHARACTER RAD5Ø)
4	TASK uic
	PROGRAMMER GROUP CODE PROGRAMMER CODE
5	1/64 LOAD IMAGE BASE ADDRESS
6	USER FUNCTION CODE
7-12	6 USER FUNCTION PARAMETER WORDS (DEVICE SPECIFIC)
13	MAXIMUM # OF RETRIES # OF RETRIES REMAINING
Ĺ	TO BE DONE
14	COUNT OF THIS TASK'S I/O IN PROGRESS
15	COUNT OF THIS TASK'S I/O WAITING IN A QUEUE

VOLUME FIELD (VLF)

THIS FIELD IS USED TO SOFTWARE IDENTIFY THE PHYSICAL MEDIA WHICH IS MOUNTED ON A DEVICE. WHICH PACK OF MULTIPLE PACKS WHICH WAS ON THE DEVICE WHICH ERRED.

VOLUME FIELD

WD	HIGH BYTE	LOW BYTE	
Ø	SIZI	E OF VOLUME FIELD	
1	(NOT USED)	$FORMAT = \emptyset$	
2-7	12 BYTH	E ASCII VOLUME NAME	
8	VOLUME OWNER uic	VOLUME OWNER uic	
	GROUP NUMBER	PROGRAMMER NUMBER	ı

STATISTICS FIELD (STF)

THIS FIELD REPRESENTS THE RATE OF ERRORS ON A DEVICE OR THE USAGE OF A DEVICE. THIS FIELD IS PRESENT FOR DEVICE ERRORS

_	₩D	HIGH B	YTE		]	LOW E	SYTE				-	
STATISTICS	Ø			SI	ZE							
FIELD	1	(NOT USED	)		FORM	$= T\ell$	2					
	2-3	NUMBER OF	TIMES	THE G	OBIT	WAS	SET	FOR	THIS	UNIT		

### ERROR LOGGING SAMPLE OUTPUT AND DATA RECORD FORMATS

BUS ACTIVITY FIELD (BAF) THIS FIELD REPRESENTS THE ACTIVITY WHICH WAS CONCURRENT ON THE UNIBUS AT THE TIME OF ERROR. THIS FIELD IS PRESENT TO ASSIST IN THE DETECTION OF BUS INTERACTION PROBLEMS.

HIGH BYTE

LOW BYTE

BUS	
ACTIVITY	
FIELD	

-	WD	
Ī	Ø	SIZE
Ī	1	(NOT USED) FORMAT = 1
	2-9	A BIT SET IN EACH WORD REPRESENTS ACTIVITY ON THAT VECTOR I.E., WORD 1 BIT Ø = VECTOR
l		100.

#### NOTE 1

ENTRY TYPES ARE DIVIDED INTO MAJOR CATEGORIES. THESE MAJOR CATEGORIES CORRESPOND TO A DIVISION WHICH DETECTED THE ERROR.

THE FOLLOWING MAJOR CATEGORIES ARE DEFINED WITH THE RANGE OF ENTRY CODE NUMBERS.

UNDEFINED

ØØ1 - Ø37 HARDWARE ERRORS

Ø4Ø - Ø77 SYSTEM

THE FOLLOWING EXPLICIT ENTRY TYPE CODES ARE DEFINED

HARDWARE ERRORS

ØØ1

DEVICE ERROR BIT SET

ØØ2 MEMORY PARITY ERRORS

SYSTEM

ø4ø Ø41

ERRLOG START-UP

Ø43

ERRLOG SHUT-DOWN SOURCE FILE RE-INITIALIZATION

NOTE 2

DEVICE CLASSES ARE DIVIDED INTO MAJOR CATEGORIES ACCORDING TO THEIR FUNCTION WITHIN THE OPERATING SYSTEM.

THE FOLLOWING MAJOR CATEGORIES AND THE CODE NUMBER RANGES ARE DEFINED

UNDEFINED

ØØ1 - Ø37 MASS STORAGE

200 - 237 CPU OPTIONS/SYSTEMS

THE FOLLOWING EXPLICIT DEVICE CLASS NUMBERS ARE DEFINED

MASS STORAGE

ØØ1

DISKS

ØØ2

TAPES

CPU OPTIONS/SYSTEMS

2ØØ

PROCESSOR

#### NOTE 3

DEVICE TYPES ARE FURTHER DEFINITION OF THE DEVICE CLASSES. THEREFORE A WHOLE SERIES OF DEVICE TYPE NUMBERS EXIST WITH EACH DEVICE CLASS.

THE FOLLOWING DEVICE TYPES ARE PRESENTLY DEFINED BY DEVICE CLASS

```
DEVICE TYPE CODES FOR DEVICE CLASS ØØ1 (DISK)
øøø
         UNDEFINED
ØØ1
          RKØ5/RKØ3
øø2
          RPØ3
ØØ3
          RF11
          RSØ4
ØØ4
øø5
          RSØ3
øø6
          RPØ4
øø7
          RPØ2 ON RP11C CONTROLLER
DEVICE TYPE CODES FOR DEVICE CLASS ØØ2 (TAPES)
øøø
         UNDEFINED
          TU56
øø1
ØØ2
          TU1Ø
          TU16
ØØ3
DEVICE TYPE CODES FOR DEVICE CLASS 200 (PROCESSOR)
          UNDEFINED
øøø
ØØ1
          11/45
          11/7Ø
øø2
```

# APPENDIX D

# 11/70 PARITY ERROR RESPONSE

### ACTION\*

SOURCE OF CYCLE	CYCLE TYPE		TYP	E OF ERROR	ALL TRAPS ENABLED	DISABLE WARNING TRAPS (CROO 1)	DISABLE ALL TRAPS (CR 01:00 3)	BITS SET IN ERROR REGISTER
CPU to	DATI/P	MAIN MAIN MAIN	BUS MEM MEM	TIMEOUT PARITY WANTED WD OTHER WD ADDRESS	ABORT ABORT ABORT TRAP TRAP	ABORT ABORT NO ACTION NO ACTION	ABORT ABORT ABORT 1 NO ACTION NO ACTION	15 11 8 15 11 1 5 11 2 or 3 11 2 or 3 11 4 or 5
CACHE		FAST	MEM	DATA	TRAP	NO ACTION	NO ACTION	11 6 or 7
	DATO/B			TIMEOUT PARITY	ABORT ABORT	ABORT ABORT	ABORT ABORT	15 11 0 15 11 1
:	DATI/P	MAIN	BUS	TIMEOUT	UNIBUS TIMEOUT	UNIBUS TIMEOUT	UNIBUS TIMEOUT	NONE
				PARITY	UNIBUS TIMEOUT	UNIBUS TIMEOUT	UNIBUS TIMEOUT	10 1
				WANTED WD	UNIBUS PB TRAP	UNIBUS PB, TRAP	UNIBUS PB	13 10 2 or 3
UNIBUS				OTHER WD	TRAP	NO ACTION	NO ACTION	10 2 or 3
(thru				ADDRESS	TRAP	NO ACTION	NO ACTION	10 4 or 5
MAP) to	0	FAST	MEM	DATA	TRAP	NO ACTION	NO ACTION	10 6 or 7
1	DATO/B	MAIN	BUS	TIMEOUT	UNIBUS TIMEOUT TRAP	UNIBUS TIMEOUT	UNIBUS TIMEOUT	10 8
		MAIN	BUS	PARITY	UNIBUS TIMEOUT TRAP	UNIBUS TIMEOUT	UNIBUS TIMEOUT	10 1
		FAST	MEM	ADDRESS	TRAP	NO ACTION	NO ACTION	10 4 or 5
1	DATI/P	MAIN	BUS	TIMEOUT	UNIBUS TIMEOUT ABORT TO 4	UNIBUS TIMEOUT ABORT TO 4	UNIBUS TIMEOUT	NONE
CPU to UNIBUS thru		MAIN	BUS	PARITY	UNIBUS TIMEOUT ABORT TO 4	UNIBUS TIMEOUT ABORT TO 4	UNIBUS TIME- OUT, ABORT TO 4	10 1
MAP to CACHE		MAIN	MEM	WANTED WD	UNIBUS PARITY ABORT TO 114	UNIBUS PARITY ABORT TO 114		13 10 9 2 or 3
		MAIN	MEM	OTHER WD	TRAP	NO ACTION	NO ACTION	10 2 or 3
				ADDRESS	TRAP	NO ACTION	NO ACTION	10 4 or 5
		FAST	MEM	DATA	TRAP	NO ACTION	NO ACTION	10 6 or 7

### 11/70 PARITY ERROR RESPONSE

SOURCE OF CYCLE	CYCLE TYPE	TYPE	OF	ERROR	ALL TRAPS ENABLED	DISABLE WARNING TRAPS (CROO 1)	DISABLE ALL TRAPS (CR 01:00 3)	BITS SET IN ERROR REGISTER
	DATAO/B	MAIN B	US	TIMEOUT	UNIBUS TIMEOUT ABORT TO 4	UNIBUS TIMEOUT ABORT TO 4	UNIBUS TIMEOUT ABORT TO 4	
		MAIN B	US	PARITY	UNIBUS TIMEOUT ABORT TO 4 TRAP	UNIBUS TIMEOUT ABORT TO 4	UNIBUS TIMEOUT, ABORT TO 4	10 1
		FAST M	EM	ADDRESS	TRAP	NO ACTION	NO ACTION	10 4 or 5

\*ALL TRAPS VECTOR TO 114

NOTES: 1. These tables only cover single errors.

- 2. The processor also will abort to 114 when doing DATI/P cycles on the UNIBUS and the device asserts PB. In these cases, bit 9 sets in the error register.
- 3. Errors on MASSBUS cycles are handled by the MASSBUS controls, and have an effect on the error register.

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