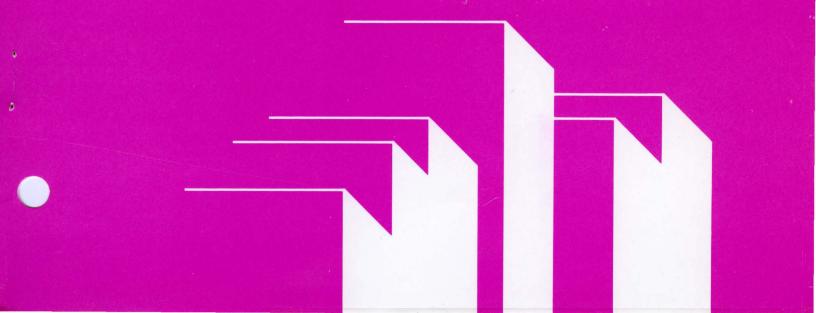
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	IBM Assists for MVS/XA
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Publication Number SA22-7092-0	File Number S370-01

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PREFACE

This publication is intended for system programmers and IBM Field Engineering personnel. The reader should be familiar with the general machine functions of 370-XA, as described in the IBM 370-XA Principles of Operation, SA22-7085, and with the MVS/SP Version 2 Licensed Program, referred to in this publication as the MVS/XA control program. The standard names for MVS/XA control blocks are used throughout the publication. The formats of these control blocks are described in the MVS/XA Debugging Handbook, Volumes 1 through 5, LC28-1164 through LC28-1168, respectively.

First Edition (March 1983)

Changes are periodically made to the information herein; before using this publication in connection with the operation of IBM equipment, refer to the latest IBM System/370 and 4300 Processors Bibliography, GC20-0001, for the editions that are applicable and current.

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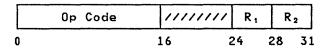
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publication describes instructions that depend on control blocks whose formats and relationships normally are established only by the These control MVS/XA program. instructions improve performance frequently used parts of the MVS/XA control program on machines which the 370-XA archi/tectural mode provide and these instructions. Four of the instructions assist the handling of control-program locks. A locking protocol is used by the MVS/XA control to achieve serial use of program control-program resources when multiprogramming and multiprocessing activities may cause temporary, multiple demands for the same resource. One instruction assists the execution of the SUPERVISOR CALL instruction, which is frequently to control-program request The services. remaining instruction helps in establishing parameters that are used by the MVS/XA control program following detection of certain failure and error conditions. The instructions are:

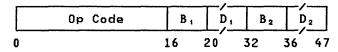
ADD FRR
OBTAIN CMS LOCK
OBTAIN LOCAL LOCK
RELEASE CMS LOCK
RELEASE LOCAL LOCK
SVC ASSIST

All of the instructions are privileged instructions.

ADD FRR uses the RRE instruction format:



The other instructions use the SSE format. The SSE instruction format is:



In addition to the operands explicitly designated by the instructions, implicit operands at fixed storage locations are also used. Sometimes operands fetched from storage are used in turn for addressing still other operands. Unless otherwise stated, all operand addresses are logical. Either 24 or 31 bits are used for addressing, depending on the value of the A-mode bit in the current PSW. The leftmost eight bits or the leftmost bit, depending on the A-mode bit in the current PSW, are ignored in a word which is a source of an operand address or a branch address. Storage

protection applies to all storage operand accesses in the usual way. A specification exception is recognized for any operand that does not meet a specified alignment requirement, and instruction execution is suppressed.

Not all possible operands are accessed for every execution of an instruction. In some cases the state of control bits or the results of comparisons of operands determine which operands are subsequently accessed. Exceptions may or may not be recognized for operands not needed for completion of a particular execution of an instruction, except where the definition of a particular instruction notes restrictions.

Program-event recording applies in the usual fashion to these instructions, except that branch events are not recognized when these instructions cause branching.

The two instructions which obtain locks fetch and test a word for zeros, and, if the fetched word is zero, cause a specified value to be stored back at the same location by means of an interlocked update. The two instructions which release locks may perform a conditional interlocked update of a doubleword location. For these instructions, if the second word of a doubleword location contains all zeros, the entire doubleword is set to zeros by means of an interlocked update. When the SVC ASSIST instruction is used to assist a type-I SVC, a conditional interlocked update of a doubleword location is performed which is similar to that performed by the COMPARE DOUBLE AND SWAP instruction.

The four lock-handling instructions may, under certain conditions, access the lock-interface-table prefix, consisting of the four words at negative offsets from the lock-interface table. The lock-interface table is located by using an address contained in the word in main storage following the second operand.

SIMPLIFIED EXECUTION PATHS

Simplified execution paths are defined for the following instructions. When a simplified path is used, specific actions defined for the corresponding instruction are performed unconditionally; that is, specific actions are taken without the prescribed tests being made to determine that those actions should be selected.

- Execution
- already held.
- RELEASE CMS LOCK. Execution proceeds as if the currently dispatched unit of work held no CMS
- already released.
- <u>SVC</u> <u>ASSIST</u>. Instruction execution is completed with normal instruc-SVC ASSIST. tion sequencing and without the performance of other actions.

Use of the simplified execution paths is not apparent to application programs using program products. In certain models, simplified execution paths are used when the available control-storage space is limited. The ADD FRR instruction is optional and may or may not be provided as part of the assists.

INSTRUCTIONS

The instructions described in this section are listed in the figure "Instruction Summary," as are their operation codes and the programinterruption conditions that can be recognized when they are executed.

In the format shown in the instruction description, the operation code is given in hex, which is signified by enclosing its value in single quotation marks ('XXXX').

Name		С	har	acte	rist	tics		Code
ADD FRR OBTAIN CMS LOCK OBTAIN LOCAL LOCK RELEASE CMS LOCK RELEASE LOCAL LOCK SVC ASSIST	RRE SSE SSE SSE SSE SSE	M M M M	A A A A A	SP SP SP SP SP	R R R R R R	0P	ST ST ST ST ST ST	B242 E506 E504 E507 E505 E503

Explanation:

Access exceptions

Privileged-operation exception

0P Operation exception

R PER general-register-alteration event RRE RRE instruction format

SP Specification exception

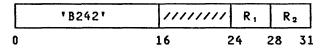
SSE SSE instruction format

PER storage-alteration event

Instruction Summary

OBTAIN CMS LOCK

[RRE]



A new entry is added to the top of the current functional-recovery-routine (FRR) stack. The entry is initialized with values provided in general registers, with the PSW A-mode bit (bit 32), and with the PSW S bit (bit 16). Optionally, the contents of control registers 3 and 4 are saved in an entry in a separate table.

The general register designated by the R_2 field provides the logical address of the FRR entry point.

Before instruction execution, the general register designated by the R₁ field provides three bytes that are stored in the new FRR entry and whose value determines if control registers 3 and 4 are to be stored as well. When instruction execution is completed, the register designated by R₁ contains the logical address of the six-word work area within the new, current FRR-stack entry.

Logical location 380 hex contains the logical address of the stack-table header. The stack-table header contains (1) a logical address which is 32 less than the address of the first dynamic entry in the stack table, (2) the logical address of the last entry in the stack table, and (3) the logical address of the current stack-table entry.

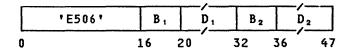
At an offset from the beginning of the stack-table header is found a table of stack-entry-extension entries. Optionally, the contents of control registers 3 and 4 are saved in an extension entry. One extension entry corresponds to each entry in the stack table. The offset to the table of extension entries, and the encoded length of an extension entry, are found in the word at logical location BA8 hex.

<u>Condition</u> <u>Code:</u> The code remains unchanged.

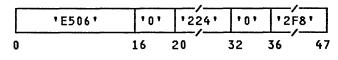
Program Exceptions:

Access (storage operands)
Operation (when the instruction is not installed)
Privileged operation
Specification

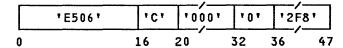
General Form



Forms Used in the Control Program



or



If the word fetched from the secondoperand location shows that the executing CPU holds the local lock and does
not hold a CMS lock, and if the CMS lock
addressed by general register 11 is not
held, then the lock is replaced, by
using an interlocked update, with the
first-operand word. Also, an indicator
is set to show that a CMS lock is held,
and zeros are placed in general register
13.

Otherwise, the updated instruction address is placed in general register 12, with a zero placed in bit position 0; the contents of the word at LIT minus 8 are placed in general register 13, and bits are selected from the word which replace the instruction-address portion of the PSW, leaving the A-mode bit unchanged.

Serialization occurs before the lock is fetched and, if the lock is obtained, again after the lock is updated.

<u>Condition Code:</u> The code remains unchanged.

Program Exceptions:

Access Privileged operation Specification

OBTAIN LOCAL LOCK

General Form

'E504'	В 1	D ,	В 2	D ₂	
0	16	20	32	36	47

Form Used in the Control Program

'E504'	'0'	12241	'0'	'2F8'	
0	16	20	32	36	47

If the local lock in the ASCB addressed by the first-operand word is not held, the lock is replaced, by using an interlocked update, with the value from the word at the second-operand location minus 4; the local-lock bit of the word fetched from the second-operand location is set to one; and zeros are placed in general register 13.

Otherwise, the updated instruction address is placed in general register 12, with a zero placed in bit position 0; the contents of the word at LIT minus 16 are placed in general register 13, and bits are selected from the word which replace the instruction-address portion of the PSW, leaving the A-mode bit unchanged.

Serialization occurs before the local lock is fetched and, if the lock is obtained, again after the lock is updated.

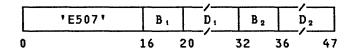
<u>Condition Code:</u> The code remains unchanged.

Program Exceptions:

Access Privileged operation Specification

RELEASE CMS LOCK

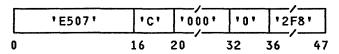
General Form



Forms Used in the Control Program

	'E507'	'0'	12241	'0'	'2F8'	
0		16	20	32	36 4	7

or



If (1) the contents of the CMS lockword addressed by general register 11 equal the contents of the first-operand word, (2) the currently dispatched unit of work holds a CMS lock, and (3) the word

after the CMS lockword is zero, then the doubleword containing the lockword is set to zero by using an interlocked update. Also, the word fetched from the second-operand location is set to show that no CMS lock is held, and zeros are placed in general register 13.

Otherwise, the updated instruction address is placed in general register 12, with a zero placed in bit position 0; the contents of the word at LIT minus 4 are placed in general register 13; and bits are selected from the word which replace the instruction-address portion of the PSW, leaving the A-mode bit unchanged.

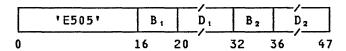
<u>Condition Code:</u> The code remains unchanged.

Program Exceptions:

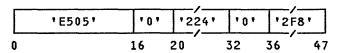
Access Privileged operation Specification

RELEASE LOCAL LOCK

General Form



Form Used in the Control Program



If the word fetched from the secondoperand location shows that the executing CPU holds the local lock and does
not hold a CMS lock, and if the word
after the local lock word in the ASCB
addressed by the first-operand word is
zero, then the doubleword containing the
lock is set to zero by using inter
locked update. Also, the local-lock bit
of the highest-lock-held-indicator word
is set to zero, and zeros are placed in
general register 13.

Otherwise, the updated instruction address is placed in general register 12, with a zero placed in bit position 0; the contents of the word at LIT minus 12 are placed in general register 13, and bits are selected from the word which replace the instruction-address portion of the PSW, leaving the A-mode bit unchanged.

<u>Condition</u> <u>Code:</u> The code remains unchanged.

Program Exceptions:

Access Privileged operation Specification

SVC ASSIST

General Form

	'E503'	В	[j	D 1	В2		2
0	1	16	20	3	32	36	47

Form Used in the Control Program

	'E503'	'0'	12241	'0'	'21C'	٦
0		16	20	32	36	47

The first and second operands are words containing the addresses of the current ASCB and TCB, respectively.

Main-storage locations which contain status for the last SVC interruption are tested to determine if SVC-assist action is to be taken. If SVC-assist action is not taken, instruction execution is completed with normal instruction sequencing. No assist action is taken unless the CPU was enabled prior to the last SVC interruption, and a task is currently dispatched which holds no locks and for which SVC screening is not activated. In addition, no assist action is taken unless (1) both the primary space and the secondary space match the dispatched space, and (2) primary-space mode is specified in the supervisor-call old PSW.

A type-1 SVC request is assisted only if the request is for an assistable type-1 function, if the only lock needed for the requested SVC function is the local lock, and if an attempt to obtain the local lock is successful.

A type-2, -3, or -4 SVC request is assisted only if the request is for an assistable function of type 2, 3, or 4, and if an attempt to dequeue an SVRB from the SVRB pool of the current address space is successful.

A type-6 SVC request is assisted only if the request is for an assistable type-6 function for which no locks are needed.

Assist action consists in copying the information stored at the last SVC interruption into the current request block, saving all 16 general registers, loading the general registers as shown below, and then loading the instruction address portion of the PSW from general register 9.

register	9.
GR No.	Contents Loaded on <u>Assist</u> <u>Action</u>
3	Address of CVT
4	Second operand
5	Address of the current RB for type 1 or type 6; address of the acquired SVRB for type 2, 3, or 4
6	SVC entry-point address
7	First operand
8	Address of the RB for the program which was being executed at SVC interruption
9	Entry address from that MPL field which is appropriate to the SVC type, that is, the address of the routine that receives control on completion of the assist
11	Address of the MBCB part of the SVRB acquired; for type 2, 3, or 4 only
12	Address of requested SVC-table entry
14	Exit address from that MPL field which is appropriate to

<u>Condition</u> <u>Code:</u> The code remains unchanged.

Program Exceptions:

Access Privileged operation Specification

the SVC type

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