

**Level 68**

**Multics  
Subsystem  
Programming**

**Student Handbook**

**Course Code F15D**

**MAY 1981**

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Subsystem  
Programming**

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## COURSE DESCRIPTION

### F15D Multics Subsystem Programming

**Duration:** Five Days

**Intended For:** Advanced Multics PL/I programmers, familiar with standard Multics subroutines, who need to use advanced Multics subsystem writer's tools.

**Synopsis:** This intensive course describes how to bypass, replace, or supplement the standard Multics user interface by using system subroutines. Interprocess communication, tailoring the command environment, the message segment facility, the Multics ring mechanism, writing gates, dialing terminals to a process and writing I/O modules are among the topics covered in this course.

Interactive workshops are included to reinforce the material presented.

**Objectives:** Upon completion of this course, the student should be able to:

1. Use subsystem writer's subroutines.
2. Use a wide variety of facilities to create an environment tailored to the needs of a particular group of users.
3. Understand the conventions compilers should follow when creating object segments.
4. Understand how gates and I/O modules are written.

**Prerequisites:** Multics Concepts and Utilization (F01), Prerequisite Concepts for Programming on Multics (F10), Introductory Multics PL/I Programming (F15A), Advanced Multics PL/I Programming (F15B), PL/I Programming with Multics Subroutines (F15C) or equivalent experience.

**Major Topics:** Writing I/O Modules  
Interprocess Communication, Locking, and Timers  
Advanced hcs Utilization  
Program Library Management  
Tailoring the Command Environment  
Dialing Terminals to a Process  
Message Segment Facility  
Rings and Gate Writing  
Data Segments, Temporary Segments  
Creating an Error Table

**Manuals:** MPM - Subsystem Writers' Guide (AK92)  
SDN - Message Segment Facility (AN69)  
PLM - Library Maintenance (AN80)

F15D TOPIC MAP

DAY	MORNING TOPICS	AFTERNOON TOPICS
1	SUBSYSTEM WRITING ----- STORAGE SYSTEM SUBROUTINES	STORAGE SYSTEM SUBROUTINES (CONTINUED) ----- WORKSHOP #1
2	WORKSHOP #1 (CONT) ----- MULTICS SECURITY	THE COMMAND ENVIRONMENT ----- WORKSHOP #2
3	ADVANCED MULTICS I/O ----- WRITING I/O MODULES ----- WORKSHOP #3	INTERPROCESS COMMUNICATION ----- WORKSHOP #4
4	INTERPROCESS DATA BASE SHARING ----- INTRAPROCESS TIMER MANAGEMENT ----- WORKSHOP #5	THE STACK AND ARGUMENT LISTS ----- SPECIAL PROGRAMMING TECHNIQUES ----- WORKSHOP #6
5	THE PROCESS ENVIRONMENT ----- WORKSHOP # 7 ----- DIAL FACILITY	DIAL FACILITY DEMONSTRATION ----- MESSAGE SEGMENT FACILITY ----- PROGRAM LIBRARY MANAGEMENT

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STUDENT BACKGROUND

Multics Subsystem Programming (F15D)

NAME: \_\_\_\_\_ PHONE: \_\_\_\_\_

TITLE: \_\_\_\_\_

COMPANY ADDRESS: \_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

MANAGER: \_\_\_\_\_ OFFICE PHONE: \_\_\_\_\_

INSTRUCTOR'S NAME: \_\_\_\_\_

1. Do you meet the prerequisite as stated in the "Course Description" of the student text? If yes, check "a" or "b". If no, check "c" or "d".

a [ ] Prerequisite satisfied by attending course indicated in "Course Description".

b [ ] Meet prerequisite by equivalent experience (explain briefly)

\_\_\_\_\_

c [ ] Elected or instructed to attend course anyway.

d [ ] Was not aware of prerequisite.

2. What related Honeywell courses have you attended? Furnish dates and instructors if possible.

\_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

(PLEASE TURN OVER)

STUDENT BACKGROUND

Multics Subsystem Programming (F15D)

3. Check the boxes for which you have any related experience. (May be other than Honeywell's)

- |                              |                                     |                                  |                                   |
|------------------------------|-------------------------------------|----------------------------------|-----------------------------------|
| <input type="checkbox"/> PL1 | <input type="checkbox"/> COBOL      | <input type="checkbox"/> FORTRAN | <input type="checkbox"/> ASSEMBLY |
| <input type="checkbox"/> JCL | <input type="checkbox"/> OPERATIONS | <input type="checkbox"/> GCOS    | <input type="checkbox"/> MULTICS  |
| <input type="checkbox"/> NPS | <input type="checkbox"/> GRTS       | <input type="checkbox"/> CP6     | <input type="checkbox"/> OTHER    |
- 
- 

4. Detail any experience you have had which is related to the material in this course.

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5. Objectives for attending this course (May check more than one).

- Require information to provide support for a Multics system
  - To maintain an awareness of this product
  - To evaluate or compare its potentials
  - Required to use or implement
  - Need update from a previous release
  - Require a refresher
  - Other: \_\_\_\_\_
- 
- 
-

HONEYWELL MARKETING EDUCATION  
COURSE AND INSTRUCTOR EVALUATION FORM

INSTRUCTOR \_\_\_\_\_  
COURSE \_\_\_\_\_  
START DATE \_\_\_\_\_  
LOCATION \_\_\_\_\_  
STUDENT NAME \_\_\_\_\_ (OPTIONAL)

In the interest of developing training courses of high quality, and then improving on that base, we would like you to complete this questionnaire. Your information will aid us in making future revisions and improvements to this course. Both the instructor and his/her manager will review these responses.

Please complete the form and return it to the instructor upon the completion of the course. In questions 1 through 14, check the appropriate box and feel free to include additional comments. Attach additional sheets if you need more room for comments. Be objective and 'concrete' in your comments -- be critical when criticism is appropriate.

1. Considering the stated objectives of this course, rate the overall length of the course.

CAN'T JUDGE	TOO SHORT				ABOUT RIGHT				TOO LONG
0	1	2	3	4	5	6	7	8	9

COMMENTS \_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

2. Considering the objectives, rate the technical level at which the course was taught.

CAN'T JUDGE	NOT TECH ENOUGH				ABOUT RIGHT				TOO TECH
0	1	2	3	4	5	6	7	8	9

COMMENTS \_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

3. Considering the objectives, rate the emphasis placed on the more important topics.

CAN'T JUDGE	POOR				GOOD				EXCELLENT
0	1	2	3	4	5	6	7	8	9

COMMENTS \_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

4. Rate the sequence in which the topics were presented.

CAN'T JUDGE	POOR				GOOD				EXCELLENT
0	1	2	3	4	5	6	7	8	9

COMMENTS \_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

5. Rate the format and quality of the learning materials (slides, student handbooks, supplementary handouts, etc.).

CAN'T JUDGE	POOR					GOOD				EXCELLENT
0	1	2	3	4	5	6	7	8	9	

COMMENTS \_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

6. Rate the amount of time given for the completion of the workshops.

CAN'T JUDGE	TOO LITTLE TIME					ABOUT RIGHT				TOO MUCH TIME
0	1	2	3	4	5	6	7	8	9	

COMMENTS \_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

7. Rate the workshops' ability to relate back to and reinforce the material presented.

CAN'T JUDGE	POOR					GOOD				EXCELLENT
0	1	2	3	4	5	6	7	8	9	

COMMENTS \_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

8. Rate the physical condition of the classroom (space available, temperature, lighting, etc.).

CAN'T JUDGE	POOR					GOOD				EXCELLENT
0	1	2	3	4	5	6	7	8	9	

COMMENTS \_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

9. Rate the physical condition of the lab or workshop room. (systems configuration, space available, learning tools, terminals, tables, etc.).

CAN'T JUDGE	POOR				GOOD				EXCELLENT	
<input type="checkbox"/> 0	<input type="checkbox"/> 1	<input type="checkbox"/> 2	<input type="checkbox"/> 3	<input type="checkbox"/> 4	<input type="checkbox"/> 5	<input type="checkbox"/> 6	<input type="checkbox"/> 7	<input type="checkbox"/> 8	<input type="checkbox"/> 9	

COMMENTS \_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

10. Rate your instructor's demonstrated knowledge of the course material.

CAN'T JUDGE	POOR				GOOD				EXCELLENT	
<input type="checkbox"/> 0	<input type="checkbox"/> 1	<input type="checkbox"/> 2	<input type="checkbox"/> 3	<input type="checkbox"/> 4	<input type="checkbox"/> 5	<input type="checkbox"/> 6	<input type="checkbox"/> 7	<input type="checkbox"/> 8	<input type="checkbox"/> 9	

COMMENTS \_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

11. Rate your instructor's ability to convey the technical aspects of the various topics.

CAN'T JUDGE	POOR				GOOD				EXCELLENT	
<input type="checkbox"/> 0	<input type="checkbox"/> 1	<input type="checkbox"/> 2	<input type="checkbox"/> 3	<input type="checkbox"/> 4	<input type="checkbox"/> 5	<input type="checkbox"/> 6	<input type="checkbox"/> 7	<input type="checkbox"/> 8	<input type="checkbox"/> 9	

COMMENTS \_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

12. Rate the classroom and workshop assistance given you by your instructor.

CAN'T JUDGE	POOR				GOOD				EXCELLENT	
<input type="checkbox"/> 0	<input type="checkbox"/> 1	<input type="checkbox"/> 2	<input type="checkbox"/> 3	<input type="checkbox"/> 4	<input type="checkbox"/> 5	<input type="checkbox"/> 6	<input type="checkbox"/> 7	<input type="checkbox"/> 8	<input type="checkbox"/> 9	

COMMENTS \_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

13. Rate the instructor's ability to create an environment in which you felt free to ask questions.

CAN'T JUDGE	POOR			GOOD			EXCELLENT		
0	1	2	3	4	5	6	7	8	9

COMMENTS \_\_\_\_\_  
\_\_\_\_\_  
\_\_\_\_\_

14. Rate the relevance of the skills learned in the course with respect to your job or further training.

CAN'T JUDGE	POOR			GOOD			EXCELLENT		
0	1	2	3	4	5	6	7	8	9

COMMENTS \_\_\_\_\_  
\_\_\_\_\_  
\_\_\_\_\_

15. What did you like most about this course?

\_\_\_\_\_  
\_\_\_\_\_  
\_\_\_\_\_  
\_\_\_\_\_  
\_\_\_\_\_  
\_\_\_\_\_

16. What did you like least about this course?

\_\_\_\_\_  
\_\_\_\_\_  
\_\_\_\_\_  
\_\_\_\_\_  
\_\_\_\_\_  
\_\_\_\_\_



17. Other comments please:

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18. Of the following job categories, check the ones which most nearly represent the bulk of your experience, and to the right of your responses indicate the number of years you have acted in that capacity.

- Applications Programmer. . . . \_\_\_\_\_ years
- Field Engineering Analyst. . . \_\_\_\_\_ years
- Manager. . . . . \_\_\_\_\_ years
- Marketing Analyst. . . . . \_\_\_\_\_ years
- Salesperson. . . . . \_\_\_\_\_ years
- Secretary. . . . . \_\_\_\_\_ years
- Systems Analyst. . . . . \_\_\_\_\_ years
- Systems Programmer . . . . . \_\_\_\_\_ years
- Other. . . . . \_\_\_\_\_ years

Please give "other" title \_\_\_\_\_

TOPIC I

Subsystem Writing

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## INTRODUCTION

- A BASIC GOAL OF THE Multics SYSTEM DESIGN PHILOSOPHY:  
TO PROVIDE A SYSTEM WHICH IS OPEN-ENDED AND CAPABLE OF SUPPORTING  
USER DESIGNED SUBSYSTEMS
  - ▮ TO ACHIEVE THIS, Multics
    - ▮ HAS BEEN HIGHLY MODULARIZED
      - ▮ FUNCTIONALITY LOCALIZED
      - ▮ COMPLEXITY OF ANY GIVEN MODULE MINIMIZED
    - ▮ IS MOSTLY WRITTEN IN PL/I
      - ▮ MORE EASILY READ THAN ALM
      - ▮ ENABLES ADOPTION OF SUBSYSTEMS FROM OTHER MACHINES USING PL/I
    - ▮ FEATURES A WEALTH OF TOOLS TO HELP DESIGN, IMPLEMENT, AND MAINTAIN SUBSYSTEMS
  
- THIS COURSE IS DESIGNED TO:
  - ▮ INTRODUCE MOST TOPICS COVERED IN THE SUBSYSTEM WRITERS' GUIDE (SWG)
  - ▮ COVER IN DETAIL SEVERAL ADVANCED TOOLS AND TECHNIQUES OFTEN USED IN WRITING SUBSYSTEMS
  - ▮ PROVIDE INSIGHT INTO HOW TO WRITE SUBSYSTEMS "THE Multics WAY"

## TERMINOLOGY

- A SUBSYSTEM CAN BE DEFINED A VARIETY OF WAYS:

- ▮ A "SYSTEM" WHICH OPERATES WITHIN THE CONFINES OF ANOTHER, LARGER SYSTEM
- ▮ PROGRAM(S) THAT PROVIDE A SPECIAL ENVIRONMENT FOR SOME PARTICULAR PURPOSE
- ▮ PROGRAM(S) THAT PROVIDE A NUMBER OF OPERATIONS ON SOME RESTRICTED UNIVERSE OF DATA

- EXAMPLES OF STANDARD Multics SUBSYSTEMS: qedx, ted, emacs, calc, probe, read\_mail, send\_mail, help, ~~abs~~

- A SUBSYSTEM IS SAID TO BE CLOSED IF:

- ▮ ALL NECESSARY OPERATIONS CAN BE HANDLED WITHIN THE SUBSYSTEM
- ▮ NO WAY EXISTS TO USE THE NORMAL MULTICS ENVIRONMENT FROM WITHIN THE SUBSYSTEM
- ▮ EXAMPLE: THE 'fast' SUBSYSTEM

DESIGN CONCERNS

● HAS THE PROBLEM ALREADY BEEN SOLVED?

● SECURITY

▮ IS SUBVERSION A REAL CONCERN?

▮ WILL ACL ALONE SUFFICE, OR MUST WE RESORT TO RINGS AND AIM?

▮ CLOSED SUBSYSTEM?

● HOW SHALL WE INTERFACE WITH THE STORAGE SYSTEM?

TEMPORARY SEGS

AREAS

PERMANENT SEGS, MSFs

NAME AND ADDRESS SPACE MANAGEMENT

USE MULTICS I/O SYSTEM?

● WILL PROCESSES NEED TO COMMUNICATE WITH EACH OTHER?

## DESIGN CONCERNS

- DOES THE SUBSYSTEM HAVE A "MULTICS FLAVOR"?
- DOCUMENTATION
- SUBSYSTEM LIBRARY MAINTENANCE

CAPABILITIES FOR SUBSYSTEM DESIGN IN MULTICS

- THE OPPORTUNITIES FOR SUBSYSTEM DESIGN IN Multics ARE VIRTUALLY UNLIMITED, AND THE SUBSYSTEM DESIGNER MAY:

- ▮ MODIFY THE COMMAND INTERFACE TO THE Multics STORAGE SYSTEM
- ▮ MANIPULATE THE ADDRESS SPACE OF A USER PROCESS
- ▮ MODIFY THE COMMAND ENVIRONMENT OF A USER PROCESS
- ▮ WRITE COMMAND AND/OR ACTIVE FUNCTION PROCEDURES
- ▮ WRITE A COMMAND PROCESSOR PROCEDURE
- ▮ HANDLE CONTROL COMMUNICATION BETWEEN ANY NUMBER OF ASYNCHRONOUS, COOPERATING PROCESSES
- ▮ CONTROL CONCURRENT ACCESS TO CRITICAL, SHARED DATA BASES
- ▮ USE TIMERS
- ▮ INTERFACE NEW I/O DEVICES, MONITOR EXISTING I/O DEVICES, ETC.
- ▮ MODIFY, RESTRICT, OR REPLACE ENTIRELY THE PROCESS ENVIRONMENT

CAPABILITIES FOR SUBSYSTEM DESIGN IN MULTICS

- ▮ DIAL TERMINALS TO A PROCESS, OR ALLOW A PROCESS TO DIAL OUT TO A TERMINAL
  
- ▮ WRITE GATES
  
- ▮ MANIPULATE MESSAGE SEGMENTS
  
- ▮ CREATE, UPDATE, AND IN GENERAL, MAINTAIN PROGRAM LIBRARIES
  
- ▮ AND MANY, MANY OTHER THINGS



## SUBSYSTEM DESIGN TOOLS

- TO ACHIEVE SOME OF THE SUBSYSTEM DESIGN TASKS MENTIONED ABOVE, THE DESIGNER HAS AVAILABLE A WIDE VARIETY OF RESOURCES INCLUDING:

- ▮ COMMANDS AND SUBROUTINES

- ▮ SOURCE PROGRAMS

- ▮ WHOSE PERUSAL SHOWS THE DESIGNER HOW Multics DOES IT

- ▮ WHICH MAY BE COPIED AND MODIFIED TO YIELD CUSTOMIZED BEHAVIOR

- ▮ PL/1 AND ALM INCLUDE FILES

- ▮ PROGRAM LIBRARY MAINTENANCE TOOLS

- ▮ EXPEDITE ACCESS TO SYSTEM SOURCE, OBJECT AND INFO SEGMENTS

- ▮ MAINTAIN USER SUBSYSTEM LIBRARIES JUST AS THEY MAINTAIN THE Multics LIBRARIES THEMSELVES

## STORAGE SYSTEM SUBROUTINES

- TOPICS 2, 3 AND 4 PRESENT THE SOFTWARE WRITERS' GUIDE (SWG) SUBROUTINES USED IN MANIPULATING THE STORAGE SYSTEM
  
- THE FOLLOWING LIST PROVIDES A COMPARISON OF THE STORAGE SYSTEM MANIPULATING SUBROUTINES COVERED IN F15C AND F15D
  - ▮ EXCEPT WHERE NOTED F15C SUBROUTINES ARE DOCUMENTED IN THE SUBROUTINES MANUAL (AG93) AND F15D SUBROUTINES ARE DOCUMENTED IN THE SWG (AK92)

STORAGE SYSTEM SUBROUTINES

F15C	F15D
<b>CREATING STORAGE SYSTEM ENTITIES</b>	
<pre> hcs_\$append_branch hcs_\$append_branchx hcs_\$append_link hcs_\$create_branch_ hcs_\$make_seg                     </pre>	
<b>DELETING STORAGE SYSTEM ENTITIES</b>	
<pre> delete hcs_\$delentry_file hcs_\$delentry_seg hcs_\$del_dir_tree (AK92)                     </pre>	
<b>OBTAINING STATUS INFORMATION</b>	
<pre> hcs_\$status_ hcs_\$status_long hcs_\$status_minf hcs_\$status_mins                     </pre>	<pre> hcs_\$get_author hcs_\$get_bc_author hcs_\$get_link_target hcs_\$get_max_length hcs_\$get_max_length_seg hcs_\$get_safety_sw hcs_\$get_safety_sw_seg (hcs_\$set_max_length) (hcs_\$set_max_length_seg) (hcs_\$set_safety_sw) (hcs_\$set_safety_sw_seg)                     </pre>

STORAGE SYSTEM SUBROUTINES

F15C	F15D
WORKING, DEFAULT, AND PROCESS DIRECTORIES	
change_wdir get_default_wdir_ (AK92) get_pdir_ get_wdir_	
MANIPULATING THE ADDRESS AND NAME SPACES	
hcs_\$fs_get_path_name hcs_\$fs_get_ref_name hcs_\$fs_get_seg_ptr hcs_\$initiate hcs_\$initiate_count hcs_\$make_seg hcs_\$terminate_file hcs_\$terminate_name hcs_\$terminate_noname hcs_\$terminate_seg term_\$refname term_\$seg_ptr term_\$single_refname term_\$term_ term_\$unsnap	
MULTISEGMENT FILES	
	msf_manager_\$acl_add msf_manager_\$acl_delete msf_manager_\$acl_list msf_manager_\$acl_replace msf_manager_\$adjust msf_manager_\$close msf_manager_\$get_ptr msf_manager_\$open

STORAGE SYSTEM SUBROUTINES

F15C	F15D
NAMING AND MOVING DIRECTORY ENTRIES	
hcs_\$chname_file hcs_\$chname_seg hcs_\$fs_move_file hcs_\$fs_move_seg	
AFFECTING LENGTH OF ENTRIES	
adjust_bit_count_ hcs_\$set_bc hcs_\$set_bc_seg hcs_\$truncate_file hcs_\$truncate_seg	
MANIPULATING PATHNAMES	
absolute_pathname_ absolute_pathname_\$add_suffix expand_pathname_ expand_pathname_\$add_suffix	
MANIPULATING THE STAR AND EQUAL CONVENTION	
	check_star_name_ get_equal_name_ hcs_\$star_ match_star_name_

STORAGE SYSTEM SUBROUTINES

F15C	F15D
<b>AREAS</b>	
<pre>get_system_free_area (AK92)</pre>	<pre>area_info area_status * create_area * define_area release_area set_system_storage * set_user_storage *</pre>
<b>SECURITY</b>	
<pre>get_group_id get_group_id \$tag star hcs \$add_acl_entries hcs \$add_dir_acl_entries hcs \$delete_acl_entries hcs \$delete_dir_acl_entries hcs \$fs get mode hcs \$list_acl hcs \$list_dir_acl hcs \$replace_acl hcs \$replace_dir_acl</pre>	<pre>cross_ring cross_ring_io \$allow cross cu \$level_get (AG93) cu \$level_set (AG93) get_ring hcs \$add_dir_inacl_entries hcs \$add_inacl_entries hcs \$delete_dir_inacl_entries hcs \$delete_inacl_entries hcs \$get_dir_ring_brackets hcs \$get_ring_brackets hcs \$get_user_effmode hcs \$list_dir_inacl_entries hcs \$list_inacl hcs \$replace_dir_inacl hcs \$replace_inacl hcs \$set_dir_ring_brackets hcs \$set_entry_bound hcs \$set_entry_bound_seg hcs \$set_ring_brackets</pre>

\* COMMANDS (INCLUDED FOR COMPLETENESS)

TOPIC II

Storage System Subroutines

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Multisegment Files . . . . .	2-8
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OBTAINING STATUS INFORMATION

● hcs\_\$get\_author

⌋ call hcs\_\$get\_author (dir\_name, entryname, chase, author, code);

⌋ RETURNS Personid.Projectid.tag OF THE CREATOR OF A SEGMENT, DIRECTORY, MULTISEGMENT FILE OR LINK

● hcs\_\$get\_bc\_author

⌋ call hcs\_\$get\_bc\_author (dir\_name, entryname, bc\_author, code);

⌋ RETURNS Personid.Projectid.tag OF THE BIT COUNT AUTHOR OF A SEGMENT OR DIRECTORY

⌋ BIT COUNT AUTHOR = LAST PERSON WHO SET THE BIT COUNT



## OBTAINING STATUS INFORMATION

### ● hcs\_\$get\_max\_length, hcs\_\$get\_max\_length\_seg

⌋ call hcs\_\$get\_max\_length (dir\_name, entryname, max\_length, code);

⌋ call hcs\_\$get\_max\_length\_seg (seg\_ptr, max\_length, code);

⌋ RETURNS THE MAXIMUM LENGTH (IN WORDS) OF A SEGMENT, DIRECTORY OR LINK TARGET

⌋ SUBROUTINES THAT CAN CHANGE THE MAXIMUM LENGTH OF A SEGMENT

⌋ hcs\_\$set\_max\_length, hcs\_\$set\_max\_length\_seg

⌋ SAME CALL ARGUMENTS AS ABOVE

⌋ A DIRECTORY CANNOT HAVE ITS MAXIMUM LENGTH CHANGED

⌋ ONCE MAX LENGTH HAS BEEN SET, AN out\_of\_bounds FAULT OCCURS WHEN REFERENCING BEYOND END OF SEGMENT

⌋ MAXIMUM LENGTH IS SET IN UNITS OF 1024 WORDS

⌋ REQUESTED LENGTH MAY NOT EXCEED sys\_info\_\$max\_seg\_size

⌋ CANNOT USE TO SHORTEN SEGMENT

⌋ DEFAULT MAX LENGTH OF A SEGMENT IS 255K

⌋ stack\_4 HAS INITIAL MAX LENGTH OF 64K

OBTAINING STATUS INFORMATION

● hcs\_\$get\_safety\_sw, hcs\_\$get\_safety\_sw\_seg

┆ call hcs\_\$get\_safety\_sw (dir\_name, entryname, safety\_sw, code);

┆ call hcs\_\$get\_safety\_sw\_seg (seg\_ptr, safety\_sw, code);

┆ RETURNS THE VALUE OF THE SAFETY SWITCH OF A DIRECTORY OR SEGMENT

┆ SUBROUTINES THAT CAN CHANGE THE VALUE OF THE SAFETY SWITCH:

┆ hcs\_\$set\_safety\_sw, hcs\_\$set\_safety\_sw\_seg

┆ SAME CALL ARGUMENTS AS ABOVE

● hcs\_\$get\_link\_target

┆ call hcs\_\$get\_link\_target (dir\_name, entryname, link\_dir\_name,  
link\_entryname, code);

┆ RETURNS THE TARGET PATHNAME OF A LINK

## OBTAINING STATUS INFORMATION

- ON THE FOLLOWING PAGES IS AN EXAMPLE USING SOME OF THE SUBROUTINES PROVIDING STATUS INFORMATION

▮ IT ALSO SERVES AS A REVIEW OF SOME ITEMS INTRODUCED IN F15C

▮ WRITING A COMMAND

▮ USING ioa\_ AND com\_err\_

▮ IN YOUR FIRST WORKSHOP YOU WILL BE ASKED TO ENHANCE THIS PROGRAM

OBTAINING STATUS INFORMATION

```
STATUS: proc;
dcl cu_$arg_count entry (fixed bin),
    cu_$arg_ptr entry (fixed bin, ptr, fixed bin, fixed bin (35)),
    hcs_$status_minf entry (char (*), char (*), fixed bin (1),
        fixed bin (2), fixed bin (24), fixed bin (35)),
    hcs_$get_safety_sw entry (char(*),char(*),bit (1),fixed bin (35)),
    hcs_$get_max_length entry (char (*), char (*), fixed bin (19),
        fixed bin (35)),
    hcs_$get_author entry (char (*), char (*), fixed bin (1), char (*),
        fixed bin (35)),
    expand_pathname_ entry (char(*),char(*),char(*),fixed bin (35)),
    (ioa_, com_err ) entry options (variable);
dcl nargs fixed bin;
dcl arg1 fixed bin;
dcl argp ptr;
dcl arg char (arg1) based (argp);
dcl dir char (168);
dcl entry char (32);
dcl code fixed bin (35);
dcl type fixed bin (2),
    bc fixed bin (24),
    author char (32),
    max_length fixed bin (19),
    safety_sw bit (1),
    ME char (6) static init ("STATUS") options (constant);
dcl error_table_$wrong_no_of_args ext fixed bin (35);

/* VERIFY NUMBER OF ARGUMENTS */
    call cu_$arg_count (nargs);
    if nargs ^= 1 then do;
        call com_err_ (error_table_$wrong_no_of_args, ME);
        return;
    end;

/* PROCESS SEGMENT NAME ARGUMENT */
    call cu_$arg_ptr (1, argp, arg1, code);
    call expand_pathname_ (arg, dir, entry, code);
    if code ^= 0 then call ERROR;

/* FIND OUT WHAT TYPE OF BRANCH IT IS */
    call hcs_$status_minf (dir, entry, 0, type, bc, code);
    if code ^= 0 then call ERROR;
```

OBTAINING STATUS INFORMATION

```
/* TELL THE USER */
    if type = 2 & bc ^= 0 then
        call ioa ("^a is a ^i component multisegment file",
            entry, bc);
    else call ioa (
        "^a is a ^[link^;segment^;directory^]
        ^[with bit count ^i^;^s^]",
        entry, type+1, (type = 1), bc);

/* GET OTHER INFORMATION AND REPORT IT TO THE USER */
    call hcs $get_author (dir, entry, 0, author, code);
    if code ^= 0 then call ERROR;
    call hcs $get_max_length (dir, entry, max_length, code);
    if code ^= 0 then call ERROR;
    call hcs $get_safety_sw (dir, entry, safety_sw, code);
    if code ^= 0 then call ERROR;
    call ioa_ ("          It was created by ^a,
        it has a max length of ^i,
        and the safety switch is ^[on^;off^].",
        author, max_length, safety_sw);
ERROR:   proc;
        call com_err_ (code, ME);
        goto FINISH; ← change
    end;

FINISH: end STATUS;
```

OBTAINING STATUS INFORMATION

!STATUS STATUS.pl1  
STATUS.pl1 is a segment  
    with bit count 23256  
It was created by NDibble.MEDmult.a,  
    it has a max length of 261120,  
    and the safety switch is off.

!STATUS <  
NDibble is a directory  
  
It was created by Initializer.SysDaemon.z,  
    it has a max length of 65536,  
    and the safety switch is on.

!STATUS test\_file  
test\_file is a 2 component multisegment file  
    It was created by NDibble.MEDmult.a,  
    it has a max length of 65536,  
    and the safety switch is off.

!lk >udd>F15D>s1 blurp

!STATUS blurp  
blurp is a link  
  
It was created by NDibble.MED.a,  
    it has a max length of 65536,  
    and the safety switch is off.

!sml STATUS.pl1 40960

!STATUS STATUS.pl1  
STATUS.pl1 is a segment  
    with bit count 23292  
It was created by NDibble.MEDmult.a,  
    it has a max length of 40960,  
    and the safety switch is off.

!STATUS \*\*  
STATUS: Entry not found.

## MULTISEGMENT FILES

- MULTISEGMENT FILES ARE:
  - ▮ FILES THAT USE MORE THAN ONE SEGMENT FOR STORAGE
  - ▮ COMPOSED OF ONE OR MORE COMPONENTS, EACH IS A SEGMENT, AND IS IDENTIFIED BY AN UNSIGNED INTEGER
  - ▮ VIEWED BY MANY MULTICS SUBROUTINES AS DIRECTORIES
  - ▮ USED FOR LARGE LISTINGS, INDEXED FILES, ETC.
  - ▮ MANAGED BY THE `msf_manager_` SUBROUTINE
- MANIPULATING A MULTISEGMENT FILE REQUIRES USE OF A MULTISEGMENT FILE CONTROL BLOCK
  - ▮ THE CONTROL BLOCK FOR A MULTISEGMENT FILE IS CREATED AND MAINTAINED BY THE `msf_manager_` IN THE USER'S PROCESS DIRECTORY
- LOCATIONS IN A MULTISEGMENT FILE ARE SPECIFIED BY A PATHNAME, COMPONENT NUMBER AND WORD OFFSET WITHIN THE COMPONENT

## MULTISEGMENT FILES

### ● msf\_manager\_\$open

- ⌋ call msf\_manager\_\$open (dir\_name, entryname, fcb\_ptr, code);
- ⌋ CREATES A FILE CONTROL BLOCK IN SYSTEM FREE STORAGE AND RETURNS A FILE CONTROL BLOCK POINTER
- ⌋ THE MSF NEEDN'T EXIST (A FCB IS STILL ALLOCATED)
- ⌋ THE fcb\_ptr IS USED BY ALL FUTURE CALLS TO msf\_manager\_

### ● msf\_manager\_\$get\_ptr

- ⌋ call msf\_manager\_\$get\_ptr (fcb\_ptr, component, create\_sw, seg\_ptr, bc, code);
- ⌋ RETURNS A POINTER TO A SPECIFIED COMPONENT IN THE MSF
- ⌋ COMPONENT IS AUTOMATICALLY CREATED, IF create\_sw = "1"b
- ⌋ IF THE FILE IS A SINGLE SEGMENT FILE AND A COMPONENT GREATER THAN 0 IS REQUESTED, THE SEGMENT IS CONVERTED INTO A MSF



MULTISEGMENT FILES

!pr MSF.pl1

```
MSF: proc;
  dcl msf_manager_$open entry (char(*), char(*),
                               ptr, fixed bin(35));
  dcl msf_manager_$get_ptr entry (ptr, fixed bin, bit(1),
                                  ptr, fixed bin(24), fixed bin(35));
  dcl hcs_$initiate entry (char(*), char(*), char(*),
                           fixed bin(1), fixed bin(2), ptr, fixed bin(35));
  dcl code fixed bin (35);
  dcl (fcb_ptr, seg_ptr) ptr;
  dcl bc fixed bin (24);
  dcl sysprint file;
  dcl ioa_entry() options(variable);

  call hcs_$initiate (">udd>MED>nd>F15D", "test_file", "",
                     0, 0, seg_ptr, code);

  call ioa ("^p^/", seg_ptr);
  /* PROBE BREAKPOINT SET HERE */

  call msf_manager_$open (">udd>MED>nd>F15D", "test_file",
                          fcb_ptr, code);

  call msf_manager_$get_ptr (fcb_ptr, 0, "0"b, seg_ptr, bc, code);
  call ioa ("~/Component 0 starts at ^p", seg_ptr);
  call msf_manager_$get_ptr (fcb_ptr, 1, "1"b, seg_ptr, bc, code);
  call ioa ("~/Component 1 starts at ^p", seg_ptr);
```

end MSF;

!create test\_file

!ls test\_file

Segments = 1, Lengths = 0.

r w 0 test\_file

!MSF

503i0

Stopped after line 17 of MSF. (level 7)

!...lrn 503

503 >udd>MED>nd>F15D>test\_file

MULTISEGMENT FILES

!continue

Component 0 starts at 503!0

Component 1 starts at 501!0

!lrn 503 501

503 >udd>MED>nd>F15D>test\_file>0

501 >udd>MED>nd>F15D>test\_file>1

!ls test\_file

Multisegment-files = 1, Lengths = 1.

r w 1 test\_file

## MULTISEGMENT FILES

### ● msf\_manager\_\$adjust

- ▮ call msf\_manager\_\$adjust (fcb\_ptr, component, bc, switch, code);
- ▮ OPTIONALLY SETS THE BIT COUNT, TRUNCATES AND TERMINATES A COMPONENT
- ▮ SWITCH HAS 3 BITS
  - ▮ IF BIT 1 IS ON THE BIT COUNT IS SET (BIT COUNT OF ALL COMPONENTS < component SET TO sys\_info\$max\_seg\_size)
  - ▮ IF BIT 2 IS ON THE COMPONENT IS TRUNCATED
  - ▮ IF BIT 3 IS ON THE COMPONENT IS TERMINATED
- ▮ ALL COMPONENTS WITH NUMBERS GREATER THAN THE GIVEN COMPONENT ARE DELETED

### ● msf\_manager\_\$close

- ▮ call msf\_manager\_\$close (fcb\_ptr);
- ▮ TERMINATES ALL COMPONENTS OF THE MSF, FREES THE FILE CONTROL BLOCK, AND SETS fcb\_ptr NULL

MULTISEGMENT FILES

msf\_manager\_ ACL ENTRY POINTS ARE SIMILAR TO hcs\_ ACL ENTRY POINTS

hcs_	msf_manager_
list_acl	acl_list
replace_acl	acl_replace
add_acl_entries	acl_add
delete_acl_entries	acl_delete

## TEMPORARY SEGMENTS

- TEMPORARY SEGMENTS

- ┆ RESIDE IN THE PROCESS DIRECTORY

- ┆ ARE MANAGED AS A POOL

- ┆ HAVE A NAME OF THE FORM:

- <unique\_name>.temp.<seg\_number>

- ┆ ARE HEAVILY USED BY MANY COMMANDS, SUCH AS qedx

- TEMPSEG POOLING ENABLES THE USE OF THE SAME TEMPSEG MORE THAN ONCE DURING THE LIFE OF A PROCESS, RESULTING IN A REDUCED COST TO THE PROCESS

- THE `list temp segments` COMMAND GIVES DETAILED INFORMATION ABOUT THE STATE OF A PROCESS' TEMPORARY SEGMENT POOL

## TEMPORARY SEGMENTS

### ● FOUR SUBROUTINES MANIPULATE TEMPORARY SEGMENTS:

▮ get\_temp\_segments\_

▮ call get\_temp\_segments\_ (program\_name, ptrs, code);

▮ RETURNS POINTERS TO TEMPORARY SEGMENTS FOR A SPECIFIED PROGRAM

▮ CALLER SUPPLIES

▮ NAME OF REQUESTING PROGRAM

▮ AN ARRAY OF POINTERS WHOSE EXTENT EQUALS THE NUMBER OF  
TEMPSEGS DESIRED

▮ SEE ALSO get\_temp\_segment\_

## TEMPORARY SEGMENTS

- ⌋ release\_temp\_segments\_
- ⌋ call release\_temp\_segments\_ (program\_name, ptrs, code);
- ⌋ USED TO RETURN TEMPORARY SEGMENTS TO THE FREE POOL (SO THAT THEY MAY BE REUSED, IF DESIRED)
- ⌋ CALLER SUPPLIES
  - ⌋ NAME OF PROGRAM "OWNING" THE TEMPSEGS
  - ⌋ ARRAY OF POINTERS TO THE TEMPSEGS TO BE RETURNED TO POOL
- ⌋ THE TEMPORARY SEGMENTS BEING 'RETURNED' ARE NOT DELETED
- ⌋ IF RELEASE IS SUCCESSFUL, POINTERS ARE NULLED
- ⌋ ANY ATTEMPT TO RELEASE TEMPSEGS NOT "OWNED" BY REQUESTOR RESULTS IN error\_table\_\$argerr; PASSED POINTERS ARE UNCHANGED.
- ⌋ SEE ALSO release\_temp\_segment\_

## TEMPORARY SEGMENTS

```
! pr DEMO_TEMP_SEGS.pl1

DEMO_TEMP_SEGS: proc;
dcl get_temp_segments_entry (char(*), (*) ptr, fixed bin(35));
dcl release_temp_segments_entry (char(*), (*) ptr, fixed bin(35));
dcl error_table $argerr fixed bin(35) ext static;
dcl ioa_entry_options (variable);
dcl p_array(3) ptr;
dcl code fixed bin(35);
call get_temp_segments ("requestor_1", p_array, code);
call ioa ("Check the following tempseg segnos:~/^(^2x^p^)", p_array);
call release_temp_segments ("requestor_2", p_array, code);
if code = error_table $argerr then call ioa
    ("requestor_2 may not free segments owned by requestor_1.");
call ioa
    ("Pointers after a bad call to release_temp_segments are:~/^(^2x^p^)",
    p_array);
call release_temp_segments ("requestor_1", p_array, code);
end DEMO_TEMP_SEGS;
```

```
! list_temp_segments
    12 Segments,  11 Free
```

```
!BBBJHmQJDkmGxW.temp.0315  command_processor_
```

```
! DEMO_TEMP_SEGS
Check the following tempseg segnos:
 344|0  354|0  355|0
requestor_2 may not free segments owned by requestor_1.
Pointers after a bad call to release_temp_segments are:
 344|0  354|0  355|0
```

```
! lrn 344 354 355
 344 >process_dir_dir>!BcDBdwpbBBBBBB>!BBBJHmQJGFjKqg.temp.0344
 354 >process_dir_dir>!BcDBdwpbBBBBBB>!BBBJHmQJkkPNPb.temp.0354
 355 >process_dir_dir>!BcDBdwpbBBBBBB>!BBBJHmQJkkkWlf.temp.0355
```

```
! list_temp_segments
    12 Segments,  11 Free
```

```
!BBBJHmQJDkmGxW.temp.0315  command_processor_
```



TOPIC III

Storage System Subroutines (cont)

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STAR AND EQUAL CONVENTIONS

● MOTIVATION

⌋ THE WRITER OF A SUBSYSTEM OR COMMAND MUST DECIDE IF HE WILL ALLOW THE USER TO SPECIFY ENTRY NAMES THAT USE THE STAR AND EQUAL CONVENTIONS

⌋ RECALL THAT ENTRY NAMES USING THESE CONVENTIONS CONTAIN THE CHARACTERS "\*", "?", "=", OR "%"

⌋ THE FOLLOWING SUBROUTINES ARE USED TO PROCESS SUCH ENTRY NAMES

⌋ hcs\_\$star\_

⌋ get\_equal\_name\_

⌋ check\_star\_name\_

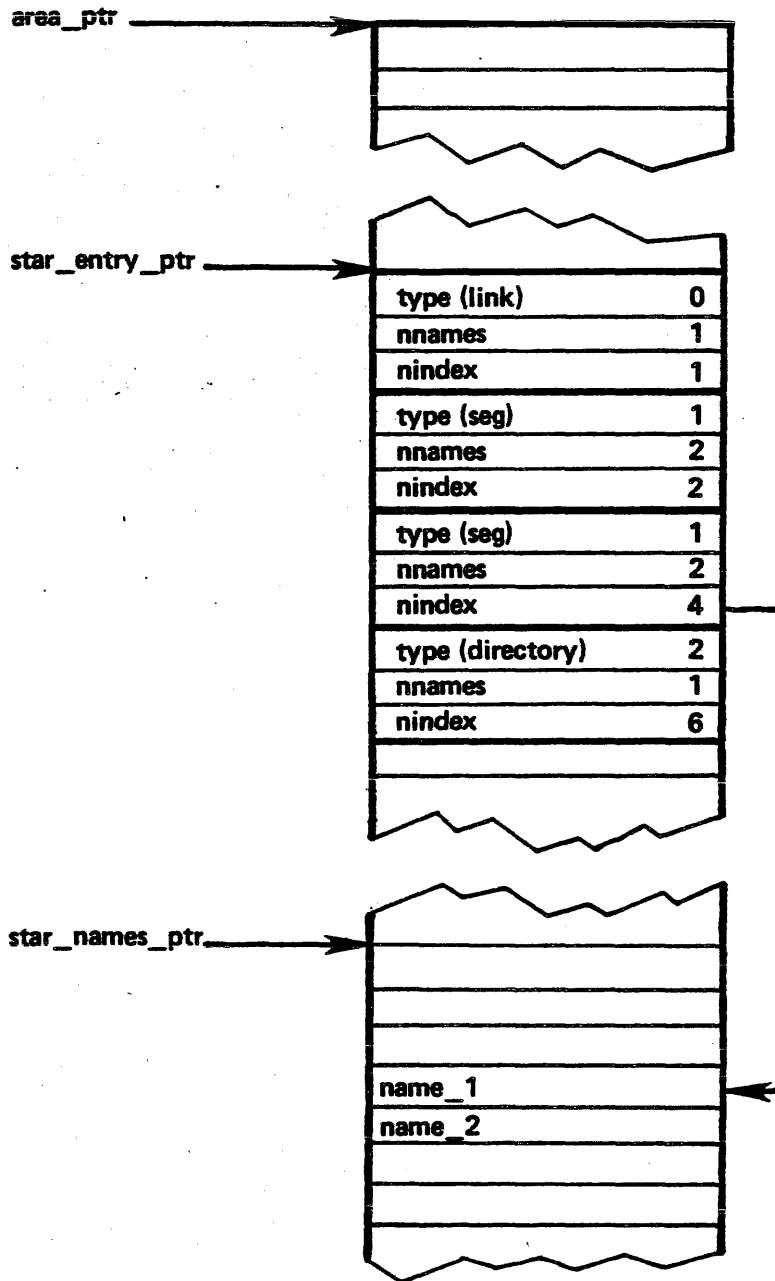
⌋ match\_star\_name\_

## STAR AND EQUAL CONVENTIONS

### ● hcs\_\$star\_

- ▮ call hcs\_\$star\_ (dir\_name, star\_name, star\_select\_sw, area\_ptr, star\_entry\_count, star\_entry\_ptr, star\_names\_ptr, code);
  
- ▮ FOR A GIVEN DIRECTORY, RETURNS AN ARRAY OF ENTRY NAMES THAT MATCH A GIVEN STARNAME
  
- ▮ star\_select\_sw DICTATES OPERATION:
  - 1 - LINK NAMES ONLY
  - 2 - SEGS AND DIRS ONLY  
(MSF'S COME BACK AS DIRS)
  - 3 - SEGS, DIRS, AND LINKS
  
- ▮ USER PROVIDES AN AREA FOR RETURNED ENTRY NAMES AND RELATED INFORMATION

STAR AND EQUAL CONVENTIONS



STAR AND EQUAL CONVENTIONS

```
!print lspl1.pl1

lspl1: proc;
dcl get_system_free_area_entry returns (ptr);

dcl hcs $star_entry (char (*), char (*), fixed bin (2),
    ntr. fixed bin, ptr, ptr, fixed bin (35));

dcl 1 star_entries (star_entry_count) aligned_based (star_entry_ptr),
    2 type fixed binary (2) unsigned unaligned,
    2 nnames fixed binary (16) unsigned unaligned,
    2 nindex fixed binary (18) unsigned unaligned;

dcl star_names (sum (star_entries (*).nnames))
    char (32) based (star_names_ptr);

dcl star_entry_count fixed binary,
    star_entry_ptr pointer,
    star_names_ptr pointer,
    code fixed bin (35),
    ioa_entry options (variable);

dcl (i, j) fixed bin;

    call hcs $star (">udd>MED>nd>F15D", "/*.pl1",
        2, get_system_free_area (), star_entry_count,
        star_entry_ptr, star_names_ptr, code);
    call ioa ("^i segments match /*.pl1:~/",
        star_entry_count);
do i = 1 to star_entry_count;
    if star_entries (i).type = 1 then do;
        do j = star_entries (i).nindex to
            star_entries(i).nindex +
            star_entries(i).nnames - 1;
            call ioa ("^[^2x]^a",
                (j = star_entries (i).nindex),
                star_names (j));
        end;
    end;
end;
free star_names_ptr->star_names;
free star_entry_ptr->star_entries;
end lspl1;
```

STAR AND EQUAL CONVENTIONS

```
!lspl1  
13 segments match **.pl1:
```

```
decls.incl.pl1  
listen_decl.incl.pl1  
listen.pl1  
put_message.pl1  
set_new_command.pl1  
command_interceptor.pl1  
process_overseer.pl1  
user_real_init_admin.pl1  
release.pl1  
  rl.pl1  
get_to_cl.pl1  
cookie.pl1  
bound_prog.pl1  
lspl1.pl1  
  list_pl1.pl1
```

- OTHER hcs \$star\_ RELATED ENTRY POINTS RETURN ADDITIONAL INFORMATION ABOUT ENTRIES

▮ hcs \$star\_dir\_list\_

▮ hcs \$star\_list\_

▮ THESE RETURN INFORMATION SUCH AS WHEN LAST MODIFIED, WHEN LAST USED, MODE, RAW MODE, RECORD LENGTH, BIT COUNT, ETC.

## STAR AND EQUAL CONVENTIONS

### ● get\_equal\_name\_

┆ call get\_equal\_name\_ (entryname, equal\_name, target\_name, code);

┆ CONSTRUCTS A TARGET NAME FROM AN ENTRYNAME AND AN EQUALNAME

┆ EXAMPLE

<u>ENTRYNAME</u>	<u>EQUAL NAME</u>	<u>TARGET NAME</u>
a.b.c	new.==	new.b.c
abc.def.ghi	=%.%.5	abc.de.5

### ● check\_star\_name\_\$path

┆ call check\_star\_name\_\$path (path, code);

┆ CHECKS THE ENTRYNAME PORTION OF A PATHNAME TO SEE IF IT HAS BEEN FORMED ACCORDING TO THE RULES FOR CONSTRUCTING STAR NAMES

┆ RETURNED CODES:

0 - ENTRYNAME VALID BUT ISN'T A STAR NAME

1 - ENTRYNAME VALID AND IS A STAR NAME

2 - ENTRYNAME IS \*\*, \*.\*\*, OR \*.\*.\*

error\_table\_\$badstar

┆ USED, FOR EXAMPLE, BEFORE CALLING hcs\_\$star\_

STAR AND EQUAL CONVENTIONS

● check\_star\_name\_\$entry

⌋ call check\_star\_name\_\$entry (entryname, code);

⌋ SAME AS check\_star\_name\_\$path, HOWEVER, ONLY REQUIRES AN ENTRYNAME AS INPUT

● match\_star\_name\_

⌋ call match\_star\_name\_ (entryname, star\_name, code);

⌋ INDICATES WHETHER OR NOT entryname MATCHES star\_name



## AREA MANIPULATION

### INTRODUCTION

#### ● AREAS ARE

- ▮ STORAGE REGIONS MANAGED BY THE AREA MANAGEMENT FACILITY
- ▮ OFTEN USED TO PASS INFO BACK AND FORTH BETWEEN USER PROCESSES AND THE SUPERVISOR
- ▮ OFTEN (BUT NOT ALWAYS) FOUND IN PROCESS DIRECTORY SEGMENTS, NAMED `<unique>.area.linker`

#### ● WHY USE AREAS?

- ▮ EMPTYING AN ENTIRE AREA (USING THE 'empty' BUILTIN) IS EASIER THAN USING SEVERAL `free` STATEMENTS
- ▮ CAN allocate IN PERMANENT SEGS AND HAVE AREA MANAGER DO ALL THE BOOK KEEPING FOR USER
- ▮ GIVES USEFUL OPTIONS LIKE EXTENSIBILITY, ZERO ON FREE, ETC.
- ▮ SOME SUBROUTINES REQUIRE POINTERS TO AREAS AS ARGUMENTS
- ▮ PL/1 OFFSETS ARE USABLE ONLY IN AREAS

## AREA FORMAT

- AREAS MAY BE DIVIDED INTO 4 TYPES BASED ON THE TWO FOLLOWING CRITERIA

- ▮ EXTENSIBILITY

- ▮ SOME AREAS ARE LIMITED TO THE SIZE OF A SEGMENT (NON-EXTENSIBLE AREAS)

- ▮ OTHERS CAN "GROW" INTO TEMP SEGMENTS IN THE PROCESS DIRECTORY (EXTENSIBLE AREAS)

- ▮ FREEING OF SPACE WITHIN AN AREA FOR REUSE

- ▮ SOME AREAS HAVE BLOCKS OF FREED SPACE MAINTAINED IN LINKED LISTS AVAILABLE FOR REUSE (FREEING AREAS)

- ▮ OTHERS DO NOT REUSE FREED SPACE IN THE AREA -- ALL ALLOCATIONS ARE DONE IN "VIRGIN AREA" (NO-FREEING AREAS)

- ▮ NO-FREEING AREAS ARE OBVIOUSLY HANDLED MUCH FASTER BY THE AREA MANAGER

- ALL AREAS HAVE 24 WORD HEADERS

- ▮ EXTENSIBLE AREAS HAVE AN ADDITIONAL 12 WORD BLOCK ALLOCATED IN THE AREA (CONTAINS INFORMATION NEEDED BY THE AREA MANAGER TO EXTEND THE AREA)

## AREA FORMAT

- FREEING AREAS ARE MADE UP OF:
  1. A HEADER THAT CONTAINS "THREAD HEADS" POINTING TO LINKED LISTS OF FREE BLOCKS (SPACE PREVIOUSLY USED AND THEN FREED)
  2. LINKED LISTS OF FREE BLOCKS (BLOCKS ARE PUT IN LIST BASED ON SIZE)
    - FIRST LIST 8 TO 14 WORDS
    - SECOND LIST 16 TO 30 WORDS
    - THIRD LIST 32 TO 62 WORDS
    - 
    - LAST LIST STARTS AT 2\*\*16 WORDS
  3. USED BLOCKS OF WORDS (EVEN WORD BOUNDARIES)
  4. VIRGIN SPACE
  
- EACH BLOCK STARTS WITH 2 WORDS OF MANAGEMENT INFORMATION SUCH AS SIZE AND A POINTER TO THE AREA HEADER
  
- THE NEXT WORD OF AN EMPTY BLOCK CONTAINS OFFSETS TO THE PREVIOUS AND NEXT BLOCKS IN THE LINKED LIST

## AREA FORMAT

- WHEN SPACE IS FREED THE AREA MANAGER:

- ▮ LOOKS AT THE FIRST 2 WORDS IN THE BLOCK TO DETERMINE SIZE OF THE BLOCK

- ▮ MERGES SMALLER ADJACENT BLOCKS IF POSSIBLE

- ▮ THREADS THE FREED BLOCKS ONTO THE APPROPRIATE LIST

- NOTE

- AREA MANAGER DOES NOT UPDATE BIT COUNT

AREA FORMAT

AREA MANIPULATING SUBROUTINES

● get\_system\_free\_area\_

▮ THIS FUNCTION RETURNS A POINTER TO THE BASE OF THE PROCESS DIRECTORY SEGMENT CONTAINING THE 'system free' AREA FOR THE RING IN WHICH IT IS CALLED

▮ USER MAY USE THIS AREA AS HE/SHE PLEASURES

```
dcl A area based (get system free area ( ) );
dcl get_system_free_area_entry returns (ptr);
dcl alpha based (beta);
dcl beta pointer;

    allocate alpha in (A) set (beta);

/* WARNING -- DO NOT SET "A" = empty(); */
```

AREA FORMAT  
AREA MANIPULATING SUBROUTINES

● define\_area\_

▮ call define\_area\_ (info\_ptr, code);

▮ INITIALIZES AN AREA

▮ USED TO CONTROL SPECIAL AREA MANAGEMENT FEATURES:

▮ EXTEND: ENABLES AREA TO GROW BEYOND MAX SIZE SET INTO TEMPSEGS  
(INSTEAD OF SIGNALLING THE area CONDITION)

▮ ZERO ON ALLOCATION

▮ ZERO ON FREEING

▮ IGNORE ALL free REQUESTS (FOR DEBUGGING PURPOSES)

▮ SET MAX SIZE TO SPECIFIED VALUE (0 modulo 8)

▮ USES AN INFORMATION STRUCTURE FOUND IN area\_info.incl.pl1

▮ REGION BEING INITIALIZED

▮ IS POINTED TO BY area\_info.areap

▮ IS AUTOMATICALLY ACQUIRED FROM PROCESS DIRECTORY TEMPSEG POOL  
IF area\_info.areap = null()

AREA FORMAT  
AREA MANIPULATING SUBROUTINES

● release\_area\_

⌋ call release\_area\_ (area\_ptr);

⌋ CLEANS UP AN AREA AFTER IT IS NO LONGER NEEDED

⌋ RETURNS ANY TEMPSEGS TO THE POOL

● area\_info\_

⌋ call area\_info\_ (info\_ptr, code);

⌋ FILLS IN THE USER-ALLOCATED area\_info STRUCTURE (CALLER MUST SET area\_info.areap)

AREA FORMAT

AREA MANIPULATING SUBROUTINES

```
dcl  area_infop ptr;

dcl  1 area_info aligned based (area_infop),
     2 version fixed bin,
     2 control aligned like area_control,
     2 owner char (32) unal,
     2 n_components fixed bin,
     2 size fixed bin (18),
     2 version_of_area fixed bin,
     2 areap ptr,
     2 allocated_blocks fixed bin,
     2 free_blocks fixed bin,
     2 allocated_words fixed bin (30),
     2 free_words fixed bin (30);

dcl  1 area_control aligned based,
     2 extend bit (1) unal,
     2 zero_on_alloc bit (1) unal,
     2 zero_on_free bit (1) unal,
     2 dont_free bit (1) unal,
     2 no_freeing bit (1) unal,
     2 system bit (1) unal,
     2 pad bit (30) unal;
```



AREA FORMAT

AREA RELATED COMMANDS

● AREA-RELATED COMMANDS (DOCUMENTED IN SWG)

|| create\_area

|| PERFORMS define area\_'S TASKS, GIVEN A VIRTUAL POINTER TO AN AREA TO BE CREATED

|| set\_system\_storage, set\_user\_storage

|| ENABLE A USER-CREATED AREA TO BE USED INSTEAD OF DEFAULT 'system free' OR 'user free' AREAS

|| AREA SPECIFIED

|| MUST BE ZERO\_ON\_FREE OR ZERO\_ON\_ALLOC

|| SHOULD BE EXTENSIBLE

|| USEFUL FOR ISOLATING BUGS WHICH ARE INADVERTENTLY DESTROYING INFORMATION IN EITHER 'system free area' OR 'user free area'

|| area\_status

|| COMMAND INTERFACE TO area\_info\_

|| YOU ARE NOW READY FOR WORKSHOP #1 ||

TOPIC IV  
Multics Security

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## INTRODUCTION

- Multics HAS THREE ACCESS CONTROL MECHANISMS

- ┆ THE ACCESS CONTROL LIST MECHANISM (ACL's)

- ┆ THE ACCESS ISOLATION MECHANISM (AIM)

SEE APPENDIX A

- ┆ THE RING MECHANISM

INITIAL ACL'S

- THERE MAY BE, ASSOCIATED WITH EVERY DIRECTORY, TWO TYPES OF INITIAL ACL CONTROL LISTS (ONE FOR INFERIOR SEGMENTS, ONE FOR INFERIOR DIRECTORIES)
- FOR EVERY SUBROUTINE THAT MANIPULATES ACL'S THERE IS A CORRESPONDING SUBROUTINE THAT MANIPULATES INACL'S

ACL ENTRY POINTS	INACL ENTRY POINTS
hcs_\$add_acl_entries	hcs_\$add_inacl_entries
hcs_\$add_dir_acl_entries	hcs_\$add_dir_inacl_entries
hcs_\$delete_acl_entries	hcs_\$delete_inacl_entries
hcs_\$delete_dir_acl_entries	hcs_\$delete_dir_inacl_entries
hcs_\$list_acl	hcs_\$list_inacl
hcs_\$list_dir_acl	hcs_\$list_dir_inacl
hcs_\$replace_acl	hcs_\$replace_inacl
hcs_\$replace_dir_acl	hcs_\$replace_dir_inacl

INITIAL ACL'S

● ALL INACL ENTRY POINTS REQUIRE SPECIFICATION OF A RING NUMBER

● SEGMENT INACL APPLIES TO MSF'S

● SEE THE COMMANDS

sis sid

dis did

lis lid

ALL OF WHICH ACCEPT A -ring CONTROL ARGUMENT

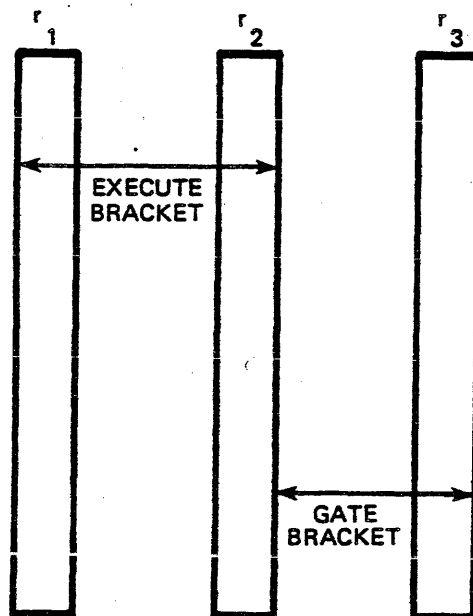
RINGS  
INTRODUCTION

- INTRAPROCESS ACCESS IS CONTROLLED BY THE RING MECHANISM
  
- TYPICAL APPLICATIONS
  - ▮ PROTECTION OF SUPERVISOR FROM USER PROGRAMS
  
  - ▮ PROTECTION OF SUBSYSTEM DATA BASE FROM DIRECT ACCESS
  
- CLARIFICATION OF MISCONCEPTION
  - ▮ ALL SEGMENTS ARE NOT "IN" JUST ONE RING
  
  - ▮ SEGMENTS MAY "SPAN" SEVERAL RINGS

RINGS  
RING BRACKETS

- EACH SEGMENT HAS ASSOCIATED WITH IT 3 RING BRACKET NUMBERS (THESE NUMBERS DETERMINE THE RING BRACKETS FOR THAT SEGMENT)
  
- RING BRACKETS DEFINE IN WHICH RING A USER CAN READ, WRITE, CALL OR EXECUTE A SEGMENT
  
- ▮ THE FIRST TWO NUMBERS DELIMIT THE EXECUTE BRACKET
  
- ▮ THE SECOND AND THIRD NUMBERS DELIMIT THE GATE BRACKET

**EXECUTE AND GATE BRACKETS**



RINGS  
RING BRACKETS

- RING BRACKET NUMBERS ARE EXPRESSED r1,r2,r3

ASSUMING A RING 4 USER, SPECIFY THE SEGMENTS FOR WHICH THE USER IS IN THE EXECUTE AND/OR GATE BRACKET

<u>r1</u>	<u>r2</u>	<u>r3</u>	<u>EXECUTE</u>	<u>GATE</u>
4	4	4	_____	_____
0	5	5	_____	_____
0	0	5	_____	_____
5	5	5	_____	_____
1	1	1	_____	_____
3	3	4	_____	_____

- TYPICAL RING BRACKETS

▮ USER SEGMENTS 4,4,4

▮ SYSTEM COMMANDS AND SUBROUTINES 1,5,5 OR 0,5,5

▮ SYSTEM GATES 1,1,5 OR 0,0,5 (hes\_)

- NOTICE THAT TO EXECUTE hes\_ THE USER MUST BE IN RING 0

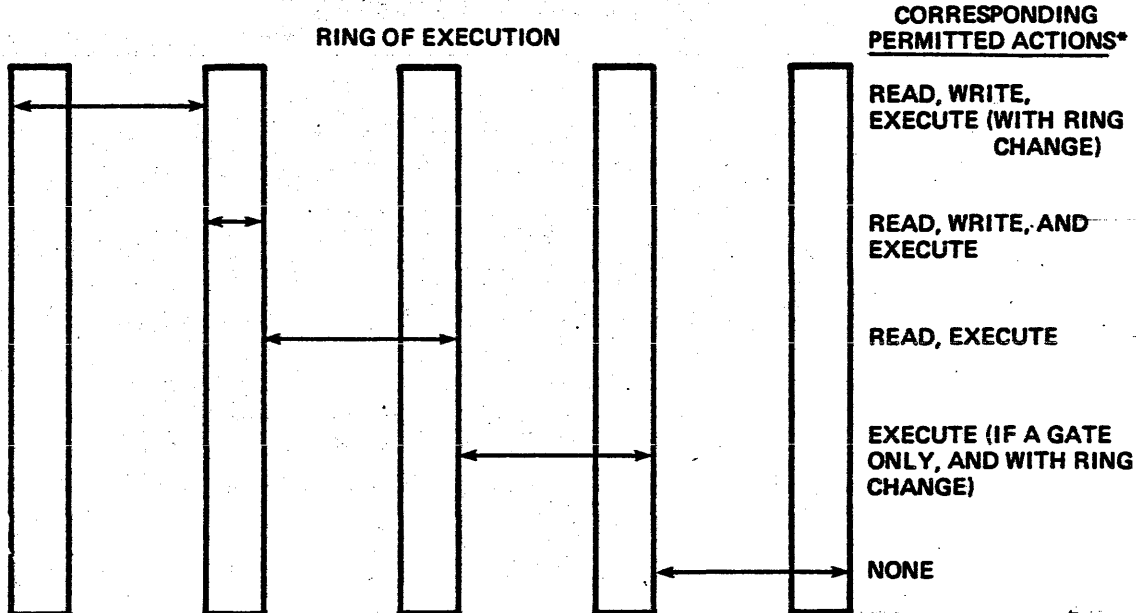
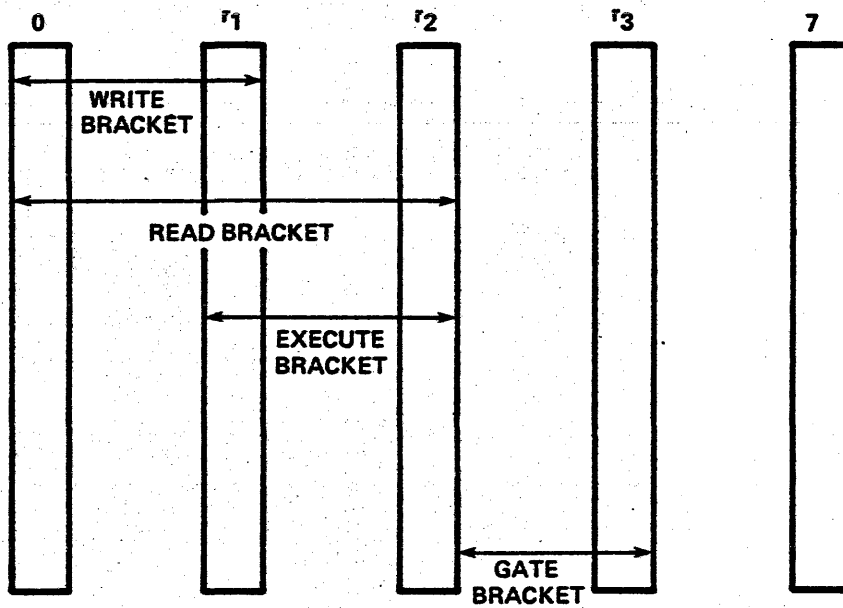
▮ HOWEVER, USERS ARE USUALLY THOUGHT OF AS BEING IN RING 4

▮ USERS ACTUALLY TRAVEL IN AND OUT OF THE RING STRUCTURE



RINGS  
RING BRACKETS

## RING MECHANISM SUMMARY



\* SUBJECT, OF COURSE, TO ACL AND AIM

## RINGS

### RING BRACKET SUBROUTINES

- `get_ring_`

- ▮ `current_ring = get_ring_();`

- ▮ RETURNS THE USER'S CURRENT RING OF EXECUTION

- `hcs_$get_ring_brackets`

- ▮ call `hcs_$get_ring_brackets (dir_name, entryname, rb, code);`

- ▮ RETURNS (r1,r2,r3) FOR A SPECIFIED SEGMENT

- `hcs_$get_dir_ring_brackets`

- ▮ call `hcs_$get_dir_ring_brackets (dir_name, entryname, drb, code);`

- ▮ RETURNS (r1,r2) FOR A SPECIFIED DIRECTORY

## RINGS

### RING BRACKET SUBROUTINES

- `hcs_$set_ring_brackets` and `hcs_$set_dir_ring_brackets`

- ▮ SAME CALL ARGUMENTS AS CORRESPONDING ENTRY POINT ON PREVIOUS PAGE

- ▮ SETS THE RING BRACKETS OF A SPECIFIED SEGMENT OR DIRECTORY

- ▮ THE RING BRACKETS MUST BE  $\geq$  THE CURRENT VALIDATION LEVEL OF THE CALLING PROCESS

- ▮ SEE ALSO THE `set ring brackets (srb)`, `set dir ring brackets (sdrb)`, `lset ring brackets (lsrb)` AND `lset_dir_ring_brackets (lsdrb)` COMMANDS

RINGS

GATES

- DEFINITION OF A GATE: ONLY POINT AT WHICH A PROCEDURE IN AN OUTER RING CAN TRANSFER TO A PROCEDURE IN AN INNER RING

- ▮ IDENTIFIED BY PRESENCE OF GATE BRACKET ( $r_2 < r_3$ )

- ▮ CHANGES USER'S RING OF EXECUTION

- GATES ARE "CREATED" BY:

- ▮ USING `alm` MACROS

- ▮ AFTER COMPILATION:

- ▮ THE RING BRACKETS ARE SET TO THAT OF A GATE

- ▮ THE ENTRY BOUND IS SET (DISCUSSED BELOW)

RINGS

GATES

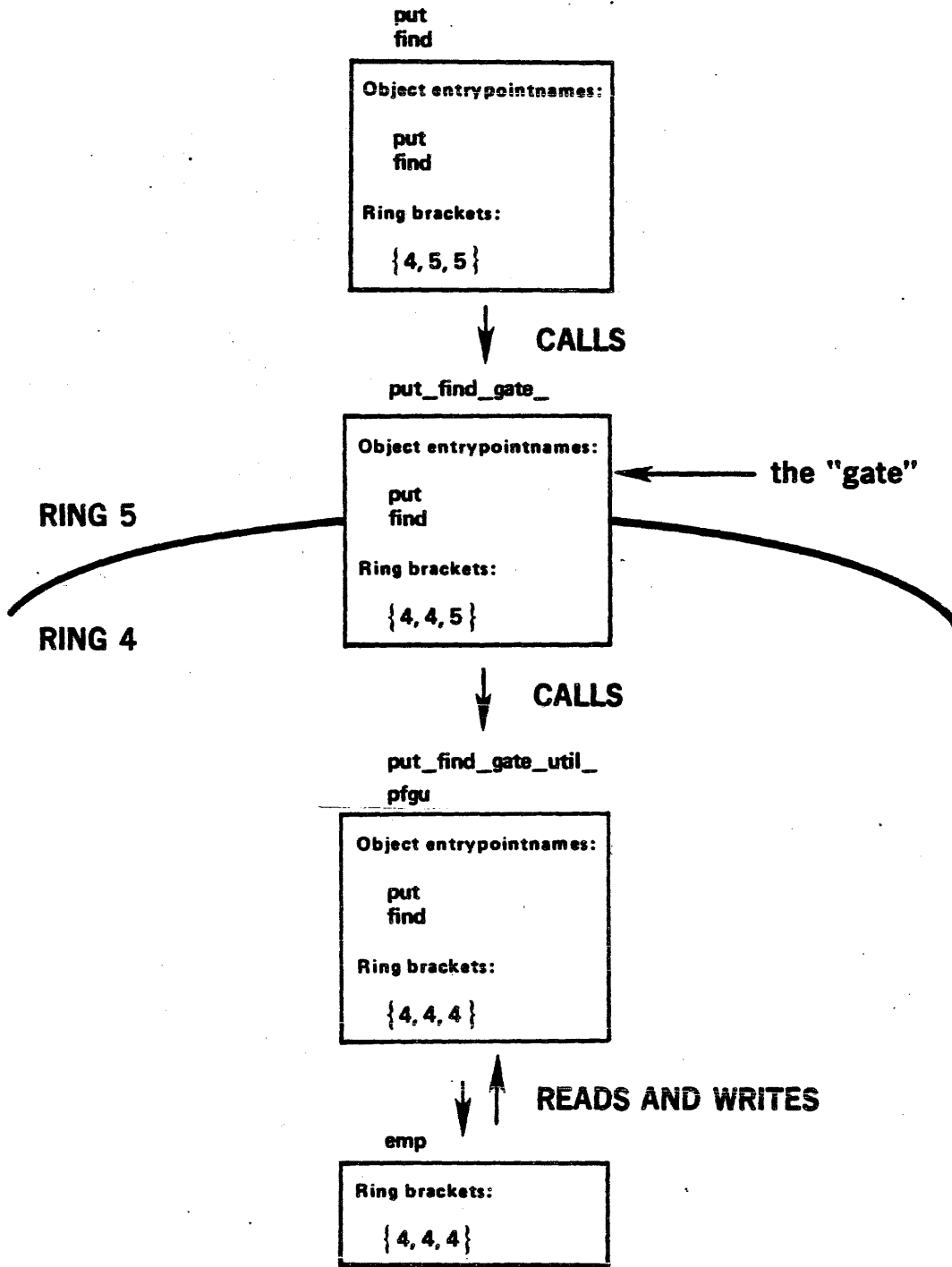
● WHY GATES SHOULD BE WRITTEN IN alm

- ▮ PREVIOUSLY IT WAS POSSIBLE TO ARTIFICIALLY JUMP INTO INNER CODE OF A GATE (POTENTIAL BREACH OF SECURITY)
  
- ▮ alm ENABLES CAREFUL CONTROL OF OBJECT SEGMENT FORMAT (ONLY TRANSFER INSTRUCTIONS ARE PLACED AT BASE OF SEGMENT)
  
- ▮ THE WRITER OF THE GATE SETS THE ENTRY POINT BOUND OF THE PROGRAM EQUAL TO THE END OF THE TRANSFER INSTRUCTIONS

RINGS

GATES

### CROSSING RING BOUNDARIES



RINGS

GATES

- USE hcs \$set\_entry\_bound OR hcs \$set\_entry\_bound\_seg TO SET ENTRY POINT BOUND

▮ call hcs \$set\_entry\_bound (dir\_name, entryname, entry\_bound, code);

▮ call hcs \$set\_entry\_bound\_seg (seg\_ptr, entry\_bound, code);

▮ SETS A HARDWARE ENFORCED LIMIT ON ENTRY POINT OFFSET

▮ IF entry\_bound IS 0 THE MECHANISM IS DISABLED

▮ ENTRY BOUND MAINTAINED IN CONTAINING DIRECTORY AND BUILT INTO THE SEGMENT DESCRIPTOR WORD (SDW) WHEN THE SEGMENT IS MADE KNOWN

▮ OBJECT SEGMENT ITSELF IS UNCHANGED

RINGS

GATES

EXAMPLE

```
!print bound_prog.pl1 1

bound_prog: proc;
dcl hcs $set_entry_bound entry (char (*), char (*),
    fixed bin (14), fixed bin (35));
dcl code fixed bin (35);

    call hcs $set_entry_bound (>udd>F15D>doodle",
        "bound_prog", 10, code);
end bound_prog;

r 14:02 0.066 2

!bound_prog
r 14:02 0.081 3

!bound_prog

Error: Attempt by cu_1373
(>system_library 1>bound_command_loop )
to access >udd>F15D>doodle>bound_prog16
which is beyond the entry bound for the gate.
r 14:02 0.164 23 level 2
```



RINGS  
VALIDATION LEVEL

● POTENTIAL PROBLEM:

- ▮ RING OF EXECUTION KEPT IN REGISTER IN THE PROCESSOR
- ▮ RING OF EXECUTION KEEPS CHANGING
- ▮ ASSUME A SEGMENT IS BEING CREATED
- ▮ HOW DO SYSTEM SUBROUTINES ASSIGN PROPER RING BRACKETS?

● VALIDATION LEVEL

- ▮ MEANS BY WHICH INNER RING (CALLED) PROCEDURE "KNOWS" THE LEVEL OF PRIVILEGE OF THE OUTER RING (CALLING) PROCEDURE
- ▮ VALIDATION LEVEL CAN BE CHANGED
- ▮ CANNOT BE SET LOWER THAN RING OF EXECUTION
- ▮ VALIDATION LEVEL CHANGE USED FOR EXAMPLE:
  - ▮ TO CREATE A MAILBOX
  - ▮ BY A SUBSYSTEM WISHING TO CREATE A SEGMENT IN INNER RING

RINGS  
VALIDATION LEVEL

● `cu_$level_get` (AG93)

⌋ `call cu_$level_get (level)`

⌋ RETURNS THE CURRENT VALIDATION LEVEL

⌋ PRIMARILY USED PRIOR TO A CALL TO `cu_$level_set` TO SAVE THE CURRENT VALIDATION LEVEL

● `cu_$level_set` (AG93)

⌋ `call cu_$level_set (level)`

⌋ ALLOWS THE CALLER TO CHANGE THE CURRENT VALIDATION LEVEL

⌋ NEW LEVEL MUST BE  $\geq$  CURRENT RING OF EXECUTION

● `hcs_$get_user_effmode`

⌋ `call hcs_$get_user_effmode (dir_name, entryname, user_id,  
ring, mode, code);`

⌋ RETURNS THE EFFECTIVE MODE FOR THE SPECIFIED RING

RINGS

CROSS RING I/O

- AN ATTEMPT TO DO "CROSS RING I/O" USUALLY RESULTS IN A FATAL PROCESS ERROR

┆ REASON: IOCB'S ARE PER RING

┆ TYPICAL EXAMPLE: CALLING `com_err_` IN AN INNER RING

- "CROSS RING I/O" IS ALLOWED USING THE FOLLOWING

┆ `cross_ring_`

┆ AN I/O MODULE WHICH ALLOWS AN OUTER RING TO ATTACH A SWITCH (BASICALLY AS A SYNONYM) TO A PREEXISTING SWITCH IN AN INNER RING, AND TO PERFORM I/O OPERATIONS BY FORWARDING I/O FROM THE ATTACHMENT IN THE OUTER RING THROUGH A GATE TO THE INNER RING

┆ AN INNER RING SWITCH MUST BE ATTACHED WHILE IN THE INNER RING BEFORE `cross_ring_` CAN BE USED TO ATTACH OUTER RING SWITCH

┆ `cross_ring_io_$allow_cross`

┆ call `cross_ring_io_$allow_cross (switch_name, ring, code);`

┆ CALL MUST BE MADE IN THE INNER RING BEFORE THE OUTER RING ATTEMPTS TO ATTACH TO THIS SWITCH WITH `cross_ring_`

RINGS

CROSS RING I/O

```
!pat
  user_i/o          tty_ -login_channel
                    stream_input_output
  user_input        syn_ user_i/o
  user_output       syn_ user_i/o
  error_output      syn_ user_i/o
```

```
!pr cross.pl1 1
```

```
cross:  proc;
  dcl gate$allow entry;
  dcl iox_$get_chars entry (ptr, ptr, fixed bin(21),
                           fixed bin(21), fixed bin(35));
  dcl iox_$attach_name entry (char(*), ptr, char(*), ptr,
                              fixed bin(35));
  dcl iox_$open entry (ptr, fixed bin, bit (1) aligned,
                      fixed bin (35));
  dcl com_err_entry() options(variable);
  dcl code fixed bin (35);
  dcl iocb ptr;
  dcl buffer char (20);

  call gate$allow;

  call iox_$attach_name ("outer", iocb, "cross_ring_file 4",
                        null(), code);

  call iox_$open (iocb, 3, "0"b, code);

end cross;
```

```
!pr gate.alm 1
```

```
include gate_macros
gate_info
gate allow,allow,allow,0
end
```

RINGS  
CROSS RING I/O

!pr allow.pl1 1

```
allow: proc;
  dcl cross_ring_io_$allow_cross entry (char(*), fixed bin,
                                         fixed bin(35));
  dcl iox_$attach_name entry (char(*), ptr, char(*), ptr,
                              fixed bin(35));
  dcl iox_$open entry (ptr, fixed bin, bit(1) aligned, fixed bin(35));
  dcl code fixed bin(35);
  dcl com_err entry() options(variable);
  dcl iocb ptr;
  dcl null builtin;
  dcl cu_$level_set entry (fixed bin);
  dcl cu_$level_get entry (fixed bin);
  dcl old_level fixed bin;
  dcl get_ring_entry() returns(fixed bin(3));

  call cu_$level_get (old_level);

  call cu_$level_set (get_ring ());

  call iox_$attach_name ("file", iocb, "vfile_ >udd>MED>nd>gate>file",
                        null(), code);

  call cross_ring_io_$allow_cross ("file", 5, code);

  call cu_$level_set (old_level);

end allow;
```

```
!st [wd] -rb
  5, 5
```

```
!st file cross gate allow -rb
  >udd>MED>NDibble>gate>file
  4, 4, 4
  >udd>MED>NDibble>gate>cross
  4, 5, 5
  >udd>MED>NDibble>gate>gate
  4, 4, 5
  >udd>MED>NDibble>gate>allow
  4, 4, 4
```

RINGS

CROSS RING I/O

!cross

!pat

```
user_i/o          tty_ -login_channel
                  stream_input_output
user_input        syn_ user_i/o
user_output       syn_ user_i/o
error_output      syn_ user_i/o
outer             cross_ring_ file 4 stream_input_output
```

```
!io put_chars outer "line 1"
!io put_chars outer "line 2"
!io position outer -1
!io get_line outer
  io_call:7 characters returned.line 1
```

!pr file

```
print: Incorrect access on entry. >udd>MED>NDibble>gate>file
```

● TWO MAJOR POINTS TO REMEMBER

▮ WORKING DIRECTORIES ARE PER RING

▮ MUST SET VALIDATION LEVEL TO INNER RING BEFORE CREATING INNER  
'IOCB'

TOPIC V

The Command Environment

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Modifying the Standard Command Environment . . . . .	5-2
Current Ready Procedure. . . . .	5-6
Current Command Processor. . . . .	5-8
Command Level Intermediary . . . . .	5-10
Some Miscellaneous cu_ Entry Points. . . . .	5-12
An Example . . . . .	5-15

## INTRODUCTION

- THE SUBSYSTEM DESIGNER HAS THE CAPABILITY OF MODIFYING SEVERAL DIFFERENT ASPECTS OF THE COMMAND ENVIRONMENT
- THE `cu` (COMMAND UTILITY) SUBROUTINE (1,5,5) PROCEDURE WRITTEN IN `alm`) IS THE TOOL USED BY SUBSYSTEM DESIGNERS TO ACCOMPLISH THE FOLLOWING BASIC TASKS:

### ▮ WRITING COMMAND OR ACTIVE FUNCTION PROCEDURES

```
cu_$arg_count  
cu_$arg_ptr  
cu_$af_arg_count  
cu_$af_return_arg  
cu_$af_arg_ptr
```

### ▮ MODIFYING THE STANDARD COMMAND ENVIRONMENT

### ▮ WRITING A COMMAND PROCESSOR

```
cu_$generate_call
```



## MODIFYING THE STANDARD COMMAND ENVIRONMENT

- THE STANDARD COMMAND ENVIRONMENT IS PROVIDED TO ALLOW THE USER TO PROCESS HIS COMMAND REQUESTS, EXECUTE HIS PROGRAMS, AND SO ON
  
- THE BASIC COMMAND ENVIRONMENT HAS THE FOLLOWING CHARACTERISTICS:
  - ▮ 'listen' IS INVOKED AT PROCESS START UP TIME AND IS ALWAYS RETURNED TO FOLLOWING THE EXECUTION OF A COMMAND (OR FOLLOWING A PROGRAM ABORT)
  
  - ▮ THIS "LISTENER" ACCEPTS INPUT FROM THE USER'S TERMINAL AND PASSES SUCH INPUT TO THE "CURRENT COMMAND PROCESSOR" FOR FURTHER PROCESSING
  
  - ▮ EVERY TIME CONTROL RETURNS BACK TO THE "LISTENER", THE 'listen\_' PROGRAM INVOKES THE "CURRENT READY PROCEDURE"
  
  - ▮ THE "CURRENT COMMAND PROCESSOR" PROCESSES THE INPUT LINE TYPED BY THE USER AND PASSED TO IT BY THE LISTENER
    - ▮ THE STANDARD COMMAND PROCESSOR (command processor ) FIRST DOES SUCH THINGS AS EXPANDING OUT ANY ITERATION LOOPS (PARENTHESES) AND EVALUATING ACTIVE FUNCTIONS (BRACKETS)
  
    - ▮ THE STANDARD COMMAND PROCESSOR THEN DEVELOPS A CALL TO THE APPROPRIATE COMMAND OR USER PROGRAM
  
  - ▮ WHENEVER A QUIT OR OTHER "UNCLAIMED" SIGNAL (CONDITION NOT HANDLED BY USER) ARISES A "NEW LEVEL" OF THE LISTENER IS INVOKED BY "REENTERING COMMAND LEVEL"
  
  - ▮ WHEN THE DEFAULT HANDLER DECIDES TO REENTER COMMAND LEVEL, THE "CURRENT COMMAND LEVEL INTERMEDIARY" WILL BE INVOKED

● STANDARD COMMAND ENVIRONMENT CONCEPTS

▮ CURRENT READY PROCEDURE

- ▮ IS INVOKED WHEN LISTENER (OR ANY OTHER PROCEDURE) CALLS `cu_$ready_proc`
- ▮ IS, BY DEFAULT, `print_ready_message_`, WHICH MERELY PRINTS THE READY MESSAGE
- ▮ MAY BE SET BY A CALL TO `cu_$set_ready_procedure`
- ▮ MAY BE DETERMINED BY A CALL TO `cu_$get_ready_procedure`

▮ CURRENT COMMAND PROCESSOR

- ▮ IS INVOKED WHEN LISTENER (OR ANY OTHER PROCEDURE) CALLS `cu_$cp`
- ▮ IS, BY DEFAULT, `command_processor_`
- ▮ MAY BE SET BY A CALL TO `cu_$set_command_processor`. (NOTE THE 'abbrev' COMMAND)
- ▮ MAY BE DETERMINED BY A CALL TO `cu_$get_command_processor`

MODIFYING THE STANDARD COMMAND ENVIRONMENT

▮ CURRENT COMMAND LEVEL INTERMEDIARY

▮ IS INVOKED WHEN DEFAULT ERROR HANDLER (OR ANY OTHER PROCEDURE) CALLS `cu_$cl`

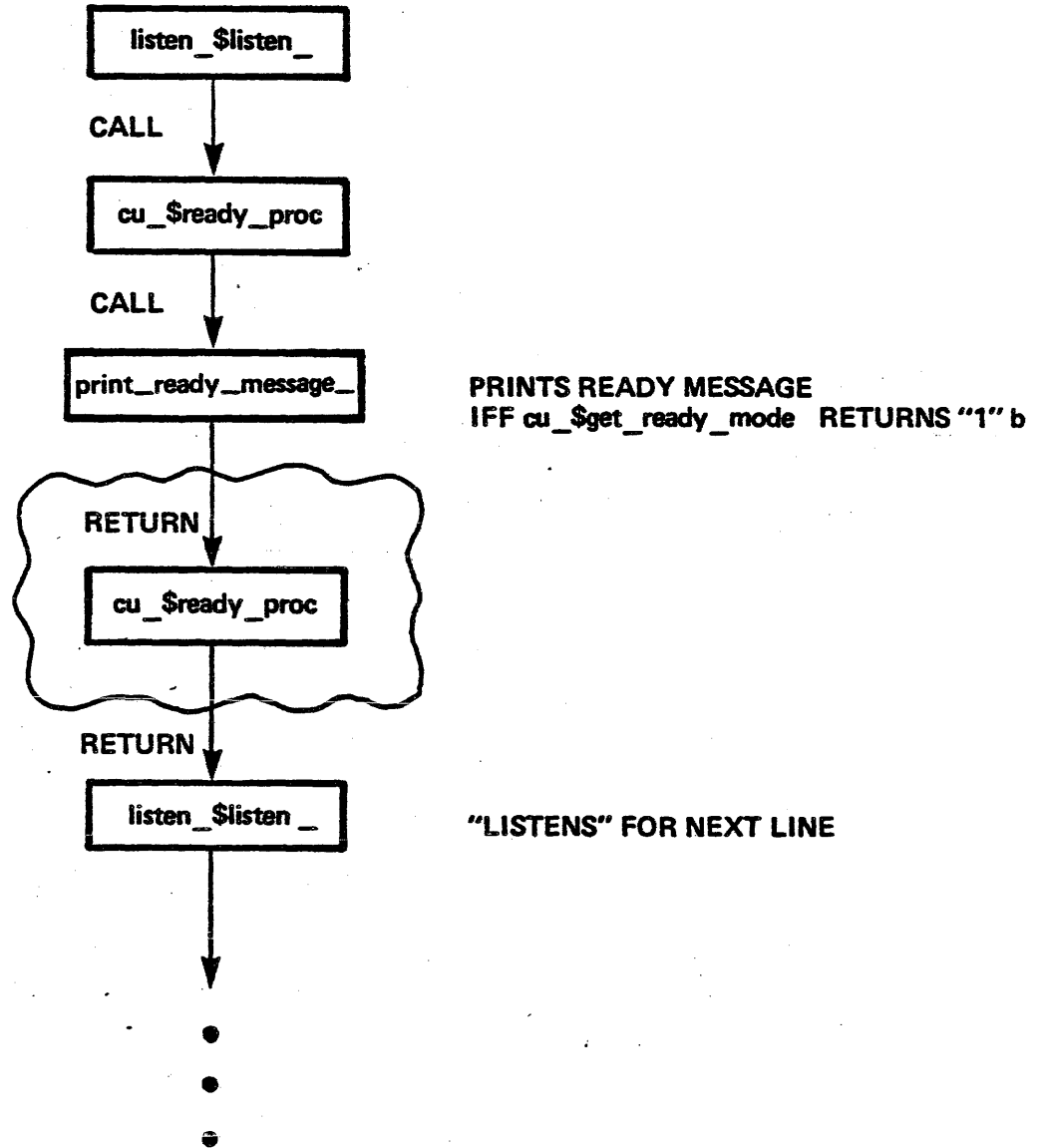
▮ IS, BY DEFAULT FOR INTERACTIVE PROCESSES, `get to cl $unclaimed` signal WHICH REENTERS COMMAND LEVEL VIA A CALL TO `listen_$release_stack`

▮ MAY BE SET BY A CALL TO `cu_$set_cl_intermediary`

▮ MAY BE DETERMINED BY A CALL TO `cu_$get_cl_intermediary`

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## WHERE THE CURRENT READY PROCEDURE FITS IN



MODIFYING THE STANDARD COMMAND ENVIRONMENT

CURRENT READY PROCEDURE

● MANIPULATING THE READY PROCEDURE

▮ cu\_\$ready\_proc (AG93)

▮ CALLS THE CURRENT READY PROCEDURE

▮ cu\_\$set\_ready\_procedure (AG93)

▮ ESTABLISHES THE SPECIFIED PROCEDURE AS THE CURRENT READY PROCEDURE

▮ cu\_\$get\_ready\_procedure (AG93)

▮ RETURNS A NULL ENTRY VALUE IF THE CURRENT READY PROCEDURE IS THE DEFAULT (print\_ready\_message\_)

▮ OTHERWISE, RETURNS THE PL/1 ENTRY VALUE OF THE CURRENT READY PROCEDURE

▮ cu\_\$set\_ready\_mode (AG93)

▮ SETS OR RESETS THE "STATIC READY MODE" SWITCH

▮ THE CURRENT READY PROCEDURE CAN (BUT NEEDN'T) CHECK THIS SWITCH TO SEE WHETHER TO PRINT OR NOT (NOTE THE 'ready\_on', AND 'ready\_off' COMMANDS)

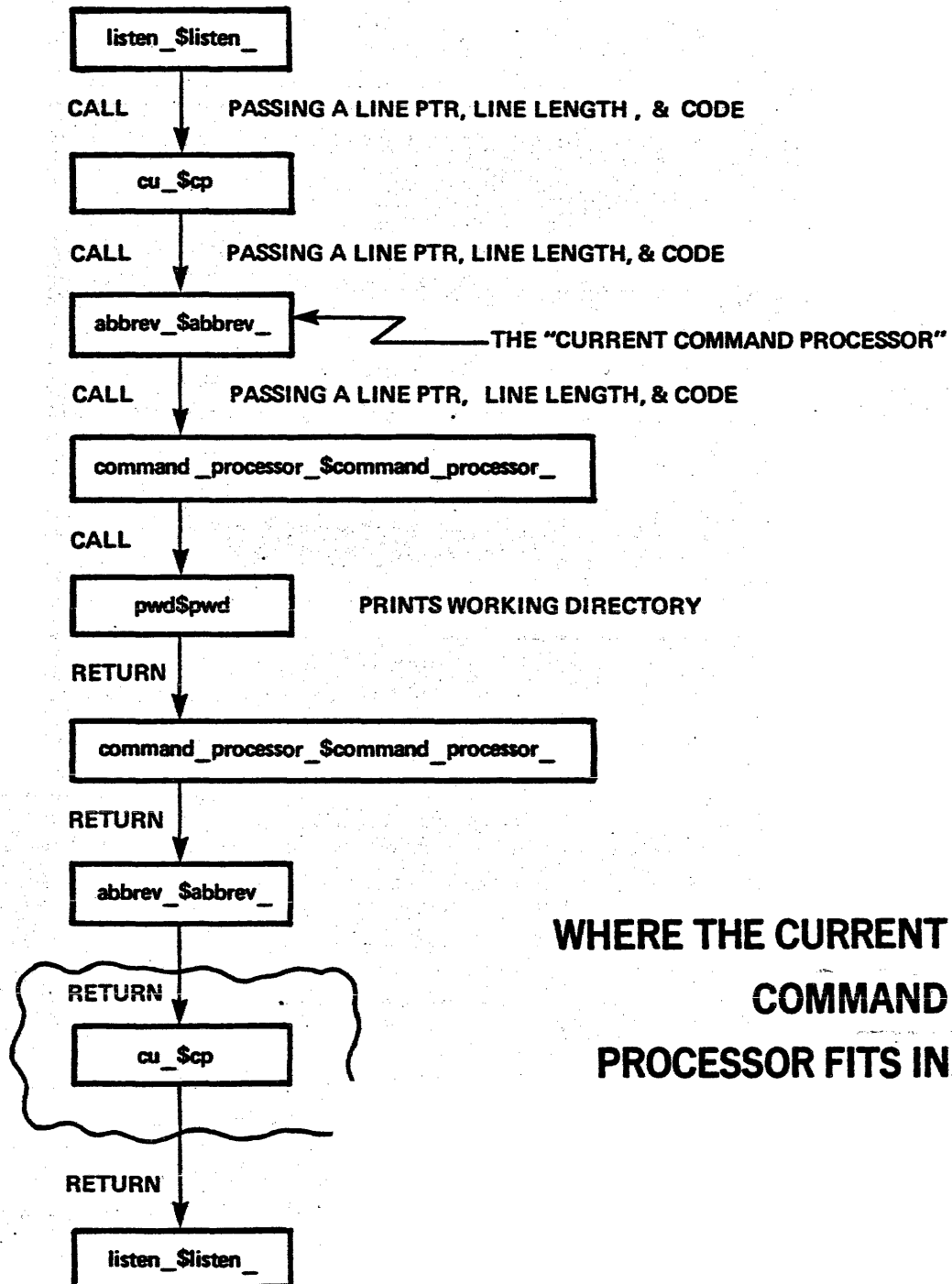
▮ cu\_\$get\_ready\_mode (AG93)

▮ RETURNS THE VALUE OF THE "STATIC READY MODE" SWITCH

MODIFYING THE STANDARD COMMAND ENVIRONMENT

CURRENT COMMAND PROCESSOR

USER TYPES "pwd"



WHERE THE CURRENT  
COMMAND  
PROCESSOR FITS IN

MODIFYING THE STANDARD COMMAND ENVIRONMENT

CURRENT COMMAND PROCESSOR

● MANIPULATING THE CURRENT COMMAND PROCESSOR

▮ cu\_\$cp (AG93)

▮ INVOKES THE CURRENT COMMAND PROCESSOR, PASSING TO IT AN INPUT LINE POINTER, LINE LENGTH AND CODE

▮ BESIDES THE LISTENER, USED ALSO BY SUBSYSTEMS HONORING AN "e" OR ".." REQUEST

▮ cu\_\$set\_command\_processor (AG93)

▮ ESTABLISHES THE SPECIFIED PROCEDURE AS THE CURRENT COMMAND PROCESSOR (BY SPECIFYING AN ENTRY VALUE)

▮ cu\_\$get\_command\_processor (AG93)

▮ RETURNS A NULL ENTRY VALUE IF THE CURRENT COMMAND PROCESSOR IS THE DEFAULT (command\_processor\_\$command\_processor\_)

▮ OTHERWISE RETURNS THE ENTRY VALUE OF THE CURRENT COMMAND PROCESSOR

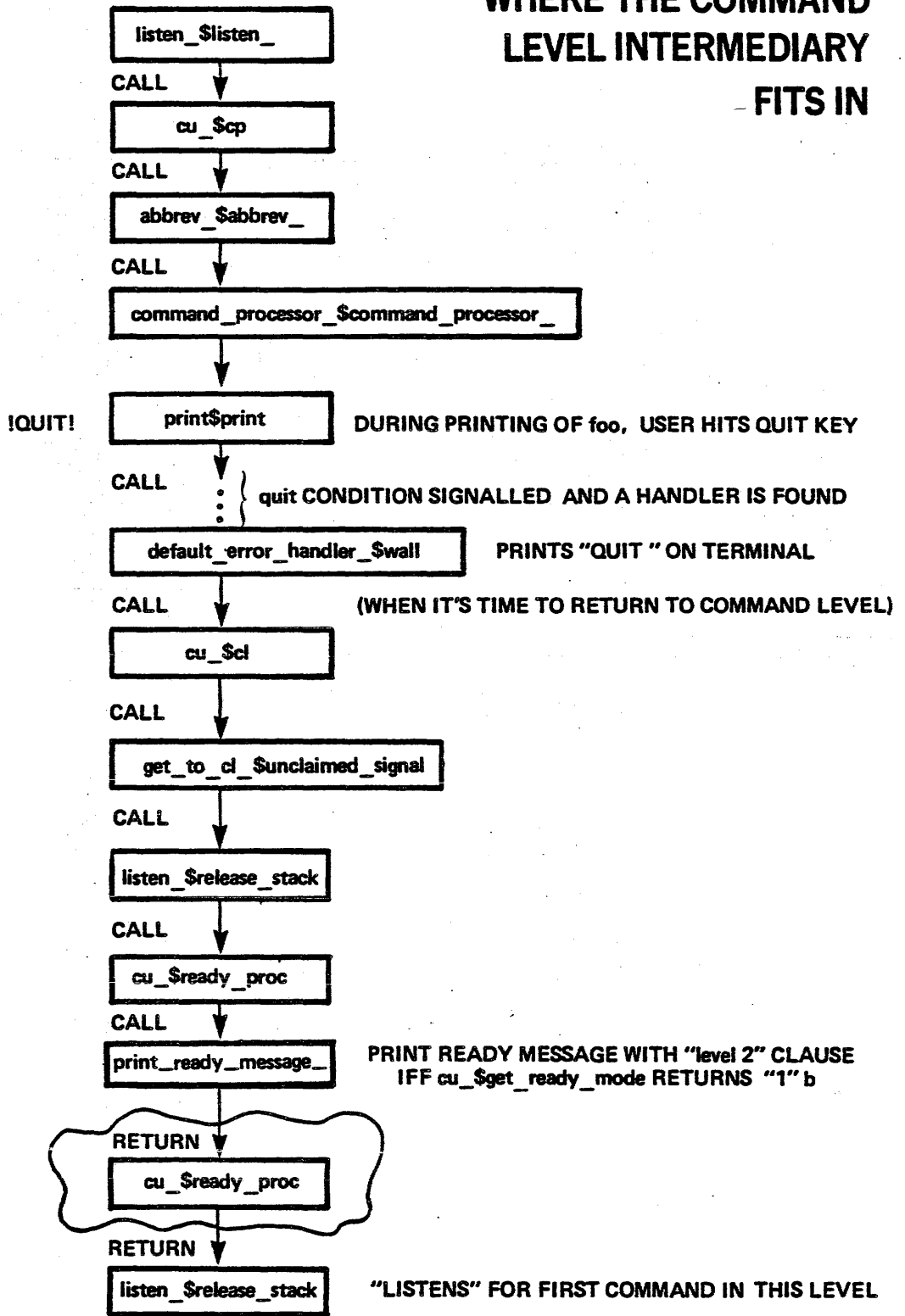


MODIFYING THE STANDARD COMMAND ENVIRONMENT

COMMAND LEVEL INTERMEDIARY

USER TYPES "print foo"

WHERE THE COMMAND  
LEVEL INTERMEDIARY  
FITS IN



MODIFYING THE STANDARD COMMAND ENVIRONMENT

COMMAND LEVEL INTERMEDIARY

● MANIPULATING THE COMMAND LEVEL INTERMEDIARY

┆ cu\_\$cl (AG93)

┆ INVOKES THE CURRENT COMMAND LEVEL INTERMEDIARY

┆ CALLED BY THE STANDARD ERROR HANDLERS

┆ cu\_\$set\_cl\_intermediary (AG93)

┆ ESTABLISHES THE SPECIFIED PROCEDURE AS THE CURRENT COMMAND LEVEL INTERMEDIARY

┆ NOTE THAT AN INTERMEDIARY IS USED IN ABSENTEE PROCESSES TO FORCE PROCESS TERMINATION "WHEN AN ATTEMPT IS MADE TO REENTER COMMAND LEVEL"

┆ cu\_\$get\_cl\_intermediary (AG93)

┆ RETURNS A NULL ENTRY VALUE IF THE CURRENT COMMAND LEVEL INTERMEDIARY IS THE DEFAULT (get\_to\_cl\_\$unclaimed\_signal)

┆ OTHERWISE RETURNS THE ENTRY VALUE OF THE CURRENT COMMAND LEVEL INTERMEDIARY

MODIFYING THE STANDARD COMMAND ENVIRONMENT

SOME MISCELLANEOUS CU ENTRY POINTS

● cu\_\$decode\_entry\_value (AG93)

┆ call cu\_\$decode\_entry\_value (entry\_value, ep\_ptr, env\_ptr);

┆ EXTRACTS THE POINTER COMPONENTS OF A PL/I ENTRY VALUE

┆ USEFUL FOR DETERMINING IF AN ENTRY VALUE IS NULL

┆ NOTE: RECENTLY REPLACED BY codeptr AND environmentptr BUILTINS

● cu\_\$arg\_list\_ptr (AG93)

┆ RETURNS A POINTER TO THE ARGUMENT LIST STRUCTURE PASSED TO THE CALLER

┆ GENERALLY USED BY SUBROUTINES WHICH ARE CALLED WITH A VARYING NUMBER OF ARGUMENTS OF VARYING DATA TYPES (ioa FOR INSTANCE), TO ALLOW EXAMINATION OF THE ARGUMENT LIST DIRECTLY

┆ SEE ALSO decode\_descriptor\_ IN AK92

MODIFYING THE STANDARD COMMAND ENVIRONMENT

SOME MISCELLANEOUS CU ENTRY POINTS

● cu\_\$arg\_ptr\_rel (AG93)

▮ REMINISCENT OF cu\_\$arg\_ptr

▮ ALLOWS A PROCEDURE TO REFERENCE THE nth ARGUMENT PASSED TO ANOTHER PROCEDURE, GIVEN A POINTER TO THAT OTHER PROCEDURE'S ARGUMENT LIST

▮ QUESTION: HOW WOULD A PROCEDURE OBTAIN THE arglist\_ptr OF ANOTHER PROCEDURE?

▮ IT COULD BE PASSED SUCH A POINTER

▮ IT COULD LOOK IN STACK FRAME OF OTHER PROCEDURE

MODIFYING THE STANDARD COMMAND ENVIRONMENT

SOME MISCELLANEOUS CU ENTRY POINTS

● GENERATING A CALL GIVEN AN ENTRY VALUE

▮ `cu_$generate_call` (AG93)

▮ GENERATES A STANDARD CALL TO THE SPECIFIED PROCEDURE (DESIGNATED BY AN ENTRY VALUE) WITH A SPECIFIED ARGUMENT LIST

▮ DESIGNED PRIMARILY TO BE USED BY COMMAND PROCESSORS THAT CALL A COMMAND WITH AN ARGUMENT LIST BUILT FROM A COMMAND LINE INPUT FROM A TERMINAL

▮ IS PREFACED BY A CALL TO `hes_$make_entry`, WHICH ACCEPTS PATHNAMES

MODIFYING THE STANDARD COMMAND ENVIRONMENT

AN EXAMPLE

```
!print change_cl.pl1 1

change_cl: proc;

dcl codeptr builtin,
    cu_$get_cl_intermediary entry (entry),
    cu_$set_cl_intermediary entry (entry),
    cu_$get_command_processor entry (entry),
    ioa_entry options (variable);

dcl var_entry entry variable,
    my_intermediary entry;

/* FIND OUT THE CURRENT COMMAND PROCESSOR */
    call cu_$get_command_processor (var_entry);
    call ioa_ ("Current command processor is ^p", codeptr (var_entry));

/* FIND OUT THE CURRENT INTERMEDIARY */
    call cu_$get_cl_intermediary (var_entry);
    call ioa_ ("Current intermediary is ^p", codeptr (var_entry));

/* NOW SET MY OWN INTERMEDIARY */
    call cu_$set_cl_intermediary (my_intermediary);

    end change_cl;
```

MODIFYING THE STANDARD COMMAND ENVIRONMENT

AN EXAMPLE

r 19:26 0.132 0

!print my\_intermediary.pl1 1

my\_intermediary: proc;

dcl get\_to\_cl \$unclaimed\_signal entry;  
dcl ioa\_entry options (variable);

call ioa ("TYPE 'start' TO RESTART PROCESS  
TYPE 'release' OR 'rl' TO DISCARD STACK HISTORY");  
call get\_to\_cl \$unclaimed\_signal;  
end my\_intermediary;

r 19:26 0.069 0

!change\_cl

Current command processor is 305|2676

Current intermediary is 77777|1

r 19:27 0.089 0

!lrn 305

305 >sss>bound\_full\_cp\_

do

response

ab

exec\_com

r 19:27 0.105 0

!den >sss>bound\_full\_cp\_ 2676

2676 abbrev|334

r 19:28 0.096 2

AN EXAMPLE

```
!(QUIT)
QUIT
TYPE 'start' TO RESTART PROCESS
TYPE 'release' OR 'rl' TO DISCARD STACK HISTORY
r 19:29 0.168 4 level 2
```

```
!probe
```

```
Condition quit raised at block|154 (level 6).
```

```
stack
```

```
13      simple_command_processor|12211
12      command_processor_|11014
11      abbrev_|7507
10      release_stack|7755
9       unclaimed_signal|27010
8       my_intermediary (line 8)
7       waIl|2602
6       block|154
5       tty_get_line|5763
4       audit_get_line|5073
3       listen_|7666
2       project_start_up_|41673
1       user_init_admin_|42376 (alm)
```

quit

```
!q
```

```
r 19:29 0.980 81 level 2
```

```
!release
```

```
r 19:29 0.044 0
```

```
!change_cl
```

```
Current command processor is 305|2676
```

```
Current intermediary is 456|26
```

```
r 19:29 0.046 0
```

```
!lrn 456
```

```
456 >udd>MED>NDibble>my_intermediary
```

```
my_intermediary
```

```
r 19:30 0.047 0
```



MODIFYING THE STANDARD COMMAND ENVIRONMENT

AN EXAMPLE

|| YOU ARE NOW READY FOR WORKSHOP ||  
#2

TOPIC VI

Advanced Multics I/O

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## REVIEW

### ● I/O SYSTEM BASIC CHARACTERISTICS:

- ▮ LOGICAL INPUT/OUTPUT REQUESTS ARE USED RATHER THAN DEVICE-SPECIFIC PHYSICAL REQUESTS
- ▮ DEVICE INDEPENDENCE IS ACHIEVED VIA THE Multics I/O SWITCH MECHANISM
- ▮ ALL I/O REQUESTS ARE DIRECTED TO A "SWITCH", WHICH IS "ATTACHED" BY A DEVICE-DEPENDENT PROGRAM, CALLED AN I/O MODULE, TO A PARTICULAR DEVICE OR FILE
- ▮ THE SUPPORTING DATA STRUCTURE OF A SWITCH IS AN I/O CONTROL BLOCK (IOCB)

### ● ALL I/O OPERATIONS CAN BE PERFORMED AT THREE BASIC LEVELS:

- ▮ LANGUAGE LEVEL - 'open', 'close', 'get', 'read', 'put', 'write'
- ▮ COMMAND LEVEL - THE 'io\_call' COMMAND
- ▮ SUBROUTINE LEVEL - THE 'iox\_' SUBROUTINE

## CONTROL ORDERS

- CONTROL OPERATIONS ARE ONE EXAMPLE OF EXTENDED POWER OF `iox_` OVER LANGUAGE I/O
- SOME I/O MODULES SUPPORT 'control' OPERATIONS AND SOME DO NOT
- THEY ARE INVOKED BY A CALL TO `iox_$control`
- COMPLETE DESCRIPTIONS OF THE I/O MODULES, AND THE CONTROL ORDERS THEY SUPPORT CAN BE FOUND IN EITHER THE "MPM - Subroutines", ORDER NUMBER AG93, THE "MPM - Peripheral Input/Output", ORDER NUMBER AX49, OR THE "Communications Input/Ouput" MANUAL, ORDER NUMBER CC92.
- THE SUBSYSTEM DESIGNER MAY WANT TO MAKE USE OF SOME OF THE 'control' OPERATIONS SUPPORTED BY THE `tty_` AND `vfile_` I/O MODULES

CONTROL ORDERS

USEFUL TTY CONTROL ORDERS

- tty\_ SUPPORTS THE FOLLOWING CONTROL ORDERS (DOCUMENTED IN CC92)

- ▮ abort  
resetread  
resetwrite

- ▮ THESE FLUSH BOTH THE INPUT AND OUTPUT INTERMEDIATE BUFFERS (abort), THE INPUT BUFFER (resetread) OR THE OUTPUT BUFFER (resetwrite) FOR AN OPENED SWITCH

- ▮ hangup

- ▮ DISCONNECTS THE TELEPHONE LINE CONNECTION OF THE TERMINAL, IF POSSIBLE

- ▮ listen

- ▮ SENDS A WAKEUP TO THE PROCESS ONCE THE LINE ASSOCIATED WITH THIS DEVICE IDENTIFIER IS DIALED UP (SEE THE DISCUSSION OF 'DIALING TERMINALS TO A PROCESS')

- ▮ terminal\_info

- ▮ RETURNS INFORMATION ABOUT THE DEVICE TO WHICH THE MODULE IS ATTACHED IN A USER-SUPPLIED STRUCTURE

- ▮ ANSWERBACK-DERIVED TERMINAL ID

- ▮ TERMINAL TYPE

- ▮ LINE TYPE

- ▮ BAUD RATE

CONTROL ORDERS

USEFUL TTY CONTROL ORDERS

▮ quit\_enable  
quit\_disable

▮ CAUSE 'quit' SIGNAL PROCESSING TO BE ENABLED OR DISABLED FOR THIS DEVICE (FOR EXAMPLE, THE STANDARD PROCESS CREATION CYCLE PROGRAMS ENABLE QUIT ONLY ONCE CONTROL HAS PASSED OUT TO THE USER RING, TO PREVENT A 'quit' FROM ALLOWING A PROCESS TO GAIN CONTROL IN AN INNER RING)

▮ NOTE: EVEN IF TERMINAL QUITTS ARE DISABLED, IT IS POSSIBLE TO SIGNAL 'quit' IN A USER PROGRAM, OR VIA THE 'signal' COMMAND

## CONTROL ORDERS

### USEFUL TTY CONTROL ORDERS

⌋ start

⌋ CAUSES A WAKEUP TO BE SIGNALLED ON THE EVENT CHANNEL ASSOCIATED WITH THIS DEVICE- THE REQUEST IS USED TO RESTART PROCESSING ON A DEVICE WHOSE WAKEUP MAY HAVE BEEN LOST OR DISCARDED (PERHAPS BECAUSE OF BEING INTERRUPTED BY AN `ipc_` OR `timer_manager_` CALLED ROUTINE)

⌋ printer\_off  
printer\_on

⌋ CAUSE THE PRINTER MECHANISM OF THE TERMINAL TO BE TEMPORARILY DISABLED OR REENABLED (IF IT IS PHYSICALLY POSSIBLE FOR THE TERMINAL TO DO SO) - THIS MAY BE USEFUL FOR SUCH THINGS AS ACCEPTING PASSWORDS

⌋ set\_delay  
get\_delay

⌋ SET, OR RETURN, THE NUMBERS OF DELAY CHARACTERS ASSOCIATED WITH THE OUTPUT OF CARRIAGE MOTION CHARACTERS

⌋ DEFAULT VALUES CAN BE USED, OR DELAY CHARACTERS CAN BE SPECIFIED FOR VERTICAL AND HORIZONTAL NEW-LINE OUTPUTS, FORMFEEDS, AND THE LIKE

⌋ set\_editing\_chars  
get\_editing\_chars

⌋ CHANGE, OR FIND OUT, WHAT CHARACTERS ARE BEING USED FOR EDITING INPUT

⌋ THE ERASE AND KILL CHARACTERS CAN BE SET TO ANY CHARACTERS DESIRED SUBJECT TO SOME RESTRICTIONS (E.G., NO CARRIAGE-MOVEMENT CHARACTERS, NOT 'NUL' OR SPACE, AND SO ON)

CONTROL ORDERS

USEFUL TTY CONTROL ORDERS

⌋ set\_input\_translation  
get\_input\_translation

⌋ SETS, OR READS, A TABLE WHICH IS USED FOR TRANSLATION OF  
TERMINAL INPUT TO ASCII

⌋ set\_output\_translation  
get\_output\_translation

⌋ SETS, OR READS, A TABLE WHICH IS USED FOR TRANSLATING ASCII  
CHARACTERS TO THE CODE TO BE SENT TO THE TERMINAL

⌋ set\_input\_conversion  
get\_input\_conversion

⌋ WRITE OR READ A TABLE WHICH IS USED IN CONVERTING INPUT TO  
IDENTIFY ESCAPE SEQUENCES AND CERTAIN SPECIAL CHARACTERS

⌋ set\_output\_conversion  
get\_output\_conversion

⌋ WRITE OR READ A TABLE USED IN FORMATTING OUTPUT TO IDENTIFY  
CERTAIN KINDS OF SPECIAL CHARACTERS (SUCH AS NEWLINE, CARRIAGE  
RETURN, BACKSPACE, AND SO ON)

⌋ set\_special  
get\_special

⌋ WRITE OR READ A TABLE THAT SPECIFIES 1-3 CHARACTER SEQUENCES  
TO BE SUBSTITUTED FOR CERTAIN OUTPUT CHARACTERS, AND CHARACTERS  
WHICH ARE TO BE INTERPRETED AS PARTS OF ESCAPE SEQUENCES ON  
INPUT



CONTROL ORDERS

USEFUL VFILE CONTROL ORDERS

● vfile\_ SUPPORTS THE FOLLOWING CONTROL ORDERS

┆ read\_position

┆ RETURNS THE ORDINAL POSITION OF THE NEXT RECORD (OR BYTE) AND THAT OF THE END OF THE FILE (RELATIVE TO THE FILE BASE)

┆ seek\_head

┆ USED FOR FILES OPENED FOR KEYED-SEQUENTIAL-INPUT OR KEYED-SEQUENTIAL-UPDATE

┆ LOCATES THE FIRST RECORD WITH A KEY WHOSE HEAD HAS THE SPECIFIED RELATION (=, >=, >) WITH A GIVEN SEARCH-KEY

┆ USEFUL FOR APPLICATIONS WHICH MUST LOCATE THE FIRST RECORD THAT, FOR INSTANCE, BEGINS WITH "B"

┆ APPLICATION: MDBM USES THIS WHEN LOCATING TUPLES GIVEN ONLY LEADING PORTION OF KEY

┆ file\_status

┆ RETURNS VARIOUS ITEMS OF INFORMATION ABOUT THE FILE

┆ SEE vfs COMMAND AND vfile\_status\_ SUBROUTINE

CONTROL ORDERS

USEFUL VFILE CONTROL ORDERS

⌋ get\_key  
add\_key  
delete\_key  
reassign\_key

⌋ MANIPULATE THE KEYS IN AN INDEXED FILE DIRECTLY (THUS ALLOWING SEVERAL KEYS TO BE ASSOCIATED WITH A GIVEN DATA RECORD, REASSIGNING THE DATA DESCRIPTOR OF A KEY TO ANOTHER DATA DESCRIPTOR, AND SO ON)

## REVIEW OF IOCB'S

### ● RECALL:

- ▮ AN 'IOCB' IS A STANDARD DATA STRUCTURE
- ▮ IT IS THE PHYSICAL REALIZATION OF A SWITCH
- ▮ THEY ARE FOUND IN THE USER'S PROCESS DIRECTORY
- ▮ AN 'IOCB' IS CREATED BY `iox` WHEN A SWITCHNAME IS USED IN AN "ATTACH STATEMENT" OR "ATTACH COMMAND" FOR THE FIRST TIME IN A PROCESS
- ▮ ONCE AN 'IOCB' IS CREATED, IT LIVES THROUGHOUT THE PROCESS (UNLESS EXPLICITLY DELETED)
- ▮ THE PRINCIPAL COMPONENTS OF AN 'IOCB' ARE 'pointer' VARIABLES AND 'entry' VARIABLES
- ▮ THERE IS ONE 'entry' VARIABLE FOR EACH I/O OPERATION, WITH THE EXCEPTION OF THE ATTACH OPERATION
- ▮ TO PERFORM AN I/O OPERATION THRU THE SWITCH, THE APPROPRIATE ENTRY VALUE IN THE CORRESPONDING 'IOCB' IS CALLED

## REVIEW OF IOCB'S

- WHEN `iox_$attach_name` IS CALLED IT:
  - ▮ INITIALIZES SOME OF THE ELEMENTS IN THE 'IOCB' STRUCTURE
  - ▮ CALLS `<module_name>$<module_name>attach`
  - ▮ THIS ENTRY POINT IN THE I/O MODULE FINISHES THE INITIALIZATION OF THE 'IOCB'
  
- IT IS THE RESPONSIBILITY OF THE I/O MODULE TO MAINTAIN THE ACCURACY OF THE 'IOCB'
  
- ONLY THE `iox_` ENTRY POINTS RESULTING IN ATTACHMENT OF A SWITCH REQUIRE THE MODULE AS AN INPUT ARGUMENT
  - ▮ AFTER THAT TIME, THE 'IOCB' "POINTS TO" THE APPROPRIATE ENTRY POINTS IN THE APPROPRIATE MODULE (THE USER NEED ONLY PROVIDE A POINTER TO THE 'IOCB')

REVIEW OF IOCB'S

```
dcl 1 iocb aligned based,      /* I/O control block. */
2  version fixed init(1), /* Version number of structure. */
2  name char (32),          /* I/O name of this block. */
2  actual_iocb_ptr ptr,    /* IOCB ultimately SYNed to. */
2  attach_descrip_ptr ptr, /* Ptr to printable attach descrip. */
2  attach_data_ptr ptr,    /* Ptr to attach data structure. */
2  open_descrip_ptr ptr,   /* Ptr to printable open description. */
2  open_data_ptr ptr,      /* Ptr to open data structure. */
2  reserved_bit (72),      /* Reserved for future use. */
2  detach_iocb entry (ptr, fixed (35)),
    /* detach_iocb (p,s) */
2  open entry (ptr, fixed, bit (1)aligned, fixed (35)),
    /* open(p,mode,not used,s) */
2  close entry (ptr, fixed (35)), /* close(p,s) */
2  get_line entry (ptr, ptr, fixed (21), fixed (21), fixed (35)),
    /* get_line(p,bufptr,buflen,actlen,s) */
2  get_chars entry (ptr, ptr, fixed (21), fixed (21), fixed (35)),
    /* get_chars(p,bufptr,buflen,actlen,s) */
2  put_chars entry (ptr, ptr, fixed (21), fixed (35)),
    /* put_chars(p,bufptr,buflen,s) */
2  modes entry (ptr, char (*), char (*), fixed (35)),
    /* modes(p,newmode,oldmode,s) */
2  position entry (ptr, fixed, fixed (21), fixed (35)),
    /* position(p,u1,u2,s) */
2  control entry (ptr, char (*), ptr, fixed (35)),
    /* control(p,order,infptr,s) */
2  read_record entry (ptr, ptr, fixed (21), fixed (21), (fixed (35))),
    /* read_record(p,bufptr,buflen,actlen,s) */
2  write_record entry (ptr, ptr, fixed (21), fixed (35)),
    /* write_record(p,bufptr,buflen,s) */
2  rewrite_record entry (ptr, ptr, fixed (21), fixed (35)),
    /* rewrite_record(p,bufptr,buflen,s) */
2  delete_record entry (ptr, fixed (35)),
    /* delete_record(p,s) */
2  seek_key entry (ptr, char (256) varying, fixed (21), fixed (35)),
    /* seek_key(p,key,len,s) */
2  read_key entry (ptr, char (256) varying, fixed (21), fixed (35)),
    /* read_key(p,key,len,s) */
2  read_length entry (ptr, fixed (21), fixed (35)),
    /* read_length(p,len,s) */;
```

REVIEW OF IOCB'S

```
/* "HIDDEN" PORTION          */
2 ios_compatibility ptr, /* Ptr to old DIM's IOS transfer vector. */
2 syn_inhibits bit(36), /* Operations inhibited by SYN. */
2 syn_father ptr, /* IOCB immediately SYNed to. */
2 syn_brother ptr, /* Next IOCB SYNed as this one is. */
2 syn_son ptr; /* First IOCB SYNed to this one. */
```

OTHER STRUCTURES OF INTEREST

```
dcl 1 attach_description based aligned,
2 length fixed bin(17),
2 string char (0 refer (attach_description.length) );

dcl 1 open_description based aligned,
2 length fixed bin(17),
2 string char (0 refer (open_description.length) );
```

● iocb.attach\_descrip\_ptr

▮ BY DEFINITION, IF THIS IS NULL, IOCB IS DETACHED

▮ THE ATTACH DESCRIPTION OF AN IOCB SYNED TO ANOTHER IS A DESCRIPTION OF THE SYNONYMIZATION, NOT A COPY OF THE OTHER IOCB'S ATTACH DESCRIPTION

REVIEW OF IOCB'S

● iocb.attach\_data\_ptr

▮ IS OPTIONALLY USED BY THE I/O MODULE TO LOCATE AN INFORMATION STRUCTURE WHOSE FORMAT AND CONTENT IS MODULE-DEPENDENT

▮ I/O MODULES AREN'T ALLOWED TO STORE SUPPLEMENTARY ATTACH DATA ANYWHERE ELSE

● iocb.open\_descrip\_ptr

▮ BY DEFINITION, IF NULL, THE IOCB IS CLOSED

● iocb.open\_data\_ptr

▮ ANALOGOUS TO attach\_data\_ptr

● iocb.ios\_compatibility

▮ POINTS TO A MODULE THAT SIMULATES ALL FUNCTIONS OF THE OBSOLETE I/O SWITCHING MECHANISM, ios\_

## SYNONYMING

- SYNONYMING IS ACCOMPLISHED THRU THE USE OF THE `syn_` I/O MODULE WHICH:
  - ▮ ATTACHES AN I/O SWITCH, `x`, AS A SYNONYM FOR ANOTHER SWITCH, `y`
  - ▮ THEREAFTER, PERFORMING AN OPERATION ON `x` (EXCEPT FOR DETACH) HAS THE SAME AFFECT AS PERFORMING IT ON `y`
  
- THE ONLY WAY TO MAKE SYNONYMING BEHAVE DIFFERENTLY THAN DESCRIBED ABOVE IS TO USE THE `-inhibit` CONTROL ARGUMENT IN THE ATTACH DESCRIPTION
  - ▮ `iocb.syn_inhibits`
    - ▮ FIRST 15 BITS MAP TO `iocb.open` THRU `iocb.read_length` (detach CAN NOT BE INHIBITED)
    - ▮ WHEN A BIT IS ON, `error_table_$no_operation` IS RETURNED



SYNONYMING

```
!io attach y vfile_my_file
!io open y stream_input
!io attach x syn_y
!io print iocb y
IOCB "y"@ 257|27320
SYN son is "x" @ 257|27160
attach description: "vfile_ >udd>MEDmult>F15D>new>my_file",
                    attach data at 257|34524
open description: "stream_input", open data at 257|27462
detach_iocb        >sl1>bound_command_loop_$err_not_closed (271|2370)
open               "
close              >sss>bound_vfile_$close_file (370|55265)
get_line          >sss>bound_vfile_$get_line_uns_file (370|722)
get_chars         >sss>bound_vfile_$get_chars_uns_file (370|542)
put_chars         >sl1>bound_command_loop_$err_no_operation (271|2360)
modes             "
position          >sss>bound_vfile_$position_uns_file (370|1142)
control           >sss>bound_vfile_$control_uns_file (370|265)
read_record       >sl1>bound_command_loop_$err_no_operation (271|2360)
write_record      "
rewrite_record    "
delete_record     "
seek_key          "
read_key          "
read_length       "
```

!dcn >sl1>bound\_command\_loop\_ 2360  
2360 iox\_|114

SYNONYMING

```
!io print iocb x
IOCB "x" @ 257|27160 (actual IOCB is "y" @ 257|27320)
SYN father is "y" @ 257|27320
attach description: "syn y", attach data at 257|34666
open description: "stream input", open data at 257|27462
detach_iocb      >sl1>bound_command_loop_$syn_detach (271|73205)
open             >sl1>bound_command_loop_$err_not_closed (271|2370)
close           >sss>bound_vfile_$close_file (370|55265)
get_line        >sss>bound_vfile_$get_line_uns_file (370|722)
get_chars       >sss>bound_vfile_$get_chars_uns_file (370|542)
put_chars       >sl1>bound_command_loop_$err_no_operation (271|2360)
modes           "
position        >sss>bound_vfile_$position_uns_file (370|1142)
control         >sss>bound_vfile_$control_uns_file (370|265)
read_record     >sl1>bound_command_loop_$err_no_operation (271|2360)
write_record    "
rewrite_record  "
delete_record   "
seek_key        "
read_key        "
read_length     "
```

SYNONYMING

!io attach inhibited syn\_ y -inhibit close

!io print iocb inhibited

IOCB "inhibited" @ 257|27020 (actual IOCB is "y" @ 257|27320)

SYN father is "y" @ 257|27320

SYN brother is "x" @ 257|27160

attach description: "syn\_ y -inh close", attach data at 257|35402

open description: "stream\_input", open data at 257|27462

detach\_iocb >sl1>bound\_command\_loop\_\$syn\_detach (271|73205)

open >sl1>bound\_command\_loop\_\$err\_not\_closed (271|2370)

close >sl1>bound\_command\_loop\_\$err\_no\_operation (271|2360) (inh)

get\_line >sss>bound\_vfile\_\$get\_line\_uns\_file (370|722)

get\_chars >sss>bound\_vfile\_\$get\_chars\_uns\_file (370|542)

put\_chars >sl1>bound\_command\_loop\_\$err\_no\_operation (271|2360)

modes "

position >sss>bound\_vfile\_\$position\_uns\_file (370|1142)

control >sss>bound\_vfile\_\$control\_uns\_file (370|265)

read\_record >sl1>bound\_command\_loop\_\$err\_no\_operation (271|2360)

write\_record "

rewrite\_record "

delete\_record "

seek\_key "

read\_key "

read\_length "

!pat y x inhibited

y vfile\_ >udd>MEDmult>F15D>new>my\_file

stream\_input

x syn\_ y

inhibited syn\_ y -inh close

!io close inhibited

io\_call: Invalid I/O operation. inhibited

!pat y

y vfile\_ >udd>MEDmult>F15D>new>my\_file

stream\_input

!io close x

!pat y x inhibited

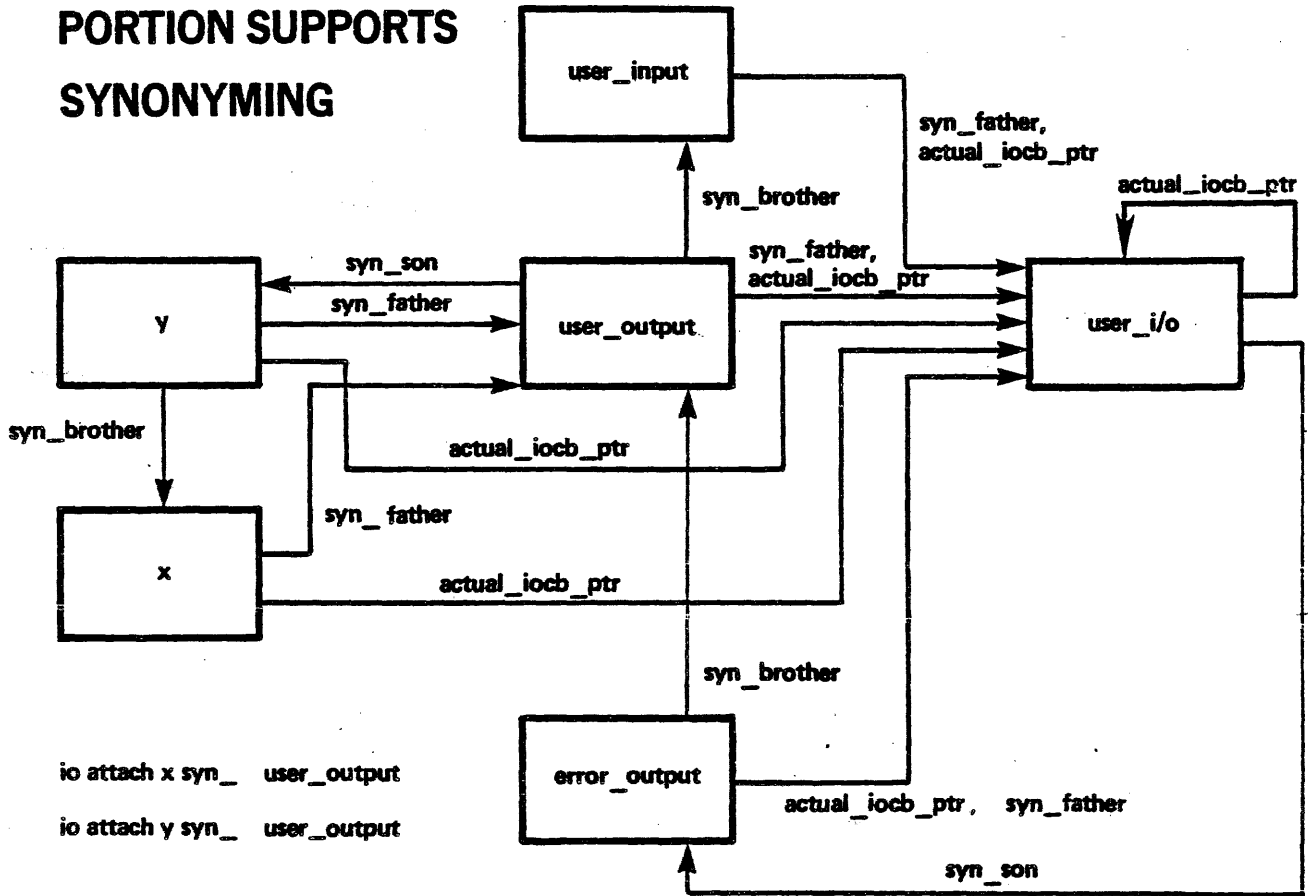
y vfile\_ >udd>MEDmult>F15D>new>my\_file

(not open)

x syn\_ y

inhibited syn\_ y -inh close

**THE HIDDEN  
PORTION SUPPORTS  
SYNONYMING**



## IOX ENTRY POINTS USED IN I/O MODULES

- THE FOLLOWING `iox_` ENTRY POINTS ARE GENERALLY OF USE TO USERS WRITING I/O MODULES:

### ▮ `iox_$propagate`

- ▮ REFLECTS MODIFICATIONS MADE TO AN ULTIMATE IOCB BACK TO ALL MEMBERS OF THE "SYNONYM FAMILY"
- ▮ MUST BE CALLED AT CERTAIN POINTS IN THE I/O MODULE AND NOT UNDER ANY OTHER CIRCUMSTANCES
- ▮ DOES INDEED PROPAGATE INHIBIT BITS BACKWARD TO THE SONS

### ▮ `iox_$find_iocb_n`

- ▮ USED TO FIND (ONE AT A TIME) ALL EXISTING IOCBS IN THE CALLING RING, WHETHER ATTACHED OR DETACHED
- ▮ SEE `print_attach_table` COMMAND

### ▮ `iox_$look_iocb`

- ▮ RETURNS A POINTER TO THE IOCB FOR THE NAMED SWITCH IF IT EXISTS

IOX ENTRY POINTS USED IN I/O MODULES

⌋ iox\_\$err\_no\_operation  
iox\_\$err\_not\_open  
iox\_\$err\_not\_closed  
iox\_\$err\_not\_attached

⌋ THESE ENTRY VALUES ARE ASSIGNED TO ENTRY VARIABLES IN THE IOCB IN ORDER TO RETURN AN ERROR CODE WHEN THAT ENTRY VARIABLE IS CALLED

⌋ THESE ENTRY POINTS SET THE VALUE OF THE 'code' ARGUMENT TO ONE OF THE FOLLOWING:

error\_table\_\$no\_operation

error\_table\_\$not\_open

error\_table\_\$not\_closed

error\_table\_\$not\_attached

TOPIC VII

Writing I/O Modules

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## INTRODUCTION

RECOMMENDED READING: THE Multics USER RING I/O SYSTEM PLM (Order No. AN57) AND CHAPTER 4 IN THE SUBSYSTEM WRITERS' GUIDE.

● SOME INSTANCES IN WHICH A USER MIGHT WISH TO CREATE A NEW I/O MODULE ARE:

▮ TO USE A PSEUDO DEVICE OR FILE

▮ AN I/O MODULE COULD BE USED TO SIMULATE I/O TO/FROM A DEVICE OR FILE (discard\_ IS AN EXAMPLE)

▮ TO SUPPORT A NEW FILE TYPE, SUCH AS ONE IN WHICH RECORDS HAVE MULTIPLE KEYS

▮ REINTERPRETING A FILE

▮ AN I/O MODULE COULD BE DESIGNED TO OVERLAY A NEW STRUCTURE ON A STANDARD TYPE OF FILE (E.G., INTERPRETING AN UNSTRUCTURED FILE AS A SEQUENTIAL FILE BY CONSIDERING 80 CHARACTERS AS A RECORD

▮ TO MONITOR A SWITCH

▮ AN I/O MODULE COULD BE DESIGNED TO PASS OPERATIONS ALONG TO ANOTHER MODULE WHILE MONITORING THEM IN SOME WAY

▮ SEE audit\_

▮ TO SUPPORT AN UNUSUAL DEVICE

▮ BY WORKING THROUGH THE tty I/O MODULE IN THE 'RAW' MODE, ANOTHER I/O MODULE MIGHT TRANSMIT DATA TO/FROM A DEVICE THAT IS NOT A STANDARD Multics DEVICE



## IMPLEMENTATION RULES

- SEVERAL IMPLEMENTATION RULES MUST BE FOLLOWED FOR PROPER OPERATION WITHIN THE Multics I/O ENVIRONMENT

- IN BRIEF, THE RULES ARE:

▮ EXCEPT FOR THE ATTACH OPERATION, THE ENTRY DECLARATION AND PARAMETERS OF A ROUTINE THAT IMPLEMENTS AN I/O OPERATION ARE THE SAME AS THAT OF THE CORRESPONDING ENTRY IN `iox_`

▮ FOR EXAMPLE, IT IS NOT PERMISSIBLE TO DECLARE THE ENTRY VALUE FOR THE 'put chars' OPERATION AS ANYTHING BUT THE DECLARATION OF `iox $put chars`, WHICH IS 'entry (ptr, ptr, fixed bin(21), fixed bin(35))'

▮ THE ATTACH OPERATION ACCEPTS FOUR ARGUMENTS:

(ptr, (\*)char(\*) varying, bit(1) aligned, fixed bin (35))

CORRESPONDING TO:

(iocb\_ptr, option\_array, com\_err\_switch, code)

▮ EXCEPT FOR ATTACH AND DETACH, THE "ULTIMATE" IOCB MUST BE REFERENCED USING THE VALUE OF `iocb_ptr->iocb.actual_iocb_ptr` - IT IS INCORRECT TO USE JUST `iocb_ptr`

▮ IF AN I/O OPERATION CHANGES ANY VALUES IN THE ULTIMATE IOCB, THE I/O MODULE MUST CALL `iox_$propagate` BEFORE RETURNING

▮ ALL I/O OPERATIONS MUST BE EXTERNAL ENTRY POINTS

## IMPLEMENTATION RULES

- WHEN MODIFYING AN IOCB IPS (INTER PROCESS SIGNAL) INTERRUPT (quit, alarm OR cput) CANNOT BE TOLERATED
  
- I/O MODULES SHOULD MASK IPS SIGNALS AT FOLLOWS:
  - ▮ ESTABLISH AN any other HANDLER THAT CALLS terminate\_process\_ IF MASKING IS IN EFFECT
  
  - ▮ CALL hcs\_\$set\_ips\_mask (0, mask)
  
  - ▮ CHANGE THE IOCB
  
  - ▮ CALL hcs\_\$reset\_ips\_mask (mask, mask)

ENTRY POINTS OF AN I/O MODULE

- AN I/O MODULE TYPICALLY HAS ENTRY POINTS FOR THE FOLLOWING:

ATTACH OPERATION

OPEN OPERATION

CLOSE OPERATION

DETACH OPERATION

ANY OTHER OPERATION ENABLED BY THE ABOVE

- THE FOLLOWING IS A SIMPLIFIED SUMMARY OF THE STEPS TAKEN BY THE FIRST FOUR OF THE ABOVE ENTRY POINTS

▮ TO AVOID ADDED CONFUSION, DETAILS ABOUT THE HANDLING OF CONTROL ORDERS AND MODES IS OMITTED

▮ COMPLETE DOCUMENTATION OF THE FOLLOWING IS FOUND IN THE SWG

## ENTRY POINTS OF AN I/O MODULE

### ● MAJOR STEPS OF THE ATTACH OPERATION

- ▮ SET iocb\_ptr->iocb.open APPROPRIATELY
- ▮ SET iocb\_ptr->iocb.detach\_iocb APPROPRIATELY
- ▮ SET iocb\_ptr->attach\_descrip\_ptr APPROPRIATELY
- ▮ CALL iox\_\$propagate

### ● MAJOR STEPS OF OPEN OPERATION

- ▮ SET actual\_iocb\_ptr->iocb.<operation> FOR EVERY OPERATION THAT IS ALLOWED APPROPRIATELY
- ▮ SET actual\_iocb\_ptr->open descrip\_ptr APPROPRIATELY
- ▮ CALL iox\_\$propagate

## ENTRY POINTS OF AN I/O MODULE

### ● MAJOR STEPS OF CLOSE OPERATION

- ▮ SET actual\_iocb\_ptr->iocb.open APPROPRIATELY
- ▮ SET actual\_iocb\_ptr->iocb.detach\_iocb APPROPRIATELY
- ▮ SET actual\_iocb\_ptr->open\_descrip\_ptr TO NULL
- ▮ CLEAN UP (SET BIT COUNTS, FREE STORAGE, ETC.)
- ▮ CALL iox\_\$propagate

### ● MAJOR STEPS OF DETACH OPERATION

- ▮ SET iocb\_ptr->iocb.attach\_descrip\_ptr TO NULL
- ▮ CALL iox\_\$propagate

EXAMPLE OF AN I/O MODULE

- THE FOLLOWING BEHAVES EXACTLY LIKE THE SYSTEM MODULE discard (IT IS NOT THE CODE FOR discard\_ ALTHOUGH MUCH OF THE CODE IS VERY SIMILAR)

```
my_discard_attach: proc(iocb_ptr,option_array,com_err_switch,code);
```

```
    dcl option_array(*) char(*) varying;  
    dcl buflen fixed(21);  
    dcl bufptr ptr;  
    dcl extend_bit bit(1) aligned;  
    dcl infptr ptr;  
    dcl iocb_ptr ptr;  
    dcl key char(256) varying;  
    dcl len fixed(21);  
    dcl com_err_switch bit(1) aligned;  
    dcl mode fixed;  
    dcl newmode char(*);  
    dcl oldmode char(*);  
    dcl order char(*);  
    dcl any_other condition;  
    dcl blkptr ptr;  
    dcl actual_iocb_ptr ptr;  
    dcl code fixed(35);  
    dcl mask bit(36) aligned;
```

```
%include iocb;
```

```
    dcl 1 block based (blkptr),  
        2 attach descrip aligned,  
        3 length fixed bin(17) init (11),  
        3 string char (11) init ("my_discard_"),  
    2 open descrip aligned,  
        3 length fixed bin (17),  
        3 string char (40);
```

```
    dcl free_area area based (get system free area ());  
    dcl get_system_free_area_entry() returns(ptr);  
    dcl com_err_ext entry options(variable);  
    dcl hcs_$set_ips_mask entry (bit(36) aligned, bit(36) aligned);  
    dcl hcs_$reset_ips_mask entry (bit(36) aligned, bit(36) aligned);  
    dcl iox_$propagate_ext entry(ptr);  
    dcl iox_$err_not_open entry() options(variable);  
    dcl error_table_$bad_mode fixed(35) ext;  
    dcl error_table_$not_detached fixed bin(35) ext;  
    dcl error_table_$no_record fixed(35) ext;  
    dcl error_table_$wrong_no_of_args fixed(35) ext;  
    dcl error_table_$no_operation fixed(35) ext;
```

EXAMPLE OF AN I/O MODULE

```
dcl stream_output_mode fixed int static init(2);
dcl sequential_output_mode fixed int static init(5);
dcl keyed_sequential_output_mode fixed int static init(9);
dcl direct_output_mode fixed int static init(12);

dcl (addr,hbound,null,size) builtin;

/* Start Executable Code */

mask = "0"b;
on any_other call handler;
call hcs $set_ips_mask ("0"b, mask);
if hbound(option_array,1)>0 then
    call error (error_table $wrong_no_of_args);
if iocb_ptr->iocb.attach_descrip_ptr^=null() then
    call error (error_table $not_detached);
allocate block in (free_area);
iocb_ptr->iocb.attach_descrip_ptr,
iocb_ptr->iocb.attach_data_ptr = addr(attach_descrip);
iocb_ptr->iocb.detach_iocb = my_discard_detach;
iocb_ptr->iocb.open = my_discard_open;
call_iox $propagate(iocb_ptr);
call hcs $reset_ips_mask(mask,mask);
return;

/* Internal procedure to handle all attach errors.
Calls "com_err" if the "com_err_switch" is set.
In any case, returns to caller of attach external
procedure with proper error code after ensuring
that the IPS interrupt mask is restored. */

error:   proc(c);
         dcl c fixed(35);
         if mask^="0"b then call hcs $reset_ips_mask(mask,mask);
         if com_err_switch then call
             com_err_ (c, "my_discard_");
         code = c;
         go to exit;
         end error;

exit:   return;

my_discard_detach: entry(iocb_ptr, code);
         code = 0;
         mask = "0"b;
         on any_other call handler;
         call hcs $set_ips_mask ("0"b, mask);
         free iocb_ptr->iocb.attach_data_ptr -> block;
         iocb_ptr->iocb.attach_descrip_ptr = null();
         call_iox $propagate(iocb_ptr);
         call hcs $reset_ips_mask(mask,mask);
         return;
```

EXAMPLE OF AN I/O MODULE

```
my_discard_open: entry(iocb_ptr,mode,extend_bit,code);
  mask = "0"b;
  on any other call handler;
  call hcs $set_ips mask ("0"b, mask);
  actual_iocb_ptr = iocb_ptr->iocb.actual_iocb_ptr;
  blkptr = actual_iocb_ptr->iocb.attach_data_ptr;

  if mode=stream_output_mode then do;
    blkptr->open_descrip.string = "stream_output";
    blkptr->open_descrip.length = 13;
    actual_iocb_ptr->iocb.put_chars = my_discard_put_chars;
    actual_iocb_ptr->iocb.modēs = my_discard_modes;
    actual_iocb_ptr->iocb.control = my_discard_control;
  end;
else if mode=sequential_output_mode then do;
  blkptr->open_descrip.string = "sequential_output";
  blkptr->open_descrip.length = 17;
  actual_iocb_ptr->iocb.write_record = my_discard_write;
end;
else if mode=keyed_sequential_output_mode then do;
  blkptr->open_descrip.string = "keyed_sequential_output";
  blkptr->open_descrip.length = 23;
  actual_iocb_ptr->iocb.write_record = my_discard_write;
  actual_iocb_ptr->iocb.seek_key = my_discard_seek_key;
end;
else if mode=direct_output_mode then do;
  blkptr->open_descrip.string = "direct_output";
  blkptr->open_descrip.length = 13;
  actual_iocb_ptr->iocb.write_record = my_discard_write;
  actual_iocb_ptr->iocb.seek_key = my_discard_seek_key;
end;
else do;
  call hcs $reset_ips mask(mask,mask);
  code = error_table_$bad_mode;
  return;
end;

  if extend_bit then blkptr->open_descrip.string
    = blkptr->open_descrip.string||"-extend";
  actual_iocb_ptr->iocb.open_descrip_ptr = addr(open_descrip);
  actual_iocb_ptr->iocb.close = my_discard_close;
  call iox $propagate(actual_iocb_ptr);
  call hcs $reset_ips_mask(mask,mask);
  return;
```



## EXAMPLE OF AN I/O MODULE

```
my_discard_close: entry(iocb_ptr, code);
    code = 0;
    mask = "0"b;
    on any other
        call handler;
    call hcs $set_ips_mask ("0"b, mask);
    actual_iocb_ptr = iocb_ptr->iocb.actual_iocb_ptr;
    blkptr = actual_iocb_ptr->iocb.attach_data_ptr;
    actual_iocb_ptr->iocb.open_descrip_ptr = null();
    actual_iocb_ptr->iocb.detach_iocb = my_discard_detach;
    actual_iocb_ptr->iocb.open = my_discard_open;
    actual_iocb_ptr->iocb.control = iox $err_not_open;
    actual_iocb_ptr->iocb.modes = iox $err_not_open;
    call iox $propagate(actual_iocb_ptr);
    call hcs $reset_ips_mask(mask,mask);
    return;

my_discard_put_chars: entry(iocb_ptr,bufptr,buflen,code);
    code = 0;
    return;

my_discard_modes: entry(iocb_ptr,newmode,oldmode,code);
    code = 0;
    oldmode = "";
    return;

my_discard_write: entry(iocb_ptr,bufptr,buflen,code);
    code = 0;
    return;

my_discard_control: entry(iocb_ptr,order,infptr,code);
    if order = "io_call" then code = error_table_$no_operation;
    else code = 0;
    return;

my_discard_seek_key: entry(iocb_ptr,key,len,code);
    len = 0;
    code = error_table_$no_record;
    return;
```

EXAMPLE OF AN I/O MODULE

```
/* Internal procedure to handle faults while IPS interrupts
are masked. While not masked, any signals are simply
passed on up the stack to their normal handlers. For a
fault while masked, the process is terminated (with the
reason "unable to do critical I/O") because the I/O
control blocks are in an inconsistent state, and we can
tolerate neither spawning a command loop with interrupts
masked nor a restart with a possibly changed mask. */
```

```
handler: procedure;
dcl continue_to_signal_entry (fixed bin(35));
dcl terminate_process_entry (char(*), ptr);
dcl error_table_unable_to_do_io fixed(35) ext;
dcl 1 ti aligned,
    2 version fixed bin init (0),
    2 code fixed bin (35);

if mask ^= "0"b
then
do;
ti.code = error_table_unable_to_do_io; /* very bad trouble */
call terminate_process ("fatal_error", addr (ti));
end;
call continue_to_signal (0);
end handler;

end my_discard_attach;
```

EXAMPLE OF AN I/O MODULE

|| YOU ARE NOW READY FOR WORKSHOP ||  
#3

## TOPIC VIII

### Interprocess Communication

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## OVERVIEW

- THE SUBSYSTEM DESIGNER IS OFTEN FACED WITH REQUIREMENTS FOR SOPHISTICATED INTERPROCESS AND INTRAPROCESS COMMUNICATIONS FACILITIES
  
- SUBROUTINES EXIST WHICH ALLOW THE DESIGNER TO HANDLE:
  - ▮ SYNCHRONIZATION OF SEVERAL COOPERATING PROCESSES
    - ▮ `ipc` FACILITATES INTERPROCESS COMMUNICATION VIA "STOP AND GO" SIGNALS BETWEEN PROCESSES
  
  - ▮ PROTECTION OF CONCURRENTLY ACCESSED DATA BASES
    - ▮ `set lock` ALLOWS COOPERATING PROCESSES TO SHARE A CRITICAL DATA BASE IN A CONTROLLED MANNER
  
  - ▮ INTRAPROCESS TIMING
    - ▮ `timer manager` ALLOWS A PROCESS TO MAKE USE OF SEVERAL CPU OR REAL-TIME TIMERS

## OVERVIEW

### ● CONCEPT:

- ▮ INTERPROCESS COMMUNICATION MEANS THAT TWO DISTINCT PROCESSES CAN PASS INFORMATION BACK AND FORTH BETWEEN THEM
- ▮ A SIMPLE FORM OF INTERPROCESS COMMUNICATION MIGHT INVOLVE THE SHARING OF A COMMON SEGMENT BETWEEN TWO PROCESSES
- ▮ THIS SIMPLE FORM OF COMMUNICATION DOES NOT ALLOW FOR DIRECT SYNCHRONIZATION, HOWEVER
- ▮ Multics SUPPORTS A FULL INTERPROCESS COMMUNICATION FACILITY WHICH ALLOWS FOR CONTROL COMMUNICATION BETWEEN PROCESSES BY MEANS OF 'STOP' AND 'GO' SIGNALS

## OVERVIEW

- THE INTERPROCESS COMMUNICATION MECHANISM

- ┆ IS IMPLEMENTED BY THE `ipc_` AND `hcs_$wakeup` SUBROUTINES

- ┆ IS ACTUALLY MANAGED/COORDINATED BY THE Multics "TRAFFIC CONTROLLER", WHICH IS RESPONSIBLE FOR THE CREATION, DELETION, AND DISPATCHING OF PROCESSES

- ┆ THE DISCUSSION WHICH FOLLOWS WILL NOT DWELL ON TRAFFIC CONTROLLER CONCEPTS (SEE F80B)

- EXAMPLE OF INTERPROCESS COMMUNICATION: THE '`send_message`' AND '`accept_message`' COMMANDS

- ┆ THEY MAKE USE OF THE IPC FACILITY TO PASS TEXT MESSAGES FROM ONE PROCESS TO ANOTHER BY HAVING THEM PRINTED UPON THE RECEIPT OF A 'WAKEUP'

## IPC TERMINOLOGY

### ● event

- ⌋ AN EVENT IS THE OCCURRENCE OF SOMETHING SIGNIFICANT
- ⌋ ONE PROCESS INFORMS ANOTHER THAT AN EVENT HAS OCCURRED BY CALLING `hcs_$wakeup`
- ⌋ A PROCESS WILL ONLY RECEIVE NOTIFICATION OF AN EVENT WHEN IT IS "BLOCKED"
- ⌋ DO NOT ASSUME THAT ONLY "SLEEPING" PROCESSES ARE SUBJECT TO "WAKING UP"

### ● event channel

- ⌋ IS THE ONE-WAY CONTROL PATH OVER WHICH NOTIFICATION OF THE OCCURRENCE OF EVENTS IS TRANSMITTED
- ⌋ IS MAINTAINED BY TRAFFIC CONTROLLER
- ⌋ IS CREATED ON BEHALF OF A USER IN A PARTICULAR RING



## IPC TERMINOLOGY

### ● event-wait channel

- ▮ AN EVENT CHANNEL WHICH, WHEN NOTIFIED OF AN EVENT, CAUSES THE PROCESS TO RESUME EXECUTION IF IT WAS "WAITING" (MORE TECHNICALLY, "BLOCKED"), OR AN EVENT CHANNEL WHICH MAY PERIODICALLY BE "POLLED" TO DETERMINE IF THE EVENT OCCURRED

### ● event-call channel

- ▮ AN EVENT CHANNEL WHICH, WHEN NOTIFIED OF THE OCCURRENCE OF AN EVENT, CAUSES THE INVOCATION OF A SPECIFIED PROCEDURE IN THE PROCESS THAT CREATED THE CHANNEL

### ● blocked

- ▮ A PROCESS CAN GO "BLOCKED" ON AN EVENT-WAIT CHANNEL, IN WHICH CASE IT WILL RESUME EXECUTION UPON EVENT NOTIFICATION
- ▮ IS A PROCESS STATE IN WHICH THE PROCESS IS INACTIVE AND "LISTENING" TO AN EVENT-WAIT CHANNEL (OR CHANNELS)

### ● wakeup

- ▮ A WAKEUP IS THE NOTIFICATION OF THE OCCURRENCE OF AN EVENT
- ▮ IS SENT TO A SPECIFIC PROCESS ACROSS A SPECIFIED EVENT-WAIT OR EVENT-CALL CHANNEL

## IPC TERMINOLOGY

### ● channel\_id

▮ IS THE fixed bin(71) VALUE WHICH IS USED TO UNIQUELY IDENTIFY A PARTICULAR EVENT CHANNEL

▮ IS DERIVED FROM THE SYSTEM CLOCK

▮ IS ASSOCIATED WITH ONE AND ONLY ONE PROCESS

### ● process\_id

▮ IS THE bit(36) VALUE WHICH UNIQUELY IDENTIFIES A PROCESS

▮ IN ORDER TO SEND A WAKEUP, ONE SPECIFIES THE channel\_id OF THE EVENT CHANNEL AND THE process\_id OF THE PROCESS OWNING THE EVENT CHANNEL

### ● message

▮ A 72-BIT VALUE OF ARBITRARY CONTENT CONTAINING INFORMATION WHOSE INTERPRETATION IS APPLICATION DEPENDENT

## IPC TERMINOLOGY

### CHANNEL TYPE

### OPTIONS

#### event-call

When a process establishes an event call channel, execution proceeds with the statement immediately following the call to create the event-call channel. The process does not (and should not) go blocked on the channel. If/when a wakeup is received, execution is momentarily interrupted and the procedure specified in the declaration is invoked. When this procedure completes, execution of the interrupted procedure continues.

#### event-wait

There are two options available:

- 1) The process can explicitly go blocked on the channel by calling `ipc_block`. In this case, the statement following the call to `ipc_block` will not be executed until/unless a wakeup is received on that channel. Notice that this effectively blocks the process itself.
- 2) The process may elect not to go blocked on this channel at all. Rather it will, from time to time, explicitly inquire as to whether or not a wakeup has been received.

This polling technique is implemented via the `ipc_read_ev_chn` entry point.

## IPC PROTOCOL

- SINCE IT IS NECESSARY TO KNOW THE `channel_id` AND `process_id` IN ORDER TO COMMUNICATE WITH ANOTHER PROCESS, SOME STANDARD PROTOCOL IS REQUIRED

- THE STEPS REQUIRED TO SET UP INTERPROCESS COMMUNICATION:

### ▮ STEP 1

- ▮ PROCESS\_1 CREATES AN EVENT CHANNEL, BY WHICH ACT PROCESS\_1 RECEIVES THE `channel_id` OF THE NEWLY CREATED EVENT CHANNEL

### ▮ STEP 2

- ▮ PROCESS\_1 STORES THE `channel_id` AND ITS OWN `process_id` IN SOME KNOWN LOCATION IN A SHARED SEGMENT, THUS ALLOWING THESE VALUES TO BE ACCESSED BY OTHER COOPERATING PROCESSES

### ▮ STEP 3

- ▮ SOME OTHER PROCESS, SAY PROCESS\_2, OBTAINS THE `process_id` AND `channel_id` VALUES FROM THE SHARED SEGMENT

- IT IS NOW POSSIBLE FOR PROCESS\_2 TO SEND WAKEUPS TO PROCESS\_1

## IPC PROTOCOL

- COMMUNICATION ON AN EVENT CHANNEL IS ONE-WAY ONLY
  
- IF PROCESS\_1 WISHES TO COMMUNICATE VIA ipc\_ WITH PROCESS\_2 IN THE ABOVE SCENARIO:
  - ▮ PROCESS\_2 WOULD NEED ITS OWN EVENT-CHANNEL
  
  - ▮ IT WOULD BE NECESSARY FOR PROCESS\_2 TO REPEAT THE STEPS OUTLINED ABOVE
  
- WHENEVER INTERPROCESS COMMUNICATION IS USED BETWEEN USER PROCESSES AND A SYSTEM PROCESS, A SUBROUTINE IS GENERALLY CALLED TO OBTAIN THE channel\_id AND process\_id OF THE SYSTEM PROCESS

IPC PROTOCOL  
SENDING WAKEUPS

- WAKEUPS ARE SENT FROM ONE PROCESS TO ANOTHER BY CALLING THE `hcs_$wakeup` ENTRY POINT

▮ USAGE

call `hcs_$wakeup (process_id, channel_id, message, code);`

▮ `process_id` SPECIFIES TARGET PROCESS

▮ `channel_id` IDENTIFIES CHANNEL THAT WAS SET UP BY TARGET PROCESS

▮ `message` IS TWO WORDS HAVING SOME AGREED-UPON MEANING FOR THE COOPERATING PROCESSES

▮ `code` IS A NON-STANDARD ERROR CODE (NOT IDEALLY SUITED TO `com_err_`)

▮ = 1 IF SIGNALLING WAS CORRECTLY DONE, BUT THE TARGET PROCESS WAS IN THE STOPPED STATE (THE STATE A PROCESS IS IN JUST BEFORE THE FINAL STEP IN PROCESS TERMINATION)

▮ = 2 IF AN INPUT ARGUMENT WAS INCORRECT, AND SIGNALLING WAS ABORTED

▮ = 3 IF THE TARGET PROCESS WAS NOT FOUND, AND SIGNALLING WAS ABORTED

▮ = `error_table_$invalid_channel` IF THE CHANNEL IDENTIFIER WAS NOT VALID

## IPC SUBROUTINES

- EVENT CHANNELS ARE CREATED, DESTROYED, AND MANIPULATED VIA THE ipc\_ SUBROUTINE

- ipc\_ ENTRY POINTS MAY BE CLASSIFIED:

- ┆ CREATING AND DESTROYING EVENT CHANNELS

- ┆ ipc\_\$create\_ev\_chn

- ┆ ipc\_\$delete\_ev\_chn

- ┆ ipc\_\$decl\_ev\_call\_chn

- ┆ ipc\_\$decl\_ev\_wait\_chn

- ┆ GOING BLOCKED ON AN EVENT\_WAIT CHANNEL

- ┆ ipc\_\$block

- ┆ READING AN EVENT-WAIT CHANNEL

- ┆ ipc\_\$read\_ev\_chn

## IPC SUBROUTINES

### ▮ CONTROL FUNCTIONS

▮ ipc\_\$drain\_chn

▮ ipc\_\$cutoff

▮ ipc\_\$reconnect

### ▮ MASKING OR ASSIGNING PRIORITY TO EVENT CHANNELS

▮ ipc\_\$set\_call\_prior

▮ ipc\_\$set\_wait\_prior

▮ ipc\_\$mask\_ev\_calls

▮ ipc\_\$unmask\_ev\_calls



IPC SUBROUTINES

IPC ERROR CODES

- ALL CALLS TO ipc\_ RETURN A NONSTANDARD STATUS CODE

<u>CODE</u>	<u>MEANING</u>
0	No error.
1	A ring violation; for instance, the event channel resides in a ring that is not accessible from the caller's ring.
2	The table that contains the event channels for a given ring was not found.
3	The specified event channel was not found.
4	A logical error in using the ipc_ subroutine was encountered; for instance, waiting on an event-call channel.
5	A bad argument was passed to the ipc_ subroutine; for instance, a zero-value event channel_id.

IPC SUBROUTINES

IPC ERROR CODES

● convert\_ipc\_code\_

▮ AN OBSCURE SUBROUTINE THAT CONVERTS A NONSTANDARD ipc\_ CODE TO A SYSTEM STANDARD CODE

▮ MAPPING

<u>ipcode</u>	<u>returned code</u>
0	0
1	error_table_\$bad_ring_brackets
2-4	error_table_\$no_message
5	error_table_\$argerr

## IPC SUBROUTINES

### CREATING AND DESTROYING EVENT CHANNELS

- `ipc` CREATES EVENT-WAIT CHANNELS BY DEFAULT; HENCE, IN ORDER TO CREATE AN EVENT-CALL CHANNEL, IT IS FIRST NECESSARY TO CREATE AN EVENT-WAIT CHANNEL, AND THEN TO CHANGE IT INTO AN EVENT-CALL CHANNEL

- ENTRY POINTS TO CREATE AND DESTROY CHANNELS

▮ `ipc_$create_ev_chn`

▮ CREATES AN EVENT-WAIT CHANNEL IN THE CURRENT RING, RETURNING THE `channel_id` OF THE NEWLY CREATED CHANNEL, AND THE NONSTANDARD STATUS CODE

▮ `ipc_$delete_ev_chn`

▮ DESTROYS AN EVENT CHANNEL PREVIOUSLY CREATED BY THE PROCESS, REQUIRING THE `channel_id` OF THE CHANNEL TO BE DESTROYED AND RETURNING THE NONSTANDARD STATUS CODE

▮ ONLY THE PROCESS CREATING AN EVENT CHANNEL (OR THE INITIALIZER PROCESS) MAY DESTROY IT

▮ EVENT-CHANNELS ARE AUTOMATICALLY DESTROYED AT PROCESS TERMINATION TIME

IPC SUBROUTINES  
CREATING AND DESTROYING EVENT CHANNELS

|| ipc\_\$decl\_ev\_call\_chn

|| CHANGES AN EVENT-WAIT CHANNEL INTO AN EVENT-CALL CHANNEL

|| call ipc\_\$decl\_ev\_call\_chn (channel\_id, proc\_entry,  
                                  data\_ptr, priority, code);

|| channel\_id IDENTIFIES EVENT-WAIT CHANNEL TO BE CHANGED

|| proc\_entry IS ENTRY VALUE OF "HANDLER" TO BE INVOKED UPON  
RECEIPT OF WAKEUP

|| data\_ptr POINTS TO A USER-DEFINED REGION CONTAINING DATA  
FOR proc\_entry TO INTERPRET

|| priority INDICATES WHICH (OF POTENTIALLY MANY) SIMULTANEOUS  
EVENT-CALL CHANNEL EVENTS IN THIS RING WILL BE HONORED  
FIRST (THE LOWEST NUMBER IS HONORED FIRST)

|| ipc\_\$decl\_ev\_wait\_chn

|| USED TO CHANGE AN EVENT-CALL CHANNEL INTO AN EVENT-WAIT CHANNEL

|| IT REQUIRES THE channel\_id OF AN EVENT-CALL CHANNEL AND IT  
RETURNS THE NONSTANDARD STATUS CODE

|| NOTE: SINCE EVENT-WAIT CHANNELS ARE CREATED BY DEFAULT, THIS  
IS USED ONLY TO CHANGE AN EVENT-CALL CHANNEL BACK INTO AN  
EVENT-WAIT CHANNEL

IPC SUBROUTINES

INVOKING AN EVENT-CALL PROCEDURE

- WHEN A PROCESS IS AWAKENED ON AN EVENT-CALL CHANNEL, CONTROL IS IMMEDIATELY PASSED TO THE "HANDLER" PROCEDURE PREVIOUSLY SPECIFIED IN THE CALL TO `ipc_$decl_ev_call_chn`

- "HANDLER" WILL BE CALLED WITH ONE ARGUMENT, A POINTER TO A STRUCTURE CONTAINING INFORMATION ABOUT THE WAKEUP

- ▮ THE HANDLER SHOULD DECLARE THE FOLLOWING STRUCTURE BASED UPON THE `event_call_info_ptr` PARAMETER

```
dcl 1 event_call_info based (event_call_info_ptr),
    2 channel_id      fixed bin(71),
    2 message         fixed bin(71),
    2 sender          bit(36),
    2 origin,
    3 dev_signal      bit(18) unal,
    3 ring            bit(18) unal,
    2 data_ptr        ptr;
```

- ▮ RECALL: THE VALUE OF `data_ptr` WAS SPECIFIED IN THE CALL TO `ipc_$decl_ev_call_chn` AND POINTS TO STRUCTURE OF THE USER'S CHOOSING

- ▮ SEE `>ldd>include>event_call_info.incl.pl1`

IPC SUBROUTINES  
GOING BLOCKED ON AN EVENT CHANNEL

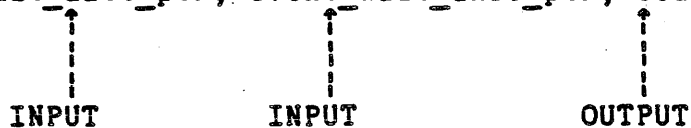
● `ipc_$block`

▮ CAUSES THE PROCESS TO GO BLOCKED ON THE SPECIFIED EVENT-WAIT CHANNEL(S)

▮ IN-LINE PROGRAM EXECUTION WILL NOT PROCEED UNTIL/UNLESS A WAKEUP IS RECEIVED ON ONE OF THE CHANNELS

▮ USAGE

```
call ipc_$block (wait_list_ptr, event_wait_info_ptr, code);
```



▮ A POINTER TO THE BASE OF A USER-ALLOCATED "WAIT-LIST" STRUCTURE IS REQUIRED

▮ THIS WAIT-LIST STRUCTURE CONTAINS THE `channel_id`S OF THE EVENT-WAIT CHANNELS TO GO BLOCKED ON (TO LISTEN TO)

```
dcl 1 wait_list          based aligned,  
    2 nchan             fixed bin,  
    2 pad               bit(36),  
    2 channel_id (n refer wait_list.nchan) fixed bin(71);
```

IPC SUBROUTINES

GOING BLOCKED ON AN EVENT CHANNEL

⌋ ipc \$block ALSO REQUIRES A POINTER TO THE BASE OF A STRUCTURE INTO WHICH IT CAN PUT INFORMATION ABOUT THE EVENT THAT FREED THE PROCESS FROM ITS BLOCKED STATE

```
dcl 1 event_wait_info aligned based(event_wait_info_ptr),
    2 channel_id      fixed bin(71),
    2 message        fixed bin(71),
    2 sender         bit(36),
    2 origin,
    3 dev_signal     bit(18) unaligned,
    3 ring          bit(18) unaligned,
    2 channel_index  fixed bin;      /* INDEX INTO
                                     wait_list.channel_id */
```

⌋ POSSIBLE USE OF event\_wait\_info.sender:

PASS sender TO get\_userid SUBROUTINE (WHICH USES THE RESTRICTED SEGMENT >sc1>answer\_table) AND IT WILL RETURN Person\_id AND Project\_id

⌋ SEE >ldd>include>event\_wait\_info.incl.pl1

IPC SUBROUTINES

READING AN EVENT-WAIT CHANNEL

- THE PROCESS CREATING AN EVENT-WAIT CHANNEL MAY SIMPLY INQUIRE AS TO WHETHER OR NOT AN EVENT HAS OCCURRED ON A SPECIFIED EVENT-WAIT CHANNEL - THIS IS REFERRED TO AS "READING" THE EVENT-WAIT CHANNEL
  
- READING THE EVENT "RESETS" IT
  
- `ipc_$read_ev_chn`
  - ⌋ READS THE INFORMATION ABOUT AN EVENT ON A SPECIFIED EVENT-WAIT CHANNEL IF THE EVENT HAS OCCURRED
  
  - ⌋ REQUIRES AS INPUT THE `channel_id` OF THE EVENT-WAIT CHANNEL TO BE READ
  
  - ⌋ RETURNS A VALUE INDICATING WHETHER OR NOT AN EVENT OCCURRED, AND IF AN EVENT HAS OCCURRED, RETURNS INFORMATION ABOUT THAT EVENT IN `event_wait_info` STRUCTURE
  
  - ⌋ ALSO RETURNS THE NONSTANDARD STATUS CODE



IPC SUBROUTINES  
READING AN EVENT-WAIT CHANNEL

```
! il 70  
! pr DEMO_READ.pl1 1
```

```
DEMO_READ: proc;
```

```
dcl
```

```
ipc_$create_ev_chn entry (fixed bin (71), fixed bin (35)),  
ipc_$delete_ev_chn entry (fixed bin(71), fixed bin(35)),  
ipc_$read_ev_chn entry (fixed bin (71), fixed bin, ptr,  
                        fixed bin(35)),  
hcs_$wakeup entry (bit (36), fixed bin (71), fixed bin (71),  
                  fixed bin (35)),  
  
addr builtin,  
ioa_entry options (variable),  
com_err_entry options (variable),  
get_process_id_entry returns (bit (36));
```

```
dcl
```

```
channel_id fixed bin (71),  
message fixed bin (71),  
code fixed bin (35),  
ev_occurred fixed bin;
```

```
%include event_info; /* My own private include file */
```

```
/* This short example illustrates the fact that reading an event-wait  
channel has the effect of 'resetting' it. */
```

```
call ipc_$create_ev_chn (channel_id, code);  
if code ^= 0 then call trouble;  
message = 1;  
call hcs_$wakeup (get_process_id_ (), channel_id, message,  
                 code);  
  
if code ^= 0 then call trouble;  
else call ioa_ ("WAKEUP successfully completed.");  
call ipc_$read_ev_chn (channel_id, ev_occurred,  
                      addr (event_info), code);  
  
if code ^= 0 then call trouble;  
call ioa_ ("read_ev_chn says the event ^[has not^;has^] occu  
\crrred, and the message is ^i.", ev_occurred+1, event_info.message);  
call ipc_$read_ev_chn (channel_id, ev_occurred,  
                      addr (event_info), code);  
  
if code ^= 0 then call trouble;  
call ioa_ ("A second call of read_ev_chn says: The event ^[h  
\cas not^;has^] occurred.", ev_occurred+1);  
  
trouble: proc;  
call com_err_ (code, "DEMO_READ", "Something unexpected  
\c occurred.");  
goto bottom;  
end trouble;
```

IPC SUBROUTINES

READING AN EVENT-WAIT CHANNEL

bottom:

```
    call ipc_$delete ev_chn (channel_id, code);  
    /* Ignore bad code, if it should occur. */  
end DEMO_READ;
```

! DEMO\_READ

WAKEUP successfully completed.

read\_ev\_chn says the event has occurred, and the message is 1.

A second call of read\_ev\_chn says: The event has not occurred.

IPC SUBROUTINES  
CONTROL FUNCTIONS

● RESETTING CHANNELS AND INHIBITING THE NOTIFICATION OF EVENTS

▮ `ipc_$drain`

- ▮ RESETS AN EVENT-WAIT CHANNEL SO THAT ANY PENDING EVENTS (EVENTS THAT HAVE BEEN RECEIVED AND QUEUED UP BUT NOT PROCESSED FOR THAT CHANNEL) ARE REMOVED
- ▮ REQUIRES THE `channel_id` OF THE EVENT-WAIT CHANNEL
- ▮ RETURNS THE NONSTANDARD STATUS CODE
- ▮ OFTEN USED IN 'cleanup' HANDLERS

▮ `ipc_$cutoff`

- ▮ INHIBITS THE "READING" (IN THE GENERAL SENSE) OF PENDING OR FUTURE EVENTS ON A SINGLE SPECIFIED EVENT (WAIT OR CALL) CHANNEL
- ▮ NOTE THAT MORE (NEW) EVENTS CAN BE RECEIVED (AND QUEUED UP), BUT THEY WILL NOT CAUSE THE PROCESS TO WAKE UP
- ▮ REQUIRES THE `channel_id` OF THE EVENT CHANNEL TO CUTOFF, AND RETURNS THE NONSTANDARD STATUS CODE
- ▮ AN ATTEMPT TO READ A CUTOFF CHANNEL RESULTS IN CODE 4 (A LOGICAL ERROR IN USING `ipc_` WAS ENCOUNTERED)

IPC SUBROUTINES  
CONTROL FUNCTIONS

▮ `ipc_$reconnect`

- ▮ ENABLES THE READING OF EVENTS ON A SINGLE SPECIFIED EVENT CHANNEL FOR WHICH READING HAD PREVIOUSLY BEEN INHIBITED BY A CALL TO `ipc_$cutoff`
  
- ▮ WHEN CALLED, ALL PENDING SIGNALS, WHETHER RECEIVED BEFORE OR DURING THE TIME READING WAS INHIBITED, ARE HENCEFORTH AVAILABLE FOR READING (IN THE GENERAL SENSE)
  
- ▮ REQUIRES THE `channel id` OF THE EVENT CHANNEL WHICH HAD BEEN CUTOFF, AND RETURNS THE NONSTANDARD STATUS CODE

IPC SUBROUTINES  
CONTROL FUNCTIONS

```
!print DEMO_CUTOFF.pl1 1  
DEMO_CUTOFF: proc;
```

```
/* This experiment demonstrates the cutting off and reconnection of  
a single channel. It accomplishes this in the following steps:
```

- 1 get wait channel
- 2 issue wakeup with msg = 1
- 3 cutoff channel
- 4 issue wakeup with message = 2
- 5 reconnect
- 6 read channel twice
- 7 delete channel

```
ioa_ is called at strategic points to confirm ipc's behavior. */
```

```
dcl
```

```
channel_id fixed bin (71),  
code fixed bin (35),  
i fixed bin,  
ev occurred fixed bin,  
addr builtin,  
get_process_id entry returns (bit (36)),  
message fixed bin (71),  
hcs_$wakeup entry (bit (36), fixed bin (71), fixed bin (71),  
fixed bin (35)),  
ipc_$create_ev_chn entry (fixed bin (71), fixed bin (35)),  
ipc_$delete_ev_chn entry (fixed bin (71), fixed bin (35)),  
ipc_$cutoff_entry (fixed bin (71), fixed bin (35)),  
ipc_$reconnect entry (fixed bin (71), fixed bin (35)),  
ipc_$read_ev_chn entry (fixed bin (71), fixed bin, ptr,  
fixed bin (35)),  
(com_err_, ioa_) entry options (variable),  
1 event_info aligned,  
2 channel_id fixed bin (71),  
2 message fixed bin (71),  
2 sender bit (36),  
2'origin,  
3 dev_signal bit (18) unaligned,  
3 ring bit (18) unaligned,  
2 channel_index fixed bin;  
call ipc_$create_ev_chn (channel_id, code);  
if code ^= 0 then call trouble;  
call hcs_$wakeup (get_process_id_ (), channel_id, 1, code);  
if code ^= 0 then call trouble;  
call ioa_ ("First wakeup successfully performed with msg = 1");  
call ipc_$cutoff (channel_id, code);  
if code ^= 0 then call trouble;  
call ioa_ ("Channel successfully cutoff.");  
call hcs_$wakeup (get_process_id_ (), channel_id, 2, code);  
if code ^= 0 then call trouble;
```

IPC SUBROUTINES  
CONTROL FUNCTIONS

```
call ioa_ ("2nd wakeup performed while channel was cutoff.");
call ipc_$reconnect (channel_id, code);
if code ^= 0 then call trouble;
do i = 1 to 2;
    call ipc_$read_ev_chn (channel_id, ev_occurred,
        addr(event_info), code);
    if code ^= 0 then call trouble;
    call ioa_ ("^[First^;Second^] reading channel after reco
\connect. The event ^[hasn't^;has^] occurred. ^[s^;^/The message is ^
\ci.^]", i, ev_occurred+1, ev_occurred+1, event_info.message);
end;
wrapup:
    call ipc_$delete_ev_chn (channel_id, code);
    return;

trouble: proc;
    call com_err_ (code, "DEMO_CUTOFF", "Error not expected.");
    goto wrapup;
end trouble;

end DEMO_CUTOFF;
```

```
!DEMO_CUTOFF
First wakeup successfully performed with msg = 1
Channel successfully cutoff.
2nd wakeup performed while channel was cutoff.
First reading channel after reconnect. The event has occurred.
The message is 1.
Second reading channel after reconnect. The event has occurred.
The message is 2.
```

## IPC SUBROUTINES

### MASKING OR ASSIGNING PRIORITY TO EVENT CHANNELS

- SINCE THERE IS ACTUALLY SOME DELAY BETWEEN THE TIME A WAKEUP IS RECEIVED AND THE TIME THE PROCESS IS NOTIFIED, IT IS POSSIBLE FOR SEVERAL WAKEUPS TO BE PRESENT BY THE TIME A PROCESS IS AWAKENED

- ▮ IT IS POSSIBLE TO SPECIFY RELATIVE PRIORITIES AMONG EVENT-CALL CHANNELS

- ▮ POSSIBLE TO SPECIFY THAT EVENT-CALL CHANNELS HAVE PRIORITY OVER EVENT-WAIT CHANNELS AND VICE VERSA

- ▮ BY DEFAULT, EVENT-CALL CHANNELS HAVE PRIORITY OVER EVENT-WAIT CHANNELS

- MANIPULATING PRIORITIES

- ▮ `ipc_$set_wait_prior`

- ▮ CAUSES EVENT-WAIT CHANNELS TO BE GIVEN PRIORITY OVER EVENT-CALL CHANNELS (THIS IS NOT THE DEFAULT)

- ▮ ONLY EVENT CHANNELS IN CURRENT RING ARE AFFECTED

## IPC SUBROUTINES

### MASKING OR ASSIGNING PRIORITY TO EVENT CHANNELS

#### ▮ `ipc_$set_call_prior`

▮ CAUSES EVENT-CALL CHANNELS TO BE GIVEN PRIORITY OVER EVENT-WAIT CHANNELS (THIS IS THE DEFAULT)

▮ ONLY EVENT CHANNELS IN CURRENT RING ARE AFFECTED

#### ▮ `ipc_$mask_ev_calls`

▮ CAUSES THE `ipc $block` ENTRY POINT TO COMPLETELY IGNORE ALL EVENT-CALL CHANNELS IN THE CALLER'S RING (I.E., TO MASK THEM) SO THAT ANY WAKEUPS SENT ACROSS EVENT-CALL CHANNELS ARE INSTEAD QUEUED UP

▮ CAUSES A "MASK COUNTER" TO BE INCREMENTED; MASKING IS IN EFFECT SO LONG AS COUNTER  $\neq$  0

#### ▮ `ipc_$unmask_ev_calls`

▮ DECREMENTS THE "MASK COUNTER"; EVENT-CALLS ARE UNMASKED (NOTICED) WHEN COUNTER = 0

▮ SEVERAL EXTERNAL PROCEDURES MAY NEED TO MASK AND UNMASK. AN INCREMENTAL COUNTER PERMITS INDISCRIMINATE CALLS WITHOUT FEAR OF PREMATURE UNMASKING



IPC SUBROUTINES

AN EXAMPLE USING EVENT-WAIT CHANNELS

```
!print abs_print_punch.pl1 1

abs_print_punch: proc;

dcl ME char (15) init ("abs_print_punch") static options (constant);
dcl hcs $initiate entry (char (*), char (*), char (*),
    fixed bin (1), fixed bin (2), ptr, fixed bin (35)),
    get_process_id_entry returns (bit (36)),
    ipc $create_ev_chn entry (fixed bin (71), fixed bin (35)),
    ipc $block entry (ptr, ptr, fixed bin (35)),
    (ioa_, com_err_) entry options (variable),
    expand_pathname_entry (char (*), char (*), char (*),
    fixed bin (35)),
    dprint_entry (char (*), char (*), ptr, fixed bin (35));

dcl code fixed bin (35);
dcl 1 ipc_info based (seg_ptr),
    2 target_process_id bit (36) unal,
    2 target_chnl_id fixed bin (71);

%include dprint_arg; /* USED BY dprint_ SUBROUTINE */

dcl 1 wait_list based (wait_list_ptr),
    2 nchan fixed bin,
    2 channel_id (0 refer (nchan)) fixed bin (71);

dcl 1 event_info,
    2 channel_id fixed bin (71),
    2 message fixed bin (71),
    2 sender bit (36),
    2 origin,
    3 dev_signal bit (18) unal,
    3 ring bit (18) unal,
    2 channel_index fixed bin;

dcl (seg_ptr, wait_list_ptr) ptr;

dcl dprint_paths$ char (168) external static;
dcl dir char (168), entry char (32);

/* SUPPLY ipc PROTOCOL INFORMATION */
call hcs $initiate (">udd>F15dw>Auerbach", "ipc_seg1", "",
    0, 1, seg_ptr, code);
if seg_ptr = null () then call ERROR;
call ipc $create_ev_chn (
    seg_ptr -> ipc_info.target_chnl_id, code);
if code ^= 0 then call ERROR;
seg_ptr -> ipc_info.target_process_id = get_process_id ();

/* NOW GO BLOCKED WAITING FOR FURTHER INSTRUCTIONS */
allocate wait_list;
```

IPC SUBROUTINES

AN EXAMPLE USING EVENT-WAIT CHANNELS

```
wait_list_ptr -> wait_list.nchan = 1;
wait_list_ptr -> wait_list.channel_id (1) =
    seg_ptr -> ipc_info.target_chnl_id;

call ipc $block (wait_list_ptr, addr (event_info), code);
if code ^= 0 then call ERROR;

/* AN EVENT HAS OCCURRED - EXAMINE */
call ioa_ ("CHANNEL ID ^i^/MESSAGE ^i^/SENDER ^.3b",
    event_info.channel_id,
    event_info.message,
    event_info.sender);
call ioa_ ("DEV SIGNAL ^.3b^/RING ^.3b^/CHANNEL INDEX ^i",
    event_info.origin.dev_signal,
    event_info.origin.ring,
    event_info.channel_index);

if event_info.message = 0 then goto CANCEL;
if event_info.message = 1 then goto PRINT;
if event_info.message = 2 then goto PUNCH;
/* INVALID MESSAGE */
call com_err (0, ME, "INVALID REQUEST CODE ^i",
    event_info.message);
return;

CANCEL:
/* NO PRINTING OR PUNCHING AT ALL */
call ioa_ ("PRINT/PUNCH REQUEST CANCELLED");
return;

PRINT:
/* PRINT REQUEST */
dprint_arg.pt_pch = 1;
dprint_arg.output_module = 1;
dprint_arg.class = "printer";
goto DPRINT;

PUNCH:
/* PUNCH REQUEST */
dprint_arg.pt_pch = 2;
dprint_arg.output_module = 3;
dprint_arg.class = "punch";
```

IPC SUBROUTINES

AN EXAMPLE USING EVENT-WAIT CHANNELS

DPRINT:

```
call expand_pathname (dprint_paths$, dir, entry, code);
if code ^= 0 then call ERROR;
dprint_arg.version = 4;
dprint_arg.copies = 1;
dprint_arg.delete = 0;
dprint_arg.queue = 3;
dprint_arg.notify = 1;
dprint_arg.heading = "";
dprint_arg.dest = "";
dprint_arg.nep = "0";
dprint_arg.single = "0";
dprint_arg.non_edited = "0";
dprint_arg.truncate = "0";
dprint_arg.center_top_label = "0";
dprint_arg.center_bottom_label = "0";
dprint_arg.lmargin = 10;
dprint_arg.line_lth = -1;
dprint_arg.page_lth = -1;
dprint_arg.top_label = "";
dprint_arg.bottom_label = "";
call dprint (rtrim (dir), rtrim (entry),
            addr (dprint_arg), code);
if code ^= 0 then call ERROR;
```

/\* REPORT ACTION AND QUIT \*/

```
call ioa ("^a>^a ^[DPRINT^;DPUNCH^] REQUEST SUBMITTED.",
         dir, entry, dprint_arg.pt_pch);
call ioa ("END ^a", ME);
```

ERROR:

```
proc;
    call com_err_ (code, ME);
    goto FINIS;
end;
```

FINIS:

```
end abs_print_punch;
```

IPC SUBROUTINES

AN EXAMPLE USING EVENT-WAIT CHANNELS

r 11:12 0.122 15

!print driver\_ipc.pl1 1

driver\_ipc: proc;

dcl ME char (10) init ("driver ipc") static options (constant);

dcl hcs\_\$initiate entry (char (\*), char (\*), char (\*),  
fixed bin (1), fixed bin (2), ptr, fixed bin (35));

dcl code fixed bin (35);

dcl seg\_ptr ptr;

dcl 1 ipc\_info based (seg\_ptr),  
2 his\_process\_id bit (36),  
2 his\_chnl\_id fixed bin (71);

dcl hcs\_\$wakeup entry (bit (36), fixed bin (71), fixed bin (71),  
fixed bin (35));

dcl dprint\_paths\$ char (168) external static;

dcl message fixed bin (71);

dcl (ioa\_, com\_err\_) entry options (variable);

/\* PICK UP ABSENTEE'S PROCESS AND CHANNEL IDS \*/

call hcs\_\$initiate (">udd>F15dw>Auerbach", "ipc\_seg1", "",  
0, 1, seg\_ptr, code);  
if seg\_ptr = null () then call ERROR;

/\* TELL ABSENTEE TO PRINT AND GIVE IT A PATH TO PRINT \*/

dprint\_paths\$ = ">udd>F15dw>Auerbach>abs\_print\_punch.pl1";  
message = 1;

/\* FIRE OFF WAKEUP SIGNAL \*/

call hcs\_\$wakeup (ipc\_info.his\_process\_id,  
ipc\_info.his\_chnl\_id,  
message,  
code);

if code ^= 0 then call ERROR;

call ioa\_ ("^[DPRINT^;DPUNCH^] REQUEST FOR ^a",  
message, dprint\_paths\$);

call ioa\_ ("END ^a", ME);

ERROR: proc;

call com\_err\_ (code, ME);  
goto FINIS;

end;

FINIS:

end driver\_ipc;

r 11:12 0.020 2

IPC SUBROUTINES

AN EXAMPLE USING EVENT-WAIT CHANNELS

```
!print ipc_example.absin 1
```

```
abs_print_punch  
logout  
&quit
```

```
r 11:12 0.020 1
```

```
!ear ipc example  
27 already requested.  
r 10:55 0.113 6
```

```
!driver ipc  
DPRINT REQUEST FOR >udd>F15dw>Auerbach>abs_print_punch.pl1  
END driver ipc  
r 11:13 0.068 4
```

```
!ldr -long
```

```
Queue 3: 1 request. 12 total requests.
```

```
Pathname: >udd>F15dw>Auerbach>abs_print_punch.pl1  
Type: print  
Copies: 1  
Time: 03/13/80 10:56 mst Thur  
Delete: no  
Notify: yes  
Options: -indent 10
```

```
r 11:13 0.138 10
```

IPC SUBROUTINES

AN EXAMPLE USING EVENT-WAIT CHANNELS

!print ipc\_example.absout 1

Absentee user Auerbach F15dw logged in: 03/13/80 11:14 mst Thur  
r 11:14 1.827 34

abs\_print\_punch

CHANNEL ID 98439864967772869469

MESSAGE 1

SENDER 004670305344

DEV SIGNAL 000000

RING 000004

CHANNEL INDEX 1

>udd>F15dw>Auerbach>abs\_print\_punch.pl1 DPRINT REQUEST SUBMITTED.

END abs\_print\_punch

r 11:15 2.314 27

logout

Absentee user Auerbach F15dw logged out 03/13/80 11:16 mst Thur  
CPU usage 2 sec, memory usage 1.0 units

r 11:17 0.084 14

From IO.SysDaemon (printer):

printed >udd>F15dw>Auerbach>abs\_print\_punch.pl1

\$0.06 queue 3 prtd 10115

IPC SUBROUTINES

AN EXAMPLE USING EVENT-CALL CHANNELS

```
!print listen.pl1 1

listen: proc;
dcl ipc $create_ev_chn entry (fixed (71), fixed (35)),
    get_process_id_entry returns (bit (36)),
    hcs $wakeup_entry (bit (36), fixed bin (71), fixed (71), fixed (35)),
    iox $control_entry (ptr, char (*), ptr, fixed (35)),
    code fixed (35),
    ipc $decl_ev_call_chn entry (fixed (71), entry, ptr, fixed,
                                fixed (35)),
    hcs $initiate_count_entry (char (*), char (*), char (*), fixed (24),
    fixed (2), ptr, fixed (35)),
    bc fixed (24),
    iox $user_output_ext ptr,
    iox $put_chars_entry (ptr, ptr, fixed (21), fixed (35)),
    (ioa_, com_err_) entry options (variable);
dcl 1 ipc_info_based (seg_ptr),
    2 process_id bit (36),
    2 channel_id fixed bin (71);
dcl seg_ptr ptr;
dcl ME_char (6) init ("listen") static options (constant);
dcl event_info_ptr ptr;
dcl 1 event_info_based (event_info_ptr),
    2 channel_id fixed bin (71),
    2 message fixed bin (71),
    2 sender bit (36),
    2 origin,
    3 dev_signal bit (18) unal,
    3 ring bit (18) unal,
    2 data_ptr ptr;

/* INITIATE PROTOCOL PASSING SEGMENT */
    call hcs $initiate_count (">udd>F15dw>Auerbach", "ipc_seg1",
        "", bc, 1, seg_ptr, code);
    if seg_ptr = null ( ) then call ERROR;

/* CREATE EVENT CHANNEL AND MAKE IT A CALL CHANNEL */
    call ipc $create_ev_chn (ipc_info.channel_id, code);
    if code ^= 0 then call ERROR;
    ipc_info.process_id = get_process_id ( );
    call ipc $decl_ev_call_chn (ipc_info.channel_id,
        print_msg, null ( ), 0, code); /* data_ptr = null */
    if code ^= 0 then call ERROR;

/* ALL DONE */
    call ioa_ ("Now listening for messages.");
    return;
```

IPC SUBROUTINES

AN EXAMPLE USING EVENT-CALL CHANNELS

```
print msg: entry (event info ptr);
/* THIS HANDLER FOR TAKING MESSAGES DOES NOT USE THE
event_info_ptr PARAMETER PASSED TO IT. */

/* GET POINTER TO MAILBOX SEGMENT */
call hes_$initiate_count (>udd>F15dw>Auerbach", "mailbox",
    "", bc, 1, seg_ptr, code);
if seg_ptr = null () then call ERROR;

/* PRINT OUT CONTENTS OF >udd>F15dw>Auerbach>mailbox */
call ioa_ ("Message is:");
call iox_$put_chars (iox_$user_output, seg_ptr,
    divide (bc, 9, 21, 0), code);
if code ^= 0 then call ERROR;

/* RESTART ANY INTERRUPTED OUTPUT AND RETURN */
call iox_$control (iox_$user_output, "start", null (),
    code);

ERROR:   proc;
        call com_err_ (code, ME);
        call iox_$control (iox_$user_output, "start", null (),
            code);
        goto FINIS;
    end;

FINIS:
    end listen;
```

r 12:36 0.383 18

!print put\_message.pl1 1

put\_message: proc;

```
%include listen decls;
dcl mailbox file;
dcl 1 ipc_info based (seg_ptr),
    2 process_id bit (36),
    2 channel_id fixed bin (71);
dcl seg_ptr ptr;
```

```
/* THE FOLLOWING PL/1 I/O STATEMENTS BUILD UP A SEGMENT
WHICH WILL BE DUMPED LATER BY THE print_msg EVENT HANDLER. */
```

```
open file (mailbox) stream output;
put file (mailbox) skip
    list ("Hello...this is the absentee process..");
put file (mailbox) skip
    list ("Just wanted to prove it works!!");
```



IPC SUBROUTINES

AN EXAMPLE USING EVENT-CALL CHANNELS

```
put file (mailbox) skip;
close file (mailbox);

/* OBTAIN process id AND channel id SO THAT WE CAN SEND WAKEUP */
call hcs $initiate_count (">udd>F15dw>Auerbach", "ipc_seg1",
    "", bc, 1, seg_ptr, code);
if seg_ptr = null ( ) then call ERROR;

/* NOW SEND THE WAKEUP */
call hcs $wakeup (ipc_info.process_id, ipc_info.channel_id,
    0, code);
if code ^= 0 then call ERROR;

ERROR:   proc;
del ME char (11) init ("put_message") static options (constant);
        call com_err (code, ME);
        goto FINISH;
end;
```

```
FINISH:
end put_message;
```

```
r 12:36 0.020 1
```

```
!print pm.absin 1
```

```
put_message
logout
&quit
```

```
r 12:36 0.029 1
```

```
!defer messages
r 12:36 0.018 5
```

```
!listen
Now listening for messages.
r 12:36 0.015 3
```

```
!ear pm
27 already requested.
r 12:36 0.126 12
```

IPC SUBROUTINES

AN EXAMPLE USING EVENT-CALL CHANNELS

!who

Multics MR8.0, load 55.0/130.0; 55 users  
Absentee users 0/3

IO.SysDaemon  
Backup.SysDaemon  
IO.SysDaemon  
IO.SysDaemon  
IO.SysDaemon  
GCOS.SysDaemon  
Volume\_Dumper.Daemon  
Opr.Operator  
MFreeman.SSF  
Susan.NCB  
Jagernauth.Multics  
Irish.Doc  
Downing.Multics  
Wardd.Multics  
Falksenj.Multics  
Coppola.HFED  
Casselman.HCRC  
Martinson.SysMaint  
Nolde.Bus-Plan  
FED.VIS  
Lombreglia.NCB  
Lutz.GSASched  
Retriever.SysDaemon  
Harrison.Rapidata  
Stryk.HCRC  
Matheson.DEBUG  
Landrum.SED  
Baryza.FORD CONV  
Auerbach.F15dw  
Johnson.SysAdmin  
Donner.Multics  
Fawcett.VIS  
Message is:

Hello...this is the absentee process..  
Just wanted to prove it works!!

Coflin.G66  
Friedman.F15aw  
Glicksman.HISCAN  
Berglund.Multics  
Arnwine.SiteSA  
NThompson.NOPS  
Chouinard.BBbench  
JWilliams.SED  
Student\_20.F01  
RBarnes.Multics

IPC SUBROUTINES

AN EXAMPLE USING EVENT-CALL CHANNELS

Bergum.HCRC  
Gildersleeve.Multics  
Tilton.MMPP  
Watts.Doc  
FED.VIS  
Sam.SRB  
Student 14.F01  
Gowans.ESask  
PHJones.BBbench  
Troost.MMPP  
Whitford.Doc  
Gintell.Multics  
Grimes.SMP

IPC SUBROUTINES  
AN EXAMPLE USING EVENT-CALL CHANNELS

|| YOU ARE NOW READY FOR WORKSHOP ||  
#4

TOPIC IX

Interprocess Data Base Sharing

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## INTRODUCTION

### ● CONCEPT

- ¶ ANOTHER FORM OF INTERPROCESS COMMUNICATION INVOLVES THE SHARING OF COMMON DATA BASE (NOT TO BE CONFUSED WITH MDBM) SEGMENTS
- ¶ BECAUSE SEVERAL PROCESSES MAY BE ATTEMPTING TO CONCURRENTLY ACCESS AND UPDATE A SHARED DATA BASE, SOME FORM OF CONTROL IS REQUIRED:
  - ¶ TO PREVENT THE DATA BASE FROM BEING LEFT IN AN INCONSISTENT STATE
  - ¶ TO PREVENT ANY PROCESS FROM OPERATING UPON PARTIALLY-UPDATED DATA

### ● A POSSIBLE SOLUTION

- ¶ A "LOCKING" FACILITY: THE `set_lock_` SUBROUTINE
- ¶ COOPERATING PROCESSES OBSERVE A LOCKING PROTOCOL, IN WHICH A GIVEN PROCESS DOES NOT ACCESS A DATA BASE UNTIL IT CAN SET A LOCK WORD, AND IN WHICH A PROCESS RESETS THAT LOCK WORD WHEN IT HAS COMPLETED A CRITICAL UPDATE OR RETRIEVAL

## THE LOCKING MECHANISM

- DATA BASE LOCKING IS IMPLEMENTED BY `set lock` SUBROUTINE (AG93)  
(BUT SEE ALSO THE PL/1 NON-STANDARD BUILTINS `stac` AND `stacq`)
  
- `set_lock` PROTOCOL
  - ▮ A CALLER-SUPPLIED LOCK WORD IS USED FOR THE MUTUAL EXCLUSION OF PROCESSES
  
  - ▮ THIS LOCK WORD IS
    - ▮ DECLARED (bit(36) aligned) BY UPDATING PROGRAM(S)
  
    - ▮ ZEROED ONCE BY SOME SPECIAL INITIALIZATION PROGRAM WHEN DATABASE FIRST CREATED (THUS INDICATING IT'S UNLOCKED)
  
    - ▮ LOCATED
      - ▮ USUALLY IN A SPECIAL, SEPARATE SEGMENT ACCESSIBLE TO ALL COOPERATING PROCESSES
  
      - ▮ SOMETIMES IN THE DATABASE ITSELF

## THE LOCKING MECHANISM

- ¶ WHEN A PROCEDURE IS ABOUT TO ENTER A CRITICAL SECTION OF CODE, IT CALLS AN ENTRY POINT IN `set_lock` WHICH ATTEMPTS TO SET THE LOCK BY PLACING THAT PROCESS'S LOCK IDENTIFIER IN THE LOCK WORD USING AN INDIVISIBLE MACHINE INSTRUCTION, `stac` (STORE-A CONDITIONAL)
  
- ¶ `stac` USES A SPECIAL MAIN MEMORY REFERENCE THAT PROHIBITS SUCH REFERENCES BY OTHER PROCESSES BETWEEN THE TEST AND THE DATA TRANSFER.
  
- ¶ IF THE LOCK WORD ALREADY CONTAINS SOME OTHER VALID LOCK IDENTIFIER, THE INTERPRETATION IS THAT THE DATA BASE IS LOCKED BY THAT OTHER PROCESS, AND THE CALLING PROCESS WAITS FOR THE LOCK TO BE UNLOCKED BY THAT OTHER PROCESS
  
- ¶ WHEN A CRITICAL SECTION OF CODE HAS BEEN COMPLETED BY THE PROGRAM, THE LOCK OUGHT TO BE UNLOCKED, ALLOWING ANOTHER PROCESS TO SET THE LOCK

### ● SUCCESS HINGES ON THE FOLLOWING CONVENTIONS:

- ¶ `set_lock` IS THE ONLY PROCEDURE THAT MAY MODIFY THE LOCK WORD (WITH THE EXCEPTION OF THE PROGRAM WHICH INITIALIZES THE DATA BASE AND LOCK WORD)
  
- ¶ ALL PROCESSES SHOULD CALL `set_lock_$lock` BEFORE ENTERING A CRITICAL SECTION OF CODE
  
- ¶ ALL PROCESSES SHOULD CALL `set_lock_$unlock` AFTER COMPLETING A CRITICAL SECTION OF CODE



THE SET LOCK SUBROUTINE

● set\_lock\_\$lock (AG93)

⌋ ATTEMPTS TO PLACE LOCK IDENTIFIER OF CALLING PROCESS IN THE GIVEN LOCK WORD

⌋ call set\_lock\_\$lock (lock\_word, wait\_time, code);

⌋ wait time INDICATES THE NUMBER OF SECONDS THAT set lock \$lock SHOULD WAIT FOR A VALIDLY LOCKED LOCK WORD TO BE UNLOCKED BEFORE RETURNING UNSUCCESSFULLY (-1 INDICATES NO TIME LIMIT)

⌋ ONE OF THE FOLLOWING CODES IS RETURNED:

⌋ 0

⌋ error\_table\_\$invalid\_lock\_reset

⌋ error\_table\_\$locked\_by\_this\_process

⌋ error\_table\_\$lock\_wait\_time\_exceeded

THE SET LOCK SUBROUTINE

● set\_lock\_unlock (AG93)

⌋ ATTEMPTS TO RESET A GIVEN LOCK WORD TO "0"

⌋ call set\_lock\_unlock (lock\_word, code);

⌋ RETURNS ONE OF THE FOLLOWING CODES:

⌋ 0

⌋ error\_table\_lock\_not\_locked

⌋ error\_table\_locked\_by\_other\_process IF lock\_word CONTAINED  
NON-ZERO VALUE NOT EQUAL TO LOCK IDENTIFIER OF THE CALLING  
PROCESS

AN EXAMPLE OF LOCKING

```
!pwd
>user_dir_dir>F15dw>Auerbach
r 92:9 0.173 1

!ls -pn >udd>F15dw>Auerbach>AJAX_db
```

Segments = 2, Lengths = 2.

```
r w 0 lock_word
re 2 book_seat
```

Multisegment-files = 1, Lengths = 7.

```
r w 7 flight_records
```

```
r 09:29 0.064 0
```

```
!print book_seat.pl1 1
```

```
book_seat: procedure;
```

```
/* THIS PROGRAM UPDATES AN AIRLINES DATABASE
WHICH IS LOCATED IN THE SAME DIRECTORY
AS THIS PROGRAM, AND WHICH ALSO CONTAINS
A LOCK SEGMENT */
```

```
dcl flight_records file;
dcl 1 flight_rec based (rec_ptr),
    2 total_seats fixed bin,
    2 seats_booked fixed bin,
    2 seat_info (0 refer (seats_booked)),
        3 name char (20) varying,
        3 address char (30) varying;
dcl flight_no char (4) varying,
    date char (6) varying;
dcl rec_ptr ptr;

dcl set_lock $lock entry (bit (36) aligned,
    fixed bin, fixed bin (35)),
    set_unlock $unlock entry (bit (36) aligned,
    fixed bin (35)),
    (ioa, com_err) entry options (variable),
    change_wdir_entry (char (168) aligned, fixed bin (35));

dcl lock_wordsegment bit (36) aligned external static;
dcl (key, cleanup) condition;

dcl (code,
    error_table $invalid_lock_reset external,
    error_table $locked_by_this_process external,
    error_table $lock_wait_time_exceeded external) fixed bin (35);

dcl ME char (9) init ("book_seat") static options (constant);
```

AN EXAMPLE OF LOCKING

```
/* ESTABLISH 'on unit' FOR 'cleanup'
DO NOT UNLOCK LOCK - DATA BASE MAY BE
IN AN INCONSISTENT STATE */
    on cleanup close file (flight_records);

/* ESTABLISH 'on unit' FOR 'key' CONDITION
REPORT ERROR AND ASK AGAIN */
    on key (flight_records) begin;
        call ioa_ ("Invalid key entered ^a", onkey ());
        goto PROMPT;
    end;

/* BEGIN UPDATE PROGRAM */
    call ioa_ ("AJAX Airlines flight booking program begins");

/* CHANGE WORKING DIR TO AIRLINES DATABASE DIRECTORY */
    call change_wdir_ (">udd>F15dw>Auerbach>AJAX_db", code);
    if code ^= 0 then do;
        call com_err_ (code, ME);
        return;
    end;

/* OPEN DATABASE */
    open file (flight_records) direct update;

/* LOCK DATABASE NOW - TRY FOR 30 SECONDS */
    call set_lock_$lock (lock_word$, 30, code);

    if code ^= 0 then
/* COULDN'T LOCK IT - FIND OUT WHY */
        if code = error_table_$lock_wait_time_exceeded
            then do;
/* DATABASE IS BUSY */
            call ioa_ ("Database busy - try again later.");
            goto WRAPUP;
        end;
        else
            if code = error_table_$invalid_lock_reset
                then do;
/* SOMEBODY DIDN'T UNLOCK BEFORE DYING */
                call ioa_ ("Database has invalid lock");
                call ioa_ ("Notify DBA - no update allowed");
                goto WRAPUP;
            end;
            else
                if code = error_table_$locked_by_this_process then do;
```

AN EXAMPLE OF LOCKING

```
/* SOMETHING IS VERY WRONG - DIE */
    call ioa ("FATAL ERROR!!");
    call ioa ("NOTIFY DBA IMMEDIATELY!!");
    goto WRAPUP;
    end;
    else;
else;

/* DATABASE IS NOW LOCKED */

PROMPT:
/* BASIC REQUEST LOOP */

    do while ("1"b);
        call ioa ("Enter flight no, date for booking");
        get list (flight no, date);
        if flight no = "0" then goto WRAPUP;
/* STOP WHEN flight no IS "0" */
/* TRY TO READ RECORD */
        read file (flight records) key (flight_no||date)
        set (rec_ptr);

/* SEE IF ANY SEATS ARE LEFT */
        if rec_ptr -> seats_booked >= rec_ptr -> total_seats
        then do;
            call ioa ("Flight is booked full.");
        end;

/* OKAY - GET REST OF INFO */
        else do;
            seats_booked = seats_booked + 1;
            call ioa ("Enter name, address of cust");
            get list (
                flight_rec.seat_info (seats_booked).name,
                flight_rec.seat_info (seats_booked).address);
            rewrite file (flight records)
                from (rec_ptr -> flight_rec);
        end;
    end;

WRAPUP:
/* UNLOCK AND CLOSE DATABASE */
    call set lock $unlock (lock word$, code);
    close file (flight records);
    call ioa ("End Update Program ^a", ME);

    end book_seat;
```

AN EXAMPLE OF LOCKING

r 09:29 0.025 2

!AJAX\_db>book\_seat

AJAX Airlines flight booking program begins

Enter flight\_no, date for booking

!112,800303

Flight is booked full.

Enter flight\_no, date for booking

!0,0

End Update Program book\_seat

TOPIC X

Intraprocess Timer Management

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## INTRODUCTION

- CERTAIN SOPHISTICATED PROGRAMS MAY REQUIRE THE USE OF ONE OR MORE CPU AND/OR REAL-TIME TIMERS
- TIMER MANAGER FACILITY (TMF) OPERATES PRIMARILY IN CONJUNCTION WITH THE INTERPROCESS COMMUNICATION FACILITY, THUS ENABLING PROGRAMS TO RUN ASYNCHRONOUSLY WITHIN A PROCESS

### TIMER - MANAGER

- ~~Multi~~cs ALLOWS A PROCESS TO:
  - ▮ BLOCK ITSELF FOR A SPECIFIED REAL TIME PERIOD (SLEEP)
  - ▮ CALL A SPECIFIED PROCEDURE WHEN A SPECIFIED TIME INTERVAL HAS ELAPSED
  - ▮ ISSUE A WAKEUP ON A SPECIFIED EVENT-WAIT CHANNEL WHEN A SPECIFIED TIME INTERVAL HAS ELAPSED
- THE SUBROUTINE INTERFACE IS timer\_manager\_ (AK92)
- FOR MORE ABOUT TIME, SEE THE MPM SUBROUTINES MANUAL:


clock	cpu_time_and_paging_	date_time
decode_clock_value_		virtual_cpu_time_



## TIMER MANAGEMENT TERMINOLOGY

### ● timer\_manager\_ MAKES USE OF SEVERAL CRITICAL CONCEPTS:

#### ▮ alarm

- ▮ "alarm" IS USED TO DESIGNATE A REAL-TIME TIMER; THAT IS, A "WALL-TIME" ELAPSED TIME - WHEN AN "alarm" TIMER GOES OFF, THE "alm" (STATIC) CONDITION IS SIGNALLED

#### ▮ cpu

- ▮ "cpu" IS USED TO DESIGNATE A VIRTUAL CPU TIMER; WHEN A "cpu" TIMER GOES OFF, THE "cput" (STATIC) CONDITION IS SIGNALLED

#### ▮ relative time

- ▮ A TIME MEASURED FROM THE CALL TO timer\_manager\_ ; THAT IS, A TIME MEASURED FROM THE TIME THE TIMER IS CREATED

#### ▮ absolute time

- ▮ A TIME MEASURED FROM THE FIXED POINT IN TIME "January 1, 1901 0000 Hours"

## TIMER MANAGER GENERIC ARGUMENTS

### ● timer\_manager\_ ENTRY POINTS ACCEPT A COMMON SET OF GENERIC ARGUMENTS

#### ▮ channel

- ▮ THE EVENT CHANNEL ID (fixed bin(71)) OVER WHICH A WAKEUP IS TO BE TRANSMITTED
- ▮ SET UP PRIOR TO INVOCATION OF A timer\_manager\_ ENTRY POINT

#### ▮ routine

- ▮ THE PROCEDURE TO BE INVOKED WHEN A "CALL" TIMER GOES OFF (SPECIFIED WHEN THE TIMER IS CREATED)
- ▮ THE PROCEDURE WILL BE PASSED TWO ARGUMENTS (UNLIKE AN EVENT-CALL PROCEDURE) AS FOLLOWS:
  - ▮ mc\_ptr
    - ▮ AN ALIGNED POINTER TO THE "MACHINE CONDITIONS" AT THE TIME THE alarm OR cput CONDITION WAS SIGNALLED (SEE SECTION 7 IN MPM REFERENCE GUIDE)
  - ▮ name
    - ▮ A CHARACTER STRING INDICATING WHETHER THE TIMER WAS AN ALARM TIMER (alarm) OR A CPU TIMER (cput)
- ▮ IS MOST OFTEN AN EXTERNAL ENTRY, BUT MIGHT BE INTERNAL (TAKE CARE!)

## TIMER MANAGER GENERIC ARGUMENTS

### ▮ time

- ▮ MANY timer\_manager ENTRY POINTS REQUIRE THAT THE TIME (fixed bin (71)) BE SPECIFIED; alm OR cput CONDITION IS SIGNALLED AT THAT TIME

### ▮ flags

- ▮ MANY timer\_manager ENTRY POINTS REQUIRE THIS bit(2) STRING, WHICH SPECIFIES HOW THE time ARGUMENT IS TO BE INTERPRETED
  - ▮ "11"b MEANS RELATIVE SECONDS
  - ▮ "10"b MEANS RELATIVE MICROSECONDS (1e-6 SECONDS)
  - ▮ "01"b MEANS ABSOLUTE SECONDS
  - ▮ "00"b MEANS ABSOLUTE MICROSECONDS (1e-6 SECONDS)

TIMER MANAGER ENTRY POINTS

● timer\_manager\_ ENTRY POINTS ALLOW A PROCESS TO:

▮ BLOCK A PROCESS FOR A SPECIFIED REAL TIME INTERVAL

▮ timer\_manager\_\$sleep

▮ CAUSE A SPECIFIED PROCEDURE TO BE INVOKED AT A SPECIFIED TIME

▮ timer\_manager\_\$alarm\_call

▮ timer\_manager\_\$cpu\_call

▮ CAUSE A WAKEUP TO BE ISSUED ON A SPECIFIED EVENT-WAIT CHANNEL AT A SPECIFIED TIME

▮ timer\_manager\_\$alarm\_wakeup

▮ timer\_manager\_\$cpu\_wakeup

## TIMER MANAGER ENTRY POINTS

### ▮ RESET AND INHIBIT TIMERS

▮ timer\_manager\_\$alarm\_call\_inhibit

▮ timer\_manager\_\$reset\_alarm\_call

▮ timer\_manager\_\$reset\_alarm\_wakeup

▮ timer\_manager\_\$cpu\_call\_inhibit

▮ timer\_manager\_\$reset\_cpu\_call

▮ timer\_manager\_\$reset\_cpu\_wakeup

TIMER MANAGER ENTRY POINTS

BLOCKING A PROCESS

- timer\_manager\_\$sleep CAUSES PROCESS TO GO BLOCKED FOR A PERIOD OF REAL TIME

▮ OTHER TIMERS THAT ARE ACTIVE ARE PROCESSED WHENEVER THEY GO OFF

▮ HOWEVER, THE PROCEDURE ISSUING THIS CALL WILL NOT RESUME (I.E., EXECUTE ITS NEXT INSTRUCTION) UNTIL THE REAL TIME HAS BEEN PASSED

▮ EXAMPLE

```

                                -----SECONDS
                                |
                                v
call timer_manager_$sleep ( 30, "11"b);
                                ^
                                |
                                RELATIVE-----
```

▮ WOULD CAUSE THIS PROCESS TO GO TO "SLEEP" FOR THIRTY SECONDS

TIMER MANAGER ENTRY POINTS

USING CALL TIMERS

- ENTRY POINTS WHICH CAUSE A SPECIFIED PROCEDURE TO BE INVOKED WHEN A TIMER GOES OFF

┆ timer\_manager\_\$alarm\_call

┆ SETS UP A REAL-TIMER

┆ A SPECIFIED ROUTINE IS CALLED WHEN THE TIMER GOES OFF

┆ IT REQUIRES THE time, flags, AND routine ARGUMENTS AS INPUT

┆ EXAMPLE

```
call timer_manager_$alarm_call ( 80, "11"b, print_usage);
```

┆ WOULD CAUSE A PROCEDURE CALLED print\_usage TO BE INVOKED AFTER 80 SECONDS OF REAL TIME HAD ELAPSED

TIMER MANAGER ENTRY POINTS

USING CALL TIMERS

▮ timer\_manager\_\$cpu\_call

▮ SETS UP A CPU TIMER WHICH WILL CAUSE A SPECIFIED PROCEDURE TO BE INVOKED WHEN A SPECIFIED INTERVAL OF CPU TIME HAS ELAPSED

▮ REQUIRES THE SAME ARGUMENTS AS THE timer\_manager\_\$alarm\_call ENTRY POINT

▮ EXAMPLE

```

                                     -----MICROSECONDS
                                     |
                                     |
                                     v
call timer_manager_$cpu_call ( 1000, "10"b, print_cpu_usage);
                                     |
                                     |
                                     ^
                                     |
                                     -----
                                     RELATIVE-----
```

▮ CAUSES THE PROGRAM print cpu usage TO BE INVOKED WHEN ONE MILLISECOND OF CPU TIME HAS ELAPSED



TIMER MANAGER ENTRY POINTS

USING WAKEUP TIMERS

- ENTRY POINTS BELOW ALLOW THE CALLER TO SPECIFY THAT A WAKEUP IS TO BE SENT ACROSS SPECIFIED EVENT-WAIT CHANNELS WHEN THE TIMER GOES OFF

┆ timer\_manager\_\$alarm\_wakeup

┆ SETS UP A REAL-TIME TIMER THAT ISSUES A WAKEUP ON THE EVENT-WAIT CHANNEL SPECIFIED WHEN THE TIMER GOES OFF

┆ CALLER MAY WISH TO GO BLOCKED ON THE CHANNEL (BUT NEEDN'T)

┆ THE EVENT MESSAGE PASSED IS THE STRING "alarm\_\_\_\_"

┆ REQUIRES THREE INPUT ARGUMENTS - THE time, flags, AND channel\_id

┆ EXAMPLE

```
call convert_date_to_binary_ ("1 hour 5 minutes", time, code);
call timer_manager_$alarm_wakeup ( time, "00"b, channel_id);
```

ABSOLUTE-----  
┆  
-----MICROSECONDS

┆ WOULD CAUSE A WAKEUP TO BE ISSUED ACROSS channel id 65 WALL CLOCK MINUTES FROM THE TIME convert\_date\_to\_binary\_ WAS CALLED

TIMER MANAGER ENTRY POINTS

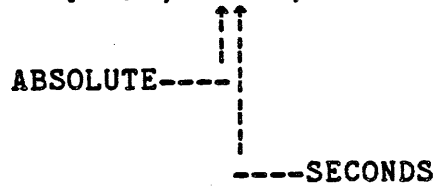
USING WAKEUP TIMERS

0 timer\_manager\_\$cpu\_wakeup

0 OPERATES EXACTLY LIKE timer\_manager \$alarm wakeup EXCEPT THAT THE TIMER IS A CPU TIMER AND THE EVENT MESSAGE IS "cpu\_time"

0 EXAMPLE

call timer\_manager\_\$cpu\_wakeup (50, "01"b, channel\_id);



0 WOULD CAUSE A WAKEUP TO BE ISSUED ON THE EVENT-WAIT CHANNEL IDENTIFIED BY channel id AFTER A TOTAL OF 50 CPU SECONDS HAD BEEN EXPENDED (MEASURED FROM PROCESS CREATION TIME)

TIMER MANAGER ENTRY POINTS  
RESETTING AND INHIBITING TIMERS

- timer manager ENTRY POINTS EXIST WHICH ALLOW A PROCESS TO RESET OR INHIBIT TIMERS

- USERS OF timer manager SHOULD BE AWARE OF THE PERILS OF ASYNCHRONOUS PROCESSING, AND PROGRAMS CREATING CALL OR WAKEUP TIMERS SHOULD PROVIDE AN 'on unit' FOR THE cleanup CONDITION TO RESET TIMERS

□ SUCH 'on units' GENERALLY DO NOTHING MORE THAN RESET THE TIMERS, THUS PREVENTING SUCH TIMERS FROM GOING OFF AT UNDESIRED TIMES

□ ENTRY POINTS WHICH RESET AND INHIBIT TIMERS:

□ timer\_manager\_\$reset\_alarm\_call

□ RESETS, OR TURNS OFF ALL REAL-TIME TIMERS (alarm) THAT CALL THE ROUTINE SPECIFIED WHEN THEY GO OFF

□ REQUIRES ONLY ONE ARGUMENT, THE NAME OF THE ROUTINE FOR WHICH TIMERS HAVE BEEN SET

□ EXAMPLE

```
call timer_manager_$reset_alarm_call (print_usage);
```

□ THIS CALL WOULD TURN OFF ANY REAL-TIME TIMERS WHICH WOULD CALL print\_usage IF THEY WENT OFF

TIMER MANAGER ENTRY POINTS  
RESETTING AND INHIBITING TIMERS

⌋ timer\_manager\_\$reset\_cpu\_call

⌋ OPERATES IN THE SAME MANNER AS timer\_manager\_\$reset\_alarm\_call EXCEPT THAT IT TURNS OFF CPU TIMERS FOR THE SPECIFIED PROCEDURE

⌋ timer\_manager\_\$reset\_alarm\_wakeup

⌋ TURNS OFF ALL REAL-TIME TIMERS THAT ISSUE A WAKEUP ON THE EVENT-WAIT CHANNEL SPECIFIED

⌋ THE ONLY INPUT ARGUMENT IS THE channel\_id OF THE CHANNEL FOR WHICH TIMER WAKEUPS ARE TO BE TURNED OFF

⌋ EXAMPLE

```
call timer_manager_$reset_alarm_wakeup (channel_id);
```

⌋ THIS CALL WOULD TURN OFF ANY REAL-TIME TIMERS WHICH WOULD OTHERWISE ISSUE A WAKEUP ON THE CHANNEL IDENTIFIED BY channel\_id WHEN THEY WENT OFF

⌋ timer\_manager\_\$reset\_cpu\_wakeup

⌋ OPERATES EXACTLY LIKE timer\_manager\_\$reset\_alarm\_wakeup EXCEPT THAT IT TURNS OFF CPU TIMERS FOR THE EVENT-WAIT CHANNEL SPECIFIED

TIMER MANAGER ENTRY POINTS  
RESETTING AND INHIBITING TIMERS

- ▮ INHIBITING INTERRUPTS WHILE USING CALL TIMERS ALLOWS THE PROCESS TO ENSURE THAT THE HANDLER (THE PROCEDURE INVOKED) WILL NOT BE INTERRUPTED BEFORE IT RETURNS
  
- ▮ IF SUCH HANDLERS DO NOT RETURN, THE PROCESS MAY MALFUNCTION, SINCE IT IS DANGEROUS TO INHIBIT INTERRUPTS FOR TOO LONG
  
- ▮ timer\_manager\_\$alarm\_call\_inhibit
  - ▮ OPERATES EXACTLY LIKE timer\_manager\_\$alarm\_call EXCEPT THAT ALL INTERRUPTS ARE INHIBITED JUST BEFORE THE HANDLER IS INVOKED
  
  - ▮ WHEN THE HANDLER RETURNS, ALL INTERRUPTS ARE REENABLED
  
- ▮ timer\_manager\_\$cpu\_call\_inhibit
  - ▮ OPERATES LIKE timer\_manager\_\$cpu\_call EXCEPT THAT ALL INTERRUPTS ARE INHIBITED WHILE THE HANDLER IS EXECUTING

TIMER MANAGER ENTRY POINTS

STANDARD SYSTEM HANDLERS

- OTHER ENTRY POINTS IN `timer_manager_` SERVE AS STATIC HANDLERS FOR TWO CONDITIONS:

- ▮ `timer_manager_$cpu_time_interrupt` (FOR 'cput' CONDITION)

- ▮ `timer_manager_$alarm_interrupt` (FOR 'alarm' CONDITION)

- SEE THE CODE FOR `user_real_init_admin_` IN APPENDIX B

TWO EXAMPLES USING TIMERS

!print cookie.pl1 1

cookie: proc;

/\* THE INFAMOUS COOKIE MONSTER PROGRAM  
THIS PROGRAM USES THE TIMER MANAGER FACILITY  
TO CAUSE THE COOKIE MONSTER PROGRAM TO BE  
REINVOKED AGAIN AND AGAIN, THUS BEWILDERING  
THE CALLER \*/

✓dcl timer\_manager\_\$alarm\_call entry(fixed bin (71), bit (2), entry),  
↳ timer\_manager\_\$reset\_alarm\_call entry (entry),  
iox\_\$control entry (ptr, char (\*), ptr, fixed bin (35)),  
(ioa\_, ioa\_\$nml) entry options (variable);

dcl (quit, cleanup) condition;  
dcl next\_time fixed bin (71);  
dcl null\_builtin;  
dcl iox\_\$user\_io ext ptr;  
dcl answer char (6);  
dcl sysin file;  
dcl code fixed bin (35);  
dcl i static init (1);  
dcl first\_time bit (1) static init ("1"b);

/\* ESTABLISH 'on unit' FOR CLEANUP  
RESET ANY TIMERS THAT MAY BE AROUND \*/  
↳ on cleanup call timer\_manager\_\$reset\_alarm\_call (handler);

/\* COME HERE IF FIRST TIME EXECUTING COOKIE \*/  
if first\_time then do;

/\* FOR INITIAL CALL, SET TIME TO 10 SECONDS \*/  
next\_time = 10;  
first\_time = "0"b;

end;

/\* WHEN COOKIE IS CALLED AGAIN,  
SET THE TIME TO 80 SECONDS \*/  
else next\_time = 80;

/\* NOW SET UP TIMER TO CALL handler AND RETURN  
IT WILL SEEM TO THE CALLER THAT ALL IS NORMAL \*/  
↳ call timer\_manager\_\$alarm\_call (next\_time, "1"b, handler);  
return;

TWO EXAMPLES USING TIMERS

```
/* COME HERE WHEN TIMER GOES OFF */
handler: entry;
/* ESTABLISH 'on unit' TO TRAP HIS QUILTS */
  on quit begin;
    call ioa_ ("QUIT");
    call ioa_ ("You wanna get yourself logged out??");
    i = i+1;
    go to ask_again;
  end;

/* BASIC PROMPTING LOOP
WE READ INPUT LOOKING FOR COOKIES */
ask_again:
  if i = 1 then call ioa_ ("I want a cookie.");
  else if i = 2 then call ioa_ ("Please give me a cookie.");
  else if i = 3 then call ioa_ (
    "You had better give me a cookie.");
  else if i = 4 then call ioa_ ("I WANT A COOKIE!");
  else if i = 5 then call ioa_ (
    "PLEASE--GIMME A COOKIE!!!!");
  else do;
    call ioa_ ("I give up. You're hopeless.");
    call ioa_$nnl ("Guess I'll have to get one myself.");
    call ioa_ ("COOKIECOOKIECOOKIE");
  end;

/* AT THIS POINT I EITHER GOT A COOKIE FROM HIM
OR FROM MYSELF - RESET I AND CALL COOKIE AGAIN */
next_set:
  i = 1;
  call cookie;
  call iox_$control (iox_$user_io, "start", null, code);
  return;
end;

/* HERE'S WHERE I GRAB HIS INPUT LINES */
get list (answer);
if answer = "cookie" | answer = "COOKIE" then do;
  call ioa_ ("Thanks. I needed that. YumYumYum...");
  go to next_set;
end;
else do;
  i = i+1;
  go to ask_again;
end;
end;
```



TWO EXAMPLES USING TIMERS

!cookie  
r 11:56 0.065 1

!hmu

Multics MR8.0, load 63.0/100.0; 63 users  
Absentee users 0/5

r 11:56 0.058 0

I want a cookie.

!who

Please give me a cookie.

!(QUIT)

QUIT

You wanna get yourself logged out??

You had better give me a cookie.

!new proc

PLEASE--GIMME A COOKIE!!!!

!logout

I give up. You're hopeless.

Guess I'll have to get one myself.COOKIECOOKIECOOKIE

!new proc

r 11:57 1.829 68

## TWO EXAMPLES USING TIMERS

```
!print usage.pl1 1
```

```
usage: proc (mc_ptr, name);
```

```
/* THIS PROGRAM PRINTS OUT PAGE FAULT AND PAGING DEVICE  
   FAULT INFORMATION EVERY HALF-SECOND OF CPU TIME */
```

```
dcl (mc_ptr ptr,  
     name char (*)) parameter;
```

```
→ dcl timer_manager_$cpu_call entry (fixed bin (71), bit (2), entry),  
    timer_manager_$reset_cpu_call entry (entry),  
    → cpu_time_and_paging_entry (fixed bin, fixed bin (71),  
    → fixed bin),  
    iox_$control entry (ptr, char (*), ptr, fixed bin (35)),  
    ioa_entry options (variable);
```

```
dcl cleanup condition,  
     code fixed bin (35),  
     iox_$user_io external ptr;
```

```
dcl (total_pf fixed bin,  
     total_pdf fixed bin,  
     first_time bit (1) init ("1"b)) static;
```

```
dcl pf fixed bin,  
     time fixed bin (71),  
     pdf fixed bin;
```

```
/* ESTABLISH cleanup HANDLER TO RESET TIMER */
```

```
→ on cleanup call timer_manager_$reset_cpu_call (usage);
```

```
/* PERFORM INITIALIZATION IF FIRST TIME */
```

```
if first_time then do;  
    → call cpu_time_and_paging_ (total_pf, time, total_pdf);  
    call ioa_ ("Usage counters initialized");  
    call ioa_ ("Total page faults since process began ^i",  
             total_pf);  
    call ioa_ ("Total pd faults since process began ^i",  
             total_pdf);  
    first_time = "0"b;  
end;
```

```
/* COMES HERE WHEN TIMER GOES OFF */
```

```
else do;  
    → call cpu_time_and_paging_ (pf, time, pdf);  
    call ioa_ ("In the half-second CPU interval:");  
    call ioa_ ("^i page faults were taken",  
             pf - total_pf);  
    call ioa_ ("^i paging device faults were taken",  
             pdf - total_pdf);  
    total_pf = pf;
```

TWO EXAMPLES USING TIMERS

```
total_pdf = pdf;
call iox $control (iox_$user_io, "start",
    null (), code);
end;
```

```
/* FIRE UP NEXT TIMER IN EITHER CASE */
    → call timer_manager_$cpu_call (500000, "10"b, usage);
```

```
end usage;
!usage
Usage counters initialized
Total page faults since process began 703
Total pd faults since process began 0
r 12:19 0.083 4
```

```
!hmu
```

```
Multics MR8.0, load 63.0/100.0; 63 users
Absentee users 1/5
```

```
r 12:19 0.055 7
```

```
!list -first 2
```

```
Segments = 62, Lengths = 127.
```

```
re 1 usage
r w 1 usage.pl1
```

```
r 12:19 0.201 21
```

```
!cookie
r 12:19 0.052 1
```

```
!pwd
>udd>F15dw>Auerbach
r 12:19 0.035 0
```

```
I want a cookie.
!cookie
Thanks. I needed that. YumYumYum...
!whome
In the half-second CPU interval:
55 page faults were taken
0 paging device faults were taken
Auerbach.F15dw
```

```
r 12:19 0.310 30
```

TWO EXAMPLES USING TIMERS

!pl1 cookie  
PL/I  
In the half-second CPU interval:  
60 page faults were taken  
0 paging device faults were taken  
In the half-second CPU interval:  
41 page faults were taken  
0 paging device faults were taken  
In the half-second CPU interval:  
63 page faults were taken  
0 paging device faults were taken  
r 12:20 1.876 163

TIMER MANAGER SUMMARY

FOR REAL-TIME TIMERS

sleep GO BLOCKED AT THIS POINT FOR SPECIFIED TIME

alarm\_call CALL SPECIFIED ROUTINE AT SPECIFIED TIME

alarm\_call\_inhibit CALL SPECIFIED ROUTINE AT SPECIFIED TIME WITH ALL INTERPROCESS SIGNALS MASKED OFF

alarm\_wakeup ISSUE A WAKEUP ON SPECIFIED EVENT-WAIT CHANNEL WHEN TIMER GOES OFF, PASSING A MESSAGE OF "alarm\_\_"

reset\_alarm\_call TURN OFF ALL TIMERS THAT CALL ROUTINE SPECIFIED WHEN MATURE

reset\_alarm\_wakeup TURN OFF ALL TIMERS THAT WAKEUP SPECIFIED EVENT-WAIT CHANNEL WHEN THEY MATURE

---

FOR CPU-TIME TIMERS

cpu\_call CALL SPECIFIED ROUTINE AT SPECIFIED TIME

cpu\_call\_inhibit CALL SPECIFIED ROUTINE AT SPECIFIED TIME WITH ALL INTERPROCESS SIGNALS MASKED OFF

cpu\_wakeup ISSUE A WAKEUP ON SPECIFIED EVENT-WAIT CHANNEL WHEN TIMER GOES OFF, PASSING A MESSAGE OF "cpu\_time"

reset\_cpu\_call TURN OFF ALL TIMERS THAT CALL ROUTINE SPECIFIED WHEN THEY MATURE

reset\_cpu\_wakeup TURN OFF ALL TIMERS THAT WAKEUP SPECIFIED EVENT-WAIT CHANNEL WHEN THEY MATURE

|| YOU ARE NOW READY FOR WORKSHOP ||  
#5

TOPIC XI

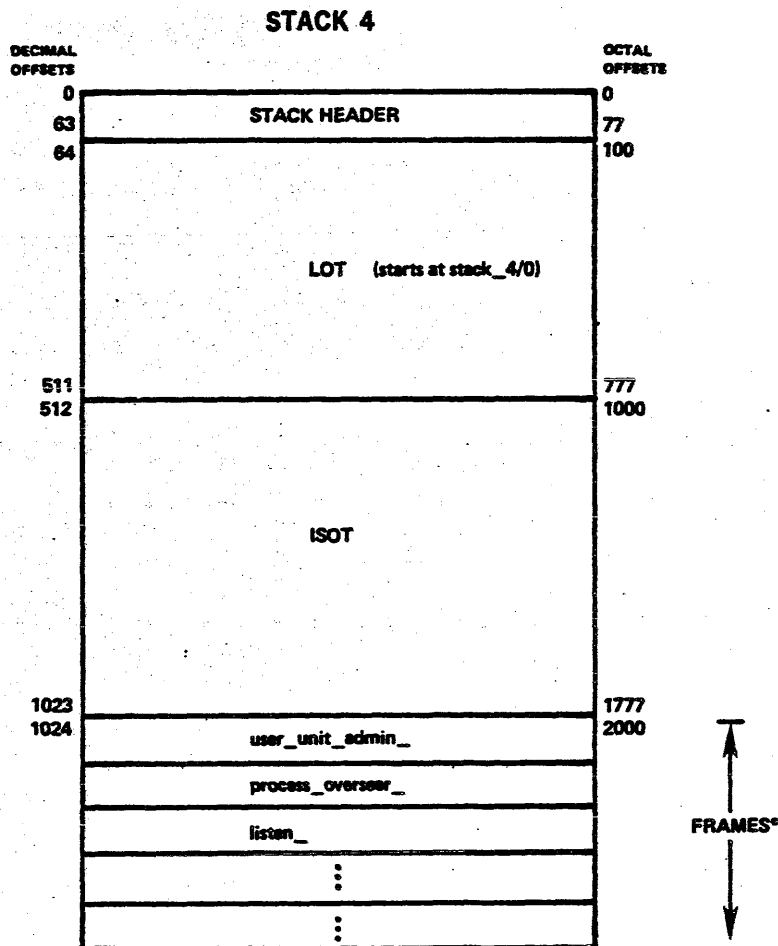
The Stack and Argument Lists

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# THE PROCESS STACK SEGMENT

## ● EXECUTION IN A STANDARD Multics PROCESS USES A STACK SEGMENT

- ▮ THERE IS ONE STACK SEGMENT PER RING WITH ENTRYNAME 'stack\_n' WHERE 'n' IS THE RING OF EXECUTION
- ▮ EACH STACK CONTAINS A STACK "HEADER" FOLLOWED BY AS MANY STACK "FRAMES" AS ARE REQUIRED BY THE PROCESS
- ▮ THE STACK DYNAMICALLY EXPANDS AND SHRINKS AS INDIVIDUAL FRAMES ARE "PUSHED" ONTO THE STACK (BY AN INVOKED PROCEDURE) AND "POPPED" OFF THE STACK (BY A RETURNING PROCEDURE)



\*Frame size is variable, but not less than 32 words

# THE PROCESS STACK SEGMENT

## THE STACK HEADER

- MOSTLY CONTAINS POINTERS TO SPECIAL CODE (STANDARD CALL, PUSH, AND RETURN OPERATORS) AND SPECIAL DATA BASES (LINKAGE OFFSET TABLE, REFERENCE NAME TABLE, ETC) WHICH ARE REQUIRED FOR THE NORMAL EXECUTION OF A PROCESS IN THE CORRESPONDING RING

```
dcl 1 stack_header      based (sb) aligned,

    2 pad1 (4)          fixed bin,
    2 old_lot_ptr       ptr, /* OBSOLETE */
    2 combined_stat_ptr ptr, /* POINTS TO AREA CONTAINING
                             SEPARATE STATIC */

    2 clr_ptr           ptr, /* AREA_PTR FOR LINKAGE SECTION ALLOCATION */
    2 max_lot_size      fixed bin(17) unal,
    2 main_proc_invoked fixed bin(11) unal,
    2 run_unit_depth    fixed bin(5) unal,
    2 cur_lot_size      fixed bin(17) unal, /* IN WORDS */

    2 system_free_ptr   ptr, /* USUALLY POINTS TO SYSTEM FREE STORAGE

    2 user_free_ptr     ptr, /* USUALLY POINTS TO
                             <unique>.area.linker|0 */

    2 null_ptr          ptr, /* THERE IS NO STACK FRAME
                             PREVIOUS TO THE HEADER */

    2 stack_begin_ptr   ptr,
    2 stack_end_ptr     ptr, /* POINTS TO NEXT USEABLE
                             STACK FRAME */

    2 lot_ptr           ptr, /* INITIALLY POINTS TO BASE OF STACK */

    2 signal_ptr        ptr, /* POINTS TO SIGNALLING PROC FOR
                             THIS RING */

    2 bar_mode_sp       ptr, /* NEEDED BECAUSE BAR MODE PROGS CAN
                             CHANGE THE STACK FRAME PTR REGISTER
                             (PR6) */

    2 pl1_operators_ptr ptr, /* POINTS TO pl1 operators_
                             $operator_table */

    2 call_op_ptr       ptr, /* POINTS TO ALM CALL OPERATOR */

    2 push_op_ptr       ptr, /* POINTS TO ALM PUSH OPERATOR */
    2 return_op_ptr     ptr, /* POINTS TO STANDARD RETURN ALM OPERATOR */
    2 return_no_pop_op_ptr ptr, /* POINTS TO SHORT RETURN ALM OPERATOR */
```



THE PROCESS STACK SEGMENT

THE STACK HEADER

```
2 entry_op_ptr      ptr, /* POINTS TO ALM ENTRY OPERATOR */
2 trans_op_tv_ptr   ptr, /* POINTS TO A VECTOR OF SPECIAL
                           LANGUAGE OPERATORS */

2 isot_ptr          ptr,
2 set_ptr           ptr, /* POINTS TO SYSTEM CONDITION
                           TABLE (SCT) */

2 unwinder_ptr      ptr, /* POINTS TO UNWINDER PROCEDURE
                           FOR THIS RING */

2 sys_link_info_ptr ptr, /* POINTS TO *system LINK NAME
                           TABLE */

2 rnt_ptr           ptr, /* POINTS TO REFERENCE NAME TABLE */
2 ect_ptr          ptr, /* OBSOLETE */
2 assign_linkage_ptr ptr, /* OBSOLETE */
2 pad3 (8)         bit (36); /* FOR FUTURE EXPANSION */
```

● NOTE:

set\_system\_storage SETS stack\_header.system\_free\_ptr

set\_user\_storage SETS stack\_header.user\_free\_ptr

WHEN THE NUMBER OF INITIATED SEGMENTS EXCEEDS 512, THE  
lot AND isot ARE COPIED TO SYSTEM FREE STORAGE

THE PROCESS STACK SEGMENT

THE STACK FRAME

- STACK FRAMES ARE VARIABLE LENGTH, AND CONTAIN BOTH CONTROL INFORMATION AND DATA FOR ACTIVE PROCEDURES
- IN GENERAL, A STACK FRAME IS ALLOCATED EXPLICITLY BY THE PROCEDURE TO WHICH IT BELONGS (ITS OWNER) WHEN THAT PROCEDURE IS INVOKED
- STACK FRAMES ARE THREADED TO EACH OTHER WITH FORWARD AND BACKWARD POINTERS, MAKING IT EASY TO TRACE THE PROCESS HISTORY

THE PROCESS STACK SEGMENT

THE STACK FRAME

```
dcl 1 stack_frame based(sp) aligned,

2 pointer_registers(0 : 7) ptr, /* FOR ALM CALL PSEUDO-OP */
2 prev_sp pointer,
2 next_sp pointer, /* IF EQUAL TO stack_end_ptr, POINTS TO
                    NEXT AVAILABLE STACK FRAME LOCATION */
2 return_ptr pointer, /* TELLS US WHERE TO RESUME EXECUTION */
2 entry_ptr pointer, /* POINTS TO THIS PROCEDURE'S ENTRY
                    POINT */
2 operator_and_lp_ptr ptr, /* POINTS TO OPERATOR SEGMENT BEING
                           USED BY THIS PROCEDURE OR,
                           IF ALM PROCEDURE, POINTS TO
                           LINKAGE SECTION */
2 arg_ptr pointer, /* POINTS TO arg_list TO BE USED BY THIS
                  PROCEDURE */
2 static_ptr ptr unaligned, /* POINTS TO INTERNAL STATIC
                           REGION */
2 fio_ps_ptr ptr unal, /* FOR FORTRAN I/O */
2 on_unit_relp1 bit(18) unaligned, /* POINTS TO A LIST OF
                                   ENABLED CONDITIONS
                                   (RELATIVE TO STACK
                                   FRAME BASE) */
2 on_unit_relp2 bit(18) unaligned, /* OBSOLETE */
2 translator_id bit(18) unaligned, /* A CODED NUMBER INDICATING
                                   WHAT GENERATED THE OBJECT */
2 operator_return_offset bit(18) unaligned, /* USED BY SOME
                                             p11_operators
                                             FUNCTIONS; IF 0,
                                             THEN A DEDICATED
                                             REGISTER CONTAINS
                                             RETURN LOCATION */

2 x(0: 7) bit(18) unaligned, }
2 a bit(36),                } USED TO SAVE REGS WHEN AND IF THIS
2 q bit(36),                } PROCEDURE DOES AN ALM CALL
2 e bit(36),                }
2 timer bit(27) unaligned,
2 pad bit(6) unaligned,
2 ring_alarm_reg bit(3) unaligned;

    /* AUTOMATIC VARIABLES */

    /* THREADED LIST OF ON UNITS */
```

THE PROCESS STACK SEGMENT

THE STACK FRAME

```
del 1 on_unit based aligned,
    2 name ptr,                /* ptr to the condition name */
    2 body ptr,                /* ptr to proc to handle condition */
    2 size fixed bin,          /* length of the condition name */
    2 next bit (18) unaligned, /* rel ptr to next on unit*/
    2 flags unaligned,
    3 pl1_snap bit (1) unaligned, /* if "1"b then call snap proc */
    3 pl1_bit (1) unaligned,     /* "1"b indicates to use
                                system condition handler */
    3 pad bit (16) unaligned,
    2 file ptr;                /* ptr to file descriptor for pl1 I/O
                                condition */
```

ARGUMENT LIST FORMAT

- AN OBJECT SEGMENT CREATES AN ARGUMENT LIST BEFORE INVOKING ANOTHER PROCEDURE

- ▮ ARGUMENT LISTS CONFORM TO STANDARD FORMAT

- ▮ ARGUMENT LISTS DO NOT CONTAIN THE ARGUMENTS THEMSELVES

- ▮ SPECIFICALLY, THE ARGUMENT LIST CONTAINS

- ▮ A TWO WORD HEADER

- ▮ AN ARRAY OF ARG POINTERS

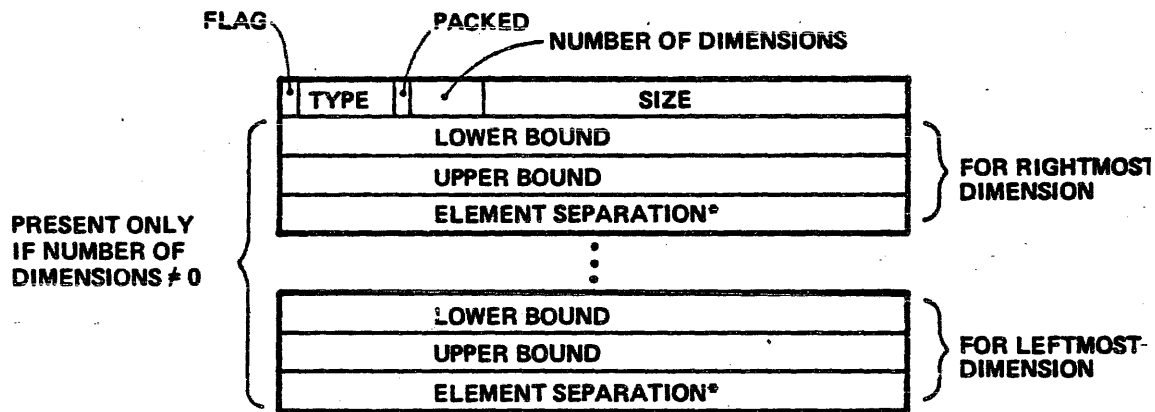
- ▮ AN OPTIONAL POINTER TO STACK FRAME OF CONTAINING BLOCK

- ▮ AN OPTIONAL ARRAY OF POINTERS TO ARGUMENT DESCRIPTORS

## STANDARD ARGUMENT LIST

0	ARGUMENT COUNT= $n$	CODE
1	DESCRIPTOR COUNT= $n$	0
2	- POINTER TO ARGUMENT 1 -	
4	- POINTER TO ARGUMENT 2 -	
	:	
	:	
$2*n$	- POINTER TO ARGUMENT $n$ -	
	- OPTIONAL POINTER TO STACK FRAME OF CONTAINING BLOCK -	
	- POINTER TO DESCRIPTOR 1 -	
	- POINTER TO DESCRIPTOR 2 -	
	:	
	:	
	- POINTER TO DESCRIPTOR $n$ -	

## STANDARD DESCRIPTOR



\* IN BITS IF PACKED; IN WORDS OTHERWISE

ARGUMENT LIST FORMAT

● IN THE PREVIOUS DIAGRAM:

n IS THE NUMBER OF ARGUMENTS PASSED TO THE CALLED PROCEDURE

code IS 4 FOR NORMAL INTERSEGMENT CALLS AND IS 10 (OCTAL)  
FOR CALLING SEQUENCES THAT CONTAIN AN EXTRA STACK FRAME  
POINTER - IT WILL BE PRESENT FOR CALLS TO PL/I INTERNAL  
PROCEDURES

descriptor count = n OR 0

## ARGUMENT DESCRIPTORS

### ● A PROCEDURE WHICH MAY RECEIVE

A VARYING NUMBER OF ARGS

ARGS WITH VARYING EXTENTS

MUST BE PASSED AN ARGUMENT LIST CONTAINING DESCRIPTORS OF THOSE ARGUMENTS, SO THAT THE CALLED PROCEDURE MAY KNOW HOW TO INTERPRET THE ARGUMENTS

- PL/1 ONLY PASSES DESCRIPTORS IF CALLED PROCEDURE IS DECLARED 'entry options (variable)' OR PARAMETERS OF CALLEE HAVE \* EXTENTS
  
- IT IS THE RESPONSIBILITY OF A PROGRAM CALLING SUCH A PROCEDURE TO BUILD DESCRIPTORS AND INCLUDE THEM IN THE ARGUMENT LIST
  
- DESCRIPTORS HAVE A STANDARD FORMAT AS DEFINED BELOW:

```
    dcl 1 descriptor      aligned,
      (2 flag            bit(1),
       2 type            bit(6),
       2 packed          bit(1),
       2 number_dims     bit(4), /* = 15 max */
       2 size            bit(24)) unaligned; /* HAS VARIOUS
                                              MEANINGS */
```

WHERE type IS ENCODED AS SHOWN ON THE NEXT PAGE.



## ARGUMENT DESCRIPTORS

```
/* BEGIN INCLUDE FILE ... std_descriptor_types.incl.pl1 */

/* This include file defines mnemonic names for the Multics
   standard descriptor types, using both pl1 and cobol terminology. */

dcl (real_fix_bin_1_dtype init (1),
     real_fix_bin_2_dtype init (2),
     realflt_bin_1_dtype init (3),
     realflt_bin_2_dtype init (4),
     cplx_fix_bin_1_dtype init (5),
     cplx_fix_bin_2_dtype init (6),
     cplxflt_bin_1_dtype init (7),
     cplxflt_bin_2_dtype init (8),
     real_fix_dec_9bit_ls_dtype init (9),
     realflt_dec_9bit_dtype init (10),
     cplx_fix_dec_9bit_ls_dtype init (11),
     cplxflt_dec_9bit_dtype init (12),
     pointer_dtype init (13),
     offset_dtype init (14),
     label_dtype init (15),
     entry_dtype init (16),
     structure_dtype init (17),
     area_dtype init (18),
     bit_dtype init (19),
     varying_bit_dtype init (20),
     char_dtype init (21),
     varying_char_dtype init (22),
     file_dtype init (23),
     real_fix_dec_9bit_ls_overp_dtype init (29),
     real_fix_dec_9bit_ts_overp_dtype init (30),
     real_fix_bin_1_uns_dtype init (33),
     real_fix_bin_2_uns_dtype init (34),
     real_fix_dec_9bit_uns_dtype init (35),
     real_fix_dec_9bit_ts_dtype init (36),
     real_fix_dec_4bit_uns_dtype init (38), /* digit-aligned */
     real_fix_dec_4bit_ts_dtype init (39), /* byte-aligned */
     real_fix_dec_4bit_bytealigned_uns_dtype init (40), /* COBOL */
     real_fix_dec_4bit_ls_dtype init (41), /* digit-aligned */
     realflt_dec_4bit_dtype init (42), /* digit-aligned */
     real_fix_dec_4bit_bytealigned_ls_dtype init (43),
     realflt_dec_4bit_bytealigned_dtype init (44),
     cplx_fix_dec_4bit_bytealigned_ls_dtype init (45),
     cplxflt_dec_4bit_bytealigned_dtype init (46),
```

## ARGUMENT DESCRIPTORS

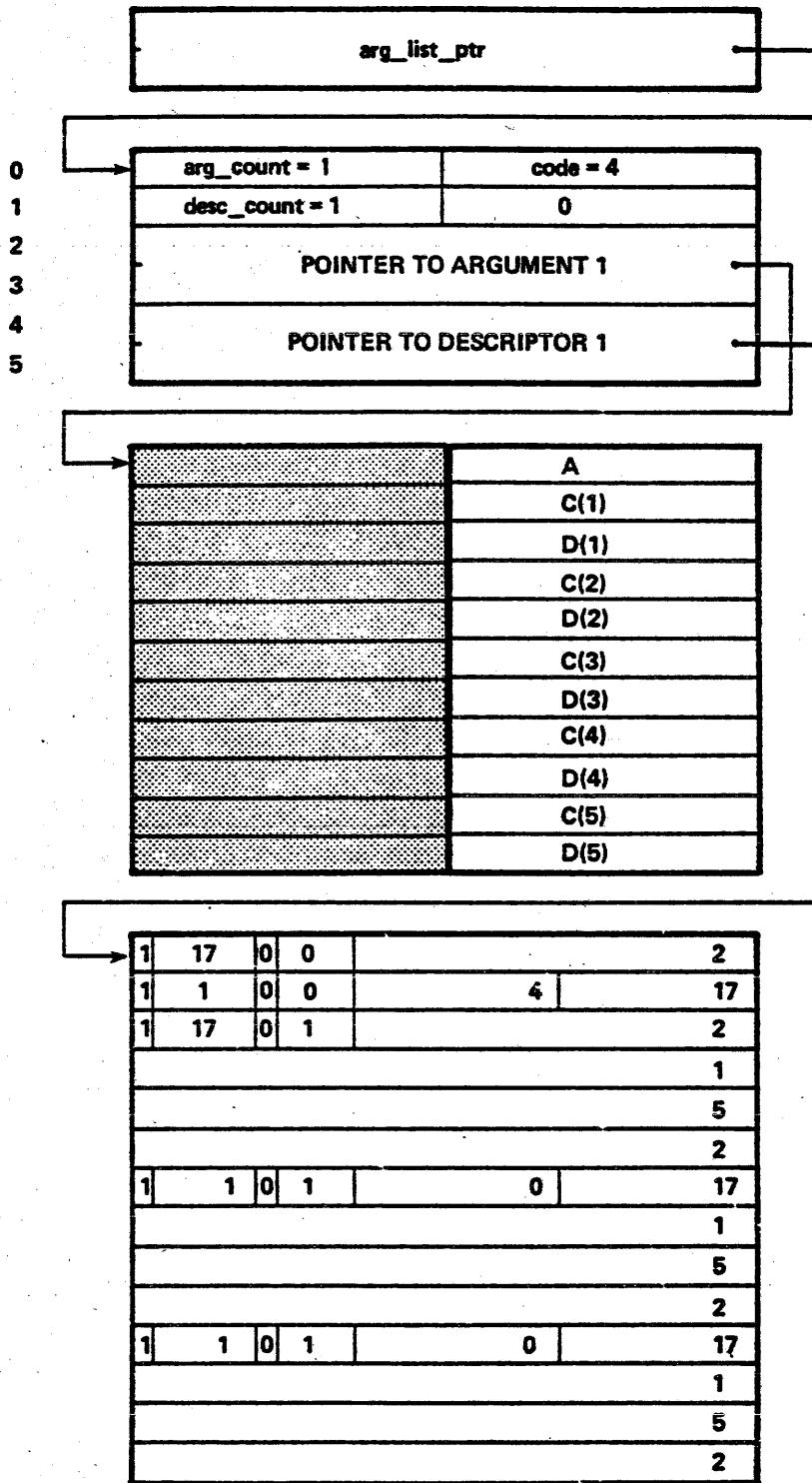
```
cobol_comp_6_dtype init (1),
cobol_comp_7_dtype init (1),
cobol_display_ls_dtype init (9),
cobol_structure_dtype init (17),
cobol_char_string_dtype init (21),
cobol_display_ls_overp_dtype init (29),
cobol_display_ts_overp_dtype init (30),
cobol_display_uns_dtype init (35),
cobol_display_ts_dtype init (36),
cobol_comp_8_uns_dtype init (38), /* digit aligned */
cobol_comp_5_ts_dtype init (39), /* byte aligned */
cobol_comp_5_uns_dtype init (40),
cobol_comp_8_ls_dtype init (41) /* digit aligned */
) fixed_bin internal static options (constant);

dcl (ft_integer_dtype init (1),
ft_real_dtype init (3),
ft_double_dtype init (4),
ft_complex_dtype init (7),
ft_external_dtype init (16),
ft_logical_dtype init (19),
ft_char_dtype init (21)
) fixed_bin internal static options (constant);

dcl (label_constant_runtime_dtype init (24),
int_entry_runtime_dtype init (25),
ext_entry_runtime_dtype init (26),
ext_procedure_runtime_dtype init (27),
picture_runtime_dtype init (63)
) fixed_bin internal static options (constant);

/* END INCLUDE FILE ... std_descriptor_types.incl.pl1 */
```

# ARGUMENT DESCRIPTORS



```

dcl 1 S,
    2 A fixed bin (17,4),
    2 B(5),
    3 C,
    3 D;
    
```

- BASIC DESCRIPTOR OF S
- BASIC DESCRIPTOR OF A
- BASIC DESCRIPTOR OF B
- LOWER BOUND OF B
- UPPER BOUND OF B
- ELEMENT SEPARATION OF B
- BASIC DESCRIPTOR OF C
- LOWER BOUND OF C
- UPPER BOUND OF C
- ELEMENT SEPARATION OF C
- BASIC DESCRIPTOR OF D
- LOWER BOUND OF D
- UPPER BOUND OF D
- ELEMENT SEPARATION OF D

TOPIC XII

Special Programming Techniques

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Building Data Segments . . . . .	12-4
Creating an Error Table. . . . .	12-9

CALLING A PROCEDURE ON THE FLY

● MOTIVATION

- ┆ CONSIDER THE COMMAND PROCESSOR, WHICH:
  - ┆ TAKES THE COMMAND LINE FROM THE USER
  - ┆ PARSES THE LINE INTO COMMAND NAME AND ARGUMENTS
  - ┆ CALLS THE COMMAND
  
- ┆ HOW CAN IT CALL ALL THE VARIOUS COMMANDS (AND USER WRITTEN OBJECT SEGMENTS) WITHOUT HAVING DECLARED THEM AS AN EXTERNAL ENTRY?

● HOW CAN THE USER DO SUCH A "CALL ON THE FLY"?

▮ hcs\_\$make\_entry

▮ call hcs\_\$make\_entry (ref\_ptr, entryname, entry\_point\_name,  
entry\_point, code);

▮ GIVEN A REFERENCE NAME AND AN ENTRY POINT NAME, RETURNS THE  
ENTRY VALUE OF THE SPECIFIED ENTRY POINT

▮ IF THE REFERENCE NAME HAS NOT BEEN INITIATED, THE SEARCH  
RULES ARE USED TO FIND A SEGMENT WITH THAT NAME, THE SEGMENT  
IS MADE KNOWN AND THE REFERENCE NAME INITIATED

▮ cu\_\$generate\_call

▮ call cu\_\$generate\_call (proc\_entry, arg\_ptr);

▮ USED TO INVOKE A PROCEDURE BY PASSING IT AN ENTRY VALUE AND  
AN ARGUMENT POINTER (THE ENTRY VALUE WAS OBTAINED BY A PREVIOUS  
CALL TO hcs\_\$make\_entry)

▮ THE USER MUST HAVE PROVIDED AN ARGUMENT LIST STRUCTURE  
(EVEN IF NO ARGUMENTS ARE PASSED TO THE PROCEDURE BEING  
INVOKED)

CALLING A PROCEDURE ON THE FLY

● SIMPLIFIED EXAMPLE:

```
generate_pwd: proc;
dcl 1 arg_list aligned,
    2 header,
        3 arg_count fixed bin (17) unsigned unal init (0),
        3 pad1 bit (1) unal,
        3 call_type fixed bin (18) unsigned unal,
        3 desc_count fixed bin (17) unsigned unal,
        3 pad2 bit (19) unal,
    2 arg_ptr ptr init (null()),
    2 desc_ptr ptr init (null());

dcl cu $generate_call entry (entry, ptr);
dcl hcs $make_entry entry (ptr, char (*), char (*), entry,
    fixed bin (35));

dcl code fixed bin (35);
dcl com_err_entry options (variable);
dcl entry_point entry variable;

    call hcs $make_entry (null(), "pwd", "pwd", entry_point, code);
    if code = 0 then do;
        call com_err_ (code, "generate");
        return;
    end;
    call cu $generate_call (entry_point, addr (arg_list));
end generate_pwd;
```

## BUILDING DATA SEGMENTS

### ● create\_data\_segment\_

⌋ call create\_data\_segment\_ (cds\_arg\_ptr, code);

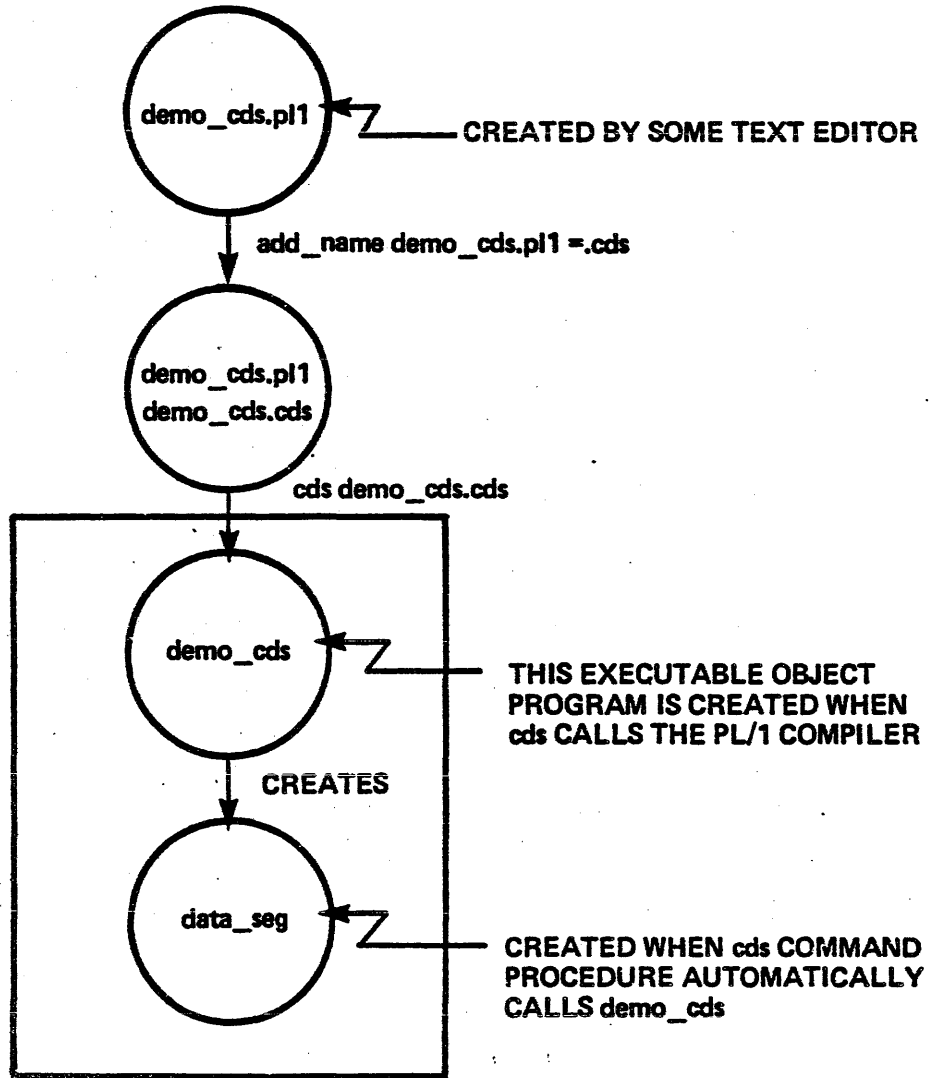
⌋ USED IN CONJUNCTION WITH THE create data segment COMMAND TO CREATE A DATA SEGMENT IN STANDARD OBJECT FORMAT

⌋ REFERENCES A PL/1 DATA STRUCTURE (SUPPLIED BY THE USER) WHEN BUILDING THE DATA SEGMENT

⌋ PERMITS THE CALLER TO CREATE A DATA SEGMENT FOR WHICH THE LEVEL-2 MEMBERS OF THE PL/1 STRUCTURE ARE ACCESSIBLE AS ENTRY POINTS (NOTE THE pds AND sys\_info SEGMENTS)



## HOW cds WORKS



BUILDING DATA SEGMENTS

! pr >ldd>include>cds\_args.incl.pl1

/\* BEGIN INCLUDE FILE cds\_args.incl.pl1 \*/

```
dcl 1 cds_args based aligned,
  2 sections (2),
    3 p ptr,          /* pointer to data for text/static section */
    3 len fixed bin (18), /* size of text/static section */
    3 struct_name char (32), /* name of declared structure */
  2 seg_name char (32), /* name to create segment by */
  2 num_exclude_names fixed bin, /* number in exclude array */
  2 exclude_array_ptr ptr, /* pointer to exclude array */
  2 switches, /* control switches */
    3 defs_in_link bit (1) unal, /* says put defs in linkage */
    3 separate_static bit (1) unal, /* separate static section
                                     wanted */
    3 have_text bit (1) unal, /* ON if text section given */
    3 have_static bit (1) unal, /* ON if static section given */
    3 pad bit (32) unal; /* must be zero */

dcl exclude_names (1) char (32) based;
    /* pointed to by cds_args.exclude_array_ptr */

/* END INCLUDE FILE cds_args.incl.pl1 */
```

BUILDING DATA SEGMENTS

! pr demo\_cds.pl1

```
demo_cds: proc;
dcl create_data_segment_entry (ptr, fixed bin (35));
%include cds_args;
dcl cds_arg_ptr ptr;
dcl (ioa, com err) entry options (variable);
dcl code fixed bin (35);
dcl (size, null) builtin;
dcl 1 entrypointnames based (cds_arg_ptr -> cds_args.sections (1).p),
    2 alpha fixed bin (35),
    2 beta char (4),
    2 gamma bit (36),
    2 delta ptr;

allocate cds_args set (cds_arg_ptr);
cds_arg_ptr -> cds_args.sections (1).len = size (entrypointnames);
cds_arg_ptr -> cds_args.sections (1).struct_name = "entrypointnames";
cds_arg_ptr -> cds_args.seg_name = "data_seg";
cds_arg_ptr -> cds_args.sections (2).len = 0;
cds_arg_ptr -> cds_args.sections (2).struct_name = "";
cds_arg_ptr -> cds_args.num_exclude_names = 0;
cds_arg_ptr -> cds_args.exclude_array_ptr = null ();
cds_arg_ptr -> cds_args.switches.defs_in_link = "0"b;
cds_arg_ptr -> cds_args.switches.separate_static = "0"b;
cds_arg_ptr -> cds_args.switches.have_text = "1"b;
cds_arg_ptr -> cds_args.switches.have_static = "0"b;
cds_arg_ptr -> cds_args.switches.pad = "0"b;

call create_data_segment (cds_arg_ptr, code);
if code ^= 0 then call com_err (code, "demo_cds");
else call ioa ("Segment creation complete.");
end demo_cds;
```

BUILDING DATA SEGMENTS

!an demo\_cds.pl1 =.cds  
r 14:32 0.096 2

!cds demo\_cds  
CDS -PL/I 26a  
Segment creation complete.  
r 14:32 1.699 53

!ls -first 3

Segments = 43, Lengths = 155.

r 1 data\_seg  
re 1 demo\_cds  
r w 1 demo\_cds.pl1  
demo\_cds.cds

r 14:33 0.151 2

!pli data\_seg

data\_seg 10/10/80 1433.9 mst Fri

Object Segment >udd>MEDmult>F15C>do>data\_seg  
Created on 10/10/80 1432.9 mst Fri  
by NDibble.MEDmult.a  
using create\_data\_segment\_, Version II of Friday, May 16, 1980

	Object	Text	Defs	Link	Symb	Static
Start	0	0	6	52	62	62
Length	224	6	44	10	126	0

6 Definitions:

segname: data\_seg

text|0 alpha  
text|1 beta  
text|4 delta  
text|2 gamma  
symb|0 symbol\_table

No Links.

## CREATING AN ERROR TABLE

- USERS MAY CREATE THEIR OWN STATUS CODE TABLE

- ▮ THEY ARE CONSTRUCTED USING ALM MACROS DEFINED IN `et_macros.incl.alm`

- ▮ THE SKELETON OF THE SOURCE CODE IS AS FOLLOWS:

```
include et_macros
et .....
ec .....
ec .....
ec .....
.
.
.
end
```

- ▮ THERE ARE THUS 2 MACROS USED

- ▮ THE "et" MACRO INITIALIZED THE CODE TABLE AND MUST APPEAR FIRST

```
et <name_of_table>
```

- ▮ THE "ec" MACRO ASSOCIATES A STATUS CODE NAME WITH A SHORT MESSAGE AND A LONG MESSAGE

```
ec <code_name>,<short_message>,(<long_message>)
```

- ▮ EXAMPLE:

```
include et_macros
et user_errors
ec too_few_arguments, toofew,(There were too few arguments.)
ec could_not_access_data,nopriv,
  (User is not sufficiently privileged to access data.
ec (fatal,disaster),disaster,
  (There was a disastrous error in the data base.)
end
```

CREATING AN ERROR TABLE

▯ THE CODE NAME:

▯ MUST BE 31 CHARACTERS OR LESS IN LENGTH

▯ MULTIPLE NAMES MAY BE GIVEN (SEPARATED BY COMMAS AND ENCLOSED IN PARENTHESIS)

▯ THE SHORT MESSAGE:

▯ MUST BE 8 OR LESS CHARACTERS IN LENGTH

▯ IF OMITTED, IT IS SET TO THE CODE NAME

▯ THE LONG MESSAGE MUST BE:

▯ 100 OR LESS CHARACTERS IN LENGTH

▯ ENCLOSED IN PARENTHESIS

CREATING AN ERROR TABLE

```
● ! pr weird_errors.alm

include et_macros
et weird_errors
ec error_a,number2,(Warning: The number has reached two.)
ec error_b,number3,(Second Warning: The number has reached three.)
end

! alm weird_errors.alm
ALM

! pr calling_program.pl1

calling_program: proc;
  dcl x fixed bin (17) external static init(0);
  dcl com_err_entry options (variable);
  dcl my_subprogram entry (fixed bin(17), fixed bin (35));
  dcl code fixed bin (35);
  dcl sysprint file;
  x = x + 1;
  put data (x);
  put skip;
  call my_subprogram (x, code);
  if code ^= 0 then call com_err_ (code, "calling_program");
end calling_program;

! pr my_subprogram.pl1

my_subprogram: proc (input, code);
  dcl input fixed bin (17);
  dcl code fixed bin (35);
  dcl weird_errors$error_a fixed bin (35) external static;
  dcl weird_errors$error_b fixed bin (35) external static;
  code = 0;
  if input = 2 then code = weird_errors$error_a;
  if input = 3 then code = weird_errors$error_b;
end my_subprogram;

! calling_program
x=      1;

! calling_program
x=      2;
calling_program: Warning: The number has reached two.

! calling_program
x=      3;
calling_program: Second Warning: The number has reached three.

! calling_program
x=      4;
```

CREATING AN ERROR TABLE

|| YOU ARE NOW READY FOR WORKSHOP ||  
#6



TOPIC XIII

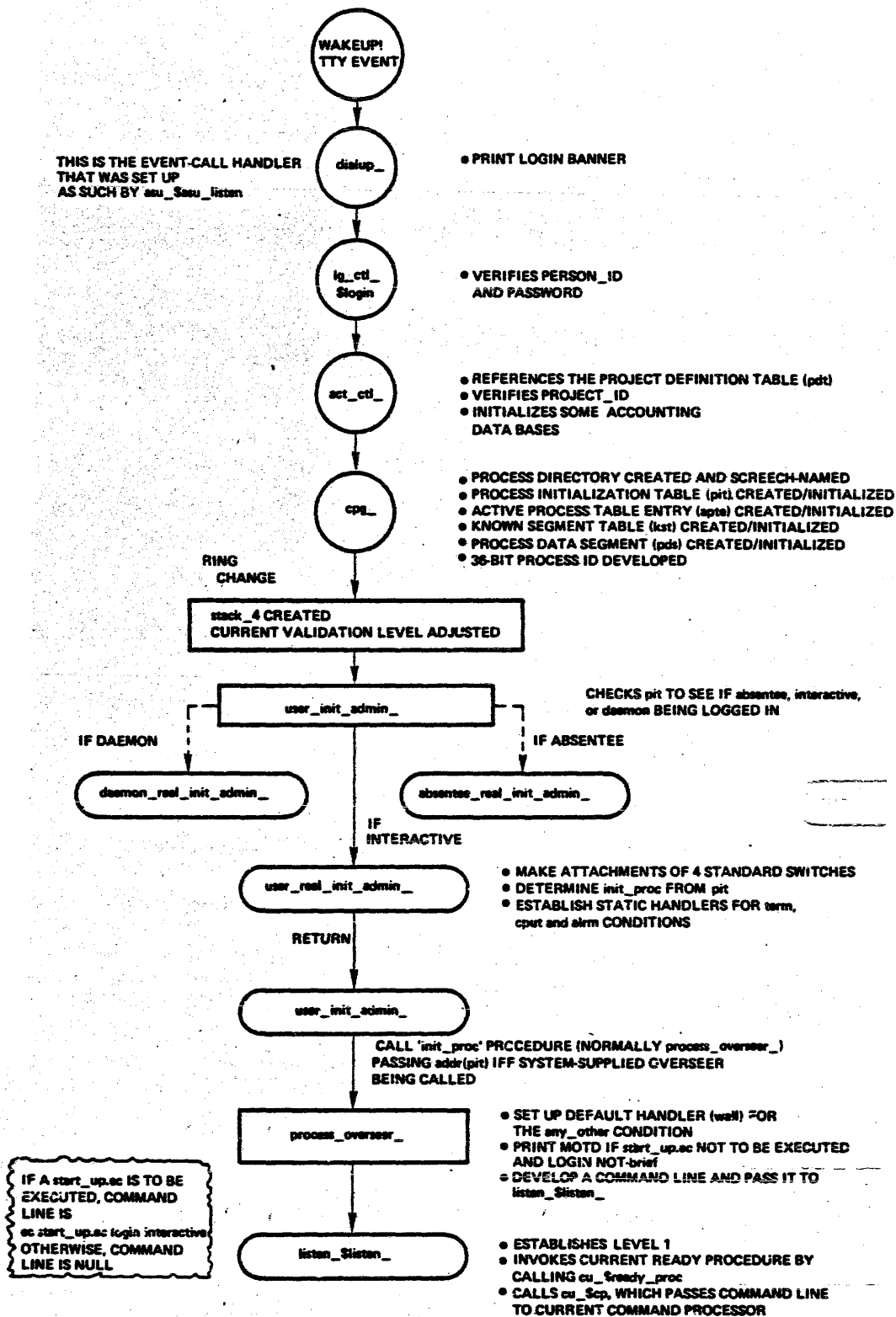
The Process Environment

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## THE STANDARD PROCESS ENVIRONMENT

- THE STANDARD PROCESS ENVIRONMENT IS ESTABLISHED FOR A PROCESS BY THE SUPERVISOR WHEN THAT PROCESS IS CREATED
- THE SERIES OF PROGRAMS INVOKED BEGINNING WITH THE TIME THAT A USER CONNECTS WITH THE FRONT-END PROCESSOR AND TERMINATING WITH THE FIRST READY MESSAGE PRINTED ON THE USER'S TERMINAL IS CONSIDERED THE "PROCESS CREATION CYCLE"
- MUCH OF THE WORK DONE TO CREATE A PROCESS IS DONE BY THE SUPERVISOR IN THE SUPERVISOR RING (RING 0), AND SOME OF THE WORK IS DONE IN THE USER'S RING
- BEFORE EXAMINING THE OPTIONS AVAILABLE TO THE DESIGNER FOR MODIFYING OR REPLACING THE STANDARD PROCESS ENVIRONMENT, AN EXAMINATION OF WHAT STEPS THE SUPERVISOR NORMALLY TAKES WILL BE WORTHWHILE

# THE STANDARD PROCESS ENVIRONMENT



THE STANDARD PROCESS ENVIRONMENT  
PROCESS CREATION IN THE SUPERVISOR RING

- THE "SYSTEM CONTROL PROCESS" (Initializer.SysDaemon) IS RESPONSIBLE FOR RESPONDING TO ATTEMPTS TO CREATE A PROCESS, AND SOME OF THE MODULES EXECUTED BY THE SYSTEM CONTROL PROCESS IMPLEMENT PROCESS CREATION AS FOLLOWS:

┆ ANSWERING SERVICE

┆ IS RESPONSIBLE FOR

┆ INTERACTIVE LOGINS

┆ HANGUPS

┆ THE LOGGING OF THESE ACTIVITIES

┆ MANAGING THE COMMUNICATION LINES CURRENTLY ATTACHED TO SYSTEM  
(USES CHANNEL DEFINITION TABLE)

┆ asu \$asu listen ESTABLISHES AN EVENT-CALL CHANNEL FOR EACH COMLINE, MAKING 'dialup\_' THE EVENT HANDLER

┆ A "CONNECT" ON A COMLINE IS CONSIDERED A "TTY EVENT", WHICH RESULTS IN A WAKEUP ON ONE OF THESE CHANNELS

┆ dialup\_ (IN >tools>bound\_user\_control\_)

┆ IS AUTOMATICALLY CALLED WHEN A TERMINAL IS CONNECTED

┆ PRINTS LOGIN BANNER ON THE TERMINAL

┆ READS INITIAL TYPED LINE (login, enter, enterp, dial, ETC) AND PASSWORD

┆ CALLS lg\_ctl\_\$login TO VERIFY PERSON\_ID AND PASSWORD

THE STANDARD PROCESS ENVIRONMENT  
PROCESS CREATION IN THE SUPERVISOR RING

- ▮ dialup\_, CONTINUED
  - ▮ CALLS act\_ctl\_\$open\_account AND act\_ctl\_\$cp TO
    - ▮ VERIFY PROJECT\_ID (VIA A CHECK OF PROJECT DEFINITION TABLE)
    - ▮ INITIALIZE SOME ACCOUNTING DATA BASES FOR THIS PROCESS
  - ▮ FINALLY CALLS cpg\_ (CREATE PROCESS GROUP) TO ACTUALLY CREATE PROCESS
    - ▮ 36-BIT PROCESS ID DEVELOPED
    - ▮ PROCESS DIRECTORY CREATED AND SCREECH-NAMED
    - ▮ PROCESS INITIALIZATION TABLE CREATED/INITIALIZED
    - ▮ ACTIVE PROCESS TABLE ENTRY CREATED/INITIALIZED
    - ▮ KST AND DSEG CREATED/INITIALIZED
    - ▮ PROCESS DATA SEGMENT (RING 0 STACK) CREATED/INITIALIZED
- ▮ FINALLY, CONTROL IS PASSED IN A RATHER UNUSUAL FASHION TO THE OUTER RING IN WHICH THE PROCESS IS TO RESIDE
  - ▮ A STACK SEGMENT IS CREATED IN THE ULTIMATE RING (USUALLY 4)
  - ▮ A STACK FRAME IS LAID DOWN FOR A USER RING INITIALIZATION PROGRAM
  - ▮ THE CURRENT VALIDATION LEVEL IS ADJUSTED TO THAT RING LEVEL
  - ▮ A PSEUDO-RETURN IS MADE TO THAT USER-RING INITIALIZATION PROGRAM

THE STANDARD PROCESS ENVIRONMENT  
PROCESS INITIALIZATION IN THE USER RING

- THE STANDARD PROCESS ENVIRONMENT IS ESTABLISHED IN THE USER RING BY THE INVOCATION OF THE USER-RING PROCESS CREATION PROGRAMS:

┆ user\_init\_admin\_ FUNCTIONS

┆ AFTER BEING INVOKED THROUGH ONE OF THREE ENTRY POINTS:

┆ user\_init\_admin\_

┆ absentee\_init\_admin\_

┆ daemon\_init\_admin\_

┆ IT CALLS EITHER

┆ user\_real\_init\_admin\_\$user\_real\_init\_admin\_

┆ OR daemon\_real\_init\_admin\_\$daemon\_real\_init\_admin\_

┆ OR absentee\_real\_init\_admin\_\$absentee\_real\_init\_admin\_

┆ - WHEN RETURNED TO BY ONE OF ABOVE, IT CALLS THE 'init\_proc' PROCEDURE

┆ THE VALUE OF 'init\_proc' IS RETURNED BY ONE OF THE ABOVE THREE

┆ A pit\_ptr IS PASSED TO init\_proc

THE STANDARD PROCESS ENVIRONMENT  
PROCESS INITIALIZATION IN THE USER RING

- ⌋ user\_real\_init\_admin\_ FUNCTIONS
  - ⌋ DEVELOP PTR TO PROCESS INITIALIZATION TABLE (pit)
  - ⌋ DETERMINE FROM pit WHICH IO\_MODULE TO USE (USUALLY tty\_)
  - ⌋ ATTACH user\_input, user\_output, error\_output AS SYNONYMS FOR user\_i/o
  - ⌋ ATTACH AND OPEN user\_i/o TO CHANNEL\_NAME FOUND IN pit
  - ⌋ DETERMINE (FROM pit) WHAT 'init\_proc' PROCEDURE IS TO BE USED (USUALLY process\_overseer\_\$process\_overseer\_)
    - ⌋ PROJECT ADMINISTRATOR MAY FORCE A PARTICULAR init\_proc TO BE USED
    - ⌋ USER MAY HAVE BEEN GIVEN PERMISSION TO USE '-po' CONTROL ARG OF THE login COMMAND, AND COULD HAVE USED IT TO SPECIFY THE DESIRED init\_proc
  - ⌋ SET UP STATIC HANDLERS FOR 'alm', 'cput', AND 'term' CONDITIONS BY CALLING sct\_manager\_\$set

THE STANDARD PROCESS ENVIRONMENT  
PROCESS INITIALIZATION IN THE USER RING

┆ process\_overseer\_ FUNCTIONS

┆ CALL condition TO SET UP 'default\_error\_handler\_\$wall' AS  
THE HANDLER FOR 'any\_other' CONDITION

┆ IMPORTANT: IN ORDER TO COMPLETELY CONTROL A USER, ONE  
MUST BE ABLE TO CHANGE THIS DEFAULT BEHAVIOR. THIS IS ONE  
OF THE MOST IMPORTANT REASONS FOR FURNISHING A SPECIAL  
init\_proc.

┆ CHECK pit TO SEE IF THIS IS new\_proc OR login

┆ CHECK pit TO SEE IF start\_up.ec SHOULD BE INVOKED

┆ RECALL: USE OF -ns LOGIN CONTROL ARG IS RESTRICTABLE

┆ BUILD ONE OF TWO INITIAL COMMAND LINES:

┆ exec\_com startup\_dir>start\_up instance type

┆ WHERE startup\_dir IS HOME DIRECTORY, PROJECT DIRECTORY,  
OR >sc1

┆ WHERE instance IS 'login' OR 'new\_proc'

┆ AND type IS 'interactive' OR 'absentee'

┆ home\_dir, instance, AND type ARE DETERMINED FROM pit

┆ "" (NULL COMMAND LINE)

┆ TERMINATE pit

┆ CALL listen\_\$listen\_ WITH INITIAL COMMAND LINE



THE STANDARD PROCESS ENVIRONMENT  
PROCESS INITIALIZATION IN THE USER RING

┆ listen\_

┆ BASIC FUNCTIONS

- ┆ 'LISTENS' FOR LINES TYPED BY THE USER
- ┆ PASSES COMMAND LINE ON TO THE CURRENT COMMAND PROCESSOR
- ┆ INVOKES CURRENT READY PROCEDURE WHEN RETURNED TO (AFTER COMMAND LINE EXECUTION OR PRINTING OF ERROR MESSAGES)
- ┆ ALSO ENABLES QUILTS

┆ ENTRY POINTS IN listen\_:

┆ listen\_

- ┆ INVOKED ONLY BY init\_proc
- ┆ ESTABLISHES A FRAME CONSIDERED TO BE "FIRST LEVEL"
- ┆ PASSES com\_line\_ptr AND com\_line\_length TO cu\_\$cp
- ┆ CALLS cu\_\$ready\_proc WHEN RETURNED TO

┆ release\_stack

- ┆ CALLED WHEN AN ATTEMPT IS MADE TO REENTER COMMAND LEVEL (LEVEL ≠ 1)
- ┆ ESTABLISHES ITS STACK FRAME AS CURRENT LEVEL OF LISTENER
- ┆ 'REMEMBERS' PREVIOUS LEVEL AND VERY FIRST LEVEL OF LISTENER (FOR PURPOSES OF release)
- ┆ PASSES com\_line\_ptr AND com\_line\_length TO cu\_\$cp
- ┆ CALLS cu\_\$ready\_proc WHEN RETURNED TO

┆ OTHER ENTRY POINTS

- ┆ USED BY OTHER PROCEDURES TO OBTAIN INFO ABOUT
  - WHERE TO RELEASE TO
  - WHERE TO 'start'

## MODIFYING THE PROCESS ENVIRONMENT

- ONE CAN MODIFY THE STANDARD PROCESS ENVIRONMENT USING SEVERAL DIFFERENT TECHNIQUES

- THE AMOUNT OF CONTROL DESIRED ON A PROCESS CAN BE CLASSIFIED AS FOLLOWS:

- ┆ SIMPLE CONTROL

- ┆ USING 'exec com' SEGMENTS, THE USER CAN BE RESTRICTED TO LABEL ENTRY POINTS IN THE ec SEGMENT ITSELF

- ┆ STANDARD PROCESS OVERSEERS

- ┆ PROVIDED BY SYSTEM TO CONTROL THE ENVIRONMENT OF THE USER IN VARYING LEVELS OF RESTRICTION

- ┆ accounts overseer IS USED BY 'REGISTRATION AND ACCOUNTING ADMINISTRATORS' TO LIMIT THE NUMBER OF THINGS THEY CAN DO (SEE APPENDIX B)

- ┆ SEE APPENDIX E FOR SOME OF THE STANDARD OVERSEERS

- ┆ CLOSED SUBSYSTEM OVERSEERS

- ┆ THE 'fst process overseer' IS AN EXAMPLE OF AN OVERSEER WHICH PLACES THE USER IN A COMPLETELY CLOSED ENVIRONMENT FROM WHICH ESCAPE IS IMPOSSIBLE

## MODIFYING THE PROCESS ENVIRONMENT

### ▮ LIMITED SUBSYSTEMS

▮ THE SYSTEM PROVIDES THREE WAYS OF FORCING USERS INTO A LIMITED SUBSYSTEM

### ▮ USER-CREATED SUBSYSTEMS

▮ BY WRITING ONE'S OWN PROCESS OVERSEER, ONE CAN ATTAIN COMPLETE, CUSTOMIZED CONTROL OVER A PROCESS

- MOST OF THESE FUNCTIONS REQUIRE INVOLVEMENT OF A PROJECT ADMINISTRATOR (BECAUSE OF NEED TO MODIFY pmf AND INSTALL NEW pdt)

- EXAMPLE OF A pmf

```
Projectid:      F15D;
Initproc:      process_overseer_;
Grace:         60;
Attributes:    vinitproc,vhomedir,multip,nostartup,dialok,
               disconnect_ok,save_on_disconnect;
Limit:         75.00;

personid:      Student_01;
personid:      Student_02;
personid:      Student_03;
personid:      Student_04;
personid:      Student_05;
personid:      Student_06;
personid:      Student_07;
end;
```

MODIFYING THE PROCESS ENVIRONMENT

PROJECT ADMINISTRATION

- THE LIST OF PERSONS WHO MAY LOG IN ON A PROJECT IS CONTAINED IN A BINARY TABLE, THE PROJECT DEFINITION TABLE (pdt), WHICH RESIDES IN THE DIRECTORY >se1>pdt
  
- ▮ ONE pdt SEGMENT EXISTS FOR EACH PROJECT
  
- ▮ ONE pdt ENTRY EXISTS FOR EACH USER, SPECIFYING THE USER'S ATTRIBUTES AND RESOURCE LIMITS ON THE PARTICULAR PROJECT
  
- ▮ USING THE cv\_pmf COMMAND, A PROJECT ADMINISTRATOR CREATES A TEMPORARY pdt FROM A SEGMENT KNOWN AS THE PROJECT MASTER FILE (pmf), WHICH IS USUALLY UNDER THE PROJECT DIRECTORY
  
- ▮ THE TEMPORARY pdt IS INSTALLED IN THE SYSTEM DIRECTORY BY THE PROJECT ADMINISTRATOR USING THE install COMMAND
  
- ▮ AT LOGIN TIME, act ctl USES THE APPROPRIATE pdt TO DETERMINE WHICH OPTIONS AND RESOURCES ARE AVAILABLE TO A USER
  
- SEE MAM PROJECT ADMINISTRATOR (Order No. AK51) FOR COMPLETE DETAILS

MODIFYING THE PROCESS ENVIRONMENT

PROJECT ADMINISTRATION

- SOME OF THE ATTRIBUTE INFORMATION MAINTAINED FOR EACH USER IN THE PDT IS GIVEN BELOW:

▮ homedir - ABSOLUTE PATHNAME OF USER'S HOME DIRECTORY

▮ initproc - NAME OF THE USER'S PROCESS OVERSEER PROCEDURE

▮ attributes:

▮ nobump - USER NOT SUBJECT TO PREEMPTION

▮ dialok - USER MAY USE THE DIAL FACILITY

▮ multip - USER MAY LOG IN MORE THAN ONE INTERACTIVE PROCESS

▮ vinitproc - USER MAY SPECIFY PROCESS OVERSEER AT LOGIN

▮ vhomedir - USER MAY SPECIFY HOME DIRECTORY AT LOGIN

▮ nostartup - USER MAY ESCAPE FROM USING HIS start\_up.ec

▮ AND SO ON...

## MODIFYING THE PROCESS ENVIRONMENT

### CLOSED SUBSYSTEMS

- A CLOSED SUBSYSTEM IS A SUBSYSTEM IN WHICH THE Multics SYSTEM IS NOT DIRECTLY AVAILABLE - RATHER, THOSE PROGRAMS IMPLEMENTING THE CLOSED ENVIRONMENT TAKE FULL CONTROL OVER A PROCESS
  
- THE SYSTEM-SUPPLIED OVERSEER 'accounts\_overseer\_' IMPLEMENTS A CLOSED SUBSYSTEM
  - ▮ accounts\_overseer\_ HANDLES ALL INPUT FROM THE TERMINAL (I.E., IT IS IT'S OWN LISTENER), AND IT SEVERELY RESTRICTS THE USER TO A SMALL SET OF COMMANDS
  
  - ▮ IT MAKES USE OF AN exec\_com SEGMENT, '>tools>master.ec' TO IMPLEMENT SPECIAL FUNCTIONS FOR THE "REGISTRATION AND ACCOUNTING ADMINISTRATOR" WHO WILL BE OPERATING UNDER THIS ENVIRONMENT
  
- THE SYSTEM-SUPPLIED 'fst process overseer\_' ALSO IMPLEMENTS A CLOSED SUBSYSTEM, A "TIME-SHARING FORTRAN" SYSTEM
  
- THE SYSTEM SUPPLIED OVERSEER 'project\_start\_up\_' MAY BE USED TO IMPLEMENT A CLOSED SUBSYSTEM
  - ▮ project\_directory>project\_start\_up.ec ALWAYS EXECUTED BEFORE start\_up.ec (WITH QUILTS DISABLED)

## MODIFYING THE PROCESS ENVIRONMENT

### LIMITED SUBSYSTEMS

- IN A LIMITED SUBSYSTEM, THE USER IS LIMITED TO A SET OF COMMANDS CONTAINED IN A "COMMAND LIST" SEGMENT
  
- THREE MEANS OF UTILIZING THE "LIMITED SERVICE SUBSYSTEM":
  - ▮ IF `init_proc = lss_login_responder_$lss_login_responder_`
  
  - ▮ LIST OF COMMANDS ARE CONTAINED IN `>sss>lss_command_list_`
  
  - ▮ `command_processor` WILL CHECK EVERY COMMAND ENTERED BY THE USER AGAINST THIS LIST - ONLY IF IT IS ON LIST WILL IT BE EXECUTED
  
  - ▮ IN ADDITION, A CPU USAGE GOVERNOR WILL BE ENABLED, LIMITING THE PROCESS TO 'ratio' CPU SECONDS PER 'interval' REAL SECONDS
    - ▮ NOTE: IF `>sss>lss_command_list_` NOT FOUND AT LOGIN TIME, MESSAGE RELAYED IS:  
  
"The system is currently unavailable".  
  
(A logout -hold IS DONE)
  
- ▮ `lss_login_responder_$limited_command_system_`
  
- ▮ SIMILAR TO ABOVE, BUT `lss_command_list_` IS SEARCHED FOR IN USER'S PROJECT DIRECTORY

MODIFYING THE PROCESS ENVIRONMENT

LIMITED SUBSYSTEMS

- ┆ THE COMMAND, `enter_1ss`
  
- ┆ USER SPECIFIES SEGMENT CONTAINING THE "COMMAND LIST"
  
- ┆ CONSIDER IF `enter_1ss` COMMAND APPEARS IN `project_start_up.ec`
  
- COMMAND TABLES ARE CREATED USING THE '`make_commands`' COMMAND
  
- ┆ `make_commands` ACCEPTS THE NAME OF AN ASCII SEGMENT WITH THE SUFFIX OF '`ct`', AND PRODUCES A COMMAND TABLE SEGMENT (THE ENTRYNAME WITHOUT THE '`ct`' SUFFIX)
  
- ┆ ASCII COMMAND LIST CONTAINS THREE TYPES OF STATEMENTS:
  - ┆ `ratio: R;`
    - ┆ SPECIFIES THE NUMBER OF CPU SECONDS MAXIMUM ALLOWED FOR THE PROCESS DURING THE SPECIFIED INTERVAL
  
  - ┆ `interval: N;`
    - ┆ SPECIFIES THE NUMBER OF REAL-TIME SECONDS WITHIN WHICH THE PROCESS IS LIMITED TO 'R' CPU SECONDS
  
  - ┆ `(command_list): pathname;`  
`command_name: pathname;`
    - ┆ SPECIFIES THAT THE COMMANDS IN THE (BLANK DELIMITED) `command_list` OR THE COMMAND SPECIFIED BY `command_name` ARE ALLOWED, AND THAT THEY SHOULD CAUSE THE PROCEDURE SPECIFIED BY `pathname` TO BE INVOKED WHEN THEY ARE ENTERED AS COMMANDS



MODIFYING THE PROCESS ENVIRONMENT

LIMITED SUBSYSTEMS

● A COMMAND LIST EXAMPLE

```
/* set ratio and interval length */
ratio:                45;
interval:             120;

/* define commands */

(addname an):        ;
calc:                ;
(delete dl):         ;
(deletename dn):    ;
(list ls):           >udd>MED>nd>list;
logout:              ;
(print pr):          ;
(program_interrupt pi): ;
(rename rn):         ;
(start sr):          ;
edit:                >sss>qedx;
```

MODIFYING THE PROCESS ENVIRONMENT

LIMITED SUBSYSTEMS

r 19:39 0.116 2

!pr xxx.ct 1

```
logout: >sss>logout;
pwd:    >sss>pwd;
ls:     >sss>ls;
new_proc: >sss>new_proc;
probe:  >sss>probe;
```

r 19:39 0.104 3

!make\_commands xxx

r 19:43 0.248 14

!ls -first 1

Segments = 72, Lengths = 80.

re 1 xxx

r 19:43 0.216 13

!enter lss xxx

r 19:44 0.065 0

!who

who is not a legal command

r 19:44 0.058 0

QUIT

r 19:44 0.118 4 level 2

!who

who is not a legal command

r 19:44 0.035 2 level 2

!pr xxx.ct

pr is not a legal command

r 19:44 0.033 0 level 2

!logout

NDibble MED logged ;out 02/26/81 1459.8 mst Thu  
CPU usage 7 sec, memory usage 23.0 units, cost \$0.48.  
hangup

MODIFYING THE PROCESS ENVIRONMENT

NEW SUBSYSTEMS

- IN THE MOST EXTREME CASE, A DESIGNER MAY IMPLEMENT HIS/HER OWN PROCESS ENVIRONMENT BY REPLACING THE STANDARD PROCESS OVERSEER WITH HIS/HER OWN PROCESS OVERSEER
  - ▮ THE DESIGNER IS WARNED TO BE SURE TO PERFORM THE CRITICAL FUNCTIONS WHICH ARE NORMALLY PERFORMED BY STANDARD OVERSEERS, BUT BEYOND THESE RESTRICTIONS, THE DESIGNER IS FREE TO IMPLEMENT ANY ENVIRONMENT DESIRED
  - ▮ 'process overseer .pl1' IS A GOOD REFERENCE FOR THE DESIGNER ATTEMPTING TO CREATE A NEW ONE (SEE APPENDIX B)

|| YOU ARE NOW READY FOR WORKSHOP ||  
#7

TOPIC XIV

Dialing Terminals to a Process

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## OVERVIEW

- NORMALLY THERE IS A ONE-TO-ONE CORRESPONDENCE BETWEEN AN INTERACTIVE PROCESS AND A TERMINAL DEVICE
  
- HOWEVER, Multics PROVIDES THE 'dial' FACILITY WHICH ENABLES A PROCESS TO CONTROL MORE THAN ONE TERMINAL DEVICE
  
- THE DIAL FACILITY IS PART OF THE ANSWERING SERVICE
  
- THE ANSWERING SERVICE IS RESPONSIBLE FOR LISTENING TO THE COMMUNICATION LINES ATTACHED TO THE FRONT-END PROCESSOR (FNP)
  
- A PROCESS HAVING THE 'dialok' ATTRIBUTE (ASSIGNABLE BY THE SYSTEM AND PROJECT ADMINISTRATORS) MAY USE THE DIAL FACILITY
  - ┆ ACCEPT DIALED TERMINALS (REQUIRES NO SPECIAL HARDWARE)
  
  - ┆ DIAL OUT TO TERMINALS (REQUIRES ACCESS TO AUTO-CALL CHANNEL)

IMPLEMENTATION OF THE DIAL FACILITY

- ipc ESTABLISHES THE EVENT CHANNEL REQUIRED FOR COMMUNICATION BETWEEN THE USER PROCESS AND THE ANSWERING SERVICE
- dial\_manager INTERFACES TO THE ANSWERING SERVICE'S DIAL FACILITY
- convert dial message INTERPRETS THE SPECIAL IPC MESSAGE SENT BY THE ANSWERING SERVICE TO A USER PROCESS
- THE 'dial' PRE-ACCESS COMMAND IS USED BY A TERMINAL OPERATOR ATTEMPTING TO DIAL IN TO AN EXISTING PROCESS (AG92)

IMPLEMENTATION OF THE DIAL FACILITY

dial manager

- A POINTER TO AN ARGUMENT STRUCTURE IS PASSED IN ALL CALLS TO `dial_manager`, AND THE MEANING OF THE STRUCTURE MEMBERS VARIES WITH EACH ENTRY POINT

```
dcl 1 dial_manager_arg based aligned,  
  2 version          fixed bin /* MUST BE SET TO 1 */  
  2 dial_qualifier   char(22),  
  2 dial_channel     fixed bin(71), /* EVENT-WAIT CHANNEL */  
  2 channel_name     char(32);
```

┆ dial\_qualifier

┆ WILL BE A 'dial\_id' TO BE SUPPLIED WHEN THE dial COMMAND IS TYPED (IF WE'RE ACCEPTING DIALS)

┆ OR WILL BE A PHONE NUMBER (IF WE'RE DIALING OUT)

┆ dial\_channel

┆ AN EVENT-WAIT CHANNEL\_ID RETURNED BY `ipc_create_ev_chn`

┆ MUST BE THE SAME FOR ALL CALLS TO `dial_manager` IN THIS PROCESS

┆ channel\_name

┆ IDENTIFIES A LINE ADAPTER AND PORT NUMBER ON THE FRONT END PROCESSOR

┆ SEE APPENDIX C

## DIALING TERMINALS TO A PROCESS

### ● STEPS INVOLVED IN DIALING TERMINALS TO A PROCESS

- ▮ PROCESS DECLARES AN EVENT-WAIT CHANNEL TO BE USED BY THE ANSWERING SERVICE TO NOTIFY THE PROCESS OF CRITICAL EVENTS (SUCH AS SOMEONE DIALING IN TO THE PROCESS, SOMEONE HANGING UP, AND SO ON)
  
- ▮ PROCESS REQUESTS THAT THE ANSWERING SERVICE NOW ALLOW DIALS FOR THE PROCESS, PASSING THE EVENT-WAIT CHANNEL\_ID AND A "DIAL QUALIFIER"
  
- ▮ PROCESS MAY NOW CONVERT WAIT-CHANNEL TO CALL CHANNEL, IF DESIRED
  
- ▮ TERMINALS ARE DIALED INTO THE PROCESS USING THE 'dial' COMMAND
  
- ▮ UPON NOTIFICATION FROM THE ANSWERING SERVICE THAT A TERMINAL HAS DIALED-IN, PROCESS MUST INTERPRET THE IPC MESSAGE PASSED, WHICH CONTAINS
  - ▮ CHANNEL-NAME (DEVICE ID) OF THE TERMINAL DIALED-IN
  
  - ▮ FLAGS INDICATING WHAT TOOK PLACE ON THE COMLINE (DIALUP, HANGUP)
  
- ▮ PROCESS NOW ATTACHES THAT DEVICE AND COMMENCES TO DO LOGICAL I/O TO THAT TERMINAL



## DIALING TERMINALS TO A PROCESS

### SUBROUTINES

#### ● dial\_manager\_\$allow\_dials

▮ REQUESTS THAT THE ANSWERING SERVICE ALLOW TERMINALS TO DIAL TO THE CALLING PROCESS

▮ THE CALLER SETS 'dial\_qualifier' IN THE dial\_manager\_arg STRUCTURE TO AN ALPHANUMERIC STRING FROM 1 TO 22 CHARACTERS

▮ THE CALLER SETS dial\_manager\_arg.dial\_channel TO THE EVENT-WAIT CHANNEL\_ID ESTABLISHED FOR COMMUNICATING WITH THE ANSWERING SERVICE (NOTE THAT FOLLOWING A CALL TO dial\_manager\_\$allow\_dials, THE CALLER MAY CHANGE THE EVENT-WAIT CHANNEL INTO AN EVENT-CALL CHANNEL IF DESIRED)

#### ● dial\_manager\_\$registered\_server

▮ SIMILAR TO dial\_manager\_\$allow\_dials

▮ PERMITS TERMINALS TO DIAL IN WITHOUT FURNISHING Personid.Projectid AS dial COMMAND ARGUMENT

▮ dial\_qualifier MUST BE REGISTERED BY SYSTEM ADMINISTRATOR

▮ CALLER MUST HAVE rw ON >sc1>rcp>dial.<dial\_qualifier>.acs

DIALING TERMINALS TO A PROCESS

SUBROUTINES

● dial\_manager\_\$shutoff\_dials

┆ INFORMS THE ANSWERING SERVICE THAT THE PROCESS WISHES TO PREVENT FURTHER DIAL CONNECTIONS, AND THAT EXISTING CONNECTIONS SHOULD BE TERMINATED

┆ ACCEPTS SAME INFORMATION AS dial\_manager\_\$allow\_dials

┆ IMPORTANT RESTRICTION: dial\_channel MUST BE AN EVENT-WAIT; CALLER MAY THEREFORE HAVE TO CALL ipc\_\$decl\_ev\_wait\_chn FIRST

● convert\_dial\_message\_\$return\_io\_module

┆ SHOULD BE INVOKED BY A PROCESS WHEN IT HAS RECEIVED A WAKEUP FROM THE ANSWERING SERVICE

┆ REQUIRES THE IPC event\_info.message AS INPUT AND RETURNS:

┆ THE DEVICE-ID (CHANNEL NAME) OF THE COMMUNICATIONS LINE THAT HAS DIALED-UP OR HUNG-UP

┆ A STRUCTURE INDICATING WHETHER THE TERMINAL IN QUESTION HAS DIALED-UP, OR HUNG-UP

DIALING TERMINALS TO A PROCESS

THE 'DIAL' COMMAND

● THE 'dial' COMMAND:

⌋ IS TYPED IN LIEU OF THE login COMMAND

⌋ IS A REQUEST TO THE ANSWERING SERVICE TO CONNECT THE TERMINAL TO AN EXISTING PROCESS AND TO NOTIFY THAT PROCESS OF THE CONNECTION

⌋ USAGE:

dial dial\_id {Person\_id.Project\_id}

⌋ THE USER MUST SPECIFY THE 'dial\_id' THAT WAS PASSED TO THE ANSWERING SERVICE BY THE PROCESS ACCEPTING DIALS

⌋ THE USER MUST ALSO SPECIFY THE 'Person\_id.Project\_id' OF THE EXISTING PROCESS, UNLESS dial\_manager\_\$registered\_server WAS ORIGINALLY USED TO ALLOW DIALS

DIALING TERMINALS TO A PROCESS

AN EXAMPLE

```
SET_UP_DIAL:      proc;
dcl ipc_$create_ev_chn entry (fixed bin (71), fixed bin (35)),
    ipc_$delete_ev_chn entry (fixed bin (71), fixed bin (35)),
    ipc_$decl_ev_call_chn entry (fixed bin (71), entry,
    ptr, fixed bin, fixed bin (35)),
    dial_manager_$allow_dials entry (ptr, fixed bin (35)),
    (ioa_, com_err_, ioa_$ioa_switch) entry options (variable);
dcl code fixed bin (35);
    ME char (12) varying init ("SET_UP_DIAL") ;
dcl 1 dial_manager_arg aligned static,
    2 version fixed bin init (1),
    2 dial_qualifier char (22) init ("astra"),
    2 dial_channel fixed bin (71),
    2 channel_name char (32) ;

    call ioa_ ("Begin ^a", ME);
    call ipc_$create_ev_chn (dial_manager_arg.dial_channel,
    code);
    if code ^= 0 then call ERROR (1);
    call dial_manager_$allow_dials (addr (dial_manager_arg),
    code);
    if code ^= 0 then call ERROR (2);
    call ipc_$decl_ev_call_chn (dial_manager_arg.dial_channel,
    DIAL_HANDLER, null(), 0, code);
    if code ^= 0 then call ERROR (3);

    call ioa_ ("Now listening for dials: ^a", ME);
    return;
```

DIALING TERMINALS TO A PROCESS

AN EXAMPLE

```
DIAL_HANDLER: entry (info_ptr);
dcl info_ptr ptr parameter;
dcl 1 event_info based (info_ptr),
    2 channel_id fixed bin (71),
    2 message fixed bin (71),
    2 sender bit (36),
    2 origin,
    3 dev_signal bit (18) unal,
    3 ring bit (18) unal,
    2 data_ptr ptr;

dcl convert_dial_message$return_io_module entry (fixed bin(71),
    char(*), char(*), fixed bin, 1 aligned, 2 bit(1) unal,
    2 bit(1) unal, 2 bit(1) unal, 2 bit(33) unal, fixed bin(35));
dcl which_channel char (32);
dcl iocb_ptr ptr;

dcl ipc_scutoff entry (fixed bin(71), fixed bin(35));
dcl ipc_reconnect entry (fixed bin(71), fixed bin(35));

dcl iox_attach_name entry (char(*), ptr, char(*), ptr, fixed bin(35));
dcl iox_open entry (ptr, fixed bin, bit(1) aligned, fixed bin(35));
dcl iox_close entry (ptr, fixed bin(35));
dcl iox_detach_iocb entry (ptr, fixed bin(35));
dcl iox_control entry (ptr, char(*), ptr, fixed bin(35));
dcl iox_get_line entry (ptr, ptr, fixed bin(21), fixed bin(21),
    fixed bin(35));

dcl buffer char(80);
dcl actually_read char(n read) based (addr(buffer));
dcl n_read fixed bin (21);

    ME = "DIAL_HANDLER";

    call ipc_scutoff (dial_manager arg.dial_channel, code);
    if code ^= 0 then call ERROR (4);

    call convert_dial_message$return_io_module (
        event_info.message, which_channel, "", 0, "0"b, code);
    if code ^= 0 then call ERROR (5);
```

DIALING TERMINALS TO A PROCESS

AN EXAMPLE

```
call iox_$attach_name ("switch", iocb_ptr,
                       "tty "||which_channel, null(),code);
if code ^= 0 then call ERROR (6);

call iox $open (iocb_ptr, 3, "0"b, code);
if code ^= 0 then call ERROR (7);

call ioa_$ioa_switch (iocb_ptr, "Welcome to my world.
Please type a line and I will echo it back.");
call iox $get_line (iocb_ptr, addr(buffer), 80, n_read, code);
call ioa_$ioa_switch (iocb_ptr, "^a", actually_read);
call ioa_$ioa_switch (iocb_ptr, "Good bye");

call iox $control (iocb_ptr, "hangup", null(), code);
if code ^= 0 then call ERROR (8);

call iox $close (iocb_ptr, code);
if code ^= 0 then call ERROR (9);

call iox $detach_iocb (iocb_ptr, code);
if code ^= 0 then call ERROR (10);
return;
```

DIALING TERMINALS TO A PROCESS

AN EXAMPLE

```
ERROR:      proc (error_number);
            /* Internal proc to report errors */
dcl  error_number;
            call com_err_ (code, ME, "Check call ^i of ERROR",
                          error_number);
            goto FINISH;
end ERROR;

SHUTOFF:  entry;

dcl  dial_manager $shutoff_dials entry (ptr, fixed bin (35));
dcl  ipc $decl_ev_wait_chn entry (fixed bin (71), fixed bin (35));

            ME = "SHUTOFF";

            call ipc $decl_ev_wait_chn (dial_manager_arg.dial_channel, code);
            if code ^= 0 then call ERROR (11);
            call dial_manager $shutoff_dials (addr (dial_manager_arg), code);
            if code ^= 0 then call ERROR (12);
            call ipc $delete_ev_chn (dial_manager_arg.dial_channel, code);
            if code ^= 0 then call ERROR (13);
            return;

FINISH:
end SET_UP_DIAL;
```

DIALING TERMINALS TO A PROCESS

AN EXAMPLE

- THE PRECEDING EXAMPLE IS VERY SIMPLE AND THEREFORE HAS LITTLE PRACTICAL APPLICATION

▮ OBVIOUS PROBLEMS:

1. IT ONLY HANDLES ONE DIALED IN TERMINAL AT A TIME
2. IT CANNOT HANDLE THE SITUATION IN WHICH THE DIALED IN TERMINAL SIMPLY "HANGS UP"
3. IT CANNOT HANDLE THE SITUATION IN WHICH THE DIALED IN TERMINAL SIGNALS QUIT (USER HITS BREAK KEY)
4. THE MASTER PROCESS GOES BLOCKED WHILE WAITING FOR INPUT FROM THE "SLAVE" TERMINAL

- A 10 PAGE EXAMPLE APPEARS IN APPENDIX G

▮ THIS EXAMPLE SOLVES ALL OF THE ABOVE PROBLEMS

▮ IT PROVIDES A STARTING POINT FOR A REALISTIC DIAL IN APPLICATION



DIALING OUT TO A TERMINAL

- STEPS INVOLVED IN DIALING OUT
  - ▮ ESTABLISH EVENT-WAIT CHANNEL
  - ▮ REQUEST ANSWERING SERVICE TO DIAL A SPECIFIED PHONE NUMBER
  - ▮ ONE CAN CHANGE EVENT-WAIT TO EVENT-CALL CHANNEL AT THIS TIME
  - ▮ AFTER NOTIFICATION OF SUCCESSFUL DIAL-OUT, USER ATTACHES DEVICE AND DOES LOGICAL I/O
  
- A PRIVILEGED PROCESS (HAVING 'rw' ON THE APPROPRIATE ACCESS CONTROL SEGMENT) MAY USE THE DIAL FACILITY TO DIAL OUT TO A TERMINAL

DIALING OUT TO A TERMINAL  
DIAL MANAGER ENTRY POINTS

● ENTRY POINTS USED TO DIAL OUT:

⌋ dial\_manager\_\$dial\_out

- ⌋ REQUESTS THAT AN AUTO-CALL CHANNEL BE DIALED TO A GIVEN TELEPHONE NUMBER, AND, IF THE CHANNEL IS SUCCESSFULLY DIALED, THAT THE CHANNEL BE ASSIGNED TO THE REQUESTING PROCESS
- ⌋ THE CALLER SETS dial\_manager\_arg.dial\_qualifier TO THE TELEPHONE NUMBER TO BE DIALED (NONNUMERIC CHARACTERS IN THE NUMBER ARE IGNORED, SO THE NUMBER MAY BE SPECIFIED AS, FOR INSTANCE, "301/977-4292")
- ⌋ dial\_manager\_arg.dial\_channel IS SET TO THE EVENT-WAIT CHANNEL CREATED TO ALLOW THE ANSWERING SERVICE TO COMMUNICATE WITH THE PROCESS
- ⌋ THE CALLER MAY SET dial\_manager\_arg.channel\_name TO A SPECIFIC CHANNEL-NAME OF AN AUTO-CALL CHANNEL - IF THE CALLER ASSIGNS THE NULL STRING TO THIS ARGUMENT, THE ANSWERING SERVICE WILL ATTEMPT TO ASSIGN ANY AVAILABLE AUTO-CALL CHANNEL AND CALLER MUST USE convert\_dial\_message\_\$return\_io\_module TO DETERMINE WHAT IT IS

⌋ dial\_manager\_\$terminate\_dial\_out

- ⌋ REQUESTS THAT THE ANSWERING SERVICE HANG UP AN AUTO-CALL LINE AND UNASSIGN IT FROM THE REQUESTING PROCESS
- ⌋ ACCEPTS THE SAME INFO AS dial\_manager\_\$dial\_out
- ⌋ HOWEVER, THE 'channel\_name' ARGUMENT MUST BE SUPPLIED WITH THE NAME OF THE AUTO-CALL CHANNEL WHICH WAS USED FOR THE DIAL OUT

TOPIC XV

Message Segment Facility

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Layered Design . . . . .	15-6
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Extended Access. . . . .	15-9
message_segment Subroutine Summary. . . . .	15-10
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WHAT IS IT?

● A SERIES OF PRIMITIVES, SUBROUTINES AND COMMANDS

● DESIGNED TO

▮ MANIPULATE RING 1 MESSAGE SEGMENTS

▮ FACILITATE PROTECTED AND ORDERED MESSAGE EXCHANGE BETWEEN AND WITHIN PROCESSES

▮ SALVAGE MESSAGE SEGMENTS CONTAINING "DAMAGED" MESSAGES

▮ MINIMIZE WRITE/UPDATE WINDOW TIME

▮ SUPPORT CHANGE-ABLE MESSAGE SIZE

## APPLICATIONS

- I/O AND ABSENTEE DAEMON QUEUES
- SUPPORT FOR MAIL AND SEND\_MESSAGE FACILITIES
- USER-DESIGNED APPLICATIONS REQUIRING THE SPECIAL CAPABILITIES OF MESSAGE SEGMENTS

THE MESSAGE SEGMENT

● PROPERTIES

▮ ACCESSIBLE ONLY IN RING 1 (AND RING 0)

▮ HAS A SUFFIX

▮ 'ms' FOR QUEUE MESSAGE SEGMENTS

▮ 'mbx' FOR MAILBOXES

▮ MUST BE A SINGLE-SEGMENT FILE

▮ HAS AN EXTENDED ACCESS CONTROL LIST

## THE MESSAGE SEGMENT

### ● STRUCTURE

#### ┆ HEADER

┆ LOCK WORD

┆ 36-BIT MESSAGE SEGMENT ID BIT PATTERN

┆ OFFSET TO FIRST MESSAGE

┆ OFFSET TO LAST MESSAGE

┆ MESSAGE COUNT

┆ SWITCHES

┆ MSEG INCONSISTENT

┆ MSEG HAS BEEN SALVAGED

┆ ALLOCATION BIT STRING SAYS WHICH BLOCKS ARE USED

┆ LENGTH (ALLOCATION BIT STRING)

┆ MESSAGE BLOCK SIZE

┆ UNUSED BLOCK COUNT

┆ ...AND OTHER INFO

THE MESSAGE SEGMENT

▮ DOUBLY-THREADED LIST OF MESSAGES

▮ EACH MESSAGE IS COMPRISED OF 1 OR MORE FIXED-LENGTH BLOCKS

▮ EACH BLOCK HAS A HEADER CONTAINING

▮ OFFSET TO NEXT BLOCK IN MESSAGE (OR ZERO)

▮ A "FIRST-BLOCK" SWITCH

▮ NUMBER OF MESSAGE BITS IN BLOCK

▮ FIRST BLOCK IN MESSAGE ALSO HAS A TRAILER

▮ EACH MESSAGE TRAILER CONTAINS

▮ MESSAGE BIT SIZE

▮ TIME MESSAGE WAS SENT

▮ VALIDATION LEVEL OF SENDER

▮ Personid.Projectid OF SENDER



## LAYERED DESIGN

- TWO HIGH-LEVEL SUBROUTINE INTERFACES ALREADY EXIST FOR MANIPULATION OF MESSAGE SEGMENTS

▮ message\_segment\_ FOR QUEUE MESSAGE SEGMENTS

▮ mailbox\_ FOR MAILBOXES

- TWO CORRESPONDING COMMAND-SETS EXIST AS WELL

▮ FOR QUEUE MSEGS WE HAVE:

ms\_add\_name, msan

ms\_create, mscr

ms\_delete, msdl

ms\_delete\_acl, msda

ms\_delete\_name, msdn

ms\_list\_acl, msla

ms\_rename, msrn

ms\_set\_acl, mssa

LAYERED DESIGN

┆ FOR MAILBOXES WE HAVE:

    mbx\_add\_name, mban

    mbx\_create, mbr

    mbx\_delete, mbd1

    mbx\_delete\_acl, mbda

    mbx\_delete\_name, mbdn

    mbx\_list\_acl, mbla

    mbx\_rename, mbrn

    mbx\_set\_acl, mbsa

● message\_segment\_ AND mailbox\_ ARE GATES INTO THE ADMINISTRATIVE RING

┆ WHICH TRANSFER CONTROL TO THE PROCEDURES queue\_mseg\_ AND mbx\_mseg\_,  
RESPECTIVELY

┆ queue\_mseg\_ AND mbx\_mseg\_ IN TURN CALL MODULES IN THE PRIMITIVE  
MESSAGE SEGMENT FACILITY

LAYERED DESIGN  
PRIMITIVE MESSAGE SEGMENT FACILITY

● THE PRIMITIVE MESSAGE FACILITY IS COMPRISED OF MODULES WHICH

▮ CREATE AND DELETE MSEGS

▮ MANIPULATE EXTENDED ACCESS

▮ LOCK AND UNLOCK MSEGS

▮ MANIPULATE MESSAGES

▮ MANIPULATE 'OWN' MESSAGES

▮ SALVAGE MSEGS

▮ CONVERT MSEGS FROM A PREVIOUS FORMAT

EXTENDED ACCESS

● BOTH QUEUE AND MAILBOX MSEGs EMPLOY THESE ATTRIBUTES:

! a

ALLOWS USER TO ADD A MESSAGE

! d

ALLOWS USER TO DELETE ANY MESSAGE

! r

ALLOWS USER TO READ ANY MESSAGE

! o

ALLOWS USER TO READ/DELETE 'OWN' MESSAGES

! s

ALLOWS USER TO DETERMINE WHETHER MSEG HAS BEEN SALVAGED AND MESSAGE COUNT

● IN ADDITION, MAILBOX MESSAGE SEGMENTS EMPLOY:

! w

ALLOWS USER TO SEND NORMAL WAKEUP WHEN ADDING MESSAGE

MESSAGE SEGMENT SUBROUTINE SUMMARY

● CREATING AND DELETING QUEUE MESSAGE SEGMENTS

message\_segment\_\$create

message\_segment\_\$delete

● MANIPULATING EXTENDED ACCESS

message\_segment\_\$ms\_acl\_add

message\_segment\_\$ms\_acl\_delete

message\_segment\_\$ms\_acl\_list

message\_segment\_\$ms\_acl\_replace

● RENAMING

message\_segment\_\$chname\_file

● OPENING AND CLOSING

message\_segment\_\$open

message\_segment\_\$close

MESSAGE SEGMENT SUBROUTINE SUMMARY

● OBTAINING HEADER STATUS INFO

message\_segment \$check\_salv\_bit\_index  
message\_segment\_\$check\_salv\_bit\_file

message\_segment \$get\_message\_count\_index  
message\_segment\_\$get\_message\_count\_file

● OBTAINING EFFECTIVE ACCESS

message\_segment \$get\_mode\_index  
message\_segment\_\$get\_mode\_file

● MANIPULATING MESSAGES

message\_segment \$add\_index  
message\_segment\_\$add\_file

message\_segment \$delete\_index  
message\_segment\_\$delete\_file

message\_segment \$read\_index  
message\_segment\_\$read\_file

message\_segment \$incremental\_read\_index  
message\_segment\_\$incremental\_read\_file

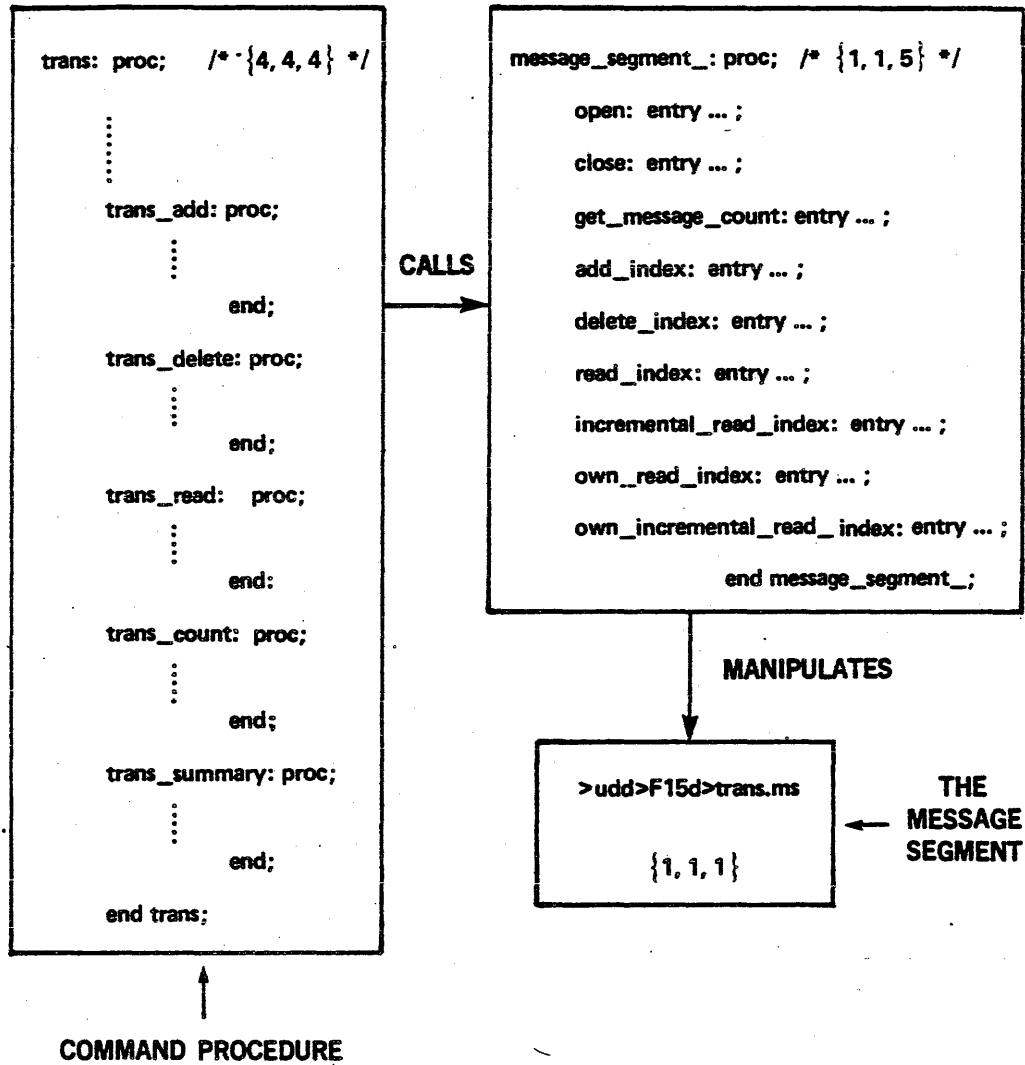
message\_segment \$update\_message\_index  
message\_segment\_\$update\_message\_file

● MANIPULATING 'OWN' MESSAGES

message\_segment \$own\_read\_index  
message\_segment\_\$own\_read\_file

message\_segment \$own\_incremental\_read\_index  
message\_segment\_\$own\_incremental\_read\_file

MESSAGE SEGMENT FACILITY ILLUSTRATIVE EXAMPLE



**MESSAGE SEGMENT FACILITY  
ILLUSTRATIVE EXAMPLE**

TOPIC XVI

Program Library Management

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## INTRODUCTION

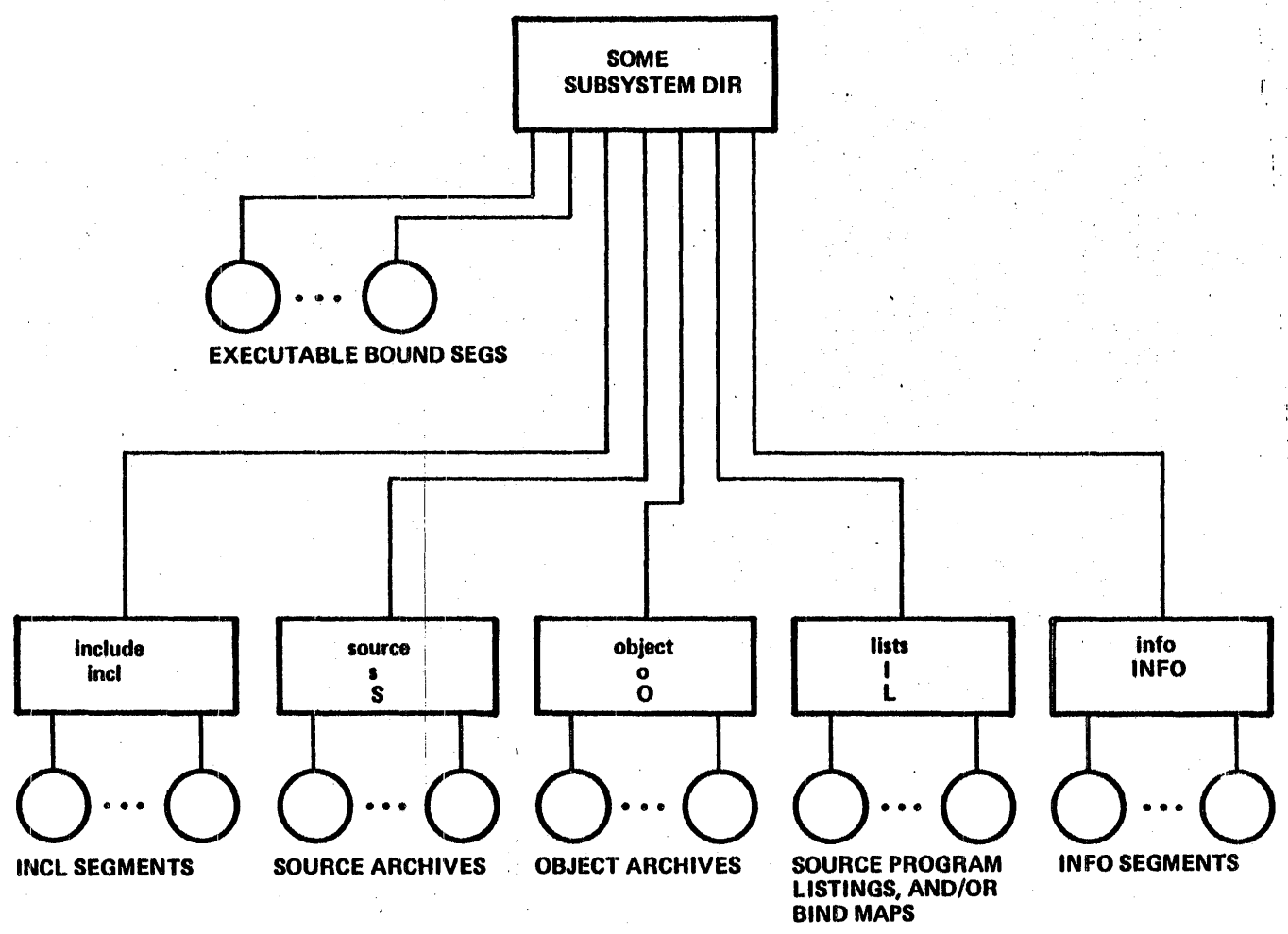
- LARGE AND COMPLEX SUBSYSTEMS REQUIRE GOOD PROGRAM LIBRARY MANAGEMENT TECHNIQUES
  - ▮ THE DESIGNER MUST BE CONCERNED WITH PROPERLY ORGANIZING THE SOURCE PROGRAMS, OBJECT PROGRAMS, BOUND SEGMENTS, LISTINGS, AND SO ON
  - ▮ THE DESIGNER COULD DEVELOP HIS OWN LIBRARY CONVENTIONS AND TOOLS, BUT:
  
- CONVENTIONS AND SYSTEM-PROVIDED TOOLS EXIST FOR
  - ▮ ORGANIZING SOURCE, OBJECT, EXECUTABLE, AND DOCUMENTATION LIBRARIES IN A CONVENIENT MANNER
  - ▮ MANIPULATING THE COMPONENTS OF A USER MAINTAINED LIBRARY IN A CONTROLLED MANNER
  - ▮ CONTROLLING THE INSTALLATION AND DE-INSTALLATION OF SUBSYSTEM MODULES IN AN ORDERLY MANNER

## INTRODUCTION

### ORGANIZATION OF PROGRAM LIBRARIES

- A USER'S PROGRAM LIBRARY FOR A GIVEN SUBSYSTEM IS GENERALLY ORGANIZED AS A DIRECTORY SUBTREE
  - ┆ A DIRECTORY SEGMENT, NAMED FOR THE LIBRARY ITSELF, SERVES AS THE ROOT OF THE SUBTREE
  - ┆ EXECUTABLE PROGRAMS, WHETHER STAND-ALONE OR BOUND, RESIDE UNDER THE LIBRARY ROOT DIRECTORY
  - ┆ SUBDIRECTORIES UNDER THIS ROOT CONTAIN:
    - ┆ SOURCE PROGRAMS, EITHER INDIVIDUALLY OR IN ARCHIVE SEGMENTS
    - ┆ OBJECT PROGRAMS, EITHER INDIVIDUALLY, OR (MORE GENERALLY) IN ARCHIVE SEGMENTS
    - ┆ LISTINGS AND/OR BIND MAPS
    - ┆ INCLUDE FILES
    - ┆ HELP FILES

### LIBRARY SUBTREE



INTRODUCTION  
NAMING CONVENTIONS

- THE LIBRARY COMPONENTS MENTIONED ABOVE ARE GENERALLY NAMED ACCORDING TO THE FOLLOWING STANDARD NAMING CONVENTIONS:

- ▮ THE SOURCE SUBDIRECTORY IS GENERALLY GIVEN THE NAMES

source  
S  
S

- ▮ THE OBJECT PROGRAM SUBDIRECTORY IS GENERALLY GIVEN THE NAMES

object  
O  
O

- ▮ THE LISTINGS SUBDIRECTORY IS GENERALLY GIVEN THE NAMES

lists  
L  
L

- ▮ THE INCLUDE FILE SUBDIRECTORY IS GENERALLY GIVEN THE NAMES

include  
incl

- ▮ THE HELP FILES SUBDIRECTORY IS GENERALLY GIVEN THE NAMES

info  
INFO

INTRODUCTION  
NAMING CONVENTIONS

● NAMING CONVENTIONS FOR BOUND SEGMENTS AND CORRESPONDING ARCHIVES

- ▮ THE BOUND SEGMENT ITSELF IS GIVEN AN ENTRYNAME "bound ????????", WHERE ???????? IS A NAME CHOSEN BY THE DESIGNER (E.G., bound\_command\_loop\_)
  
- ▮ THE ARCHIVE WHICH CONTAINS THE SOURCE PROGRAMS USED TO GENERATE THE INDIVIDUAL COMPONENTS OF THE BOUND SEGMENT IS NAMED "bound ????????.s.archive" (E.G., bound\_command\_loop\_.s.archive)
  
- ▮ THE ARCHIVE WHICH CONTAINS THE OBJECT PROGRAMS AND WHICH WAS INPUT TO THE BINDER IS NAMED "bound ????????.archive" (E.G., bound\_command\_loop\_.archive)

INTRODUCTION

A TYPICAL PROGRAM LIBRARY

```
!list >udd>F15dw>Auerbach>user_library_1 -all
```

Directories = 1.

```
sma user_library_1
    ul1
```

```
!list -pn >udd>F15dw>Auerbach>ul1 -all
```

Segments = 3, Lengths = 4.

```
re      2 bound_cde_
         c
         d
         e
re      1 b
re      1 a
```

Directories = 5.

```
sma include
    incl
sma info
    INFO
sma lists
    l
    L
sma object
    o
    O
sma source
    s
    S
```

```
!cwd >udd>F15dw>Auerbach>ul1
```

INTRODUCTION

A TYPICAL PROGRAM LIBRARY

!ls -pn source -all

Segments = 3, Lengths = 3.

```
r w    1  bound_cde_.s.archive
          c.pl1
          d.pl1
          e.pl1
r w    1  b.fortran
r w    1  a.pl1
```

!ls -pn object -all

Segments = 3, Lengths = 4.

```
r w    2  bound_cde_.archive
          c
          d
          e
          bound_cde_.bind
r w    1  b
r w    1  a
```

!ls -pn L -all

Segments = 6, Lengths = 6.

```
r w    1  bound_cde_.list
r w    1  a.list
r w    1  b.list
r w    1  c.list
r w    1  d.list
r w    1  e.list
```

INTRODUCTION

A TYPICAL PROGRAM LIBRARY

```
!ls -pn INFO -all
```

```
Segments = 6, Lengths = 6.
```

```
r w    1  user_library_1.gi.info  
r w    1  e.info  
r w    1  d.info  
r w    1  c.info  
r w    1  b.info  
r w    1  a.info
```

```
!ls -pn include -all
```

```
Segments = 3, Lengths = 3.
```

```
r w    1  DATABASE_STRUCTURE.incl.pl1  
r w    1  REC2.incl.pl1  
r w    1  REC1.incl.pl1
```



PROGRAM LIBRARY MANAGEMENT TOOLS

- LIBRARY ADMINISTRATOR TOOLS EXIST TO PROPERLY UPDATE AND MANIPULATE LIBRARIES IN A STRICTLY CONTROLLED AND CONSISTENT MANNER

- THE MAJOR TOOLS CAN BE CLASSIFIED AS FOLLOWS:

- ▮ INSTALLATION TOOLS

- ▮ PROGRAM LIBRARY MANIPULATION TOOLS

PROGRAM LIBRARY MANAGEMENT TOOLS

INSTALLATION TOOLS

● THE "INSTALLATION PROBLEM"

- ▮ ARISES FROM ATTEMPTS TO DYNAMICALLY INSTALL A NEW OR REPLACEMENT VERSION OF A HEAVILY USED SUBSYSTEM MODULE (OR MODULES)
  
- ▮ THOSE USERS CURRENTLY EXECUTING THE (NOW) OBSOLETE MODULES MUST CONTINUE TO EXECUTE THEM UNTIL THEY HAVE COMPLETED THEIR SESSION - IN ADDITION, ANY USERS WHO SUBSEQUENTLY ATTEMPT TO EXECUTE THE MODULE SHOULD RECEIVE THE NEW, UPDATED VERSION
  
- ▮ ANY MODIFICATIONS TO THE PROGRAM LIBRARY SHOULD BE CAREFULLY DOCUMENTED OR LOGGED

PROGRAM LIBRARY MANAGEMENT TOOLS

INSTALLATION TOOLS

- SOLUTION TO PROBLEM: Multics INSTALLATION SYSTEM (MIS)

▮ MIS SUBROUTINES ARE

- ▮ RESTARTABLE ACROSS A SYSTEM OR PROCESS FAILURE (AS LONG AS STORAGE SYSTEM IS INTACT)
- ▮ REVERSIBLE, ALLOWING FOR "DE-INSTALLATION" IF TROUBLE ARISES MIDSTREAM

▮ MIS FEATURES

- ▮ PLANNED AUTOMATIC RECOVERY (VIA DE-INSTALL ENTRY POINTS) FOR ERRORS LIKE `record_quota_overflow`, `namedup`, `entry_not_found`
- ▮ AUTOMATIC DOCUMENTATION OF AN INSTALLATION
- ▮ A COMMAND INTERFACE: `update_seg`

PROGRAM LIBRARY MANAGEMENT TOOLS

INSTALLATION TOOLS

- update seg IS USED TO DEFINE THE CONTENTS OF A MODIFICATION, AND TO INSTALL OR DE-INSTALL THAT MODIFICATION IN A LIBRARY
  
- ▮ A MODIFICATION IS A GROUP OF PHYSICALLY OR LOGICALLY RELATED SEGMENTS WHICH MUST BE INSTALLED IN A LIBRARY AT THE SAME TIME IN ORDER TO MAINTAIN LIBRARY CONSISTENCY AND INTEGRITY
  
- ▮ SOURCE AND OBJECT ARE PHYSICALLY RELATED
  
- ▮ OBJECT AND OBJECT ARE LOGICALLY RELATED
  
- ▮ A MODIFICATION IS INSTALLED THUSLY:
  - ▮ THE INSTALLATION OF EACH SEGMENT IS DIVIDED INTO A SERIES OF STEPS (GETTING A UNIQUE ID, NAMES, AND ACL OF THE NEW AND OLD SEGMENTS, COPYING THE TARGET SEGMENT, ADDING TO AND DELETING FROM THE TARGET SEGMENT'S NAMES, FREEING NAMES ON THE OLD SEGMENT, ETC.)
  
  - ▮ ONE STEP AT A TIME IS PERFORMED FOR ALL SEGMENTS OF THE MODIFICATION BEFORE MOVING ON TO THE NEXT STEP
  
  - ▮ THE EXECUTABLE SEGMENTS ARE INSTALLED LAST, AS A GROUP, AFTER INSTALLING THE OTHER SEGMENTS IN THE MODIFICATION (SOURCE SEGMENTS, ARCHIVES, ETC.)
  
- ▮ THE INSTALLATION WINDOW CAN BE REDUCED TO LESS THAN ONE MINUTE PER MODIFICATION, AND IS USUALLY ABOUT FIVE SECONDS

PROGRAM LIBRARY MANAGEMENT TOOLS

INSTALLATION TOOLS

● OPERATIONS PERFORMED BY update\_seg:

┆ CREATING MODIFICATIONS

initiate  
set\_defaults  
print\_defaults

┆ DEFINING OPERATIONS TO BE PERFORMED DURING THE MODIFICATION

add  
delete  
move  
replace

┆ LISTING THE DEFINED MODIFICATION

print  
list

┆ INSTALLING/DE-INSTALLING THE MODIFICATION

install  
de\_install

┆ CLEARING THE CURRENT MODIFICATION

clear

PROGRAM LIBRARY MANAGEMENT TOOLS

INSTALLATION TOOLS

- NONPRIVILEGED USERS OF update\_seg SHOULD FIRST TYPE:

initiate [wh hcs\_] installation\_tools\_

OTHERWISE ENTRY POINTS IN installation\_tools\_ WILL BE CALLED BY  
update\_seg AND MOST USERS HAVE NULL ACCESS TO THIS SEGMENT

PROGRAM LIBRARY MANAGEMENT TOOLS

INSTALLATION TOOLS

```
!list -first 3
```

```
Segments = 74, Lengths = 62.
```

```
r w    0  04/06/81.audit
rew    1  test
rew    1  test.pl1
```

```
!cwd junk
```

```
!list
```

```
Directory empty: >user_dir_dir>MED>NDibble>junk
```

```
!cwd <
```

```
!us print_defaults
```

```
Global defaults
ring brackets:
  1,5,5
ACL:
  re      *.*.*
```

```
!us initiate example -rb 4 4 4
```

```
!us print_defaults
```

```
Defaults for >user_dir_dir>MED>NDibble>example.io
ring brackets:
  4,4,4
ACL:
  re      *.*.*
```

```
Global defaults
ring brackets:
  1,5,5
ACL:
  re      *.*.*
```

```
!list -first 4
```

```
Segments = 75, Lengths = 126.
```

```
r w    64  example.io
r w    0  04/06/81.audit
rew    1  test
rew    1  test.pl1
```

PROGRAM LIBRARY MANAGEMENT TOOLS

INSTALLATION TOOLS

!us add test.pl1 junk>==

!us move test junk>==

!us print

```
Add      >user_dir_dir>MED>NDibble>test.pl1
as        >user_dir_dir>MED>NDibble>junk>test.pl1
Set ring brackets:
      4,4,4
Access control list:
re      *.*
Names:
      test.pl1
```

```
Move      >user_dir_dir>MED>NDibble>test
to        >user_dir_dir>MED>NDibble>junk>test
Access control list:
rew     NDibble.MED.*
rew     NDibble.*.*
rw      *.SysDaemon.*
Names:
      test
```

!list -first 2

Segments = 75, Lengths = 126.

```
r w 64 example.io
r w 0 04/06/81.audit
```

!us list

!list -first 2

Segments = 76, Lengths = 127.

```
r w 1 example.il
r w 64 example.io
```



PROGRAM LIBRARY MANAGEMENT TOOLS

INSTALLATION TOOLS

!print example.il 1

INSTALLATION OBJECT SEGMENT >user\_dir\_dir>MED>NDibble>example.io

Listed on: 04/06/81 0728.3 mst Mon  
Created by: NDibble.MED.\*  
Created with: update seg (MIS Version 1.5)  
Created on: 04/06/81 0727.1 mst Mon

SUMMARY OF THE INSTALLATION:

Add >user\_dir\_dir>MED>NDibble>test.pl1  
as >user\_dir\_dir>MED>NDibble>junk>test.pl1

Move >user\_dir\_dir>MED>NDibble>test  
to >user\_dir\_dir>MED>NDibble>junk>test

INSTALLATION OBJECT SEGMENT HAS NOT BEEN INSTALLED.

A DESCRIPTION OF THE INSTALLATION FOLLOWS.

INSTALLATION DESCRIPTION:

Add >user\_dir\_dir>MED>NDibble>test.pl1  
as >user\_dir\_dir>MED>NDibble>junk>test.pl1  
Set ring brackets:  
4,4,4  
Access control list:  
re \*.\*.\*  
Names:  
test.pl1

Move >user\_dir\_dir>MED>NDibble>test  
to >user\_dir\_dir>MED>NDibble>junk>test  
Access control list:  
rew NDibble.MED.\*  
rew NDibble.\*.\*  
rw \*.SysDaemon.\*  
Names:  
test

INSTALLATION TOOLS

```
!us install
Beginning installation of example.io
Error: Linkage error by upd_ring_task_$set!1000
(>system_library_tools>bound_mis_)
referencing installation_tools_|set_ring_brackets
Incorrect access on entry.
```

```
!list -first 5
```

```
Segments = 76, Lengths = 127.
```

```
r w    1  example.il
r w   64  example.io
r w    0  04/06/81.audit
rew    1  test
rew    1  test.pl1
```

```
!cwd junk
```

```
!list
```

```
Segments = 2, Lengths = 2.
```

```
r w    1  !BBBJKzgHHgZMDK
r w    1  !BBBJKzgHHFjDLd
```

```
!in [wh hcs_] installation_tools_
```

```
!us de_install
```

```
update_seg: The lock was already locked by this process.
Non-fatal error encountered while locking
  >user_dir_dir>MED>NDibble>example.io.
update_seg will continue performing the de_install function.
Non-special target segments deleted.
De-installation complete.
```

```
!list
```

```
Directory empty: >user_dir_dir>MED>NDibble>junk
```

```
!us install
```

```
Beginning installation of example.io
Installation complete.
```

PROGRAM LIBRARY MANAGEMENT TOOLS

INSTALLATION TOOLS

!list

Segments = 2, Lengths = 2.

```
rew 1 test
re 1 test.pl1
```

!cwd <

!list -first 5

Segments = 76, Lengths = 127.

```
r w 1 example.il
r w 64 example.io
r w 0 04/06/81.audit
rew 1 test.1
      !BBBJKzgHHgZM11
rew 1 test.pl1
```

!cob test.1 junk>test

```
>user dir dir>MED>NDibble>test.1: (segment 1)
03/14781 -1121.1 mst Sat PL/I
```

```
>user dir dir>MED>NDibble>junk>test: (segment 2)
03/14781 -1121.1 mst Sat PL/I
```

The 2 segments match.

!us list

!dp example.il

PROGRAM LIBRARY MANAGEMENT TOOLS

LIBRARY DESCRIPTOR TOOLS

- THE 'library descriptor' COMMANDS ARE A COLLECTION OF TOOLS ALLOWING THE SUBSYSTEM DESIGNER OR LIBRARY ADMINISTRATOR TO MANIPULATE LIBRARY STRUCTURES

┆ ALL REFERENCE 'library\_descriptors', WHICH ARE

┆ SPECIAL SEGMENTS THAT

┆ DESCRIBE THE STRUCTURE OF LIBRARIES IN THE HIERARCHY

┆ LIST THOSE LIBRARY DESCRIPTOR COMMANDS WHICH MAY BE USED ON THE DESCRIBED LIBRARIES

┆ NAME THE PROCEDURES WHICH "KNOW" HOW TO OPERATE ON THE DESCRIBED LIBRARIES

┆ CREATED IN A TWO STEP OPERATION

┆ ASCII DESCRIPTOR SOURCE SEGMENT IS TRANSLATED INTO AN alm SEGMENT BY library\_descriptor\_compiler COMMAND PROCEDURE

┆ alm ASSEMBLER GENERATES BINARY LIBRARY DESCRIPTOR

┆ THE COMMANDS ARE

┆ library\_fetch, lf

┆ COPIES SPECIFIED ENTRIES FROM A LIBRARY DEFINED BY THE "CURRENT LIBRARY DESCRIPTOR" INTO THE USER'S WORKING DIRECTORY

┆ ACCEPTS THE STAR CONVENTION

┆ HAS SOME USEFUL OPTIONS

┆ CAN TELL YOU WHERE MATCHING ENTRY WAS FOUND (-long)

┆ CAN BE TOLD WHERE TO PUT FETCHED ENTRIES AND WHAT TO CALL THEM (-into)

PROGRAM LIBRARY MANAGEMENT TOOLS

LIBRARY DESCRIPTOR TOOLS

▮ CAN COPY THE ENTIRE ARCHIVE CONTAINING THE MATCHING ENTRIES, AS OPPOSED TO JUST SOME OF THE ARCHIVE COMPONENTS (-container)

▮ CAN INDIVIDUALLY COPY ALL COMPONENTS OF ARCHIVES CONTAINING THE MATCHING ENTRIES (-components)

▮ library\_print

▮ SELECTS PRINTABLE ENTRIES FROM A LIBRARY DEFINED BY THE CURRENT LIBRARY DESCRIPTOR AND WRITES THEM TO A FILE SUITABLE FOR DPRINTING

▮ DPRINT CONTAINS AN INDEX

▮ ACCEPTS THE STAR CONVENTION

▮ USEFUL OPTIONS

▮ -container

▮ -components

▮ CAN PRINT CUSTOMIZED PAGE FOOTINGS (-footer) AND FIRST PAGE HEADING (-header)

▮ library\_info, li

▮ RETURNS TO THE TERMINAL STATUS INFORMATION ABOUT SPECIFIED ENTRIES IN LIBRARY DEFINED BY CURRENT LIBRARY DESCRIPTOR

▮ ACCEPTS THE STAR CONVENTION

▮ USEFUL OPTIONS

▮ -container

▮ -components

PROGRAM LIBRARY MANAGEMENT TOOLS

LIBRARY DESCRIPTOR TOOLS

- ┌ library\_map
  - ┌ LIKE library\_info, BUT GENERATES A MAP FILE SUITABLE FOR DPRINTING
  - ┌ USEFUL OPTIONS
    - ┌ -header
    - ┌ -footer
  
- ┌ library\_cleanup, lcln
  - ┌ LISTS LIBRARY ENTRIES THAT HAVEN'T BEEN MODIFIED WITHIN THE SPECIFIED "GRACE" PERIOD
  - ┌ OPTIONALLY DELETES SUCH "OLD" SEGMENTS, LINKS, AND MULTISEGMENT FILES
  - ┌ ACCEPTS THE STAR CONVENTION
  
- ┌ library\_descriptor, lds
  - ┌ PRINTS INFORMATION ABOUT LIBRARY DESCRIPTORS, AND CONTROLS USE OF LIBRARY DESCRIPTORS BY THE OTHER LIBRARY DESCRIPTOR COMMANDS
    - ┌ RETURNS NAME OF CURRENT LIBRARY DESCRIPTOR BEING USED
    - ┌ CHANGES CURRENT LIBRARY DESCRIPTOR

PROGRAM LIBRARY MANAGEMENT TOOLS

LIBRARY DESCRIPTOR TOOLS

● lib\_descriptor\_ SUBROUTINE

▮ CONTAINS ENTRY POINTS THAT ARE CALLED BY ABOVE COMMANDS TO ACHIEVE THEIR GOALS

▮ REFERENCES THE LIBRARY DESCRIPTORS

● WHY ALL THIS INDIRECTION?

▮ AVOIDS REPLICATION OF COMMON CODE IN LIBRARY DESCRIPTOR COMMANDS

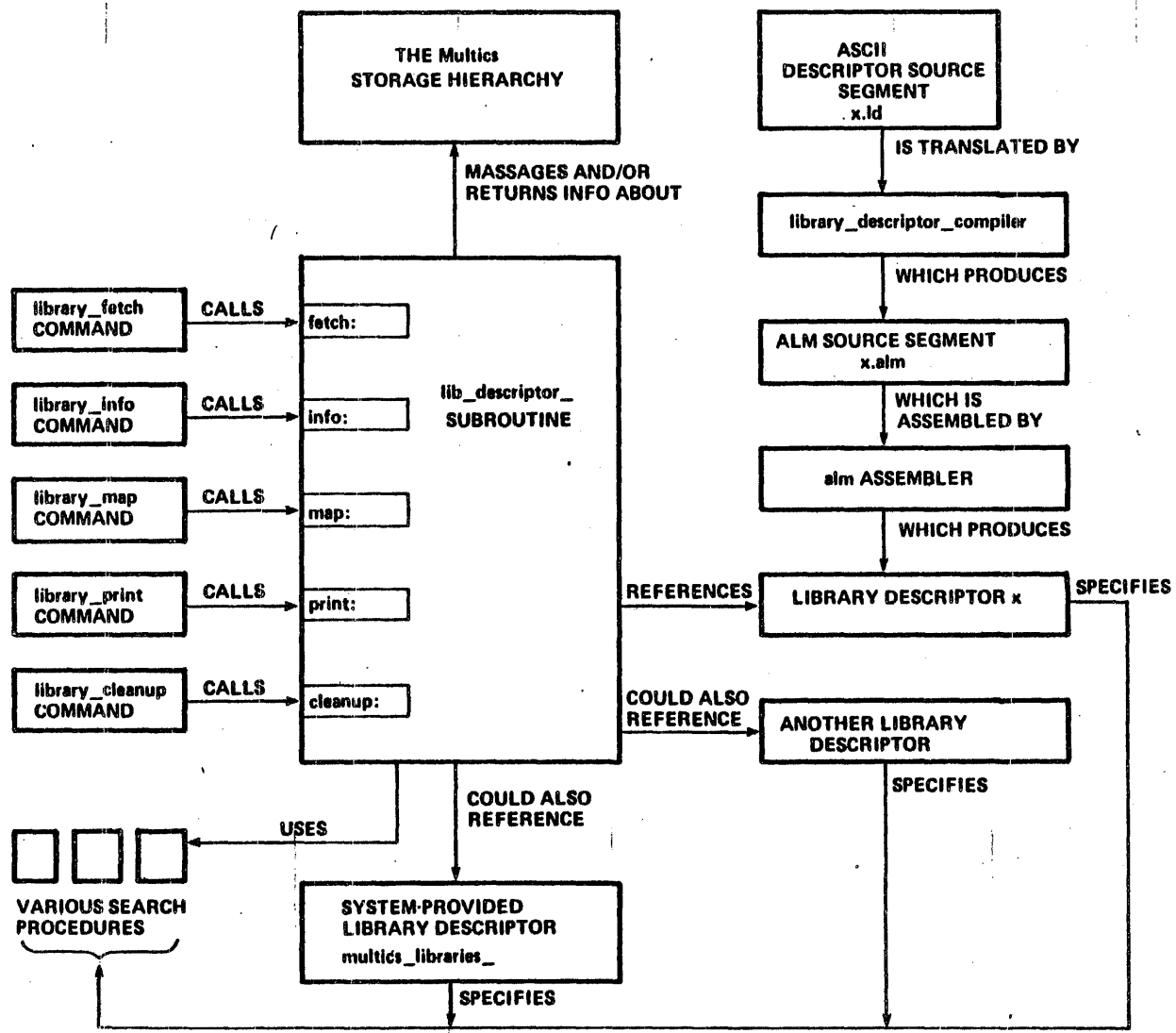
▮ AVOIDS MODIFICATION OF MANY, SEPARATE COMMANDS WHEN

▮ A NEW LIBRARY IS ADDED

▮ A NEW LIBRARY ORGANIZATION IS INSTITUTED

▮ PERMITS LIBRARY DESCRIPTOR COMMANDS TO WORK ON NON-SYSTEM LIBRARIES WITHOUT REWRITING THEM

# LIBRARY DESCRIPTOR SUMMARY





PROGRAM LIBRARY MANAGEMENT TOOLS

LIBRARY DESCRIPTOR TOOLS

!list -first 1

Segments = 3, Lengths = 80.

r w 1 handout\_desc.ld

!pr handout\_desc.ld 1

Descriptor: handout\_desc;

Define: commands;

command: library\_print;  
library name: handout;  
search names: \*\*;  
command: li;  
library name: handout;  
search names: \*\*;

Root: handout;

path: >udd>F15D>s1>handout;

search procedure: multics\_library\_search\_\$hardcore\_bc\_dir;

End: handout\_desc;

!ldc handout\_desc

!list -first 2

Segments = 4, Lengths = 82.

r 2 handout\_desc.alm

r w 1 handout\_desc.ld

!alm handout\_desc.alm

ALM

!list -first 3

Segments = 5, Lengths = 83.

re 1 handout\_desc

r 2 handout\_desc.alm

r w 1 handout\_desc.ld

PROGRAM LIBRARY MANAGEMENT TOOLS

LIBRARY DESCRIPTOR TOOLS

```
!li gw.archive
library_info: Use of star convention resulted in no match.
While searching for entries in the library.
Descriptor:      multics_libraries
library name:    online_libraries
search name:     gw.archive

!lds set handout_desc

!li gw.archive

1 gw.archive                                type: archive
      path: >udd>F15D>s1>handout
contents modified: 02/18/81 1648.4         system id: 34-32

!lpr gw.archive -components

!list -first 4

Segments = 6, Lengths = 119.

r w 36 library.print
re  1 handout_desc
r   2 handout_desc.alm
r w  1 handout_desc.ld
```

APPENDIX A

AIM

	Page
Concepts . . . . .	A-1
Commands and Subroutines . . . . .	A-3

CONCEPTS

- SOME BASIC TERMINOLOGY AND PROPERTIES SHOULD BE UNDERSTOOD:

⌋ AIM IS A NONDISCRETIONARY ACCESS CONTROL MECHANISM

⌋ SENSITIVITY (AS MANY AS 8)

⌋ CATEGORY SET (AS MANY AS 18)

⌋ ACCESS CLASS OF A SEGMENT IS COMPRISED OF A SENSITIVITY AND A CATEGORY SET

⌋ ACCESS AUTHORIZATION OF A PROCESS IS LIKEWISE COMPRISED OF A SENSITIVITY AND A CATEGORY SET

⌋ RELATIONSHIPS BETWEEN AUTHORIZATIONS AND ACCESS CLASSES

	SEGS	DIRS
AUTH $\geq$ ACCESS CLASS	re	s
AUTH = ACCESS CLASS	rew	sma
OTHERWISE	null	null

⌋ SYSTEMS "NOT RUNNING AIM" USE A SENSITIVITY OF "SYSTEM\_LOW" WITH NO CATEGORIES

CONCEPTS

- ▮ DETERMINING THE "PROCESS MAXIMUM AUTHORIZATION"
  - ▮ TAKE THE MINIMUM OF THE FOLLOWING 3:
    - ▮ PERSON MAXIMUM AUTHORIZATION ON ANY PROJECT
    - ▮ PERSON MAXIMUM AUTHORIZATION ON THE GIVEN PROJECT
    - ▮ PROJECT MAXIMUM AUTHORIZATION
  - ▮ A SEGMENT RECEIVES ITS ACCESS CLASS FROM THE CONTAINING DIR, NOT FROM THE ACCESS AUTHORIZATION OF THE CREATING PROCESS
  - ▮ A DIRECTORY'S ACCESS CLASS DEFAULTS TO ITS CONTAINING DIRECTORY, BUT CAN BE "UPGRADED" UP TO THE PROCESS MAXIMUM AUTHORIZATION

## COMMANDS AND SUBROUTINES

- THE FOLLOWING COMMANDS AND SUBROUTINES DEAL WITH THE ACCESS ISOLATION MECHANISM:

┆ print\_auth\_names (AG92)

┆ THIS COMMAND PRINTS THE SHORT AND LONG NAMES OF THE AIM SENSITIVITIES AND CATEGORIES SET FOR THIS SITE

┆ get\_authorization\_ (AG93)

┆ RETURNS THE AUTHORIZATION VALUE FOR THE CALLING PROCESS AS 'bit(72)'

┆ print\_proc\_auth (AG92)

┆ THIS COMMAND RETURNS CHAR-STRING REPRESENTATION OF THE PROCESS' AUTHORIZATION

┆ get\_max\_authorization\_ (AG93)

┆ RETURNS THE MAXIMUM AUTHORIZATION VALUE OF THE CALLING PROCESS AS THE 'bit(72)' VALUE

┆ convert\_authorization\_ (AG93)

┆ PROVIDES SEVERAL ENTRY POINTS FOR CONVERTING AN AUTHORIZATION BACK AND FORTH BETWEEN ITS BINARY AND ITS CHARACTER-STRING REPRESENTATION

## COMMANDS AND SUBROUTINES

▮ hcs\_\$get\_access\_class AG93)  
hcs\_\$get\_access\_class\_seg (AG93)

▮ THESE RETURN THE ACCESS CLASS OF A SEGMENT OR DIRECTORY GIVEN EITHER A DIRECTORY PATHNAME AND ENTRYNAME, OR GIVEN A POINTER TO THE SEGMENT ITSELF

▮ aim\_check\_ (AK92)

▮ PROVIDES SEVERAL ENTRY POINTS WHICH ALLOW THE CALLER TO DETERMINE THE AIM RELATIONSHIP ("EQUAL", "GREATER", "GREATER-OR-EQUAL") BETWEEN TWO ACCESS ATTRIBUTES (AUTHORIZATION OR ACCESS CLASS)

▮ read\_allowed\_ (AK92)  
write\_allowed\_ (AK92)  
read\_write\_allowed\_ (AK92)

▮ DETERMINE WHETHER THE SUBJECT OF A SPECIFIED AUTHORIZATION HAS ACCESS TO READ, WRITE, OR READ-AND-WRITE AN OBJECT OF SPECIFIED ACCESS CLASS

▮ get\_privileges\_ (AK92)

▮ THIS FUNCTION RETURNS THE ACCESS PRIVILEGES OF THE CALLING PROCESS (E.G., ipc ALLOWED, ETC.)

FOR A DISCUSSION OF AIM, SEE CHAPTER 6 OF THE MPM REFERENCE GUIDE.

APPENDIX B  
Program Listings

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user_real_init_admin . . . . .	B-3
process_overseer . . . . .	B-6
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listen . . . . .	B-15



user init admin

```
" *****
" *
" *
" * Copyright (c) 1972 by Massachusetts Institute of *
" * Technology and Honeywell Information Systems, Inc. *
" *
" *
" *****
```

```
name      user_init_admin_
entry     user_init_admin_
entry     daemon_init_admin_
entry     absentee_init_admin_
```

```
tempd    pit_ptr,po_ptr,arg(3)
include   stack_header
```

```
user_init_admin_:
  eppab   <user_real_init_admin_>|[user_real_init_admin_]
  tra    join
```

```
daemon_init_admin_:
  eppab   <daemon_real_init_admin_>|[daemon_real_init_admin_]
  tra    join
```

```
absentee_init_admin_:
  eppab   <absentee_real_init_admin_>|[absentee_real_init_admin_]
  tra    join
```

```
join:      push
           eppbp      po_ptr      prepare argument list
           spribp     arg+2
           eppbp      pit_ptr
           spribp     arg+4
           fld        =2b24,d1
           staq       arg
```

```
"
" call real_init_admin_ (po_ptr, pit_ptr)
```

```
" call      ab|0(arg)
```

```
" call the process overseer ... always pass it the pit pointer
"
```

```
use_arglist:
    fld      =1b24,d1          set up argument list (1 arg)
    staq     sp|0
    eppbp    pit_ptr,*        retrieve the pit pointer
    spribp   sp|4
    eppbp    sp|4
    spribp   sp|2
    eppbp    po_ptr,*         save pointer to process overseer
    eppap    sp|0             get pointer to argument list
    short_call      bp|0
    eppbp    process_killing_pointer
"    get a pointer to -2|0 which will blow us away
    arg      0
    even
process_killing_pointer:
    its      -2,0

    end
```



```
%include pitmsg;
%include static_handlers;
%include iox_modes;

    call hcs $make_seg ("", "pit", "", 01000b, pp, code);
        /* get pointer to PIT */
    type = pp -> pit.outer_module; /* Get DIM name */
    call ios $ios_quick_init; /* initialize syn attachments */

/* now do things that need doing before working dir exists */

    if ^pp -> pit.at.vinitproc then call find_po_and_dim;

    call change_wdir ((pp -> pit.homedir), code);
    /* put home_dir into search path */
    /* ignore code-- if no wdir we do the best we can */

    if pp -> pit.at.vinitproc then call find_po_and_dim;

/* Now set up static handlers for "alarm", "cput", and "term" */

    call sct_manager $set (cput_sct_index,
        timer_manager $cpu_time_interrupt, code);
    call sct_manager $set (alarm_sct_index,
        timer_manager $alarm_interrupt, code);
    call sct_manager $set (term_sct_index,
        term_signal_handler, (0));
    call sct_manager $set (wkp_sct_index,
        wkp_signal_handler, code);
    call sct_manager $set (susp_sct_index,
        sus_signal_handler, (0));

    return;

find_po_and_dim:
    procedure ();

        if type = "tty" then do;
            call iox $attach_iocb (iox $user_io,
                "tty_login_channel", code);
            if code = 0 then go to open;
        end;
        call iox $attach_iocb (iox $user_io,
            type || " " || pp -> pit.tty, code);
        /* attach primary input/output stream */
        if code = 0 then do;
open:
            call iox $open (iox $user_io,
                Stream_input_output, "0"b, code);
            if code ^= 0 then call login_abort ("io_attach",
                code);
        end;

end;
```

user real init admin

```
else do;
    string = pp -> pit.tty;
    call ios $attach ("user i/o", type,
        string, "", status);
    if based_code ^= 0 then call login_abort
        ("io_attach", based_code);
end;

pit_ptr = pp;

call find_command $fc_no_message
    (addr (pp -> pit.login_responder),
    length (pp -> pit.login_responder), po_ptr, code);
if code ^= 0 then do;
    call ioa ("Could not find specified initial
\nprocedure: ^a", pp -> pit.login_responder);
    call login_abort ("no_initproc", code);
end;
process_overseer_ptr = po_ptr;
return;
end /* find_po_and_dim */;

login_abort: proc (why, fatal_code);
    /* this procedure logs out the process with a special
    message indicating that an initialization error occurred */

dcl why char (*); /* reason we can't go */
dcl fatal_code fixed bin (35); /* code indicating fatal error */

dcl 1 term_structure aligned static,
    2 version fixed bin init (0), /* version of structure */
    2 status_code fixed bin (35); /* fatal error code */

    status_code = fatal_code;
    /* transmit code to terminate routine */
    call terminate_process (why, addr (term_structure));
    /* terminate the process */
end login_abort;

end user_real_init_admin;
```

process overseer

```
/* *****  
*  
* Copyright (c) 1972 by Massachusetts Institute of *  
* Technology and Honeywell Information Systems, Inc. *  
*  
*  
***** */
```

```
process_overseer_: proc (pit_ptr);
```

```
/* process_overseer_ is the standard process overseer on the system.  
It has four responsibilities:
```

setting up an unclaimed signal handler, otherwise known as an any other handler. This handler catches otherwise uncaught conditions. The supplied handler, default\_error\_handler\_\$wall prints any message provided for the error condition, establishes a condition wall, and calls the listener to get a new listener level. A condition wall is just another any other handler; this intercepts conditions that might otherwise be caught by other handlers present on the stack.

setting up a static handler for the mme2 condition. The mme2 condition is raised when the mme2 instruction is executed. It is used by debug to establish breakdots. The handler transfers control to debug when the condition is signalled.

allowing the "." escape to command query. This is enabled by calling command\_query\_\$set\_cp\_escape with the appropriate bits.

finding the start\_up.ec. It looks in the homedir, projectdir, and finally >sc1 to try to find it. It ends by calling listen\_ with "ec start\_upName" as the initial command line.

The code is written for time rather than space efficiency, so that operations that might look prettier in a do loop are done with inline code. \*/

process overseer

```
/* Automatic */

dcl initial_command_line char (104) var init ("");
dcl pit_ptr_ptr;
dcl code fixed bin (35);
dcl unaligned_homedir char (64) unaligned based
    (addr (pit_ptr_ptr -> pit.homedir));
dcl bc fixed bin (24);
dcl entry_type fixed bin (2);
dcl first_process bit (1);

/* Constants */

dcl process_type (0 : 3) character (12) varying internal
    static_options (constant) initial ("initializer", "interactive",
    "absentee", "daemon");
dcl down_sc1 char (4) internal static_options (constant) init (">sc1");
dcl start_up_dot_ec char (11) internal static_options (constant)
    init ("start_up.ec");
/* Entries */

dcl hcs $terminate_noname entry (ptr, fixed bin (35));
dcl condition_entry (char (*), entry);
dcl command_query $set_cp_escape_enable entry (bit (1) aligned,
    bit (1) aligned);
dcl listen_ext entry (char (*) var);
dcl default_error_handler $wall entry;
dcl hcs $status_minf entry (char (*), char (*), fixed bin (1),
    fixed bin (2), fixed bin (24), fixed bin (35));
dcl sct_manager $set entry (fixed bin, ptr, fixed bin (35));
dcl process_overseer $mme2_fault_handler_entry (ptr, char (*),
    ptr, ptr, bit (1));

/* External variables */

dcl iox_$user_output ptr ext;

/* Builtins */

dcl (codeptr, divide, null, rtrim) builtin;

#include pitmsg;
#include static_handlers;

/*set up the unclaimed signal handler */
    call condition_ ("any_other", default_error_handler_$wall);

/* turn on "." */
    call command_query_$set_cp_escape_enable ("1"b, ("b));

    first_process = (pit_ptr_ptr -> pit.n_processes = 1);
    /* see if new_proc or login */
```

```
if ^ pit_ptr -> pit.at.nostartup then do;
    /* start_up is allowed */
    initial_command_line = "exec_com ";

/* First try homedir */

    call hcs $status_minf (unaligned homedir,
        start_up_dot_ec, 1, entry_type, bc, code);

/* note that we assume any error is cause to look elsewhere to
    give best chance of success */

    if code = 0 & entry_type = 1 then
        initial_command_line = initial_command_line ||
            rtrim (pit_ptr -> pit.homedir);

/* now try projectdir */

    else do;
        call hcs $status_minf (">udd>" ||
            rtrim (pit_ptr -> pit.project),
            start_up_dot_ec, 1, entry_type, bc, code);

        if code = 0 & entry_type = 1 then
            initial_command_line = initial_command_line ||
                ">udd>" || rtrim (pit_ptr -> pit.project);
        else do;
            call hcs $status_minf (down_sc1,
                start_up_dot_ec, 1, entry_type, bc, code);
            if code = 0 & entry_type = 1 then
                initial_command_line =
                    initial_command_line || down_sc1;
            else do;
                initial_command_line = "";
                goto no_start_up;
            end;
        end;
    end;

    initial_command_line = initial_command_line || ">";
    initial_command_line = initial_command_line ||
        start_up_dot_ec;

/* the piecemeal assemble makes faster code */

    if first_process
        then initial_command_line =
            initial_command_line || " login ";
    else initial_command_line =
        initial_command_line || " new_proc ";

    initial_command_line = initial_command_line
        || process_type (pit_ptr -> pit.process_type);
end; /* the block that checked pit.nostart*/
```



process overseer

```
no_start_up:
    call hcs_terminate_noname (pit_ptr, code);

/* set up the mme2 handler */
/* this is done here rather than in xxx_real_init_admin so that
   process overseers for limited subsystems can leave it out */

    call sct_manager$set (mme2_sct_index,
        codeptr (process_overseer_$mme2_fault_handler_), code);

    do while ("1"b);
        call listen_ (initial_command_line);
    end;

    return;

mme2_fault_handler : entry (mcp, cname, cop, infop, cont);
dcl (mcp_ptr, /* to machine conditions */
    cname char (*), /* name of condition being signalled */
    cop_ptr,
    infop_ptr,
    cont bit (1)) parameter;

dcl debug$mme2_fault entry (ptr);

    call debug$mme2_fault (mcp);
    cont = "0"b; /* do not continue searching for handlers */
    return;
end process_overseer_;
```

project start up

```
/* *****  
*  
* Copyright (c) 1972 by Massachusetts Institute of *  
* Technology and Honeywell Information Systems, Inc. *  
*  
*  
***** */
```

```
project_start_up :  
  procedure (pit_ptr);
```

```
  dcl pit_ptr ptr;  
  dcl initial_command_line char (256) varying;  
  dcl listen_entry (char (*) var);  
  dcl terminate_process_entry (char (*), ptr);  
  dcl com_err_entry () options (variable);  
  dcl ioa_$ioa_switch entry options (variable);  
  dcl any_other_handler entry variable;  
  dcl any_other_condition;  
  dcl (length, null, unspec) builtin;
```

```
%include iox_dcls;
```

```
  any_other_handler = error_handler;  
  on any_other call any_other_handler;  
  /* Set up any_other handler outside the begin block */
```

```
  begin options (non_quick);
```

```
    dcl saved_cl_intermediary entry variable;  
    dcl home_dir char (168);  
    dcl project_dir char (168);  
    dcl mme2_handler ptr;  
    dcl saved_mme2_handler ptr;  
    dcl ps_ec_cl character (256) aligned;  
    dcl code fixed bin (35);  
    dcl bc fixed bin (24);  
    dcl entry_type fixed bin (2);  
    dcl first_process bit (1);  
    dcl (first_ec_arg, second_ec_arg) char (12);  
    dcl wall_entry entry variable;
```

```
    dcl process_type (0:3) character (12)  
      internal static options (constant) initial  
      ("initializer", "interactive", "absentee", "daemon");  
    dcl down_sc1 char (19) internal static  
      options (constant) init (">system control_dir");  
    dcl start_up_dot_ec char (11) internal static  
      options (constant) init ("start_up.ec");
```

project start up

```
dcl cu $cp entry (ptr, fixed bin (21), fixed bin (35));
dcl convert_status_code_entry (fixed bin (35),
char (8) aligned, char (100) aligned);
dcl change_wdir_entry (char (168), fixed bin (35));
dcl hcs $make_entry entry (ptr, char (*), char (*),
entry, fixed bin (35));
dcl hcs $terminate_noname entry (ptr, fixed bin (35));
dcl command_query_$set_cp_escape_enable entry
(bit (1) aligned, bit (1) aligned);
dcl default_error_handler $wall entry;
dcl hcs $status_minf entry (char (*), char (*),
fixed bin (1), fixed bin (2), fixed bin (24),
fixed bin (35));
dcl sct_manager $set entry (fixed bin, ptr,
fixed bin (35));
dcl sct_manager $get entry (fixed bin, ptr,
fixed bin (35));
dcl process_overseer $mme2_fault_handler
entry (ptr, char (*), ptr, ptr, bit (1));
dcl cu $set_cl_intermediary entry (entry);
dcl cu $get_cl_intermediary entry (entry);

dcl command_error condition;

dcl (addr, codeptr, length, null, rtrim, unspec)
builtin;

#include pitmsg;
#include static_handlers;

home_dir = pit_ptr -> pit.homedir;
project_dir = ">user_dir_dir>" ||
rtrim (pit_ptr -> pit.project);

call sct_manager $get (mme2_sct_index,
saved_mme2_handler, (0));

call hcs $status_minf (project_dir,
"project_start_up.ec", 1, entry_type, bc, code);
if ^(entry_type = 1 & code = 0)
then call abort_handler (rtrim (project_dir) ||
">project_start_up.ec was not found or is not a segment.",
code);

call change_wdir_ (project_dir, code);
if code ^= 0
then call abort_handler
("Could not set working directory to project directory.", code);
```

project start up

```
first_process = (pit_ptr -> pit.n_processes = 1);
if first_process
then first_ec_arg = "login";
else first_ec_arg = "new_proc";

second_ec_arg = process_type (pit_ptr -> pit.process_type);

call hcs $make_entry (null (), "default_error_handler_",
    "wall", wall_entry, code);
if code ^= 0
then wall_entry = default_error_handler_$wall;

any_other_handler = wall_entry;

call cu $get_cl_intermediary (saved_cl_intermediary);
call cu $set_cl_intermediary (error_handler);

on command_error call com_err_handler;
/* die on com_err */

ps_ec_cl = "exec com " || rtrim (project_dir) ||
    ">project_start_up " || rtrim (first_ec_arg) ||
    " " || rtrim (second_ec_arg);
call cu $cp (addr (ps_ec_cl),
    length (rtrim (ps_ec_cl)), (0));

revert command_error;

call cu $set_cl_intermediary (saved_cl_intermediary);

call change_wdir_ (home_dir, code);
if code ^= 0
then call com_err_ (code, "project_start_up_",
"Could not set working directory to ^a.", home_dir);

call command_query_$set_cp_escape_enable ("1"b, ("b));

if ^pit_ptr -> pit.at.nostartup
then do;
    initial_command_line = "exec com ";
    call hcs $status_minf (home_dir,
        start_up_dot_ec, 1, entry_type, bc, code);
    if code = 0 & entry_type = 1
    then initial_command_line =
        initial_command_line || rtrim (home_dir);
```

project start up

```
else do;
    call hcs $status_minf (project_dir,
        start_up_dot_ec, 1, entry_type, bc, code);

    if code = 0 & entry_type = 1
        then initial_command_line =
            initial_command_line || project_dir;
    else do;
        call hcs $status_minf (down_sc1,
            start_up_dot_ec, 1, entry_type,
            bc, code);
        if code = 0 & entry_type = 1
            then initial_command_line =
                initial_command_line || down_sc1;
        else do;
            initial_command_line = "";
            goto no_start_up;
        end;
    end;
end;

initial_command_line = initial_command_line || ">";
initial_command_line = initial_command_line ||
    start_up_dot_ec;
initial_command_line = initial_command_line || " " ||
    first_ec_arg;
initial_command_line = initial_command_line || " " ||
    second_ec_arg;
end;

no_start_up:
    call hcs $terminate_noname (pit_ptr, code);
    call sct_manager $get (mme2_sct_index, mme2_handler, (0));
    if mme2_handler = saved_mme2_handler
        then call sct_manager $set (mme2_sct_index, codeptr
            (process_overseer_$mme2_fault_handler_), code);
    end;
call listen_ (initial_command_line);
do while ("1"b);
    call listen_ ("");
end;
return;

com_err_handler:
    procedure;
%include condition_info_header;
%include condition_info;
%include com_af_error_info;
declare 1 CI aligned like condition_info;
declare find_condition_info_entry (pointer, pointer,
    fixed binary (35));
declare code fixed bin (35);
```

project start up

```
unspec (CI) = ""b;

call find_condition_info_ (null (), addr (CI), code);
if code ^= 0
then call abort_handler ("Can't get error message.", code);
call ioa $ioa_switch (iox $error_output, "^a",
    CI.info_ptr -> com_af_error_info.info_string);
call abort_handler ("Error in project start up.", 0);
end;
error_handler:
entry;

    call abort_handler ("Error in project start up.", 0);

abort_handler:
proc (reason, code);

    dcl code fixed bin (35);
    dcl reason char (*);
    dcl 1 term structure aligned,
        2 version fixed bin init (0),
        2 status_code fixed bin (35);

    status_code = code;
    call ioa $ioa_switch (iox $error_output, reason);
    call terminate_process_ ("init_error", addr (term_structure));
end;

end project_start_up_;
```

listen

```
/* *****  
*  
* Copyright (c) 1972 by Massachusetts Institute of *  
* Technology and Honeywell Information Systems, Inc. *  
*  
*  
***** */  
  
listen_ : procedure (initial_command_line);  
  
/* Multics Listener */  
dcl iox_$user_input ptr ext static;  
dcl iox_$user_io ptr ext static;  
dcl (buffer_ptr,  
    read_ptr ptr,  
    dummy_ptr ptr,  
    pct internal static initial (null),  
    old_sp  
    ) pointer aligned;  
  
dcl (input_length, buffer_length) fixed bin (21);  
dcl total_input_length fixed bin (21);  
dcl entry, /* 0->$listen_, 1->$release_stack */  
    i fixed bin aligned;  
dcl code fixed bin (35) aligned;  
  
dcl initial_command_line char (*) var,  
    /* first command line to be executed */  
    command_line char (input_length) aligned based (buffer_ptr);  
  
dcl spno bit (18) aligned, /* used to store stack segno */  
    (should_restore_attachments,  
    (first, /* "1"b means control structure not initialized */  
    quits_not_enabled) int static init ("1"b)  
    ) bit(1) aligned;  
  
dcl 1 x based (buffer_ptr) aligned,  
    2 ch (0:65536) char (1) unaligned;  
  
dcl 1 label_var aligned based, /* overlay for a label */  
    2 target_ptr, /* target of entry/label variable */  
    2 stackp_ptr; /* stack offset of entry/label variable */  
  
dcl 1 ct aligned, /* automatic structure containing control info */  
    2 prev_ptr ptr, /* ptr to last listener stack frame (if any) */  
    2 release_all_label, /* label in "top" level to release to */  
    2 release_label, /* label to release to */  
    2 new_release_label, /* label next invocation is to release to */  
    2 start_label, /* label for start command */  
    2 flags aligned,  
    3 dont_restore bit (1) unal, /* "1"b causes io attachments  
        not to be restored on start */
```

listen

3 pad bit (35) unal,  
2 frame fixed bin, /\* stack frame of current invocation \*/  
2 level fixed bin; /\* level of current invocation (from 1) \*/

dcl 1 bct aligned based (pct) like ct;

dcl ios\_signal\_entry (char (32) aligned, fixed bin (35)),  
iox\_\$get\_line entry (ptr, ptr, fixed bin (21), fixed bin (21))  
returns (fixed bin (35)),  
iox\_\$control entry (ptr, char (\*), ptr) returns (fixed bin (35)),  
com\_err\_entry options (variable),  
cu\_\$cp\_ext entry (ptr, fixed bin (21), fixed bin (35)),  
cu\_\$ready\_proc\_ext entry (),  
cu\_\$grow\_stack\_frame entry (fixed bin (21), ptr, fixed bin (35)),  
get\_system\_free\_area\_entry returns (ptr),  
cu\_\$stack\_frame\_ptr\_ext entry () returns (ptr);

dcl (addr,  
baseno,  
divide,  
fixed,  
length,  
min,  
null,  
ptr,  
rel  
) builtin;

dcl cleanup condition;

dcl error\_table \$long\_record ext static fixed bin (35);

%include stack\_frame;



listen

```
/* Establish this frame as the "top" of the listener frame thread,
   so that this frame cannot be "released" around. */
   entry = 0;
   go to re_enter;

/* Entry called after processing quit or unclaimed signal */
release_stack: entry (should_restore_attachments);

   entry = 1;

/* Save pointer to previous listener control information, save return
\c
point for subsequent invocations of the listener,
and initialize switches */
re_enter:
   if first then do; /* no previous invocation to work from */
       ct.prev_ptr = null;
       ct.level = 1; /* this is first invocation */
       sp = cu $stack frame_ptr (); /* find stack frame */
       spno = baseno (sp); /* get segno for comparing */
       i = 0;
       do while (baseno (sp -> stack_frame.prev_sp) = spno);
           i = i + 1;
           sp = sp -> stack_frame.prev_sp;
       end;
       ct.frame = i;
   end;
   else do; /* can use info from previous invocation */
       ct.prev_ptr = pct;
       ct.level = bct.level + 1;
       old_sp = addr (bct.start) -> label_var.stackp;
       /* find previous frame */
       sp = cu $stack_frame_ptr ();
       i = bct.frame;
       do while (sp ^= old_sp);
           /* find number of intervening frames */
           i = i + 1;
           sp = sp -> stack_frame.prev_sp;
       end;
       ct.frame = i;
   end;

/* fill in labels for release and start */
   if (entry = 0) | first then do;
       ct.release_all,
       ct.release,
       ct.new_release = readyt;
       first = "0"b;
   end;
   else do;
       /* will want to release to invocation before this one */
       ct.release_all = bct.release_all; /* dont change it */
       ct.release = bct.new_release;
```

listen

```
        ct.new_release = readyt;
    end;
    ct.start = start_return_point;

    pct = addr (ct);
    /* have finished getting info from old frame */

    ct.flags.dont_restore = "0"b;

/* set ptrs to current control info and to buffer
   in which to read in command line */
    buffer_length = 32; /* start with 128 char input buffer */
    call cū $grow_stack frame (buffer_length,buffer_ptr, code);
    /* get storage for initial buffer */

/* Establish cleanup procedure to restore control structure thread */
    on condition (cleanup) begin;
        pct = bct.prev_ptr; /* pop structure of interest */
        if pct = null then first = "1"b;
    end;

/* Check for entering via "release_stack" entry without having first
   entered via "listen". This can happen, e.g., if user takes fault
   before standard process_overseer_ calls listen_. If this happens,
   enable quits. */
    if quits_not_enabled then do;
        quits_not_enabled = "0"b;
        code = iox_$control (iox_$user_io,"quit_enable",null);
    end;
/* If called at the listen_ entry, set up initial command
   line and enable quits */
    if entry = 0 then do;
        if initial_command_line ^= "" then do;
            if length (initial_command_line) >
                buffer_length * 4 then do;
                call com_err (0, "listen_",
                    "Initial command line is too long.!!
                    "Max=^d chars.",buffer_length*4);
                go to readyt;
            end;
            input_length = length (initial_command_line);
            command_line = initial_command_line;
            total_input_length = 0;
            go to CALL_CP;
        end;
    end;
end;
```

listen

```
/* *****START OF BASIC LISTENER LOOP***** */
```

```
/* Call the "ready procedure". */
```

```
readyt: call cu_$ready_proc ();
```

```
/* Read the next command line */
```

```
readnew: read_ptr = buffer_ptr;
```

```
total_input_length = 0;
```

```
/* extra input line character count */
```

```
read:
```

```
code = iox_$get_line (iox_$user_input, read_ptr,  
buffer_length*4-total_input_length,input_length);
```

```
if code ^= 0 then do;
```

```
if code ^= error_table$long_record then
```

```
call ios_signal("user_input", code);
```

```
else do;
```

```
if input_length <
```

```
buffer_length * 4 - total_input_length
```

```
then goto CALL_CP;
```

```
call cu_$grow_stack_frame (buffer_length,
```

```
dummy_ptr, code); /* double size of buffer */
```

```
buffer_length = buffer_length + buffer_length;
```

```
read_ptr = addr (read_ptr -> ch (input_length));
```

```
total_input_length = total_input_length +
```

```
input_length;
```

```
end;
```

```
goto read;
```

```
end;
```

```
CALL_CP: call cu_$cp (buffer_ptr, total_input_length + input_length,  
code);
```

```
if code = 100 then go to readnew;
```

```
/* ignore null command line */
```

```
go to readyt;
```

```
/* *****END OF BASIC LISTENER LOOP***** */
```

## listen

```
start_return_point: /* start command goes here */
    if ct.flags.dont_restore then
        should_restore_attachments = "0"b;
        pct = bct.prev_ptr;
        return;

get_ptr: entry (ct_ptr);
dcl ct_ptr ptr;

/* Return pointer to control structure */
    ct_ptr = pct;
    return;

get_level: entry (level_no, frame_no);
/* Return command level number and stack frame number of caller's
   caller */

dcl (level_no, frame_no) fixed bin;

    if pct = null then do; /* no previous invocation */
        level_no = 0;
        old_sp = ptr (addr (old_sp), 0) ->
            stack_header.stack_begin_ptr;
    /* in case we're not in highest ring */
        frame_no = 0;
    end;
    else do; /* count only up to previous listener */
        level_no = pct -> bct.level;
        old_sp = addr (bct.start) -> label_var.stackp;
        frame_no = bct.frame;
    end;

    sp = cu $stack_frame_ptr () -> stack_frame.prev_sp ->
        stack_frame.prev_sp;
    /* want frame no of caller's caller */
    do while (sp ^= old_sp);
        frame_no = frame_no + 1;
        sp = sp -> stack_frame.prev_sp;
    end;
    return;

get_area: entry returns (ptr);
    return (get_system_free_area ());

#include stack_header;

    end;
```

**APPENDIX C**

**Encoding of Channel Names**

**Page**

The name used to designate an MCS communications channel is a character string of up to 32 characters. The name is composed of components separated by periods, where each component represents a level of multiplexing. The first two components identify the physical channel on an FNP; further components (if present) identify the subchannels of a concentrator (such as a VIP 7700 controller).

Format of physical channel name: The physical channel name (which corresponds to the old-style name of the form ttyXXX) has the following format:

F.ANSS

where:

F is an FNP identifier (a, b, c, or d)  
A is an adapter type (h for an HSLA channel, l for an LSLA channel)  
N is the number of the particular adapter (0-2 for an HSLA, 0-5 for an LSLA)  
SS is the decimal number of the subchannel on the specified adapter.

Examples:

Name	Description	Old form
a.l000	FNP a, LSLA 0, subchannel 0	tty000
a.h108	FNP a, HSLA 1, subchannel 8	tty708
b.h016	FNP b, HSLA 0, subchannel 16	ttyG16

Multiplexed channels: The format of the additional components of the names of subchannels of a concentrator or "multiplexer" depends on the particular multiplexer; it may be a station id, or a sequential number, etc. For example:

Name	Description
b.h016.01	FNP b, HSLA 0, subchannel 16, concentrator subchannel 1
b.h016.09	same physical channel, concentrator subchannel 9

ARPANET channels: The names of ARPANET channels are of the form netXXX for user telnet channels or ftpXXX for file-transfer channels, where XXX is an arbitrary 3-digit number.

APPENDIX D

Instructor Code for IPC Workshop

Page

The following segments set up the environment such that students may complete the interprocess communication workshop. Note in ipc\_driver.pl1 the call to get\_userid. For this call to successfully return, it is required that your instructor obtain read access to >sc1>answer\_table.

```
ipc_report: proc;

dcl 1 fixed bin,
    get_wdir_entry returns (char (168)),
    stud_ptr_ptr,
    ioa_entry options (variable),
    hcs_$initiate entry (char (*), char (*), char (*),
    fixed bin (1), fixed bin (2), ptr, fixed bin (35)),
    code fixed bin (35),
    clock_entry returns (fixed bin (71)),
    date_time_entry (fixed bin (71), char (*)),
    my_time char (24);

dcl 1 stud_ipc based (stud_ptr) aligned,
    2 index fixed bin,
    2 studs (0 refer (index)),
    (3 codes char (8),
    3 name char (22),
    3 proj char (9),
    3 time char (16)) unal;

    call date_time (clock (), my_time);
    call hcs_$initiate (get_wdir_ ( ), "ipc_status", "",
    0, 1, stud_ptr, code);

    call ioa_ ("REPORT FOR F15D WORKSHOP #3 ^a", my_time);
    call ioa_ ("^3/      user_id      time^/");
    do i = 1 to index;
        if name (i) ^= "" then
            call ioa_ ("^a.^a      ^a",
            name (i), proj (i), time (i));
    end;

end ipc_report;
```



```

ipc_driver: proc;

dcl  send_mail_entry (char (*), char (*), ptr, fixed bin (35)),
    get_wdir_entry returns (char (168)),
    hcs_$make_seg entry (char (*), char (*), char (*), fixed bin (5),
    ptr, fixed bin (35)),
    hcs_$wakeup entry (bit (36), fixed bin (71), fixed bin (71),
    fixed bin (35)),
    ipc_$create_ev_chn entry (fixed bin (71), fixed bin (35)),
    ipc_$decl_ev_call_chn entry (fixed bin (71), entry, ptr, fixed bin,
    fixed bin (35)),
    get_process_id entry returns (bit (36)),
    (ioa_, com_err_) entry options (variable),
    get_userid_entry (bit (36), char (*), char (*), fixed bin, fixed bin,
    fixed bin (35)),
    iox_$control entry (ptr, char (*), ptr, fixed bin (35)),
    date_time_entry (fixed bin (71), char (*)),
    clock_entry returns (fixed bin (71)),
    unique_bits_entry returns (bit (70)),
    unique_chars_entry (bit (*)) returns (char (15));

dcl 1 send_mail_info aligned,
    2 version fixed bin init (1),
    2 sent_from char (32) aligned init ("Mr. Wonderful"),
    2 switches,
    (3 wakeup bit (1) init ("1"b),
    3 mbz1 bit (1),
    3 always_add bit (1) init ("1"b),
    3 never_add bit (1) init ("0"b),
    3 mbz2 bit (1),
    3 acknowledge bit (1) init ("0"b),
    3 mbz bit (30)) unal;

dcl  congrats char (40) internal static options (constant)
    init ("Congratulations - mission accomplished!!");
dcl  ipc_status_full_msg char (66) internal static options (constant) init
    ("Instructor's table has overflowed. Please notify him immediately.");
dcl  destination char (32);
dcl  code fixed bin (35);

dcl 1 set_up based (su_ptr),
    2 my_pid bit (36),
    2 my_chid fixed bin (71);

dcl  me char (10) init ("ipc_driver") static options (constant);

dcl 1 event_info based (ei_ptr),
    2 channel_id fixed bin (71),
    2 message fixed bin (71),
    2 sender bit (36),
    2 origin,
    3 dev_signal bit (18) unal,
    3 ring bit (18) unal,
    2 data_ptr ptr;

```

```

dcl (su_ptr, ei_ptr, sptr) ptr;
dcl io_x$user_io ext ptr;
dcl error_table$invalid_channel ext fixed bin (35);
dcl string0 static fixed bin (71);
dcl string1 char (8) based (sptr),
    string2 char (8) based (mptr);
dcl mptr ptr;
dcl person char (22),
    project char (9),
    (type, anon) fixed init (0) bin,
    stud_ptr static ptr;

dcl 1 stud_ipc based (stud_ptr) aligned,
    2 index fixed bin,
    2 studs (0 refer (index)),
    (3 codes char (8),
     3 name char (22),
     3 proj char (9),
     3 time char (16)) unal;

dcl i fixed bin;

/* Set it up */
    call hcs $make_seg (get_wdir_ (), "channel_info", "", 10,
        su_ptr, code);
    call ipc $create_ev_chn (my_chid, code);
    if code ^= 0 then do;
        call com_err_ (code, "ipc_driver");
        return;
    end;
    my_pid = get_process_id ();
    call ipc $decl_ev_call_chn (my_chid, wakeme, null (), 0, code);
    if code ^= 0 then do;
        call com_err_ (code, "ipc_driver");
        return;
    end;
    call ioa ("End ^a$^a", me, me);
/* *****T***** */

dcl timer_manager_$sleep entry (fixed bin (71), bit (2));
    do while ("1"b);
        call timer_manager_$sleep (900, "11"b);      /* 15 min. */
    end;
    return;

/* Come here when wakeup received */
wakeme:    entry (ei_ptr);

    call hcs $make_seg (get_wdir_ (), "ipc_status", "", 10,
        stud_ptr, code);

```

```

if stud_ptr = null () then do;
    call com_err_ (code, "ipc_driver");
    return;
end;
call get_userid_ (sender, person, project, type, anon, code);
if code ^= 0 then do;
    call com_err_ (code, "ipc_driver",
        "Need 'r' access on >sc1>answer_table");
    return;
end;

call ioa_ ("A wakeup from ^a.^a was just received.", person,
    project); /* But that does not imply he'll wake me
                up again with the proper reversed msg */

mptr = addr (message);
do i = 1 to 250 while (codes (i) ^= "");
    if string2 = codes (i) then goto got_one;
end;

if i = 251 then do;
    call com_err_ (0, "ipc_driver",
        "Table in the segment 'ipc_status' is full.
System will not function properly.");
    destination = rtrim (person) || "." || rtrim (project);
    call send_mail_ (rtrim (destination), ipc_status_full_msg,
        addr (send_mail_info), code);
    if code ^= 0 then do;
        call com_err_ (code, "ipc_driver", "Bad call to send_ma
occurred while trying to complain about full table in 'ipc_status'.");
        return;
    end;
    return;
end;

sptr = addr (string0); /* Overlay string1 onto string0 */
string1 = substr (unique_chars_ (unique_bits_ ()), 8, 15);
index = index + 1 /* index reflects the true size of
                    stud_ipc.studs array */
codes (index) = reverse (string1);
call hcs $wakeup (sender, message, string0, code);
if code ^= 0 then do;
    call com_err_ (code, "ipc_driver");
    return;
end;
goto finis;

got_one:
name (i) = person;
proj (i) = project;
call date_time_ (clock (), time (i));
call ioa_ ("^a.^a completes assignment", person, project);
destination = rtrim (person) || "." || rtrim (project);

```

```
call send_mail_ (rtrim (destination), congrats,  
                _addr_ (send_mail_info), code);  
if code ^= 0 then do;  
    call com_err_ (code, "ipc_driver", "Bad call to send_mail_");  
    return;  
end;
```

finis:

```
end ipc_driver;
```

```

init_ipc: proc;

dcl get_wdir_entry returns (char (168));
dcl 1 stud_ipc based (stud_ptr) aligned,
  2 index fixed bin,
  2 studs (250),
  (3 codes char (8),
  3 name char (22),
  3 proj char (9),
  3 time char (16)) unal;
dcl hcs_$make_seg entry (char (*), char (*), char (*),
  fixed bin (5), ptr, fixed bin (35));
dcl hcs_$add_acl_entries entry (char (*), char (*), ptr,
  fixed bin, fixed bin (35));
dcl 1 seg_acl aligned,
  2 access name char (32) init ("*.F15d.*"),
  2 modes bit (36) init ("1"b),
  2 zero_pad bit (36) init ("0"b),
  2 status code fixed bin (35);
dcl code fixed bin (35),
  ioa_entry options (variable),
  stud_ptr ptr,
  i fixed bin;

  call ioa_ ("Begin init_ipc");
  call hcs_$make_seg (get_wdir_ (), "ipc_status", "",
    10, stud_ptr, code);
  call hcs_$add_acl_entries (get_wdir_ (), "ipc_status",
    addr_(seg_acl), 1, code);

  index = 0;
  do i = 1 to 250;
    time (i), name (i), proj (i), codes (i) = "";
  end;
  call ioa_ ("End init_ipc");
end init_ipc;

```

ABSENTEE SCRIPT ipc.absin

```

&ready off
cwd >udd>F15d>s1
ipc_driver
& To prevent the absout segment from growing inordinately large,
& the instructor has inserted the following command, which will guarantee
& an absout segment of 10k or less.
if [greater [st ipc.absout -bc] 1474560] -then "tc ipc.absout"
logout

```

**APPENDIX E**  
**Standard Process Overseers**

**Page**

On April 23, 1981, the following command was typed:

```
li *overseer_* -library source
```

The terminal output that resulted appears below.

```
1 accounts_overseer_.pl1                                type: arch comp
  path: >ldd>tools>source>bound_admin_rtnes_.s.archive
  component updated: 01/29/75 1711.6

1 cards_overseer_.pl1                                  type: arch comp
  path: >ldd>tools>source>bound_card_input_.s.archive
  component updated: 09/04/79 1718.4

1 dfast_process_overseer_.pl1                          type: arch comp
  path: >ldd>unb>source>bound_dfast_.s.archive
  component updated: 09/01/76 1342.7

1 fst_process_overseer_.pl1                            type: arch comp
  path: >ldd>unb>source>bound_fast_.s.archive
  component updated: 06/07/77 1504.7

1 ftp_server_overseer_.pl1                             type: arch comp
  path: >ldd>net>source>bound_ftp_server_.s.archive
  component updated: 09/23/77 1031.5

1 iod_overseer_.pl1                                    type: arch comp
  path: >ldd>tools>source>bound_iodec_.s.archive
  component updated: 03/13/81 1038.7

1 terminals_overseer_.pl1                              type: arch comp
  path: >ldd>tools>source>bound_admin_rtnes_.s.archive
  component updated: 03/23/81 1014.8

1 tolts_overseer_.pl1                                  type: segment
  path: >ldd>tools>source
  contents modified: 12/01/80 1136.8
```

Another important overseer: project\_start\_up\_

**APPENDIX F**

**Gate and Message Segment Examples**

**Page**



**APPENDIX G**

**Advanced Dial Facility Example**

**Page**

```

set_up_dial: proc;

/* The set_up_dial entry point initializes the dialing environment

1) An event-call channel is established so that
the answering service can notify this process
of dialins, hangups, etc.

2) Dials are enabled
*/
/* 'dialok' attribute essential if this procedure is to succeed */

dcl ipc_$create_ev_chn entry (fixed bin (71), fixed bin (35)),
ipc_$delete_ev_chn entry (fixed bin (71), fixed bin (35)),
ipc_$decl_ev_call_chn entry (fixed bin (71), entry,
ptr, fixed bin, fixed bin (35)),
hcs_$assign_channel entry (fixed bin (71), fixed bin (35)),
dial_manager_$allow_dials entry (ptr, fixed bin (35)),
(ioa_, com_err_, ioa_$ioa_switch) entry options (variable);

dcl wasted_channel fixed bin (71);
dcl time char (24);
dcl code fixed bin (35),
iox_$user_output external static ptr,
ME char (12) varying init ("set_up_dial") ;

dcl 1 dial_manager_arg aligned static,
2 version fixed bin init (1),
2 dial_qualifier char (22) init ("astra"),
2 dial_channel fixed bin (71),
2 channel_name char (32) ;

/* ***** */

call ioa_ ("Begin ^a", ME);

/* The following code is inserted to fake out tty_, which will
attempt to give us fast channels we can not deal with (see the
'read_status' control order). By consuming all the fast channels now
(and there aren't many available to us), we'll force tty_ to use
garden_variety channels, which we can easily handle. */

code = 0;
do while (code = 0);
call hcs_$assign_channel (wasted_channel, code);
end;

```

```

/* Channel must be obtained for notifying this process of */
/* all hangups and dialups. */
    call ipc $create_ev_chn (dial_manager_arg.dial_channel,
        code);
    if code ^= 0 then call ERROR (1);

/* Next let the answering service know we will accept
dials - we must do this before changing the event-wait
channel to an event-call channel */
    call dial_manager_$allow_dials (addr (dial_manager_arg),
        code);
    if code ^= 0 then call ERROR (2);

/* Make the event-wait channel an event-call channel and
specify that my 'dial_handler' will be invoked whenever
the answering service wakes my process on this channel */
    call ipc $decl_ev_call_chn (dial_manager_arg.dial_channel,
        dial_handler, null (), 0, code);

    if code ^= 0 then call ERROR (3);

/* Okay...now return and wait for something to happen */
    call ioa_ ("Now listening for dials: ^a", ME);
    return;

/* ***** */
dial_handler: entry (info_ptr);

/* Handler for dial messages - this entry point will
be invoked whenever something happens that the answering
service notifies me about */

dcl info_ptr ptr parameter;
dcl 1 event_info based (info_ptr),
    2 channel_id fixed bin (71),
    2 message_fixed bin (71),
    2 sender bit (36),
    2 origin,
    3 dev_signal bit (18) unal,
    3 ring_bit (18) unal,
    2 data_ptr ptr;

dcl listen_to_dial entry(ptr); /* proc to dialog with terminals */

```

```

dcl convert_dial_message $return_io module entry
  (fixed bin (71), char (*), char (*), fixed bin,
   1 aligned like status_flags, fixed bin (35));

dcl 1 status_flags aligned,
  (2 dialed_up bit (1),
   2 hung_up bit (1),
   2 control bit (1),
   2 pad bit (33)) unal;

dcl 1 dialed static,          /* This structure works for
                             max_num_allowed <= 10 */
  2 no_dialed fixed bin init (0),
  2 sw (10),
  3 sname char (6) init (
    "dial01", "dial02",
    "dial03", "dial04",
    "dial05", "dial06",
    "dial07", "dial08",
    "dial09", "dial10"),
  3 iocb_ptr ptr init ((10)null ()), /* An available switch is
                                     characterized by null ptr */
  3 devname char (32);

/* We will only talk to 2 terminals at a time, hanging up the third */
dcl max_num_allowed internal static options (constant) init (2);

dcl nomore_ptr internal static ptr init (null ()); /* This iocbptr is
  used for talking to a doomed tty when system full */

dcl iox_$find_iocb entry (char (*), ptr, fixed bin (35)),
  iox_$attach_ptr entry (ptr, char (*), ptr, fixed bin (35)),
  iox_$destroy_iocb entry (ptr, fixed bin (35)),
  iox_$control_entry (ptr, char (*), ptr, fixed bin (35)),
  error_table $io_no_permission external static fixed bin (35),
  iox_$open_entry (ptr, fixed bin, bit (1) aligned, fixed bin (35)),
  iox_$detach_iocb entry (ptr, fixed bin (35)),
  iox_$close_entry (ptr, fixed bin (35)),
  clock_entry returns (fixed bin (71)),
  date_time_entry (fixed bin (71), char (*)),
  ipc_$cutoff entry (fixed bin (71), fixed bin (35)),
  ipc_$reconnect entry (fixed bin (71), fixed bin (35));

dcl i fixed bin; /* an index */

dcl (which_channel automatic, nomore_channel static) char (32),
  io module char (32),
  n_dialed fixed bin;

```

```

ME = "dial handler";          /* For com_err */
/* First of all, interpret the event message sent from
   the answering service - it should either be that
   someone has dialed in or hung up */
call convert_dial_message $return_io_module (
    event_info.message, which_channel,
    io_module, n_dialed, status_flags, code);
/* n_dialed = -1 If this is an informative
   message (which it is)*/
if code ^= 0 &
    code ^= error_table_$io_no_permission then call ERROR (5);

/* Log in event */

call date_time (clock (), time);
call ioa $ioa_switch (iox $user_output,
    "a: [DIALED UP^;HUNG UP^] AT ^a.",
    which_channel, status_flags.dialed_up, time);

/* Restart any interrupted io to the master terminal */
call iox $control (iox $user_output, "start", null (), code);

if ^status_flags.dialed_up then do; /* Then it must
    be a hang up and we have to find out which
    device, mark it available, & detach the switch. */

    do i = 1 to max_num_allowed
        while (dialed.sw (i).devname ^=
            which_channel);
    end;

    if ^(i > max_num_allowed) then do;
        dialed.no_dialed = dialed.no_dialed - 1;
        call ioa $ioa_switch (iox $user_output,
            "9xAt this instant, ^i [terminal^;terminals^] [is^;are^] logged on.",
            dialed.no_dialed, (dialed.no_dialed = 1),
            (dialed.no_dialed = 1));

        /* Close and detach the switch */
        call iox $close (dialed.sw (i).iocb_ptr, code);
        if code ^= 0 then call ERROR (6);
        call iox $detach_iocb (dialed.sw (i).iocb_ptr,
            code);
        if code ^= error_table $io_no_permission
            & code ^= 0 then call ERROR (7);
        dialed.sw (i).iocb_ptr = null (); /* free
            this switch */

    end;
return;
end;

```

```

else if status flags.dialed up then do;
/* Somebody dialed in - get to work on attaching and
listening to him. As we must update a critical database
(the structure 'dialed'), let us prevent interruption
during the update period by 'masking' wakeups on the
event-call channel 'dial_manager_arg.dial_channel'. */
    call ipc $cutoff (dial_manager_arg.dial_channel, code);
    if code ^= 0 then call ERROR (4);

    /* Loop until we find a free iocb
    (indicated by a null iocb_ptr), OR
    until we exceed max_num_allowed. */
do i = 1 to max_num_allowed
    while (dialed.sw (i).iocb_ptr
        ^= null ());
end;

if ^(i > max_num_allowed) then do;
/* if there's a switch available */
dialed.devname (i) = which_channel;
dialed.no_dialed = dialed.no_dialed + 1;
call ioa $ioa switch (iox $user output,
    "9xAt this instant, ^i ^[terminal^;terminals^] ^[is^;are^] logged on.",
    dialed.no_dialed, (dialed.no_dialed = 1),
    (dialed.no_dialed = 1));

/* Find an iocb for the user and attach user's device
via tty I/O module */
    call iox $find_iocb (dialed.sw (i).swname,
        dialed.sw (i).iocb_ptr, code);
    call ipc $reconnect
        (dial_manager_arg.dial_channel, code);
        /* Safe to unmask */

    call iox $attach_ptr (dialed.sw (i).iocb_ptr,
        "tty "||dialed.sw (i).devname, null (), code);
    if code ^= 0 then call ERROR (9);

    call iox $open (dialed.sw (i).iocb_ptr,
        3 /* stream_io */,
        "0"b, /* unused must be zero */ code);
    if code ^= 0 then call ERROR (10);

```

```

/* Now that we have attached the user's terminal, call the main
program that handles these users; it will handle all
processing for these terminals, and when the user is
done, it will simply return to me */
    call listen_to_dial (dialed.sw (i).iocb_ptr);
    return;

end;
else do;          /* we've run out of switches */
    if nomore_ptr = null () then do;
        call iox $find_iocb ("nomore", nomore_ptr, code);
        /* Nope, we can not avoid this call. */
        if code ^= 0 then call ERROR (14);
    end;
    nomore_channel = which_channel;
    call iox $attach_ptr (nomore_ptr,
        "tty_" || which_channel,
        null (), code);
    if code ^= 0 then call ERROR (15);
    call iox $open (nomore_ptr, 2, "0"b, code);
        /* Stream output suffices for doomed tty */
    if code ^= 0 then call ERROR (16);
    call ioa $ioa_switch (nomore_ptr,
        "DIAL SYSTEM astra FULL WITH ^i USERS.
TRY AGAIN LATER.", max num allowed); /* SORRY FELLA */
    call iox $control (nomore_ptr, "hangup", null (), code);
    if code ^= 0 then call ERROR (17);
    call iox $close (nomore_ptr, code);
    if code ^= 0 then call ERROR (24);
    call iox $detach_iocb (nomore_ptr, code);
    if code ^= 0 &
        code ^= error_table $io no_permission
        then call ERROR (25);
    call ipc $reconnect (dial_manager_arg.dial_channel,
        code); /* Safe to unmask now */
    return;
end;
end;

ERROR:  proc (error_number);
        /* Internal proc to report errors */
dcl error_number;
        call com_err (code, ME, "Check call ^i of ERROR", error_number);
/* Restart any interrupted io to the master terminal */
        call iox $control (iox_$user_output, "start", null (), code);
        goto FINISH;
end;

```

```
shutoff: entry;
```

```
/* This entry point resets the environment:  
1) shuts off dials  
2) zeroes dialed.no_dialed  
3) wipes out iocbs and nulls ptrs. */
```

```
dcl dial_manager_$shutoff_dials entry (ptr, fixed bin (35));  
dcl ipc_$decl_ev_wait_chn entry (fixed bin (71), fixed bin (35));
```

```
ME = "shutoff";
```

```
/* IT IS FIRST NECESSARY TO CHANGE THE EVENT-CALL BACK INTO AN  
EVENT-WAIT CHANNEL, BECAUSE THE shutoff_dials ENTRY POINT WILL  
NOT WORK ON ANYTHING BUT. */
```

```
call ipc $decl ev_wait_chn (dial_manager_arg.dial_channel, code);  
if code ^= 0 then call ERROR (21);  
call dial_manager_$shutoff_dials (addr (dial_manager_arg), code);  
if code ^= 0 then call ERROR (22);  
call ipc $delete_ev_chn (dial_manager_arg.dial_channel, code);  
if code ^= 0 then call ERROR (23);  
dialed.no_dialed = 0; /* In case we 'set_up_dial' again */  
do i = 1 to max_num allowed;  
if dialed.sw(i).iocb_ptr ^= null() then do;  
call iox $close (dialed.sw(i).iocb_ptr, code);  
call iox $detach_iocb (dialed.sw(i).iocb_ptr, code);  
call iox $destroy_iocb (dialed.sw(i).iocb_ptr, code);  
/* The above call automatically NULLS the  
iocb_ptr. */
```

```
end;
```

```
end;
```

```
return;
```

```
FINISH:
```

```
end set_up_dial;
```



```

listen_to_dial: proc (iocb_ptr);
/*          Procedure to dialog with a dial-up terminal          */
#include dcl_iox_entries;
dcl
    iocb_ptr ptr parameter,
    code fixed bin (35),
    hcs $initiate count entry (char (*), char (*), char (*),
    fixed bin (24), fixed bin (2), ptr, fixed bin (35)),
    (seg_ptr_1, seg_ptr_2) internal static ptr init (null ()),
    (bit_count_1, bit_count_2) fixed bin (24) internal static,
    buff_ptr ptr,
    buffer char (256),
    n_read fixed bin (21),
    com_err_entry options (variable),
    get_wdir_entry returns (char (168));

dcl 1 info structure aligned,
    2 ev_chan fixed bin (71),
    2 input_available bit (1);
dcl (addr, after) builtin;
dcl ipc $decl_ev_call_chn entry (fixed bin (71), entry,
    ptr, fixed bin, fixed bin (35));
dcl ipc $delete_ev_chn entry (fixed bin (71), fixed bin (35));

dcl ioa $ioa_switch entry options (variable);

#include iocb; /* We need this structure to get the
    attach_descrip_ptr */

dcl 1 attach description unaligned based
    (iocb_ptr -> iocb.attach_descrip_ptr),
    2 descrip_length fixed bin (35),
    2 desc char (0
        refer (attach_description.descrip_length)) unal;

    call ioa $ioa_switch (iocb_ptr,
        "Dialed to 'astra' on channel ^a.",
        after (attach_description.desc, " "));
    if seg_ptr_1 = null () then /* Efficiency technique */
        call hcs $initiate count (get_wdir (), "msg!",
            "", bit_count_1, 1, seg_ptr_1, code);
    call iox $put_chars (iocb_ptr, seg_ptr_1,
        bit_count_1/9, code);
    if code ^= 0 then call ERROR;

```

```

call iox_$control (iocb_ptr, "read_status",
  addr (info_structure), code);
if code ^= 0 then call ERROR;
if ^info_structure.input available then do;
/* If no_line is available, use 'garden variety'
event channel provided to us by 'read_status' order */
  call ipc $decl ev call chn (info_structure.ev_chan,
    respond_to_line_later, iocb_ptr, 0, code);
/* iocb_ptr in above call is really our 'data_ptr' */
  if code ^= 0 then call ERROR;
end;
else call respond_to_line_now (iocb_ptr, 0); /* Else
there is something out there to read NOW */
return; /* so much for listen_to_dial */

```

```

respond_to_line_later: entry (information_ptr);

```

```

dcl information_ptr ptr parameter;

```

```

dcl 1 event_info based (information_ptr),
2 channel_id fixed bin (71),
2 message fixed bin (71),
2 sender bit (36),
2 origin;
  3 dev_signal bit (18) unal,
  3 ring bit (18) unal,
2 data_ptr ptr;

```

```

  call respond_to_line_now (event_info.data_ptr,
    event_info.channel_id);
/* But we are not sure whether there is a quit or
null line out there. */
return;

```

```

respond_to_line_now: entry (iocb_ptr, read_status_channel);

```

```

dcl read_status_channel fixed bin (71) parameter;
/* Channel ultimately to be deleted */

```

```

  call iox_$control (iocb_ptr, "read_status",
    addr (info_structure), code);
  if code ^= 0 then call ERROR;

```

```

if info_structure.input_available then do;
  n_read = 0;
  buff_ptr = addr (buffer);
  call iox_$get_chars (iocb_ptr, buff_ptr,
    256, n_read, code);
  if code ^= 0 then call ERROR;
  if n_read > 1 then do; /* If there's something more
    than just a newline, then do... */
    call iox_$put_chars (iocb_ptr, buff_ptr,
      n_read, code);
    if code ^= 0 then call ERROR;
    if seg_ptr_2 = null () then
      call hcs_$initiate_count (get_wdir_(),
        "msg2", "", bit_count_2, 1,
        seg_ptr_2, code);
  /* Ignore code, it should be fine */

  call iox_$put_chars (iocb_ptr, seg_ptr_2,
    bit_count_2/9, code);

  if read_status_channel ^= 0 then call
    ipc_$delete_ev_chn (read_status_channel,
      code); /* We no longer need it */
  call iox_$control (iocb_ptr, "hangup",
    null (), code);
  if code ^= 0 then call ERROR;
end;
return; /* There really is nothing else to do */
end;
else return; /* If no input_available */

ERROR: proc;

  call com_err_ (code, "listen_to_dial",
    "Truly unexpected.");
  goto return_point; /* Intentionally non-local */
end ERROR;

return_point: return;

  end listen_to_dial;

```

Quiz 1

1. Subsystem requests that accept pathnames as input may allow the final entryname in the pathname to be a star name. To first determine whether the starname is valid (ie. does not begin or end with a period etc.), the following routine should be called:

- a. hcs\_\$star
- b. hcs\_\$star\_list
- c. check\_star\_name\_
- d. match\_star\_name\_
- e. none of the above

2. A free pool of temporary segments available to each user makes it possible to use the same temporary segment more than once via the "set\_temp\_segments\_" and "release\_temp\_segments\_" routines, without having to create one when needed.

- a. temp segs belong to a single procedure when in use
- b. the maximum size of a temp seg is 16k
- c. the above mentioned free pool is system wide
- d. if all temp segs have been used, an appropriate error is returned
- e. all of the above

3. A multisegment file is composed of one or more components each the size of a segment.

- a. like single segment files, any word in a MSF can be specified by a pathname and a word offset
- b. the first component of a MSF is named component 0
- c. components are stored in a relational data base with the pathname of the MSF
- d. components are identified by consecutive upper case letters
- e. all of the above

Quiz 1

4. A call to the routine `hcs_$status_minf` is commonly used to distinguish between a segment, a directory and a MSF. Specifically the return arguments indicate a MSF when:

- a. type is 0
- b. type is 1
- c. type is 2 and bit count is 0
- d. type is 2 and bit count is nonzero
- e. none of the above

5. The working directory is the directory in which the user's activity is centered and which identifies the user's location within the storage system. The directory that becomes the working dir when the "cwd" command is given without arguments is known as the:

- a. home directory
- b. referencing directory
- c. process directory
- d. default working directory
- e. none of the above

6. The subroutine "`hcs_$set_max_length`" is often used to define the size limit of a segment so that an "out\_of\_bounds" condition will be obtained with any too-large offset.

- a. the normal maximum length of a segment is 255k
- b. the maximum length of a directory is 64k
- c. the maximum length of a directory cannot be changed
- d. system wide maximum lengths can be imposed by the system administrator
- e. all of the above

Quiz 1

7. The "msf\_manager" Multics subroutine creates a File Control Block (fcb) in order to keep track of manipulations on the MSF. Which of the followings is true?

- a. the MSF must exist for the fcb to be allocated
- b. the fcb is allocated on the stack frame of the caller
- c. the fcb is good throushout the life of the the process
- d. the address of the fcb must be siven for any operation on the MSF
- e. all of the above

8. Many Multics subroutines require an "area\_ptr" to an area in which data and information is returned to the user. This is so for example with "hcs\_\$star" which returns star name matches. Areas can be formatted by:

- a. a PL1 declaration specifying the "area" attribute
- b. the "define\_area\_" subroutine
- c. the Multics command "create\_area"
- d. use of the function "set\_system\_free\_area\_ ()"
- e. all of the above

9. The subroutine call to "hcs\_\$star" returns 2 pointers within a user provided area, star\_entry\_ptr and star\_names\_ptr. The star\_names\_ptr addresses a name array of matching names:

- a. for links
- b. for segments
- c. for directories
- d. for entries dependings upon the value of "star\_select\_sw"
- e. none of the above

Quiz 1

10. The "define\_area\_" subroutine must be provided with the address of some place to be formatted into an area. This address in the form of a pointer:

- a. is provided as argument 1 of the call to "define\_area\_"
- b. must not be null
- c. is provided within the information structure that "define\_area\_" must know about
- d. must point to a permanent segment in the storage hierarchy
- e. none of the above

Quiz 2

1. If the rins brackets for a particular object segment are (x y z) the segment is regarded as a safe when:

- a.  $x = y = z$
- b.  $x < z$
- c.  $y < z$
- d.  $x > z$
- e. none of the above

2. Inner rins procedures are often called by outer rins procedures to perform some service. It is necessary therefore for the inner rins procedure to know which rins it is working for. This rins information known as the "VALIDATION LEVEL" can be obtained via:

- a. hcs\_\$set\_rins\_brackets
- b. cu\_\$level\_set
- c. . set\_rins\_ ()
- d. cross\_rins\_
- e. none of the above

3. Many subsystems honor an "e" request (sedx calc probe etc.) What entrypoint is used in the implementation of this feature?

- a. cu\_\$generate\_call
- b. cu\_\$cp
- c. cu\_\$set\_command\_processor
- d. cu\_\$row\_stack\_frame
- e. none of the above



## Quiz 2

4. Since the "hcs\_" subroutine is the means whereby the user can manipulate rings 0 directory segments, likely rings brackets are:

- a. 0 0 0
- b. 0 1 1
- c. 0 0 5
- d. 4 4 4
- e. none of the above

5. The command level intermediary is usually not an active procedure unless an abnormal event (es. condition) occurs and it becomes invoked. The procedure responsible for invoking it is:

- a. default error handler
- b. cu\_\$cl
- c. cu\_\$set\_cl\_intermediary
- d. set\_to\_cl\_\$unclaimed\_signal
- e. none of the above

6. Saving the attachments of the standard I/O switches, restoring these attachments to their default state and entering a new loop of reading and executing command lines is part of establishing a new command level. This is accomplished by:

- a. the current command level intermediary
- b. the current command processor
- c. the current process overseer
- d. the listener (listen\_)
- e. none of the above

Quiz 2

7. The routine responsible for actually printing the ready message which indicates command level to the user is:

- a. listen\_
- b. cu\_\$ready\_proc
- c. cu\_\$set\_ready\_procedure
- d. cu\_\$set\_ready\_procedure
- e. none of the above

8. Consider the case of a user (Rins 4) attempting to execute a segment with rins brackets of {3, 3, 4}

- a. a segment can be in one rins only
- b. the user can only read the segment subject of course to ACL and AIM
- c. the user's process cannot execute in rins 3
- d. the user can execute the segment with a rins change
- e. none of the above

9. The "hcs\_\$set\_entry\_bound" routine provides the user with a method of limiting which locations of a segment may be targets of a call. If the entry-bound is to be set for a given object segment:

- a. the user must have "modify" permission on the containing directory
- b. all calls to the segment must be made to an entrypoint with offset less than the entry-bound
- c. the segment itself remains unchanged
- d. the default entry bound is 0
- e. all of the above

Quiz 2

10. The "cross\_rins\_" io\_module which allows cross rins attachments of switches:

- a. must be given the outer rins switchname as an argument
- b. is used to attach the inner rins switch to a previously existing outer rins switch
- c. enables use of the "cross\_rins\_io\_\$allow\_cross" subroutine to do cross rins io
- d. the inner rins switch must be open
- e. all of the above

### Quiz 3

1. There is actually a delay between the time a wakeup is received and the time the process is notified. It is possible therefore for several wakeups to be queued by the time a process is awakened. The priority disposition is as follows:

- a. event wait channels have priority by default
- b. event call channels have priority by default
- c. no priority exists and priority cannot be assigned
- d. no priority exists but priority can be assigned
- e. none of the above

2. When a process establishes an event call channel:

- a. the process must be blocked on the channel and wait for a wakeup to be received
- b. the channel should be polled via "ipc\_\$read\_ev\_chn" to determine whether a wakeup has been received
- c. the process should call "timer\_manaser\_\$sleep" for at least 900 seconds
- d. the process may continue executing until interrupted by a wakeup on that channel
- e. none of the above

3. On occasion, io to the user's terminal may be interrupted by an invoked "ipc\_" or "timer\_manaser\_" routine. In order to get things going again, the user should call the "iox\_\$control" subroutine with the following order:

- a. printer\_on
- b. quit\_disable
- c. start
- d. set\_delay
- e. none of the above

### Quiz 3

4. The IOCB which is the supporting structure for a switch, is created and partially initialized by either "iox\_\$attach\_name" or "iox\_\$attach\_ptr". Thereafter the IOCB is updated and maintained by:

- a. the io\_module
- b. the user program
- c. iox\_\$look\_iocb
- d. continue\_to\_signal
- e. none of the above

5. When "pointer" or "entry" variables are to be changed in a chain of synonymously attached IOCB's, the change must take place in the actual IOCB and then reflected in the other IOCB's via:

- a. <module\_name>\$<module\_name>attach
- b. iox\_\$find\_iocb
- c. iox\_\$move\_attach
- d. iox\_\$propagate
- e. none of the above

6. If a switch has been opened for "stream\_input", obviously record\_io is not supported. This implies that the "entry" value for "read\_record" in the IOCB would be:

- a. a null pointer
- b. a null character string
- c. iox\_\$err\_no\_operation
- d. error\_table\_\$no\_operation
- e. none of the above

Quiz 3

7. In order that a process communicate with another, it must know the event channel identifier of a channel created by the other process. The former process sets this info from:

- a. user\_info\_\$terminal\_data
- b. ipc\_\$read\_ev\_chn
- c. ipc\_\$decl\_ev\_call\_chn
- d. ipc\_\$decl\_ev\_wait\_chn
- e. none of the above

When a process is awakened on an event call channel, control is immediately passed to the procedure specified by the "ipc\_\$decl\_ev\_call\_chn" with one argument. This argument is a pointer to a structure that specifies:

- a. the channels on which events are being awaited
- b. information about the event that caused it to return
- c. data passed from the procedure that set up the event call channel
- d. the machine conditions at the time of the interrupt
- e. none of the above

8. The one way control path over which notification of the occurrence of events is transmitted is called an event channel. The user obtains an event channel for the purposes of inter-process communication by:

- a. looking up the channel master file (CMF)
- b. looking up the channel definition table (CDT)
- c. making one up
- d. calling ipc\_\$create\_ev\_chn
- e. none of the above

### Quiz 3

10. If a user wishes to inhibit reading of events on a particular channel but would like to have them queued for later handling, the following procedure should be invoked:

- a. ipc\_\$cutoff
- b. create\_ips\_mask\_
- c. ipc\_\$mask\_ev\_calls
- d. ipc\_\$block
- e. none of the above

## Quiz 4

1. Often to ensure that only one process at a time can execute critical section of code, that section of code is associated with a so-called "lockword" which must be zero (unlocked) and into which a process places an identifier (thus locking it) via "set\_lock\_" routine. If many critical sections of code share the same lockword:

- a. an execution time error will result
- b. each section can be executed by a different process at a given time, thus multiplexing the lockword
- c. only one process can execute in any section at a given time
- d. the lockword will be reset by any process encountering a locked lock from another section
- e. none of the above

2. The "lock identifier" placed in the lockword by a process to indicate a locked status is a:

- a. bit strings of 36 binary 1's
- b. bit strings of 36 binary zeroes
- c. unique value generated on the fly
- d. a special lock identifier kept in the active process table entry for this process
- e. none of the above

3. If a procedure that sets up timers by calling entrypoints in "timer\_manager\_" terminates abnormally, those timers which have been set will:

- a. go off at some undesired time
- b. be discarded as the procedure is popped off the stack
- c. cause the process to hang
- d. be reset by the system default condition handler
- e. none of the above



Quiz 4

4. When an "alarm" or "cpur" condition is signalled:
- a. the stack is searched for a user defined on-unit
  - b. a static handler is invoked that determines which user-specified procedure should be called
  - c. some handler executes and the user's process is returned to command level
  - d. either internal static storage for the "timer\_manager\_" subroutine has been destroyed or the system is about to crash
  - e. none of the above
5. In order that a process so blocked for a certain period of time the following subroutine should be called.
- a. timer\_manager\_\$sleep
  - b. ipc-\$block
  - c. create\_ips\_mask\_
  - d. sus\_signal\_handler\_
  - e. none of the above
6. The user\_free\_ptr stored in the process stack header points to the place where based and controlled variables etc. are stored. This is by default the:
- a. process data segment
  - b. descriptor segment
  - c. [unique].area.linker
  - d. user ring stack
  - e. none of the above

Quiz 4

7. The process stack segment is a:

- a. circular linked list
- b. single forward linked list
- c. double threaded list with forward and backward pointers
- d. stack of controlled variables
- e. none of the above

8. A null pointer has a segment number of:

- a. negative 2
- b. negative 1
- c. zero
- d. "garbage value" that causes a fault tag 2
- e. none of the above

9. The standard descriptor that complements the standard argument list indicates the argument data type to be an "offset" if the TYPE field has a value of

- a. 16
- b. 17
- c. 18
- d. 19
- e. none of the above

Quiz 4

10. The "cde" software used to build data segments for the user:

- a. creates a file in standard object format
- b. enables users to reference data using the format  
<data\_seg>\${data\_item}
- c. requires as input a pointer addressing among other things the name of the data segment
- d. invokes the PL1 compiler when the user issues the "cde" command
- e. all of the above

APPENDIX W

Workshops

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## WORKSHOP ONE

1. Copy the segment >udd>F15D>s1>STATUS.pl1. This is the source code for the program that was discussed in Topic 2. Alter this program so that it will:

-- accept the star convention

-- print out the link target if it is a link

2. When calling hcs \$star you must pass it a pointer to an area. Rather than using get\_system\_free\_area, create an area using the define\_area subroutine. Have define\_area obtain a segment for the area from the temporary segment pool. Make the area freeing, zero on free, extensible with a size of 50 words. Pass hcs \$star a pointer to that area.

Part of this workshop is to look at the contents of the area you have created. Therefore, contrary to good programming practice, do not free the structures allocated in the area by hcs \$star.

3. Test out your program. Be sure to test its ability to accept star names. As a final test give your program the argument, \*\*. Use the list\_temp\_segments command. The last temp segment will be the area you just created. Use the area\_status command with the -long argument to examine the contents of the area. (The area\_status command will accept a segment number for an argument, therefore, you can avoid typing the screech name).

You will probably want to use the following declarations which are in >udd>F15D>s1>include>w1.incl.pl1.

WORKSHOP ONE

```
dcl hcs_$get_link_target entry (char(*), char(*), char(*),
                                char(*), fixed bin(35));
dcl check_star_name_sentry entry (char(*), fixed bin(35));
dcl hcs_$star_entry (char(*), char(*), fixed bin(2), ptr,
                    fixed bin, ptr, ptr, fixed bin(35));
dcl star_entry_count fixed bin;
dcl star_entry_ptr pointer;
dcl star_names_ptr pointer;
dcl 1 star_entries (star_entry_count) aligned based (star_entry_ptr),
    2 type fixed binary (2) unsigned unaligned,
    2 nnames fixed bin (16) unsigned unaligned,
    2 nindex fixed bin (18) unsigned unaligned;

dcl star_names (sum (star_entries (*).nnames)) char (32)
    based (star_names_ptr);

dcl define_area_entry (ptr, fixed bin(35));
dcl area_info_entry (ptr, fixed bin(35));

dcl area_infop ptr;

dcl 1 area_info aligned based (area_infop),
    2 version fixed bin,
    2 control aligned like area_control,
    2 owner char (32) unal,
    2 n_components fixed bin, /* returned only */
    2 size fixed bin (18),
    2 version_of_area fixed bin, /* returned only */
    2 areap ptr,
    2 allocated_blocks fixed bin,
    2 free_blocks fixed bin,
    2 allocated_words fixed bin (30),
    2 free_words fixed bin (30);

dcl 1 area_control aligned based,
    2 extend bit (1) unal,
    2 zero_on_alloc bit (1) unal,
    2 zero_on_free bit (1) unal,
    2 dont_free bit (1) unal,
    2 no_freeing bit (1) unal,
    2 system bit (1) unal,
    2 pad bit (30) unal;
```

Changing the Command Environment

The command environment can be shaped in several ways by using cu entry points. In this workshop, you will write a procedure which will 'intercept' commands before they get to your usual command processor, and by so doing, restrict yourself to a limited set of commands. In addition, your own version of a ready procedure will be installed.

In order to use some internal static variables it is recommended that you write just one procedure with alternate entry points. Your procedure should have the following skeleton:

```
change_env: proc;
    <declarations>
    <code to save entry value of current command processor
      and install interceptor and new ready procedure>
restore_cp: entry;
    <code to reinstall the command processor whose entry
      value you saved>
my_ready: entry;
    <code for new ready procedure>
command_interceptor: entry (a_line_ptr, a_line_len, code);
    <code for your command interceptor>
end change_env;
```

## WORKSHOP TWO

1. The 'my\_ready' entry point should determine if you are in 'ready on' mode. If so, it should output the following line:

'Next Command: '

You should not append a new-line character at the end of this 'ready message'.

2. The 'command\_interceptor' entry point will intercept commands before they get to your usual command processor (typically 'command\_processor' or 'abbrev'). It should examine the a line to determine whether the command line contains any of the following commands:

```
new_proc
logout
probe, pb
exec_com, ec
ready_on, rdn
ready_off, rdf
restore_cp
```

If one of these commands is input to command\_interceptor, it should pass the parameters along to whatever was the previous command processor. If any other command is input, the command\_interceptor should print out a message informing the user that the command is not allowed, and should then return.

3. The 'change\_env' entry point should establish command\_interceptor as the current command processor and the my\_ready procedure as the current ready procedure. Note that this procedure should first determine the value of the current command processor and make that value available to the 'command\_interceptor' and 'restore\_cp' entry points.
4. The 'restore\_cp' entry point should restore your usual command processor environment (ie. disable the activity of command\_interceptor).

Test out your solutions by executing 'change\_env'. You should now be running under the modified command environment. Try typing some disallowed commands and some allowed commands. Especially observe the behavior of ready on and ready off. Verify your ability to return to your usual command processor environment by executing 'restore\_cp'.

You will probably want to use the following declarations which are in >udd>F15D>s1>include>w2.incl.pl1.



WORKSHOP TWO

```
dcl a_line_ptr ptr;
dcl a_line_len fixed bin (21);
dcl a_line_char (a_line_len) based (a_line_ptr);
dcl com_err_entry_options (variable);
dcl (ioa $nrl, ioa_) entry options (variable);
dcl (ltrim, rtrim, substr, index, null, codeptr) builtin;
dcl cu_$set_ready_procedure entry (entry);
dcl cu_$get_command_processor entry (entry);
dcl cu_$set_command_processor entry (entry);
dcl cu_$get_ready_mode entry (1 aligned, 2 bit (1) unaligned,
                               2 bit (35) unaligned);
dcl old_command_processor entry variable options (variable)
                               internal static;
dcl command_processor_$command_processor entry (ptr,
                                                  fixed bin (21), fixed bin (35));
dcl 1 mode aligned,
    2 ready_sw bit (1) unaligned,
    2 mbz bit (35) unaligned;
dcl code fixed bin (35);
dcl end fixed bin (21);
dcl test_string char (80) init (" new_proc logout probe pb exec_com
ec ready_on_rdn ready_off rdf restore_cp")
                               internal static options (constant);
```

General Outline for Workshop 2, F15D

```
change_env:  proc;

    /* Variable and subroutine declarations */

    /* Get old command processor and old ready procedure and
       store them in internal entries */

    /* Set the new command processor and new ready
       procedure.  These both are entries which are internal
       to procedure 'change_env' */

    /* Pop the 'change_env' frame off the stack via a
       'return' statement */

my_ready:  entry;

    /* 'my_ready' is the entry specified above in the
       subroutine call to set the new ready procedure */

    /* check the value of the 'ready_sw'.  If ready_sw is
       true ("1"b) then print your ready message */

    /* pop the frame off the stack */

restore_ready:  proc;

    /* 'restore_ready' should simply make a call to the
       appropriate subroutine to set the ready message back
       to the entry value that was stored above */

    /* Pop the frame off the stack */

command_interceptor:  entry(a_line_ptr, a_line_len, code);

    /* 'command_interceptor' is the entry specified above in
       the call to the subroutine to set the new command
       processor */

    /* Process the command.  If it is a valid command, i.e.
       is a command contained in the variable 'test_string',
       then pass the command to the subroutine which
       accesses the default command processor.  Otherwise,
       print an error message */

    /* Pop the frame off the stack */

restore_cp:  entry;

    /* 'restore_cp' should simply make a call to the
       appropriate subroutine to set the command processor
       back to the entry value that was stored above */

    /* Pop the frame off the stack */

end change_env;
```

## WORKSHOP THREE

### Synonyming and IOCB's

1. Use the command 'io\_call print iocb <switchname>' to examine the IOCB's for user\_i/o, user\_input, user\_output and error\_output. Verify that the diagram showing synonyming at the end of Topic 6 is correct.
2. Look at the declaration of an 'IOCB' in Topic 6. List the structure members where a synonymed switch's IOCB may differ from the IOCB of the switch to which it is synonymed, assuming no inhibition in effect. There are 7 of them.

3. Referring back to part 1, look carefully at the entry value for iocb.put\_chars for all four switches examined. Try doing a 'io call put\_chars <name> hello' using each of the four switches. You should be able to explain exactly what happened, in terms of what IOCB's were referenced, what arguments were passed to which entry point of which I/O module and how the code in that module caused the observed result. If not, ask your instructor.

What will happen if you execute 'io\_call close error\_output'? Try it.

4. Detach user\_output. Attach it using the syn\_ module, however, this time inhibit the close I/O operation.

```
'io_call attach user_output syn_user_i/o -inhibit close'
```

Use io\_call to examine the IOCB for user\_output. What two places do you see evidence that the close operation has been inhibited?

Try to close user\_output.

Using io\_call is an excellent way to learn about and test the Multics I/O mechanism. In any remaining time, feel free to experiment, but be aware of the fact that you may get yourself into some situations which require that you hang up the phone to get out of them.

## WORKSHOP FOUR

### An Interprocess Communication Workshop

The interprocess communication facility is used to coordinate processing between separate processes. In this workshop, you will design and implement a procedure which will communicate with the instructor's process as follows:

Write a pl1 procedure called 'get\_message.pl1' which will:

- 1) Create an event-call channel specifying that the procedure entry point 'reverse message' is to be invoked when a wakeup is received on the event-call channel. The handler called 'reverse message' will be an entry point in the 'get\_message' procedure itself (see step 3 below).
- 2) Obtain the process id and channel id of the instructor's event call channel by initiating the segment >udd>F15D>s1>channel\_info and using the following structure:

```
{ dcl 1 channel_info_overlay based(channel_info_ptr),  
  2 process_id    bit(36),  
  2 channel_id    fixed bin (71); for base of segment
```

Using this information, send a wakeup to the instructor's process providing the channel id of your event call channel as the 'message'.

- 3) Your 'reverse message' handler will be invoked when the instructor's process sends your process a wakeup. This handler should interpret the 'message' from the instructor as a 'char(8)' string; the handler should reverse this string (the pl1 'reverse' builtin function can be used) and finally, your process should send another wakeup to the instructor's process using the reversed string as the 'message'.

Note that the instructor's process will be able to determine whether or not you have successfully accomplished the task. If you have, you will receive immediate notification from the instructor's process (via 'send\_mail'). Therefore, make sure you are in 'accept message' mode prior to testing your solution. In addition, use the 'who' command to verify that the instructor's absentee process is indeed running prior to testing your solution.

You will probably want to use the following declarations which are in >udd>F15D>s1>include>w4.incl.pl1.

*PSOEA >WKSPS > F15D>s1 >incl > w4.incl.pl1*

WORKSHOP FOUR

```
dcl (ioa_, com_err_) entry options (variable),  
hcs_$wakeup entry (bit (36), fixed bin (71), fixed bin (71),  
fixed bin (35)),  
ipc_$create_ev_chn entry (fixed bin (71), fixed bin (35)),  
ipc_$decl_ev_call_chn entry (fixed bin (71), entry, ptr, fixed bin,  
fixed bin (35)),  
hcs_$initiate entry (char (*), char (*), char (*),  
fixed bin (1), fixed bin (2), ptr, fixed bin (35)),  
hcs_$terminate_noname entry (ptr, fixed bin (35));  
  
dcl my_chid fixed bin (71),  
c_ptr ptr static,  
i_ptr ptr,  
code fixed bin (35);  
  
dcl 1 channel_info based (c_ptr),  
2 his_pid bit (36),  
2 his_chid fixed bin (71);  
  
dcl 1 event_info based (i_ptr),  
2 channel_id fixed bin (71),  
2 message fixed bin (71),  
2 sender bit (36),  
2 origin,  
3 dev_signal bit (18) unal,  
3 ring bit (18) unal,  
2 data_ptr ptr;  
  
dcl (null, reverse) builtin;
```

## WORKSHOP FIVE

### Timers

Write a program called lock.pl1 that implements the following fictitious command:

*no args = default = 5 mins*

USAGE: lock [-min minutes]

*lock -min 5*

FUNCTION: Prompts the user for a password and locks the user's terminal for the number of minutes specified (default = 10 minutes). To regain control of the terminal the user hits the break key and is again prompted for the password. If he does not supply the correct password, the terminal remains locked. *= incorrect password*

When the time specified expires, the user is logged out.

NOTES: This command prints out the following message:

- It is MM/DD/YY hhmm.m zzz www.
- This terminal is locked, please find another terminal.

This message is printed out when lock is first invoked and every 5 minutes thereafter.

Essentially your program should go to sleep for a specified length of time. However, during that period, you must also repeat the message every 5 minutes.

This command takes either 0 or 2 arguments.

For the sake of debugging you should probably use a time interval shorter than 5 minutes and rather than logging out the user, simply print some message.

WORKSHOP FIVE

Sample terminal session:

```
! lock -min 20
  Password:
!           <- user supplies password

  It is 07/21/81 0812.2 mst Tue.
  This terminal is locked, please find another terminal.

  It is 07/21/81 0817.2 mst Tue.
  This terminal is locked, please find another terminal.

! <QUIT>
  Password:
!           <- user gives wrong password

  Incorrect password given.

  It is 07/21/81 0822.2 mst Tue.
  This terminal is locked, please find another terminal.

  It is 07/21/81 0827.2 mst Tue.
  This terminal is locked, please find another terminal.

! <QUIT>
  Password:
!           <- user gives correct password

  Terminal unlocked.
```

---

```
! lock -min 1
  Password:
!

  It is 07/21/81 0905.1 mst Tue.
  This terminal is locked, please find another terminal.

  The lock time for your terminal has expired.
  You will be logged out.

  NDibble MED logged out 07/21/81 0906.3 mst Tue.
  CPU usage 12 sec, memory usage 97.3 units, cost $3.55.
  hangup
```

You will probably want to use the following declarations which are in  
>udd>F15D>s1>include>w5.incl.pl1.

WORKSHOP FIVE

~~dcl prompt char(6) init ("Enter password:");~~

dcl cu \$arg\_count entry (fixed bin, fixed bin(35));

dcl cu \$arg\_ptr entry (fixed bin, ptr, fixed bin(21), fixed bin(35));

dcl nargs;

dcl aptr ptr;

dcl arg char(arg\_length) based (aptr);

dcl arg\_length fixed bin (21);

~~dcl password char(20) varying;~~

dcl read\_password\_entry (char(\*), char(\*));

\* → dcl timer\_manager\_\$sleep entry (fixed bin(71), bit(2));

dcl timer\_manager\_\$alarm\_call entry (fixed bin(71), bit(2), entry);

dcl timer\_manager\_\$reset\_alarm call entry (entry);

dcl time fixed bin(71) init (10);

dcl cv\_dec\_check\_entry (char(\*), fixed bin(35)) returns(fixed bin(35));

dcl quit condition;

dcl clock\_entry() returns(fixed bin(71));

dcl date\_time\_entry (fixed bin(71), char(\*));

dcl (password1, password2) char(8);

dcl code fixed bin(35);

dcl (com\_err\_, ioa\_) entry() options(variable);



## WORKSHOP SIX

### Generating Calls and Creating Status Tables

1. Write a procedure called 'please\_execute.pl1' which will ask the user, "What do you want me to execute for you? ". The user may then respond with an entryname. The entryname could be a command or a user written program. To keep things simple, no arguments are allowed (ie. the user may respond 'list', but not 'list -d'). Also, you may assume the user's response is both the entryname and the entry point name. Your program should then generate a call to execute the user's request.

To save time, you will probably want to make a copy of the segment, >udd>F15D>s1>generate\_pwd.pl1 and simply edit in the necessary changes.

2. Try out your procedure responding with commands such as 'pwd', 'hmu' and 'list'. Then use it to execute some simple programs you have written or copied into your working directory during this course. You might even like to ask 'please\_execute' to execute 'please\_execute'.
3. Ask 'please\_execute' to execute the 'list' command and when it starts to list your segments, hit the break/quit/interrupt key. Use the 'stack' request from within 'probe' to examine the user stack.

#### NUMBERS 4 AND 5 ARE OPTIONAL

4. Next ask 'please\_execute' to execute some program that does not exist (make up a name). Notice the error message returned when your program calls com\_err\_.
5. Create a status table and alter your program so that when you ask your program to execute a program that does not exist, it will still call com\_err\_, however, it will not behave as in part 4. Instead it should print some other error message (one you have made up and put in a status table you have built).

## WORKSHOP SEVEN

### The Limited Service Subsystem

In workshop two you modified the user's environment such that only certain commands were acceptable. Utilize the 'enter\_lss' command to achieve the same result. The following list of commands should be accepted (you may expand upon this list if you want to):

```
new_proc
logout
probe, pb
exec_com, ec
ready_on, rdn
ready_off, rdf
```

Test your solution.

Large Information Systems Div, Honeywell Information Systems, Inc.  
Multics Computer Center, Phoenix Az.

Printout of the 5 Entries  
of the  
handout Library

Which Match the Search Name  
gw.archive

Printed on: 04/06/81 0804.0  
Printed by: NDibble.MED.a  
Descriptor: handout\_desc

## COMPILATION LISTING OF SEGMENT pfgu

Compiled by: Multics PL/I Compiler, Release 25c, of February 18, 1980

Compiled at: Honeywell LISD Phoenix, System M

Compiled on: 03/19/80 1043.5 mst Wed

Options: map

```
1 find: proc (last_name, first_name, emp_struct, code);
2
3 /* Modified by R. Frommer on June 12, 1979 to do elegant ORing, reference mode bits directly, and use -extend option */
4 /* THESE DECLARATIONS SHARED BY BOTH ENTRYPOINTS */
5
6 dcl acs_names_array (4) char (32) internal static options (constant)
7     init ("soc_sec_no.acs", "manager.acs", "mail_station.acs", "salary.acs");
8 dcl access_allowed (4) bit (1);
9 dcl access_allowed_overlay bit(4) unaligned defined (access_allowed);
10 dcl dir_name char (168);
11 dcl user_id char (32);
12 dcl ring_fixed bin (17);
13 dcl mode_fixed bin (5);
14 dcl 1 bit mode_struct aligned based (mode_ptr),
15     2 pad_bit (30) unaligned,
16     2 two_bits bit(2) unaligned, /* don't care */
17     2 read_bit (1) unaligned, /* On if read effective access */
18     2 fourth_bit (1) unaligned, /* don't care */
19     2 write_bit(1) unaligned; /* On if write effective access */
20 dcl ME char (4);
21 dcl find_bit (1);
22 dcl 1 emp_struct,
23     2 soc_sec_no pic "(9)9",
24     2 manager_pic "(9)9",
25     2 mail_station char (3),
26     2 salary_fixed dec (8, 2);
27 dcl (code, dummy_code) fixed bin (35);
28 dcl (first_name, last_name) char (15);
29 dcl (substr, null, addr, size, rtrim) builtin;
30 dcl error_table_moderr ext static fixed bin (35);
31 dcl error_table_no_record ext static fixed bin (35);
32 dcl 1 fixed bin;
33 dcl (ioobp, mode_ptr) ptr init (null ());
34 dcl unused_rec_len fixed bin (21);
35 dcl old_level fixed bin;
36
1 dcl
1 2 iox_attach_ptr entry (ptr, char(*), ptr, fixed bin (35)),
1 3 iox_attach_name entry (char(*), ptr, char(*), ptr, fixed bin (35)),
1 4 iox_close entry (ptr, fixed bin (35)),
1 5 iox_control entry (ptr, char(*), ptr, fixed bin(35)),
1 6 iox_delete_record entry (ptr, fixed bin(35)),
1 7 iox_detach_ioob entry (ptr, fixed bin(35)),
```

```

1   8 iox_$find_iocb entry (char(*), ptr, fixed bin(35)),
1   9 iox_$get_chars entry (ptr, ptr, fixed bin(21), fixed bin(21), fixed bin(35)),
1  10 iox_$get_line entry (ptr, ptr, fixed bin(21), fixed bin (21), fixed bin (35)),
1  11 iox_$modes entry (ptr, char(*), char(*), fixed bin (35)),
1  12 iox_$move_attach entry (ptr, ptr, fixed bin(35)),
1  13 iox_$open entry (ptr, fixed bin, bit(1) aligned, fixed bin(35)),
1  14 iox_$position entry (ptr, fixed bin, fixed bin(21), fixed bin (35)),
1  15 iox_$put_chars entry (ptr, ptr, fixed bin(21), fixed bin(35)),
1  16 iox_$read_key entry (ptr, char(256) varying, fixed bin(21), fixed bin (35)),
1  17 iox_$read_length entry (ptr, fixed bin(21), fixed bin(35)),
1  18 iox_$read_record entry (ptr, ptr, fixed bin(21), fixed bin(21), fixed bin (35)),
1  19 iox_$rewrite_record entry (ptr, ptr, fixed bin (21), fixed bin(35)),
1  20 iox_$seek_key entry (ptr, char(256) varying, fixed bin(21), fixed bin(35)),
1  21 iox_$write_record entry (ptr, ptr, fixed bin(21), fixed bin (35)),
1  22 iox_$destroy_iocb entry (ptr, fixed bin(35));
37
38
39 dcl cu_$level_get entry (fixed bin);
40 dcl cu_$level_set entry (fixed bin);
41 dcl get_ring_entry () returns (fixed bin (3));
42 dcl hos_$get_user_effmode entry (char (*), char (*), char (*), fixed bin, fixed bin (5), fixed bin (35));
43 dcl get_group_id_$tag_star entry () returns (char (32));
44
45     ME = "find";
46     find = "1"b;
47     goto COMMON_CODE;
48
49 put:   entry (last_name, first_name, emp_struct, code);
50     ME = "put";
51     find = "0"b;
52
53 COMMON_CODE:
54
55     call cu_$level_get (old_level);
56     ring = get_ring ();
57     call cu_$level_set (ring);
58
59     code = 0;
60
61     dir_name = ">udd>F15d>s1";
62     user_id = get_group_id_$tag_star ();
63     call hos_$get_user_effmode (dir_name, "emp", user_id, ring, mode, code);
64     if code = 0 then go to return_point;
65     mode_ptr = addr (mode);
66     if
67     (^find & ^(bit_mode_struct.read & bit_mode_struct.write)) | /* if put, we must have read and write on emp */
68     (find & ^bit_mode_struct.read) then do; /* If find, we must have read on emp */
69         code = error_table_$moderr;
70         goto return_point;
71     end;
72
73 /* If we get here, then it must be the case that we have status on
74 * our current working directory. */
75
76 /* Obtain the effective access on the access control segments of interest */

```

```

77
78 do i = 1 to 4; /* Once for each access control seg */
79
80 call hcs $get_user_effmode (dir_name, aoc_names_array (i), user_id, ring, mode, code);
81 if find then access_allowed (i) = bit_mode_struct.read;
82 else access_allowed (i) = bit_mode_struct.write;
83 end;
84
85 /* ONE OF TWO BLOCKS OF CODE, DEPENDING UPON ENTRYPOINT TAKEN... */
86
87 if find then do;
88
89 if access_allowed_overlay then do; /* elegant ORing */
90 call iox $find iocb ("emp_sw", iocbp, code);
91 call iox $attach ptr (iocbp, "vfile_" || rtrim (dir_name) || ">emp", null (), code);
92 call iox $open (iocbp, 11 /* direct input */, "0*b", code);
93 call iox $seek key (iocbp, (last_name || first_name), 4 * size (emp_struct), code);
94 if code = 0 then do;
95 call iox $read_record (iocbp, addr (emp_struct), 4 * size (emp_struct), unused_rec_len, code);
96 if ^access_allowed (1) then soc_sec_no = 0;
97 if ^access_allowed (2) then manager = 0;
98 if ^access_allowed (3) then mail_station = "";
99 if ^access_allowed (4) then salary = 0;
100 end;
101 else; /* else nothing */
102 call iox $close (iocbp, dummy_code);
103 call iox $detach iocb (iocbp, dummy_code);
104 goto return_point; /* To reset validation level */
105 end;
106
107 else code = 1; /* This special code indicates that no fields were accessible
*/
108
109 end;
110 else do;
111
112 if access_allowed (1) & access_allowed (2) & access_allowed (3) & access_allowed (4) then do; /* If all of
/*them are write-able */
113 call iox $find iocb ("emp_sw", iocbp, code);
114 call iox $attach ptr (iocbp, "vfile_" || rtrim (dir_name) || ">emp -extend", null (), code);
115 call iox $open (iocbp, 12 /* direct output */, "0*b", code);
116 call iox $seek key (iocbp, (last_name || first_name), 4 * size (emp_struct), code);
117 if code = error_table $no_record then call iox $write_record (iocbp, addr (emp_struct), 4 * size (emp_str
/*cuc), code);
118 call iox $close (iocbp, dummy_code);
119 call iox $detach iocb (iocbp, dummy_code);
120 goto return_point;
121
122 end;
123
124 else code = 1; /* This special code indicates that some fields were inaccessi
/*ble */
125
126 /* FALL THROUGH TO return_point */
127 end;

```

```
128
129 return_point:
130     call cu_$level_set (old_level);
131     end find;
```

SOURCE FILES USED IN THIS COMPILATION.

LINE	NUMBER	DATE MODIFIED	NAME	PATHNAME
	0	07/19/79 1110.0	pfgu.pl1	>user_dir_dir>F15D>Student_01>pfgu.pl1
37	1	02/13/79 1714.4	dcl_iox_entries.incl.pl1	>user_dir_dir>F15D>Student_01>dcl_iox_entries.incl.pl1



NAMES DECLARED IN THIS COMPILATION.

IDENTIFIER	OFFSET	LOC	STORAGE CLASS	DATA TYPE	ATTRIBUTES AND REFERENCES (* indicates a set context)
NAMES DECLARED BY DECLARE STATEMENT.					
NE	000165		automatic	char(4)	unaligned dcl 20 set ref 45* 50*
access_allowed	000100		automatic	bit(1)	array unaligned dcl 8 set ref 81* 82* 89 89 96 97 98 99 112 112 112 112
access_allowed_overlay			defined	bit(4)	unaligned dcl 9 ref 89
acs_names_array	000000		constant	char(32)	initial array unaligned dcl 6 set ref 80*
addr				builtin function	dcl 29 ref 65 95 95 117 117
bit_mode_struct			based	structure	level 1 dcl 14
code			parameter	fixed bin(35,0)	dcl 27 set ref 1 49 59* 63* 64 69* 80* 90* 91* 92* 93* 94 95* 107* 113* 114* 115* 116* 117 117* 124*
cu_level_get	000034		constant	entry	external dcl 39 ref 53
cu_level_set	000036		constant	entry	external dcl 40 ref 57 129
dir_name	000101		automatic	char(168)	unaligned dcl 10 set ref 61* 63* 80* 91 114
dummy_code	000167		automatic	fixed bin(35,0)	dcl 27 set ref 102* 103* 118* 119*
emp_struct			parameter	structure	level 1 unaligned dcl 22 set ref 1 49 93 95 95 95 116 117 117 117
error_table_moderr	000010		external static	fixed bin(35,0)	dcl 30 ref 69
error_table_no_record	000012		external static	fixed bin(35,0)	dcl 31 ref 117
find	000166		automatic	bit(1)	unaligned dcl 21 set ref 46* 51* 66 66 81 87
first_name			parameter	char(15)	unaligned dcl 28 ref 1 49 93 116
get_group_id_tag_star	000044		constant	entry	external dcl 43 ref 62
get_ring	000040		constant	entry	external dcl 41 ref 56
hcs_get_user_effmode	000042		constant	entry	external dcl 42 ref 63 80
i	000170		automatic	fixed bin(17,0)	dcl 32 set ref 78* 80 81 82*
ioebp	000172		automatic	pointer	initial dcl 33 set ref 33* 90* 91* 92* 93* 95* 102* 103* 113* 114* 115* 116* 117* 118* 119*
iox_attach_ptr	000014		constant	entry	external dcl 1-1 ref 91 114
iox_close	000016		constant	entry	external dcl 1-1 ref 102 118
iox_detach_ioeb	000020		constant	entry	external dcl 1-1 ref 103 119
iox_find_ioeb	000022		constant	entry	external dcl 1-1 ref 90 113
iox_open	000024		constant	entry	external dcl 1-1 ref 92 115
iox_read_record	000026		constant	entry	external dcl 1-1 ref 95
iox_seek_key	000030		constant	entry	external dcl 1-1 ref 93 116
iox_write_record	000032		constant	entry	external dcl 1-1 ref 117
last_name			parameter	char(15)	unaligned dcl 28 ref 1 49 93 116
mail_station	4(18)		parameter	char(3)	level 2 packed unaligned dcl 22 set ref 98*
manager	2(09)		parameter	picture(9)	level 2 packed unaligned dcl 22 set ref 97*
mode	000164		automatic	fixed bin(5,0)	dcl 13 set ref 63* 65 80*
mode_ptr	000174		automatic	pointer	initial dcl 33 set ref 33* 65* 66 66 66 81 82
null				builtin function	dcl 29 ref 33 33 91 91 114 114
old_level	000177		automatic	fixed bin(17,0)	dcl 35 set ref 53* 129*
read	0(32)		based	bit(1)	level 2 packed unaligned dcl 14 ref 66 66 81
ring	000163		automatic	fixed bin(17,0)	dcl 12 set ref 56* 57* 63* 80*
rtrim				builtin function	dcl 29 ref 91 114
salary	6		parameter	fixed dec(8,2)	level 2 dcl 22 set ref 99*
size				builtin function	dcl 29 ref 93 95 116 117
soc_sec_no			parameter	picture(9)	level 2 packed unaligned dcl 22 set ref 96*
unused_rec_len	000176		automatic	fixed bin(21,0)	dcl 34 set ref 95*
user_id	000153		automatic	char(32)	unaligned dcl 11 set ref 62* 63* 80*
write	0(34)		based	bit(1)	level 2 packed unaligned dcl 14 ref 66 82

NAMES DECLARED BY DECLARE STATEMENT AND NEVER REFERENCED.

iox_\$attach_name	000000	constant	entry	external dcl 1-1
iox_\$control	000000	constant	entry	external dcl 1-1
iox_\$delete_record	000000	constant	entry	external dcl 1-1
iox_\$destroy_iocb	000000	constant	entry	external dcl 1-1
iox_\$get_chars	000000	constant	entry	external dcl 1-1
iox_\$get_line	000000	constant	entry	external dcl 1-1
iox_\$modes	000000	constant	entry	external dcl 1-1
iox_\$move_attach	000000	constant	entry	external dcl 1-1
iox_\$position	000000	constant	entry	external dcl 1-1
iox_\$put_chars	000000	constant	entry	external dcl 1-1
iox_\$read_key	000000	constant	entry	external dcl 1-1
iox_\$read_length	000000	constant	entry	external dcl 1-1
iox_\$rewrite_record	000000	constant	entry	external dcl 1-1
substr			builtin function	dcl 29

NAMES DECLARED BY EXPLICIT CONTEXT.

COMMON_CODE	000145	constant	label	dcl 53 ref 47
find	000117	constant	entry	external dcl 1
put	000134	constant	entry	external dcl 49
return_point	001200	constant	label	dcl 129 ref 64 70 104 120

THERE WERE NO NAMES DECLARED BY CONTEXT OR IMPLICATION.

STORAGE REQUIREMENTS FOR THIS PROGRAM.

	Object	Text	Link	Symbol	Defs	Static
Start	0	0	1414	1462	1214	1424
Length	1700	1214	46	201	200	0

BLOCK NAME	STACK SIZE	TYPE	WHY NONQUICK/WHO SHARES STACK FRAME
find	233	external procedure	is an external procedure.

STORAGE FOR AUTOMATIC VARIABLES.

STACK FRAME	LOC IDENTIFIER	BLOCK NAME
find	000100 access_allowed	find
	000101 dir_name	find
	000153 user_id	find
	000163 ring	find
	000164 mode	find
	000165 ME	find
	000166 find	find
	000167 dummy_code	find
	000170 i	find
	000172 iocbp	find
	000174 mode_ptr	find
	000176 unused_rec_len	find
	000177 old_level	find

THE FOLLOWING EXTERNAL OPERATORS ARE USED BY THIS PROGRAM.

alloc_cs	cat_realloc_cs	call_ext_out_desc	call_ext_out	return	shorten_stack
ext_entry					

THE FOLLOWING EXTERNAL ENTRIES ARE CALLED BY THIS PROGRAM.

cu_\$level_get	cu_\$level_set	get_group_id_\$tag_star	get_ring_
hcs_\$get_user_effmode	iox_\$attach_ptr	iox_\$close	iox_\$detach_iocb
iox_\$find_iocb	iox_\$open	iox_\$read_record	iox_\$seek_key
iox_\$write_record			

THE FOLLOWING EXTERNAL VARIABLES ARE USED BY THIS PROGRAM.

error_table_\$moderr	error_table_\$no_record
----------------------	-------------------------

LINE	LOC	LINE	LOC	LINE	LOC	LINE	LOC	LINE	LOC	LINE	LOC	LINE	LOC
33	000105	1	000112	45	000125	46	000127	47	000131	49	000132	50	000142
51	000144	53	000145	56	000154	57	000165	59	000174	61	000176	62	000201
63	000210	64	000250	65	000254	66	000256	69	000303	70	000306	78	000307
80	000314	81	000352	82	000362	83	000367	87	000371	89	000373	90	000376
91	000422	92	000507	93	000532	94	000570	95	000574	96	000617	97	000630
98	000641	99	000651	102	000661	103	000672	104	000703	107	000704	109	000707
112	000710	113	000732	114	000756	115	001043	116	001066	117	001124	118	001152
119	001163	120	001174	124	001175	129	001200	131	001207				

## COMPILATION LISTING OF SEGMENT put

Compiled by: Multics PL/I Compiler, Release 25c, of February 18, 1980

Compiled at: Honeywell LISD Phoenix, System M

Compiled on: 03/19/80 1039.0 mst Wed

Options: map

```
1 /* Command interface to the put_find_gate_gate entrypoints */
2
3 /* Modified on June 12, 1979 by R. Frommer to produce better diagnostic messages. */
4 find: proc;
5
6 /* DECLARATIONS */
7
8 del ME char (4) varying;
9 del nargs fixed bin;
10 del argl fixed bin (21);
11 del (first_name, last_name) char (15);
12 del code fixed bin (35);
13 del arg char (argl) based (argp);
14 del argp ptr;
15 del 1 emp_struct,
16     2 soc_sec_no pic "(9)9",
17     2 manager_pic "(9)9",
18     2 mail_station char (3),
19     2 salary fixed dec (8, 2);
20 del (error_table $bigarg, error_table $wrong_no_of_args) ext static fixed bin (35);
21 del (fixed_overflow, conversion, size) condition;
22
23 del cu $arg_ptr entry (fixed bin, ptr, fixed bin (21), fixed bin (35));
24 del cu $arg_count entry (fixed bin);
25 del (com_err, ioa) entry options (variable);
26 del (put_find_gate $find, put_find_gate $put) entry
27     (char (15), char (15), 1 structure, 2 pic "(9)9" member, 2 pic "(9)9" member,
28     2 char (3) member, 2 fixed dec (8, 2) member, fixed bin (35));
29 del cv_dec_check_entry (char(*), fixed bin(35)) returns (fixed bin(35));
30
31     ME = "find";
32     call cu $arg_count (nargs);
33     if nargs ^= 2 then do;
34         call com_err (error_table $wrong_no_of_args, ME,
35             "~/Correct invocation: find First_name Last_name");
36     return;
37 end;
38 call cu $arg_ptr (1, argp, argl, code);
39 if argl > 15 then do;
40     call com_err (error_table $bigarg, ME, "~a exceeds 15 characters.", arg);
41 return;
42 end;
43 else first_name = arg;
```

```

44     call ou $arg_ptr (2, argp, arg1, code);
45     if arg1 > 15 then do;
46         call com_err_ (error_table_$bigarg, ME, "a exceeds 15 characters.", arg);
47     return;
48     end;
49     else last_name = arg;
50
51 /* Initialize emp_struct prior to calling the gate endpoint */
52
53     soc_sec_no, manager, salary = 0;
54     mail_station = "";
55
56     call put_find_gate $find (last_name, first_name, emp_struct, code);
57     if code = 1 /* this is a special case */ then
58         call com_err_ (0, ME, "You don't have proper access on the fields in the employee file.");
59     else if code = 0 then call com_err_ (code, ME);
60     else call ioa ("^[^s^;SOCIAL SECURITY NUMBER = ^i
61 ^]^[^s^;MANAGER'S SOCIAL SECURITY NUMBER = ^i
62 ^]^[^s^;MAIL STATION = ^a
63 ^]^[^s^;SALARY = $^i
64 ^]", (soc_sec_no = 0), soc_sec_no,
65         (manager = 0), manager,
66         (mail_station = ""), mail_station,
67         (salary = 0), salary);
68     return; /* END OF find PROCEDURE */
69
70 put:   entry ();
71
72     ME = "put";
73     call ou $arg_count (nargs);
74     if nargs ^= 6 then do;
75         call com_err_ (error_table $wrong_no_of_args, ME,
76             "^/Correct invocation: put First_name Last_name Employee's_ss_no Manager's_ss_no Mail_Station Salary")
77     return;
78     end;
79
80 /* OBTAIN SIX REQUIRED ARGUMENTS */
81
82     call ou $arg_ptr (1, argp, arg1, code);
83     if arg1 > 15 then do;
84         call com_err_ (error_table_$bigarg, ME, "a", arg);
85     return;
86     end;
87     else first_name = arg;
88
89     call ou $arg_ptr (2, argp, arg1, code);
90     if arg1 > 15 then do;
91         call com_err_ (error_table_$bigarg, ME, "a", arg);
92     return;
93     end;
94     else last_name = arg;
95
96     call ou $arg_ptr (3, argp, arg1, code);
97     if arg1 ^= 9 | verify (arg, "0123456789") ^= 0 then do;

```

```

98         call com_err_ (0, ME, "a is not a suitable social security number for a.",
99             arg, last_name);
100        return;
101    end;
102    else soc_sec_no = cv_dec_check_(arg, code);
103
104    call cu $arg_ptr (4, argp, argl, code);
105    if argl ^= 9 | verify (arg, "0123456789") ^= 0 then do;
106        call com_err_ (0, ME, "a is not a suitable social security number for a's manager.", arg, last_name);
107        return;
108    end;
109    else manager = cv_dec_check_(arg, code);
110
111    call cu $arg_ptr (5, argp, argl, code);
112    if argl ^= 3 then do;
113        call com_err_ (0, ME, "a is not a suitable 3-character mail drop.", arg);
114        return;
115    end;
116    else mail_station = arg;
117
118    call cu $arg_ptr (6, argp, argl, code);
119
120    on fixedoverflow, size begin;
121        call com_err_ (0, ME, "a is not a proper salary. It must resemble 999999.99", arg);
122        goto bottom;
123    end;
124    salary = cv_dec_check_(arg, code);
125    if code ^= 0 then do;
126        call com_err_ (0, ME, "a is not a proper salary. It must resemble 999999.99", arg);
127        goto bottom;
128    end;
129
130    call put_find_gate $put (last_name, first_name, emp_struc, code);
131    if code = 1 then call com_err_ (0, ME, "You lack access on one or more of the fields in the database.");
132    else if code ^= 0 then call com_err_ (code, ME);
133    else call ioa_ ("Record written.");
134
135 bottom:
136    return;
137
138    end find;

```

SOURCE FILES USED IN THIS COMPILATION.

LINE	NUMBER	DATE MODIFIED	NAME	PATHNAME
	0	07/19/79 1109.5	put.pl1	>udd>F15D>Student_01>put.pl1

NAMES DECLARED IN THIS COMPILATION.

IDENTIFIER	OFFSET	LOC	STORAGE CLASS	DATA TYPE	ATTRIBUTES AND REFERENCES (* indicates a set context)
NAMES DECLARED BY DECLARE STATEMENT.					
ME		000100	automatic	varying char(4)	dcl 8 set ref 31* 34* 40* 46* 57* 59* 72* 75* 84* 91* 98* 106* 113* 121* 126* 131* 132*
arg			based	char	unaligned dcl 13 set ref 40* 43 46* 49 84* 87 91* 94 97 98* 102* 105 106* 109* 113* 116 121* 124* 126*
argl		000103	automatic	fixed bin(21,0)	dcl 10 set ref 38* 39 40 40 43 44* 45 46 46 49 82* 83 84 84 87 89* 90 91 91 94 96* 97 97 98 98 102 102 104* 105 106 106 109 109 111* 112 113 113 116 118* 121 121 124 124 126 126
argp		000116	automatic	pointer	dcl 14 set ref 38* 40 43 44* 46 49 82* 84 87 89* 91 94 96* 97 98 102 104* 105 106 109 111* 113 116 118* 121 124 126
code		000114	automatic	fixed bin(35,0)	dcl 12 set ref 38* 44* 56* 57 59 59* 82* 89* 96* 102* 104* 109* 111* 118* 124* 125 130* 131 132 132*
com_err_		000020	constant	entry	external dcl 25 ref 34 40 46 57 59 75 84 91 98 106 113 121 126 131 132
ou_\$arg_count		000016	constant	entry	external dcl 24 ref 32 73
cu_\$arg_ptr		000014	constant	entry	external dcl 23 ref 38 44 82 89 96 104 111 118
cv_dec_check		000030	constant	entry	external dcl 29 ref 102 109 124
emp_struc		000120	automatic	structure	level 1 unaligned dcl 15 set ref 56* 130*
error_table_\$bigarg		000010	external static	fixed bin(35,0)	dcl 20 set ref 40* 46* 84* 91*
error_table_\$wrong_no_of_args		000012	external static	fixed bin(35,0)	dcl 20 set ref 34* 75*
first_name		000104	automatic	char(15)	unaligned dcl 11 set ref 43* 56* 87* 130*
fixedoverflow		000132	stack reference	condition	dcl 21 ref 120
ioa		000022	constant	entry	external dcl 25 ref 60 133
last_name		000110	automatic	char(15)	unaligned dcl 11 set ref 49* 56* 94* 98* 106* 130*
mail_station	4(18)	000120	automatic	char(3)	level 2 packed unaligned dcl 15 set ref 54* 60 60* 116*
manager	2(09)	000120	automatic	picture(9)	level 2 packed unaligned dcl 15 set ref 53* 60 60* 109*
nargs		000102	automatic	fixed bin(17,0)	dcl 9 set ref 32* 33 73* 74
put_find_gate_\$find		000024	constant	entry	external dcl 26 ref 56
put_find_gate_\$put		000026	constant	entry	external dcl 26 ref 130
salary	6	000120	automatic	fixed dec(8,2)	level 2 dcl 15 set ref 53* 60 60* 124*
size		000140	stack reference	condition	dcl 21 ref 120
soc_sec_no		000120	automatic	picture(9)	level 2 packed unaligned dcl 15 set ref 53* 60 60* 102*
NAME DECLARED BY DECLARE STATEMENT AND NEVER REFERENCED.					
conversion		000000	stack reference	condition	dcl 21
NAMES DECLARED BY EXPLICIT CONTEXT.					
bottom		002025	constant	label	dcl 135 ref 122 127
find		000315	constant	entry	external dcl 4
put		000753	constant	entry	external dcl 70
NAME DECLARED BY CONTEXT OR IMPLICATION.					
verify				builtin function	ref 97 105



STORAGE REQUIREMENTS FOR THIS PROGRAM.

	Object	Text	Link	Symbol	Defs	Static
Start	0	0	2364	2416	2237	2374
Length	2610	2237	32	155	125	0

BLOCK NAME	STACK SIZE	TYPE	WHY NONQUICK/WHO SHARES STACK FRAME
find	226	external procedure	is an external procedure.
on unit on line 120	98	on unit	

STORAGE FOR AUTOMATIC VARIABLES.

STACK FRAME	LOC IDENTIFIER	BLOCK NAME
find	000100 ME	find
	000102 nargs	find
	000103 arg1	find
	000104 first_name	find
	000110 last_name	find
	000114 code	find
	000116 argp	find
	000120 emp_struct	find

THE FOLLOWING EXTERNAL OPERATORS ARE USED BY THIS PROGRAM.

r_e_as	r_ne_as	call_ext_out_desc	call_ext_out	return	tra_ext
enable	ext_entry	int_entry	unpack_pic		

THE FOLLOWING EXTERNAL ENTRIES ARE CALLED BY THIS PROGRAM.

com_err_	cu \$arg count	cu \$arg ptr	cv_dec_check_
ioa_	put_find_gate_\$find	put_find_gate_\$put	

THE FOLLOWING EXTERNAL VARIABLES ARE USED BY THIS PROGRAM.

error_table_\$bigarg	error_table_\$wrong_no_of_args
----------------------	--------------------------------

LINE	LOC	LINE	LOC	LINE	LOC	LINE	LOC	LINE	LOC	LINE	LOC	LINE	LOC
4	000314	31	000322	32	000326	33	000334	34	000337	36	000363	38	000364
39	000403	40	000406	41	000440	43	000441	44	000445	45	000464	46	000467
47	000521	49	000522	53	000526	54	000541	56	000544	57	000561	59	000612
60	000632	68	000751	70	000752	72	000760	73	000764	74	000773	75	000776
77	001022	82	001023	83	001042	84	001045	85	001076	87	001077	89	001103
90	001122	91	001125	92	001156	94	001157	96	001163	97	001202	98	001224
100	001263	102	001264	104	001315	105	001334	106	001356	107	001415	109	001416
111	001447	112	001466	113	001471	114	001524	116	001525	118	001531	120	001550
121	001564	122	001620	120	001623	124	001630	125	001661	126	001663	127	001717
130	001720	131	001735	132	001766	133	002006	135	002025				

---

put\_find\_gate\_.alm

arch comp in: >udd>F15D>s1>handout>gw.archive  
component updated: 03/18/80 1609.5

system id: 34-14.C

---

" Gate procedure written by Raphael Frommer of Marketing Education  
" to be used in conjunction with the "find" and "put" command  
" procedures. Transfers control to the pl1 procedure "put\_find\_gate\_util\_."  
"

```
include gate_macros
gate_info
gate find,put find_gate_util_,find,4
gate put,put_find_gate_util_,put,4
end
```

put\_find\_gate\_.list

arch comp in: >udd>F15D>s1>handout>gw.archive  
component updated: 03/18/80 1609.5

system id: 34-14.C

ASSEMBLY LISTING OF SEGMENT >udd>F15D>s1>put\_find\_gate\_.alm  
ASSEMBLED ON: 03/18/80 1557.0 mst Tue  
OPTIONS USED: list  
ASSEMBLED BY: ALM Version 6.2, June 1979  
ASSEMBLER CREATED: 01/28/80 0927.8 mst Mon

```
1 *      Gate procedure written by Raphael Frommer of Marketing Education
2 *      to be used in conjunction with the "find" and "put" command
3 *      procedures.  Transfers control to the pl1 procedure "put_find_gate_util_
4 *
5        include gate_macros
1-1
1-2 * BEGIN INCLUDE FILE ..... gate_macros.incl.alm
1-3 *   Last modified 6/77 by N. Morris, B. Greenberg, & T. VanVleck
1-4 *   Modified 741212 by PG to inhibit while computing virtual CPU time
1-5
1-6 * This file contains several macros used when generating gate segments.
1-7 * The major macros that are defined are:
1-8 *
1-9 *      gate_info          general setup code for normal gates
1-10 *     hardcore_gate_info general setup code for hardcore gates
1-11 *     bad_dir_handler    setup and handler for bad_dir condition (goes at bottom)
1-12 *     bad_dir_handler_entry entrypoint for above (goes at top)
1-13 *     gate                to define a normal gate
1-14 *     hgate              to define a hardcore gate
1-15 *     fgate              to define a fast hardcore gate
1-16 *
1-17
1-18
1-19 *      maclist off
1-20
1-21
```

000000

```

1-22 " HGATE - define a hardcore gate entry
1-23 "
1-24 "         hgate      gatename,procedure,entry,args{,bad_dir_trap)
1-25 "
1-26 " The endpoint gatename is defined in the gate segment.  If args
1-27 " is nonzero, the number of arguments passed to gatename must be
1-28 " equal to args.  When gatename is called, it will in turn call
1-29 " procedure$entry.
1-30 "
1-31 "
1-32 macro      hgate
1 1-33         gentry      &1,&4*2,&1.t
2 1-34         tsx2        .setup
3 1-35         &=&5,bad_dir_trap&[ tsx0        .set_dir_trap
4 1-36         &]         short_call &2&&3(ap|0)
5 1-37         eppbp       lp|&1.t
6 1-38         tra         .return_code
7 1-39         maclist     restore
8 1-40         use         linkage
9 1-41         even
10 1-42         maclist    on,save
11 1-43         &1.t:
12 1-44         bss         ,6
13 1-45
14 1-46         maclist    restore
15 1-47         &end
1-48
1-49
1-50 " FGATE - define a fast hardcore gate
1-51 "
1-52 "         fgate      gatename,procedure,entry
1-53 "
1-54 "
1-55 macro      fgate
1 1-56         gentry      &1,0,0
2 1-57         epplp      .my lp,*
3 1-58         tra         &2&&3
4 1-59
5 1-60         maclist     off
6 1-61         &end
1-62
1-63
1-64 " GATE - define a normal gate entry
1-65 "
1-66 "         gate       gatename,procedure,entry,args
1-67 "
1-68 "
1-69 macro      gate
1 1-70         gentry      &1,&4*2,0
2 1-71         tsx2        .setup
3 1-72         short_call &2&&3(ap|0)
4 1-73         return
5 1-74
6 1-75         madlist    restore

```

7 1-76 &end  
1-77  
1-78

```

1-79
1-80 macro gentry
1 1-81 maclist on,save
2 1-82 segdef &1
3 1-83 maclist restore
4 1-84 use transfer_vector
5 1-85 maclist on,save
6 1-86 &1:
7 1-87 tra &1.e
8 1-88 maclist restore
9 1-89 use main
10 1-90 maclist on,save
11 1-91 zero &2,&3
12 1-92 &1.e:
13 1-93 &end
1-94
1-95

```

```

1-96 " HARDCORE_GATE_INFO - general info for hardcore gates
1-97
1-98 macro    hardcore_gate_info
1 1-99      maclist  on,save
2 1-100     name     &1
3 1-101
4 1-102     include  stack_header
5 1-103
6 1-104     include  stack_frame
7 1-105
8 1-106     maclist  restore
9 1-107     eject
10 1-108
11 1-109    tempd   .temp
12 1-110    tempd   .label variable(0)
13 1-111    tempd   .time1,.time2
14 1-112    tempd   .unwinder_arglist(0)
15 1-113    tempd   .on unit(5)
16 1-114    temp    .pf,.entryp
17 1-115    tempd   .vfl arglist(2)
18 1-116    tempd   .moptr
19 1-117
20 1-118    use     transfer_vector
21 1-119    equ     .tv_begin,*
22 1-120
23 1-121    tra     .actor
24 1-122
25 1-123    use     tv_end
26 1-124
27 1-125    segdef  .tv_end
28 1-126    .tv_end: vfd   14/(#-.tv_begin)
29 1-127
30 1-128    use     main
31 1-129
32 1-130    segdef  .my_lp
33 1-131    even
34 1-132    .my_lp:  bss    ,2
35 1-133
36 1-134    join    /text/transfer_vector,tv_end,main
37 1-135
38 1-136    use     linkage
39 1-137    join    /link/linkage
40 1-138
41 1-139    use     main
42 1-140
43 1-141    maclist  restore
44 1-142    eject
45 1-143
46 1-144    .actor:  epplp   .my_lp,*
47 1-145    maclist  restore
48 1-146    gate_actor
49 1-147
50 1-148

```

```

51 1-149
52 1-150 .setup: push
53 1-151 epplp .my_lp,*
54 1-152 sprilp sp|stack_frame.lp_ptr
55 1-153
56 1-154 maclist restore
57 1-155 gcheck
58 1-156
59 1-157 inhibit on <+><+><+><+><+><+><+><+><+><+><+>
60 1-158 rocl sys_info$clock,* calculate times.
61 1-159 sbaq pds$cpu_time
62 1-160 staq .time1
63 1-161 sbaq pds$virtual_delta
64 1-162 staq .time2
65 1-163 lda pds$page_waits
66 1-164 sta .pf
67 1-165 inhibit off <-><-><-><-><-><-><-><-><-><-><->
68 1-166 tra 0,2
69 1-167
70 1-168 maclist restore
71 1-169 eject
72 1-170
73 1-171 .return_code:
74 1-172 inhibit on <+><+><+><+><+><+><+><+><+><+><+><+>
75 1-173 rocl sys_info$clock,* calculate times
76 1-174 sbaq pds$cpu_time
77 1-175 staq .temp
78 1-176 sbaq .time1
79 1-177 adaq bp|0
80 1-178 staq bp|0
81 1-179 ldaq .temp
82 1-180 sbaq pds$virtual_delta
83 1-181 sbaq .time2
84 1-182 adaq bp|2
85 1-183 staq bp|2
86 1-184 lda pds$page_waits
87 1-185 sbla .pf
88 1-186 asa bp|4
89 1-187 aos bp|5
90 1-188 inhibit off <-><-><-><-><-><-><-><-><-><-><->
91 1-189
92 1-190 return
93 1-191
94 1-192 maclist restore
95 1-193 eject
96 1-194 maclist restore
97 1-195 &end
1-196
1-197

```



```

1-198 " BAD_DIR_HANDLER - code to setup and handle bad_dir condition
1-199 " put this after the last hgate macro
1-200
1-201 macro bad_dir_handler
1 1-202 maclist on,save
2 1-203
3 1-204 include on_unit
4 1-205
5 1-206 use transfer_vector
6 1-207 .handler_entry:
7 1-208 tra .handler
8 1-209 .handler_restart_entry:
9 1-210 tra .handler_restart_point
10 1-211
11 1-212 use main
12 1-213 .set_dir_trap:
13 1-214 stx0 .entryp save for restart
14 1-215
15 1-216 mlr (),(pr),fill(000)
16 1-217 desc9a 0,0
17 1-218 desc9a .on_unit,10*4
18 1-219 eppbp .bad_dir_name
19 1-220 spribp .on_unit+on_unit.name
20 1-221 eppbp .handler_entry
21 1-222 spribp .on_unit+on_unit.body
22 1-223 lxll .bad_dir_desc
23 1-224 sxll .on_unit+on_unit.size
24 1-225 eaa .on_unit set up on-unit for bad_dir_
25 1-226 sbll sp|0,du .. make rel to sp
26 1-227 sta sp|stack_frame.on_unit_rel_ptrs
27 1-228 lda stack_frame.condition_bit,d1
28 1-229 orsa sp|stack_frame.flag_word
29 1-230 tra 0,0
30 1-231
31 1-232 string bad_dir_
32 1-233
33 1-234

```

```

34 1-235
35 1-236 .handler: epaq      sp|0      verify that call came from ring 0
36 1-237   cana      -1,d1     check ring number in AL
37 1-238   tze      #+2
38 1-239   zero     0          go way kid you bother me
39 1-240
40 1-241   push
41 1-242   ep|p      .my lp,*  " ok, we like the call
42 1-243   ld|x0     ap|0     get display
43 1-244   ep|bp     ap|2,0*
44 1-245   lda      bplstack_frame.prev_sp
45 1-246   cana     =o700000,d1 from another ring?
46 1-247   tze      .continue_signal if not, back to signal_
47 1-248   ep|ap     ap|2,*   Get mcptr
48 1-249   ep|ap     ap|0,*   ..
49 1-250   spr|ap    bpl.mcptr .. save in gate frame
50 1-251   spr|bp    .label_variable+2
51 1-252   ep|bp     .handler_restart_entry
52 1-253   spr|bp    .label_variable
53 1-254   ep|bp     .label_variable
54 1-255   spr|bp    .unwinder_arglist+2
55 1-256   fld      =1b24,d1
56 1-257   staq     .unwinder_arglist
57 1-258   call    unwinder_$unwinder_(.unwinder_arglist)
58 1-259
59 1-260 .continue_signal:
60 1-261   lda      =o400000,du  *1*b
61 1-262   sta      ap|10,*     set continue bit
62 1-263   return
63 1-264
64 1-265 .handler_restart_point:
65 1-266   epaq      sp|0      check that call came from ring 0
66 1-267   cana     -1,d1
67 1-268   tze      #+2
68 1-269   zero     1
69 1-270   ep|p      .my lp,*
70 1-271   lca      stack_frame.condition_bit+1,d1 Vanish on-unit
71 1-272   ansa     sp|stack_frame.flag_word
72 1-273   ep|bp     .mcptr
73 1-274   spr|bp    .vfl_arglist+2
74 1-275   fld      =1b24,d1
75 1-276   staq     .vfl_arglist
76 1-277   short_call verify_lock$verify_lock_bad_dir(.vfl_arglist)
77 1-278   ld|x0     .entryp
78 1-279   ep|ap     sp|stack_frame.arg_ptr,*
79 1-280   tra      0,0      retry the call
80 1-281
81 1-282   maclist  restore
82 1-283   eject
83 1-284   maclist  restore
84 1-285   &end
   1-286
   1-287

```

```

1-288 " GATE_INFO - general info for non-hardcore gates
1-289
1-290 macro gate_info
1 1-291 maclist on,save
2 1-292 use transfer_vector
3 1-293 tra .actor
4 1-294
5 1-295 use main
6 1-296 join /text/transfer_vector,main
7 1-297
8 1-298 maclist restore
9 1-299 eject
10 1-300
11 1-301 .actor: getlp
12 1-302 maclist restore
13 1-303 gate_actor
14 1-304
15 1-305 maclist restore
16 1-306 eject
17 1-307
18 1-308 .setup: push
19 1-309 getlp
20 1-310 maclist restore
21 1-311 gcheck
22 1-312 tra 0,2
23 1-313
24 1-314 maclist restore
25 1-315 eject
26 1-316
27 1-317 maclist restore
28 1-318 &end
1-319
1-320

```

```

1-321 " Macro to generate gate actor.
1-322
1-323 macro gate_actor
1 1-324 maclist on,save
2 1-325 eppbb ap|2,#
3 1-326 lda bp|-1 get length of string
4 1-327 tze .return_name zero length => get name
5 1-328
6 1-329 adla 1,d1 include length of acc
7 1-330 stz ap|4,#
8 1-331 tsx0 .search_defs
9 1-332
10 1-333 cmpc (pr,r1),(pr,r1) compare name
11 1-334 desc9a bp|-1(3),al
12 1-335 desc9a bb|0,al
13 1-336 tnz .next_def
14 1-337
15 1-338 lda ab|1,2 return location
16 1-339 arl 18
17 1-340 sta ap|4,#
18 1-341
19 1-342 short_return
20 1-343
21 1-344 .return_name:
22 1-345 lxl3 ap|4,# get location
23 1-346 tsx0 .search_defs
24 1-347
25 1-348 cmpx3 ab|1,2 compare location
26 1-349 tnz .next_def
27 1-350
28 1-351 lda bb|0 get length of name
29 1-352 arl 27
30 1-353 sta bp|-1 set length of varying string
31 1-354 mlr (pr,r1),(pr,r1) return string
32 1-355 desc9a bb|0(1),al
33 1-356 desc9a bp|0,al
34 1-357
35 1-358 short_return
36 1-359
37 1-360 .search_defs:
38 1-361 eax2 0
39 1-362 eppab lp|0,# ab -> defs
40 1-363 .defs_loop:
41 1-364 lxl1 ab|1,2 get class and flags
42 1-365 cmpx1 =o400000,du must be class 0
43 1-366 tnz .next_def
44 1-367
45 1-368 ldx7 ab|2,2
46 1-369 eppbb ab|0,7 bb -> name
47 1-370 tra 0,0 test definition
48 1-371 .next_def:
49 1-372 ldx2 ab|0,2 chain to next def
50 1-373 tnz .defs_loop
51 1-374

```

```
52 1-375      short_return
53 1-376
54 1-377  &end
    1-378
    1-379
```

```

1-380 " Miscellaneous macros.
1-381
1-382 macro      gcheck
1      1-383      maclist  on,save
2      1-384      ldx1     -2,2          get number of args expected
3      1-385      tze       .no_gate_error  if zero, none or doesn't matter
4      1-386      cmpx1    ap!0         compare against number given
5      1-387      tze       .no_gate_error  args match, call procedure
6      1-388
7      1-389      call     signal_$signal_(signal_arglist)
8      1-390      oct      0
9      1-391
10     1-392      even
11     1-393      signal_arglist:
12     1-394      zero     2,4
13     1-395      zero     2,0
14     1-396      arg      .gate_errorname
15     1-397      arg
16     1-398      arg      .gate_errordesc
17     1-399      arg
18     1-400
19     1-401      string   gate_error
20     1-402
21     1-403
22     1-404      .no_gate_error:
23     1-405      &end
1-406
1-407 macro      string
1      1-408      .&iname:   aci          "&1"
2      1-409
3      1-410      .&ldesc:  vfd         o9/525,o27/&11
4      1-411
5      1-412
6      1-413      &end
1-414
1-415 macro      eject
1      1-416      maclist  on,save
2      1-417

```

```
3 1-418 &end
1-419
1-420 " END INCLUDE FILE ..... gate_macros.incl.asm
1-421
6
gate_info
use transfer_vector
tra .actor

use main
join /text/transfer_vector,main

000000 0a 000004 7100 00
```

```

000004 aa 7 00046 2721 20
000005 aa 0 00002 3521 20
000006 aa 2 77777 2351 00
000007 0a 000023 6000 00

000010 aa 000001 0350 07
000011 aa 0 00004 4501 20
000012 0a 000036 7000 00

000013 aa 0 00140 1065 40
000014 aa 277777 600005
000015 aa 300000 000005
000016 0a 000046 6010 00

000017 aa 1 00001 2351 12
000020 aa 000022 7710 00
000021 aa 0 00004 7551 20

000022 aa 7 00044 7101 20

000023
000023 aa 0 00004 7231 20
000024 0a 000036 7000 00

000025 aa 1 00001 1031 12
000026 0a 000046 6010 00

000027 aa 3 00000 2351 00
000030 aa 000033 7710 00
000031 aa 2 77777 7551 00
000032 aa 0 00140 1005 40
000033 aa 300000 200005
000034 aa 200000 000005

000035 aa 7 00044 7101 20

000036
000036 aa 000000 6220 00
000037 aa 4 00000 3515 20
000040
000040 aa 1 00001 7211 12
000041 aa 400000 1010 03
000042 0a 000046 6010 00

000043 aa 1 00002 2271 12
000044 aa 1 00000 3535 17
000045 aa 000000 7100 10
000046
000046 aa 1 00000 2221 12
000047 0a 000040 6010 00

000050 aa 7 00044 7101 20

```

```

.actor:  getl  ap|2,*
         eppbp ap|2,*
         lda   bp|-1      get length of string
         tze   .return_name zero length => get name

         adla  1,d1      include length of acc
         stz   ap|4,*
         tsx0  .search_defs

         cmpc  (pr,r1),(pr,r1) compare name
         desc9a bp|-1(3),al
         desc9a bb|0,al
         tnz   .next_def

         lda   ab|1,2     return location
         arl   18
         sta   ap|4,*

         short_return

.return_name:
         lxl3  ap|4,*     get location
         tsx0  .search_defs

         cmpx3 ab|1,2     compare location
         tnz   .next_def

         lda   bb|0      get length of name
         arl   27
         sta   bp|-1     set length of varying string
         mlr   (pr,r1),(pr,r1) return string
         desc9a bb|0(1),al
         desc9a bp|0,al

         short_return

.search_defs:
         eax2  0
         eppab lp|0,*     ab -> defs

.defs_loop:
         lxl1  ab|1,2     get class and flags
         cmpx1 so400000,du must be class 0
         tnz   .next_def

         ldx7  ab|2,2
         eppbb ab|0,7     bb -> name
         tra   0,0       test definition

.next_def:
         ldx2  ab|0,2     chain to next def
         tnz   .defs_loop

         short_return

```





```

000051 aa 000060 6270 00
000052 aa 7 00040 2721 20
000053 aa 7 00046 2721 20
000054 aa 777776 2210 12
000055 0a 000102 6000 00
000056 aa 0 00000 1011 00
000057 0a 000102 6000 00

```

```

000060 aa 6 00000 2541 00
000061 0a 000070 3500 00
000062 4a 4 00010 3521 20
000063 aa 6 00040 7531 00
000064 aa 7 00036 6701 20
000065 aa 6 00000 1731 00
000066 aa 6 00040 0731 00
000067 aa 000000 000000

```

```

000070
000070 aa 000002 000004
000071 aa 000002 000000
000072 0a 000076 0000 00
000073 aa 000000 0000 00
000074 0a 000101 0000 00
000075 aa 000000 0000 00

```

```

000076
000076 aa 147 141 164 145
000077 aa 137 145 162 162
000100 aa 157 162 000 000
000101
000101 aa 525000 000010

```

```

000102
000102 aa 000000 7100 12

```

```

.setup: push
        getlp
        ldx1      -2,2          get number of args expected
        tze       .no_gate_error if zero, none or doesn't matter
        cmpx1    ap|0         compare against number given
        tze       .no_gate_error args match, call procedure

        call     signal_$signal_(signal_arglist)

        oct      0

        even
signal_arglist:
        zero     2,4
        zero     2,0
        arg      .gate_errortname
        arg      .gate_errordesc
        arg

        string   gate_error
.gate_errortname:
        aci      "gate_error"

.gate_errordesc:
        vfd      09/525,027/10

.no_gate_error:
        tra     0,2

```

```

000001
000001 0a 000104 7100 00
000103 aa 000010 000000
000104
000104 0a 000051 7020 00
000105 4a 4 00012 3521 20
000106 aa 7 00036 6701 20
000107 aa 6 00030 3701 20
000110 aa 7 00042 7101 20

```

```

000002
000002 0a 000112 7100 00
000111 aa 000010 000000
000112
000112 0a 000051 7020 00
000113 4a 4 00014 3521 20
000114 aa 7 00036 6701 20
000115 aa 6 00030 3701 20
000116 aa 7 00042 7101 20

```

NO LITERALS

```

7      gate    find,put_find_gate_util_,find,4
      segdef  find
find:   tra     find.e
      zero    4*2,0
find.e: tsx2     .setup
      short_call put_find_gate_util_$find(ap|0)

      return

8      gate    put,put_find_gate_util_,put,4
      segdef  put
put:    tra     put.e
      zero    4*2,0
put.e:  tsx2     .setup
      short_call put_find_gate_util_$put(ap|0)

      return

9      end

```

NAME DEFINITIONS FOR ENTRY POINTS AND SEGDEFS

```

000120 5a 000003 000000
000121 5a 000032 600000
000122 aa 000000 000000
000123 55 000012 000002
000124 5a 000002 400003
000125 55 000006 000012
000126 aa 016 160 165 164
000127 aa 137 146 151 156
000130 aa 144 137 147 141
000131 aa 164 145 137 000
000132 55 000016 000003
000133 0a 000002 400000
000134 55 000015 000003
000135 aa 003 160 165 164
000136 55 000023 000012
000137 0a 000001 400000
000140 55 000021 000003
000141 aa 004 146 151 156
000142 aa 144 000 000 000
000143 55 000002 000016
000144 6a 000000 400002
000145 55 000026 000003
000146 aa 014 163 171 155
000147 aa 142 157 154 137
000150 aa 164 141 142 154
000151 aa 145 000 000 000

```

put

find

symbol\_table

DEFINITIONS HASH TABLE

```

000152 aa 000000 000015
000153 aa 000000 000000
000154 aa 000000 000000
000155 5a 000012 000000
000156 aa 000000 000000
000157 aa 000000 000000
000160 aa 000000 000000
000161 5a 000016 000000
000162 5a 000023 000000
000163 aa 000000 000000
000164 aa 000000 000000
000165 aa 000000 000000
000166 aa 000000 000000
000167 aa 000000 000000

```

EXTERNAL NAMES

```

000170 aa 023 160 165 164
000171 aa 137 146 151 156
000172 aa 144 137 147 141
000173 aa 164 145 137 165
000174 aa 164 151 154 137
000175 aa 007 163 151 147
000176 aa 156 141 154 137

```

put\_find\_gate\_util\_

signal\_

NO TRAP POINTER WORDS

TYPE PAIR BLOCKS

000177	aa	000004	000000
000200	55	000050	000015
000201	aa	000004	000000
000202	55	000050	000021
000203	aa	000004	000000
000204	55	000055	000055
000205	aa	000001	000000
000206	aa	000000	000000

INTERNAL EXPRESSION WORDS

000207	5a	000057	000000
000210	5a	000061	000000
000211	5a	000063	000000

LINKAGE INFORMATION

000000	aa	000000	000000
000001	0a	000120	000000
000002	aa	000000	000000
000003	aa	000000	000000
000004	aa	000000	000000
000005	aa	000000	000000
000006	22	000010	000016
000007	a2	000000	000000
000010	9a	777770	0000 46
000011	5a	000071	0000 00
000012	9a	777766	0000 46
000013	5a	000070	0000 00
000014	9a	777764	0000 46
000015	5a	000067	0000 00

signal\_|signal\_

put\_find\_gate\_util\_|find

put\_find\_gate\_util\_|put

## SYMBOL INFORMATION

## SYMBOL TABLE HEADER

000000	aa	000000	000001
000001	aa	163171	155142
000002	aa	164162	145145
000003	aa	000000	000004
000004	aa	000000	106730
000005	aa	414344	011300
000006	aa	000000	107027
000007	aa	564756	217737
000010	aa	141154	155040
000011	aa	040040	040040
000012	aa	000024	000040
000013	aa	000034	000040
000014	aa	000044	000100
000015	aa	000002	000002
000016	aa	000064	000000
000017	aa	000000	000143
000020	aa	000000	000117
000021	aa	000000	000130
000022	aa	000134	000117
000023	aa	000064	000000
000024	aa	101114	115040
000025	aa	126145	162163
000026	aa	151157	156040
000027	aa	066056	062054
000030	aa	040112	165156
000031	aa	145040	061071
000032	aa	067071	040040
000033	aa	040040	040040
000034	aa	106162	157155
000035	aa	155145	162056
000036	aa	115105	104141
000037	aa	144155	151156
000040	aa	056141	040040
000041	aa	040040	040040
000042	aa	040040	040040
000043	aa	040040	040040
000044	aa	154151	163164
000045	aa	040040	040040
000046	aa	040040	040040
000047	aa	040040	040040
000050	aa	040040	040040
000051	aa	040040	040040
000052	aa	040040	040040
000053	aa	040040	040040
000054	aa	040040	040040
000055	aa	040040	040040
000056	aa	040040	040040
000057	aa	040040	040040
000060	aa	040040	040040
000061	aa	040040	040040
000062	aa	040040	040040

```
000063 aa 040040 040040
000064 aa 000000 000001
000065 aa 000000 000002
000066 aa 000076 000037
000067 aa 063453 431404
000070 aa 000000 107027
000071 aa 564153 400000
000072 aa 000106 000041
000073 aa 057660 372535
000074 aa 000000 105766
000075 aa 053324 600000
000076 aa 076165 144144
000077 aa 076106 061065
000100 aa 104076 163061
000101 aa 076160 165164
000102 aa 137146 151156
000103 aa 144137 147141
000104 aa 164145 137056
000105 aa 141154 155040
000106 aa 076154 144144
000107 aa 076151 156143
000110 aa 154165 144145
000111 aa 076147 141164
000112 aa 145137 155141
000113 aa 143162 157163
000114 aa 056151 156143
000115 aa 154056 141154
000116 aa 155040 040040
```

```
>udd>F15D>s1>put_find_gate_.alm
```

```
>ldd>include>gate_macros.incl.alm
```



MULTICS ASSEMBLY CROSS REFERENCE LISTING

Value	Symbol	Source file	Line number
4	.actor	put_find_gate_:	6.
40	.defs_loop	put_find_gate_:	6.
101	.gate_errordesc	put_find_gate_:	6.
76	.gate_errorname	put_find_gate_:	6.
46	.next_def	put_find_gate_:	6.
102	.no_gate_error	put_find_gate_:	6.
23	.return_name	put_find_gate_:	6.
36	.search_defs	put_find_gate_:	6.
51	.setup	put_find_gate_:	6, 7, 8.
1	find	put_find_gate_:	7.
104	find.e	put_find_gate_:	7.
4	main	put_find_gate_:	6, 7, 8.
2	put	put_find_gate_:	8.
112	put.e	put_find_gate_:	8.
	put_find_gate_util_	put_find_gate_:	7, 8.
	signal	put_find_gate_:	6.
70	signal_arglist	put_find_gate_:	6.
0	transfer_vector	put_find_gate_:	6, 7, 8.

NO FATAL ERRORS

## COMPILATION LISTING OF SEGMENT trans

Compiled by: Multics PL/I Compiler, Release 25c, of February 18, 1980

Compiled at: Honeywell LISD Phoenix, System M

Compiled on: 03/19/80 1043.7 mst Wed

Options: map

```
1 trans: proc;
2
3 /*
4 *
5 *
6 * Name: trans
7 *
8 * The trans command adds, counts, deletes, reads, and summarizes
9 * transactions resident in a message segment whose pathname is
10 * >udd>F15d>trans.ms. The amount of information it may return and the
11 * ability to add "transactions" into the segment are governed by
12 * extended access.
13 *
14 *
15 * Usage
16 *
17 * trans key {args}
18 *
19 * where:
20 *
21 * 1. key
22 * is one of the functions listed below.
23 *
24 * add, a
25 * adds a transaction to the message segment. Three
26 * additional arguments must be furnished:
27 * 1. part_name
28 * _16 characters or less
29 * 2. unit_price
30 * not to exceed 9999.99
31 * 3. how_many_sold
32 * not to exceed 99999
33 *
34 * The caller must have 'append' extended access to
35 * the segment trans.ms.
36 *
37 * count, c
38 * returns the number of transactions (messages) in
39 * the segment. The caller must have status extended
40 * access to the segment trans.ms.
41 *
42 * delete, d
43 * deletes the "current" transaction. A transaction
```

44\* is made current by some previous read operation.  
45\* If the last transaction read was not added by the  
46\* caller, the caller must have delete extended  
47\* access on trans.ms. Otherwise, the caller must at  
48\* least have 'own' extended access.  
49\*  
50\* read, r  
51\* reads one or more transactions. If an additional  
52\* argument is furnished, it must be one of the  
53\* following:  
54\*  
55\* 1. all  
56\* If user has 'read' extended access,  
57\* every transaction is dumped.  
58\*  
59\* If user has 'own', but not 'read', only  
60\* 'own' transactions are dumped.  
61\*  
62\* 2. first  
63\* If user has 'read' extended access, the  
64\* first transaction is dumped.  
65\*  
66\* If user has 'own', but not 'read', the  
67\* first 'own' transaction is dumped.  
68\*  
69\* 3. last  
70\* analogous to 'trans read first'  
71\*  
72\* 4. next  
73\* If user has 'read' extended access, the  
74\* next transaction in trans.ms is dumped.  
75\*  
76\* If user has 'own', but not 'read', the  
77\* caller's next transaction is dumped.  
78\*  
79\* 5. prior  
80\* analogous to 'trans read next'  
81\*  
82\* 6. <argument missing>  
83\* reads the current message.  
84\*  
85\* summary, s  
86\* If the caller has 'read' extended access, the  
87\* dollar amount grand total of all transactions in  
88\* the segment is returned. If the caller has 'own',  
89\* but not 'read', the dollars amount grand total of  
90\* 'own' transactions is returned instead.  
91\*  
92\* 2. args  
93\* is dependent upon key as documented above.  
94\*\*/  
95 /\*

```

*/
96 dol ME internal static options (constant) char (5) init ("trans");
97 dol TARGET_SEGMENT int static options (constant) char (8) init ("trans.ms");
98 dol TARGET_DIR int static options (constant) char (9) init (">udd>F15d");
99 dol (error_table $moderr, error_table $no_message, error_table $bigarg,
100 error_table $bad_arg, error_table $too_many_args, error_table $wrong_no_of_args)
101 ext static fixed bin (35);
102 dol areap ptr;
103 dol (argp, messagep, arg_list_ptr) ptr;
1 1 /* BEGIN INCLUDE FILE . . . mseg_return_args_v3.incl.pl1 */
1 2
1 3
1 4 /* structure returned when message is read from a message segment */
1 5
1 6
1 7 dol ms_arg_ptr ptr;
1 8
1 9 dol 1 mseg_return_args based (ms_arg_ptr) aligned,
1 10 2 ms_ptr ptr, /* pointer to message */
1 11 2 ms_len fixed bin (18), /* length of message in bits */
1 12 2 sender_id char (32), /* process-group ID of sender */
1 13 2 level fixed bin, /* validation level of sender */
1 14 2 ms_id bit (72), /* unique ID of message */
1 15 2 sender_authorization bit (72), /* access authorization of message sender */
1 16 2 access_class bit (72); /* message access class */
1 17
1 18
1 19 /* END INCLUDE FILE . . . mseg_return_args_v3.incl.pl1 */
104
105 dol (code, binary_number_sold) bin (35);
106 dol index fixed bin;
107
108 dol message_segment $open entry (char (*), char (*), fixed bin, fixed bin (35));
109 dol message_segment $close entry (fixed bin, fixed bin (35));
110 dol message_segment $get_message_count index entry (fixed bin, fixed bin, fixed bin (35));
111 dol message_segment $add_index entry (fixed bin, ptr, fixed bin (24), bit (72) aligned, bin (35));
112 dol message_segment $delete_index entry (fixed bin, bit (72) aligned, fixed bin (35));
113 dol message_segment $read_index entry (fixed bin, ptr, bit (1) aligned, ptr, fixed bin (35));
114 dol message_segment $incremental_read_index entry (fixed bin, ptr, bit (2) aligned, bit (72) aligned, ptr, fixed bin (35)
\o);
115 dol message_segment $own_read_index entry (fixed bin, ptr, bit (1) aligned, ptr, fixed bin (35));
\o(35));
116 dol message_segment $own_incremental_read_index entry (fixed bin, ptr, bit (2) aligned, bit (72) aligned, ptr, fixed bin
117
118 dol 1 trans_msg based (ms_arg_ptr -> mseg_return_args.ms_ptr), /* Sometimes pointed to by messagep */
119 2 widget_name char (16),
120 2 unit_price pic "$$$9.v99",
121 2 how_many_sold fixed dec (5, 0),
122 2 total_cost pic "(7)$9.v99";
123 dol grand_total_dollars internal static pic "(7)$9.v99";
124 dol silent bit (1) internal static init ("0b");
125 dol read_option (2:6) char (5) internal static options (constant)
126 init ("all", "next", "prior", "first", "last");
127 dol current_message_id bit (72) internal static aligned init ("0b");
128 dol looping_index fixed bin (17) init (1);

```

```

129 dcl own bit (1) init ("0"b);
130 dcl read_options internal static options (constant) char (121)
131 init ("Your read options are:
132     trans read
133     trans read all
134     trans read next
135     trans read prior
136     trans read first
137     trans read last");
138 dcl (com_err, ioa) entry options (variable);
139 dcl cu_$arg_count entry (fixed bin);
140 dcl cu_$arg_ptr entry (fixed bin, ptr, fixed bin(21), fixed bin (35));
141 dcl cu_$arg_list_ptr entry (ptr);
142 dcl cu_$arg_ptr_fcl entry (fixed bin, ptr, fixed bin(21), fixed bin (35), ptr);
143 dcl nargs fixed bin, arg1 fixed bin(21);
144 dcl arg char (arg1) based (argp);
145 dcl auto_area area (2048); /* Should be large enough for any and all allocated structures */
146 dcl string char (24);
147 dcl date_time entry (fixed bin (71), char (*));
148 dcl clock_entry returns (fixed bin (71));
149 dcl cv_dec_check_entry (char (*), fixed bin (35)) returns (fixed bin (35));
150 dcl conversion_condition;
151 dcl binary_time fixed bin (71);
152 dcl alternate_binary_time bit (72) aligned based;
153 dcl (size, addr) builtin;
154 dcl message count;
155 dcl iox_$control entry (ptr, char(*), ptr, fixed bin(35));
156 dcl iox_$user_io ext static ptr;
157
/*

```

```

158                                     /* COMMON BEGINNING POINT */
159
160     call ou_$arg_count (nargs);
161     if nargs = 0 then do;
162         call com_err_ (error_table $wrong_no_of_args, "trans",
163             ""/You have not invoked 'trans' properly.
164 Please type 'trans key [args]', where 'key' is either
165 'add', 'delete', 'read', 'summary', or 'count'.
166 'args' is a function of what key has been supplied.");
167         return;
168     end;
169     call ou_$arg_ptr (1, argp, argl, oode);
170     call ou_$arg_list_ptr (arg_list_ptr); /* Sets an automatic pointer */
171     if arg = "add" | arg = "a" then call trans_add;
172     else if arg = "delete" | arg = "d" then call trans_delete;
173     else if arg = "read" | arg = "r" then call trans_read;
174     else if arg = "summary" | arg = "s" then call trans_summary;
175     else if arg = "count" | arg = "c" then call trans_count;
176     else call com_err_ (error_table $bad_arg, ME, arg);
177     return;
178                                     /*

```

```

179 trans_read: proc;
180
181     if nargs > 2 then do;
182         call com_err_ (error_table_$too_many_args, ME, read_options);
183         return;
184     end;
185
186     /* Goal of the following do group is to set looping_index to
187      * 1,2,3,4,5, or 6 for a subsequent 'goto'. */
188     if nargs = 2 then do;
189         call cu $arg_ptr_rel (2, argp, arg1, code, arg_list_ptr);
190         do looping_index = 2 to 6 while (arg ^= read_option(looping_index));
191         end;
192
193         if looping_index = 7 then do;
194             call com_err_ (error_table_$bad_arg, ME, read_options);
195             return;
196         end;
197     end;
198
199     call message_segment_$open (TARGET_DIR, TARGET_SEGMENT, index, code);
200     if code ^= 0 then do;
201         call com_err_ (code, ME, "While attempting to open ^a^a.", TARGET_DIR, TARGET_SEGMENT);
202         return;
203     end;
204
205     areap = addr (auto_area);
206     allocate mseg_return_args in (auto_area) set (ms_arg_ptr);
207     /* No need to worry about freeing above, because area itself is automatic */
208     goto read_label (looping_index);
209
210 read_label (1):
211                                     /* 'trans read' */
212     if current_message_id = "0"b then do;
213         call com_err_ (0, ME, "There is no current message.");
214         return;
215     end;
216     else do;
217
218         call message_segment_$incremental_read_index (index, areap, "00"b,
219             current_message_id, ms_arg_ptr, code); /* User must have 'r' extended access */
220         if code = error_table_$moderr then call message_segment_$own_incremental_read_index (index,
221             areap, "00"b, current_message_id, ms_arg_ptr, code);
222         if code ^= 0 then do;
223             call com_err_ (code, ME);
224             call message_segment_$close (index, code);
225             return;
226         end;
227
228         call print_message ();
229
230         call message_segment_$close (index, code);
231         return;
232     end;

```

```

233 read_label (2):
234                                     /* 'trans read all' */
235 call message_segment_read_index (index, areap, "0"b /* from the first */,
236 ms_arg_ptr, code);
237 if code = error_table$moderr then do;
238 call message_segment_own_read_index
239 (index, areap, "0"b, ms_arg_ptr, code);
240 if code = error_table$moderr then do;
241 call com_err (code, ME);
242 call message_segment_close (index, code);
243 return;
244 end;
245 own = "1"b;
246 end;
247 if code = error_table$no_message then do;
248 call com_err (code, ME, "[You have no^;No] messages in ^a^a.", own, TARGET_DIR, TARGET_SEGMENT);
249 call message_segment_close (index, code);
250 return;
251 end;
252
253 if own then do;
254
255 do while (code ^= error_table$no_message);
256
257 call print_message ();
258 current_message_id = ms_arg_ptr -> mseg_return_args.ms_id;
259 call message_segment_own_incremental_read_index (index, areap, "01"b /* next */,
260 current_message_id, ms_arg_ptr, code);
261 end;
262
263 end;
264 else do;
265
266 do while (code ^= error_table$no_message);
267
268 call print_message ();
269 current_message_id = ms_arg_ptr -> mseg_return_args.ms_id;
270 call message_segment_incremental_read_index (index, areap, "01"b /* next */,
271 current_message_id, ms_arg_ptr, code);
272 end;
273 end;
274
275 call message_segment_close (index, code);
276 return;
277 read_label (3):
278                                     /* 'trans read next' */
279
280
281 if current_message_id = "0"b then do;
282 call com_err (0, ME, "There is no current message.");
283 return;
284 end;
285 else do;
286
287 call message_segment_incremental_read_index (index, areap, "01"b,

```



```

288     current_message_id, ms_arg_ptr, code);
289 if code = error_table $moderr then call message_segment $own_incremental_read_index (index,
290 areap, "01"b, current_message_id, ms_arg_ptr, code);
291 if code = error_table $moderr then do;
292     call com_err (code, ME);
293     call message_segment $close (index, code);
294     return;
295 end;
296 if code = error_table $no_message then do;
297     call com_err (code, ME, "There is no next message.");
298     call message_segment $close (index, code);
299     return;
300 end;
301
302     call print_message ();
303
304     current_message_id = ms_arg_ptr -> ms_id; /* Update what is now the current msg */
305     call message_segment $close (index, code);
306     return;
307 end;
308 read_label (4):
309
310                                     /* 'trans read prior' */
311
312 if current_message_id = "0"b then do;
313     call com_err (0, ME, "There is no current message.");
314     return;
315 end;
316 else do;
317     call message_segment $incremental_read_index (index, areap, "10"b,
318     current_message_id, ms_arg_ptr, code);
319 if code = error_table $moderr then call message_segment $own_incremental_read_index (index,
320 areap, "10"b, current_message_id, ms_arg_ptr, code);
321 if code = error_table $moderr then do;
322     call com_err (code, ME);
323     call message_segment $close (index, code);
324     return;
325 end;
326 if code = error_table $no_message then do;
327     call com_err (code, ME, "There is no prior message.");
328     call message_segment $close (index, code);
329     return;
330 end;
331
332     call print_message ();
333
334     current_message_id = ms_arg_ptr -> mseg_return_args.ms_id; /* Update what is now the current msg */
335     call message_segment $close (index, code);
336     return;
337 end;
338 read_label (5):
339
340                                     /* 'trans read first' */
341
342     call message_segment $read_index (index, areap, "0"b /* the first */,
ms_arg_ptr, code);

```

```

343     if code = error_table $moderr then do;
344         call message_segment $own_read_index (index, areap, "0"b /* the first */,
345             ms_arg_ptr, code);
346         if code = error_table $moderr then do;
347             call com_err_ (code, ME);
348             call message_segment $close (index, code);
349             return;
350         end;
351         own = "1"b;
352     end;
353     if code = error_table $no_message then do;
354         call com_err_ (code, ME, "[You have no^;No] messages in ^a>^a.", own, TARGET_DIR, TARGET_SEGMENT);
355         call message_segment $close (index, code);
356         return;
357     end;
358
359     call print_message ();
360     current_message_id = ms_arg_ptr -> mseg_return_args.ms_id; /* For a possible
361 *                                     'trans read' in the future */
362     call message_segment $close (index, code);
363     return;
364 read_label (6):
365                                     /* 'trans read last' */
366
367     call message_segment $read_index (index, areap, "1"b /* the last */,
368         ms_arg_ptr, code);
369     if code = error_table $moderr then do;
370         call message_segment $own_read_index (index, areap, "1"b /* the last */,
371             ms_arg_ptr, code);
372         own = "1"b;
373     end;
374     if code = error_table $no_message then do;
375         call com_err_ (code, ME, "[You have no^;No] messages in ^a>^a.", own, TARGET_DIR, TARGET_SEGMENT);
376         call message_segment $close (index, code);
377         return;
378     end;
379
380     call print_message ();
381     current_message_id = ms_arg_ptr -> mseg_return_args.ms_id; /* For a possible
382 *                                     'trans read' in the future */
383     call message_segment $close (index, code);
384     return;
385 end trans_read;
386                                     /*

```

```

*/
387 trans_add: proc;
388     if nargs ^= 4 then do;
389         call com_err_ (error_table $wrong_no_of_args, ME, "Proper invocation:
390 trans add widget_name unit_price how_many_sold");
391         return;
392     end;
393                                     /* Assign the elements in the structure trans_msg one by one
\c*/
394 allocate trans_msg in (auto_area) set (messagep);
395 call cu $arg_ptr_rel (2, argp, argl, code, arg_list_ptr);
396 if argl > 16 then do;
397     call com_err_ (error_table $bigarg, ME, "Widget_name mustn't exceed 16 characters.");
398     return;
399 end;
400 else messagep -> trans_msg.widget_name = arg;
401
402 call cu $arg_ptr_rel (3, argp, argl, code, arg_list_ptr);
403 on conversion begin;
404     call com_err_ (error_table $bad_arg, ME, "Unit price specified does not look something like 9999.99");
405     goto add_return_point;
406 end;
407 messagep -> trans_msg.unit_price = arg;
408 revert conversion; /* If we get here, conversion went fine */
409 call cu $arg_ptr_rel (4, argp, argl, code, arg_list_ptr);
410 binary_number_sold = cv_dec_check_ (arg, code);
411 if code ^= 0 then do;
412     call com_err_ (error_table $bad_arg, ME, "Number sold could not be converted into an integer.");
413     return;
414 end;
415 else if binary_number_sold > 99999 then do;
416     call com_err_ (error_table $bad_arg, ME, "Number sold exceeds 99999.");
417     return;
418 end;
419 else messagep -> trans_msg.how_many_sold = binary_number_sold;
420
421 messagep -> trans_msg.total_cost = messagep -> trans_msg.unit_price *
422     messagep -> trans_msg.how_many_sold; /* trans_msg structure ready for addition to >udd>F15d>trans.
\cms */
424
425 call message_segment $open (TARGET_DIR, TARGET_SEGMENT, index, code);
426 if code ^= 0 then do;
427     call com_err_ (code, ME, "While attempting to open ^a>^a.", TARGET_DIR, TARGET_SEGMENT);
428     return;
429 end;
430 binary_time = clock (); /* Need a unique id for message id, hence clock */
431 call message_segment $add_index (index, messagep, 36 * size (messagep -> trans_msg),
432     addr (binary_time) -> alternate_binary_time, code);
433 if code ^= 0 then call com_err_ (code, ME, "While attempting to add a transaction.");
434 else do;
435     current_message_id = addr (binary_time) -> alternate_binary_time;
436     call date_time_ ((current_message_id), string);
437     call ioa_ ("Transaction added at "a", string);
438 end;

```

```
439         call message_segment_$close (index, code);
440 add_return_point:
441         end trans_add;
442
```

/\*

```

*/
443 trans_delete: proc;
444
445     if nargs ^= 1 then do;
446         call com_err_ (error_table_$wrong_no_of_args, ME, "Proper invocation:
447 trans delete ");
448         return;
449     end;
450     if current_message_id = "0"b then do;
451         call com_err_ (0, ME, "There is no current transaction to delete.
452 Try some kind of read first.");
453         return;
454     end;
455     else do;
456         call message_segment_$open (TARGET_DIR, TARGET_SEGMENT, index, code);
457         if code ^= 0 then do;
458             call com_err_ (code, ME, "While attempting to open ^a>^a.", TARGET_DIR, TARGET_SEGMENT);
459             return;
460         end;
461         call message_segment_$delete_index (index, current_message_id, code);
462         if code = 0 then do;
463             current_message_id = "0"b;
464             call ioa_ ("Transaction deleted.");
465         end;
466         else call com_err_ (code, ME, "While attempting to delete the current message.");
467         call message_segment_$close (index, code);
468     end;
469 end trans_delete;
470
/*

```

```

*/
471 trans_summary: proc;
472
473 decl trans entry options (variable);
474
475     if nargs ^= 1 then do;
476         call com_err_ (error_table_$wrong_no_of_args, ME, "Proper invocation:
477 trans summary");
478         return;
479     end;
480     /* We can not tolerate a quit, so ... */
481     call iox_$control (iox_$user_io, "quit_disable", null (), code);
482     silent = "1"b;
483     grand_total_dollars = 0;
484
485     call trans ("read", "all");
486
487     call ioa_ ("Grand total = ^a", grand_total_dollars);
488
489     silent = "0"b;
490     grand_total_dollars = 0;
491     /* Now safe to unmask quits */
492     call iox_$control (iox_$user_io, "quit_enable", null (), code);
493
494 end trans_summary;
495
/*

```

```

*/
496 trans_count: proc;
497
498 /* User must have status extended access on trans.ms */
499
500     if nargs ^= 1 then do;
501         call com_err_ (error_table_$wrong_no_of_args, ME, "Proper invocation:
502 trans count");
503         return;
504     end;
505
506     call message_segment_$open (TARGET_DIR, TARGET_SEGMENT, index, code);
507     if code ^= 0 then do;
508         call com_err_ (code, ME, "While attempting to open ^a>^a.", TARGET_DIR, TARGET_SEGMENT);
509         return;
510     end;
511
512     call message_segment_$get_message_count_index (index, message_count, code);
513     if code = 0 then call ioa_ ("There are ^d messages in ^a", message_count, TARGET_SEGMENT);
514     else call com_err_ (code, ME);
515
516     call message_segment_$close (index, code);
517
518 end trans_count;
519
*/

```

```
*/
520 print_message: proc ();
521
522         /* Can tally dollars OR print transaction info */
523
524         call date_time_ ((ms_arg_ptr -> ms_id), string);
525         if ^silent then call ioa_ ("At ^a, ^a sold ^d ^a, ^/ at a total cost of ^a, (unit price = ^a)",
526             string, ms_arg_ptr -> mseg_return_args.sender_id, how_many_sold, widget_name, total_cost, unit_price);
527         else grand_total_dollars = grand_total_dollars + total_cost;
528
529         end print_message;
530
531 bottom_of_trans:
532
533         end trans;
```



SOURCE FILES USED IN THIS COMPILATION.

LINE	NUMBER	DATE MODIFIED	NAME	PATHNAME
	0	03/11/80 1122.9	trans.pl1	>user_dir_dir>F15D>Student_01>trans.pl1
104	1	08/03/77 1651.9	mseg_return_args_v3.incl.pl1	>ldd>include>mseg_return_args_v3.incl.pl1

NAMES DECLARED IN THIS COMPILATION.

IDENTIFIER	OFFSET	LOC	STORAGE CLASS	DATA TYPE	ATTRIBUTES AND REFERENCES (* indicates a set context)
NAMES DECLARED BY DECLARE STATEMENT.					
ME	000062		constant	char(5)	initial unaligned dcl 96 set ref 176* 182* 194* 201* 213* 223* 241* 248* 282* 292* 297* 312* 322* 327* 347* 354* 375* 389* 397* 404* 412* 416* 427* 433* 446* 451* 458* 466* 476* 501* 508* 514*
TARGET_DIR	000054		constant	char(9)	initial unaligned dcl 98 set ref 199* 201* 248* 354* 375* 425* 427* 456* 458* 506* 508*
TARGET_SEGMENT	000060		constant	char(8)	initial unaligned dcl 97 set ref 199* 201* 248* 354* 375* 425* 427* 456* 458* 506* 508* 513*
addr				builtin function	dcl 153 ref 205 431 435
alternate_binary_time			based	bit(72)	dcl 152 set ref 431* 435
areap	000100		automatic	pointer	dcl 102 set ref 205* 218* 220* 233* 238* 259* 270* 287* 289* 317* 319* 338* 344* 364* 370*
arg			based	char	unaligned dcl 144 set ref 171 171 172 172 173 173 174 174 175 175 176* 190 400 407 410*
arg_list_ptr	000106		automatic	pointer	dcl 103 set ref 170* 189* 395* 402* 409*
arg1	000120		automatic	fixed bin(21,0)	dcl 143 set ref 169* 171 171 172 172 173 173 174 174 175 175 176 176 189* 190 395* 396 400 402* 407 409* 410 410
argp	000102		automatic	pointer	dcl 103 set ref 169* 171 171 172 172 173 173 174 174 175 175 176 189* 190 395* 400 402* 407 409* 410
auto_area	000122		automatic	area(2048)	dcl 145 set ref 145* 205 206 394
binary_number_sold	000113		automatic	fixed bin(35,0)	dcl 105 set ref 410* 415 419
binary_time	004130		automatic	fixed bin(71,0)	dcl 151 set ref 430* 431 435
clock_	000072		constant	entry	external dcl 148 ref 430
code	000112		automatic	fixed bin(35,0)	dcl 105 set ref 169* 189* 199* 200 201* 218* 220 220* 222 223* 224* 230* 233* 237 238* 240 241* 242* 247 248* 249* 255 259* 266 270* 276* 287* 289 289* 291 292* 293* 296 297* 298* 305* 317* 319 319* 321 322* 323* 326 327* 328* 335* 338* 343 344* 346 347* 348* 353 354* 355* 362* 364* 369 370* 374 375* 376* 383* 395* 402* 409* 410* 411 425* 426 427* 431* 433 433* 439* 456* 457 458* 461* 462 466* 467* 481* 492* 506* 507 508* 512* 513 514* 516*
com_err_	000054		constant	entry	external dcl 138 ref 162 176 182 194 201 213 223 241 248 282 292 297 312 322 327 347 354 375 389 397 404 412 416 427 433 446 451 458 466 476 501 508 514
conversion	000000		stack reference	condition	dcl 150 ref 403 408
cu_\$arg_count	000060		constant	entry	external dcl 139 ref 160
cu_\$arg_list_ptr	000064		constant	entry	external dcl 141 ref 170
cu_\$arg_ptr	000062		constant	entry	external dcl 140 ref 169
cu_\$arg_ptr_rel	000066		constant	entry	external dcl 142 ref 189 395 402 409
current_message_id	000014		internal static	bit(72)	initial dcl 127 set ref 210 218* 220* 258* 259* 269* 270* 278 287* 289* 304* 308 317* 319* 334* 360* 381* 435* 436 450 461* 463*
cv_dec_check_	000074		constant	entry	external dcl 149 ref 410
date_time	000070		constant	entry	external dcl 147 ref 436 524
error_table_\$bad_arg	000024		external static	fixed bin(35,0)	dcl 99 set ref 176* 194* 404* 412* 416*

error_table \$bigarg	000022	external static	fixed bin(35,0)	dcl 99 set ref 397*
error_table \$moderr	000016	external static	fixed bin(35,0)	dcl 99 ref 220 237 240 289 291 319 321 343 346 369
error_table \$no_message	000020	external static	fixed bin(35,0)	dcl 99 ref 247 255 266 296 326 353 374
error_table \$too_many_args	000026	external static	fixed bin(35,0)	dcl 99 set ref 182*
error_table \$wrong_no_of_args	000030	external static	fixed bin(35,0)	dcl 99 set ref 162* 389* 446* 476* 501*
grand_total dollars	000010	internal static	picture(11)	unaligned dcl 123 set ref 483* 487* 490* 527* 527
how_many_sold	6	based	fixed dec(5,0)	level 2 dcl 118 set ref 419* 421 525*
index	000114	automatic	fixed bin(17,0)	dcl 106 set ref 199* 218* 220* 224* 230* 233* 238* 242* 249* 259* 270* 276* 287* 289* 293* 298* 305* 317* 319* 323* 328* 335* 338* 344* 348* 355* 362* 364* 370* 376* 383* 425* 431* 439* 456* 461* 467* 506* 512* 516*
ioa	000056	constant	entry	external dcl 138 ref 437 464 487 513 525
iox \$control	000076	constant	entry	external dcl 155 ref 481 492
iox \$user io	000100	external static	pointer	dcl 156 set ref 481* 492*
looping_index	000115	automatic	fixed bin(17,0)	initial dcl 128 set ref 128* 190* 190* 193 208
message_count	000132	automatic	fixed bin(17,0)	dcl 154 set ref 512* 513*
message_segment \$add_index	000040	constant	entry	external dcl 111 ref 431
message_segment \$close	000034	constant	entry	external dcl 109 ref 224 230 242 249 276 293 298 305 323 328 335 348 355 362 376 383 439 467 516
message_segment \$delete_index	000042	constant	entry	external dcl 112 ref 461
message_segment \$get_message_count_index	000036	constant	entry	external dcl 110 ref 512
message_segment \$incremental_read_index	000046	constant	entry	external dcl 114 ref 218 270 287 317
message_segment \$open	000032	constant	entry	external dcl 108 ref 199 425 456 506
message_segment \$own_incremental_read_index	000052	constant	entry	external dcl 116 ref 220 259 289 319
message_segment \$own_read_index	000050	constant	entry	external dcl 115 ref 238 344 370
message_segment \$read_index	000044	constant	entry	external dcl 113 ref 233 338 364
messagep	000104	automatic	pointer	dcl 103 set ref 394* 400 407 419 421 421 421 431* 431
ms_arg_ptr	000110	automatic	pointer	dcl 1-7 set ref 206* 218* 220* 233* 238* 258 259* 269 270* 287* 289* 304 317* 319* 334 338* 344* 360 364* 370* 381 524 525 525 525 525 527
ms_id	14	based	bit(72)	level 2 dcl 1-9 ref 258 269 304 334 360 381 524
ms_ptr		based	pointer	level 2 dcl 1-9 ref 525 525 525 525 527
msg_return_args		based	structure	level 1 dcl 1-9 set ref 206
nargs	000117	automatic	fixed bin(17,0)	dcl 143 set ref 160* 161 181 188 388 445 475 500
own	000116	automatic	bit(1)	initial unaligned dcl 129 set ref 129* 245* 248* 253 351* 354* 372* 375*
read_option	000045	constant	char(5)	initial array unaligned dcl 125 ref 190
read_options	000006	constant	char(121)	initial unaligned dcl 130 set ref 182* 194*
sender_id	3	based	char(32)	level 2 dcl 1-9 set ref 525*
silent	000013	internal static	bit(1)	initial unaligned dcl 124 set ref 482* 489* 525
size			builtin function	dcl 153 ref 431
string	000122	automatic	char(24)	unaligned dcl 146 set ref 436* 437* 524* 525*
total_cost	10	based	picture(11)	level 2 packed unaligned dcl 118 set ref 421* 525* 527
trans	000102	constant	entry	external dcl 473 ref 485
trans_msg		based	structure	level 1 unaligned dcl 118 set ref 394 431
unit_price	4	based	picture(8)	level 2 packed unaligned dcl 118 set ref 407* 421 525*
widget_name		based	char(16)	level 2 packed unaligned dcl 118 set ref 400* 525*

NAMES DECLARED BY EXPLICIT CONTEXT.

add_return_point	003465	constant	label	dcl 440 ref 405
bottom_of_trans	001025	constant	label	dcl 531
print_message	004330	constant	entry	internal dcl 520 ref 228 257 268 302 332 359 380
read_label	000000	constant	label	array(6) dcl 210 ref 208
trans	000570	constant	entry	external dcl 1
trans_add	002644	constant	entry	internal dcl 387 ref 171
trans_count	004122	constant	entry	internal dcl 496 ref 175
trans_delete	003466	constant	entry	internal dcl 443 ref 172
trans_read	001026	constant	entry	internal dcl 179 ref 173
trans_summary	003727	constant	entry	internal dcl 471 ref 174

NAMES DECLARED BY CONTEXT OR IMPLICATION.

empty			builtin function	ref 145
null			builtin function	ref 481 481 492 492

STORAGE REQUIREMENTS FOR THIS PROGRAM.

	Object	Text	Link	Symbol	Defs	Static
Start	0	0	5024	5130	4501	5034
Length	5462	4501	104	315	322	6

BLOCK NAME	STACK SIZE	TYPE	WHY NONQUICK/WHO SHARES STACK FRAME
trans	2555	external procedure	is an external procedure.
trans_read		internal procedure	shares stack frame of external procedure trans.
trans_add	262	internal procedure	enables or reverts conditions.
on unit on line 403	94	on unit	
trans_delete		internal procedure	shares stack frame of external procedure trans.
trans_summary		internal procedure	shares stack frame of external procedure trans.
trans_count		internal procedure	shares stack frame of external procedure trans.
print_message		internal procedure	shares stack frame of external procedure trans.

STORAGE FOR INTERNAL STATIC VARIABLES.

LOC IDENTIFIER	BLOCK NAME
000010 grand_total_dollars	trans
000013 silent	trans
000014 current_message_id	trans

STORAGE FOR AUTOMATIC VARIABLES.

STACK FRAME	LOC IDENTIFIER	BLOCK NAME
trans	000100 areap	trans
	000102 argp	trans
	000104 messagep	trans
	000106 arg_list_ptr	trans
	000110 ms_arg_ptr	trans
	000112 code	trans
	000113 binary_number_sold	trans
	000114 index	trans
	000115 looping_index	trans
	000116 own	trans
	000117 nargs	trans
	000120 argl	trans
	000122 auto_area	trans

```

004122 string          trans
004130 binary_time    trans
004132 message_count  trans

```

THE FOLLOWING EXTERNAL OPERATORS ARE USED BY THIS PROGRAM.

```

r_e_as      call_ext_out_desc  call_ext_out  call_int_this  return  tra_ext
enable      ext_entry          int_entry     any_to_any_tr  unpack_pic  alloc_based
empty

```

THE FOLLOWING EXTERNAL ENTRIES ARE CALLED BY THIS PROGRAM.

```

clock      com_err          cu $arg_count  cu $arg_list_ptr
cu $arg_ptr  cu $arg_ptr_rel  cv_dec_check  date_time
ioa        iox $control  message_segment_$add_index  message_segment_$close
message_segment_$delete_index  message_segment_$get_message_count_index
message_segment_$incremental_read_index  message_segment_$open
message_segment_$own_incremental_read_index  message_segment_$own_read_index
message_segment_$read_index  trans

```

THE FOLLOWING EXTERNAL VARIABLES ARE USED BY THIS PROGRAM.

```

error_table $bad_arg  error_table $bigarg  error_table $moderr  error_table $no_message
error_table $too_many_args  error_table $wrong_no_of_args  iox $user_io

```

WARNING 235 ON LINE 407

"arg" has been converted from a string value to an arithmetic value.

WARNING 235 ON LINE 436

"current\_message\_id" has been converted from a string value to an arithmetic value.

WARNING 235 ON LINE 524

"ms\_id" has been converted from a string value to an arithmetic value.

LINE	LOC	LINE	LOC	LINE	LOC	LINE	LOC	LINE	LOC	LINE	LOC	LINE	LOC
1	000567	128	000575	129	000577	145	000601	160	000604	161	000613	162	000615
167	000643	169	000644	170	000663	171	000672	172	000715	173	000732	174	000747
175	000764	176	001001	177	001024	531	001025	179	001026	181	001027	182	001032
183	001053	188	001054	189	001055	190	001076	191	001112	193	001114	194	001117
195	001140	199	001141	200	001166	201	001170	202	001224	205	001225	206	001227
208	001234	210	001236	213	001241	214	001265	218	001266	220	001310	222	001336
223	001340	224	001355	225	001366	228	001367	230	001370	231	001401	233	001402
237	001423	238	001427	240	001447	241	001453	242	001467	243	001500	245	001501
247	001503	248	001505	249	001544	250	001555	253	001556	255	001561	257	001565
258	001566	259	001573	261	001615	263	001616	266	001617	268	001623	269	001624
270	001631	272	001653	276	001654	277	001664	278	001665	282	001670	283	001714
287	001715	289	001737	291	001765	292	001771	293	002005	294	002016	296	002017
297	002021	298	002044	299	002055	302	002056	304	002057	305	002064	306	002074
308	002075	312	002100	313	002124	317	002125	319	002147	321	002175	322	002201
323	002215	324	002226	326	002227	327	002231	328	002254	329	002265	332	002266
334	002267	335	002274	336	002304	338	002305	343	002326	344	002332	346	002352
347	002356	348	002372	349	002403	351	002404	353	002406	354	002410	355	002447
356	002460	359	002461	360	002462	362	002467	363	002477	364	002500	369	002521
370	002525	372	002545	374	002547	375	002553	376	002612	377	002623	380	002624
381	002625	383	002632	384	002642	387	002643	388	002651	389	002655	391	002700

394 002701	395 002710	396 002731	397 002735	398 002761	400 002762	402 002767
403 003010	404 003024	405 003047	407 003052	408 003071	409 003072	410 003113
411 003137	412 003142	413 003166	415 003167	416 003172	417 003216	419 003217
421 003223	425 003244	426 003271	427 003274	428 003330	430 003331	431 003340
433 003362	435 003412	436 003415	437 003432	439 003453	440 003465	443 003466
445 003467	446 003472	448 003516	450 003517	451 003522	453 003546	456 003547
457 003573	458 003575	459 003631	461 003632	462 003645	463 003647	464 003652
465 003670	466 003671	467 003715	469 003726	471 003727	475 003730	476 003733
478 003757	481 003760	482 004012	483 004015	485 004021	487 004041	489 004062
490 004064	492 004070	494 004121	496 004122	500 004123	501 004126	503 004152
506 004153	507 004200	508 004202	509 004236	512 004237	513 004252	514 004301
516 004316	518 004327	520 004330	524 004331	525 004352	527 004425	529 004453

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