## OS/VS-DOS/VSE-VM/370

 Assembler LanguagePage of GC33-4010-5
As Updated 28 Dec 1981
By TNL GN20-9372

## Sixth Edition (March 1979)

This edition, as amended by technical newsletter GN20-9372, applies to Release 4 of OS/VS1, Release 3 of OS/VS2, Release 2 of VM/370, DOS/VSE, and to any subsequent releases until otherwise indicated in new editions or technical newsletters.

Changes are periodically made to the information herein; before using this publication in connection with the operation of IBM systems, consult the latest IBM System/370 and 4300 Processors Bibliography, GC20-0001, for the editions that are applicable and current.

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## Read This First

This manual describes the OS/VS - DOS/VSE - VM/370 assembler language.

The OS/VS - VM/370 assembler language offers the following improvements over the 0S/360 assembler language as processed by the $F$ assembler:

1. New instructions and functions
2. Relaxation of language restrictions on character string lengths, attribute usage, SET symbol dimensions, and on the number of entries allowed in the External Symbol Dictionary
3. New system variable symbols
4. New options: for example, for the printing of statements in the program listings or for the alignment of constants and areas.

The figure on the following pages lists in detail these assembler language improvements and indicates the sections in the manual where the instructions and functions incorporating these improvements are described. If you are already familiar with the OS/360 assembler language as processed by the $F$ assembler, you need only read those sections. Also included in the figure on the following pages are the improvements of the DOS/VS assembler language over the DOS/360 assembler language as processed by the D assembler.

NOTE: Sections I through L, describing the macro facility and the conditional assembly language, have been expanded to include more examples and detailed descriptions.

## Note for VM/370 Users

The services provided by the OS Linkage Editor and Loader programs are paralleled in VM/370 by those provided by the CMS Loader. Therefore, for any reference in this publication to those OS programs, you may assume that the CMS Loader performs the same function.

Certain shaded notes in this publication refer to "OS only" information. Nhere you see these notes you may assume the information also applies for VM/370 users.

Note for DOS/VSE Users
All references to DOS and DOS/VS are also applicable to DOS/VSE.




1 This is a reference manual for the OS/VS - DOS/VSE - VM/370 assembler language. It will enable you to answer specific questions about language functions and specifications. In many cases it also provides information about the purpose of the instruction you refer to, as well as examples of its use.

The manual is not intended as a text for learning the assembler language.

## Who This Manual Is For

| This manual is for programmers coding in the OS/VS - VM/370 or DOS/VSE assembler language.

## Major Topics

This manual is divided into four main parts (aside from the "Introduction" and the Appendixes) :

PART I (Sections $B$ and $C$ ) describes the coding rules for, and the structure of, the assembler language. PART II (Section D) describes the machine instruction types and their formats. PART III (Sections E through H) describes the assembler instructions. PART IV. (Sections I through L) describes the macro facility and the conditional assembly language.

How To Use This Manual

Since this is a reference manual, you should use the Index or the Table of Contents to find the subject you are interested in.

Complete specifications are given for each instruction or feature of the assembler language (except for the machine instructions, which are documented in Principles of Operation -- see "References You May Need"). In many cases, a "Purpose" section suggests why you might use the feature; a "Specifications" section explains use of a complex feature; and one or more figures give examples of coding an instruction.

If you are a present user of the OS Assembler $F$ or the DOS Assembler D, you need only read those sections listed in the table preceding this "Preface" which indicates those language features that are different from the DOS or OS System/360 languages.

TABS: Tabs mark the beginning of the specifications portion of the language descriptions. Use the tabs for quick referencing.

Tab -

## USING

OS-DOS DIFFERENCES: Wherever the OS/VS and DOS/VS assembler languages differ, the specifications that apply only to one assembler or the other are so marked. The 'OS only' markings also apply for the VM/370 assembler.
os only

KEYS: The majority of figures are placed to the right of the text that describes them. Numbered keys within a figure are duplicated to the left of the text describing the figure. Use the numbered keys to tie the underlined passages in the text to specific parts of the figure.

## Key - 3

GLOSSARY: The glossary at the back of the manual contains terms that apply to assembler programming specifically and to allied terms in data processing in general. You can use the Glossary for terms that are unfamiliar to you.

IBM is grateful to the American National Standards Institute (ANSI) for permission to reprint its definitions from the American National Standard Vocabulary for Information Processing, which was prepared by Subcommittee X3.5 on Terminology and Glossary of American National Standards Committee X3.

## References You May Need

You may want to refer to
IBM System/370 Principles of Operation, GA22-7000, or IBM 4300
Processors Principles of Operation for ECPS:VSE Mode, GA22-7070, the definitive publications for machine instructions, and to

OS/VS - VM/370 Assembler Programmer's Guide, GC33-4021
for detailed information about the OS/VS - VM/370 Assembler.
Guide to the DOS/VSE Assembler, GC33-4024
for detailed information about the DOS/VSE Assembler.

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## Section A: Introduction

## What the Assembler Does

A computer can understand and interpret only machine language. Machine language is in binary form and, thus, very difficult to write. The assembler language is a symbolic programming language that you can use to code instructions instead of coding in machine language.

Because the assembler language allows you to use meaningful symbols made up of alphabetic and numeric characters instead of just the binary digits 0 and 1 used in the machine language, you can make your coding easier to read, understand, and change.

The assembler must translate the symbolic assembler language into machine language Lefore the computer can execute ycur program, as shown in the figure below.


> Assume that your program, written in the assembler language, has been punched into a deck of cards called the scurce deck. This deck, also known as a source module, is the input to the assembler. (You can also enter a source module as input to the assembler through a terminal.)

The assembler processes your source module and produces an object module in machine language (called object ccde). Assume that the assembler punches this object module into a deck of cards called the okject deck.

The object deck or object module can be used as input to be processed by another processing progran, called the linkage editor. The linkage editor produces a load module that can be loaded later into the main storage of the computer, which then executes the program. Your source module and the object code produced is printed, alcng with other information on a program listing.

## Al -- The Assembler Language

The assembler language is the symbolic programming language that lies closest to the machine language in form and content. You will, therefore, find the assembler language useful when:

```
- You need to control your program closely, down to the
byte and even bit level or
- You must write subroutines for functions that are not
provided by other symbolic programming languages such as:
ALGOL, COBOL, FORTRAN, or PL/I.
The assembler language is made up of statements that
represent instructions or comments. The instruction
statements are the working part of the language and are
divided into the following three groups:
1. Machine instructions
2. Assembler instructions
3. Macro instructions.
```


## Machine Instructions

A machine instruction is the symbolic representation of a machine language instruction of the IBM Syster/ 370 instruction set. It is called a machine instruction because the assembler translates it into the machine language code which the computer can execute. Machine instructions are described in PART II; SECTION C of this manual.

```
An assembler instruction is a request to the assembler
program to perform certain operations during the assembly
of a source module, for example, defining data constants,
defining the end of the source module, and reserving storage
areas. Except for the instructions that define constants,
the assembler does not translate assembler instructions
into object code. The assembler instructions are described
in PART III; SECTIONS E, F,G, and H and PART IV; SECTIONS
J, K, and L of this manual.
```


## Macro Instructions

A macro instruction is a request to the assembler program to process a predefined sequence of code called a macro definition. From this definition, the assembler gener ates machine and assembler instructions which it then processes as if they were part of the original input in the source module.

IBM supplies macro definitions for input/output, data management, and supervisor operations that you can call for processing by coding the required macro instruction. (These IBM-supplied macro instructions are not described in this manual.)

You can also prepare your own macro definitions and call them by coding the corresponding macro instructions. This macro facility is introduced in PART IV; SECTION I. A complete description of the macro facility, including the macro definition, the macro instruction and the conditional assembly language, is given in PART IV; SECTIONS $J, K$, and L.

## A2 -- The Assembler Program

The assembler program, also referred to as the "assembler", processes the machine, assembler, and macro instructions you have coded in the assembler language and produces an object module in machine language.

The assembler processes the three types of assembler language instructions at different times during its processing sequence. You should be aware of the assembler's processing sequence in order to code your program correctly. The figure below relates the assembler processing sequence to the other times at which your program is processed and executed.


The assembler processes most instructions on two occasions;
(1)
first at pre-assembly time and later at assembly time.
However, it does some processing, for example, macro processing, only at pre-assembly time.

The assembler also produces information for other processors. The linkage editor uses such information at
(3) linkage-edit time to combine object modules into load modules. The loader loads your program (combined load modules) into virtual storage (see GLOSSARY) at program (4) fetch time. Finally, at execution time, the computer assembly time.

The assembler processes all machine instructions and translates them into object code at assembly time, as shown in the figure belcw.


## Assembler Instruction Processing

Assembler instructions are divided into two main types:

1. Ordinary assembler instructions
2. Conditional assembly instructions and the macro processing instructions MACRO, MENE, MEXIT and MNOTE .

The assembler processes ordinary assembler instructions
(1) at assembly time, as shown in the figure below.


NOTES:

1. The assembler evaluates absolute and relocatable expressions at assembly time; they are sometimes called assembly time expressions.
(3) 2. Some instructions produce output for processing after assembly time.

The assembler processes conditional assembly instructions 1 and macro processing instructions at pre-assembly time, as shown in the figure below.


NOTES:
21. The assembler evaluates the conditional assembly expressions (arithmetic, logical, and character) at preassembly tire.
2. The assembler processes the machine and assembler instructions generated from pre-assembly processing at assembly time.

Nacro Instruction Processing

The assembler processes macro instructions at pre-assembly time, as shown in the figure below.


2
NOTE: The assembler processes the machine and ordinary assembler instructions generated from a macro definition called by a macro instruction at assembly time.

The assembler prints in a program listing all the information it produces at the various processing times described in the above figures.

## A3 - Relationship of Assembler to Operating System

The assembler is a programming component of the OS/VS, VM/370, or DOS/VS. These system control programs provide the assembler with the services:

- For assembling a source module and
- For running the assembled orject module as a progran.

In writing a source module you must include instructions that request the desired service functions from the operating system.

## Services Frovided by the Operating System

OS NS and DOS/VS provide the following services:

1. For assembling the source module:
a. A control program
b. Libraries tc contain source code and macro definitions
c. Utilities
2. For preparing for the execution of the assembler program as represented by the okject module:
a. A control program
b. Storage allocation
c. Input and output facilities
d. A linkage editor
e. A loader.

VM/370 provides the following services:

1. For assembling the source module:
a. An interactive control program
b. Files to contain source code and macro definitions
c. Utilities.
2. For preparing for the execution of the assembler programs as represented by the object modules:
a. An interactive control program
b. Storage allocation
c. Input and output facilities
d. The CMS Loader.

## A4 -- Coding Aids

It can be very difficult to write an assembler language program using only machine instructions. The assembler provides additional functions that make this task easier. They are summarized belcw.

Symbolic Representation of Program Elements

Symbols greatly reduce programming effort and errcrs. You can define symbols to represent storage addresses, displacements, constants, registers, and almost any element that makes up the assembler language. These elements include operands, operand subfields, terms, and expressions. Symbols are easier to rememker and code than numbers; moreover, they are listed in a symkol cross-reference table which is printed in the program listings. Thus, you can easily find a symbol when searching for an error in your code.

## Variety of Lata Representation

You can use decimal, binary, hexadecimal or character representation which the assembler will convert fcr you into the binary values required by the machine language.

Controlling Address Assignment

If you code the apfropriate assembler instruction, the assembler will compute the displacement from a base address of any symbclic addresses you specify in a machine instruction. It will insert this displacement, along with the base register assigned ky the assembler instruction, into the object code of the machine instruction.

At execution time, the object code of address references must be in the base-displacement form. The computer obtains the required address ky adding the displacement to the base address contained in the base register.

## Relocatability

```
The assembler produces an object module that can be
relocated from an originally assigned storage area to any
other suitable virtual storage area without affecting
program execution. This is made easier because most
addresses are assembled in their base-displacement form.
```


## Segmenting a Program

You can divide a source module into one or more control sections. After assembly, you can include or delete individual control sections from the resulting object module before you load it for execution. Control sections can be loaded separately into storage areas that are not contiguous.

## Linkage Between Source Modules

You can create symbolic linkages ketween separately assembled scurce modules. This allows you to refer symbolically from one source module to data defined in another scurce module. You can also use symbolic addresses to branch between modules.

Program Listings
The assembler produces a listing of your source module, including any generated statements, and the object code assembled from the source module. You can control the form and content of the listing to a certain extent. The assembler also prints messages akout actual errors and warnings abcut potential errors in your source module.

## Part I: Coding and Structure

SECTION B: CODING CONVENTIONS
SECTION C: ASSEMBLER LANGUAGE STRUCTURE

## Section B: Coding Conventions

```
This section describes the coding conventions that you must
follow in writing assembler language programs. Assembler
language statements are usually written on a coding form
before they are punched onto cards, or entered as source
statements through other forms of input (for example,
through terminals or directly onto tape).
```

Standard Assembler Coding Form
You can write assembler language statements on the standard
coding form (Order No. GX28-6509) shown kelow. The cclumns
on this form correspend to the columns on a punched card
or positions on a source statement entered through a
terminal. The form has space for program identification
and instructions to keypunch operatcrs.


## Bl -- Coding Specifications

B1A - FIELD BOUNDARIES

Assembler language statement usually occupy one 80-column line on the standard form (for statements occupying more than 80 columns, see $B 1 B$ below). Note that any printable character punched into any column of a card, or otherwise entered as a position in a source statement, is reproduced in the listing printed by the assembler. All characters are placed in the line by the assembler. Whether they are printed or not depends on the printer. Each line of the coding form is divided into three main fields:
(1) The Statement field,

2 The Identification -Sequence field, and

The Statement Field

The instructions and comments statements must be written in the statement field. The statement field starts in the "begin" column and ends in the "end" column. Any continuation lines needed must start in the "continue" column and end in the "end" column. The assembler assumes the following standard values for these columns:
(4) - The "begin" column is column 1
(5) - The "end" column is column 71, and

6 - The "continue" column is column 16.
These standard values can be changed by using the ICTL instruction. However, all references to the "begin", "end", and "continue" columns in this manual refer to the standard value described above.


The identification-sequence field can contain identification characters or sequence numbers or both. If the ISEQ instruction has been specified to check this field, the assembler will verify whether or not the source statements are in the correct sequence.

NOTE: The field the assembler normally checks lies in columns 73 through 80. However, if the ICTL instruction has been used to change the begin and end columns, the boundaries for the identification-sequence field can be affected.

## The Continuation Indicator Field

The continuation indicator field occupies the column after the end column. Therefore, the standard position for this field is column 72. A non-blank character in this column indicates that the current statement is continued on the next line. This column must be blank if a statement is completed on the same line; otherwise the assembler will treat the statement that follows on the next line as a continuation line of the current statement.

Field Positions

The statement field always lies between the begin and the end columns. The continuation indicator field always lies in the column after the end column. The identificationsequence field usually lies in the field after the continuation indicator field. However, the ICTL instruction, by changing the standard begin, end, and continue columns can create a field before the begin column. This field can then contain the identification-sequence field.

To continue a statement on another line, the following applies:

1. Enter a non-blank character in the continuation indicator field (column 72). This non-blank character must not be part of the statement coding. When more than one continuation line is needed, a non-blank character must be entered in column 72 of each line that is to be continued.
2. Continue the statement on the next line, starting in the continue column (column 16). Columns to the left of the continue column must be blank. Comments may be continued after column 16.

Note that if an operand is continued after column 16 it is taken to be a comment. Also if the continuation indicator field is filled in on one line and the user tries to start a totally new statement after column 16 on the next line, this statement will be taken as a comment belonging to the previous statement.

Only two continuation lines are allowed for a single assembler language statement. However, macro instruction statements and the prototype statement of macro definitions can have as many continuation lines as needed.


Comments statements are nct assembled as part of the ckject module, but are only printed in the assembly listing. As many comments statements as needed can ke written, subject to the following rules:

1. 2. Comments statements require an asterisk in the begin column.

NOTE: Internal macro definition comments statements require a period in the begin column, followed by an asterisk (for details see J6A).
2. Any characters, including blanks and special characters, of the IBM System/370 Character Set (see C3) can be used.
3. Comments statements must lie in the statement field

2 and not run over into the continuation indicator field; otherwise the statement following the comments statement will be considered as a continuation line of that comments statement.
4. Comments statements must not appear ketween an instruction statement and its continuation lines.


The statement field of an instruction statenent must be formatted tc include frcm one to four of the following entries:

1. A narfe entry
2. An operaticn entry
3. An operand entry
4. A remarks entry.

## Fixed_Fcrmat

The standard coding form is divided into fields that provide fixed positions for the first thré entries, as fcllcws:

(1) An 8-character name field starting in cclumn 1.
(2) A 5-character operation field starting in cclumn 10.
(3) An operand field that begins in column 16.
(4) Note that with this fixed format one klank separates each field.

## Free Format

It is not necessary to code the nare, cperaticn, and cperand entries acccrding tc the fixed fields on the standard coding form. Instead, these entries can ke written in any fcsition, subject tc the formatting specifications below.

## Formatting Specifications

Whether using fixed or free format, the following general rules apply to the coding of an instruction statement:

1. The entries must be written in the following order: name, operation, operand, and remarks.
2. The entries must be contained in the begin column
through the end column (71) of the first line and; if needed, in the continue column (16) through the end cclumn (71) of any continuation lines.
(1) 3. The entries must be separated from each other by one or more blanks.

2 4. If used, the name entry must start in the begin column.
5. The name and operation entries, each followed by at

3 least one blank, must be contained in the first line of an instruction statement.

4 6. The operation entry must start at least one column tc the right of the begin column.


THE NAME ENTRY: The name entry identifies an instruction statement.

The following applies to the name entry:

1. It is usually optional.
2. It must be a valid symbol at assembly time (after
substitution for variable symbols, if specified); for an exception see the TITLE instruction (H3E).

THE OPERATION ENTRY: The operation entry provides the symbolic operation code that specifies the machine, assembler, or macro instruction to be processed. The following applies to the operation entry:

1. It is mandatory.
2. For machine and assembler instructions it must be a valid symbol at assembly time (after substitution for variable symbols, if specified). The standard symbolic operation codes are five characters or less (see Appendixes IV and V).

OS NOTE: The standard set of codes can be changed by OPSYN only instructions (as described in H5).
3. For macro instructions it can be any valid symbol that is not identical to the operation codes described in 2 above.

THE OFERAND ENTRY: The operand entry has one or more operands that identify and describe the data used by an instruction. The following applies to operands:

1. One or more operands are usually required, depending on the instruction.
2. Operands must be separated by commas. No blanks are allowed between the operands and the commas that separate them.
3. Operands must not contain embedded blanks, because a blank normally indicates the end of the operand entry. However, blanks are allowed if they are included in character strings enclosed in apostrophes (for example, C'J N') or in logical expressions (see L4C).

THE REMARKS ENTRY: The remarks entry is used tc descrike the current instruction. The following applies to the remarks entry:

1. It is ceticnal.
2. It can contain any of the 256 characters for punch combinations) of the IEM Syster/370 character set, including klanks and special characters.
(13) 3. It can follow any operand entry.
3. If an cpticnal cperand entry is omitted, remarks are allowed if the absence of the operand entry is indicated


## Section C: Assembler Language Structure

This section describes the structure of the assembler language, that is, the various statements which are allowed in the language and the elements that make up those statements.

## Cl -. The Source Module

```
A source module is a sequence of assembler language
statements that constitute the input to the assembler.
The figure on the orfosite page shows an overall picture
of the structure of the assembler language.
```


## C2 - Instruction Statements

The instruction statements of a source module are composed of one to four entries that are contained in the statement field. Other entries outside the statement field are discussed in B1A. The four statement entries are:

1. A name entry (usually optional)
2. An operation entry (mandatory)
3. An operand entry (usually required)
4. A remarks entry (optional).

NOTES:

1. The figures in this subsection show the overall structure of the statements that represent the assembler language instructions and are not specifications for these instructions. The individual instructions, their purposes, and their specifications are described in other sections of this manual (as cross-referenced in the figures). Model statements, used to generate assembler language statements, are described in $\mathrm{J}^{4}$.
2. The remarks entry is not processed by the assembler, but only copied into the listings of the program. It is therefore not shown except in the overview opposite.


The machine instruction statements are described in the figure below.

The instructions themselves are discussed in Part II cf this manual and summarized in Appendix IV.


## C2B -- ASSEMBLER INSTRUCTIONS

The assembler instruction statements can be divided into two main groups: ordinary assembler instructions and conditional assembly instructions.

## Ordinary Assembler Instructions

Ordinary assembler instruction statements are described in the figure on the opposite page.

These instructions are discussed in Part III of this manual and summarized in Appendix $V$.


## Conditional Assembly Instructions

Conditional assembly instruction statements and the macro processing statements (MACRO, MEND, MEXIT, MNOTE) are described in the figure below.

The conditional assembly instructions are discussed in Section $L$ and macro processing instructions in Section $J$; both types are summarized in Appendix $V$.


## Macro instruction statements are described in the figure below; the prototype statement of a macro definition, which serves as a model for the macro instruction staterent, is also shown. <br> Macro instruction statements are discussed in Section $k$ of this manual and the prototype statement is discussed in Section J2.



## C3 - Character Set

Terms, expressions, and character strings used to build source statements are written with the fcllowing characters:

1. Alphameric Characters

Alphabetic characters (or letters): A through $Z$, and
\$, \#, a
Ligits (or numerals): 0 through 9
2. Special characters

+     - , = * () • / E Elank
Examples, showing the use of the above characters are given in the figure below.

Normally, ycu would use strings of alphameric characters to represent data (terms, see C4), and special characters as:
a. Arithmetic cperators in expressions
b. Data or field delimiters
c. . Indicators to the assembler for specific handling.

Characters are represented ky the card-funch combinaticns and internal bit configurations listed in Appendix I.
In addition to the printable characters listed abcve, any of the 256 combinations for punched cards listed in Affendix I can be used:

1. Between paired apcstrophes
2. As statement remarks
3. In comments statements
4. In macrc instructicn operands (for restrictions see K5).

Char. Set

\begin{tabular}{|c|c|c|c|}
\hline Characters \& Usage \& Example \& Constituting \\
\hline Alphameric \& In symbols \& LABEL NINE\#01 \& Terms \\
\hline Digits \& As decimal self-defining terms \& 019 \& Terms \\
\hline \begin{tabular}{l}
Special Characters \\
\(+\) \\
- \\
* \\
/ + or -
\end{tabular} \& \begin{tabular}{l}
As Operators \\
Addition \\
Subtraction \\
Multiplication \\
Division \\
(Unary)
\end{tabular} \& \begin{tabular}{l}
\(\left.\begin{array}{l}\text { NINE+FIVE } \\ \text { NINE-5 } \\ \text { } 9 * F I V E \\ \text { TEN } / 3\end{array}\right\}\) \\
+NINE -FIVE
\end{tabular} \& \begin{tabular}{l}
Expressions \\
Terms
\end{tabular} \\
\hline \begin{tabular}{l}
Blanks \\
Comma \\
Apostrophes \\
Parentheses
\end{tabular} \& \begin{tabular}{l}
As Delimiters \\
Between fields \\
Between operands \\
Enclosing character strings \\
Enclosing subfields or subexpressions
\end{tabular} \& \begin{tabular}{l}
LABEL AR 3,4 \\
OPND1,OPND2 \\
C'STRING' \\
MOVE MVC TO (80), FROM
\[
(\mathrm{A}+\mathrm{B} *(\mathrm{C}-\mathrm{D}))
\]
\end{tabular} \& \begin{tabular}{l}
Statement \\
Operand field \\
String \\
Statement \\
Expression
\end{tabular} \\
\hline Ampersand
Period

Asterisk

Equal sign \& \begin{tabular}{l}
As indicators for <br>
Variable symbol <br>
Sequence symbol <br>
Comments statement in Macro definition Concatenation <br>
Bit-length specification <br>
Decimal point <br>
Location counter reference <br>
Comments statement <br>
Literal reference <br>
Keyword

 \& 

\&VAR <br>
. SEQ <br>

- *THIS IS A COMMENT <br>
\&VAR.A <br>
DC CL. $7^{\prime} \mathrm{AB}^{\prime}$ <br>
DC $F^{\prime} 1.7 E 4^{\prime}$ <br>
$*+72$ <br>
* THIS IS A COMMENT <br>
L $\quad 6,=F^{\prime} 2^{\prime}$ <br>
$\& \mathrm{KEY}=\mathrm{D}$

 \& 

Term <br>
(label) <br>
Statement <br>
Term <br>
Operand <br>
Operand <br>
Expression <br>
Statement <br>
Statement <br>
Keyword <br>
Parameter
\end{tabular} <br>

\hline
\end{tabular}

## C4 -- Terms

A term is the smallest element of the assembler language that represents a distinct and separate value. It can therefore te used alone or in combination with other terms to form expressions. Terms have absolute or relocatable values that are assigned by the assembler or are inherent in the terms themsel ves.

A term is absolute if its value does not change upon program relocation and is relocatable if its value changes upon relocation. The various types of terms described below are summarized in the figure to the right.

|  |  |  | Terms |  |
| :---: | :---: | :---: | :---: | :---: |
| Terms | Term Can Be |  | Value Is |  |
|  | Absolute | Relocatable | Assigned by Assembler | Inherent in Term |
| Symbols | X | x | X |  |
| Location Counter Reference |  | X | X |  |
| Symbol Length Attribute | X |  | X |  |
| Other Data <br> Attributes | x |  | x |  |
| Self-Defining Terms | x |  |  | x |

## Furpose

You can use a symbcl tc represent storage locations or arbitrary values.

SYMBOLIC REFRESENTATICN: You can write a symbol in the name field of an instruction. Ycu can tren specify this symbcl in the cperands cf other instructions and thus refer to the former instruction symbclically. This symbol represents a relocatable address.

You can alsc assign an absolute value to a symbol ky coding it in the name field of an EQU instruction with an crerand whose value is absclute. This allows you to use this symbol in instruction operands tc represent registers, displacements in explicit addresses, imediate data, lengths, and implicit addresses with atsclute values. Fcr details of these frcgram elements, see [5. The advantages of symbolic over numeric recresentation are:

1. Symbcls are easier tc remember and use than numerical values, thus reducing programming errors and increasing crogramming efficiency.
2. Ycu can use meaningful symbcls to descrike the prcgram elements they represent; for example, INFUT can name a field that is to contain infut data, cr INDEX can name a register to te used for indexing.
3. You can change the value of one symbol (through an EQU instructicn) more easily than you can change several numerical values in many instructions.
4. Symbols are entered into a cross-reference table that the assembler frints in the frogram listing. This table helps you to find a symkol in a prcgram listing, because it lists (1) the number of the staterent in which the symbel is defined (that is, used as the name entry) and (2) the numbers of all the statements in which the synkcl is used in the cperands.

THE SYMEOL TABLE: The assembler maintains an internal takle called a symbcl takle. When the assembler frccesses your source statements for the first time, the assembler assigns an aksolute or relocatable value tc every symbcl that apfears in the name field of an instruction. The assembler enters this value, which normally reflects the setting cf the location counter, into the symbol table; it also enters the attrikutes asscciated with the data represented by the symbcl. The values of the symbol and its attrikutes are availakle later when the assemtlex finds this symbcl cr attribute reference used as a term in an operand or expression (Attrikute references used as terms are discussed in C4C and C4D below).

## Specifications

The three types of symbol recognized by the assembler are:

1. Ordinary symbols

Symbols
2. Sequence symbcls
3. Variable symbols.

ORDINARY SYMBOLS: Ordinary symbols can be used in the name and operand field of machine and assembler instruction statements. They must be coded in the format shown in the figure to the right.

NOTES:

1. No special characters are allowed in an ordinary symbol.
2. No blanks are allowed in an ordinary symbcl


Examples:

| HERE | \#01 | X |
| :--- | :--- | :--- |
| READER | \#12 | Y |
| A001 | @33 | Z |
| B002 | \$OPEN | F2A |

## Var. Sym.

VARIABLE SYMBOLS: Variable symbols can only be used in macro processing and conditional assembly instructions. They must be coded in the format shown in the figure to the right.

SEQUENCE SYMBOLS: Sequence symbols can only be used in macro processing and conditional assembly instructions. They must be coded in the format shown in the figure to the right.


Examples:
. SEQ
.LOOP11
.EXIT20
.T0001

Symbol Definition
An ordinary symbol is considered defined when it appears as:

1. The name entry in a machine or assembler instruction of the assembler language.
2. One of the operands of an EXTRN or WXTRN instruction.

NOTE: Ordinary symbols that appear in instructions generated from model statements at pre-assembly time are also considered defined.

The assembler assigns a value to the ordinary symbol in the name fields as follows:

1. According to the address of the leftmost byte of the storage field that contains one of the following:

1 in Any machine or assembler instruction (except the EQU or CFSYN instructions)

> b. A storage area defined by the cs instruction
(3)
c. Any constant defined by the DC instruction
d. A channel command word defined by the CCW instruction.

The address value thus assigned is relocatable, because the object code assembled from these items is relocatable; the relocatability of addresses is descriked in L5B.
2. According to the value of the first or only expression specified in the operand of an FQU instruction. This expression can have a
(4) relocatable or absolute value, which is then assigned to the ordinary 5
 symbol. The value of an ordinary symbol must lie in the range $-2^{31}$ through $+2^{31-1 .}$

UNIQUE DEFINITION: A symbol must be defined only once in a source module:
(1) either in the name field of a source statement

2
or in the operand field of an EXTRN or WXTRN instruction.

This is true even for a source module which contains two or more control sections.

NOTE: The ordinary symbol that appears in the name field of an OPSYN or TITLE instruction does not constitute a definition of that symbol. It can therefore be used in the name field of any other statement in a source module.

CONTRCL SECTION NAMES: A duplicate symbol can, however, be used as the name entry of a START, CSECT, DSECT, or COM instruction. The
(3) first time a symbol is used to name these instructions, it identifies the beginning of the control section;
4 a duplicate use of the symbol identifies the resumption of an interrupted control section.

PREVIOUSLY DEFINED SYMBOL: In some instructions the symbols used in
(5) their operands must have been defined in a previous instruction. Previously defined symbols are
 required for the operands of the

EQU
CNOP
ORG
DC and DS (in modifier and duplication factor expressions).

## Furpose

The assembler runs a location counter to assign storage addresses to your program statements. It is the assembler's equivalent of the instruction counter in the computer. You can refer to the current value of the location counter at any place in a source module by specifying an asterisk as a term in an operand.

THE LOCATION COUNTER: As the instructions and constants of a source module are being assembled, the location counter has a value that indicates a location in storage. The assembler increments the location counter according to the following:

1. After an instruction or constant has been assembled, the location counter indicates the next available location.
2. Before assembling the current instruction or constant, the assembler checks the koundary alignment required for it and adjusts
the location counter, if necessary, to indicate the frcper boundary.
3. While the instruction or constant is being assembled, the location counter value does not change. It indicates the lccation of the current data after boundary alignment

| Location in Hex |  | Source <br> Statements |  |
| :---: | :---: | :---: | :---: |
| 000004 | DONE | DC | CL3'SOB' |
| 1000007 | BEFORE | EQU | * |
| 000008 | (3) DURING | DC | $F^{\prime} 200{ }^{\prime}$ |
| 00000C | AFTER | EQU |  |
| 000010 | NEXT | DS | D | and is the value assigned to the symbol, if present, in the name field of the statement.

4. After assembling the instruction or constant, the assembler increments the location counter ky the length of the assembled data to indicate
the next available location.
The assembler maintains a location counter for each control section in a source module; for complete details about the location counter setting in control sections, see E2C. The assembler carries an internal location ccunter value as a 4-byte, 32-bit value, kut it only uses the low-crder 3 bytes, which are printed in the program listings. However, if you specify addresses greater than 224-1, you cause overflow into the high-order kyte, and the assembler issues the error message "LOCATION COUNTER OVERFLOW".

NOTE: In the figure below, an example of a location counter overflow (or wrap-around) is shown.
(1) The internal address value of the symbol $B$ is carried as a 4 -byte value, but the printed location only includes the low-order 3 bytes.
(3) The location counter value for instructions or constants is usually printed as a 3-byte value. However, the 4-byte value, with up to 3 leading zeros suppressed, is frinted for the addresses specified in the operands of the fcllowing instructions: EQU, ORG, and USING. Only 3-byte values are printed for the operands in the above instructions.

You can control the setting of the location counter in
(5) a particular control section by using the START or ORG instructions.


## Specifications

Loc. Ctr Ref
The lccation counter reference is specified Ly an asterisk (*). The asterisk can be specified as a relocatable term according to the follcwing rules:

1. It can only be specified in the cperands cf:
a. Machine instructions
b. The IC and IS instructions
c. The FQU, ORG, and USING instructions.
(1)
2. It can alsc ke specified in
literal constants (see C5).
The value of the location ccunter
 reference (*) is the current value of the lccation counter of the control section in which the asterisk
(*) is specified as a term. The
asterisk has the same value as the address of the first byte of the
instruction in which it appears (for the value of the asterisk in address constants with duplicaticn factcrs, see G3J).

Purpose

When you specify a symbol length attribute reference, you obtain the length of the instruction or data referred to by a symbol. You can use this reference as a term in instruction operands to:

1. Specify unknown storage area lengths
2. Cause the assembler to compute length specifications for you
3. Build expressions to be evaluated by the assembler.

## Specifications

The symbol length attribute reference must be specified according to the following rules:

1. The format must be $L^{\prime}$ immediately followed by a valid symbol or the location counter reference (*).
2. The symbol must be defined in the same source module in which the symbol length attribute reference is specified.
3. The symbol length attribute reference can be used in the operand of any instruction that requires an absolute term. However, it cannot be used in the form $L^{\prime} *$ in any instruction or expression that requires a previously defined symbol.

The value of the length attribute is normally the length in bytes of the storage area required by an instruction, constant, or field represented by a symbcl. The assembler stores the value of the length attribute in the symbol table along with the address value assigned to the symbcl.

When the assembler encounters a symbol length attribute reference, it substitutes the value of the attribute from the symbol table entry for the symbol specified.

The assembler assigns the length attribute values to symbols in the name field of instructions as follows:
(1) For machine instructions, it assigns either 2,4 , or 6 , depending on the format of the instruction.

2 For the DC and DS instructions, it assigns either the implicit or explicitly specified length. The length attribute is not affected by a duplication factcr.

For the ECU instruction, it assigns the length attribute value of the
(3) leftmost or only term of the first expression in the first operand, unless a specific length attribute is supplied in a second oferand.

Only one ccerand is allowed in the EQU instruction.

Note the lergth attricute values of the following terms in an EQU instructicn:
(4) - self-defining terrs
(5) lecaticn ccunter reference
(6) $L^{1 *}$

The length attribute of the location
(7)
counter reference (L'*) is equal
to the lergth attribute of the
instruction in which the $\mathrm{L}^{* *}$ appears.
For the remaining assembler
instructions, see the specifications for the individual instructions.

Length Attr.

| Source Module |  |  | Value of Symbol Length Attribute (at assembly time) |
| :---: | :---: | :---: | :---: |
| MACHA | MVC | TO,FROM | L'MACHA 6 |
| MACHB | L | 3,ADCON | L'MACHB 1 |
| MACHC | LR | 3,4 | L'MACHC |
| то | DS | CL80 | L'TO ( 80 |
| FROM | DS | CL240 | L'FROM 240 |
| ADCON | DC | A (OTHER) | L'ADCON 2 4 |
| CHAR | DC | C'YUKON' | L'CHAR 5 |
| DUPL | DC | 3F'200' | L'DUPL ( 4 |
|  |  | (3) |  |
| RELOC1 | EQU | T0 | L'RELOCl 80 |
| RELOC2 | EQU | T0 ${ }^{\text {+ }} 0$ | L'RELOC2 80 |
| ABSOLI | EQU | FROM-TO | L'ABSOL1 240 |
| ABSOL2 | EQU | ABSOL1 | L'ABSOL2 240 |
| SDT1 | EQU | 102 | L'SDTl |
| SDT2 | EQU | $X^{\prime} \mathrm{FF}^{\prime}+\mathrm{A}-\mathrm{B}$ | L'SDT2 4 l |
| SDT3 | EQU | C'YUK' | L'SDT3 |
| ASTERISK | EQU | *+10 | L'ASTERISK 51 |
| LOCTREF | EQU | $L^{1} *$ | L'LOCTREF (6) |
| LENGTHI DC | A ( $L^{\prime} *$ ) |  | L'* ${ }^{\text {¢ }}$ |
|  | TO ( | *), FROM | L'LENGTHl 7 L'* $\begin{aligned} & 4 \\ & 6\end{aligned}$ |
| LENGTH3 MVC | TO(L | TO-20), FROM | L'TO 80 |

There are other attributes which describe the characteristics and structure of the data you define in a program. For example, the kind of constant you specify or the number of characters you need to represent a value. These other attributes are the type ( $T^{\prime}$ ), scaling ( $S^{\prime}$ ), integer ( $I^{\prime}$ ), count ( $K^{\prime}$ ), and number ( $N^{\prime}$ ) attributes.

NOTE: You can refer to these attributes only in conditional assembly instructions and expressions; for full details, see L 1 B .

C4E - SELF-DEFINING TERMS

## Purpose

A self-defining term allows you to specify a value explicitly. with self-defining terms, you can specify decimal, binary, hexadecimal, or character data. These terms have absolute values and can te used as absclute terms in expressions to represent bit configurations, absolute addresses, displacements, length or other modifiers, or duplication factors.

## Specificaticns

Self-Defining

GENERPL RULES: Self-defining terms:
(1) Recresent machine language binary values

- Are absclute terms; their values do not change upon program relccaticn.

The assembler carries the values represented ky self-defining terms to 4 kytes or 32-bits; the high-
order bit is the sign bit.
Values are carried to 3 lytes or 24 kits.

CECINAL: A decinal self-defining term is an unsigned decimal number. The assembler allows:
(1) High-crder zercs
(2) A maximum of 10 decimal digits

- A range of values from 0 througr:

3) $2,147,483,647$.

DOS - A maximum of 8 decimal digits.

- A range of values from 0 through 16,777,215.

EINARY: A Linary self-defining term must be coded in the format shown in the figure to the right. The assembler:

1 - Assembles each kinary digit as it is sfecified
2. Allows a maximum of 32 tinary digits

- Allows a range of values from
(3) $-2,147,483,648$ thrcugh
$2,147,483,647$.
dos - Allows a maximum cf 24 binary digits.
- Allows a range of values from 0 through 16,777,215.

NOTE: When used as an absolute term in expressions, a tinary selfdefining term has a negative value if the high-order tit is 1.


HEXALECIMAL: A hexadecimal selfdefining term must be ccded as shown in the figure to the right. The assembler:

- Assembles each hexadecimal digit
(1) intc its 4-bit binary equivalent (listed in the figure to the right)
(2) Allows a maximum cf 8 hexadecimal
- Allows a range cf values from
(3) $-2,147,483,648$ through $2,147,483,647$.

Dos Allows a maximum cf 6 hexadecimal digits.

- Allows a range of values from

0 through 16,777,215.
NOTE: When used as an absolute term in an expression, a hexadecinal self-defining term has a negative
(4) value if the high-order tit is 1.


Conversion Table:

| Hexadecimal <br> Digit | Decimal <br> Equivalent | 4-bit <br> Binary <br> Representation |
| :---: | :---: | :---: |
| 0 | 0 | 0000 |
| 1 | 1 | 0001 |
| 2 | 2 | 0010 |
| 3 | 3 | 0011 |
| 4 | 4 | 0100 |
| 5 | 5 | 0101 |
| 6 | 6 | 0110 |
| 7 | 7 | 0111 |
| 8 | 8 | 1000 |
| 9 | 9 | 1001 |
| A | 10 | 1010 |
| B | 11 | 1011 |
| C | 12 | 1100 |
| D | 13 | 1101 |
| E | 14 | 1110 |
| F | 15 | 1111 |

Examples: Binary Value


CHARACTER: A character self-defining tern must be coded as shown in the figure to the right. The assembler:

- Allows any of the 256 punch combinations when using punched cards as infut. This includes the printable characters, that is, blanks and special characters.
- Assembles each character into its 8-bit binary equivalent. (A table of characters and their kinary equivalents can be found in Appendix I) -
(2) Requires that twc ampersands or apostrophes be specified in the character sequence for each ampersand
(3) or apostrophe required in the assembled term.
(4) Allows a maximum of 4 characters.

Dos • Allows a maximum of 3 characters.

## C5 - Literals

## Purpose

You can use literals as operands in order to introduce data into your program. However, you cannot use a literal as a term in an expression. The literal represents data rather than a reference to data. This is convenient, kecause

1. The data ycu enter as numbers for computation, addresses, or messages to be printed is visible in the instruction in which the literal appears, and
2. You avoid defining constants elsewhere in your source module and then using their symbolic names in machine instruction oper ands.

| Character self-defining term | Characters Assembled | Hexadecimal Value | Binary Value |
| :---: | :---: | :---: | :---: |
| $C^{\prime} A^{\prime}$ |  | X'C1' | 11000001 |
| C'l' | 1 | X'Fl' | 11110001 |
| C' ' | (blank) | X'40' | 01000000 |
| C'\#' | \# | X'7B' | 01111011 |
| C'@' | @ | $\mathrm{X}^{\prime} 7 \mathrm{C}^{\prime}$ | 01111100 |
|  | 3 | X'50' | 01010000 |
| c'r' | 1 | X'7D' | 01111101 |
| C'L'A' | L'A | X'D37DCl' |  |
| C'''י' | " | X'7D7D' |  |
| C'EOUR' | FOUR | $\mathrm{X}^{\prime} \mathrm{C} 6 \mathrm{D} 6 \mathrm{E} 4 \mathrm{D} 9^{\prime}$ |  |

The assembler assembles the data specified in a literal into a "literal pool" (fully described in H1B). It then assembles the address of this literal data in the pool into the object code of the instruction that contains the literal specification. Thus the assembler saves you a programming step by storing your literal data for you. The assembler also organizes literal pools efficiently so that the literal data is aligned on the proper boundary alignment and occupies the minimum amount of space.

LITERALS, CONSTANTS, AND SELFDEFINING TERMS: Do not confuse literals with constants or selfdefining terms. They differ in three important ways:

1. In where you can specify them in machine instructicns, that is, whether they represent data or an address of data.
2. In whether they have relocatable or aksolute values.
3. In what is assembled into the object code of the machine instruction in which they appear.

The figure to the right illustrates the first two points.
(1) A literal represents data.

- A constant is represented ky

2 its relocatable address. Note that a symbol with an absolute value does not refresent the address of
(3) a constant, but represents immediate data (see D5D) or an absclute
4 address.

- A self-defining term represents

5 data and has an absclute value.

## Compare:



A Literal with a self-defining term and a symbol with an absolute value


A symbol having an absolute address value with a self-defining term


The figure to the right illustrates the third point.
(1) The address of the literal, rather than the literal data itself is assembled into the cbject code.
2. The address of a constant is assembled intc the object code. Note that when a symbol with an
3 absolute value represents immediate data, it is the absolute value that is assembled into the cbject code.

4 - The absolute value of a selfdefining term is assembled into the object code.


## Specifications

A literal must be coded as shown in the figure to the right.
(1) The literal is specified in the same way as the operand of a DC instruction (for restrictions see G3C) -

GENERAL RULES FOR LITERAL USAGE: A literal is not a term and can be specified only as a complete operand in a machine instruction. In instructions with the RX format they must not be specified in operands in which an index register is also specified.

Because literals provide "read-only" data, they must not be used:

1. In operands that represent the receiving field of an instruction that modifies storage
2. In any shift or $1 / O$ instructions.

## C6 - Expressions

C6A -- PURPOSE

You can use an expression to specify:
(1) An address
(2) An explicit length
(3) A modifier
(4) A duplication factor
(5) A complete cperand

You can write an expression with a simple term or as an arithmetic combination of terms. The assembler reduces multiterm expressions to single values. Thus, you do not have to compute these values yourself.


Expressions have absolute or relocatable values. Whether an expression is absolute or relocatable depends on the value of the terms it contains. You can use the absclute or relocatakle expression described in this subsection in a machine instruction or any assembler instruction cther than a conditional assembly instruction. The assembler evaluates relocatable and absolute expressions at assembly time. Throughout this manual, the word "expression" refers to these types of expression.

NOTE: There are three types of expression that you can use only in conditional assembly instructions: arithmetic, logical, and character expressions. They are evaluated at pre-assembly time. In this manual they will always be referred to by their full names; they are described in detail in L4.

The figure below defines both absolute and relocatable expressions.NOTE: The relocatable values that are paired must have the opposite sign after the resolution of all unary
2 operators.



## Source Module

| FIRST | CSECT |  |
| :---: | :---: | :---: |
| A | ${ }_{\text {DC }}$ | $F^{\prime} 2^{\prime}$ |
| B | DC | $\mathrm{F}^{\prime \prime}$ |
| C | DC | $F^{\prime \prime}{ }^{\prime}$ |
| ABSA | EQU | 100 |
| ABSB | EQU | X'FF' |
| ABSC | EQU | B-A |
| ABSD | - | Paired |
|  | EQU | \%-A |
|  | END |  |

1. 2. Is composed of a symbol with an absolute value, a self-defining term, or a symbol length attritute

Absolute
The assembler reduces an absolute

Expressions
(3) 2. If it contains relocatable terms, alone or in combination with absolute terms, and if all these relocatable
reference, or any arithmetic expression to a single absolute value if the expression:

> combination of absolute terms.
(4) terms are paired.

## Absolute Expressions

ABSA

1) 1

L'A
(2) $\{$ ABSA $+A B S C-A B S C * 15$

3


A relocatable expression is one whose value changes, for example. ky a 1000, if the cbject module into which it is assembled is relocated 1000 bytes away from its originally assigned storage area. The assembler reduces a relocatable expression to a single relocatable value if the expression:

1. Is composed of a single relocatable term, or
2. Contains relocatable terms, alone or in combination with absolute terms, and:
a. All the relocatable terms
but one are paired. Note that
the unpaired term gives the
expression a relocatable value;
the paired relccatable terms
and other absolute terms
constitute increments or
decrements to the value of the
unpaired term.
b. The relocatakility attrikute of the whole expression is that of the unpaired term.
c. The sign preceding the unpaired relocatable term must be positive, after all unary operators have been resolved.

COMPLEX RELOCATABLE EXPRESSIONS: Complex relocatable expressions, unlike relocatable expressions, can contain:
a. Two or more unpaired relocatable terms or
b. An unpaired relocatable term preceded by a negative sign.

Complex relocatable expressions can be used only in A-type and $Y$ type address constants (see G3J).

The rules fcr coding an absolute or relocatable expression are:

1. Both unary (ocerating on one value) and binary (operating on two values) operators are allowed in expressions.
(1)
2. An expression can have one or more unary operators preceding any term in the expression or at the beginning of the expression.

2 3. An expression must not begin with a binary operatcr, nor can it contain two binary oper ators
(3) in successicn.
4. An expression must not contain two terms in succession.
5. No blanks are allowed between an operator and a term nor between two successive operators.
6. An expression can contain up to 19 unary and binary operators and up to 6 levels of parentheses. Note that parentheses that are part of an operand specification do not count toward this limit.
. An expression can contain up to 15 unary and binary operators and up to 5 levels of parentheses.
7. A single relocatable term is not allowed in a multiply or divide operation. Note that paired relocatable terms have absolute values and can be multiplied and divided if they are enclosed in parentheses.

8. A literal is not a valid term and is therefore not allowed in an expression.

The assembler reduces a multiterm expression to a single value as follows:

1. It evaluates each term.
2. It performs arithmetic operations from left tc right. Hcwever:
(1) a. It performs unary operations before binary crerations, and
b. It performs the binary
(3) operations of multiplication and division befcre the binary operations of addition and subtraction.
(4) 3. In division, it gives an integer result; any fractional portion is dropped. Livision by zero gives 0 .

3. In parenthesized expressions, the assembler evaluates the inner most expressions first and then considers them as terms in the next outer level of expressions. It 3 continues this process until the outermost expression is evaluated.
4. A term or expression's intermediate value and computed result must lie in the range of $-2^{31}$ through $+2^{31-1}$.
5. The computed result is then

Dos truncated tc a 24-bit value that lies between 0 and $16,777,215$.

NOIE: It is assured that the assembler evaluates paired relccatable terms at each level of expression nesting.


# Part II: Functions and Coding of Machine Instructions 

 SECTION D: MACHINE INSTRUCTIONS
## Section D: Machine Instructions

This section introduces the main functions of the machine instructions and provides general rules for coding them in their symbolic assembler language format. For the complete specifications of machine instructions, their object code format, their coding specifications, and their use of registers and virtual storage (see GLOSSARY) areas see the Principles of Operation manuals:

- IEM System/360 Principles of Operation, Order No. GA226821
- IBN System/370 Principles of Operation, Order No. GA227000


## D1 - Functions

```
At assembly time, the assembler converts the symbolic `
assembler language representation of the machine
instructions to the corresponding object code. It is this
object code that the computer processes at execution time.
Thus, the functions described in this section can be called
execution time functions.
Also at assembly time, the assembler creates the object code of the data constants and reserves storage for the areas you specify in your EC and LS assembler instructions (see G3). At execution time, the machine instructions can refer to these constants and areas, kut the constants themselves are not executed.
```


## Purpose

You use fixed-point instructions when you wish to perform arithmetic operations on data represented in binary form. These instructions treat all numbers as integers. If they are to operate upen data representing mixed numbers (such as 3.14 and 0.235 ) ycu must keep track of the decimal point yourself. For your constants you must provide the necessary number of binary positions to represent the fractional portion of the number specified ky using the scale mcdifier (see G3B) -

## Operations Performed

Fixed-point instructions allow you to perform the operations listed in the figure to the right.

## Data Constants Used

In fixed-point instructions, you can refer to the constants listed in the figure to the right.

NOTE: Except for the conversion operations, fixed-point arithmetic is performed on signed binary values.

| Fixed - Point Operations | Mnemonic Operation Codes |
| :---: | :---: |
| Add | AR, A, AH, ALR, AL |
| Subtract | SR, S, SH, SLR, SL |
| Multiply | MR, M, MH |
| Divide | DR, D |
| Arithmetic Compare (taking sign into account) | CR, C, CH |
| Load into registers | LR, L, LH, LTR, LCR, LPR, LNR, LM |
| Store into areas | ST, STH, STM |
| Arithmetic Shift of binary contents of registers to left or right (retaining sign) | SLA, SRA, SLDA, SRDA |
| Convert (packed) decimal data to binary | CVB |
| Convert binary data to (packed) decimal data | CVD |
| Constants Used | Type |
| Fixed-Point | H and F |
| Binary | B |
| Hexadecimal | X |
| Character | C |
| Decimal (packed) | P |
| Address | Y, A, S, V and Q |

## Furpose

You use the decimal instructions when you wish to perform arithmetic cperations on data that has the binary equivalent of decimal representaticn, either in packed or zoned form. These instructions treat all numbers as integers. For example, 3.14, 31.4, and 314 are all processed as 314. You must keep track of the decimal pcint yourself.

## Operaticns Perfcrmed

Decimal instructicns allow you to perform the operations listed in the figure tc the right.

## Lata Constants Used

In decimal instructions ycu can refer tc the constants listed in the figure to the right.

NOTE: Exceft for the conversion operations, decimal arithretic is performed on signed packed decimal values.

| Decimal Operations | Mnemonic Operation Codes |
| :---: | :---: |
| Add | AP |
| Subtract | SP |
| Multiply | MP |
| Divide | DP |
| Arithmetic Compare (taking sign into account) | CP |
| Move decimal data with a 4-bit offset | mvo |
| Shift decimal data in fields to left or right | SRP |
| Set a field to zero and add contents of another field | ZAP |
| Convert zoned to packed decimal data <br> Convert packed to zoned decimal data | PACK <br> UNPK |
| Constants Used | Type |
| Decimal (packed) <br> (zoned) | P $Z$ |

Furpose

You use floating-point instructicns when you wish tc ferform arithmetic operations on binary data that represents kcth integers and fractions. Thus, you do not have to keep track cf the decimal foint in your corrcutaticns. Flcatingpoint instructions also allow you to ferfcrm arithmetic cperations on both very large numbers and very small numbers, with greater precision than with fixed-point instructions.

## Operaticns_Performed

Floating-fcint instructicns allow you to perform the operations listed in the figure tc the right.

Cata Constants Used

In floating-point instructions, you can refer tc the constants listed in the figure to the right.

NOTE: Flcating-Fcint arithmetic

| Floating - Point Operations | Mnemonic Operation Codes |
| :---: | :---: |
| Add 1 | ADR, AD, AER, AE, AWR AW, AUR, AU, AXR |
| Subtract | SDR, SD, SER, SE, SWR, SW, SUR, SU, SXR |
| Multiply | MDR, MD, MER, ME, MXR, MXDR,MXD |
| Divide | DDR, DD, DER, DE |
| Halve (division by 2) | HDR, HER |
| Arithmetic Compare (taking sign into account) | CDR, CD, CER, CE |
| Load into floating point registers | LDR, LD, LER, LE, LTDR, LTER, LCDR, LCER, LPDR, LPER̄,LNDR,LDER,LRDR, LRER |
| Store into areas | STD, STE |
| Constants Used | Type |
| Floating - Point | $E, D$ and $L$ | is performed on signed values that must have a special floating-foint format. The fracticnal pcrtion of floating-point numbers, when 1 used in addition and subtraction, can have a normalized (no leading zercs) cr unncrmalized format.

## Purpose

You can use the logical instructions to introduce data, move data, or inspect and change data.

## Operations Performed

The logical instructions allow you to perform the operations listed in the figure to the right.

| Logical Operations | Mnemonic Operation Codes |
| :---: | :---: |
| Move | MVI, MVC, MVN, MVZ, MVCL |
| Logical Compare (unsigned binary values) | $\begin{aligned} & \text { CLR, CL, CLI, CLC, CLCL, } \\ & \text { CLM } \end{aligned}$ |
| AND (logical multiplication) | NR, N, NI, NC |
| OR (logical addition) | OR, O, OI, OC |
| Exclusive OR (either........ or but not both) | XR, X, XI, XC |
| Testing binary bit patterns | TM |
| Inserting characters into registers | IC, ICM |
| Store characters into areas | STC, STCM |
| Load address into register | LA |
| Logical Shift of unsigned binary contents of registers to left or right | SLL, SRL, SLDL, SRDL |
| Replace argument values by corresponding function values from table (translate) | TR, TRT |
| Edit (packed and zoned decimal data) values in preparation for printing | ED, EDMK |

## Purpose

You can use several types of branching instructions, combined with the logical instructions listed in D1D, to code and control loops, subroutine linkages, and the sequence of processing.

## Operations Performed

The kranching instructions allow you to perform the operations listed in the figure to the right.

## NOTE: Additional mnemonics for

kranching on condition are described in section I 1 H below.

| Branching <br> Operations | Mnemonic Operation <br> Codes |
| :--- | :--- |
| Branch depending <br> on the results of <br> the preceding <br> operation (that <br> sets the condition <br> code) | BCR, BC |
| Branch to a <br> subroutine with a <br> return link to <br> current code | BALR, BAL |
| Branch according <br> to a count con- <br> tained in a register <br> (count is decremented <br> by one before deter- <br> mining course of <br> action) | BCTR, BCT |
| Branch by comparing <br> index value to fixed | BXH, BXLE |
| comparand, (index <br> incremented or de- <br> cremented before <br> determining course <br> of action) | Temporary Branch in <br> order to execute a <br> specific machine <br> instruction |
| EX |  |

## E1F -- STATUS SWITCHING

## Purpose

You can use the status switching instructions to communicate between your program and the system control program. However, some of these instructions are privileged
instructions and you can use them only when the CPU is in the supervisor state, but not when it is in the problem state. The privileged instructions are marked with a "p" in the figure to the right.

## Operations Performed

The status switching instructions allow you to perform the operations listed in the figure to the right.


## Purpose

You can use the input/output instructions, instead of the IBMsupplied system macro instructions, when you wish to control your input and output cperations more closely.

## Operations Ferformed

The input or output instructions allow you to identify the channel, or the device on which the input or output operation is to be performed. The operations performed are listed in the figure to the right. However, these are privileged instructions, and you can only use them when the CPU is in the supervisor state, but not when it is in the problem state.

| Input or Output <br> Operations | Mnemonic Operation <br> Codes |
| :--- | :--- |
| Start 1/O | SIO, SIOF |
| Halt I/O | HIO |
| Test state of channel <br> or device being used <br> Halt Device | TIO,TCH |

## Purpose

The branching instructions described below allow you to specify a mnemonic code for the condition on which a kranch is to occur. Thus, you avoid having to specify the mask value required by the EC and ECR branching instructions. The assembler translates the mnemonic code that represents the condition into the mask value, which is then assembled in the object code of the machine instruction.

## Specifications

The extended mnemonic codes are given in the figure on the opposite page.

They can be used as operation codes for kranching
1 instructions, replacing the $B C$ and $B C R$ machine instruction codes. Note that the first operand of the BC and BCR
3 instructions must not be present in the operand field of the extended memonic branching instructions.

4 NOTE: The addresses represented are explicit addresses; however, implicit addresses can also be used in this type of instruction.


Used After Arithmetic Instructions

| BO | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | Branch on Overflow | RX | BC | 1, D2 ( $\mathrm{X} 2, \mathrm{~B} 2)$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| BOR | R2 |  | RR | BCR | 1,R2 |
| BP | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | Branch on Plus | RX | BC | 2,D2 (X2, B2) |
| BPR | R2 |  | RR | BCR | 2,R2 |
| BM | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | Branch on Minus | RX | BC | 4, D2 (X2, B2) |
| BMR | R2 |  | RR | BCR | 4,R2 |
| BNP | D2 ( $\mathrm{X} 2, \mathrm{~B} 2)$ | Branch on Not Plus | RX | BC | 13,D2 (X2, B2) |
| BNPR | R2 |  | RR | BCR | 13,R2 |
| BNM | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | Branch on Not Minus | RX | BC | 11,D2 (X2, B2) |
| BNMR | R2 |  | RR | BCR | 11,R2 |
| BNZ | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | Branch on Not zero | RX | BC | 7,D2 (X2, B2) |
| BNZR | R2 |  | RR | BCR | 7,R2 |
| BZ | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | Branch on Zero | RX | BC | 8,D2 (X2,B2) |
| B2R | R2 |  | RR | BCR | 8,R2 |
| BNO | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | Branch on No Overflow | RX | BC | 14, D2 (X2, B2) |
| BNOR | R2 |  | RR | BCR | 14,R2 |

Used After Test Under Mask Instructions

| BO | D2 ( $\mathrm{X} 2, \mathrm{~B} 2)$ | Branch if | Ones | RX | BC | 1,D2 (X2,B2) |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| BOR | R2 |  |  | RR | BCR | 1, R2 |
| BM | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | Branch if | Mixed | RX | BC | 4,D2 (X2, B2) |
| BMR | R2 |  |  | RR | BCR | 4,R2 |
| BZ | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | Branch if | Zeros | RX | BC | 8,D2 (X2,B2) |
| BZR | R2 |  |  | RR | BCR | 8,R2 |
| BNO | D2 ( $\mathrm{X} 2, \mathrm{~B} 2)$ | Branch if | Not Ones | RX | BC | 14,D2 (X2, B2) |
| BNOR | R2 |  |  | RR | BCR | 14,R2 |
| BNM | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | Branch if | Not Mixed | RX | BC | 11,D2 (X2, B2) |
| BNMR | R2 |  |  | RR | BCR | 11,R2 |
| BNZ | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) | Branch i | Not Zeros | RX | BC | 7,D2 (X2, B2) |
| BNZR | R2 |  |  | RR | BCR | 7,R2 |

D2=displacement, $\mathrm{X} 2=$ index register, $\mathrm{B} 2=$ base register, $\mathrm{R} 2=$ register containing branch address

## Purpose

You use the relocation instructions in connnection with the relocate feature of IBM System/370.

## Operations Performed

The relocation instructions allow you to perform the crerations listed in the figure to the right. However, these instructions are privileged instructions, and you can use them only when the CPU is in the supervisor state, but not when it is in the problem state.

| Relocation <br> Operations | Mnemonic Operation <br> Code |
| :--- | :--- |
| Load Real Address | LRA |
| Purge Translation <br> Lookaside Buffer | PTLB |
| Reset Reference Bit | RRB |
| Set Clock Comparator | SCKC |
| Store Clock Comparator | STCKC |
| Set CP U Timer | STPT |
| Store CP U Timer | STNSM |
| Store and AND System <br> Mask | STOSM |
| Store and OR System <br> Mask |  |

## Purpose

The assembler automatically aligns the object code of all machine instructions on halfword toundaries. For execution of the IBM System/ 370 machines, the constants and areas do not have to lie on specific toundaries to be addressed by the machine instructions.

However, if the assembler option ALIGN is set, you can cause the assembler to align constants and areas; for example, on fullword boundaries. This allows faster execution of the fullword machine instructions.

If the NOALIGN option is set, you do not need to align constants and areas. They will ke assembled at the next available byte, which allows you to save space (no bytes are skipped for alignment).

## Specifications

MACHINE INSTRUCTIONS: When the assembler aligns machine instructions on halfword boundaries, it sets
any bytes skipped to zero.
CONSTANTS AND AREAS: One of the assembler options that can be set in the job control language (that initiates execution of the assembler program) concerns the alignment of constants and areas; it can be specified as ALIGN or NOALIGN.

If AIIGN is specified, the following applies:

2 - The assembler aligns constants and areas on the boundaries implicit in their type, if no length specification is supplied.

- The assembler checks all

3 expressions that represent storage addresses to ensure that they are aligned on the boundaries required by the instructions. If they are not, the assembler issues a warning message.

If NCALIGN is specified, the following applies:

- The assembler does not align constants and areas on special boundaries, even if the length specification is cmitted. Note that the CCW instruction, however, always causes the alignment of the channel command word on a doubleword boundary.
- The assembler does not check storage addresses for boundary alignment.

NOTE 1: The assembler always forces alignment if a duplication factor of 0 is specified in a constant or area without a length modifier (for an example, see G3N). Alignment occurs when either ALIGN or NOALIGN is set.


NOTE 2: When NOALIGN is specified, the CNOF assembler instruction can be used to ensure the correct alignment of data referred to ky the privileged instructions that require specific boundary alignment. The mnemonic operation codes for these instructions are listed in the figure to the right.

| Mnemonic Operation Codes for Privileged Operations | Meaning |
| :---: | :---: |
| LPSW | Load program status word |
| ISK | Insert Storage Key. |
| SSK | Set Storage Key. |
| LCTL | Load Control registers. |
| SCK | Set Clock. |
| STIDP | Store CPU Identification |
| STCTL | Store Control registers. |
| - |  |
| (Diagnose - not handled by assembler) |  |

## D3 -- Statement Formats

Machine instructicns are assembled into object code according to one of the six formats given in the figure to the right.

When you code machine instructions you use symbolic formats that correspond to the actual machine language formats. Within each tasic format, you can also code variations of the symbolic representation (Examples of coded machine instructions, divided into groups according to the six basic formats, are illustrated in 56 below).

The assembler converts only the operation code and the operand entries of the assembler language statement into object code. The assembler assigns tc the symbcl you code as a name entry the value of the address of the leftmost
3 byte of the assembled instruction. When you use this same symbol in the operand of an assembler language statement, the assembler uses this address value in converting the symbolic operand intc its object code form. The length attribute assigned to the symbcl depends on the basic machine language format of the instruction in which the symbol appears as a name entry (for details on the length attribute see C4C) .

5 A remarks entry is not converted into object code.

| Format | Length of Object Code <br> Reguired for the Assembled <br> Instruction in Bytes |
| :--- | :--- |
| RR | 2 |
| RX | 4 |
| RS | 4 |
| SI | 4 |
| S | 4 |
| SS | 6 |

Example:
Assembler Language Statement


## Object Code

(machine language) ot Assembled Instruction in Hex

## D4 - Mnemonic Operation Codes

## Furpose

You must specify an cperation code for each machine instruction staterent. The mnemonic cperation code indicates the type of cperation to ke perfcrmed; for example, "A" indicates the "addition" cperaticn. Appendix IV contains a complete list of mnemonic operation codes and the fcrmats of the corresponding machine instructions.

## Specificaticns

The general format of the machine instruction cFeration code is shown in the figure to the right.
(1) The verk must always be present. It usually consists of on $\in$ or two characters and specifies the operation to ke performed. The cther iters in the cperation code are not always present. They include:
(2) The modifier which further defines the cperaticn
(3) The type qualifier, which indicates the type of data used ky the instruction in its operation, and
(4*) The fcrmat gualifier, $R$ cr $I$, which indicates that an RR or SI machine instruction fcrmat is assembled.

VERB

Examples:


## D5 -- Operand Entries

Furpose

You must specify cne cr mare operands in each machine instruction statenent to ercvide the data cr the location of the data upon which the machine operation is to be ferformed. The operand entries consist of one cr more fields cr subfields depending on the format of the instruction keing coded. They can specify a register, an address, a length, and immediate data.

You can code an operand entry either with symbcls cr with self-defining terms. You can omit length fields or subfields, which the assembler will compute for you from the ctrer cperand entries.

General Specifications for Coding Operand Entries

The rules for coding ocerand entries are as fcllcws:
(1)

A comma must separate operands.
2 Farentheses must enclcse subfields.
(3) A comma must separate subfields enclosed in parentheses.

If a subfield is critted because it is in implicit in a symbclic address, the farentheses that would 4 have enclosed the sukfield must ke critted.


If twc sukfields are enclcsed in parentheses and separated ky commas, the fcllcwing apflies:

If both subfields are oritted kecause they are implicit in a symbclic
(1) the parentheses that hould have been needed must also $k \in$ omitt $\in d$.

If the first subfield is cmitted,
the comma that separates it from
the second subfield must be written as well as the enclosing parentheses.

If the seccra subfield is omitted, the comma that separates it from the first subfield must be omitted, however, the enclosing farentheses must te written.

NOTE: Elanks must not appear within the cferand field, except as part of a character self-defining $t \in r \pi$ or in the specification cf a

## (5)

 character literal.

Purpose and Usage

You can specify a register in an operand for use as an arithmetic accumulator, a base register, an index register, and as a general depository for data to which you wish to refer over and over.

You must be careful when specifying a register whose contents have been affected by the execution of another machine instruction, the control program, or an IEM-supplied system macro instruction.

For some machine instructions you are limited in which registers you can specify in an operand.

## Specifications

The expressions used to specify registers must have absolute values; in general, registers 0 through 15 can be specified for machine instructions. However, the following restrictions on register usage apply:

1. The floating-point registers
(0, 2, 4, or 6) must be specified
for floating-point instructions:
2. The even numbered registers $(0$,

2, 4, 6, 8, 10, 12, 14) must be specified for the following groups of instructions:
a. The double-shift instructions

3 b. The fullword multiply and divide instructions

4 C. The move long and compare logical long instructicns.
3. The floating-point registers 0 and 4 must be specified for the instructions that use extended
(5) floating-point data:

AXR, SXR, LRDR, MXR, MXDR, MXL.
NOTE: The assembler checks the registers specified in the instruction statements of the above groups. If the specified register does not comply with the stated restrictions, the assembler issues a diagnostic message and does not assemble the instruction.

REGISTER USAGE BY MACHINE
INSTRUCTIONS: Registers that are not explicitly coded in the symbolic assembler language representation of machine instructions, but are nevertheless used by the assembled machine instructions, are divided into two categories:

1. The base registers that are implicit in the symbolic addresses specified. These implicit addresses are described in detail in D5B.
The registers can be identified
by examining the object ccde of the assembled machine instruction 2 or the USING instruction(s) that assigns base registers for the source module.
2. The registers that are used ky machine instructions in their operations, but do not appear even in the assembled object ccde. They are as follows:
a. For the double shift and fullword multiply and divide instructions, the odd-numbered register whose number is one greater than the even-numbered register specified as the first operand.
b. For the Move Long and Compare Logical Long instructions, the odd-numbered reqisters whose number is one greater than the even numbered registers specified in the two operands.
c. For the Branch on Index High (BXH) and the Branch on Index Low or Equal (EXIE) instructions; if the register specified for the second operand is an evennumbered register, the next
higher odd-numbered register is used to contain the value to be used for comparison.
d. For the Translate and Test (TRT) instruction, registers 1 and 2 are also used.
e. For the Load Multiple (LM) and Store Multiple (STM)
instructions, the registers that lie between the registers specified in the first two operands.

REGISTER USAGE BY SYSTEM: The control program of the IBM System/370 uses registers $0,1,13,14$, and 15.


Purpose and Definition

You can code a symbol in the name field of a machine instruction statement to represent the address of that instruction. Ycu can then refer to the symbol in the operands. of other machine instruction statements. The object code for the IBM System/370 requires that all addresses be assembled in a numeric kase-displacement format. This format allows you to specify addresses that are relocatable or absolute.

You must not confuse the concept of relocatakility with the actual addresses that are coded as relocatable, nor with the format of the addresses that are assembled.

EEFINING SYMBOLIC ADDRESSES: You define symbols to represent either relocatable or absolute addresses. You can define relocatable addresses in two ways:

By using a symbol as the label in the name field of an assembler language statement or

By equating a symbol to a relocatakle expression.

You can define absolute addresses (or values) by equating a symbcl to an absolute expression.

REFERRING TC ADLRESSES: You can refer to relocatable and absolute addresses in the operands of machine instruction statements. Such address references are also called addresses in this manual. The two ways of coding addresses are:
(4) Implicitly: that is, in a form that the assembler must first convert into an explicit base-displacement form before it can be assembled into object ccde.
(5) Explicitly: that is, in a form that can be directly assembled into object code.


## Relocatability of Addresses

Addresses in the base-displacement form are relocatable, because:

- Each relocatable address is assembled as a displacement from (2) a base address and a base register.
- The base register contains the kase address.
- If the object module assembled from your source module is relocated, only the contents of the base register need reflect this relocation. This means that the location in virtual storage of your base has changed and that your base register must contain this new base address.
- Your addresses have been assembled as relative to the base address; therefore, the sum of the displacement and the contents of the kase register will point to the correct address after relocation.

NOTE: Absolute addresses are also assembled in the base-displacement form, but always indicate a fixed location in virtual storage. This means that the contents of the base register must always be a fixed absolute address value regardless of relocation.


MACHINE OR OBJECT COLE FORMAT: All addresses assembled into the object code of the IBM System/370 machine

Addresses instructions have the format given in the figure kelow.


The addresses represented have a value which is the sum of:
(1) A displacement and
2) The contents of a base register.

NOTE: In RX instructions, the address represented has a value which is the sum of a displacement, the contents
3 of a base register, and the contents of an index register.

## Implicit Address

An implicit address is specified by coding one expression. The expression can be relocatable or absolute: The assembler converts all implicit addresses into their
(3) kase-displacement form before it assembles them into object code. The assembler converts implicit addresses into explicit addresses only if a USING instruction has been specified. The USING instruction assigns both a base address, from which the assembler computes displacements, and a base register, to contain the base address. The base register must ke loaded with the correct base address at execution time. For details on how the USING instruction is used when establishing addressability, thus allowing implicit references, see $F 1$.

## Explicit Address

An explicit address is specified
by coding two absolute expressions
as follows:
The first is an absolute expression
for the displacement, whose value must lie in the range 0 through 4095 (4095 is the maximum value that can be represented by the 12 binary bits available for the displacement in the object code).

The second (enclosed in parentheses) is an absolute expression for the 2 base register, whose value must lie in the range 0 through 15.

If the base register contains a value that changes when the program is relocated, the assembled address is relocatakle. If the base register contains a fixed absolute value that is unaffected by program relocation, the assembled address is aksolute.

NOTES (for implicit and explicit addresses) :

1. An explicit base register designation must not accompany an implicit address.
2. However, in $R X$ instructions an index reqister can be coded with an implicit address as well as with an explicit address.
3. When two addresses are required, one address can be coded as an
explicit address and the cther as an implicit address.


## Purpose

You can specify the length field in an SS-type instruction. This allows you to indicate explicitly the number of bytes of data at a virtual storage location that is to be used by the instruction. However, you can omit the length specification, because the assembler computes the number of bytes of data to be used from the expression that represents the address of the data.

## Specifications

IMPLICIT LENGTH: When a length subfield is omitted from an SS-type machine instruction an implicit length is assembled into the object code of the instruction. The implicit length is either of the following:

1. For an implicit address (see D5B above), it is the

length attribute of the first or only term in the expression representing the implicit address.
2. For an explicit address (see D5B above), it is the 2 length attribute of the first or only term in the expression that represents the displacement.

For details on the length attribute of symbols and other terms see C4C.

EXPLICIT LENGTH: when a length sutfield is specified in 3 an SS-type machine instruction, the explicit length thus defined always overrides the implicit length.

NOTES:
4. An implicit or explicit length is the effective length. The length value assembled is always one less than the effective length. If an assembled length value of 0 is 5 desired, an explicit length of 0 or 1 can be specified.
2. In the $S S$ instructicns requiring one length value, the allowable range for explicit lengths is 0 through 256. In the $S S$ instructions requiring two length values, the allowable range for explicit lengths is 0 through 16.


Purpose

In addition to addresses, registers, and lengths, some machine instruction operands require immediate data. Such data is assembled directly into the object code of the machine instructions. You use immediate data to specify the bit patterns for masks or other absolute values you need.

You should be careful to specify immediate data only where it is required. Lo not confuse it with address references to constants and areas or with any literals you specify as the operands of machine instruction (for a compariscn between constants, literals, and immediate data, see C5).

## Specifications

Immediate data must be specified as absolute expressions whose range of values depends on the machine instruction for which the data is required. The immediate data is assembled into its 4-bit or 8-bit Linary representaticn 2 according to the figure on the opposite page.

| Machine Instructions in which immediate data is required (Op codes in Appendix IV) | Range of Values allowed for immediate data | Examples Object Code <br> in Hex |
| :---: | :---: | :---: |
| SRP (SS) | 0 through 9 | $\begin{aligned} & \text { SRP A, B, 3, } 1 \\ & \qquad \begin{array}{\|l\|l\|l\|l\|l\|} \hline \text { F0 } & 7 & 3 & \mathrm{xxxx} & \mathrm{xxxx} \\ \hline \end{array} \\ & \hline \begin{array}{l} \text { Length of } \\ \text { Field } \mathrm{A} \end{array} \\ & \hline \end{aligned}$ |
| $\begin{array}{lll} \text { All } & \mathrm{BCR} & \text { (RR) } \\ \text { All } & \mathrm{BC} & \text { (RX) } \end{array}$ | 0 through 15 0 through 15 | BCR 8,3 <br> $\cdot 1$07 8 3 BC 11, AAA $\square$$\square$ <br> 4747 B 0 <br> Address |
| ICM $\quad$ (RS) STCM CLM | 0 through 15 |  |
| NI (SI) <br> CLI  <br> XI  <br> MVI  <br> OI  <br> TM  <br> RDD  <br> WRD  | 0 through 255 |  |
| SVC (RR) | 0 through 255 | $\text { SVC } \quad 128$ $\square$ <br> 0A <br> 80 <br> 2 |

## D6 - Examples of Coded Machine Instructions

```
The examples in this subsection
are grouced according to machine
instruction format. They illustrate
the varicus ways in which you can
code the operands of machine
instructions. Both symbolic and
numeric representation of fields
and subfields are shown in the
examples. You must therefore assume
that all symbcls used are defined
elsewhere in the same source module.
The okject ccde assembled from at
least one coded statement per group
is also included. A complete summary
of machine instruction formats with
the coded assembler language variants
can be found in Appendix III and
IV.
```

RR Format

You use the instructions with the RR format mainly to move data tetween registers. The operand fields must thus designate registers, with the following exceptions:

In BCR branching instructions when
(1) a 4 -kit kranching mask replaces the first register specification

In SVC instructions, where an
2 immediate value (between 0 and 255) replaces both registers.

3 NOTE: Symbcls used in RR
instructions are assumed to be equated tc absclute values ketween 0 and 15.

RX Format

You use the instructions with the RX format mainly to move data between a register and virtual storage. By adjusting the contents of the index register in the RX-instructions you can change the location in
virtual storage being addressed. The operand fields must therefore designate registers, including index registers, and virtual storage addresses, with the following exception:

1. In BC branching instructions a 4kit kranching mask, with a value between 0 and 15, replaces the first register specification.

NOTES:

1. Symbols used tc represent

2 registers are assumed to be equated to absolute values between 0 and 15.
(3)
2. Symbols used to represent implicit addresses can be either relocatakle or absolute.
3. Symbols used to represent

4 displacements in explicit addresses are assumed to be equated to absolute values between 0 and 4095.


## RS Format

You use the instructions with the RS format mainly to move data between one or more registers and virtual storage or to compare data in one or more registers (see the BXH and BXLE operations in Appendix IV).

In the Insert Characters under Mask (ICM) and the Store Characters Under
© Mask (STCM) instructions, when a 4-bit mask, with a value between 0 and 15 , replaces the second register specification.

NOTES:

1. Symbols used to represent
 registers are assumed to be equated to absolute values between 0 and 15.

2. Symbols used to represent implicit addresses can be either relocatable or absolute.
3. Symbols used to represent
displacements in explicit addresses are assumed to be equated to absolute values between 0 and 4095.


You use the instructions with the SI format mainly to move immediate data into virtual stcrage. The operand fields must therefore designate immediate data and virtual storage addresses, with the following exception:
(1) An immediate field is not needed in the statements whose operation codes are: LPSW, SSM, TS, TCH, and TIO.

NOTES:

1. Symbols used to represent

2 immediate data are assumed to be equated to absolute values ketween 0 and 255.
(3)
2. Symbols used to represent implicit addresses can be either relocatable or absolute.
4. Symbols used to represent

4 displacements in explicit addresses are assumed to be equated to absolute values between 0 and 4095.


Assembly Examples:
Assembler Language Statement
Object Code of Machine Instruction In Hex

SI Format


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|  | Mnemonic <br> Operation <br> Codes | Assembled <br> Operation <br> Code in |
| :--- | :--- | :--- | :--- |
| Hex |  |  |$\quad$| Description |
| :--- |

You use the instructions with the SS format mainly to move data between two virtual storage locations.
The operand fields and subfields must therefore designate virtual storage addresses and the explicit data lengths you wish to include. However, note the following exception:

In the Shift and Round Lecimal (SRP)
(1)
instruction a 4-bit immediate data
field, with a value between 0 and 9, is specified as a third operand.

NOTES:
(2)

1. Symbols used to represent kase
registers in explicit addresses are assumed to be equated to absolute values between 0 and 15.
(3)
2. Symbols used to represent explicit lengths are assumed to be equated to absolute values between 0 and 256 for SS instructions with one length specification and ketween 0 and 16 for $S S$ instructicns with two length specifications.

4
3. Symbols used to represent implicit addresses can be either relocatable or absolute.
(5)
4. Symbols used to represent

5 displacements in explicit addresses are assumed to be equated to absolute values between 0 and 4095.

# Part III: Functions of Assembler Instructions 

SECTION E: PROGRAM SECTIONING
SECTION F: ADDRESSING
SECTION G: SYMBOL AND DATA DEFINITION
SECTION H: CONTROLLING THE ASSEMBLER PROGRAM

## Section E: Program Sectioning

This section explains how you can subdivide a large program into smaller parts that are easier to understand and maintain. It also explains how you can divide these smaller parts into convenient sections: for example, one section to contain your executable instructions and another section to contain your data constants and areas.

You should consider two different subdivisions when writing an assembler language program:

1. The source module
2. The control section.

You can divide a program into two or more source modules. Each source module is assembled into a separate object module. The object modules can then te combined into load modules to form an executable
program.

You can also divide a source module into two or more control sections.
(1) Each control section is assembled as part of an object module. By writing the proper linkage edit control statements, you can select 2 a complete object module or any individual control section of the 3 object module to be linkage edited and later loaded as an executable program.

SIZE OF PROGRAM PARTS: If a source module becomes so large that its logic is not easily comprehensible, break it up into smaller modules.

Unless you have special programming reasons, you should write each control section so that the resulting object code is not larger than 4096 bytes. This is the largest number of bytes that can be ccvered by one base register (for the assignment of base registers to control sections, see $F 1 \mathrm{~A}$ ).

## COMMUNICATION BETWEEN PROGRAM PARTS:

You must be able to communicate between the parts of your program: that is, be able to refer to data in a different part or be able to branch to another part.


To communicate between two or more source modules, ycu must symbolically link them together; symbolic linkage is described in $F 2$.

To communicate between two or more control sections within a source module, you must establish the addressability of each control section; establishing addressability is descriked in F 1 .

Source Mod.

## El -- The Source Module

A source module is composed of source statements in the assembler language. You can include these statements in the source module in two ways:

1. You write them on a coding form

1 and then enter them as input, for example, through a terminal or, using punched cards, through a card reader.
2. You specify one or more COPY instructions among the source statements being entered. When the assembler encounters a COPY instruction, it replaces the CCPY instruction with a predetermined set of source statements from a library. These statements then become a part of the source module.

## The Beginning of a Scurce Module



The first statement of a source module can be any assembler language statement, except MEXIT and MENL, that is described in this manual. You can initiate the first control section of a source module by using the START instruction. However, you can or must write some source statements before the beginning of the first contrcl section (for a list of these statements see E2L).

The End of a Source Module
os The END instruction usually marks
only The
(1) the end of a source module. However,
you can code several ENL
instructions. The assembler stops
assembling when it processes the
first END instruction. If no END
instruction is found, the assembler
will generate one.


E1A -- THE COFY INSTRUCTICN

## Purpose

1 The CCPY instruction allows you 1 to copy predefined source statements
(2) from a library and include them
in a source module. Ycu thereby avoid:

1. Writing the same, often-used sequence of code over and over
2. Keypunching and handling the punched cards for that code.


Specifications

The format of the COPY instruction statement is shown in the figure to the right.

The symbol in the operand field must identify a part of a library called:

A member of a partitioned data set

DOS

## A book in the source statement library

This member (or bock) contains the coded source statements to be copied.

The source coding that is copied into a source module:
(1). Is inserted immediately after the CCPY instruction

- Is inserted and processed according to the standard instruction statement coding format (described in B1D), even if an ICTL instruction has keen specified
- Must not contain either an ICTL or ISEQ instruction
(3) Can contain a COPY instruction. Up to 5 levels of nesting of the COPY instruction are allowed.

Dos Up to 3 levels of nesting are allowed.
(4)

- Can contain macro definitions (see Section J) -

If a source macro definition is copied into the beginning of a source module, both the MACRO and MEND statements that delimit the definition must be contained in the same level of copied code.

NOTES:

1. The COPY instruction can also be used to cory statements into source macro definiticns (see J5C).
2. The rules that govern the occurrence of assembler language statements in a source module also govern the statements copied into the source module.

| Name | Operation | Operand |
| :--- | :--- | :--- |
| Blank | COPY | One ordinary <br> Symbol |



|  | COPY | CODE2 |
| :---: | :---: | :---: |
|  | L <br> A <br> ST <br> B <br> COPY | 3, AREA <br> 3, FTWO <br> 3,SUM <br> PAST <br> CODE3 |
| $\begin{aligned} & \text { AREA } \\ & \text { FTWO } \\ & \text { SUM } \end{aligned}$ | DS <br> DC <br> DS | $\begin{aligned} & \mathrm{F}=1 \\ & \mathrm{~F}^{\prime} \mathbf{F}^{\prime} \end{aligned}$ |
| PAST | C | 4 , SUM |

$\stackrel{\rightharpoonup}{\bullet}$

END

## Furpose

You use the END instruction to mark the end of a source module. It indicates to the assembler where to stop assembly processing. You can also supply an address in the operand field to which control can be passed when your program is loaded. This is usually the address of the first executable instruction in a source module.

Specifications
END
The format of the ENL instruction statement is shown in the figure to the right.

If specified, the operand entry can be generated by substitution into variable symbols. However, after sutstitution, that is, at assembly time:

1. It must be a relocatable

1 expression representing an address in the source module delimited ky the END instruction, or
2) 2. If it contains an external symbol, the external symbol must be the only term in the expression, or
(3) the remaining terms in the expression must reduce to zero.
3. It must not be a literal.

| Name | Operation | Operand |
| :--- | :---: | :---: |
| A sequence <br> symbol or <br> blank | END | A relocatable <br> expression or <br> blank |

## Source Module A

| A | START | 0 |
| :---: | :---: | :---: |
| ENTERA | BALR | 12,0 |
|  | USING | *,12 |
|  | ENTRY | ENTERA |
|  | - |  |
|  | - |  |
|  |  | (1) |
|  | END | ENTERA |

Source Module B


This page left blank intentionally.

E2 - General Information About Control Sections
Contrl Sect.

A control section is the smallest subdivision of a prcgram that can be relocated as a unit. The assembled control sections contain the object code for machine instructions, data constants, and areas.

Consider the concept of a control section at different processing times.
(1) AT CODING TIME: You create a control section when you write the instructions it contains. In addition, you establish the addressability of each control section within the source module, and provide any symbolic linkages between control sections that lie in different source modules. You also write the linkage editor control statements to combine the desired control sections into a load rodule, and to provide an entry point address for the beginning of program execution.
(2) AT ASSEMBLY TIME: The assembler translates the source statements in the control section into object code. Each source module is assembled into one object module. The entire object module and each of the control sections it contains is relocatable.

AT LINKAGE EDITING TIME: According to linkage editor control statements, the linkage editor combines the object code of one or more control sections into one load module. It also calculates the linkage addresses necessary for communication between two or more control sections from different okject modules. In addition, it calculates the space needed to accommodate external dummy sections (see E4) -
(4) AT PROGRAM FETCH TIME: The control program loads the load module into virtual storage. All the relocatable addresses are converted to fixed locations in storage.

AT EXECUTION TIME: The control program passes control to the load module now in virtual storage and your program is executed.

NOTE: You can specify the relocatakle address of the starting point for program execution in a linkage editor control statement or in the operand field of an END statement.


## Executable Contrcl Sections

An executable control section is one you initiate by using the START
or CSECT instructions and is
assembled into object code. At
execution time, an executable control
section contains the binary data assembled from your coded
instructions and constants and is therefore executable.

An executable control section can also be initiated as "private code", without using the START or CSECT instruction (see E2E).

Reference Control Sections

A reference control section is one COM, or DXD instruction and is not assembled into object code. You can use a reference control section 3 either to reserve storage areas or to describe data to which you can refer from executable control sections. These reference control sections are considered tc be empty at assembly time, and the actual

4 binary data to which they refer is not entered until execution time.


The assembler maintains a separate location counter for each control section. The location counter setting for each control section starts at 0. The location values assigned to the instructions and other data in a control section are therefore relative to the location counter setting at the beginning of that control section.However, for executable control sections, the location values that appear in the listings do not restart at 0 for each subsequent executable control section. They carry on from the end of the previous contrcl section. Your executable control sections are usually loaded into storage in the order you write them. You can therefore match the source statements and object code produced from them with the contents of a dump of your program.

DOS For executable contrcl sections. the location values that appear in the listings always start from 0 , except the control section initiated by a START instruction with a non-zero operand entry.

3 For reference control sections, the location values that appear in the listings always start from (4) 0.



The specifications below apply to the first executable control section, and not to a reference control section.

INSTRUCTICNS THAT ESTAELISH THE FIRST CONTRCL SECTION: Any
instruction that affects the location counter or uses its current value establishes the beginning of the first executable control section. The instructions that establish the first control section are listed in the figure to the right.

The statements copied into a source module by a COPY instruction, if specified, determine whether or not it will initiate the first control section.

OS NOTE: The LSECT, COM, and LXD
only instructions initiate reference control sections and do nct establish the first executable control section.

First Contrl Sect.

WHAT MUST CCME BEFORE THE FIRST CONTROL SECTION: The fcllowing instructions or groups of instructions, if specified, must appear before the first control section, as shown in the figure to the right.

- The ICTL instructicn, which, if specified, must be the first statement in a source module


## OS - The OPSYN instruction

- Any source macro definitions (see J1B)
- The COPY instructicn, if the code to be copied contains only OPSYN instructions or complete macro definitions.

Source Module


WHAT CAN CFIICNALLY CONE EEFORE THE FIRST CONTROL SECIION: The instructions or groups of
instructions that can cftionally
be scecified before the first control section are shown in the figure to the right.

Any instructions copied by a COPY
instruction or generated by the
processing of a macro instruction before the first control section must kelong exclusively to one of the groups of instructions shown in the figure to the right.

NOTES:

1. The EJECT, ISEQ, PRINT, SPACE, or TITLE instructions and comments statements must follow the ICTL instruction, if specified. However, they can precede or appear between source macro definitions. The OPSYN instruction must (1) follow the ICTL instruction, if specified, and (2) precede any source macro definition specified.
2. All the other instructions of the assembler language must follow any source macro definitions specified.
3. All the instructions or groups of instructions listed in the figure to the right can also appear as part of a control section.

The unnamed control section is an executable control section that can be initiated in cne of the following two ways:

1. By coding a START or CSECT instruction without a name entry
2. By coding any instruction, other than the START or CSECT instruction, that initiates the first executable control section.

The unnamed control section is sometimes referred tc as frivate code.

All control sections ought to be provided with names so that they can te referred to symbolically:

1. Within a source module
2. In EXTRN and WXTRN instructions and linkage editor control statements for linkage between source modules.

NOTE: Unnamed common control sections or dummy control sections can be defined if the name entry is omitted from a COM or DSECT instruction.

Only unnamed common control sections (initiated by the COM instruction) and named dummy control sections (initiated by the ESECT instruction) are allowed.

## E2F - - IITERAL POOLS IN CONTROL SECTIONS

Literals, collected into pools by the assembler, are assembled as
part of the executable control section to which the pools belong. If a ITORG instruction is specified at the end of each control section, the literals specified for that section will te assembled into the pool starting at the LTORG instruction. If no LTORG instruction is specified, a literal pool containing all the literals used in the entire source module is assembled at the end of the first control section. This literal pool appears in the listings after the 3 END instruction.

NOTE: If any control section is divided into segments, a LTORG instruction should be specified at the end of each segment to create a separate literal pool for that segment. (For a complete discussion of the literal pool see H1B.)

| Type Code Assigned for External Symbol Dictionary | Unnamed Control Sections in separate Source Modules | Notes |
| :---: | :---: | :---: |
| PC <br> PC | CSECT <br> END | Unnecessary unless dictated by specific programming purpose |
| PC <br> PC sign | BALR 12,0 USING $*, 12$ $\cdot$ END <br> fies "private | Inadvertent and inadvisable initiation of first control section: instead, precede with a named START instruction |



The assembler keeps a record of each control section and prints the following information about it in an External Symbol Lictionary.

1. Its symbolic name, if cne is specified
2. Its type code
3. Its individual identification
4. Its starting address.

The figure to the right lists:

1. The assembler instructions that define control sections and dummy control sections or identify entry and external symbols,
2. The type code that the assembler assigns to the control sections or dummy control sections and to the entry and external symbols.

NOTE: The total number of entries identifying separate control sections, dummy control sections, entry symbols, and external symbols in the external symbol dictionary must not exceed 399. External symbols identified in a Q-type address constant and specified as the
(3) name entry of a DSECT instruction are counted twice in determining this total.

DOS The maximum number of external symbol dictionary entries (control sections, dumy control sections, and external symbols) allowed is 511. The maximum allowable number of symbols identified by the ENTRY instruction is 200.

| Name Entry | Instruction | Type code entered into external symbol dictionary |
| :---: | :---: | :---: |
| optional | START <br> CSECT <br> START <br> CSECT <br> Any instruction that initiates the unnamed control section | $\begin{array}{c\|c} S D & \begin{array}{l} \text { if name } \\ \text { entry is } \end{array} \\ \text { SD } \\ \text { present } \end{array}$ <br> PC |
| optional <br> Dos <br> blank <br> optional <br> DOS <br> mandatory | COM <br> DSECT | CM <br> none |
| OS only mandatory mandatory | $\left\|\begin{array}{ll} \text { DXD } & \\ & 3 \\ \text { (external } & \text { DSECT) } \end{array}\right\|$ | XD <br> XD |
|  | $2\left\{\begin{array}{l}\text { ENTRY } \\ \text { EXTRN } \\ \text { DC(V-type ad- } \\ \text { dress constant) } \\ \text { WXTRN }\end{array}\right.$ | LD <br> ER <br> ER <br> wx |

## E3 - Defining a Control Section

You must use the instructions described kelow to indicate to the assembler:

- Where a control section begins and
- Which type of contrcl section is being defined.


## E3A -- THE START INSTRUCTION

## Purpose

The START instruction can be used only to initiate the first or only executable control section of a source mcdule. You should use the START instruction for this purpose, because it allows you:

1. To determine exactly where the first control section is to begin; you thereby avoid the accidental initiation of the first control section by some other instruction.
2. To give a symbolic name to the first control section, which can then be distinguished from the other contrcl sections listed in the external symbol dictionary.
3. To specify the initial setting of the location counter for the first or only control section.

## Specifications

The START instruction must be the first instruction of the first executable control section of a source module. It must not be preceded by any instruction that affects the location counter and thereby causes the first control section to be initiated.

The format of the START instruction statement is given in the figure to the right.

|  |  |  |
| :--- | :--- | :--- |
| Name | Operation | Operand |
| Any Symbol <br> or blank | START | A self-defining <br> term, or blank |



The symbol in the name field, if specified, identifies the first control section. It must be used in the name field of any CSECT
(2)
instruction that indicates the continuation of the first control section. This symbcl represents the address of the first kyte of the control section and has a length attribute value of 1.

The assembler uses the value of the self-defining term in the operand field, if specified, to set the location counter to an initial value for the source module. All control sections are aligned on a doubleword boundary. Therefcre, if the value specified in the operand is not divisible by eight, the assembler sets the initial value of the
 location counter to the next higher doubleword boundary. If the operand
5 entry is omitted, the assembler sets the initial value to 0 .


The source statements that follow the START instruction are assembled into the first control section. If a CSECT instruction indicates the continuation of the first control section, the source statements that
(2)
follow this CSECT instruction are also assembled into the first control section.

Any instruction that defines a new or continued control section marks
3 the end of the preceding control section or portion of a control section. The END instruction marks the end of the control section in effect.

## E3B -- THE CSECT INSTRUCTION

## Furpose

The CSECT instruction allows you to initiate an executable control section or indicate the continuation of an executable control section.

## Specifications

The CSECT instruction can be used anywhere in a source module after any source macro definitions that are specified. If it is used to initiate the first executable control section, it must not be preceded ky any instruction that affects the location counter and thereby causes the first control section to be initiated.

The format of the CSECT instruction statement is shown in the figure to the right.

| Name | Operation | Operand |
| :--- | :--- | :--- |
| Any Symbol <br> or blank | CSECT | Not required |

The symbol in the name field, if specified, identifies the control section. If several CSECT instructions within a source module have the same symbol in the name
1 field, the first occurrence initiates the control section and the rest
(2) indicate the continuation of the control section. If the first control section is initiated by a START instruction, the symbol in the name field must be used to indicate any continuation of the first control section.

NOTE: A CSECT instruction with a blank name field either initiates or indicates the continuation of the unnamed control section (see E2E) -

The symbol in the name field represents the address of the first byte of the control section and has a length attribute value of 1.

The beginning of a control section is aligned on a doubleword boundary. However, the continuation of a

control section begins at the next available location in that control section.

The source statements that follow a CSECT instruction that either initiates or indicates the continuation of a control section are assembled into the object code of the control section identified by that CSECT instruction.

NOTES:

1. The end of a control section or portion of a control section is marked by:
a. Any instruction that defines a new or continued control section or
b. The END instruction.

DOS 2. The location counter is reset to zero each time the DOS/VS assembler encounters a CSECT instruction. (The figure on the right illustrates location counter settings when using the OS/VS assembler.)

## Purpose

You can use the ESECT instruction
to initiate a dummy control section
or to indicate its continuation.

A dummy control section is a reference control section that allows you to describe the layout of data in a storage area without actually reserving any virtual storage.

## How to Use a Lummy Control Section

The figure to the right illustrates a dummy control secticn.

A dummy control section (dummy section) allows you to write a
1 sequence of assembler language statements to describe the layout of
2 unformatted data located elsewhere in your program. The assembler produces no object code for statements in a dummy control section and it reserves no storage for the dummy section. Rather, the dummy section provides a symbolic format that is empty of data. However, the assembler assigns location values to the symbols you define in a dummy section, relative to the beginning of that dummy section.

Therefore, to use a dummy section you must:
4. Reserve a storage area for the unformatted data

- Ensure that this data is loaded into the area at execution time
- Ensure that the locations of the symbols in the dummy section actually correspond to the locations of the data being described

5. Establish the addressability of the dummy section in combination with the storage area (see F1A).

6 You can then refer to the unformatted data symbolically by using the symbols defined in the durmy section.

## Specifications

The DSECT instruction identifies
the beginning or continuation of a dummy control section (dummy section) . One or more dumm sections can be defined in a source module.

The DSECT instruction can be used anywhere in a source module after the ICTL instruction, or after any source macro definitions that may be specified.

The format of the ISECT instruction statement is given in the figure to the right.

The symbol in the name field, if specified, identifies the dummy section. If several LSECT instructions within a source module have the same symbol in the name
field, the first occurrence initiates the dummy section and the rest
(2) indicate the continuation of the dummy section.

NOTE: A DSECT instruction with a blank name field either initiates or indicates the continuation of the unnamed dummy section.

The symbol in the name field represents the first location in the dummy section and has a length attribute value of 1.

The location counter for a dummy
section is always set to an initial value of 0 . However, the continuation of a dummy section kegins at the next available location in that dummy section.

(1)

The source statements that follow a DSECT instruction belong to the dummy section identified by that DSECT instruction.

## NOTES:

1. The assembler language statements that appear in a dummy control section are not assembled into object code.
2. When establishing the

2 addressability of a dummy section,
2 the symbol in the name field of
(3)
the ESECT instruction or any symbol defined in the dummy section can be specified in a USING instruction.
3. A symbol defined in a dummy section can be specified in an address constant only if the symbol is paired with another symbol from the same durmy section, and if the symbols have the opposite sign.


## Purpose

You can use the COM instruction to initiate a common control section or to indicate its continuation. A common control section is a reference control section that allows you to reserve a storage area that can be used by two or more source modules.

How to Use a Common Control Section

The figure to the right illustrates a common control section.

A common control section (common section) allows you to describe a common storage area in one or more source modules.

When the separately assembled object modules are linked as cne program, the required storage space is reserved for the common control section. Thus, two or more modules share the common area.

Only the storage area is provided; the assembler does not assemble the source statements that make up a common control section into object code. You must provide the data for the common area at execution time.

3 The assembler assigns locations to the symbcls you define in a common section relative to the beginning of that common section.
4 This allows you to refer symbolically to the data that will be loaded at execution time. Note that you
5 must establish the addressability of a common control section in every source module in which it is specified (see $F 1 A$ ). If you code identical common sections in two or more source modules, you can communicate data symbolically between these modules through this common


NOTE: You can alsc code a common control section in a source module written in the FORTRAN language. This allows you to communicate ketween assembler language modules and FCRTRAN modules.

## Specifications

The COM instruction identifies the beginning or continuation of a common control section (common section) .

One or more commen sections can be defined in a source module.

Only one common section can be defined.

The COM instruction can be used anywhere in a source module after the ICTL instruction, or after any source macro definitions that may ke specified.

The format of the COM instruction statement is given in the figure to the right.

os The symbol in the name field, if only specified, identifies the common control section. If several COM instructions within a source module have the same symbol in the name field, the first occurrence initiates the common section and the rest indicate the continuation of the common section.

NOTE: A COM instruction with a blank name field either initiates or indicates the continuation of the unnamed common section.

The symbol in the name field represents the address of the first byte in the common section and has a length attribute value of 1.

The location counter for a common section is always set to an initial value of 0. However, the continuation of a common section begins at the next availakle location in that common secticn.

If a common section with the same name (or unnamed) is specified in two or more source modules, the amount of storage reserved for this common section is equal to that required by the longest common section specified.


The source statements that follow a COM instruction belong to the common section identified by that COM instruction.

NOTES:

1. The assembler language statements that appear in a common control section are not assembled into object code.
2. When establishing the addressability of a common section, the symbol in the name field of the COM instruction or any symbol defined in the common section can be specified in a USING instruction.

DOS Because the name entry of the com instruction must be blank, a symbol defined in the common section must be used as the base address in a USING instruction.

## Source Module

| FIRST | START | 0 |
| :---: | :---: | :---: |
|  | 2 |  |
|  | USING COMMON, 11 |  | or USING | A, 11 |
| :--- |

END

## E4 - External Dummy Sections

os
only

## Purpose

An external dummy section is a reference control section that allows you to describe storage areas for one or rore source modules, to be used as:

1. Work areas for each source module or
2. Communication areas between two or more source modules.


#### Abstract

When the assembled object modules are linked and loaded, you can dynamically allocate the storage required for all your external dummy sections at one time from one source module (for example, by using the GETMAIN macro instruction). This is not only convenient but you save space and prevent fragmentation of virtual storage.

To generate and use external dummy sections, you need to specify a combination of the following:


1. The DXD or DSECT instruction
2. The Q-type address constant
3. The CXD instruction.

Generating an External Lummy Section

An external dummy section is generated when you specify a EXL 2. instruction or a DSECT instruction in combination with a $Q$-type address constant that contains the name of the DSECT instruction.

4 You use the $Q$-type address constant 5 to reserve storage for the offset to the external dummy section whose name is specified in the operand. This offset is the distance in bytes from the beginning of the area allocated for all the external dummy sections to the beginning of the external durmy section specified. You can use this offset value to address the external dummy section. The g-type address constant is described in G3M.

How to Use External Eumry Sections

To use an external dummy section, you must dc the fcllcwirg (as illustrated in the figure kelow):
(1) Identify and define the external dummy section. The assembler will compute the length and alignment required.

2 Frovide a Q-type constant for each external dummy section defined.

Use the CXD instructicn tc reserve a fullword area intc which the linkage editor or loader will insert the tctal
(3) length of all the external dummy secticns that are specified in the scurce modules of your program. The linkage editcr computes this length frcm the lengths of the individual external dummy sections supplied ky the assembler.

Allocate a stcrage area using the computed tctal length.
5 Load the address of the allocated area into a register (for this example, register 11). Ncte that register 11 must contain this address throughout the whole prcgrar.

Add, to the address in register 11, the cffset into the allocated area of the desired external durny secticn. The linkage editcr inserts this cffset intc the fullword area reserved by the acpropriate $\varepsilon$-type address constar.t.
(1) Estaklish the addressakility of the external dumy section in combination with the portion of the allocated area reserved for the external dumy section.

8 You can now refer symbolically to the locations in the external dummy secticn.

Note that the source statements in an external dummy secticn are not assembled intc cbject code. Thus, at execution time you must insert the data descriked into the area reserved for the external dummy sections.


E5-Defining an External Dummy Section<br>os<br>only<br>E5A - THE DXD INSTRUCTION

## Purpose

The DXD instruction allows you to identify and define an external dummy section.

## Specifications

The DXD instruction defines an external dummy section. The DXD instruction can be used anywhere in a source module, after the ICTL instruction or after any source macro definitions that may be specified.

NOTE: The DSECT instruction also defines an external dummy section, but only if the symbol in the name field appears in a Q-type address constant in the same source module. Otherwise, a DSECT instruction defines a dummy section.

The format of the EXL instruction is given in the figure to the right.
(1) The symbol in the name field must appear in the operand of a Q-type address constant. This symbol represents the address of the first byte of the external dummy section defined and has a length attribute value of 1.

2 The subfields in the operand field are specified in the same way as in the DS instruction. The assembler computes the amount of storage and the alignment required for an external dummy section from the area specified in the operand field.

The linkage editor or loader uses the information provided by the assembler to compute the total length of storage required for all external dummy sections specified in a program.

NOTE: If two or more external dummy sections for different source modules have the same name, the linkage editor uses the most restrictive alignment and the largest section to compute the total length.

## Purpose

The CXD instruction allows you to reserve a fullword area in storage. The linkage editor or loader will insert into this area the total length of all external dummy sections specified in the source modules that are assembled and linked together into one prcgram.

## Specifications

The CXD instruction reserves a fullword area in storage, and it can appear in one or more of the source modules assembled and combined by the linkage editor into one program.

The format of the CXI instruction statement is given in the figure to the right.


The symbol in the name field, if specified, represents the address 1 of a fullword area aligned on a fullword boundary. This symbol has a length attribute value of 4. The linkage editor or loader

2 inserts into this area the total length of storage required for all the external dummy sections specified in a program.


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## Section F: Addressing

This section describes the techniques and instructions that allow you to use symbolic addresses when referring to data. You can address data that is defined within the same source module or data that is defined in another source module. Symbolic addresses are more meaningful and easier to use than the corresponding object code addresses required for machine instructions. Also, the assembler can convert the symbolic addresses you specify into their object code form.

Fl - Addressing Within Source Modules: Establishing Addressability

By establishing the addressability of a control section, you can refer to the symbolic addresses defined in it in the operands cf machine instructions. This is much easier than coding the addresses in the base-displacement form required by the System/370. The symbolic addresses you code in the instruction
 operands are called implicit addresses, and the addresses in the base-displacement form are
called explicit addresses, both of which are fully described in L5B.

The assembler will convert these implicit addresses for you into the explicit addresses required for the assembled object code of the machine instruction. However, you must supply the assembler with:

3 1. A base address from which it can compute displacements to the addresses within a control section and 2. A base register to hold this kase address.


How to Estaklish Addressability

To establish the addressability of a control section, you must, at coding time:

1. Specify a base address from which the assembler can compute displacements
(2) Assign a base register to contain this base address
(3) - Write the instruction that loads the base register with the base address.

At assembly time, the implicit addresses you code are converted
4 into their explicit tase-displacement form: then, they are assembled intc the crject code of the machine instructicns in which they have been coded.

At execution tire, the kase address
(5) is lcaded into the kase register and should remain there thrcughcut the execution of your program.


FIA - THE USING INSTRUCTION

## Purpose

The USING instruction allows you to specify a base address and assign one or more base registers. If you also load the base register with the base address, you have established addressability in a control section.

To use the USING instruction correctly you should:

1. Know which locations in a control section are made addressable by the USING instruction
2. Know where in a source module you can use these estatlished addresses as implicit addresses in instruction operands.

The range of a USING instruction (called the USING range) is the

## (1)

 4,096 bytes beginning at the base address specified in the USING instruction. Addresses that lie within the USING range can ke(2.) converted from their implicit to their explicit form; those outside
3. the USING range cannct be converted.

The USING range does not depend upon the position of the USING instruction in the source module; rather, it depends upon the location of the base address specified in the USING instruction.

NOTE: The USING range is the range of addresses in a control section that is associated with the base register specified in the USING instruction. If the. USING instruction assigns more than one base register, the composite USING range is the sum of the USING ranges that would apply if the base registers were specified in separate USING instructions. If register 0 is specified as the base register, the USIING range will be location 0 through 4095 regardless of the base address specified.

The Eomain of a USING Instruction

The domain of a USING instruction (called the USING domain) begins where the USING instruction appears in a source module and continues to the end of the source module. (Exceptions are discussed later in this subsection, under NOTES ABOUT THE USING DOMAIN.) The 5 assembler converts implicit address references into their explicit form:

1. If the address reference appears in the domain of a USING instruction and

2. If the addresses referred to lie within the range of the same USING instruction.
3. The assembler does not convert address references that are cutside the USING domain. The USING domain depends on the position of the USING instruction in the source module after conditional assembly, if any, has been performed.

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## How to Use the USING Instruction

You should specify your USING instructions so that:

1. All the addresses in each control section lie within a USING range and
2. All the references for these addresses lie within the corresponding USING domain.

You should therefore place all USING instructions at the beginning of the source module and specify a base address in each USING instruction that lies at the beginning of each control section.

FOR EXECUTABLE CONTROL SECTIONS: The figure to the right illustrates a way of estaklishing the addressability of an executable control section (defined by a START or CSECT instruction) . You specify a base address and assign a base
2 At execution time the base register is loaded with the correct base address.

Note that for this particular combination of the BALR and USING instructions, you should code them exactly as shown in the figure to the right.

If a control section is longer than 4096 bytes, you must assign more than one base register. This allows you tc establish the addressability of the entire control section with one USING instruction as shown in the figure on the opposite page.

The assembler assumes that the base registers that you assign contain the correct base addresses. The address of HERE is loaded into the first base register. The addresses HERE+4096 and HERE+8192 are loaded into the second and third base registers respectively.

2 Note that you must define the address, EASES, within the first part of the total USING range, that is, the addresses covered by base register 9. This is because the explicit
3 address converted from the implicit address reference, is assembled into the LM instruction. At execution time, the assembled address must have a kase register which already contains a base address at this point; the only base register loaded with its base address is register 9.

The addressability of addresses in the USING range covered by the second and third base registers is not completely established until after the LM instruction.

NOTE: Addresses specified in address constants (except the S-type) are not converted to their base-displacement form.


FOR REFERENCE CONTROL SECTIONS:
The figure to the right illustrates how to establish the addressability of a dummy section. A dummy section is a reference control section defined by the DSECT instructions. Examples of establishing addressability for the other reference control sections are given in $E 3 \mathrm{E}$ and E 4 .

1 As the kase address, ycu shculd specify the address of the first kyte of the dummy secticn, sc that all its addresses lie within the Fertinent USING range.


The address you load into the base register must be the address of the storage area being formatted by the dummy section.

Note that the assembler assumes that you are referring to the symbolic addresses of the dummy section, and it computes
displacements accordingly. However, at execution time, the assembled addresses refer to the location of real data in the storage area.


The USING instruction must ke coded as shown in the figure to the right.

The operand, EASE, specifies a base address, which can be a relocatable or absolute expression. The value of the expression must lie between -224 and 224-1.

The remaining operands specify from 1 to 16 base registers. The operands must ke atsclute expressions whose values lie in the range 0 through 15.

The assembler assumes that the first base register (EASREG1) contains the base address BASE at execution time. If present, the subsequent operands, BASREG2, BASREG3,.... represent registers that the assembler assumes will contain the address values, BASE+4096. BASE+8192,..., respectively.

NOTES ABOUT THE USING DOMAIN: The domain of a USING instruction continues until the end of a scurce module except when:
(1) A subsequent LROP instruction specifies the same base register or registers assigned by the preceding USING instruction.
2. A subsequent USING instruction specifies the same register or registers assigned by the preceding USING instruction.



NOTES AROUT THE USING RANGE: TwC USING ranges coincide when the same tase address is specified in twic different USING instructions, even though the lase registers used are different. When two USING ranges coincide, the assembler uses the
(1) higher numbered register for assembling the addresses within
2 the common USING range. In the example, this apflies cnly to the
(3) implicit addresses that arpear after the second USING instruction. In effect, the first USING domain is terminated after the seccnd USING instruction.


Two USING ranges overlap when the base address of cne USING instruction lies within the range of another
 USING instructicn. When twc ranges
) overlap, the assembler computes
2 displacements from the kase address that gives the smallest displacement; it uses the corresponding base
register when it assembles the addresses within the range cverlap. This applies only to implicit
(4) addresses that apcear after the second USING instruction.


EASE REGISTERS FOR AESOIUTE
ADLRESSES: Absclute addresses used in a source mcdule must alsc be made addressable. Absolute addresses require a kase register cther than the base register assigned to relocatakle addresses (as descriked above).

However, the assembler does nct need a USING instruction to convert absolute implicit addresses in the range 0 thrcugh 4,095 to their explicit form. The assembler uses register 0 as a base register. Cisplacements are computed from the base address 0 , because the assembler assumes that a kase cr index of 0 implies that a zero quantity is tc ke used in forming the address, regardless of the contents of register 0. The USING domain for this automatic base register assignment is the whcle of a source mcdule.


For absolute implicit addresses greater than 4095, a USING instruction must be specified according to the fcllcwing:
(1) With a base address representing an atsolute expressicn, and
(2) With a base register that has not been assigned by a USING instruction in which a relocatable base address is specified.

3 This kase register must be lcaded with the base address specified.


## Purpose

You can use the $\operatorname{DROP}$ instruction to indicate to the assembler that one or more registers are no longer available as base registers. This allows you:

1. To free kase registers for other programming purposes
2. To ensure that the assembler uses the base register you wish in a particular coding situation, for example, when two USING ranges overlap or coincide (as described above in F1A, Notes about the USING range).

The LROP instruction must be coded as shown in the figure to the right.

Up to 16 operands can be specified. They must be absolute expressions whose values represent the general registers 0 through 15. A LROP instruction with a blank cperand field causes all currently active base registers assigned by USING instructions to be dropped.
(1)

After a DROF instruction, the assembler will not use the registers specified in a DROP instruction as base registers. A register made unavailable as a base register by a DROF instruction can be reassigned as a kase register by a subsequent USING instruction.


A LROF instruction is nct needed:

- If the base address is being
changed ky a new USING instructicn, and the same base register is assigned. However, the new kase address must be loaded into the kase register. Ncte that the
(3) implicit address "E" lies within the first USING dcrain, and that
(4) the base address to which it refers lies within the first USING range.
- At the end of a source module.



## F2 - Addressing Between Source Modules: Symbolic Linkage

This section describes symbclic linkage, that is, using symbols to communicate between different source modules that are separately assembled and then linked together by the linkage editor.

How to Establish Symbolic Linkage

You must establish symbolic linkage between source modules so that you can refer or kranch to symbolic locations defined in the control sections of external source modules. To establish symbolic linkage with an external source module you must do the following:

1. In the current source module, you must identify the symbols that are not defined in that scurce module, if you wish to use them in instruction operands. These symbols are called external symbols. because they are defined in another (external) source module. You identify external symbols in the EXTRN or WXTRN instruction or the V-type address constant.
2. In the external source modules, you must identify the symbols that are defined in those source modules and to which you refer from the current source module. These symbols are called entry symbols kecause they provide points of entry to a control section in a source module. You identify entry symbols with the ENTRY instruction.
3. You must provide the A-type or Y-type address constants needed ky the assembler tc reserve storage for the addresses represented by the external symbcls.

The assembler places information about entry and external symbols in the External Symbol Lictionary. The linkage editor uses this

information to resolve the linkage addresses identified by the entry and external symbols.

TO REFER TC EXTERNAL LATA: You should use the EXTRN instruction 1 to identify the external symbol that represents data in an external source module, if you wish to refer to this data symbclically.

For example, you can identify the address of a data area as an external symbol and load the address constant specifying this symbol into a kase register. Then, ycu use this kase register when establishing the
(3) addressability cf a duminy section that formats this external data. You can now refer symbclically tc the data that the external area contains.

You must also identify, in the source module that contains the data area, the address cf the data as an entry symbol.


TO BRANCH TC AN EXTERNAL ALERESS: You should use the $V$-type address constant to identify the external symbol that represents the address in an external source module to which you wish tc kranch. Fcr the specifications of the $v$-type address constant, see G3L.

For example, you can lcad intc a register the $V$-type address constant that identifies the external symbcl. Using this register, you can then
3 kranch to the external address represented by the symbol.

If the symbcl is the name entry cf a START or CSECT instruction in the other source rodule, and thus names an executakle control section, it is autcratically identified as an entry symbol. If the symbcl represents an address in the middle of a control section, you must, however, identify it as
(5) an entry symbol for the external source module.

You can alsc use a combination of an EXIRN instructicn tc identify and an A-type address constant to contain the external kranch address. However, the $V$-type address constant is more convenient kecause:

1. Ycu do nct have to use an EXTRN instruction.
2. The symbcl identified is not considered as defined in the source module and can ke used as the name entry for any other statement in the same source ncdule.

## Furpose

The entry instruction allows you to identify symkcls defined in a source module so that they can $\mathrm{k} \in$ referred to in ancther scurce module. These symbols are entry symkols.

## Specifications

The format of the ENTRY instructicn is shown in the figure to the right.

ENTRY SYNECIS: The fcllcwing applies to the entry symbcls identified in the ccerand field:

- They must ke valid symkols.
- They must be defined in an
(1) exccutable control section.
- They must not ke defined in a

2 dummy contrcl section, a common contrcl section, or an external contrcl section.

- The length attribute value cf entry symbols is the same as the length attribute value of the symbol at its point cf definition.

A symbol used as the name entry of a STARI cr CSECI instructicn is also autcmatically considered
(3) an entry symbcl and dces not have to be identified by an ENTRY instruction.

The assembler lists each entry symbol of a scurce module in an
(4) External Symbol Lictionary along with entries for external symbcls, commen control sections, and external contral sections. The maximum number of External Symbel Dictionary entries for each source module is 399.

DOS The maximum nuxber of external symbol dictionary entries (contrcl sections and external symbols) allowed is 511. The maximum allowable number of entry symbols identified ky the ENTRY instruction is 200.
(5) NOTE: A symkcl identified in an ENTRY instruction counts towards this maximum, even thcugh it may not be used in the name field of a statement in the scurce module nor constitute a valid entry point.

| Source Module |  |  | Entry in External 4 Symbol Dictionary |  |
| :---: | :---: | :---: | :---: | :---: |
|  |  |  | Symbol | Type Code |
| FIRST | StART |  | FIRST | SD |
|  | ENTRY | SUBRTN,INVALID | SUBRTN | LD |
|  | - |  | INVALID | LD |
| DUMMY <br> 2 <br> INVALID | DSECT |  | DUMMY | none |
|  | DS | F | INVALID | - |
|  | END |  |  |  |

## Purpose

The EXTRN instruction allows you to identify symbcls referred tc in a source module but defined in another scurce mcdule. These symbcls are external symbols.

## Specifications

The format of the EXIRN instructicn staterent is shown in the figure to the right.

EXTERNAL SYMEOLS: The follcwing applies to the external symkols identified in the cFerand field:
(1) They must be valid symkols.

- They must not be used as the name entry cf a scurce statement in the source module in which they are identified.
- They have a length attrikute value of 1.

2

- They must ke used alcne and cannct be paired when used in an expression (fcr pairing of terms see C6).

The assembler lists each external symbcl identified in a source module in the External Symbcl Eicticnary along with entries for entry symkols, common control secticns, and external contrcl sections. The maximum number of External Symbcl Dicticnary entries for each source module is 399.

The maximum number of external symbel dictionary entries (contrcl sections and external symbols) allowed is 511. The naximum allowable number of entry symbols identified ly the ENTRY instruction is 200 .

NOTE: The symbcl sfecified in a v-type address constant is implicitly identified as an external symbcl and ccunts towards this maximum.

Purpose
The WXTRN instruction allows you to identify symbcls referred tc in a source module kut defined in another scurce module.
The KXIRN instruction differs from the EXTRN instruction as follcws:
The EXTRN instruction causes the linkage editcr to make an automatic search of libraries to find the module that contains the external symbcls that you identify in its operand field. If the module is found, linkage addresses are resolved; then the module is linked to your module, which contains the EXTRA instruction.
The KXTRN instruction supfresses this automatic search cf likraries. The linkage editor will only resolve the linkage addresses if the external symbcls that you identify in the WXTRN operand field are defined:

1. In a module that is linked and loaded along with the ckject ricdule assembled from your source module or
2. In a module brought in from a likrary due to the fresence of an EXIRN instruction in another module linked and lcaded with ycurs.

## Specifications

The format of the WXIRN instruction statement is shown in the figure to the right.
EXTERNAL SYMEOLS: The external symbcls identified ky a WXTRN instruction have the same prcperties as the external symbcls identified
(2) ty the EXTRN instructicn. However, the type code assigned to these external symbcls differs.
NOTE: If a symbol, specified in
(3) a v-type address constant, is also identified by a wXTRN instruction in the same scurce module, it is assigned the same type code as the symkol in the WXTRN instructicn.
If an external symbol is identified ky koth an EXIRN and WXIRN instruction in the same source
module, the first declaration takes frecedence, and suksequent declarations are flagged with warning

| Source Module | Entry in External Symbol Dictionary |  |
| :---: | :---: | :---: |
|  | Symbol | Type Code |
| FIRST START 0 <br> 2 EXTRN OUT,A | FIRST | SD |
|  | OUT | ER |
|  |  | ER |
| 1 WXTRN WOUT, A | WOUT | WX |
| VCON DC V (WOUT) | WOUT | WX |
| END |  |  | messages.

## Section G: Symbol and Data Definition

This section describes the assembly time facilities which you can use to:

1. Assign values to symbols
2. Define constants and storage areas
3. Define channel command words.

By assigning an absolute value to a symbol and then using that symbol to represent, for example, a register or a length, you can code machine instructions entirely in symbolic form.

## G1 - Establishing Symbolic Representation

You define symbols tc be used as elements in your programs. This
(1) symkolic representation is superior
2 to numeric representation because:

- You can give meaningful names to the elements;
- You can debug a program more easily, because the symbols are cross-referenced to where they are defined and used in your program. The cross-referenced statement numbers containing the symbols are printed in your assembly listing.
- You can maintain a program more easily, because you can change a symbolic value in one place and its value will be changed throughout a program.
(3) Some symbols represent absolute values, while others represent (4) relocatable address values. The relocatable addresses are of:
(5) instructions
(6) constants

7 storage areas.
You can use these defined symbols in the operand fields of instruction statements to refer to the
8 instructions, constants, or areas represented by the symbol.


## Assigning Values

You can create symbcls and assign them absolute or relocatable values anywhere in a source module with an EQU instruction (see G2A). You can use these symbols instead of the numeric value they represent in the operand of an instruction.

## Defining and Naming Data

DATA CONSTANTS: You can define a data constant at assembly time that will be used by the machine instructions in their operations at execution time. The three steps for creating a data constant and introducing it into your program in symbolic form are:

## (1) define the data

2) provide a label for the data
(3) refer to the data by its label.

The symbol used as a label represents 4 the address of the constant; it is not to be confused with the
5 assembled object code of the actual constant.

Defining data constants is discussed in G3.

LITERALS: You can also define data at its point of reference in the operand of a machine instruction
6 by specifying a literal.
Literal constants are discussed in G3C.

STORAGE AREAS: You must usually reserve space in virtual storage at assembly time for insertion and manipulation of data at execution time. The three steps for reserving virtual storage and using it in your program are:
(1) define the space
(2) provide a label for the space
(3) refer to the space by its label.

Defining storage areas is discussed in G3N.


CHANNEL COMNAND WORDS: When you define a channel command word at assembly time you create a command for an input or output operation to be performed at execution time. You should:

- define the channel command word
- provide a label fcr the word.

Channel cormand words are discussed in subsection G3C.

## G2 -- Defining Symbols

G2A -- THE ECQU INSTRUCTION

## Purpose

The EÇU instruction allows you to assign absolute or relccatable values to symbols. You can use it for the fcllowing purfcses:
(1) 1. To assign single absolute values to symbols
(2)
2. To assign the values of previously defined symbols or expressicns to new symbols, thus allowing you to use different mnemonics for different purpcses.
3. 3. To compute expressicns whose values are unknown at coding time or difficult to calculate. The value of the expression is then


The EQU instruction can be used anywhere in a source module after the ICTL instruction, or after any source macro definitions that may be specified. Note, however, that the FCU instruction can initiate an unnamed control section (private code) if it is specified before the first control section (initiated by a START or CSECT instruction).

The format of the FQU instruction statement is given in the figure to the right.

Dos Only one operand (expression 1) is allowed.

| Name | Operation | EQU |
| :--- | :--- | :--- |
| An ordinary <br> symbol or <br> a variable <br> symbol | EQU | Operand |

Expression 1 represents a value. It must always be specified and can have a relocatable or absolute value. The assembler carries this value as a signed four-byte (32-bit) number; all four bytes are printed in the program listings opposite the symbol.

OS Expression 2 represents a length only attribute. It is optional, but, if specified, it must have an absolute value in the range of 0 through 65,535. Expression 3 represents a type attribute. It is optional, but, if specified, must be a self-defining term with a value in the range of 0 through 255.

Any symbols appearing in these three expressions must have been previously defined.

EXPRESSION 1 (VALUE): The assembler assigns the relocatable or absolute value of expression 1 to the symbol in the name field at assembly time.

If expression 2 is omitted, the assembler also assigns a length attribute value to the symbol in the name field according to the length attribute value of the leftmost (or only) term of expression 1. The length attribute value (described in C4C) thus assigned is as follows (see figure on following page) :

1. If the leftmost term is a lccation counter reference (*) , a self-defining term or a symkol length attribut (1) value reference, the length attribute value is 1. Note that this also applies if the leftmost term is a symbcl that is equated tc any cf these values.
2. If the leftmost term is a symbol that is used in the name field of a DC or DS instruction, the length attribute value is equal to the implicit or explicit length of the first (or only) constant specified in the DC or DS operand field.
3. If the leftmost term is a symbol that is used in the name field of a machine instruction, the length attribute value is equal to the length of the assembled instruction.
4. Symbols that name assembler instructions, except the DC and DS instructions, have a length attribute value of one. However, the name of a CCW instruction has a length attribute value of eight.

NOTE: The length attribute value assigned in cases 2-4 only applies to the assembly-time value of the attribute. Its value at pre-assembly time, during conditional assembly processing, is always 1.

Further, if expression 3 is omitted, the assembler assigns a type attribute value cf "U" to the symbcl in the name field.

| Value assigned to symbol is: | Source Module |  |  | Length Attribute Value assigned to symbol in name field: |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | SECTA | START | 0 | At Assembly Time | At Pre-assembly Time |
|  | SECTA | - |  | (5) | $6$ |
|  | RR | LR | 3,4 |  |  |
|  | RX | A | 3,FULL |  |  |
|  | SS | MVC | TO,FROM |  |  |
|  | FULL | DC | F'33' |  |  |
|  | AREA | DS | XL2000 |  |  |
|  | TO | DS | CL240 |  |  |
|  | FROM | DS | CL80 |  |  |
|  | ADCONS | DC | AL1 (A) , AL2 (B) , AL3 ( C ) |  |  |
|  | ADCCW | CCW | 2,READER, ${ }^{\prime}$ '48',80 |  |  |
| Absolute | A | EQU | $X^{\prime} F^{\prime}$ | 1 | 1 |
| Absolute | B | EQU | L'FROM | 1 | 1 |
| Relocatable | C | EQU | *+4 | 1 | 1 |
| Absolute | D | EQU | A $* 10$ | 1 | 1 |
| Relocatable | E | EQU | FULL |  |  |
| Relocatable | F | EQU | AREA +1000 | 2000 | 1 |
| Relocatable | G | EQU | TO 2 | 240 | 1 |
| Absolute | H | EQU | FROM-TO | 80 | 1 |
| Relocatable | I | EQU | ADCONS | 1 | 1 |
| Relocatable | J | EQU | RR 3 | 2 | 1 |
| Relocatable | K | EQU | RX 3 | 4 | 1 |
| Relocatable | L | EQU | SS | 6 | 1 |
| Relocatable | M | EQU | SECTA | 1 | 1 |
| Relocatable | N | EQU | ADCCW | 8 | 1 |

OS EXPRESSION 2 (LENGTH-ATTRIBUTE VALUE): If expression 2 is only specified, the assembler assigns its value as a length

1 attribute value to the symbol in the name field. This value overrides the normal length attribute value implicitly assigned from expression 1.

If expression 2 is a self-defining term, the assembler also assigns the length attribute value to the symbol at pre-assembly time (during conditional assembly processing).

OS EXPRESSION 3 (TYPE-ATTRIBUTE VALUE): If expression 3 is only specified, it must be a self-defining term. The assembler
 assigns its EBCDIC value as a type attribute value to the symbol in the name field. This value overrides the normal type attribute value implicitly assigned from expression 1. Note that the type attribute value is the EBCDIC character
(4) equivalent of the value of expression 3 .

| Value assigned | Source Module |  |  | Length Attribute Value assigned |  | Type Attribute <br> Value assigned |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | $\mathrm{F}^{\prime} \perp \mathrm{KS} \mathrm{S}^{\prime} \mathrm{C}$ <br> AREA <br> SDT <br> ASTERISK | $\begin{aligned} & \text { START } \\ & \text { DS } \\ & \cdot \\ & \text { EQU } \\ & \cdot \\ & \text { EQU } \end{aligned}$ | XL2000 <br> Implicit <br> $\begin{array}{ll}X^{\prime} F^{\prime} & \begin{array}{l}\text { Attribute } \\ \\ \\ \text { Values }\end{array}\end{array}$ | At <br> Assembly <br> Time <br> 2000 <br> 1 <br> 1 | At Preassembly Time <br> 2000 | $\begin{aligned} & \mathrm{X} \\ & \mathrm{U} \\ & \mathrm{U} \end{aligned}$ |
| $\left.\begin{array}{l}\text { Value of } \\ \text { AREA } \\ 255 \\ \text { Value of } \\ \text { Location } \\ \text { Counter }\end{array}\right\}$ | $\begin{aligned} & \mathrm{A} \\ & \mathrm{~B} \\ & \mathrm{C} \end{aligned}$ | EQU <br> EQU <br> EQU | $\begin{array}{lr} \text { AREA }, 1000 \\ \text { SDT }, 4 \\ \text { ASTERISK, } 4 \end{array} \quad 1$ | $\begin{array}{r} 1000 \\ 4 \\ 4 \end{array}$ | $\begin{array}{r} 1000 \\ 4 \\ 4 \end{array}$ | $\begin{aligned} & \mathrm{U} \\ & \mathrm{U} \\ & \mathrm{U} \end{aligned}$ |
|  | $\begin{aligned} & \hline \mathrm{D} \\ & \mathrm{E} \\ & \mathrm{~F} \end{aligned}$ | $\begin{aligned} & \text { EQU } \\ & \text { EQU } \\ & \text { EQU } \end{aligned}$ | $\begin{aligned} & \text { AREA, , C'F' } \\ & \text { SDT, C'N' } \\ & \text { ASTERISK, , }{ }^{\prime}{ }^{\prime}{ }^{\prime} \end{aligned}$ | $\begin{array}{r} 2000 \\ 1 \\ 1 \end{array}$ | 1 1 1 | $\left.\begin{array}{l}\text { F } \\ \mathrm{N} \\ \mathrm{A}\end{array}\right\} 3$ |
|  | $\begin{gathered} \mathrm{G} \\ \mathrm{H} \\ \mathrm{I} \\ \\ \mathrm{~J} \end{gathered}$ | $\begin{aligned} & \text { EQU } \\ & \text { EQU } \\ & \text { EQU } \\ & \\ & \text { EQU } \end{aligned}$ | $\begin{aligned} & \text { AREA, 1000,C'1' } \\ & \text { SDT,4,C'F' } \\ & \text { ASTERISK, } 4, C^{\prime} A^{\prime} \\ & \text { AREA, } 100,198 \end{aligned}$ | $\begin{array}{r} 1000 \\ 4 \\ 4 \\ \hline 100 \\ \hline \end{array}$ | 1000 <br> 4 <br> $\quad 4$ <br> $\quad 100$ | $\begin{aligned} & \hline \mathrm{l} \\ & \mathrm{~F} \\ & \mathrm{~A} \\ & \\ & \hline \mathrm{~F} \end{aligned}$ |

## Using Preassembly Values

You can use the preassembly values assigned by the assembler
in conditional assembly processing.
If only expression 1 is specified, the assembler assigns a
preassembly value of 1 to the length attribute and a
preassembly value of $U$ to the type attribute of the symbol.
These values can be used in conditional assembly (although
references to the length attribute of the symbol will be
flagged). The absolute or relocatable value of the symbol,
however, is not assigned until assembly, and thus may not be
used at preassembly.
If you include expressions 2 and 3 and wish to use the
explicit attribute values in preassembly processing, then
? The symbol in the name field must be an ordinary symbol
?
Q Expression 2 and expression 3 must be single

THE SYMECL IN THE NPME FIELL: The assembler assigns an absolute cr relccatakle value, a length attribute value, and a type attrikute value tc the symbcl in the name field.

The absolute or relocatable value of the symkol is assigned at assembly time, and is therefore not available for conditional assembly processing at pre-assembly tirre.

OS The type and length attribute values of the symbol are only available for conditional assembly processing under the following conditicns:

1. The symbcl in the name field must be an ordinary symbol.
2. Expression 2 and Expression 3 must ke single selfdefining terms.

## G3 - Defining Data

This section descrikes the IC, IS, and CCW instructicns; these instructions are used to define constants, reserve storage and specify the contents of channel command words respectively. Ycu can alsc provide a lakel for these instructions and then refer to the data symbclically ir the operands cf rachine and assembler instructions. this data is generated and storage is reserved at assembly time, and used by the machine instructions at execution tire.

## Purpose

You specify the DC instruction to define the data constants you need for program execution. The LC instruction causes the assembler
to generate the binary representation of the data constant you specify, into a particular location in the assembled source module; this is done at assembly time.

TYPES OF CONSTANTS: The LC instruction can generate the following types of constants:

1 Binary constants -- to define bit patterns

2 Character constants -- to define character strings or messages
(3) Hexadecimal constants -- to define large bit patterns
4. Fixed-Point constants -- for use by the fixed-point and other instructions of the standard set
(5) Decimal constants -- for use by the decimal instructions

6 Floating-Foint constants -- for use ky the floating-point instruction set

| (1) FLAG | DC | $B^{\prime} 0001000{ }^{\prime}$ |  |
| :---: | :---: | :---: | :---: |
| 2 CHAR | DC | C'STRING OF | CHARACTERS' |
| (3) PATTERN | DC | X'FFOOFF00' |  |
| (4) $\left\{_{\text {FCON }}\right.$ | $\mathrm{L}_{\mathrm{L}} \mathrm{C}$ | $\begin{aligned} & 3, F C O N \\ & \text { F'IO }^{\prime} \end{aligned}$ |  |
| (5) $\begin{aligned} & \text { PCON } \\ & \text { AREA }\end{aligned}$ | $\begin{aligned} & \text { AP } \\ & \text { DC } \\ & \text { DS } \end{aligned}$ | $\begin{aligned} & \text { AREA, PCON } \\ & P^{\prime} 100^{\prime} \\ & P^{\prime} \end{aligned}$ |  |
| (6) $\left\{_{\text {ECON }}\right.$ | $\begin{aligned} & \text { LE } \\ & \text { DC } \end{aligned}$ | $\begin{aligned} & 2, E C O N \\ & \text { E' }^{\prime} 100.50^{\prime} \end{aligned}$ |  |
| $1\left\{_{A D C O N}\right.$ | $\begin{aligned} & \mathrm{L} \\ & \mathrm{DC} \end{aligned}$ | $\begin{aligned} & 5, \text { ADCON } \\ & \text { A (SOMWHERE) } \end{aligned}$ |  |

7. Address constants -- to define addresses mainly for the use of the fixed-point and cther instructions in the standard instruction set.

The general format of the LC instructions staterents is shown in the figure to the right. represents the address cf the first byte of the assembled constant.
1 If several cperands are specified, the first constant defined is
(2) addressakle ky the symbcl in the name field. The other constants can $k \in$ reached ky relative addressing.

Each cperand in a IC instruction staterent consists cf fcur subfields. The fcrmat of a IC instruction operand is given in the figure tc the right.

The first three subfields describe the constant, and the fourth subfield specifies the nominal value of the constant to be generated.


Rules for the DC Operand
(1) 1. The type subfield and the nominal (2) value must always be specified.
2. The duplication factor and modifier subfields are optional.
(3) 3. When multiple operands are specified, they can be of different types.
(4) 4. When multiple nominal values are specified in the fourth subfield. they must be separated by commas and be of the same type.
5. 5. The descriptive subfields apply to all the nominal values.

NOTE: Separate constants are generated for each separate operand and nominal value specified.
6. No blanks are allowed:
a. Between subfields
b. Between multiple operands
c. Within any subfields -unless they occur as part of the nominal value cf a character constant or as part of a character self-defining term in a modifier expression or in the duplication factor subfield.


Information akout Constants

SYMBOLIC ADLRESSES OF CONSTANTS: Constants defined by the DC instruction are assembled into an object module at the location where the instruction is specified. However, the type of constant being defined will determine whether the constant is to be aligned on a particular storage boundary or not. (see kelow under Alignment of Constants). The value of the symbol that names the DC instruction is the address of the leftmost byte (after alignment) of the first or only constant.


| Type of constant | Implicit Length ${ }^{1}$ <br> 1 | Examples | Value of Length Attribute ${ }^{2}$ |
| :---: | :---: | :---: | :---: |
| B | as needed | DC B'10010000' | 1 |
| C | as needed | DC C'WOW' | 3 |
|  |  | DC CL8'WOW' | 8 |
| x | as needed | DC X'FFEEOO' | 3 |
|  |  | DC XL2'FFEE' | 2 |
| H | $2$ | DC H'32' | 2 |
| $F$ | 4 | DC FL3'32' | 3 |
| P | as needed | DC P'123' | 2 |
|  |  | DC PL4'l23' | 4 |
| Z | as needed | DC z'l23' | 3 |
|  |  | DC zLl0'123'」 | 10 |
| E | 4 |  |  |
| D | 8 |  |  |
| L | 16 |  |  |
| Y | 2 | DC Y (HERE) | 2 |
| A | 4 | DC ALI (THERE) | 1 |
| S | 2 |  |  |
| V | 4 |  |  |
| Q | 4 |  |  |
| ${ }^{1}$ Depends on type |  |  |  |
| ${ }^{2}$ Depends on whether or not an explicit length is specified in constant |  |  |  |

ALIGNNENT OF CONSTANIS: The assembler aligns constants on different boundaries according to the following:
(1) On boundaries implicit to the type of constant, when no length specification is supplied.

2 On byte koundaries when an explicit length specification is made.

Bytes that are skipped to align a constant at the proper boundary are not considered part of the constant. They are filled with zeros. Note that the automatic alignment of constants and areas does not occur if the NOALIGN assembler oftion has been specified in the job control language which invoked the assembler.

NOTE: Alignment can be forced to any boundary by a preceding DS (or DC) instruction with a zero duplication factor (see G3N). This occurs when either the ALIGN or NOALIGN option is set.

| Type of Constant | Implicit <br> Boundary <br> Alignment | Examples | Boundary Alignment |
| :---: | :---: | :---: | :---: |
| B | byte |  |  |
| C | byte |  |  |
| X | byte |  |  |
| H F | halfword <br> fullword | $\begin{aligned} & \text { DC H'25' } \\ & \text { DC HL'3'25' } \\ & \text { DC F'225' } \\ & \text { DC FL7'225' } \end{aligned}$ | halfword byte <br> fullword <br> byte |
| P Z | byte byte | $\left\lvert\, \begin{array}{ll} D C & P^{\prime} 2934^{\prime} \\ D C & Z^{\prime} 1235^{\prime} \\ D C & Z L 2^{\prime} 1235^{\prime} \end{array}\right.$ $2$ | byte <br> byte byte |
| E D L | fullword <br> doubleword <br> doubleword | DC E'1.25' DC EL5'l.25' DC 8D'95' DC 8DL7'95' DC L'2.57E65' | fullword byte doubleword byte doubleword |
| $\begin{aligned} & \mathrm{Y} \\ & \mathrm{~A} \\ & \mathrm{~S} \\ & \mathrm{~V} \\ & \mathrm{Q} \end{aligned}$ | halfword <br> fuliword <br> halfword <br> fullword <br> fullword | DC Y(HERE) <br> DC AL3(THERE) | halfword <br> byte |
| ${ }^{1}$ Depends on type 1 |  |  |  |

The nominal values sfecified for constants are assembled into storage. The amount of space availakle for the romiral value of a constant is determined:

1. Ey the explicit length specified in the second cperand sukfield, cr
2. If nc explicit length is specified, by the implicit length according to the type of constant defined (see Appendix VI).

PADEING: If more space is available than is needed to accommodate the binary representation of the nominal value, the extra space is padded:
(1) With binary zeros on the left for the kinary (B), hexadecimal (X), fixed-point ( $\mathrm{H}, \mathrm{F}$ ) , packed decimal (P), and all address (A,Y,S,V,Q) constants

With EBCDIC zeros on the left (X'FO') for the zoned decimal (Z) constants

3 With EBCDIC blanks on the right (X'40') for the character (C) constant

NOTE: Floating-point constants (E,D,I) are also padded on the right with zeros (see G3I).


TRUNCATION: If less space is available than is needed to accommodate the nominal value, the nominal value is truncated and part of the constant is lost. Truncation of the nominal value is:

On the left for the binary ( $B$ ), hexadecimal (X), decimal ( $P$ and 2) , and address ( $A$ and $Y$ ) constants.
(2) On the right for the character (C)
(3)

However, the fixed-pcint constants ( $H$ and $F$ ) will not be truncated, but flagged if significant bits would be lost through truncation.

NOTE: Floating-point constants (E, L,L) are not truncated; they are rounded (see G3I).

NOTE: The akove rules for padding and truncation also apply when the kit-length specification is used (see below under Subfield 3: Modifiers).

## Subfield 1: Duplication Factor

The duplication factor, if specified, causes the nominal value cr multiple nominal values specified in a constant to be generated the number of times indicated by the factor.
It is applied after the nominal
2 value or values are assembled into the constant.

The factor can be specified by an
(3) unsigned decimal self-defining term or by an absolute expression enclosed in parentheses.

The expression should have a positive value or be equal tc zero.

Any symbols used in the expression must ke previously defined.

Truncation


## NOTES:

1. The value of a location counter reference in a duplication factor is the value before any align to boundaries, according to the type of constant specified.
2. A duplication factor of zero is permitted with the following results:
a. No value is assembled.
(1) E. Alignment is fcrced according to the type of constant scecified, if no length attribute is present (see acove under Alignment of Constants) .
c. The length attrikute of the symbol naming the constant is established according to the
implicitly or explicitly specified

3. If duplication is specified for an address constant containing a location counter reference, the value of the location counter reference is incremented by the length of the constant before each duplication is performed (for examples, see G3J).

Subfield 2: Type

The type subfield must be specified. It defines the type cf constant to be generated and is specified ky a single letter code as in the figure to the right.

The type specification indicates to the assembler:

1. How the nominal value(s) scecified in subfield 4 is to be assembled; that is, which binary representation or machine format the cbject code of the constant must have.
2. At what koundary the assembler aligns the constant, if no length specification is present.
3. How much storage the constant is to occupy, according to the implicit length of the constant, if no explicit length specification is present (for details see above, under padding and Truncation of Constants) -

## Type

| Code | Type of Constant | Machine Format |
| :---: | :---: | :---: |
| C | Character | 8-bit code for each Character |
| $\times$ | Hexadecimal | 4 -bit code for each hexadecimal digit |
| B | Binary | Binary format |
| F | Fixed-point | Signed, fixed-point binary format; normally a fullword |
| H | Fixed-point | Signed, fixed-point binary format; normally a halfword |
| E | Floating-point | Short floating-point format ; normally a fullword |
| D | Floating-point | Long floating-point format; normally a doubleword |
| L | Floating-point | Extended floating-point format; normally two doublewords |
| P | Decimal | Packed decimal format |
| z | Decimal | Zoned decimal format |
| A | Address | Value of address; normally a fullword |
| Y | Address | Value of address; normally a halfword |
| S | Address | Base register and displacement value; a halfword |
| v | Address | Space reserved for external symbol addresses; each address normally a fullword |
| a | Address | Space reserved for external dummy section offset |
| $\begin{aligned} & \text { os } \\ & \text { only } \end{aligned}$ |  | $1$ |
|  |  | Object Code in hex |
|  | Examples: | $D C P^{\prime}+234{ }^{\prime} \quad 234 \mathrm{C}$ |
|  |  | DC C'ABC' C1C2C3 |
|  |  | DC X'FO' FO |
|  |  | DC H'2' 0002 |

## Subfield 3: Modifiers

The three modifiers that can be specified to describe a constant are:
(1) The length modifier (I), which explicitly defines the length in bytes desired for a constant.

2 The scale modifier (S), which is only used with the fixed-point or floatingpoint constants (for details see below under Scale Modifier).

3 The exponent modifier (E), that is only used with fixed-point or floating-point constants, and which indicates the power of 10 by which the constant is to be multiflied before conversion to its internal kinary format.

4 If multiple modifiers are used, the $\bar{y}$ must appear in the sequence: length, scale, exponent.

LENGTH MOLIFIER: The length modifier indicates the number of bytes of storage into which the constant is to be assembled. It is written as Ln , where n is either of the following:
(1) A decimal self-defining term
(2) An absolute expression enclosed in parentheses. It must have a positive value and any symbols it (3) contains must be previously defined.

Dos/ NOTE:. location counterf reference VSÆ nust not be used in the modifiler. subfield.


Length


BIT-IENGTH SEECIFICATICN: The length modifier can ke specified tc indicate the number cf bits intc which a constant is to be assembled. The kit-length specificaticn is written as L. $n$, where $n$ is either of the following:

A decimal self-defining term
An absolute expression enclosed in parentheses. It must have a positive value and any symbcls it contains must be previously defined.

The value of $n$ must lie ketween 1 and the number of bits (a multiple of 8) that are required to make up the maximum number of bytes allowed in the type of constant keing defined. The kit length-specification cannot be used with the $S, V$, and c-type constants.

When only one operand and one nominal value are specified in a rC instruction, the following rules apply:

1. The bit-length specification allocates a field intc which a constant is to be assembled.

2 The field starts at a byte boundary. and can run over one or more kyte boundaries, if the bit-length specified is greater than 8.

If the field does not end at a byte boundary, if the bit-length specified is not a multiple of 8 , the remainder of the last byte is filled with zeros.


1) 2. The nominal value cf the constant is assembled into the field:

2 Starting at the high order end for the $C, E, E$, and $L$ type constants.

3 Starting at the low order end for the remaining types of constants that allow bit-length specification.

The nominal value is padded or truncated to fit the field (see above under Padding or Truncation of Constants).
4. Padding of character constants is with hexadecimal blanks, X'40'; other constant types are padded with zeros.

NOTE: The length attribute value of the symbol naming a DC instruction with a specified bit-length is equal to the minimum number of integral bytes needed to contain the bitlength specified for the constant. I'TRUNCF is equal to 2. Thus, a
5 reference to TRUNCF would address the entire two bytes that are assembled.


L'VALUES=2
Source VALUES DC FL.10'161,21,57'
Code


Source OPERANDS DC FL.7'8',CL. $10^{\prime} \mathrm{AB}^{\prime}, \mathrm{XL} .14^{\prime} \mathrm{C} 4$
Code


STCRAGE REQUIREMENT FOR CONSTANTS:
The total amount of storage required to assemble a DC instruction is the sum of:

1) 2. The requirements for the individual DC operands specified in the instruction.

The requirement of a $L C$ operand is the product of:

2 a. The length (implicit or explicit).
(3) $b$. The number of nominal values. and

4 C. The duplication factor, if specified.

5 2. The number of bytes skipped for the koundary alignment between different operands.

SCALE MODIFIER: The scale modifier specifies the amount $c f$ internal scaling that is desired:

Binary digits for fixed-point (H,F) constants

Hexadecimal digits for flcatingpoint (E,D,I) constants

It can only be used with the above types of constant.

The scale modifier is written as Sn, where $n$ is either:
(1) A decimal self-defining term or
(2) An absolute expression enclosed in parentheses.

Dos Any symbols used in the expression must be previously defined.

Both types of specification can be preceded by a sign; if no sign is present, a plus sign is assumed.

Scale

assemkled in its final "scaled"
form. Scaling causes the binary
point to move from its assumed fixed
position at the right of the
rightmost bit position.


## NOTES:

1. When the scale modifier has a positive value, it indicates the number of binary positions to ke occupied ky the fracticnal pertion of the binary number.
2. When the scale medifier has a negative value, it indicates the number of binary positions to ke deleted from the integer forticn of the binary number.
3. When pcsitions are lcst kecause of scaling (or lack of scaling), rounding cccurs in the leftmostkit of the lost farticn. The rounding is reflected in the
(4)
rightmost fcsition saved.


SCALE MODIFIER FOR FLOATING-POINT CONSTANTS: The scale modifier for floating-point constants must have a positive value. It specifies the number of hexadecimal positions that the fractional portion of the kinary representation of a floatingpoint constant is to be shifted
2 to the right. The hexadecimal point is assumed to be fixed at the left of the leftrost position in the fractional field. when scaling is specified, it causes an unnormalized hexadecimal fraction to be assembled (unnormalized is when the leftmost positions of the fraction contain hexadecimal zeros). The magnitude of the constant is retained kecause the exponent in
3 the characteristic portion of the constant is adjusted upward accordingly. When hexadecimal positions are lost, rounding occurs in the leftmost hexadecimal position of the lost portion. The rounding is reflected in the rightmost position saved.

Source Code Object Code in Hex

DC E'4'

DC ES2'4'


DC E'3.3'
4134 CCCD



EXPONENT MOLIFIER: The exponent modifier specifies the power of 10 by which the nominal value of a constant is to be multiplied before it is converted to its internal binary representation. It can only be used with the fixedpoint ( $\mathrm{H}, \mathrm{F}$ ) and floating-point (E, D, L) constants. The exponent modifier is written as Fn , where $n$ can be either of the following:

## (1) A decimal self-defining term.

An absolute expression enclosed in parentheses.

Dos Any symbols used in the expression must be previously defined.

The decimal self-defining term or the expression can be preceded by a sign: if no sign is present, a plus sign is assumed. The range for the exponent modifier is - 85 through +75 .

NOTES:

1. The exponent modifier is not to be confused with the expcnent that can be specified in the nominal
2 value sutfield of fixed-coint and floating-point constants (see sections G3G and G3I).
(3)

The exponent modifier affects each nominal value specified in the operand, whereas the exponent written as part of the nominal value subfield only affects the ncrinal value it follows. If both types of exponent specification are present in a DC operand, their values are
(5) algekraically added together before the nominal value is converted to kinary form. However, this sum must lie within the permissible
(6) range -85 thrcugh +75 .
2. The value of the constant, after any exponents have been applied, must be contained in the implicitly or explicitly specified length of the constant to be assembled.

## Subfield 4: Nominal Value

The nominal value subfield must always be specified. It defines the value of the constant for constants) described and affected ky the subfields that precede it. It is this value that is assembled into the internal binary representation of the constant. The formats for specifying nominal values are described in the figure to the right.

Dosonly one nominal value is allowed in binary ( $B$ ) and hexadecimal ( $X$ ) constants.

How nominal values are specified and interpreted by the assembler is explained in the subsections that descrike each individual constant, beginning at G3I.


Nom. Value
Nom. Value

| Constant Type | Formats of Nominal Value Subfields |  |
| :---: | :---: | :---: |
|  | Single Nominal Values | Multiple <br> Nominal <br> Values |
| c | 'Value' | Not allowed |
| $\left.\begin{array}{l} B \\ X \\ H \\ F \\ P \\ Z \\ Z \\ E \\ D \\ L \end{array}\right\}$ | 'Value' | 'Value, value,......val ue,' <br> multiple values must be separated by commas |
| $\left.\begin{array}{l}\text { A } \\ \text { A } \\ \text { Y } \\ \text { S } \\ \text { Q } \\ \text { V }\end{array}\right\}$Address | (Value) | (Value, value,.......value) |

## Purpose

Literal constants allow you to
define and refer to data directly
in machine instruction operands.
You do not need to define a constant separately in another part of your source module. The difference between a literal, a data constant, and a self-defining term is described in C5.

## Specifications

A literal constant is specified in the same way as the operand of a DC instruction. The general rules for the operand subfields of a DC instruction (as described in G3B above) also apply to the subfield of a literal constant. Moreover, the rules that apply to the
individual types of constants, as described in G3D through G3M, apply to literal constants.

However, literal constants differ from LC operands in the fcllowing ways:

- Literals must be preceded by an equal sign.
(2) Multiple operands are nct allowed.
(3) The duplication factor must not be zero.
bos - q-type and s-type address constants are not allowed.



## Furpcse

The tinary constant allcws ycu tc specify the precise kit pattern ycu want assembled into storage.
Specifications

The constants of the sukfields defining a binary constant are described in the figure kelow.

1 NOTE: Each kinary constant is assembled intc the integral B number of bytes required to contain the lits specified.

| Subfield | Binary Constants |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | 3. Constant Type |  |  |  |
| 1. Duplication Factor | Binary (B) |  |  |  |
|  | Yes |  |  |  |
| 2. Modifiers <br> Implicit Length: (Length Modifier not present) | As needed <br> B DC B'lol01111' <br> C DC B'l01' |  | $\begin{aligned} & \mathrm{L} \cdot \mathrm{~B}=1 \\ & \mathrm{~L} \cdot \mathrm{C}=1 \end{aligned}$ |  |
| Alignment: <br> (Length Modifier not present) | Byte |  |  |  |
| Range for Length: | 1 through 256 (byte length) . 1 through . 2048 (bit length) |  |  |  |
| Range for Scale: | Not allowed |  |  |  |
| Range for Exponent: | Not allowed |  |  |  |
| 4. Nominal Value Represented by: | Binary digits (0 or 1) |  |  |  |
| Enclosed by: | Apostrophes |  |  |  |
| Exponent allowed: | No |  |  |  |
| Number of Values per Operand: | Multiple | Only one DOS |  |  |
| Padding: | With zer at left |  |  |  |
| Truncation of Assembled Value : | At left |  |  |  |

## Furpcse

The character constant allows ycu tc specify character strings such as error messages, identifiers, or cther text, that the assembler will convert into their kinary (EEC[IC) representation.

## Specificaticns

The contents of the sukfields defining a character corstant are descrited in the figure on the cppcsite page.
(1) Each character scecified in the norinal value subfield is assembled into one kyte.

Multifle nominal values are nct allcwed, because if a comma is specified in the nominal value sutfield, the assemtler
2 considers the comma a valid character and therefore assembles it intc its kinary (FBCDIC) representaticn.

NOTE: When apostrophes or ampersands are to be included in the assembled constant, doukle apostrcphes or dcutle ampersands must te specified. They are assembled as single (3) apostrophes and ampersands.

| Subfield | Character Constants |  |  |
| :---: | :---: | :---: | :---: |
|  | 3. Constant Type |  |  |
| 1. Duplication Factor | Character (C) |  |  |
|  | Yes |  |  |
| 2. Modifiers <br> Implicit Length: (Length Modifier not present) | As needed <br> C DC C'LENGTH' <br> 1 | $L^{\prime} \mathrm{C}=6$ |  |
| Alignment: <br> (Length Modifier not se present) | Byte |  |  |
| Range for length: | 1 through 256 (byte length) .1 through .2048 (bit length) |  |  |
| Range for Scale: | Not allowed |  |  |
| Range for Exponent: | Not allowed |  |  |
| 4. Nominal Value Represented by: | Characters (All 256 8 -bit combinations) | DC C'A' ${ }^{\prime} B^{\prime}$ <br> Assembled $A^{\prime}$ B $A \& B$ 3 <br> DC C'A\&\&B' | Object Code (hex). |
| Enclosed by: | Apost rophes |  |  |
| Exponent allowed: | No |  |  |
| Number of values per Operand: | One | $D C C^{\prime} A, B^{\prime}$ <br> Assembled A, B $\qquad$ 2 | $\begin{array}{\|l\|l\|l\|} \hline \mathrm{C} 1 & 6 \mathrm{~B} & \mathrm{C} 2 \\ \hline \end{array}$ |
| Padding: | With blanks at right ( $X^{\prime} 40^{\prime}$ ) |  |  |
| Truncation of Assembled value: | At right |  |  |

## Furcose

You can use hexadecimal constants to generate large rit patterns more conveniently than with binary constants. Also, the hexadecimal values you specify in a scurce ncdule allow ycu tc compare them directly with the hexadecimal values generated for the object code and address locaticns printed in the prcgram listing.

## Specificaticns

The contents of the sulfields defining a hexadecimal constant are descriked in the figure on the cpposite fage.
(1) Each hexadecimal digit specified in the nominal value sutfield is assembled intc four bits (their kinary patterns can be found in C4F). The implicit length in kytes cf a hexadecimal constant is then half the number of hexadecimal digits specified (assuming that a hexadecimal zerc is added to an odd number cf digits).


```
Purpose
Fixed-point constants allow you to introduce data that
is in a form suitable fcr the cperations cf the fixed-point
machine instructions of the standard instruction s\int.
The constants ycu define can alsc be autcmatically aligned
to the froper fullword or halfword koundary for the
instructions that refer tc addresses on these boundaries
    (unless the NCALGN option has keen specifitd; see [2).
You can ferform algekraic functicns using this type of
constant because they can have positive or negative values.
```


## Specifications

```
The contents cf the subfields defining fixed-point constants are described in the figure on the opposite fage.
1 The nominal value can te a signed (Flus is assumed if the number is unsigned) integer, fraction, or mixed number 2 followed by an exponent (Fositive or negative). The exponent must lie within the permissikle range. If an exponent modifier (see G3E) is alsc specified, the algekraic sum cf the exponent and the exponent modifier must lie within the permissikle range.
```

H or F

| Subfield | Fixed-Point Constants |  |  |
| :---: | :---: | :---: | :---: |
|  | 3. Constant Type |  |  |
| 1. Duplication Factor | Fullword(F) | Halfword (H) |  |
|  | Yes | Yes |  |
| 2. Modifiers <br> Implicit Length: (Length <br> Modifier not present) | 4 bytes | 2 bytes |  |
| Alignment: <br> (Length Modifer not present) | Full word | Half word |  |
| Range for Length: | 1 through 8 (byte length) .1 through . 64 (bit length) | 1 through 8 (byte length) .1 through . 64 (bit length) |  |
| Range for Scale: | -187 through +346 | - 187 through + 346 |  |
| Range for Exponent: | $-85 \text { through }+75$ | $\begin{array}{cc}-85 \text { through }+75 & \text { DC } \mathrm{HE}+90^{\prime} 2 \mathrm{E}-88^{\prime} \\ \text { value }=2 \times 10^{2}\end{array}$ |  |
| 4. Nominal Value Represented by: | Decimal digits ( 0 through 9 ) $\begin{aligned} & \text { DC F' } 200^{\prime} \\ & \text { DC FS4'2.25' } \end{aligned}$ | Decimal digits ( 0 through 9) $\begin{aligned} & \text { DC H' }+200^{\prime} \\ & \text { DC HS } 4^{\prime} .25^{\prime} \end{aligned}$ |  |
| Enclosed by: | Apostrophes | Apostrophes |  |
| Exponent allowed: | $\begin{aligned} & \text { Yes } \\ & \text { DC F'2E6' } \end{aligned}$ | $\begin{aligned} & \text { Yes } \\ & \text { DC H '2E-6' } \end{aligned}$ |  |
| Number of Values per Operand: | Multiple | Multiple |  |
| Padding: | With zeros at left | With zeros at left |  |
| Truncation of Assembled value: | (error message issued) |  |  |

Some examples of the range of values that can be assembled intc fixedpoint constants are given in the figure to the right.

The range of values depends on the implicitly or explicitly specified length (if scaling is disregarded). If the value specified for a particular constant does not lie within the allowable range for a given length, the constant is not assembled but flagged as an error.

| Length <br> 1 | Range of Values that <br> can be Assembled |  |  |
| :---: | :---: | :---: | :---: |
| 8 | $-2^{63}$ through $2^{63}-1$ |  |  |
| 4 | $-2^{31}$ | $"$ | $2^{31}-1$ |
| 2 | $-2^{15}$ | $"$ | $2^{15}-1$ |
| 1 | $-2^{7}$ | $"$ | $2^{7}-1$ |

A fixed-pcint constant is assembled as fcllows:

1. The specified number, multiplied ky any expenents, is converted to a kinary numker.
2. Scaling (see G3E) is perfcrmed, if specified. If a scale modifier is not provided the fracticnal portion of the number is lost.
3. The kinary value is rcunded, if necessary. The resulting number will not differ from the exact number specified ky more than cne in the least significant bit Fosition at the right.
4. A negative number is carried in $2^{\prime \prime}$ s complement form.
5. Luplication is apflied after the constant has been assembled.

G3H -- DECIMAL CONSTANTS (P AND Z)
Furccse

The decimal constants allow you tc intrcduce data that is in a form suitable for the operations of the decimal feature machine instructicns. The facked decimal constants (P-tyFe) are used for processing ky the deciral instructicn set. The zoned decimal constants (z-type) are in the form (EBCLIC representation) that you can use as a print inage (except the digits in the rightmost kyte).

## Specifications

The contents cf the sukfields defining decimal constants are described in the figure on the opposite fage.

The nominal value can ke a signed (clus is assumed if the number is unsigned) decimal number. A decimal foint can be written anywhere in the number, kut it does nct affect the assembly of the constant in any way. The specified 1 digits are assumed tc constitute an integer. Decimal constants are assemkled as follows:

PACKEL LECINAI CCNSTANIS: Each digit is converted into
2 its 4-bit binary equivalent. The sign indicator is assembled intc the rightmost four tits of the constant.

ZONEL DECIMAL CONSIANTS: Each digit is converted intc 4 its 8 -bit EECLIC representation. The sign indicatcr
replaces the first fcur bits cf the low-crder byte of the constant.

|  | Decimal Constants |  | Por Z |
| :---: | :---: | :---: | :---: |
| Subfield | 3. Constant Type |  |  |
| 1. Duplication Factor | Packed (P) | Zoned (Z) |  |
|  | Yes | Yes |  |
| 2. Modifiers <br> Implicit Length: (Length Modifier not present) | $\begin{aligned} & \text { As needed } \\ & \text { P DC P' }+593^{\prime} \\ & L^{\prime} P=2 \end{aligned}$ | $\begin{aligned} & \text { As needed } \\ & Z \text { DC } Z^{\prime}-593^{\prime} \\ & L^{\prime} Z=3 \end{aligned}$ |  |
| Alignment: <br> (Length Modifer not present) | Byte | Byte |  |
| Range for Length: | 1 through 16 (byte length) .1 through . 128 (bit length) | 1 through 16 (byte length) .1 through . 128 (bit length) |  |
| Range for Scale: | Not allowed | Not allowed |  |
| Range for Exponent: | Not allowed | Not allowed |  |
| 4. Nominal Value <br> Represented by: | Decimal digits (0 through 9) | Decimal digits (0 through 9) | $\begin{aligned} & \hline \text { DC } P^{\prime} 5.5^{\prime} \\ & \begin{array}{\|l\|l\|l\|l\|} \hline 0 & 5 & 5 & C \\ \hline \end{array} \\ & \text { DC } P^{\prime} 55^{\prime} \end{aligned}$ |
| Enclosed by: | Apostrophes | Apostrophes |  |
| Exponent allowed: | No | No |  |
| Number of Values per Operand: | Multiple | Multiple |  |
| Padding: | With Binary zeros at left | With EBCDIC zeros (X'FO') at left |  |
| Truncation of Assembled value: | At left | At left |  |

The range of values that can be assembled into a decimal constant is shown in the figure to the right.

| Type of Decimal <br> Constant | Range of Values that <br> can be Specified |
| :--- | :--- |
|  |  |
| PACKED | $10^{31}-1$ through $-10^{31}$ |
|  | $10^{16}-1$ through $-10^{16}$ |


#### Abstract

Furfcse

Floating-foint constants allow you to introduce data that is in a form suitable for the ccerations of the floatingpoint feature instruction set. These constants have tre following advantages cver fixed-fcint constants. 1. You do not have to consider the fractional portion cf a value ycu specify, ncr worry akcut the fcsiticn cf the decimal pcint when algetraic oferations are to be perfcrmed. 2. Ycu can specify koth much larger and much smaller values. 3. You retain greater processing precision, that is, ycur values are carried in more significant figures.


## Specificaticns

The contents of the subfields defining floating-pcint constants are descriked in the figure on the opposite page.
(1) The nominal value can be a signed (clus is assumed if the number is unsigned) integer, fracticn, cr mixed number 2 follcwed by an exponent (Fositive or negative). The
(3) exponent must lie within the permissitle range. If ar exponent nodifier (see G3E under Ncdifiers) is also specified, the algebraic sum of the exponent and the exponent modifier must lie within the $f \in r m i s s i k l e ~ r a n g \epsilon . ~$

D, E or L

| Subfield | Floating Point Constants |  |  |
| :---: | :---: | :---: | :---: |
|  | 3. Constant Type |  |  |
| 1. Duplication Factor | SHORT (E) | LONG (D) | EXTENDED (L) |
|  | Yes | Yes | Yes |
| 2. Modifiers <br> Implicit Length: <br> (Length Modifier Not Precent) | $4 \text { Bytes }$ | 8 Bytes | 16 Bytes |
| Alignment: <br> (Length Modifier Not Present) | Full Word | Double Word | Double Word |
| Range for Length: | 1 through 8 (byte length) . 1 through . 64 (bit length) | 1 through 8 (byte length) . 1 through . 64 (bit length) | 1 through 16 (byte length) .1 through . 128 (bit length) |
| Range for Scale: | 0 through 14 | 0 through 14 | 0 through 28 |
| Range for Exponent: | -85 through +75 | -85 through +75 | -85 through +7.5 |
| 4. Nominal Value <br> Represented by: | Decimal Digits <br> (0 through 9) <br> DC . $E^{\prime}+525^{\prime}$ <br> DC E'5.25! | $\begin{aligned} & \begin{array}{l} \text { Decimal Digits } \\ \text { (0through } 9) \end{array} \\ & \text { DC } D^{\prime-}-525^{\prime} \\ & D C D^{\prime}+.001^{\prime} \end{aligned}$ | Decimal Digits ( 0 through 9) DC L'525' DC L'3.414' |
| Enclosed by: | Apostrophes | Apostrophes | Apostrophes |
| Exponent Allowed: | $\begin{aligned} & \text { Yes } \\ & \text { DC } E^{\prime} 1 E+60^{\prime} 3 \end{aligned}$ | $\begin{aligned} & \text { Yes } \\ & \text { DC } D^{\prime}-2.5 E 10^{\prime} \end{aligned}$ | $\begin{aligned} & \text { Yes } \\ & \text { DC L'3.712E-3 } \end{aligned}$ |
| Number of Values per Operand: | Multiple | Multiple | Multiple |
| Padding: | With hexadecimal zeros at right | With hexadecimal zeros at right | With hexadecimal zeros at right |
| Truncation of Assembled Value: | Not applicable (Values are rounded) | Not Applicable (Values are Rounded) | Not applicable (Values are Rounded) |

The range of values that can be assembled into floating-point constants is given in the figure to the right.

If the value specified for a particular constant does not lie within these ranges, the constant is not assembled but flagged as an error.

FORMAT: The format of the floatingpoint constants is descriked below. The value of the constant is represented by two parts:

1) 2. An exponent portion, followed by
1) 2. A fractional portion.
(3) A sign bit indicates whether a positive or negative number has keen specified. The number specified must first be converted into a hexadecimal fraction, befcre it can be assembled into the proper internal format. The quantity expressed is the product of the
4 fraction and the number 16 raised 5 to a power.

| Type of <br> Constant | Range of Magnitude (M) <br> of Values (Positive and <br> Negative) |
| :---: | :--- |
| E | $16^{-65} \leq \mathrm{M} \leq\left(1-16^{-6}\right) \times 16^{63}$ |
| D | $16^{-65 \leq M \leq\left(1-16^{-14}\right) \times 16^{63}}$ |
| L | $16^{-65} \leq M \leq\left(1-16^{-28}\right) \times 16^{63}$ <br> (For all Three) <br> Approximately |
| $5.4 \times 10^{-79} \leq M \leq 7.2 \times 10^{75}$ |  |


| Type | Called | Format |
| :---: | :---: | :---: |
| E | Short <br> Floating- <br> Point <br> Number |  |
| D | Long <br> Floating- <br> Point <br> Number |  |
| L | Extended <br> Floating- <br> Point <br> Number |  |
| where $a, b, c$. . . are hexadecimal digits, and $E$ is an exponent that has a positive or negative value indicated by the characteristic |  |  |

BINARY REPRESENTATION: The assembler assembles a floating-point constant into its binary representation as follows:

The specified number, multiplied by any exponents, is converted to the required two-part format. The value is translated intc:
(1) 1. A fractional portion represented ky hexadecimal digits and the sign
2 indicator. The fraction is then entered into the leftmost part of the fraction field of the constant (after rounding).
(3)
2. An exponent portion represented
by the excess 64 binary notation, which is then entered into the characteristic field of the constant.

The excess 64 binary notation is when the value of the characteristic between +127 and +64 represents the exponents of 16 between +63 and 0 (by subtracting 64) and the value of the characteristic between +63 and 0 represents the exponents of 16 between -1 and -64 .

NOTES:

1. The I-type floating-point constant resembles two contiguous $D$-type constants. The sign of the second doubleword is assumed to be the same as the sign of the first.

The characteristic for the second doubleword is equal to the characteristic of the first doubleword minus 14 (the number of hexadecimal digits in the fractional portion of the first doukleword).
2. If scaling has been specified, hexadecimal zeros are added to the left of the normalized fraction (causing it to become unnormalized) and the exponent in the characteristic field is adjusted accordingly. (For further details on scaling see G3E under Modifiers).
3. Rounding of the fraction is performed according to the implicit or explicit length of the constant. The resulting number will not differ from the exact number specified ty more than one in the last place.
4. Negative fractions are carried in true representation, not in the 2's complement form.
5. Duplication is applied after the constant has been assembled.

| Binary Representation |  |  |
| :---: | :---: | :---: |
|  |  |  |
| Hexadecimal | Decimal | Decimal |
|  | $\left.\begin{array}{c}127 \\ 126 \\ ! \\ 65 \\ 64 \\ 63 \\ 62 \\ \cdot \\ 0\end{array}\right\}$ | $\begin{array}{r} +63 \\ +62 \\ \vdots \\ \vdots \\ +1 \\ \hline 0 \\ -1 \\ -2 \\ \vdots \\ -64 \end{array}$ |
| Values Expressed in Characteristic |  | Exponent of 16 expressed by Characteristic |
| Excess 64 Binary Notation |  |  |

This subsection and the three following suksections describe how the different types of address constants are assembled from expressions that usually represent storage addresses, and how the constants are used for addressing witrin ard ketw $\in \in$ n scurce mcdules.

## Furpose

In the A-type and Y-type address constant, ycu can scecify any of the three types of assembly-time expressions (see C6). whose value the assembler then computes and assembles into cbject code. You use this expression computaticr:
as follows:

1. Relocatable expressions for addressing
2. Absolute expressions for addressing and value computation.
3. Complex relocatable expressions to relate addresses in different scurce ricdules.

## Specificaticns

The contents of the sutfields defining the $A-t y p e$ and $y-$ type address constants are described in the figure on the opposite page.

NOTES:

1. 2. Ne bit-length specification is allowed when a relccatakle 2 or complex relocatakle expression is specified. The cnly explicit lengths that can be specified with these addresses are:
a. 3 or 4 bytes for A-type constants
b. 2 bytes for $Y$-type constants.
1. The value of the lcation counter reference (*) when specified in an address constant varies from constant tc constant, if any cf the fcllcwing cr a combination of the following are specified:
```
    a. Nultiple cFerands
```

(3) E. Multiple ncrinal values

4 C. A duplication factor.
The lccation counter is incremented with the length of the previcusly assembled constant.
3. When the location counter reference occurs in a literal address constant, the value of the location counter is the address of the first kyte of the instruction.

A or $\mathbf{Y}$

| Subfield | Address Constants ( A and Y ) |  |  |
| :---: | :---: | :---: | :---: |
|  | 3. Constant Type |  |  |
| 1. Duplication Factor | A - Type | Y - Type |  |
|  | Yes | Yes <br> Object Code in Hex | $\begin{aligned} & \text { A DC 5ALl }(*-\mathrm{A}) \\ & 0001020304 \end{aligned}$ |
| 2. Modifiers |  |  |  |
| Implicit Length: (Length Modifer not present) | 4 bytes | 2 bytes |  |
| Alignment: <br> (Length Modifier not present) | Full word | Half word |  |
| Range for Length: | 1 through 4 (byte length) . 1 through . 32 (bit length) | 1 through 2 (byte length) .1 through . 16 (bit length) |  |
| Range for Scale: | Not allowed | Not allowed |  |
| Range for Exponent: | Not allowed | Not allowed |  |
| 4. Nominal Value Represented by: | Absolute, relocatable, or $\left.\begin{array}{l}\text { complex relocatable } \\ \text { expressions }\end{array}\right\}$ DC A (ABSOL+10) | $\begin{aligned} & \left\{\begin{array}{l} \text { Absolute, relocatable, or } \\ \text { complex relocatable } \\ \text { expressions } \end{array}\right. \\ & \text { DC Y(RELOC }+32) \end{aligned}$ | 3 $\begin{gathered} A \text { DC } Y(*-A, *+4) \\ 0 \quad A+6 \\ \text { values } \end{gathered}$ |
| Enclosed by: | Parentheses | Parentheses |  |
| Exponent allowed: | No | No |  |
| Number of Values per Operand: | Multiple | Multiple |  |
| Padding: | With zeros at left | With zeros at left |  |
| Truncation of Assembled value: | At left | At left |  |

CAUTION: Specificaticn of $Y$-type address constants with relocatable expressions should be avoided in frograms that are to be executed on machines having more than 32,767 kytes of stcrage capacity. In any case, Y-type relccatable address constants should not te used in programs tc $k \in$ executed under IEM System/370 contrcl.

The A-type and Y-type address constants are processed as follows: If the ncrinal value is an aksolute expression, it is computed to its $32-$ bit value and then truncated on the left to fit the implicit or explicit length of the constant. If the ncrinal value is a relccatable cr complex relocatable expression, it is not completely evaluated until linkage edit time when the cbject mcdules are transformed into load modules. The 24 -上it (cr smaller) relocated address values are then placed in the fields set aside far then at assemkly time by the $A-t y p e$ and $y$ type constants.

Furcose

You can use the s-type address constant tc assemble an explicit address that is, an address in kase-disflacenent form). You can specify the explicit address yourself cr allow the assembler to compute it from an implicit address, using the current base register and address in its computation (for details cn implicit and explicit addresses, see [5B) .

## Specificaticns

The contents cf the sutfields defining the S-type address constants are descriked in the figure on the oppcsite fage.

The ncminal values can ke specified in two ways:
(1) 1. As one absolute or relocatakle expression representing an implicit address
22. As two absolute expressions, the first of which represents the displacement and the second, the base 4 register.

| Subfield | Address Constants (S) |  |  |
| :---: | :---: | :---: | :---: |
|  | 3. Constant Type |  |  |
| 1. Duplication Factor | S - Type |  |  |
|  | Yes |  |  |
| 2. Modifiers <br> Implicit Length: <br> (Length Modifier not present) | 2 bytes |  |  |
| Alignment: <br> (Length Modifier not present) | Half word |  |  |
| Range for length: (in bytes) | 2 only (no bit length) |  |  |
| Range for Scale: | Not allowed |  |  |
| Range for Exponent: | Not allowed |  |  |
| 4. Nominal Value <br> Represented by: | Absolute or relocatable expression <br> Two absolute expressions | ```DC S(RELOC) DC S(1024)} DC S(512(12))``` |  |
| Enclosed by: | Parentheses |  |  |
| Exponent allowed: | No |  |  |
| Number of Values per operand: | Multiple |  |  |
| Padding: | Not applicable |  |  |
| Truncation of Assembled value: | Not applicable |  |  |

## Purfcse

The V-type address constant allows you tc reserve stcrage
for the address cf a lccation in a contrcl section that lies in another source module. You should use the v-type address constant cnly tc branch to the external address specified. This use is contrasted with anotrer methcd, that is: of specifying an external symbcl, identified by an EXTRN instruction, in an A-type address constant (for a comparison, see F2).

Because you specify a symbol in a V-type address constant, the assembler assumes that it is an external symbol. A value of zero is assembled into the space reserved for the v-type constant; the correct relocated value of the address is inserted into this space by the linkage editor before your object program is loaded.

| Source Module | Object Module in Hex |  | Load Module in Hex |
| :---: | :---: | :---: | :---: |
|  | A |  | A. |
| A START 0 |  |  |  |
| DC V (OUTSIDE) | 00000000 | 3 | 00003000 |
| END |  | 2 |  |
|  | OUTSIDE | X'3000' | OUTSIDE |

## Specificaticns

The contents of the subfields defining the $V$-type address constants are described in the figure on the oppcsite page.

The symbcl specified in the nominal value subfield dces not constitute a definition of the symkol for the scurce module in which the v-type address constant appears.

The symbcl specified in a V-type constant rust nct refresent external data in an cverlay frcgram.

| Subfield | Address Constants (V) |  |  |
| :---: | :---: | :---: | :---: |
|  | 3. Constant Type |  |  |
| 1. Duplication Factor | V - Type |  |  |
|  | Yes |  |  |
| 2. Modifiers <br> Implicit I.ength: (Length <br> Modifier not present) | 4 bytes |  |  |
| Alignment: (Length Modifier not present) | Full word |  |  |
| Range for Length: (in bytes) | 4 or 3 only (no bit length) |  |  |
| Range for Scale: | Not allowed |  |  |
| Range for Exponent: | Not allowed |  |  |
| 4. Nominal Value Represented by: | A single relocatable symbol | DC V(MODA) 1 <br> DC V(EXTADR) |  |
| Enclosed by: | Parentheses |  |  |
| Exponent allowed: | No |  |  |
| Number of values per Operand: | Multiple |  |  |
| Padding: | With zeros at left |  |  |
| Truncation of assembled value: | Not applicable |  |  |

## Furpcse

You use this constant tc reserve stcrage for the cffset
into a storage area of an external dummy secticn. The
offset is entered intc this scace $k y$ the linkage editcr.
When the offset is added to the address of an overall blcck
of storage set aside fcr external dummy sections, it allows
you tc address the desired section. (For a descripticr.
of the use cf the Q-type address constant in combination
with an external dummy section, sé F 4.$)$
Specificaticns
The contents cf the surfields defining the $\mathcal{C}$-type address constant are descriked in the figure kelcw.
(1) The symkol specified in the nominal value subfield must be freviously defined as the latel of a IXI or ESECT statement.

| Subfield | Address Constants (Q) |  |  |
| :---: | :---: | :---: | :---: |
|  | 3. Constant Type |  |  |
| 1. Duplication Factor allowed | Q-Type |  |  |
|  | Yes |  |  |
| 2. Modifiers <br> Implicit Length: (Length Modifier not present) | 4 bytes |  |  |
| Alignment: (Length <br> Modifier not present) | Fullword |  |  |
| Range for Length: (in bytes) | $1-4$ bytes (no bit length) |  |  |
| Range for Scale: | Not allowed |  |  |
| Range for Exponent: | Not allowed |  |  |
| 4. Nominal Value <br> Represented by | A single relocatable symbol | $\begin{aligned} & \text { DC } Q(D U M M Y E X T) \\ & D C \text { Q(DXDEXT) } \end{aligned}$ |  |
| Enclosed by: | Parentheses |  |  |
| Exponent allowed: | No |  |  |
| Number of Values per Operand: | Multiple |  |  |
| Padding: | With zeros at left |  |  |
| Truncation of Assembled Value | At left |  |  |

## Purpose

The LS instruction allows you to:

1. Reserve areas cf stcrage
2. Provide latels fcr these areas
3. Use these areas ky referring to the symbols defined
as latels.
The LS instruction causes no data tc be assembled. Unlike the LC instruction (see G3E), you do nct have tc specify the nominal value (fcurth subfield) of a $L S$ instruction operand. Therefore, the $I S$ instruction is the best way of symbolically defining storage for work areas, input/output buffers, etc.

How to Use the DS Instruction

TO RESERVE STORAGE; If ycu wish to take advantage of automatic boundary alignment (if the ALIGN option is specified) and implicit
2 length calculation, you should not supply a length modifier in your operand specifications. You should specify a type subfield that corresponds to the type of area you need for your instructions (See individual types in sections G3D through G3M).

| Named (Mnemonic) Areas for FixedPoint Instructions | Areas Aligned on Boundary <br> 1 | Length Attribute of Symbols Naming Areas same as Implicit Length of Areas |
| :---: | :---: | :---: |
| FAREA DS F | Full word | 4 |
| HAREA DS H | Half word | 2 |
| AAREA DS A | Full word | 4 |
| DUPF DS 10F | Eull word | L'DUPF=4 |
| 10 full words of storage reserved |  | Juplication has no effect on implicit length |
| Named Areas for Floating-Point Instructions |  |  |
| EAREA DS 3E | Full word | 4 |
| DEAREAS DS 9D | Double word | 8 |
| 9 double words reserved |  |  |
| LAREA DS L | Double word | 16 |

Using a length modifier can give you the advantage of explicitly specifying the length attribute value assigned to the label naming the area reserved. However. your areas will not be aligned automatically according to their
 type. If you omit the nominal value in the operand, you should use a length modifier for the binary (E). character (C), hexadecimal (X). and decimal ( F and Z) type areas; otherwise their labels will be given a length attribute value of 1 .


1. When you need to reserve large areas ycu can use a
duplication factor. However, you can only refer tc tre
2 first area ky the lakel in this case. Ycu can also use the character (C) and hexadecimal (X) fíld tyfes tc sfecify
(3) large areas using the length rcdifier.


Although the nominal value is optional for a DS instruction, you can put it to good use by letting the assembler compute the length for areas of the $E, C, X$, and decimal ( P or Z ) type areas. You achieve this by specifying the general
2) format of the nominal value that will be placed in the area at execution time.

| Area Specified | Area Reserved in bytes | Length Atribute or computed implicit length of area (duplication disregarded) <br> 1 |
| :---: | :---: | :---: |
| Cl DS C'THIS IS AN ERROR' | 16 | 16 |
| XI DS X'OAA' ${ }^{\prime}$ | 2 | 2 |
| X2 DS 30X'F1F2' | 60 | 2 |
| P1 DS P'99999' | 3 | 3 |
| P2 DS 5P'99999' 2 | 15 | 3 |
| Z1 DS Z'99999' | 5 | 5 |

TO FORCE ALIGNMENT: Ycu can use the DS instruction to force alignment to a boundary that otherwise would not be provided. You do this ky using a duplication factor of zero. No space is reserved for such an instruction, yet the data that follows is aligned on the desired
(3) koundary.

NOTE: Alignment is forced when either the ALIGN or NOALIGN assembler option is set (see [2).


TO NAME FIELES OF AN AREA: Using a duplication factcr of zero in a LS instruction also allows you to ercvide a lakel for an area cf stcrage withcut actually reserving
the area. You can use LS or LC instructions to reserve storage for and assign lakels tc fields within the area. These fields can then ke addressed symbolically. (Ancther way of accomplishing this is described in E3C.) The whole area is addressable ky its lakel. In addition, the symbclic
label will have the length attrikute value of the whole
area. Within the area each field is addressable by its
label. The LATE field has the same address as the sutfield
4 LAY. However, LATE addresses 6 kytes, while DAY addresses only 2 bytes.


Specificaticns

The format of the $I S$ instruction statement is given in the figure to the right.

1 The format of the cferand of a DS instruction is identical to that of the $L C$ crerand (sfe G3E).


The two differences in the specification of sukfields are:

The nominal value sukfield is opticnal in a LS operand, but it is mandatcry in a LC cperand. If
2 a nominal value is specified in a LS operand, it must ke valid.
(3) The maximum length that can be specified in a IS operand for the character (C) and hexadecimal (X) type areas is 65,535 bytes, rather than 256 bytes for the same IC operands.

| OPTION1 | DS | $3 F L 3$ |
| :--- | :--- | :--- |
| OPTION2 | DS | $3 F L 3: 31$ |
| AMUST | DC | $3 F L 3: 31$ |
|  |  | 2 |
| LONGC | DS | CL65535 |
| LONGX | DS | XL65535 |

$\begin{array}{ll}\text { LIMITEDC DC } & \text { CL256'A' } \\ \text { LIMITEDX DC } & \text { XL256'00' }\end{array}$

The label used in the name entry of a 5 IS instruction, like the label fcr a $L C$ instruction (see G3E):

1) 2. Has an address value of the leftmost byte of the area reserved, after any kcundary alignment is performed
2. Has a length attrikute value, decending on the implicit or explicit length of the type of area reserved.

If the IS instruction is specified with more than one cperand or more than one nominal value in the operand, the label addresses the area reserved for the field that
 corresponds to the first nominal value of the first operand.
(4) The length attribute value is equal to the length explicitly specified or implicit in the first ccerand.


NOTE: Unlike the DC instruction,
(1) bytes skipped for alignment are not set to zero. Also, ncthing
2 is assembled into the storage area reserved by a DS instruction. No assumption should be made as to the contents of the reserved area.

The size of a storage area that can te reserved by a DS instruction is limited only by the size of virtual storage or by the maximum value of the location counter, whichever is smaller.


## Furpcse

You can use the CCW instruction to define and generate an eight-kyte channel comrand word for infut/output operations.

The channel command word is an eight-byte field aligned at a doukleword koundary, and contains the information described in the figure to the right.


Specifications

The format of the CCW instruction statement is given in the figure to the right.

(1) All fcur cperands must ke specified in the order descrited in the figure to the right. The generated channel command word is aligned on a
2 doubleword koundary. Any bytes skipped are set tc zerc.

The symbol in the name field, if present, is assigned the value of
(3) the address of the leftmost lyte of the channel command word generated. The length attribute value of the symbel is 8.

## Section H: Controlling the Assembler Program

This section describes the assembler instructions that request the assembler to perform certain functions that it would otherwise perform in a standard predetermined way. You can use these instructions to:

1. Change the standard coding format for writing your source statements
2. Control the final structure of your assembled program
3. Alter the format of the source module and object code printed on the assembler listing
4. Produce punched card output in addition to the object deck
5. Substitute your cwn mnemonic operation codes for the standard codes of the assembler language
6. Save and restore frogramming environments, such as the status of the PRINT options and the USING base register assignment.

## Hl -- Structuring a Program

The instructions described in this subsection affect the location counter and thereby the structure of a control section. You can use them to interrupt the normal flow of assembly and redefine portions of a control section or to reserve space to receive literal constants. Alsc, you can use them to align data on any desired boundary.

## Purpose

You use the ORG instruction to alter the setting of the location counter and thus control the structure of the current control section. This allows you to redefine portions of a control section.

For example, if you wish to build a translate table (to convert EBCDIC character code intc some other internal code) :

1) 2. You define the table as keing filled with zeros.
2. You use the ORG instruction to alter the location counter so that its counter value indicates a desired location within the table.
3. 3. You redefine the data to be assembled intc that location.
1. After repeating the first three steps until your translate table is complete, you use an ORG instruction with a blank operand field to alter the location counter so that the counter value indicates
(5) the next available location in the current control section (after the end of the translate table).

Both the assembled object code for the whole table filled with zeros and the object code for the portions of the table you redefined are printed in the program listings. However, the data defined later is loaded over the previously defined zeros and becomes part of your object program, instead of the zeros.

In other words, the ORG instruction can cause the location counter to point to any part of a control section, even the middle of an instruction, into which you can assemble desired data. It can also cause the location counter to point to the next available location so that your program can continue to be assembled in a sequential fashion.

Specifications

The format of the ORG instruction is shown in the figure to the right.

The symbols in the expression in the operand field must be previously defined. The unpaired relocatable term of the expression (see C6E) must be defined in the sare control section in which the ORG statement appears.
(1)

The location counter is set to the value of the expression in the operand. If the operand is omitted, the location counter is set to the
(3) next available location for the current control section.

| Name | Operation | Operand |
| :--- | :--- | :--- |
| OS <br> Any symbol <br> or blank <br> DOS <br> Sequence symbol <br> or blank | ORG | A relocatable <br> expression <br> or blank |



The expression in the ocerand of an ORG instruction must not specify a location before the beginning of the control section in which it appears. In the example to the
(1) right, the CRG instruction is invalid if it appears between the beginning of the current control section and 500 kytes from the beginning of the same control section. This is because the expression specified is then negative and will set the location counter to a value larger than the assembler can process. The location counter will "wrap around" (the location counter is discussed in detail in section C4B).

NOTE: Using the ORG instruction to insert data assembled later at the same location as earlier data will not always work.

In the example to the right, it
1 appears as if the character constant will te loaded over the address constant. However, after the character constant is loaded into
(3) the constant, the relocation factor required for the address constant
4 is added to the value of the constant. This sum then constitutes the object code that resides in the four bytes with the address ADER.

H1B -- THE ITORG INSTRUCTION


## Purpose

You use the LTORG statement so that the assembler can collect and assemble literals into a literal pool. A literal pccl contains the literals you specify in a source module either:
(1)

After the preceding LTORG instruction or

After the beginning of the source module.

The assembler ignores the borders between control sections when it collects literals into pools. Therefore, you must be careful to include the literal pools in the control sections to which they belong (for details see Addressing Considerations belcw).

The creation of a literal pcol gives the following advantages:

1. Automatic organization of the literal data into sections that are properly aligned and arranged so that no space is wasted
2. Assembling of duplicate data into the same area
3. Because all literals are crossreferenced, you can find the literal constant in the pool into which it has keen assembled.

The Literal Pool

A literal pool is created immediately after a LTORG instruction or, if no LTCRG instruction is specified, at the end of the first control section.

Each literal pool has four segments into which the literals are stored (1) in the order that the literals are specified and (2) according to their assembled lengths, which, for each literal, is the total explicit or implicit length, as described below.
(1) The first segment contains all literal constants whose assembled lengths are a multiple of eight.
(2) The second segment contains those whose assembled lengths are a multiple of four, but not of eight.
(3) The third segment contains those whose assembled lengths are even, but not a multiple of four.
(4) The fourth segment contains all the remaining literal constants whose assembled lengths are odd.


The keginning of each literal pocl is aligned on a doubleword boundary. Therefore, the literals in the first segment are always aligned on a doubleword koundary, those in the second segment on a fullword
boundary, and those in the third segment on a halfword boundary.

## Addressing Considerations

If you specify literals in source modules with multicle control sections, you should:

1. Write a LTORG instruction at the end of each control section, so that all the literals specified in the section are assembled into the one literal pool for that section. If a contrcl section is divided and interspersed among other control sections, you should write a LTCRG instruction at the end of each segment of the interspersed control section.
2. When establishing the addressability of each control section, make sure (a) that the entire literal pool for that section is also addressable, by including it within a USING range, and (b) that the literal specifications are within the corresponding USING domain. The USING range and domain are described in F1A.

NOTE: All the literals specified after the last LTORG instruction, or, if no LTORG instruction is specified, all the literals in a source module are assembled into a literal pool at the end of the first control section. You must then make this literal pool addressable along with the addresses in the first control section. This literal pool is printed in the program listing after the END instruction.


If you specify duplicate literals within the part of the source module that is controlled by a LTORG instruction, only one literal constant is assembled into the pertinent literal pccl. This also applies to literals assembled into the literal pool at the end of the first or only control section of a source module that contains no LTORG instructions.

Literals are duplicates only if their specifications are identical, 3 not if the object code assembled happens to be identical.

When two literals specifying identical A-type (or Y-type) address constants contain a reference to the value of the location counter (*) , both literals are assembled into the literal pool. This is because the value of the location counter is different in the two literals.


The format of the LTORG instruction is given in the figure to the right.

If an ordinary symbol is specified in the name field, it represents the first byte of the literal pool; this symbol is aligned on a doubleword koundary and has a length attribute value of one. If kytes are skipped after the end of a literal pool to achieve alignment for the next instruction, constant, or area, the bytes are not filled with zeros.

LTORG

| Name | Operation | Operand |
| :--- | :--- | :--- |
| Any symbol <br> or blank | L'TORG | Not required |

H1C - THE CNOP INSTRUCTION

## Purpose

You can use the CNOP instruction to align any instruction cr other $(1$ data on a specific halfword coundary. The CNOP instruction ensures an unbroken flow of executable instructions by generating no-
2 operation instructions to fill the bytes skipped to perform the alignment that you specified.

For example, when you code the linkage to a subroutine, you may
 wi sh to pass parameters to the subroutine in fields immediately following the branch and link
4 instruction. These carameters, for instance, channel command words (see G30), can require alignment on a specific boundary.


The subroutine can then address the parameters you pass through the register with the return address.


## Specificaticns

The CNOF instruction forces the alignment of the lcation counter to a halfword, fullword, or doubleword scundary. It does not affect the location counter if the counter is already frcperly aligned. If the specified alignment requires the location counter tc be incremented. one to three nooperation instructicns ( $B C R$ 0,0 occupying two bytes each) are generated to fill the skipped kytes. Any single kyte skipced to achieve alignment to the first no-operation instruction is filled with zeros.

The format of the CNOP instruction statement is given in the figure to the right.

The operands must be absolute expressions, and any symbols must have been previously defined. The first operand, b , specifies at which even-numbered byte in a fullword or doubleword the location counter is set. The second operand, w , specifies whether the byte is in a fullword ( $w=4$ ) or a doubleword ( $w=8$ ) . Valid pairs of $b$ and $w$ are as indicated in the figure to the right.

NOTE: Both 0,4 and 2,4 specify two locations in a doubleword.

| Name |  |  | Operation |  |  | Operand |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\begin{gathered} \text { os } \\ \text { oos } \end{gathered}$ | Any sym or blank <br> Sequence <br> or blank | bol <br> symb | CNOP |  |  |  |  |
| $0^{0,4}$ |  |  |  |  |  |  |  |
| HALFWORD |  | HALFWORD |  | HALFWORD |  | HALFWORD |  |
| Byte | Byte | Byte | Byte | Byte | Byte | Byte | Byte |
| ${ }_{\text {¢ }}^{1}$ | 2,8 | $\qquad$ |  |  | $\underset{6,8}{\uparrow}$ |  |  |

## H2 -- Determining Statement Format and Sequence

You can change the standard coding conventions for the assembler language statements or check the sequence of source statements by using the fcllowing instructions.

H2A -- THE ICTL INSTRUCTION

## Purpose

The ICTL instruction allows you to change the begin, end, and continue columns that establish the coding format cf the assembler language source statements.

For example, with the ICTL instruction, you can increase the number of columns to be used for the identificaticn or sequence checking of your source statements. By changing the begin column, you can even create a field before the begin column to contain identification or sequence numbers.

You can use the ICTL instruction only once, at the very beginning of a source module. If you do not use it, the assembler recognizes the standard values for the, kegin, end, and continue columns.


## Specifications

The ICTL instructicn, if specified, must be the first statement in a source module.

The format of the ICTL instruction statement is shown in the figure to the right.

The operand entry must be one to three decimal self-defining terms. There are only three possible ways of specifying the operand entry.

The operand $b$ must always be specified. The operand e, when not specified, is assumed to be 71. If the operand $c$ is not specified, or if $e$ is specified as 80 , the assembler assumes that continuation lines are not allowed. The values specified for the three operands depend on each other.

NOTE: The ICTL instruction does not affect the format of statements brought in ty a COPY instruction or generated from a library macro definition. The assembler processes these statements according to the standard kegin, end, and continue columns described in Section E1A.

## purpose

You can use the ISEQ instruction to cause the assembler to check if the statements in a source module are in sequential order. In the ISEQ instruction you specify the

1) columns between which the assembler is to check for sequence numbers.

The assembler begins sequence checking with the first statement line following the ISEQ instruction. The assembler also checks


Sequence numbers on adjacent
4 statements or lines are compared according to the 8-bit internal EBCDIC collating sequence. When the sequence number on one line is not greater than the sequence number on the preceding line, a sequence error is flagged, and a warning message is issued, but the assembly is not terminated.

NOTE: If the sequence field in the preceding line is blank, the assembler uses the last preceding line with a non-blank sequence field to make its comparison.

## Specifications

The ISEC instruction initiates or terminates the checking of the sequence of statements in a source module.

The format of the ISEQ instruction is shown in the figure to the right.
(1) The first option in the operand entry must be two decimal selfdefining terms. This format of the ISEQ instruction initiates sequence checking, beginning at the statement or line fcllowing the ISEQ instruction. Checking begins at the column represented by 1 and ends at the column represented by $r$. The second option of the ISEQ format terminates the sequence checking operation.


ISEQ
$\left.\begin{array}{|c|c|l|}\hline \text { Name } & \text { Operation } & \text { Operand } \\ \hline \text { Blank } & \text { ISEQ } & 1, r \quad 1 \\ & & \text { or blank }\end{array}\right\}$

\begin{tabular}{|c|c|c|c|}
\hline Column \& Specifies \& \multicolumn{2}{|l|}{Rules for interaction} \\
\hline  \& \begin{tabular}{l}
leftmost column of field to be checked \\
I and \(r\) not allowed to lie between begin and end columns \\
rightmost column of field to be checked
\end{tabular} \& \(1 \leq r\)

$r \geq 1$ \& | I must not be greater than $\mathbf{r}$ |
| :--- |
| r must not be less than I | <br>

\hline
\end{tabular}

NOTE: The assembler checks only those statements that are specified in the coding of a scurce module. This includes any COPY instruction statement or macro instruction.

However, the assembler does not check:

1. Statements inserted by a COPY instruction
2. Statements generated from model statements inside macro definitions or from model statements in open code (statement generation is discussed in detail in Section J)
3. Statements in library macro definitions.

## H3 .- Listing Format and Output

The instructions described in this section request the assembler to produce listings and identify output cards in the object deck according to your special needs. They allow you to determine printing and page formatting options other than the ones the assembler program assumes ky default. Among other things, you can introduce your own page headings, control line spacing, and suppress unwanted detail.

H3A -- THE PRINT INSTRUCTION

## Purpose

The FRINT instruction allows you to control the amount of detail you wish printed in the listing of your programs. The three options that you can set are given in the figure to the right.

They are listed in hierarchic order; if OFF is specified, GEN and DATA will not apply. If NOGEN is specified, $L A T A$ will not apply to constants that are generated. The standard options inherent in the assembler program are $\mathrm{CN}, \mathrm{GEN}$, and NODATA.

$\sim$

| Hierarchy | Description | PRINT options |
| :--- | :--- | :--- |
| 1 | A listing is printed | ON |
| No listing is printed | OFF |  |
| 2 | All statements generated by the <br> processing of a macro instruction <br> are.printed | GEN |
| Statements generated by the <br> processing of a macro instruction <br> are not printed (Note: The <br> MNOTE instruction always causes <br> a message to be printed) | NOGEN |  |
| 3 | Constants are printed in full in <br> the listing | DATA |
| Only the leftmost eight bytes of <br> constants are printed in the <br> listing | NODATA |  |


| Name | Operation | Operand |
| :--- | :--- | :--- |
| A sequence <br> symbol or <br> blank | PRINT | $\left[\begin{array}{cc}{\left[\begin{array}{l}\text { ON } \\ \text { OFF }]\left[\begin{array}{c}\text { GEN } \\ \text { NOGEN }\end{array}\right]\left[\begin{array}{c}\text { NODATA } \\ \text { DATA }\end{array}\right]\end{array}\right.} \\ \hline\end{array}\right.$ |
| Any sequence of <br> specification allowed |  |  |

The format of the PRINT instruction statement is shown in the figure to the right.
(1) At least one of the operands must be specified, and at most one of the options from each group. The PRINT instruction can be specified any number of times in a source module, but only those print options actually specified in the instruction change the current print status.

PRINT options can be generated by macro processing, at pre-assembly time. However, at assembly time, all options are in force until the assembler encounters a new and opposite option in a PRINT instruction.

OS The FUSH and FOP instructions,
only described in 46 , also influence the FRINT options by saving and restoring the PRINT status.

NOTE: The option specified in a PRINT instruction takes effect after the FRINT instruction. If PRINT OFF is specified, the PRINT instruction itself is printed, but not the statements that follow it. If the NOLIST assembler oftion is specified in the job control language, the entire listing for the assembly is suppressed.

TITLE 'THIS IS A HEADING'


THIS IS A HEADING

Purpose

The TITLE instruction allows you to:
(1) 1. Provide headings for each page of the assembly listing of your source modules.
2. Identify the assembly output cards of your object modules. You can specify up to 8 identification
(2) characters that the assembler will punch into all the output cards, 3 keginning at column 73.

Dos Up to 4 identification characters are allowed.
(4.

The assembler punches sequence numbers into the colums that are left, up to column 80.

NOTE: The name field of the TITLE instruction is generated throughout the assembly listing, preceding the generation of the operand data from any TITLE instruction.

## Specifications

The format of the TITLE instruction statement is given in the figure to the right.

Any of the five options can be specified in the name field.
(1) The first three options for the name field have a scecial significance only for the first TITLE instruction in which they are specified. For subsequent TITLE instructions, the first three options do not apply.

TITLE

|  | Name | Operation | Operand |
| :---: | :---: | :---: | :---: |
| $\text { (1) }\left\{\begin{array}{l} 1 \\ 2 \\ 3 \end{array}\right.$ | A string of alphameric characters A variable symbol A combination of 1 and 2 <br> A sequence symbol blank | TITLE | A character string up to 100 characters, enclosed in apostrophes |

For the first TITLE instruction of a source module that has a nonblank name entry that is not a sequence symbol, the following applies:
(1) Up to eight alphameric characters can ke specified in any combination in the name field.

DOS Up to four alphameric characters can be specified.

These characters are punched as identification, beginning at column 73, into all the output cards from the assembly, except those produced by the PUNCH and REFRO instructions. The assembler substitutes the current
(2) value into a variable symbol and uses the generated result as identification characters.
(3) If a valid ordinary symbol is specified, its appearance in the name field does not constitute a definition of that symbol for the source module. It can therefore be used in the name field of any other statement in the same source module.


The character string in the operand
(1)
the top of each page of the assembly listing. The heading is printed beginning on the page in the listing following the page on which the TITLE instruction is specified. A new heading is printed when a subsequent TITLE instruction appears in the source module.

Each TITLE statement causes the listing to be advanced to a new page (before the heading is printed) except when PRINT NOGEN is in use.

Any printable character specified will appear in the heading, including blanks. Variable symbols are allowed. However, the following rules apply to ampersands and apostrophes:

- A single ampersand initiates an attempt to identify a variable
(2) symbol and to substitute its current value.

3. Double ampersands or apostrophes specified, print as single ampersands or apostrophes in the heading.
4. A single apostrophe followed by one or more blanks simply terminates the heading prematurely. If a non-blank character follows a single apostrophe, the assembler issues an error message and prints no heading.

Only the characters printed in the heading count toward the maximum of 100 characters allowed.

NOTE: The TITLE statement itself is not printed in an assembly listing.


## Purpose

The EJECT instruction allows you to stop the printing of the assembly listing on the current page and continue the printing on the next page.

## Specifications

The format of the EJECT instruction statement is shown in the figure to the right.

The EJECT instruction causes the next line of the assembly listing to be printed at the top of a new 1 page. If the line before the EJFCT instruction appears at the bottom of a page, the EJECT instruction has no effect An EJECT instruction immediately following another EJECT instruction causes a blank page in the listing.

NOTE: The EJECT instruction statement itself is not printed in the listing.


## Purpose

You can use the SPACE instruction to insert one or more blank lines in the listing of a source module. This allows you to separate sections of code on the listing fage.

SPACE

## Specifications

The format of the SPACE instruction statement is given in the figure to the right.

The operand entry specifies the number of lines to be left blank. A blank operand entry causes one blank line to be inserted. If the operand specified has a value greater than the number of lines remaining on the listing page, the instruction will have the same effect as an EJECT statement.

NOTE: The SPACE instruction itself is not listed.

## H4 -- Punching Output Cards

The instructions described in this section produce punched cards as output from the assembly in addition to those produced for the object module (object deck).

H4A -- THE FUNCH INSTRUCTION

Purpose

The PUNCH instruction allows you to punch source or other statements into a single card. With this feature you can:

1. Code PUNCH statements in a source module to produce control statements for the linkage editor. The linkage editor uses these control statements to process the orject module.
2. Code PUNCH statements in macro definitions to produce, for example, source statements in other computer languages or for other processing phases.

The card that is punched has a physical position immediately after the PUNCH instruction and before any other TXT cards of the object decks that are to follow.

## Specifications

The PUNCH instruction causes the data in its operand to be punched into a card. One PUNCH instruction produces one punched card, but as many FUNCH instructions as necessary can ke used.

The PUNCH instruction statement can appear anywhere in a source module except before and ketween source macro definitions. If a PUNCH instruction occurs before the first control section, the resultant card punched will precede all other cards in the object deck.

The cards punched as a result of a PUNCH instruction are not a logical part of the object deck, even though they can be physically interspersed in the object deck.


PUNCH
The format of the PUNCH instruction statement is shown in the figure to the right.

All 256 punch combinations of the IBM System/370 character set are allowed in the character string of the operand field. Variable symbols are also allowed.
(1)

The position of each character specified in the PUNCH statement
2 corresponds to a column in the card to be punched. However, the following rules apply to ampersands and apostrophes:

1. A single ampersand initiates an attempt to identify a variable
(3) symbol and to substitute its current value.

4 2. Double ampersands or apostrophes are punched as single ampersands or apostrophes.
3. A single apostrophe followed by one or more blanks simply terminates the string of characters punched. If a non-blank character follows a single apcstrophe, an error message is issued and nothing is punched.

Only the characters punched, including blanks, count toward the maximum of 80 allowed.

NOTES:

1. No sequence number cr identification is punched into the card produced.
2. If the NCLECK option is specified in the EXEC statement of the job control language for the assembler program, no cards are punched: neither for the PUNCH or REPRO instructions, nor for the object deck of the assembly.
Examples:

| Source Statement | Value of Variable Symbol | Characters <br> Punched |
| :---: | :---: | :---: |
| PUNCH 'CHARS \&VAR' | ABC | CHARS ABC |
|  |  | CHARS \&' |
| PUNCH <br> 'CHARS |  | CHARS A |
| PUNCH 'CHARS A'REMARKS <br> * * * ERROR * * * |  |  |
| PUNCH 'CHARS A' REMARKS |  | CHARS A |

## Purpose

The REPRO instruction causes the data specified in the statement that follows to be punched into a card. Unlike the PUNCH instruction, the REPRO instruction does not allow values to be substituted into variable symbols tefore the card is punched.

## Specifications

The REPRC instruction causes data on the statement line that follows it to be punched into the corresponding columns of a card. One REPRO instruction produces one punched card.

The REPRO instruction can appear anywhere in a source module except before and between source macro definitions. The punched cards are not part of the object deck, even though they can be physically interspersed in the object deck.


The format of the REPRO instruction statement is shown in the figure to the right.

The line to be reproduced can contain any of the 256 punch characters, including blanks, ampersands, and apostrophes. No substitution is performed for variable symbcls.


NOTES:

1. No sequence numbers or identification is punched in the card.
2. If the NODECK option is specified in the job control language for the assembler program, no cards are punched: neither for the PUNCH or REPRO instructions, nor for the object deck of the assembly.

## Purpose

The OPSYN instruction allows you to define your own set of symbols to represent operation codes for:

1. Machine and extended mnemonic branch instructions.
2. Assembler instructions including conditional assembly instructions.

You can also prevent the assembler from recognizing a symbol that represents a current operation code.

OPSYN

## Specifications

The OPSYN instruction must be written after the ICTI instruction and can be preceded only by the EJECT, ISEQ. PRINT, SPACE, and TITLE instructions. The CFSYN instruction must precede any source macro definitions that may be specified.

The OPSYN instruction has two basic formats as shown in the figure to the right.

The operation code specified in
(1) the name field or the oferand field 2 must represent either:

1. The operation code of cne of the machine or assembler instructions as described in PARTS II, III, and PART IV of this manual, or
2. The operation code defined by a previous CFSYN instruction.
(3)

The OPSYN instruction assigns the properties of the operation code specified in the operand field to the symbol in the name field. A
(4) klank in the operand field causes the operation code in the name field to lose its properties as an operation code.

NOTE: The symbol in the name field can represent a valid creration code. It loses its current properties as if it had been defined in an OPSYN instruction with a blank operand field. Further, when the same symbol appears in the name field of two OPSYN instructions the latest definition takes
precedence.


```
H6 - Saving and Restoring Programming Environments
    OS
only
The instructions described in this subsection can save and restore the status of PRINT options and the base register assignment of your program.
```


## H6A -- THE FUSH INSTRUCTION

## Purpose

The PUSH instruction allows you to save the current PRINT or USING status in "push-down" storage on a last-in, firstout basis. You can restore this PRINT and USING status later, also on a last-in, first-out basis, by using a corresponding POP instruction.

## Specifications

The format of the PUSH instruction statement is shown in the figure to the right.

One of the four options for the operand entry must be specified. The PUSH instruction does not change the status of the current PRINT

| Name | Operation | Operand |  |
| :--- | :--- | :--- | :--- |
|  |  |  | Options |
| A sequence | PUSH | PRINT | 1 |
| symbol or |  | USING | 2 |
| blank |  | PRINT, USING 3 |  |
|  |  | USING, PRINT 4 |  | or USING instructions; the status is only saved.

NOTE: When the PUSH instruction is used in combination with the POP instruction, a maximum of four nests of PUSH PRINT - POP PRINT or PUSH USING - POP USING are allowed.

H6B -- THE FOF INSTRUCTION

## Purpose

The PCP instruction allows you to restore the PRINT or USING status saved by the most recent PUSH instruction.

## Specifications

The format of the POP instruction is given in the figure to the right.

One of the four options for the operand entry must be specified. The FCP instruction causes the status of the current PRINT or USING instruction to be overridden by the PRINT or USING status saved by the last PUSH instruction.

NOTE: When the POP instruction is used in combination with the PUSH instruction, a maximum of four nests of PUSH PRINT - POP PRINT or PUSH USING - POP USING are allowed.


# Part IV: The Macro Facility 

SECTION I: INTRODUCING MACROS
SECTION J: THE MACRO DEFINITION
SECTION K: THE MACRO INSTRUCTION
SECTION L: THE CONDITIONAL ASSEMBLY LANGUAGE

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## Section I: Introducing Macros

This section introduces the basic macro concept; what you can use the macro facility for, how you can prepare ycur own macro definitions, and how you call these macro definitions for processing by the assembler.

Read this section straight through before referring to the detailed descriptions identified by the cross-reference arrows.

NOTE: IBM supplies macro definitions in system libraries for input/output and other control program services, such as the dynamic allocation of main storage areas. To prccess these macro definitions ycu only have to write the macro instruction that calls the definition.

## Using Macros

FOR TEXT INSERTION: The main use of macros is to insert assembler language statements into a source program.


You call a named sequence of statements the macro
definition) by using a macro instruction, or macro call.
The assembler replaces the macro call ky the statements
from the macro definition and inserts them into the source
module at the point of call. The process of inserting
the text of the macro definition is called macro generation or macro expansion. The assembler expands a macro at preassembly time.

The expanded stream of code then Lecomes the input for processing at assembly time, that is, the time at which the assembler translates the machine instructions intc object code.

FOR TEXT MODIFICATION: You may want to modify the statements in a macro definition kefore they are generated.

1 You can do this by supplying character string values as operands in a macro call. These values replace parameters in the statement to be generated. This means that you can change the content of the generated statements each time you call the macro definition.

1 See K 3
2 See J3


FOR TEXT MANIFULATION: You can also select and reorder the statements to be generated from a macro definition by using the conditional assembly language described later in this section.

(1) The conditional assembly language allows you to manipulate text generation, for example, by branching upon the result of a condition test. You can choose exactly which

```
To use the complete macro facility provided by the assembler
you must:
```

- Prepare a macro definition and
- Call this definition using a macro instruction.


You can create a macro definition ky enclosing any seguence of assembler language statements between the MACRO and
(1) MEND statements, and by writing a prototype statement in which you give your definition a name. This name is then
(2) the operation code that you must use in the macro instruction to call the definition.

When you write a macro instruction in your source module, you tell the assembler to process a particular macro definition. The assembler produces assembler language statements from this macro definition for each macro instruction that calls the definition.


By using the macro facility you reduce programming effort, because:

1. You write and test the code a macro definition contains once. You and other programmers can then use the same code as often as you like by calling the definition; which means that you do not have to reconstruct the coding logic each time you use the code.
2. You need write only one macro instruction to call for the generation of many assembler language statements from the macro definition.

When you are designing and writing large assembler language programs, the above features allow you to:

- Prepare macro definitions, containing difficult code, for your less experienced colleagues. They can then call your definitions to generate the appropriate statements, without having to learn the code in the definition.
- Change the code in one place when updating or making corrections, that is, in the macro definition. Each call gets the latest version automatically, thus providing standard coding conventions and interfaces.
- Describe the functions of a complete macro definition rather than the function of each individual statement it contains, thus providing more comprehensible documentation for your source module.

Defining a Macro

Defining a macro means preparing the statements that constitute a macro definition. To define a macro you must:

1. Give it a name
2. Declare any parameters to be used
3. Write the statements it contains.
4. Establish its boundaries


1 The MACRO and MENL instructions establish the boundaries of a macro definition.

2
3
4
You use the prototype statement to name the macro and to declare its parameters. In the operand field of the macro instruction, you can assign values to the parameters declared for the called macro definition.

5 The body of a macro definition contains the statements that will be generated when you call the macro. These statements are called model statements; they are usually
 interspersed with conditional assembly statements or cther processing statements.

WHERE YOU CAN PLACE A MACRO DEFINITION: You can include a macro definition at the beginning of a source module. This type of definition is called a source macro definition.

You can also insert a macro definition in a system or user library (located, for example, on disk) by using the appropriate utility program. This type of definition is called a library macro definition. The IBM-supplied macro definitions mentioned earlier are examples of library macro definitions.


Calling a Macro

3 You can call a source macro definition only from the source module in which it is included. You can call a likrary macro definition from any source module.

WHERE YOU CAN CALL A MACRO EFFINITION: You can call a
1 macro definition by specifying a macro instruction anywhere
in a source module, except before or between any source macro definitions that may be specified.


You can also call a macro definition from within another macro definition. This type of call is an inner macro call; it is said to be nested in the macro definition.

2
$\xrightarrow{\text { See K6 A }}$

The body of a macro definition can contain a combination of model statements, processing statements, and comments statements.


MODEL STATENENTS: You can write assembler language statements as model statements. The assembler copies them exactly as they are written when it expands the macro. You can also use variable symbols as points of substitution in a model statement. The assembler will enter values in place of these points of substitution each time the macro is called.

The three types of variable symbol in the assembler language are

1. Symbolic parameters, declared in the prototype statement
2. System variable symbols (see J7)
3. SET symbols, which are part of the conditional assembly language (see IIA).

The assembler processes the generated statements, with or without value substitution, at assembly time.

PROCESSING STATEMENTS: Processing statements perform functions at pre-assembly time when macros are expanded, but they are not themselves generated for further processing at assembly time. The processing statements are:

1. Conditional assembly instructions
2. Inner macro calls
3. The MNOTE instruction
4. The MEXIT instruction.

(1) The MNOTE instruction allows you to generate an error
 message with an error condition code attached, or to generate comments in which you can display the results of pre-assembly computation.

2 The MEXIT instruction tells the assembler to stop processing
 a macro definition. The MEXIT instruction therefore provides an exit from the middle of a macro definition.
(3) The MEND instruction not cnly delimits the contents of a macro definition but also provides an exit from the definition.

COMMENTS STATEMENTS: One type of comments statement describes pre-assembly operations and is not generated. The other type describes assembly-time operations and is therefore generated (for details see J6).

## The Conditional Assembly Language

The conditional assembly language is a programming language with most of the features that characterize such a language. For example, it provides:

1. Variables
2. Data attributes
3. Expression computation
4. Assignment instructions
5. Labels for branching
6. Branching instructions
7. Sukstring operators that select characters from a string.

You can use the conditional assembly language in a macro definition to receive input from a calling macro instruction. You can produce output from the conditional assembly language by using the MNOTE instruction.

You can use the functions of the conditional assembly language to select statements for generation, to determine their order of generation, and to perform computations that affect the content of the generated statements.

The conditional assembly language is fully described in Section L.

## Section J: The Macro Definition

This section describes macro definitions: where they can ke placed in order to be available to call, how they are specified, and what they can contain.

Jl -- Using a Macro Definition

J1A -- PURPCSE

A macro definition is a named sequence of statements which you can call with a macro instruction. When it is called, the assembler processes and usually generates assembler language statements from the definition into the source module. The statements generated can ke:

1. Copied directly from the definition,
2. Modified by parameter values before generation, or
3. Manipulated by internal macro processing to change the sequence in which they are generated.

You can define your own macro definitions in which any combination of these three processes can occur. Some racro definitions do nct generate assembler language statements, but perform only internal processing, like some of the macro definitions used for system generation.

## Where to Define a Macro In a Source Module

A macro definition within a source module must be specified at the beginning of that source module. This type of macro definition is called a source macro definition. A macro definition can also reside in a system library; this type of macro is called a library macro definition. Either type can be called from the source module by. the appropriate macro instruction.

NOTE: A source macro definition can be entered into a library and thereby become a library macro definition. A library macro definition can be included at the beginning of a source module and thereby become a source macro definition.

Some control and comments statements can appear at the beginning of a source module along with the source macro definitions. They can be used:

## Eefore all macro definitions.

2. Between macro definitions.
(3) After macro definitions and before open code

All other statements of the assembler language must appear after any source macro definitions that are specified.

Open Code

Open code is that part of a source module that lies outside of and after any source macro definition. Open code is initiated by any statement of the assembler language that appears outside of a macro definition, except the ICTL, OPSYN, ISEQ, EJECT, FRINT, SPACE, or TITLE instruction, or a comments statement.

At coding time, it is important to distinguish between source statements that lie in open code and those that lie inside macro definitions.

NOTES:

1. The ISEQ, EJECT, PRINT, SPACE, and TITLE instructions, and one or more comments statements, can appear between source macro definitions and the start of open code. However, in this position, the above instructions must not contain any variable symbols.
2. After the start of open code, variable symbols are allowed in any statement.
3. A macro definition must not be specified after the start of open code.

The Format of a Macro Definition

The general format of a macro definition is shown in the figure to the right.

The four parts are described in detail below.


## J2 -- Parts of a Macro Definition

J2A -- THE MACRO DEFINITION HEADER

## Purpose

The macro definition header instruction indicates the beginning of a macro definition.

## Specifications

The MACRO instruction is the macro definition header; it must be the first statement of every macro definition. Its format is given in the figure to the right.

| Name | Operation | Operand |
| :--- | :---: | :---: |
| Not used, <br> must not be <br> present | MACRO | Not required |

J2B -- THE MACRO DEFINITION TRAILER

## Furpose

The macro definition trailer instruction indicates the end of a macro definition. It also provides an exit when it is processed during macro expansion.

## Specifications

The MEND instruction statement is the macro definition trailer; it must be the last statement of every macro definition. Its format is given in the figure to the right.

| Name | Operation | Operand |
| :--- | :--- | :--- |
| A sequence <br> symbol, or <br> not used | MEND | Not required |

## Purpose

The prototype statement in a macro definition serves as a model
(prototype) of the macro instruction you use to call the macro definition.

## Specifications

The prototype statement must be the second statement in every macro definition. It comes immediately after the MACRO instruction.

The format of the prototype statement statement is given in the figure to the right.

The maximum number of symbolic parameters allowed in the operand field is not fixed. It depends on the amount of virtual storage available to the program.

|  |  | Prototype |
| :--- | :--- | :--- |
| Name | Operation | Operand |
| A name <br> field <br> parameter <br> or blank | A symbol | Zero or more <br> symbolic <br> parameters <br> separated by <br> commas |

dos only 200 parameters are allowed in the operand field.

If no parameters are specified in the operand field, remarks are allowed, if the absence of the operand entry is indicated by a comma preceded and fcllowed by one or more blanks.

## Alternate Ways of Coding the Prototype Statement

The prototype statement can be specified in one of the following three ways:
1 The normal way, with all the symbolic parameters preceding any remarks.

2 An alternate way, allowing remarks for each parameter.
(3) A combination of the first two ways. NOTES:

1. Any number of continuation lines is allowed. However, each
(4) continuation line must be indicated by a nonblank character in the column after the end column on the preceding card.
2. For each continuation line, the operand field entries (symbolic parameters) must begin in the continue column otherwise the whole line and any lines that follow will be considered to contain remarks.


## Purpose

You can write a name-field parameter similar to the symbclic parameter, as the name entry of a macro prototype statement. You can then assign a value to this parameter from the name entry in the calling macro instruction.

## Specifications

If used, the name entry must be a variable symbol. If this parameter
(2) also appears in the body of a macro, 3 it will be given the value assigned to the parameter in the name field of the corresponding macrc
instruction. Note that the value assigned to the name field parameter has special restrictions that are listed in K2A.

The Operation Entry

Purpose

The operation entry is a symbol
that identifies the macro definition. When you specify it in the operation field of a macro instruction, the appropriate macro definition is called and processed by the assembler.

## Specifications

4 The symbol in the operation field of the prototype statement establishes the name by which a macro definition must be called.


This name becomes the operation code required in any macro instruction that calls the macro.
os NOTE: Unless operation codes have
only been changed by the OPSYN
instruction, the operation code specified in the prototype statement must not be the same as that specified in:

1. A machine instruction.
2. An assembler instruction.
3. The prototype statement of another source (or library) macro definition.

Purpose

The operand entry in a prototype statement allows you to specify positional or keyword parameters. These parameters represent the values you can pass from the calling macro instruction to the statements within the body of a macro definition.

## Specifications

The operands of the macro prototype statement must be symbolic parameters separated by commas. They can be positional parameters or keyword parameters or both (see J3)

NOTE: The operands must be symbolic parameters; parameters in sublists are not allowed. For a discussion of sublists in macro instruction operands, see K 4 .

## Purpose

The body of a macro definition contains the sequence of statements that constitutes the working part of a macro. You can specify:

1. Model statements to be generated.
2. Processing statements that, for example, can alter the content and sequence of the statements generated or issue error messages.
3. Comments statements, some of which are generated and others which are not.
4. Conditional assembly instructions to compute results tc be displayed in the message created by the MNOTE instruction; without causing any assembler language statements to ke generated.

## Specifications

The statements in the body of a macro definition must appear between the macro prototype statement and the MEND statement of the definition. The three main types of statements allowed in the body of a macro are:

1. Model statements (see J4),
(2) Processing statements (see J5). and
(3) Comments statements (see J6).

NOTE: The body of a macro definition can be empty, that is, contain no statements.

## J3 -. Symbolic Parameters

## Purpose



Symbolic parameters allow you to pass values into the body of a macro definition from the calling macro instruction. You declare these parameters in the macro prototype statement. They can serve as goints mabstitution in the body of the 4 macro definition and are replaced the calling macro instruction.

By using symbolic parameters with meaningful names you can indicate the purpose for which the parameters (or substituted values) are used.

## General Specifications

Symbolic parameters must be valid variable symbols, as shown in the figure to the right.They have a local scope: that is, the value they are assigned only applies to the macro definition in which they have been declared. The value of the parameter remains constant throughout the processing of the containing macro definition for every call on that definition.

NOTE: Symbolic parameters must not be multiply defined or identical to any other variable symbols within

$\qquad$
MEND

The two kinds of symbolic parameters are:

- Positional parameters
- Keyword parameters.

Each positional or keyword parameter 3 used in the body of a macro
 definition must be declared in the prototype statement.

## Subscripted Symbolic Parameters

Subscripted symbolic parameters must be coded in the format shown in the figure to the right.
(1.)

The subscript can be any arithmetic expression allowed in the operand field of a SETA instruction (arithmetic expressions are discussed in L4A). The arithmetic expression can contain subscripted variable symbols. Subscripts can be nested up to 5 levels of nesting.

The value of the subscript must be greater than or equal to one. The subscript indicates the position of the entry in the sublist that is specified as the value of the subscripted parameter (sublists as values in macro instruction operands are fully described in K4).


## Purpose

You should use a positional parameter
in a macro definition if you wish to change the value of the parameter each time you call the macro definition. This is because it is easier to supply the value for a positional parameter than for a keyword parameter. You only have to write the value you wish the parameter to have in the proper position in the operand of the calling macro instruction.

For keyword (described below) parameters, you must write the entire keyword and the equal sign that precedes the value to be passed. However, if you need a large number of parameters, you should use keyword parameters. The keywords make it easier to keep track of the individual values you must specify at each call, by reminding you which parameters are being given values.

## Specifications

The general specifications for symbolic parameters described in J3 also apply to positional parameters. Note that the specification for each positional parameter declared in the prototype statement definition must be a valid variable symbol. Values are assigned to the positional parameters by
2 the corresponding positional operands specified in the macro instruction that calls the definition.

## Purpose

You should use a keyword parameter in a macro definition for a value that changes infrequently. Ey specifying a standard default value to be assigned to the keyword parameter, you can omit the corresponding keyword operand in the calling macro instruction.

Keyword parameters are also convenient kecause:

1. You can specify the corresponding keyword operands in any order in the calling macro instruction.
2. The keyword, repeated in the operand, reminds you which parameter is being given a value and for which purpose the parameters is being used.

## Specifications

The general specifications for symbolic parameters described in J3 also apply to keyword parameters. Each keyword parameter must be in the format shown in the figure to the right.
1.

The actual parameter must be a valid variable symbol.

A value is assigned to a keyword
 parameter by the corresponding keyword operand through the name of the keyword as follows:
3) 1. If the corresponding keyword operand is omitted, the standard 4 value specified in the prototype statement becomes the value of the parameter for that call (for full details on values passed see K5).

5
2. If the corresponding keyword operand is specified, the value after the equal sign overrides the standard value in the prototype and becomes the value of the parameter for that call (see k5).

Key. Param.

(1) NOTE: A null character string can be specified as the standard value of a keyword parameter, and will be generated if the corresponding keyword operand is omitted.


## Purpose

By using positional and keyword parameters in a prototype statement, you combine the benefits of both. You can use positional parameters in a macro definition for passing values that change frequently and keyword parameters for passing values that do not change often.

## Specifications

Fositional and keyword parameters can be mixed freely in the macro prototype statement. The same applies to the positional and keyword operands of the macro instruction (see K3C) : Note, however, that
(3) the order in which the positional parameters appear determines the order in which the positional operands must appear. Interspersed
(4) keyword parameters or operands do not affect this order.

All positional parameters must precede any keyword parameters, if specified. The same applies to positional and keyword operands of a macro instruction (see K3C).


## J4 - Model Statements

J4A -- PURPCSE

> Model statements are statements from which assembler language statements are generated at pre-assembly time. They allow you to determine the form of the statements to be generated. Ey specifying variable symbols as pcints of substitution in a model statement, you can vary the content of the statements generated from that model statement. You can also use model statements into which you substitute values in open code.

## J4B - SPECIFICATIONS

The following specifications also apply to model staterents in open code. Exceptions are noted where applicable.

Model Stmnt

Format of Model Statements

A model statement consists of one or more fields separated by one or more blanks.

Each field or subfield can consist of:
(1) An ordinary character string
(2) A variable symbol as a point of substitution
(3) Any combination of ordinary character strings and variable symbols to form a concatenated string.

The statements generated at preassembly time from model statements must be valid machine or assembler instructions, but must not be conditional assembly instructions. They must obey the coding rules described in Section $E$ or they will be flagged as an error at assembly time.


## Variable Symbols as Points of Substitution

Values can be substituted for variable symbols that appear in the name, operation, and operand fields of model statements; thus, variable symbols represent points of substitution. The three main types of variable symbol are:
(1) Symbolic parameters (described in J3 above) ,
(2) System variable symbols (described in J7 below) , and
(3) SET symbols (described in L1A). NOTES:

1. Symbolic parameters, SET symbols, and the system variable symbol, ESYSIIST, can all be subscripted. The remaining system variable symbols ESYSNLX, ESYSECT, ESYSPARM, ESYSDATE, and ESYSTIME cannot be subscripted.

2. The fields in a statement generated from a model statement 1) appear in the listings in the same columns as in the model statement.
(2) However, when values are substituted for variable symbols the generated
0 fields can be displaced to the right.

When variable symbols are concatenated to ordinary character strings, the following rules apply to the use of the concatenation character (a period):

The concatenation character is mandatory when:
(1) An alphameric character is to follow a variable symbol.
2 - A left parenthesis that does not enclose a subscript is to follow a variable symbol.
(3) A period (.) is to be generated.

4 Two periods must be specified in the concatenated string following a variable symbol.

The concatenation character is not necessary when:
(5) - An ordinary character string precedes a variable symbol.
6) A special character, except left parenthesis or period, is to follow a variable symbol.
(1) - A variable symbol follows another variable symbol.

The concatenation character must
(8) not be used between a variable symbol and its subscript; otherwise, the characters will be considered a concatenated string and not a subscripted variable symbol.

## Rules for Model Statement Fields

The fields that can be specified in model statements are the same fields that can be specified in an ordinary assembler language statement. They are the name, operation, operand and remarks fields. It is also possible to specify a continuation - indicator field, an identification - sequence field, and a field before the begin column, if the appropriate ICTL instruction has been specified. Character strings in the last three fields (in the standard format only columns 72 through 80) are generated exactly as they appear in the model statement, and no values are substituted for variable symbols.

Model statements must have an entry in the operation field, and, in most cases, an entry in the operand field in order to generate valid assembler language instructions.

THE NAME FIELL: The entries allowed in the name field of a model statement are given in the figure to the right, including the allowable results of generation.

Variable symbols must not be used to generate comments statement indicators.

NOTE: Restrictions on the name entry are further specified where each individual assembler language instruction is described in this manual.

| Name <br> Field | Allowed | Not Allowed |
| :---: | :---: | :---: |
| In Model Statements <br> (before generation) | blank <br> - ordinary symbol <br> - sequence symbol <br> - variable symbol <br> - any combination of variable symbols and other character strings concatenated together |  |
| In Generated Statements (generated results) | - blank <br> - valid ordinary symbol | $\stackrel{*}{\cdot *}\} 1$ |

THE OPERATION FIELE: The entries allowed in the operation field of a model statement are given in the figure to the right, including the allowable results of generation.

1 The operation codes ICTL and OPSYN are not allowed inside a macro
2 definition. The MACRO and MEND operation codes are not allowed in model statements; they are used only for delimiting macro definitions.

DOS The END operation code is not allowed 3 Inside a macro definition.

4 If the REPRO operation code is specified in a model statement, no substitution is performed for
5 the variable symbols in the statement line following the REPRO statement. Variable symbols can be used alone or as part of a concatenated string to generate operation codes for:
(6) - Any machine instruction, or
(1) The assembler instructions listed.

8 NOTE: The MNOTE and MEXIT statements are not model statements; they are described in J5D and J5E respectively.

The generated operation code must not be an operation code for the following (or their OPSYN equivalents) :
(9) - A macro instruction,

10 - A conditional assembly instruction, or
(11) The assembler instructions listed. DOS - The END operation code must not be generated.

| Operation <br> Field | Allowed | Not Allowed |
| :---: | :---: | :---: |
| In Model Statements (Before Generation) | An ordinary symbol that represents the operation code for: <br> - any machine instruction <br> - a macro instruction <br> - the following Assembler instructions: | blank <br> The assembler operation codes: <br> (1) $\left\{\begin{array}{l}\text { ICTL }\end{array}\right.$ lopsyn <br> 2 macro MEND <br> 3 END DOS |
| In Generated <br> Statements <br> (Generated <br> Results) | An ordinary symbol that represents the operation code for: - any machine instruction - the following assembler instructions: $\left(\begin{array}{lll}\text { CCW } & \text { EJECT } & \text { SPACE } \\ \text { CNOP } & \text { END } & \text { TITLE } \\ \text { COM } & \text { ENTRY } & \text { USING } \\ \text { CSECT } & \text { EQU } & \text { WXTRN } \\ \text { CXD } & \text { EXTRN } & \\ \text { DC } & \text { LTORG } & \text { (MNOTE) } \\ \text { DROP } & \text { ORG } & 8 \\ \text { DS } & \text { POP } & \\ \text { DSECT } & \text { PRINT } & \\ \text { DXD } & \text { PUNCH } & \text { PUSH } \\ & & \\ \text { OS } & \end{array}\right.$ only | blank a macro instruction operation code a conditional assembly operation code: ACTR |

THE OFERAND FIFLD: The entries allowed in the operand field of a model statement are given in the figure to the right, including the allowable results of generation.

NOTE: Variable symbcls must not
be used in the operand field of a COPY, ICTL, ISEQ, or OPSYN instruction.

| Operand Field | Allowed | Not Allowed |
| :---: | :---: | :---: |
| In Model <br> Statements <br> (Before Generation) | Blank (if valid) <br> - An ordinary_symbol <br> - A character string. combining alphameric and special characters (but not variable symbols) <br> A variable symbol <br> - A combination of variable symbols and other character strings concatenated together |  |
| In Generated Statements <br> (Generated Results) | blank (if valid) <br> Character String. that represents a valid assembler or machine instruction operand field | operand field of a: COPY, ICTL, ISEQ or OPSYN statement <br> 1 |

THE REMARKS FIELD: Any combination of characters can be specified in the remarks field of a model
1 statement. No values are substituted into variable symbols in this field.

NOTE: One or more blanks must be used in a model statement to separate the name, operation, crerand, and remarks fields from each other. Blanks cannot be generated between fields in order to create a complete assembler language statement.

OS The exception to this rule is that only a combined operand-remarks field 3 can be generated with one or more blanks to separate the two fields.

## J5 -- Processing Statements

J5A - - CONDITIONAL ASSEMBLY INSTRUCTIONS

Conditional assembly instructions allow you to determine at preassembly time the content of the generated statements and the sequence in which they are generated. The instructions and their functions are given in the figure to the
right.

Conditional assembly instructions can ke used both inside macro definitions and in open code. They are fully described in Section $L$.

| Conditional Assembly <br> Instruction | Function Performed |
| :--- | :--- |
| GBLA, GBLB , GBLC <br> LCLA, LCLB, LCLC | Declaration of initial values <br> of variable symbols (global <br> and local SET symbols) <br> SETA, SETB , SETC <br> AIFAssignment of values to <br> variable symbols (SET <br> symbols) <br> AGO <br> ANOPBranching <br> - Conditional (based on <br> logical test) |
| - Unconditional |  |
| - To next Sequential |  |
| instruction (No |  |
| operation) |  |

## J5B -- INNER MACRO INSTRUCTIONS

Macro instructions can be nested inside macro definitions, allowing you to call other macros from within your own definitions. Nesting of macro instructions is fully described in K 6 .

## J5C -- THE COPY INSTRUCTION

## Purpose

The COPY instruction, inside macro definitions, allows you to copy into the macro definition any sequence of statements allowed in the body of a macro definition. These statements become part of the body of the macro before macro processing takes place. You can also use the CCPY instruction to copy complete macro definitions into the beginning of a scurce module.

The specifications for the COPY instruction, which can also be used in open code, are described in E1A.

## Purpose

You can use the MNOTE instruction to generate your own error messages or display intermediate values of variable symbols computed at preassembly time.

## Specifications

The MNOTE instruction can be used OS inside macro definitions or in open only code, and its operation code can be created by substitution. The MNOTE instruction causes the generation of a message which is given a statement number in the printed listing.

The format of the MNOTE instruction statement is given in the figure to the right.

1 The $n$ stands for a severity code. The rules for specifying the contents of the severity code subfield are as follows:

1. The severity code can be specified

2 as any arithmetic expression allowed
in the operand field of a SETA
instruction. The expression must have a value in the range 0 through 255.
(3) 2. If the severity code is omitted, kut the comma separating it from the message is present, the assembler assigns a default value of 1 as the severity code.
(4) 3. An asterisk in the severity code subfield causes the message and the asterisk to be generated as a comments statement.

5 4. If the entire severity code subfield is omitted, including the comma separating it from the message, the assembler generates the message as a comments statement.

| Name | Operation |  | Operand |
| :---: | :---: | :---: | :---: |
| A sequence symbol or blank | MNOTE |  | One of four options allowed: <br> n, 'message' error <br> $\left.\begin{array}{rl}\text {,'message' } & \text { message } \\ \text { \%,'message' } \\ \text { 'message }\end{array}\right\} \begin{aligned} & \text { comments }\end{aligned}$ |
| Examples: |  |  |  |
| Source Statements |  |  | Generated Result |
| 2 <br> MNOTE 2,'ERROR IN SYNTAX' <br> 3 <br> MNOTE ,'ERROR, SEV 1' <br> 4 <br> MNOTE *,'NO ERROR' <br> 5 |  |  | 2,ERROR IN SYNTAX <br> ,ERROR, SEV 1 <br> *,NO ERROR <br> NO ERROR |

NOTES:

1. An MNOTE instruction causes a message to be printed, if the current PRINT option is ON, even if the PRINT NOGEN option is specified.
2. The statement number of the message generated from an MNOTE instruction with a severity code is listed among any other error messages for the current source module. However, the message is printed only if the severity code specified is greater than or equal to the severity code "nnn" in the assembler option, FLAG (nnn), contained in the EXEC statement that invokes the assembler.

DOS The assembler option FLAG does not exist, and the severity code is not used by the DOS control program.
3. The statement number of the comments generated from an MNOTE instruction without a severity code is not listed among other error messages.

Any combination of up to 256 characters enclosed in apostrophes can be specified in the message subfield. The rules that apply to this character string are as follows:
(1) Variable symbols are allowed (NOTE: variable symbols can have a value that includes even the enclosing apostrophes).
(2) Double ampersands and double apostrophes are needed to generate one ampersand or one apostrophe. If variable symbols have ampersands or apostroches as values, the values must have double ampersands or apostrophes.

NOTE:
Any remarks for the MNOTE instruction statement must be separated from the apostrophe that ends the message by one or more blanks.

|  | Value of Variable Symbol | Generated Result |
| :---: | :---: | :---: |
| 3,'THIS IS A MESSAGE' <br> 3, \&PARAM | $\text { \&PARAM }=\text { ERROR' }$ $\& A=10$ | 3,THIS IS A MESSAGE $3 \text {,ERROR }$ <br> 3,VALUE OF \&A IS 10 |
|  |  | $\begin{aligned} & \text { 3,L'FIELD1 } \\ & \text { 3,DOUBLE \& } \\ & \text { 3,DOUBLE L'FIELD1 } \end{aligned}$ |
|  |  | 3,MESSAGE STOP RMRKS |

## Purpose

The MEXIT instruction allows you to provide an exit for the assembler from any point in the body of a macro definition. The MEND instruction provides an exit only from the end of a macro definition (see J2B) .

## Specifications

The MEXIT instruction statement can be used only inside macro definitions. It has the format given in the figure to the right.

The MEXIT instruction causes the assembler to exit from a macro

1) definition to the next sequential instruction after the macro instruction that calls the definition. (This also applies to nested macro instructions, which are described in K6.)

| Name | Operation | Operand |
| :---: | :---: | :---: |
| Sequence symbol or blank | MEXIT | Not requir |
|  |  |  |
| MACRO |  |  |
|  | ITS <br> A <br> B <br> C <br> XIT <br> $\cdot$ $\dot{D}$ $E$ $F$ |  |
|  | ND <br> ART 0 <br> ITS |  |

## J6 -- Comments Statements

J6A - - INTERNAL MACRO COMMENTS STATEMENTS

## Purpose

You write internal macro comments in the kody of a macro definition, to describe the operations performed at preassembly time when the macro is processed.
Specifications
Internal macro comments statements can be used only inside macro definitions. An example of their correct use is given in the figure to the right.
No values are substituted for any variable symbols that are specified in macro comments statements.


J6B -- ORDINARY COMMENTS STATEMENTS

## Purpose

Ordinary comments statements (described in E1C) allow you to make descriptive remarks about the generated output from a macro definition.

## Specifications

Ordinary comments statements can be used in macro definitions and in open code. An example of their correct use is shown in the figure to the right.

Even though this type of statement is generated along with the model
 statements of a macro definition, values are not substituted for any variable symbols specified.

## J7 -- System Variable Symbols

## Purpose

System variable symbols are variable symbols whose values are set by the assembler according to specific rules. You can use these symbcls as points of substitution in model statements and conditional assemkly instructions.

## General Specifications for System Variable Symbols

The system variable symbols: छSYSDATE, ESYSPARM, and GSYSTIME, can be used as points of substitution both inside macro definitions and in open code. The remaining system variable symbols: ESYSECT, ESYSLIST, and ESYSNDX, can ke used only inside macro definitions. All system variable symbols are subject to the same rules of concatenation and substitution as other variable symbols (see J4B).

System variable symbols must not be used as symbolic parameters in the macro prototype statement. Also, they must not be declared as SET symbols (see L2).

The assembler assigns read-only values to system variable symbols; they cannot be changed by using the SETA, SETB, or SETC instructions (see L3).

THE SCOPE OF SYSTEM VARIAELE SYMEOLS: The system variable symbols: ESYSDATE, ESYSPARM, and ESYSTIME, have a global sccpe. This means that they are assigned a read-only value for an entire source module; a value that is the same throughout open code and inside any macro
2 definitions called. The system variable symbols: छSYSECT, छSYSLIST, and ESYSNDX, have a local scope. They are assigned a read-only value each time a macro is called, and have that value only within the expansion of the called macro.

## Purpose

\&SYSDATE
You can use ESYSDATE tc obtain the date on which your source module is assembled.

## Specifications

The global system variable symbol gSYSDATE is assigned a read-only value of the format given in the figure to the right.

NOTE: The value of the type attribute of ESYSDATE (T'ESYSDATE) is always U and the value of the count attrikute (K'ESYSDATE) is always eight. (Attributes are fully described in L1B.)

## Purpose

You can use ESYSECT in a macro definition to generate the name of the current contrcl section. The current control section is the control section in which the macro instruction that calls the definition appears.

## \&SYSECT

## Specifications

The local system variable symbol ESYSECT is assigned a read-only value each time a macro definition is called.

1 The value assiqned is the symbol that represents the name of the current control section from which the macro definition is called. Note that it is the contrcl section in effect when the macro is called. A control section that has been
substitution does nct affect the value of ESYSECT for the expansion of the current macro. However, it does affect ESYSECT for a
(4) subsequent macro call. Nested
(5) macros cause the assembler to assign 5 a value to ESYSECT that depends on the control section in force inside the outer macro when the inner macro is called (see K6).

NOTES:

1. The control section whose name
is assigned to ESYSECT can be defined by a START, CSECT, LSECT, or COM instruction.
2. The value of the type attrikute
of ESYSECT, T'ESYSECT, is always
U , and the value of the count attribute (K'ESYSECT) is equal to the number of characters assigned as a value to ESYSECT. (Attributes are fully described in L1E.)


## Purpose

You can use ESYSLIST instead of a positional parameter inside a macro definition, for example, as a point of substitution. Ey varying the subscripts attached to ESYSLIST, you can refer to any positional operand or sublist entry in a macro call. ESYSLIST allows you to refer to positional operands for which no corresponding positional parameter is specified in the macro prototype statement.

## Specifications

The local system variable symbol ESYSLIST is assigned a read-only value each time a macro definition is called.

ESYSLIST refers to the complete list of positional operands specified in a macro instruction. ESYSLIST does not refer to keyword operands.

However, ESYSLIST cannot be specified as ESYSLIST alone. One of the two forms given in the figure to the right must be used as a point of substitution:
(1.) To refer to a positional operand
(2) 2. To refer to a sublist entry of a positional operand (sublists are fully described in K 4 below).
(3) The subscript $n$ indicates the position of the operand referred
(4) to The sukscript m, if specified, indicates the position of an entry in the suklist specified in the operand whose position is indicated by the first subscrift $n$.


The subscripts $n$ and $m$ can be any arithmetic expression allowed in the operand of a SETA instruction (see L3A) . The subscript $n$ must be greater than or equal to 0 . The subscript m must be greater than or equal to 1.

The figure to the right shows examples of the values assigned to ESYSLIST according to the value of its subscript, $m$ and $n$.

If the position indicated by $n$ (1) refers to an omitted operand or refers past the end of the list
2 of positional operands specified, the null character string is substituted for ESYSLIST (n) . If the position (in a sublist) indicated by the second subscript, $m$, refers to an omitted entry or refers past the end of the list of entries specified in the sublist referred to by the first subscript, n, the null character string is substituted for ESYSLIST $(n, m)$. Further, if the nth positional operand is not a sublist, $\in S Y S L I S T(n, 1)$ refers to the operand but ESYSLIST ( $n, m$ ), where m is greater than 1, will cause the null character string to be substituted.
(6)

NOTE: If the value of subscript n is zero, then ESYSLIST (n) is assigned the value specified in the name field of the macro instruction, except when it is a sequence symbol.


Attribute references can be made to the previously described forms of ESYSLIST. The attributes will be the attributes inherent in the positional operands or sublist entries to which you refer.
(Attributes are fully described in L1B.) However, the number attribute of ESYSLIST, N'ESYSLIST, is different from the number attribute described in L1B. One of the two forms given in the figure to the right can be used for the number attribute:
(1)

- To indicate the number of positional operands specified in a call

2. To indicate the number of sublist entries that have been specified in a positional operand indicated
(3) by the subscript.

NOTES:

1. For N'ESYSLIST, positicnal
(4) operands are counted if specifically omitted by specifying the comma that would normally have followed the operand.
(5)
2. For $\mathrm{N}^{\prime} \varepsilon S Y S L I S T(n)$, sublist entries are counted if specifically omitted by specifying the comma that would normally have followed the entry.
3. If the operand indicated by $n$ is not a sublist, N'ESYSLIST ( $n$ )
is 1. If it is omitted,
N'ESYSLIST ( n ) is zero.

| (1) $\mathrm{N}^{\prime} \&$ SYSLIST |  |  |
| :---: | :---: | :---: |
| Macro Instruction |  | Value of N'\&SYSLIST |
| MACLST | 1,2,3,4 | 4 |
| MACLST | $A, B, D, E$ | 5 ) 4 |
| MACLST | , A, B, C, D |  |
| MACLST | $(A, B, C),(D, E, F)$ | 2 Counts sublists |
| MACLST MACLST MACLST |  | $0 \sim$ |
|  | $\mathrm{KEY} 1=\mathrm{A}, \mathrm{KEY} 2=\mathrm{B}$ A | ${ }_{0}^{0}$ Keyword operands |
|  | A, B, |  |
| $2 \mathrm{~N}^{\prime} \& \operatorname{SYSLIST}^{(n)}$ |  |  |
| Macro Instruction |  | Value of $N^{\prime} \& S Y S L I S T(2)$ |
| $(\mathrm{n}=2)$ |  |  |
| MACSUB $\quad$, $(1,2,3,4,5), \mathrm{B}$ |  |  |
| MACSUB | A, $(1,1,3,5), B$ |  |
| MACSUB | A, $(, 2,3,4,5), B$ |  |
| MACSUB | A, B, C |  |
| MACSUB | A), C | 0 |
| MACSUB | $A, \operatorname{KEY}=(\mathrm{A}, \mathrm{B}, \mathrm{C})$ | 0 |
|  |  | $\begin{aligned} & \begin{array}{l} \text { Keyword sublists } \\ \text { are not counted } \end{array} \\ & \hline \end{aligned}$ |

## J7D -- ESYSNDX

## Purpose

You can attach ESYSNDX to the end of a symbol inside a macro definition to generate a unique suffix for that symbol each time you call the definition. Although the same symbol is generated by two or more calls to the same definition, the suffix provided by ESYSNDX produces two or more unique symbols. Thus you avoid an error being flagged for multiply defined symbcls.

## Specifications

The local system variable symbol ESYSNDX is assigned a read-only value each time a macro definition is called from a source module.

The value assiqned to ESYSNDX is a 4 -digit number, starting at 0001 for the first macro called by a
program. It is incremented by one for each subsequent macro call (including nested macro calls, described in K6).

NOTES:

1. ESYSNDX does not generate a valid symbol, and it must:
(3) a. Follow the symbol to which it is concatenated
b. Be concatenated to a symbol containing four characters or less.
2. The value of the type attribute of ESYSNDX (T'ESYSNDX) is always N , and the value of the count attribute ( $\mathrm{K}^{\prime}$ ESYSNDX) is always four.
(Attributes are fully described in L 1 B .)

## J7E -- ESYSPARM

## Furpose



You can use ESYSPARM to communicate with an assembler source module through the job control language. Through ESYSPARM, you pass a character string into the source module to be assembled from a job control language statement or from a program that dynamically invokes the assembler. Thus, you can set a character value from outside a source module and then examine it as part of the source module at pre-assembly time, during conditional assembly processing.

The global system variable symbol
1 ESYSFARN is assigned a read-only value in a job contrcl statement or in a field set up by a program that dynamically invokes the assembler. It is treated as a global SETC symbol in a scurce module, except that its value cannot be changed.

The largest value that ESYSPARM can hold when you code your own procedure is 91 characters, which can be specified by an invoking program. However, if the PARM field of the EXEC statement is used to specify its value, the PARM field restrictions reduce its maximum possible length.

CMS Under CMS, the option line of the ASSEMBLE command cannot exceed 100 characters, thus limiting the number of characters you can specify for ESYSPARM.

Dos The largest value GSYSPARM can hold is 8 characters.

NOTES:

1. No values are substituted for variable symbols in the specified value; however, double ampersands must be used to represent single ampersands in the value.

CMS Since CMS does not strip ampersands from the variable symbol, you need not specify double ampersands for CMS.
2. Double apostrophes are needed to represent single apostrophes because the entire PARM field specification is enclosed in apostrophes.

CMS Since CMS does not strip single apostrophes from the variable symbol, you need not specify double apostrophes for CMS.

3. If SYSPARM is not scecified in a job control statement outside the source module, ESYSPARM is assigned a default value of the null character string.
4. The value of the type attribute
of ESYSPARM (T'ESYSPARM) is always
U. while the value of the count
attribute ( $\mathrm{K}^{\prime}$ ESYSPARM) is the number of characters specified for SYSPARM in a job control statement or in a field set up by a program that dynamically invokes the assembler. Double apostrophes and double ampersands count as one character.

CMS 5. CMS parses the command line, breaking the input into eightcharacter tokens; therefore, the SYSPARM option field under VM/370 is limited to an eight-character field. If you want to enter larger fields or if you want to enter parentheses or embedded blanks, you must enter the special symbol "?" (the question mark symbol) in the option field. When CMS encounters this symbol in the command line, it will prompt you with the message ENTER SYSPARM:, after which you may enter any characters you want up to the option line limit of 100 characters. The following code is an example of how to use the ? symbol in the SYSPARM field:
assemble test (load deck sysparm (?) ENTER SYSPARM: 6 Eam, 'bo) . fy

R;

Purpose

You can use ESYSTIME to obtain the time at which your source module is assembled.

## Specifications

The global system variable symbol eSYSTIME is assigned a read-only value of the format given in the figure to the right.

NOTES:

1. The value of the type attribute of ESYSTIME (T'ESYSTIME) is always $U$ and the value of the count attribute (K'ESYSTIME) is always 5.
2. For systems without the internal timer feature, ESYSTIME is a 5character string of blanks.

Format: 5-Character String
$L$
hh.mm
Where: hh gives the hours mm gives the minutes

Example: 22.1


Corresponds to the time printed in the page heading of listings, remains constant for each assembly

In addition to the PRINT options that you can set from inside a source module, you can set other listing options from outside a source module by using the jot control language. These options can be specified in the PARM field of the EXEC statement or by a program that dynamically invokes the assembler.

J8A -- LIBMAC

## Purpose

The LIBMAC option allows you to print in the program listings the library macro definitions called from your source module, and any statements in open code following the first END statement (coded or generated) that is processed by the assembler.

## Specifications

The LIBMAC option, when set, causes:
1 Any statements in open code that follow the first END statement and
(2) All library macro definitions called to be printed in the program listings after the first (or cnly) END statement of the source module.

3 NOTE: Multiple END statements can be coded or generated and are printed, but the first ENL statement processed ends the assembly.

The option NOLIEMAC suppresses the listing of the items mentioned above. It is the default option that applies to the assembling of source modules.

J8B -- MCALL

## Purpose

The MCALL option allows ycu to list all the inner macro instructions that the assembler processes.

## Specifications

The MCALL option, when set, causes all inner macro instructions processed by the assembler to be listed. The NOMCALL option suppresses the listing of inner macro instructions. It is the default option that applies to the assembling of source modules.

NOTE: The MLOGIC and ALOGIC options concern the listing of conditional assembly statements. They are discussed in L8.

## Section K: The Macro Instruction

```
This section describes macro instructions: where they
can be used and how they are specified, including details
on the name, operation, and operand entries, and what will
be generated as a result of that macro call.
After studying this section, you should be able to use
the macro instructions correctly to call the macro
definitions that you write. You will also have a better
understanding of what to specify when you call a macro
and what will be generated as a result of that call.
```


## K1 -- Using a Macro Instruction

K1A - PURPCSE

The macro instruction provides the assembler with:

1. The name of the macro definition to ke processed.
2. The information or values to be passed to the macro definition. This information is the input to a macro definition. The assembler uses the information either in processing the macro definition or for substituting values into a model statement in the definition.

The output from a macro definition, called by a macro instruction, can be:

1. A sequence of statements generated from the model statements of the macro for further processing at assembly time.
2. Values assigned to global SET symbols. These values can be used in other macro definitions and in open code (see L1A).


A macro instruction can be written anywhere in the open code portion of a source module. However, the statements generated from the called macro definition must be valid assembler language instructions and allowed where the calling macro 2 instruction appears. A macro instruction is not allowed before or between any source macro definitions, if specified, but it can be nested inside a macro
4 definition (see K6).

## Macro Instruction Format

The format of a macro instruction statement is given in the figure to the right.

The maximum number of operands allowed is not fixed. It depends on the amount of virtual storage available to the program.

DOS Only 200 operands are allowed in the operand field.

If no operands are specified in the operand field, remarks are allowed if the absence of the operand entry is indicated by a comma preceded and followed by one or more blanks.

The entries in the name, operation, and operand fields correspond to entries in the prototype statement of the called macro definition (see K2) .


Macro Inst.

## Alternate Ways of Coding a Macro Instruction

A macro instruction can be specified in one of the three following ways:
(1) The normal way, with the cperands preceding any remarks.

2 The alternate way, allowing remarks for each operand.
(3) A combination of the first two ways.

NOTES:

1. Any number of continuation lines 4 are allowed. However, each continuation line must ke indicated ky a non-klank character in the column after the end column of the previcus staterment line (see B1B).
2. Operands on continuation lines must begin in the continue column, or
3. Otherwise, the assembler assumes that any lines that follow contain remarks.

NOTE: If any entries are made in the columns before the continue column in continuation lines, the assembler issues an error message and the whole statement is not processed.


## K2 -- Entries

## Purpose

You can use the name entry of a macro instruction:

1. Either to generate an assemblytime label for a machine or assembler instruction.
2. Or to provide a conditional assembly label (see sequence symbol in L1C) so that you can branch to the macro instruction at pre-assembly time if you want the called macro definition expanded.

## Specifications

The name entry of a macro instruction can be:

1 an ordinary symbol
2 a variable symbol
3 a character string in which a variable symbol is concatenated to other characters
(4) a blank
(5) ge sequence symbol, which is never


Purpose

The symbolic operation code you specify identifies the macro definition you wish the assembler to process.

## Specifications

The operation entry for a macro
(1) instruction must be a valid symkol that is identical to the symbolic operation code specified in the prototype statement of the macro definition called.

NOTE: If a source macro definition
(2) with the same operation code as a library macro definition is called. the assembler processes the source (3) macro definition.

```
K2C -- THE CFERANL ENIRY
```


## Purpose

You can use the operand entry of a macro instruction to pass values into the called macro definition. These values can be passed through:

1. The symbolic parameters you have specified in the macro prototype, or
2. The system variable symbol ESYSLIST if it is specified in the body of the macro definition (see J7C).

The two types of operands allowed in a macro instruction are the positional cperand and the keyword operand (see K3). You can specify a sublist with multiple values in both types of operands (see K4). Special rules for the various values you can specify in operands are given in K 5 .


## K3- Operands

K3A -- POSITIONAL OPERANDS

## Purpose

You can use a positional operand to pass a value into a macro definition through the corresponding positional parameter declared for the definition. you should declare a positional parameter in a macro definition when you wish to change the value passed at every call to that macro definition.

You can also use a positional operand to pass a value to the system variable symbol $\varepsilon S Y S L I S T$. If ESYSLIST, with the appropriate subscripts, is specified in a macro definition, you do not need to declare positional parameters in the prototype statement of the macro definition. Ycu can thus use $\varepsilon S Y S L I S I$ to refer to any positional operand. This allows you to vary the number of operands you specify each time you call the same macro definition. The use of ESYSLIST is described in J7C.

## Specifications

1 The positional operands of a macro instruction must be specified in the same order as the positional
0 parameters declared in the called macro definition.

Each positional operand constitutes
a character string. It is this character string that is the value passed through a positional parameter into a macro definition.


Examples of Macro Instructions:


The figure to the right illustrates what happens when the number of positional operands in the macro instruction differs from the number of positional parameters declared in the prototype statement of the called macro definition.

| Number of positional parameters in Prototypt of macro definition | Number of Positional Operands in macro instruction |  |  |
| :---: | :---: | :---: | :---: |
|  | EQUAL | GREATER <br> THAN | $\begin{aligned} & \text { LESS } \\ & \text { THAN } \end{aligned}$ |
|  | Valid, if Operands are correctly specified |  |  |
|  |  | Meaningless, unless \&SYSLIST is specified in definition to refer to excess operands | $\cdots$ |
|  |  |  | Omitted operands give null character values to corresponding parameters (or \&SYSLIST specification) |

## Purpose

You can use a keyword operand to pass a value through a keyword parameter into a macro definition. The values you specify in keyword operands override the default values assigned to the keyword parameters. The default value should be a value you use frequently. Thus, you avoid having to write this value every time you code the calling macro instruction.

When you need to change the default value, you must use the corresponding keyword operand in the macro instruction. The keyword can indicate the purpose for which the passed value is used.

## Specifications

Any keyword operand specified in a macro instruction must correspond to a keyword parameter in the macro definition called. However, keyword operands do not have to be specified in any particular order.

A keyword operand must be coded in the format shown in the figure to the right. If a keyword operand 1 is specified, its value overrides the default value specified for
the keyword parameter.
The standard default value obeys the same rules as the value specified in the keyword operand (see K5).


The following examples describe the relationship between keyword operands and keyword parameters and the values that the assembler assigns to these parameters under different conditions.

The keyword of the operand corresponds to a keyword parameter. The value in the operand overrides the default value of the parameter.

2 The keyword operand is not specified. The default value of the parameter is used.
(3)

The keyword of the operand does not correspond to any keyword parameter. The assembler issues an error message, but the macro is generated using the default values of the other parameters.

NOTE: The default value specified for a keyword parameter can be the 4 null character string. The null character string is a character string with a length of zero; it is not a blank, because a blank occupies one character position.


Purpose

You can use positional and keyword operands in the same macro instruction: use a cositional operand for a value that you change often and a keyword operand for a value that you change infrequently.

## Specifications

Positional and keyword operands can be mixed in the macro instruction operand field. However, the
(1) positional operands must be in the same order as the corresponding positional parameters in the macro prototype statement.

Dos All positional operands must precede any keyword operands, if specified.



END

NOTE: The system variable symbol
1 ESYSLIST ( $n$ ) refers only to the positional operands in a macro instruction.

Dos All keyword operands must follow 3 any positional operands specified.

## K4 -- Sublists in Operands

## Purpose

You can use a sublist in a positional or keyword operand to specify several values. A sublist is one or more entries separated by commas and enclosed in parentheses. Each entry is a value to which you can refer in a macro definition by coding:

1. The corresponding symbolic parameter with an appropriate subscript or
2. The system variable symbol \&SYSLIST with appropriate subscripts, the first to refer to the positional operand and the second to refer to the sublist entry in the operand.
\&SYSLIST can refer only to sublists in positional operands.

## Specifications

The value specified in a positional or keyword operand can be a sublist.

A symbolic parameter can refer to the entire sublist or to an individual entry of the sublist. To refer to an individual entry,
2 the symbolic parameter must have a subscript whose value indicates the position of the entry in the sublist. The subscript must have a value greater than or equal to one.


The format of a sublist is given in the figure to the right. A sublist, including the enclosing parentheses, must not contain more than 255 characters.

Format:

Examples:
One or more entries separated by commas and enclosed in parentheses
$A, B, C, D, E$


| Parameter | Sublist specified in corresponding operand (or as default value of keyword parameter) | Value generated (or used in computation) |
| :---: | :---: | :---: |
| \&PAR (3) | $(1,2,4)$ | Null character string |
| \& PAR (5) | $(1,2,3,4)$ | Null character string |
| \&PAR <br> \&PAR (1) <br> \&PAR (2) |  | A <br> A <br> Null character string |
| \&PAR <br> \& PAR (1) <br> \&PAR (2) <br> $\& P A R$ <br> \&PAR (1) <br> \&PAR (3) | $\left.\begin{array}{l}\text { (A) } \\ \text { (A) } \\ \text { (A) } \\ \text { () } \\ \text { () } \\ \text { () }\end{array}\right\}$Considered as <br> Sublists$\$$. | (A) <br> A <br> Null character string <br> () <br> Null character string <br> Null character string |
| \&PAR (2) <br> \&PAR (1) |  | $\begin{aligned} & \text { Nothing } \\ & \left\{\begin{array}{l} \text { *ERROR } * \\ \text { Unmatched left } \\ \text { parentheses } \end{array}\right. \\ & \text { Nothing } \end{aligned}$ |
| \&POSPAR (3) <br> \&SYSLIST $(2,3)$ | $\begin{aligned} & \text { Positional Operands } \\ & \hline \text { A, }(1,2,3,4) \\ & \mathrm{A},(1,2,3,4) \end{aligned}$ | $\begin{aligned} & 3 \\ & 3 \end{aligned}$ |

## K5 -- Values in Operands

Purpose

You can use a macro instruction operand to pass a value into the called macro definition. The two types of value you can pass are:

1. Explicit values or the actual character strings you specify in the operand.
2. Implicit values, or the attributes inherent in the data represented by the explicit values.

Attrikutes are fully described in L1B.

Specifications

The explicit value specified in a macro instruction operand is a character string that can contain one or more variable symbols.

The character string must not be greater than 255 characters after substitution of values for any variable symbols. This includes a character string that constitutes a sublist (see K4).

The character string values, including sublist entries, in the operands are assigned to the corresponding parameters declared in the prototype statement of the called macro definition. A sublist entry is assigned to the corresponding subscripted parameter.



APOSTROPHES: A single apostrophe is used: (1) to indicate the beginning and end of a quoted string, and (2) in a length attribute notation that is not within a quoted string.

QUOTED STRINGS: A quoted string is any sequence of 1 characters that begins and ends with a single apostrophe (compare with conditional assembly character expressions 2 described in L4B). Couble apostrophes must ke specified inside each quoted string. This includes substituted
apostrophes.
Macro instruction operands can have values that include one or more quoted strings. Each quoted string can be separated from the following quoted string by one or more
(4) characters, and each must contain an even number of apostrophes.

|  |  |  |
| :---: | :---: | :---: |
| Value specified in Operand | Value of <br> Variable <br> Symbol | Value Passed |
| '\&\&NOTATION' <br> '\&MESSAGE' | BLANKS OK | '\&\&NOTATION' <br> 'BLANKS OK' |
| $2\left\{\begin{array}{l}\text { 'L''SYMBOL' } \\ \text { 'L''\&VAR' }\end{array}\right.$ <br> ' \&QUOTES' |  | $\begin{aligned} & \text { 'L''SYMBOL' } \\ & \text { 'L''SYMBOL' } \\ & \hline \text { ITr'I' } \end{aligned}$ |
|  |  | INVALID OPERAND VALUE |
| 'QUOTE1 'AND 'QUOTE2 ' 4 <br> No apostrophes, single ampersands, commas, blanks, or equal signs allowed between quoted strings in one operand <br> $\underbrace{\text { 'AB' }{ }^{\prime} C D ' E '}_{\text {Quoted strings }}$ |  | 'QUOTEL'AND 'QUOTE2 <br> 'AB' $C D$ 'E'FGH\&\&' |

LENGTH ATTRIBUTE NOTATION: In macro instruction operand values, the length attribute notation with ordinary symbols can be used outside of quoted strings, if the length attribute notation is preceded by any special character except the ampersand.

PARENTHESES: In macro instruction operand values, there must be an equal number of left and right parentheses. They must be paired, that is, to each left parenthesis belongs a following right parenthesis at the same level of nesting. An unpaired (single) left or right parenthesis can appear only in a quoted string.


BLANKS: One or more blanks outside a quoted string indicates the end of the entire operand field of a macro instruction. Thus blanks should only be used inside quoted strings.

COMMAS: A comma outside a quoted string indicates the end of an operand value or sublist entry. Commas that do not delimit values
 can appear inside quoted strings paired parentheses that do not enclose sublists.

EQUAL SIGNS: An equal sign can appear in the value of a macro instruction operand or sublist entry:
(1) As the first character,
2. Inside quoted strings,
(3) Between paired parentheses, or

In a positional parameter, provided that the parameter does not resemble a keyword parameter. In a keyword parameter


Examples of Macro Instructions:
MACCALL A,B,C,D
MACCALL ( $\mathrm{A}, \mathrm{B}, \mathrm{C}, \mathrm{D}$ )


Examples of Macro Instructions:


PERIODS: A period (.) can be used in the value of an operand or sublist entry. It will be passed as a period.

However, if it is used immediately after a variable symbol it becomes a concatenation character. Then, two periods are required if one is to be passed as a character.

## K6 - Nesting in Macro Definitions

## K6A -- PURPOSE

A nested macro instruction is a macro instruction that you specify as one of the statements in the body of a macro definition. This allows you to call for the expansion of a macro definition from within another macro definition.

Inner and Outer Macro Instructions

Any macro instruction you write
in the open code of a source module is an outer macro instruction or call. Any macro instruction that appears within a macro definition
2 is an inner macro instruction or call.

| Character String specified as value of Operand or Sublist Entry | Value of <br> Variable <br> Symbol | Value <br> Passed |
| :---: | :---: | :---: |
| $\left\lvert\, \begin{aligned} & 3.4 \\ & (3.4,3.5,3.6) \end{aligned}\right.$ |  | $\begin{aligned} & 3.4 \\ & 3.4 \\ & 3.5 \quad 3.6 \end{aligned}$ |
| $\begin{aligned} & \& A!1 \\ & \& A \cdot 1 \\ & \& A!.1 \\ & \& A \& B \\ & \& A!\& B \end{aligned}$ | $\left.\begin{array}{c} \text { FIELD } \\ 3 \\ 3 \\ \& A=A R E A \\ \& B=200 \end{array}\right\}$ | $\begin{aligned} & \text { FIELD1 } \\ & 31 \\ & 3!1 \\ & \text { AREA200 } \\ & \text { AREA200 } \end{aligned}$ |
| \&DISP.(\&BASE) | $\begin{gathered} \& D I S P=1000 \\ \& B A S E=10 \end{gathered}$ | 1000(10) |



## Levels of Nesting

The code generated by a macro definition called by an inner macro call is nested inside the code generated by the macro definition that contains the inner macro call. In the macro definition called by an inner macro call, you can include a macro call to another macro definition. Thus, you can nest macro calls at different levels.

2 The zero level includes outer macro calls, calls that appear in open code; the first level of nesting includes inner macro calls that appear inside macro definitions called from the zero level; the second level of nesting includes inner macro calls inside macro definitions that are called from the first level, etc.


## Recursion

You can also call a macro definition recursively, that is, you can write macro instructions inside macro definitions that are calls to the containing definition. This allows you to define macros to process recursive functions.


## General Rules and Restrictions

Macro instruction statements can
be written inside macro definitions. Values are substituted in the same way as they are for the model statements of the containing macro definition. The assembler processes the called macro definition, passing to it the operand values (after substitution) from the inner macro instruction. In addition to the operand values described in K 5 above, nested macro calls can specify values that include:

1 Any of the symbolic parameters specified in the prototype statement of the containing macro definition

2 Any SET symbols declared in the containing macro definition

Any of the system variable symbols OS
only
The number of nesting levels permitted depends on the complexity and size of the macros at the different levels, that is: the number of operands specified, the number of local and global SET symbols declared (see L1A) and the number of sequence symbols used.


## Source Module



Passing Values thrcugh Nesting Ievels

The value contained in an outer macro instruction operand can be passed through one cr more levels of nesting. However, the value specified in the inner macro instruction operand must be identical to the corresponding symbclic parameter declared in the prototype of the containing racrc definition.
3 Thus, a sublist can be passed and referred to as a sublist in the macro definition called by the inner macro call. A1so, any symbol that is passed will carry its inherent attribute values through the nesting levels.

Values can be passed from open code through several levels cf macro nesting if inner macro calls at each level are specified with symbolic parameters as operand values.


Source Module

NOTE: If a symbolic parameter is only a part of the value specified in an inner macro instruction
2 operand, only the character string value given tc the parameter by an outer call is passed through
(3) the nesting level. Inner sublist entries and attributes of symbols 4 are nct available fcr reference in the inner macro.


System Variable Symbols in Nested Macros

The global read-only system variable symbols: ESYSPARM, GSYSDATE, and gSYSTIME are not affected by the nesting of macros. The remaining system variable symbols are given local read-only values that depend on the position of a macro instruction in code and the operand value specified in the macro instruction.

If ESYSLIST is specified in a macro definition called by an inner macro instruction, then ESYSLIST refers to the positional operands of the inner macro instruction.

(1) The assembler increments ESYSNDX by one each time it encounters a macro call. It retains the incremented value 2
3 throughout the expansion of the macro definition that is called, that is, within the local scope of the nesting level.


The assembler gives $\begin{gathered}\text { SYSECT } \\ \text { the character string value }\end{gathered}$ of the name of the control section in force at the point

2where a macro call is made. For a macro definition called by an inner macro call, the assembler will assign 6 SYSECT 3 the name of the control section generated in the macrc definition that contains the inner macro call. The control processed.

If no control section is generated within a macro definition, the value assigned to छSYSECT does not change. It is the same for the next level of macro definition called by an inner macro instruction. (6)
7. ESYSECT has a local scope; its read-only value remains constant throughout the expansion of the called macro definition.


## Section L: The Conditional Assembly Language

```
This section describes the conditional assembly language.
With the conditional assembly language, you can perform
general arithmetic and logical computations as well as
many of the other functions you can perform with any cther
programming language. In addition, by writing conditional
assembly instructions in combination with other assembler
language statements you can:
1. Select sequences of these source statements, called
model statements, from which machine and assembler
instructions are generated
2. Vary the contents of these model statements during
generation
The assembler processes the instructions and expressions
of the conditional assembly language at pre-assembly time.
Then, at assembly time, it processes the generated
instructions. Conditional assembly instructions, however,
are not processed after pre-assembly time.
The conditional assembly language is more versatile when used to interact with symbolic parameters and the system variable symbols inside a macro definition. However, you can also use the conditional assembly language in open code as described in \(L 7\) below.
```


## Ll -- Elements and Functions

```
The elements of the conditional assembly language are
1. SET symbols that represent data (see L1A)
2. Attributes that represent different characteristics
of data (see L1B)
3. Sequence symbols that act as labels for branching to
statements at pre-assembly time (see L1C).
The functions of the conditional assembly language are:
```

```
    1. Declaring SET symbols as variables for use by the
        conditional assembly language in its computations (see
        L2)
    2. Assigning values to the declared SET symbols (see L3)
    3. Evaluating conditional assembly expressions used as
    values for substitution, as subscripts for variable symbols,
    or as condition tests for branch instructions (see L4)
    4. Selecting characters from strings for substitution in
    and concatenation to other strings, or for inspection in
    condition tests (see L5)
    5. Branching and exiting from conditional assembly loops
    (see L6).
```

L1A - SET SYMBOLS
Purpose
SET symbols are variable symbols that provide you with
arithmetic, binary, or character data, whose values you
can vary at pre-assembly time.
You can use SET symbols as:

1. Terms in conditional assembly expressions
2. Counters, switches, and character strings
3. Subscripts for variable symbols
4. Values for substitution.
Thus, SET symbols allow you to control your conditional
assembly logic and to generate many different statements
from the same model statement.
SUBSCRIPTED SET SYMBOLS: You can use a SET symbol to
represent an array of many values. You can then refer
to any one of the values of this array by subscripting
the SET symbol.

You must declare a SET symbol before you can use it. The scope of a SET symbol is that part of a program for which the SET symbol has been declared.


If you declare a SET symbol to have a local scope, you can use it only in the statements that are part of:


- The same macro definition or
(3) Open code.

If you declare a SET symbol to have (4) a global scope, you can use it in the statements that are part of:

- The same macro definition, and
- A different macro definition, and
- Open code.

You must, however, declare the SET symbol as global for each part of 5 the program (a macro definition or open code) in which you use it.

You can change the value assigned to a SET symbol without affecting the scope of this symbol.

THE SCOPE OF OTHER VARIABLE SYMBOLS: A symbolic parameter has a local scope. You can use it only in the statements that are part of the macro definition for which the parameter is declared. You declare a symbolic parameter in the prototype statement of a macro definition.


The system variable symbols, ESYSLIST, ESYSECT, and ESYSNDX have a local scope; you can use them only inside macro definitions.
os However, the system variable symbols, only ESYSPARM, ESYSDATE, and ESYSTIME
have a global scope; you can use them in both open code and inside any macro definition.

## Specifications

SET symbols can be used in model statements from which assembler language statements are generated, and in conditional assembly instructions. The three types of SET symbols are: SETA, SETE, and SETC. A SET symbol must be a valid variable symbol, as shown in the figure to the right.

A SET symbol must be declared before
2 it can be used. The instruction that declares a SET symbol determines its scope and type (see L2).


The features of SET symbols and other types of variable symbol are compared in the figure to the right.

The value assigned to a SET symbol
can be changed by using the SETA, SETB, or SETC instruction within the declared scope of the SET symbol. However, a symbolic parameter and the system variable symbols are
(2) assigned values that remain fixed throughout their scope. Wherever a SET symbol appears in a statement, the assembler replaces the symbol with the last value assigned to the symbol.

| Feature | Types of Variable Symbol |  |  |
| :---: | :---: | :---: | :---: |
|  | SETA, SETB, or SETC Symbols | Symbolic Parameters | System <br> Variable <br> Symbols |
| Can be used : In open code | YES | NO | $\begin{aligned} & \text { only: \&SYSPARM } \\ & \text { OS \&SYSDATE } \\ & \text { only \&SYSTIME } \end{aligned}$ |
| In macro definitions | YES | YES | All |
| Scope: <br> Local or | YES | YES | \&SYSLIST \&SYSECT \&SYSNDX |
| Global | YES | NO | $\begin{gathered} \text { \&SYSPARM } \\ \text { OS SSYSDATE } \\ \text { \&SYSTIME } \end{gathered}$ |
| Values can be changed within scope of symbol | 1 YES | 2 <br> NO: <br> read only value | 2 <br> NO: <br> read only <br> value |

NOTE: SET symbols can be used in the name and operand field of macro ${ }^{2}$ instructions. However, the value
(3) thus passed through a symbolic parameter into a macro definition
4 is considered as a character string and is generated as such.

DOS The "LCLC ELIST, ELAEEL" instruction must precede the START instruction.


Subscripted SET Symbols -
Specifications

A subscripted SET symbol must be specified as shown in the figure to the right.

The subscript can be any arithmetic expression allowed in the operand field of a SETA instruction (see L4A) .

A subscripted SET symbol can be used anywhere an unsubscripted SET symbol is allowed. However, subscripted SET symbols must be declared as subscripted by a previous local or global declaration instruction.
(1) The subscript refers to one of the many positions in an array of values identified by the SET symbol. The value of the subscript must not
(3) exceed the dimension declared for the array in the corresponding LCLA, LCLB, LCLC, GBLA, GBLE, or GELC instruction.

NOTE: The subscrict can be a
4 subscripted SET symbcl. Five levels of subscript nesting are allowed.


## What Attributes Are

The data, such as instructions, constants, and areas, which you define in a.source module can ke described in terms of:

1. Type, which distinguishes one form of data from ancther: for example, fixed-point constants from floating-point constants, or machine instructions from macro instructions.
2. Length, which gives the number of bytes occupied by the object code of the data.
3. Scaling, which indicates the number of positions occupied by the fractional portion of fixed-point and decimal constants in their object code form.
4. Integer, which indicates the number of positions occupied
by the integer portion of fixed-point and decimal constants
in their object code form.
5. Count, which gives the number of characters that wculd be required to represent the data, such as a macro instruction operand, as a character string.
6. Number, which gives the number of sublist entries in a macro instruction operand.

These six characteristics are called the attributes of the data. The assembler assigns attribute values to the ordinary symbols and variable symbols that represent the data.

## Purpose

Specifying attributes in conditional assembly instructions allows you to control conditional assembly logic, which in turn can control the sequence and contents of the statements generated from model statements. The specific purpose for which you use an attribute depends on the kind of attribute being considered. The attributes and their main uses are shown in the figure to the right.
(1) NOTE: The number attribute of ESYSLIST ( m ) and ESYSLIST $(\mathrm{m}, \mathrm{n})$ is described in J7C.

## Specifications

FORMAT: The format for an attribute reference is shown in the figure to the right.

The attribute notation indicates the attribute whose value is desired.
2 The ordinary or variable symbol represents the data which possesses the attribute. The assembler substitutes the value of the attribute for the attribute reference.

WHERE ALLOWED: An attribute reference to the type, scaling, integer, count, and number attrikutes can be used only in a conditional assembly instruction. The length attribute reference can be used both in a conditional assembly instruction and in a machine or assembler instruction (for details on this use see C4C).

| Attribute | Purpose | Main Uses |
| :---: | :---: | :---: |
| Type | Gives a letter that identifies type of data represented | - In tests to distinguish between different data types <br> - For value substitution <br> - In macros to discover missing operands |
| Length | Gives number of bytes that data occupies in storage | - For substitution into length fields <br> - For computation of storage requirements |
| Scaling | Refers to the position of the decimal point in decimal, fixed-point and floating-point constants | - For testing and regulating the position of decimal points <br> - For substitution into a scale modifier |
| Integer | Is a function of the length and scaling attributes of decimal, fixedpoint, and floatingpoint constants | - To keep track of significant digits (integers) |
| Count | Gives the number of characters required to represent data | - For scanning and decomposing of character strings <br> - As indexes in substring notation |
| Number <br> 1 | Gives the number of sublist entries in a macro instruction operand sublist | - For scanning sublists <br> - As counter to test for end of sublist |

COMBINATION WITH SYMBOLS: The figure below shows the six kinds of attributes and the type of symbol with which the attributes can be combined.

NOTE: Whether or not an attribute reference is allowed in open code, in macro definitions, or in toth, depends on the type of symbol specified.

|  | Symbols Specified | ATtRIbUTES SPECIFIED |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $\begin{gathered} \text { Type } \\ \text { T' }^{\prime} \end{gathered}$ | Length L' | Scaling S' | Integer $I^{\prime}$ | Count $K^{\prime}$ | Number $\mathrm{N}^{\prime}$ |
|  | Ordinary Symbols | YES | YES | YES | YES | YES | YES |
| IN OPEN CODE | SET Symbols | YES | NO | NO | NO | YES | NO |
| $1$ | System Variable Symbols: \&SYSPARM,\&SYSDATE, \&SYSTIME | YES | NO | NO | NO | YES | NO |
|  | Ordinary Symbols | YES | YES | YES | YES | NO | NO |
|  | SET Symbols | YES | NO | No | No | YES | No |
|  | Symbolic Parameters | YES | YES | YES | YES | YES | YES |
| IN MACRO DEFINITIONS | System Variable Symbols. \&SYSLIST | yes | YES | YES | YES | YES | YES |
|  | \&SYSNDX,\&SYSPARM, \&SYSDATE, \&SYSECT, \&SYSTIME | YES | NO | NO | NO | YES | NO |
| IN OPEN CODE \{ | Ordinary Symbols | YES | YES | YES | YES | No | No |
|  | Ordinary Symbols | No | YES | No | No | No | NO |
| IN MACRO DEFINITIONS | Symbolic Parameters | YES | YES | YES | YES | YES | YES |
|  | System Variable Symbol \&SYSLIST | YES | YES | YES | YES | YES | YES |

ORIGIN OF VALUES: The value of
1 an attribute for an ordinary symbol specified in an attribute reference comes from the data represented by the symbol, as shown in the figure to the right.

The symbol must appear in the name field of an assembler or machine instruction, or in the operand field of an EXTRN or WXTRN instruction. The instruction in which the symbol is specified:

2. Must not contain any variable symbols, and
3. Must not be a generated instruction.

1
The value of an attribute for a variable symbol specified in an attribute reference comes from the value substituted for the variable symbol as follows (see also the figure to the right):

OS 1. For SET symbols and the system only variable symbols छSYSECT, ESYSNDX, ESYSPARM, ESYSDATE, and ESYSTIME, the attribute values come from the current data value of these symbols.
2. For symbolic parameters and the systern variable symbol, ESYSLIST, the values of the count and number attrioutes come from the operands of macro instructions.

The values of the type, length, scaling, and integer attributes, however, come from the values represented by the macro instruction operands, as follows:
a. If the operand is a sublist, the sublist as a whole has attributes; all the individual entries and the whole sublist have the same attributes as those of the first suboperand in the sublist (except for 'count', which can be different, and 'number', which is relevant only for the whole sublist).
b. If the first character or characters of the operand for sublist entry) constitute an ordinary symbol, and this symbol is followed by either an arithmetic operator (+,-,*, or /), a left parenthesis, a comma, or a blank, then the values of the attributes for the operand are the same as for the ordinary symbol.
c. If the operand (or sublist entry) is a character string other than a sublist or the character string described in b. above, the type attribute is undefined (U) and the length, scaling, and integer attributes are invalid.


VALUES: Because attribute references are allowed only in conditional assembly instructions, their values are available only at fre-assembly time, except for the length attribute, which can be referred to cutside conditional assembly instructions and is therefore also available at assembly time (see C4C).

NOTE: The system variable symbol, ESYSLIST, can be used in an attribute reference to refer to a macro instruction operand, and, in turn, to an ordinary symbol. Thus, any of the attribute values for macro instruction operands and ordinary symbols listed below can also be substituted for an attribute reference containing ESYSLIST.

THE TYPE ATTRIBUTE (T'): The type attribute has a value of a single alphabetic character that indicates the type of data represented by:

## (1) - An ordinary symbol

Dos NOTE: An ordinary symbol outside a only macro cannot be used as the operand of $T^{\prime}$ inside a macro in DOS assembler.
2. A macro instruction operand

NOTE: The type attribute of a sublist is set to the same value as the type attribute of the first element of the sublist.
os
only

- A SET symbol. 3

The type attribute reference can be used only in the operand field of the SETC instruction or as one of the values used for comparison in the operand field of a SETB or AIF instruction.

NOTE: Ordinary symbcls used in the name field of an ECU instruction have the type attribute value "U".

However, the third operand of an EQU instruction can be used explicitly to assign a type attribute value to the symbol in the name field.

T'

| Type Attribute | Data Characterized |
| :---: | :---: |
|  | 1 For ordinary symbols and outer macro instructions that are symbols |
|  | : Defined as labels for DC and DS instructions |
| A | A-type constant, implicit length, aligned (also CXD OS instruction label) |
| B | Binary Constant |
| C | Character Constant |
| D | Long floating-point constant, implicit length, aligned |
| E | Short floating-point constant, implicit length, aligned |
| F | Full-word fixed-point constant, implicit length, aligned |
| G | Fixed-point constant, explicit length |
| H | Half-word fixed-point constant, implicit length, aligned |
| K | Floating-point constant, explicit length |
| L | Extended floating-point constant, implicit length, aligned Packed decimal constant |
| os 0 | Q-type address constant, implicit length, aligned |
| only ${ }_{R}$ | A-, S-, Q-, V- or Y-type address constant, explicit length |
| S | S-type address constant, implicit length, aligned |
| V | V-type address constant, implicit length, aligned |
| X | Hexadecimal constant |
| Y | Y-type address constant, implicit length, aligned |
| z | Zoned decimal constant |
|  | : Defined as labels for assembler language statements |
| 1 | Machine instruction |
| M | Macro Instruction |
| w | CCW instruction |
| J | : Identified as control section name |
| $\begin{aligned} & \mathrm{T} \\ & \$ \end{aligned}$ | : Identified as external symbol by EXTRN or WXTRN instruction |
|  | A macro Instruction Operand that is: <br> A self-defining term <br> Omitted (has a value of a null character string) |
| $\begin{array}{ll} \hline \text { Os N } \\ \text { only } \end{array}$ | The value of a SETA or SETB variable |

When a symbol or macro instruction operand cannot be assigned any of the type attribute values listed in the preceding figure, the data represented is considered to be undefined and its type attribute is U. Specific cases of where $U$ is assigned as a type attribute value are given in the figure to the right.
(1)

The type attribute will be set to $U$ when the same ordinary symbol is used to define a label in more than one place, even though only one label will be generated by conditional assembly statements.

THE LENGTH ATTRIRUTE ( $\mathrm{L}^{\circ}$ ): The
(1) length attribute has a numeric value equal to the number of bytes occupied by the data that is represented by the symbol specified in the attritute reference.

If the length attribute value is desired for pre-assembly processing, the symbol specified in the attribute reference must ultimately represent the name entry of a statement in
open code. In such a statement, the length modifier (for LC and LS instructions) or the length field (for a machine instruction), if specified, must be a self-defining term. The length modifier or length field must not be coded as a multiterm expression, because the assembler does not evaluate this expression until assembly time.

The length attribute can also be specified outside conditional assembly instructions. Then, the length attribute value is not available for conditional assembly processing, but is used as a value at assembly time.

At pre-assembly time, an ordinary symbol used in the name field of an EQU instruction has a length attribute value of 1. At assembly time, the symbol has the same length attribute value as the first symbol of the expression in the first operand of the EQU instruction.

OS However, the second operand of an EQU only instruction can be used to assign a length attribute value to the symbol in the name field.

The Type Attribute Value=U is assigned to the following:
Ordinary symbols that are used as tabels:

- for the LTORG instruction
- for the EQU instruction without a third operand
- that are defined more than once
- for DC and DS statements that contain variable symbols
Example: U1 DC \& X'I'

DOS only

- for DC and DS statements that contain expressions as duplication factors

Example: DC (AA BBIF'15'

The SETC variable symbol
OS only
The system variable symbols: \&SYSPARM, \&ŞYSDATE,
and \&SYSTIME
Macro instruction operands that specify literals.
Inner macro instruction operands that are ordinary symbols.

1. The length attribute reference, when used in conditional assembly processing, can be specified only in arithmetic expressions (see L4).
2. A length attribute reference to a symbol with the type attribute value of $M, N, O, T, U$, or $\$$ will be flagged. The length attribute for the symbol will be given the default value of 1.

THE SCALING ATTRIBUTE (S'): The scaling attribute can be used only when referring to fixed-point, floating-point, or decimal, constants. It has a numeric value that is assigned as shown in the figure to the right.

## NOTES:

1. The scaling attribute reference can be used only in arithmetic expressions (see L4).
2. When no scaling attribute value

| Constant <br> Types <br> Allowed | Type <br> Attributes <br> Allowed | Value of Scaling <br> Attribute Assigned |
| :--- | :--- | :--- |
| Fixed-Point | H,F, and G | Equal to the value of the <br> scale modifier <br> (-187 through +346 |
| Floating- <br> Point | D,E,L, and K | Equal to the value of the <br> scale modifer <br> (0 through 14-D,E) <br> (0 through 28-L) |
| Decimal | Pand Z | Equal to the number <br> of decimal digits <br> specified to the right <br> of the decimal point <br> (0 through 31-P) <br> (0 through 16-Z) | can be determined, the reference is flagged and the scaling attribute is given the value of 1 .

THE INTEGER ATTRIBUTE (I'): The integer attribute has a numeric value that is a function of (depends on) the length and scaling attribute values of the data being referred to by the attribute reference. The formulas relating the integer attribute to the length and scaling attrikutes are given in the figure below.

NOTE: The integer attribute reference can be used only in arithmetic expressions (see L4).

| Constant <br> Type Allowed (attribute value) | Formula <br> Relating the Integer to the Length and Scaling Attributes | Examples | Values Of the Integer Attribute |
| :---: | :---: | :---: | :---: |
| Fixed-point ( $H, F$, and G) | I' $=8 * L^{\prime}-S^{\prime}-1$ | $\begin{aligned} & \text { HALFCON DC HS6' }-25.93^{\prime} \\ & 8 * 2-6-1 \\ & \text { ONECON DC FS8'100.3E-2' } \\ & 8 * 4-8-1 \end{aligned}$ | $\begin{gathered} 9 \\ 23 \end{gathered}$ |
| Floating-point ( $D, E, L$, and $K$ ) | $\begin{aligned} & \text { when } L^{\prime} \leq 8 \\ & I^{\prime}=2 *\left(L^{\prime}-I\right)-S^{\prime} \end{aligned}$ | ```SHORT DC ES2'46.415' \(2 *(4-1)-2\) LONG DC DS5'-3.729' \(2 *(8-1)-5\)``` | $4$ |
| Only for L-Type | $\begin{aligned} & \text { when } L^{\prime}>8 \\ & I^{\prime}=2 *\left(L^{\prime}-1\right)-S^{\prime}-2 \end{aligned}$ | EXTEND DC LSLO'5.312 ${ }^{\prime}$ $2 *(16-1)-10-2$ | 18 |
| Decimal equal to the number of decimal digits to the left of the assumed decimal point after the number is assembled <br> Packed (P) <br> Zoned (Z) | $I^{\prime}=2 * L^{\prime}-S^{\prime}-1$ $I^{\prime}=L^{\prime}-S^{\prime}$ | ```PACK DC P'+3.513' 2*3-3-1 03/513C ZONE DC Z'3.513' 4-3``` |  |

THE CCUNT ATTRIBUTE (K'): The count attribute applies only to macro instruction operands, to SET symbols, and to the system variable symbols. It has a numeric value that is equal to the number of characters:

1

- That constitute the macro
instruction operand, or
os a That would be required to
only represent as a character string
2 the current value of the SET symbol
3 or the system variable symbol.


## NOTES:

1. The count attribute reference can be used only in arithmetic expressions (see L4).
2. The count attribute of an omitted macro instruction operand has a default value of 0 .

THE NUMBER ATTRIBUTE (N'): The number attribute applies only to the operands of macro instructions. It has a numeric value that is equal 1 to the number of sublist entries in the operand.

NOTES:

1. The number attribute reference can be used only in arithmetic expressions (see L4).
2. N'ESYSLIST refers to the number of positional operands in a macro instruction, and N'ESYSLIST (m) refers to the number of sublist entries in the m-th cferand (for further details on the number attribute of ESYSLIST, see J7C).

| Macro Instruction <br> Operand Sublist | Vatue of <br> Number Attribute |
| :--- | :--- |
|  | 1+ number of commas <br> separating the entries |

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As Updated 28 Dec 1981
By TNL GN20-9372
L1C - SEQUENCE SYMBOLS

Purpose

You can use a sequence symbol in the name field of a statement to branch to that statement at pre-assembly time, thus altering the sequence in which the assembler processes your conditional assembly and macro instructions. You can thereby select the model statements from which the assembler generates assembler language statements for processing at assembly time.

Seq. Sym.

## Specifications

Sequence symbols must be specified as shown in the figure to the right.

Format


Examples: • SEQ

- Al234
-\#924

Statements in which
sequence symbols must not
be used as name entries

The following assembler instructions:

COPY
EQU
GBLA
Sequence symbols can be specified
in the name field of assembler
language statements and model
statements, except as noted in the figure to the right.

GBLB
GBLC
ICTL
ISEQ
LCLA
LCLB
LCLC
MACRO
OPSYN
DOS DSECT

The Macro prototype
instruction

Any instruction that already contains an ordinary symbol or variable symbol

Sequence symbols can be specified in the operand field of an AIF or AGO instruction to branch to a statement with the same sequence symbol as a label.
(1) A sequence symbol has a local scope. Thus, if a sequence symbol is used in an AIF or AGO instruction, the sequence symbol must be defined as a label in the same part of the program in which the AIF or AGO instruction appears; that is, in
2 the same macro definition or in
3 open code.

NOTE: A sequence symbol in the name field of a macro instruction
4 is not substituted for the parameter, if specified, in the name field of the corresponding prototype statement (for specifications, about the name entry of macro instructions see K 2 A ).


## L2 -- Declaring Set Symbols

```
You must declare a SET symbol before
you can use it. In the declaration,
you specify whether it is to have
a glocal or local sccpe. The
assembler assigns an initial value
to a SET symbol at its point of
declaration.
```

L2A -- THE LCLA, LCLB, AND LCLC
INSTRUCTIONS

Purpose

You use the LCLA, LCLB, and LCLC instructions to declare the local SETA, SETB, and SETC symbols you need.

Specifications
LCLA
LCLB
LCLC
The format of the LCLA, LCLB, and LCLC instruction statements is given in the figure to the right.

These instructions can be used anywhere in the body of a macro definition or in the open code portion of a source module.

| Name | Operation | Operand |
| :--- | :--- | :--- |
| Blank | LCLA, <br> LCLB, or <br> LCLC | One or more variable <br> symbols separated <br> by commas |

DOS The LCLA, LCLE, and LCLC
instructions, if specified, must appear immediately following any GBLA, GBLB, or GBLC instructions that may be specified.

If specified inside a macro definition, the global declaration instructions must appear immediately following the macro prototype statement. If specified outside a macro definition, the global declarations must appear first in open code; that is, they must follow any source macro definitions specified and precede the beginning of the first control section.


The assembler assigns initial values to these SET symbols as shown in the figure to the right.

LOCAL VARIAELE SYMBOLS MUST NOT BE MULTIPLY DEFINEL: A local SET variable symbol declared by the ICLA, LCLB, or LCLC instruction must not be identical to any other variable symbol used within the same local scope. The following rules apply to a local SET variable symbol:

1. Within a macro definition, it must not be the same as any symbolic 1 parameter declared in the prototype statement.
2. It must not be the same as any

2 global variable symbol (see L2E) declared within the same local scope.
3. The same variable symbol must not be declared or used as two different types of SET symbols, for example, as a SETA and a SETB symbol, within the same local scope.

NOTE 1: A local SET symbol should not begin with the four characters ESYS, which are reserved for system variable symbols (see J7).

DOS NOTE 2: The global declarations must precede the local declarations.


SUBSCRIPTED LOCAL SET SYMBOLS:
A local subscripted SET symbol is declared by the LCLA, LCLB, or LCLC instruction. This declaration must be specified as shown in the figure to the right.
(1) The maximum dimension allowed is 32,767.

Dos The maximum dimension allowed is 4095.

The dimension indicates the number of SET variables associated with the subscripted SET symbcl. The
(3. assembler assigns an initial value. to every variable in the array thus declared.

NOTE: A sutscripted local SET symbol can be used only if the declaration has a subscript, which represents a dimension; a nonsubscripted local SET symbol can be used only if the declaration had no subscript.


You use the GELA, GELE, and GBLC instructions to declare the global SETA, SETB, and SETC symbols you need.

## Specifications

The format of the GBLA, GELE, and GBLC instruction statements is given in the figure to the right.

These instructions can be used anywhere in the body of a macro definition or in the open code portion of a source module.

If specified inside a macro definition, the GBLA, GELE, and GBLC instructions must appear immediately following the macro prototype statement. If specified outside a macro definition, the global declarations must appear first in open code; that is, they must follow any source macro definitions specified and precede the beginning of the first control section.

Any variable symbols declared in the operand field have a global scope. They can be used as SET symbols anywhere after the pertinent GBLA, GBLB, or GBLC instructions. However, they can be used only within those parts of a program in which they have been declared as global SET symbols, that is in any macro definition and in open 3 code.

NOTE: Values can be passed between:
(4)

- The macro definitions, MAC1, and MAC2, only by using the variable symbols EB and EC.
- The macro definition, MAC2, and open code, only by using the variatle symbol EC.
- The macro definition, MAC1, and open code, only by using the variakle symbol EC.

DOS NOTE: The "GBLC EC" instruction


The assembler assigns initial values to these SET symbols as shown in the figure to the right.

The assembler assigns this initial value to the SET symbol only when it processes the first GBLA, GBLB, or GBLC instruction in which the symbol appears. Subsequent GBLA, 2 GBLB, or GBLC instructions do not reassign an initial value to the SET symbol.

DOS NOTE: The "GBLA EA" instruction (3) must precede the START instruction.

| Instruction | Initial Value assigned <br> to SET variable symbols <br> in operand field |
| :---: | :---: |
| GBLA | 0 |
| GBLB | 0 |
| GBLC | Null character string |



GLOBAL VARIABLE SYMECLS MUST NOT BE MULTIPLY DEFINED: A global SET variable symbol declared by the GBLA, GBLB, or GBLC instruction must not be identical to any other variable symbol used in open code or within the same macro definition. The following rules apply to a global SET variable symbol:

1. Within a macro definition, it must not be the same as any symbolic
 parameter declared in the prototype statement.
2. It must not be the same as any local variable symbol (see L2A) declared within the same local scope.
3. The same variable symbol must not ke declared or used as two
different types of global SET symbol, for example, as a SETA or SETB symbol.

NOTE 1: A global SET symbol should not legin with the four characters ESYS, which are reserved for system variable symbols (see J7).

DOS NOTE 2: The qlobal declarations must precede the local declarations.

SUBSCRI PTED GLOBAL SFT SYMBOLS: A glokal subscripted SET symbol is declared by the GBLA, GBLB, or GBLC instruction. This declaration must be specified as shown in the figure to the right.
(1)

The maximum dimension allowed is $32,7 \overline{67}$.

DOS The maximum dimension allowed is 4095.

The dimension indicates the number of SET variables associated with
2. the subscripted SET symbol. The assembler assigns an initial value
3 to every variable in the array thus declared.


NOTES:
(1) A subscripted global SET symbol can be used only if the declaration has a subscript, which represents
2 a dimension; a nonsubscripted global SET symbol can be used only if the declaration had no subscript.
2. Wherever a particular global SET symkol is declared with a
dimension as a subscript, the
0 dimension must be the same in each declaration.

DOS NOTE: The "GBLB ESWITCH (50)" instruction must precede the START
4 instruction.


## L3 -- Assigning Values to Set Symbols

## L3A -- THE SETA INSTRUCTION

## Purpose

The SETA instruction allows you to assign an arithmetic value to a SETA symbcl. You can specify a single value or an arithmetic expression from which the assembler will compute the value to assign.

You can change the values assigned to an arithmetic or SETA symbol. This allows you to use SETA symbols as counters, indexes, or for other repeated computations that require varying values.

## Specifications

The format of the SETA instruction statement is given in the figure to the right.

The variable symbcl in the name field must have been previously
(1) declared as a SETA symbol in a GBLA or LCLA instruction.

OS The variable symbol is assigned only a type attribute value of $N$.

The assembler evaluates the arithmetic expression in the operand
 field as a signed 32-bit arithmetic
value and assigns this value to the SETA symbol in the name field. An arithmetic expression is descrived in L4A.

SUBSCRIPTED SETA SYMBOLS: The SETA symbol in the name field can be subscripted, but only if the same SETA symbol has been previously declared in a GBLA or LCLA instruction with an allowable dimension.

The assembler assigns the value of the expression in the operand field to the position in the declared array given by the value of the subscript. The subscript expression must not be 0, or have a negative 4 value, or exceed the dimension actually specified in the declaration.

## purpose

The SETC instruction allows you to assign a character string value to a SETC symbol. You can assign whole character strings or concatenate several smaller strings together. The assembler will assign the composite string to your SETC symbol. You can also assign parts of a character string to a SETC symbol by using the substring notation (see L5) .

You can change the character value assigned to a SETC symbol. This allows you to use the same SETC symbol with different values for character comparisons in several places or for substituting different values into the same model statement.

Specifications

The format of the SETC instruction statement is given in the figure to the right.

The variable symbol in the name field must have been previously declared as a SETC symbol in a GBLC or LCLC instruction.
os The variable symbol is assigned
only a type attribute value of $U$.
The four opticns that can be specified in the operand field are:

## 1 <br> 1. A type attribute reference

(2) 2. A character expression (see L4B)
(3) 3. A substring notation (see L5)
(4) 4. A concatenation of substring notations, or character expressions, or both.

5 The assembler assigns the character string value represented in the operand field to the SETC symbol in the name field. The string length must be in the range 0 (null character string) through 255 characters.


NOTE: When a SETA or SETB symbol is specified in a character expression, the unsigned decimal value of the symbol (with leading zeros removed) is the character value given to the symbol.

03
city
4 A duplication factor can prectade any of the first three options, or any of the parts (character expression or substring notation)
(3) that make up the fourth option of the SETC instruction operamd. The duplication factor can be any
(44) arithmetic expression allowed in the operand of a SETA instruction.

NOTE: The assembler evaluates the character string represented (in particular the substring) before
(3) applying the duplication factor. The resulting character string is then assigned to the SETC symbol in the name field.

| Examples: |  |  | Value of \&A1 | Character Value Assigned to SETC symbol |
| :---: | :---: | :---: | :---: | :---: |
| $\begin{aligned} & \& C 1 \\ & \& C 2 \\ & \& C 3 \end{aligned}$ | $\begin{aligned} & \text { SETC } \\ & \text { SETC } \\ & \text { SETC } \end{aligned}$ | '\&Al' '\&A1' '\&A1' | $\begin{array}{r} 200 \\ 00200 \\ -200 \end{array}$ | (1) $\begin{aligned} & 200 \\ & 200 \\ & 200\end{aligned}$ |
| $\begin{aligned} & 8 C 4 \\ & 8 C 5 \end{aligned}$ | $\begin{aligned} & \text { SETC } \\ & \text { SETC } \end{aligned}$ | $\begin{aligned} & '-200 ' \\ & \text { '\&A1' } \end{aligned}$ |  | $\begin{array}{r} -200 \\ 0 \end{array}$ |
| \& 66 | SETC |  |  | 00200 |
| \&C7 | SETC | ' 8 Al+1' | 30 | 30+1 |
| \& C 8 | SETC | '1-\&Al' | -30 | 1-30 |


(1)

SUBSCRIPTED SETC SYMBOLS: The SETC symbol in the name field can be
subscripted, kut only if the same
2. SETC symbol has been previously declared in a GBLC or LCLC instruction with an allowable dimension.

The assembler assigns the character
3 value represented in the operand field to the position in the declared array given by the value of the subscript. The subscript expression must not be 0 , or have a negative value, or exceed the dimension actually specified in the declaration.


L3C - THE SETB INSTRUCTION

Purpose

The SETB instruction allows you to assign a kinary bit value to a SETB symbcl. You can assign the bit values, 0 or 1 , to a SETB symbol directly and use it as a switch.

If you specify a logical expression (see L4C) in the operand field, the assembler evaluates this expression to determine whether it is true or false and then assigns the values 1 or 0 respectively to the SETB symbol. You can use this computed value in condition tests or for substitution.

## Specifications

The format of the SETB instruction statement is given in the figure to the right.

The variable symbol in the name field must have been previously declared as a SETE symbol in a GELE or LCL.B instruction.

OS The variable symbol is assigned only a type attribute value of $N$.

The three options that can be specified in the operand field are:

1) 2. A binary value ( 0 or 1 )
2. A binary value enclosed in parentheses

OS NOTE: An arithmetic value enclosed only in parentheses is allowed. This value can be represented by an unsigned decimal self-defining term, a SETA symbol, or an attribute reference other than the type attribute reference. If the value is 0 , the assembler assigns a value of 0 to the symbol in the name field. If the value is not 0, the assembler assigns a value of 1.
3. A. A logical expression enclosed in parentheses (see L4C).

The assembler evaluates the logical expression, if specified; to determine if it is true or false. If it is true, it is given a value of 1 ; if it is false, a value of
0. The assembler assigns the explicitly specified binary value (0 or 1) or the computed logical value ( 0 or 1) to the $\operatorname{SETB}$ symbol in the name field.

1 (1. SUBSCRIPTED SETE SYMEOLS: The SETB symbol in the name field can be subscripted, but only if the same SETB symbol has been previously
declared in a GBLE or LCLE instruction with an allowable dimension.

The assembler assigns the binary value explicitly specified or implicit in the logical expression present in the operand field to 3 the position in the declared array given by the value of the subscript. The subscript expression must not be 0 , or have a negative value, 4 or exceed the dimension actually specified in the declaration.

SETB

| Format: <br> Name | Operation | Operand |  |
| :---: | :---: | :---: | :---: |
| A variable symbol | SETB | One of three options, exemplified below |  |
| Examples: |  |  |  |
| \& Bl | SETB | 0 (1) | 0 |
| \& B 2 | SETB | (1) 2 | 1 |
| \& B3A | SETB | $\begin{gathered} (2 . \mathrm{GT} \\ \substack{\text { Greater } \\ \text { Ghan } \\ \text { the }} \\ \text { false } \end{gathered}$ | 0 |
| \& B3B | SETB | $\left.\begin{array}{c} \left(\begin{array}{cc} 2 & \mathrm{LT}_{3} \end{array}\right) \text { true } \\ \text { Less than } \end{array}\right)$ | 1 |



## L4 - Using Expressions

> There are three types of expressions that you can use only in conditional assembly instructions: arithmetic, character, and logical. The assembler evaluates these conditional assembly expressions at pre-assembly time.
> Do not confuse the conditional assembly expressions with the absolute or relocatable expressions used in other assembler language instructions and descriked in c6. The assembler evaluates absolute and relocatable expressions at assembly time.

L4A -- ARITHMETIC (SETA) EXPRESSIONS

## Purpose

You can use an arithmetic expression for assigning an arithmetic value to a SETA symbol, or for computing a value used during conditional assembly processing.

An arithmetic expression can contain one or more SET symbols, which allows you to use arithmetic expressions wherever you wish to specify varying values, for example as:

1. Sukscripts for SET symbols, symbolic parameters, and ESYSLIST, and in substring notation.
```
OS 2. Duplication factors in the operand of the SETC
only instruction.
You can then control loops, vary the results of computations, and produce different values for sulstitution into the same model statement.
```


## Specifications

Arithmetic expressions can be used as shown in the figure to the right.

NOTE: When an arithmetic expression is used in the operand field of
(ii) a SETC instruction, the assembler assigns the character value representing the arithmetic expression to the SETC symbol, after
(2) substituting values into any variable symbols. It does not evaluate the arithmetic expression.

| Can be Used In | Used As | Example |
| :---: | :---: | :---: |
| SETA instruction | operand | \&A1 SETA \&Al+2 |
| AlF instruction or <br> SETB instruction | comparand in arithmetic relation | AIF (\&A*10 GT 30).A |
| Subscripted SET symbols | subscript | \&SETSYM (\&A $+10-\& C)$ |
| Substring notation <br> (See L6) | subscript | ${ }^{\prime} \& S T R I N G '(\& A * 2, \& A-1)$ |
| Sublist notation | subscript | sublist ( $A, B, C, D$ ) <br> when $\& A=1$ <br> $\& \operatorname{PARAM}(\& A+1)=B$ |
| \&SYSLIST | subscript | \&SYSLIST (\&M+1,\&N-2) <br> \&SYSLIST(N'\&SYSLIST) |
| SETC instruction | character string in operand | $\begin{aligned} & \& C \operatorname{SETC} C^{\prime} 5-10 * \& A^{\prime} \\ & \text { if \&A=10 } \\ & \text { then } \& C=5-10 * 10 \end{aligned}$ |

The figure below defines an arithmetic expression (selfdefining terms are described in C4E).

Arith Exp.
時 -


The variable symbols that are allowed as terms in an arithmetic expression are given in the figure to the right.

| Variable Symbol | Restrictions | Example | Value |
| :---: | :---: | :---: | :---: |
| SETA | none | - | - |
| SETB | none | - | - |
| $\left.\begin{array}{l} \text { SETC } \\ \text { \&SYSPARM } \end{array}\right\}$ | value must be an unsigned decimal self-defining term in the range 0 through <br> 2,147,483,647 <br> 0-99,999,999 | $\& \mathrm{C}$ <br> \&SYSPARM | $\begin{aligned} & 123 \\ & 2000 \end{aligned}$ |
| Symbolic <br> Parameters | value must be a self-defining term | \&PARAM <br> \&SUBLIST (3) | $\begin{aligned} & X^{\prime} A I^{\prime} \\ & C^{\prime} Z^{\prime} \end{aligned}$ |
| \&SYSLIST(n) <br> \&SYSLIST $(\mathrm{n}, \mathrm{m})$ | corresponding operand or sublist entry must be a self-defining term | $\begin{aligned} & \& \operatorname{SYSLIST}(3) \\ & \& \operatorname{SYSLIST}(3,2) \end{aligned}$ | $\begin{aligned} & 24 \\ & B^{\prime} 101 \end{aligned}$ |
| \&SYSNDX | none | - | - |

RULES FOR CODING ARITHMETIC
EXPRESSIONS: The following is a summary of coding rules for arithmetic expressions:

1. Both unary (operating on one value) and binary (operating on two values) operators are allowed in arithmetic expressions.
2. An arithmetic expressicn can have one or more unary operators preceding any term in the expression or at the beginning of the expression.

2 3. An arithmetic expression must not begin with a binary operator, and it must not contain two binary 3 operators in succession.
4. An arithmetic expression must 4 not contain two terms in succession.
5. An arithmetic expression must not contain blanks between an operator and a term nor between two successive operators.
6. An arithmetic expression can contain up to 24 unary and binary operators and up to 11 levels of parentheses.

DOS An arithmetic expression can contain up to 16 unary and binary operators and up to 5 levels of parentheses.

Note that the parentheses required for sublist notation, substring notation, and subscript notation count toward this limit.

Operators

| Unary | ,+- |
| :--- | :--- |
| Binary | $+,-, \%, l$ |

Examples

(2. $\left\{\begin{array}{l}* 3 \\ / \& A\end{array}\right.$

INVALID
INVALID
(3)
$\begin{cases}\& C * / \& D & \text { INVALID } \\ \& C+\# \& D & \text { INVALID }\end{cases}$

Leftmost operator between two terms is Binary


EVALUATION OF ARITHMETIC EXPRESSIONS:
The assembler evaluates arithmetic expressions at pre-assembly time as follows:

1. It evaluates each arithmetic term.
2. It performs arithmetic operations from left to right. However:
a. It performs unary operations before binary operations, and
b. It performs the binary operations of multiplication and division before the binary operations of addition and subtraction.
3. 3. In division, it gives an integer result; any fractional portion is dropped. Division by zero gives a 0 result.

1. In parenthesized arithmetic expressions, the assembler evaluates 1. the innermost expressions first and then considers them as arithmetic terms in the next cuter level of 3 expressions. It continues this process until the outermost expression is evaluated.
2. The computed result, including intermediate values, must lie in the range $-2^{31}$ through $+2^{31}-1$.


## Purpose

The main purpose of a character expression is to assign a character value to a SETC symbol. You can then use the SETC symbol to substitute the character string into a model statement.

You can also use a character expression as a value for comparison in condition tests and logical expressions (see L4C) . In addition, a character expression provides the string from which characters can be selected by the substring notation (see L5).

```
Substitution of one or more character values into a
character expression allows you to use the character
expression wherever you need to vary values for substitution
or to control loops.
```


## Specifications

Character (SETC) expressions can be used only in conditional assembly instructions as shown in the figure to the right.


A character expression consists of any combination of characters enclosed in apostrophes. Variable symbols are allowed. The assembler substitutes the representation of their values as character strings into the character expression before evaluating the expression.

Up to 255 characters are allowed in a character expression.

NOTE: Attribute references are not allowed in character expressions.

| Variable Symbol | Must not contain more than 255 characters (including blanks) |  | Value Substituted |
| :---: | :---: | :---: | :---: |
|  |  |  |  |
|  | Restrictions | Example |  |
| SETA | sign and leading zeros are suppressed stand alone zero is used |  | 201 0 |
| SETB | none | \& B SETB 1 | 1 |
| SETC | none |  | ABC |
| Symbolic Parameters | none | \&PARAM $=(\operatorname{ABC})$ $\& C l$ ¢ | (ABC) |
| System Variable symbols | none |  | 0201 |

EVALUATION CF CHARACTER EXPRESSIONS:
The value of a character expression is the character string within the enclosing apostrophes, after the assembler performs any substitution for variable symbcls.
(2)

Character strings, including variatle symbols, can be concatenated to

3 each other within a character
3 expression. The resultant string is the value of the expression used in conditional assembly operations: for example, the value assigned to a SETC symbol.

4
A double apostrophe must be used to generate a single apostrophe as part of the value of a character expression.

A double ampersand will generate a double ampersand as part of the value of a character expression. To generate a single ampersand in a character expression, use the substring notation, for example, ('\&\&' $(1,1)$ ).

NOTE: To generate a period, two periods must be specified after a variable symbol, or the variable symbol must have a period as part of its value.

| Examples <br> Concatenation operator is a period (.) | Value of Variable Symbols Used | Value of Character Expression |
| :---: | :---: | :---: |
| $\begin{aligned} & \text { 'ABC' } \\ & \text { '\&PARAM' } \\ & \text { 'A+B-C*D' } \\ & \text { '\&A+10' } \\ & \text { '\&A\&A' } \end{aligned}$ | SYMBOL <br> 10 <br> 15 | 1 ABC SYMBOL $\begin{aligned} & A+B-C * D \\ & 10+10 \\ & (\text { Not } 20) \\ & 1515 \end{aligned}$ |
|  | DEF <br> DEF $\begin{aligned} & \& A=200 \\ & \& C=A R E A \\ & \& C=. \end{aligned}$ | $\left.\begin{array}{l} \text { DEFABC } \\ \text { DEFDEF } \end{array}\right\} \begin{aligned} & 3 \\ & A R E A+10 \div 200 \\ & A B C . \end{aligned}$ |
| ${ }^{\prime} \& C^{\prime}$ <br> 'ABC\&C.DEF' | null $\& C=n u l l$ | null character string <br> ABCDEF |
| 4 'L"SYMBOL' <br> 5 $\begin{aligned} & \text { '\&C..505' } \\ & ' \& C .505^{\prime} \end{aligned}$ | 2 <br> 2. | L'SYMBOL <br> 2.505 <br> 2.505 <br> Resultant Value must be in the range 0 through 255 characters <br> DOS 0.127 |

CONCATENATION OF CHARACTER STRING VALUES: Character expressions can be concatenated to each other or to substring notations in any order. This concatenated string can then be used in the operand field of a SETC instruction or as a value for comparison in a logical expression.
(1) The resultant value is a character string composed of the concatenated parts.
(2) Note: The concatenation character (a period) is needed to separate the apostrophe that ends one character expression from the apostrophe that begins the next.


## L4C - LOGICAI (SETE) EXPRESSIONS

## Purpose

You can use a logical (Boolean) expression to assign the binary value 1 or 0 to a SETB symbcl.

You can also use a logical expression to represent the condition test in an AIF instruction. This use allows you to code a logical expression whose value ( 0 or 1 ) will vary according to the values substituted into the expression and thereby determine whether or not a branch is to be taken.


## Specifications

Logical (SETB) expressions can be used only in conditional assembly instructions as shown in the figure. to the right.

| Can be used in | Used As | Example |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| SETB instruction | operand | \& BI | SETB | ( \& B2 OR 8 | 8 GT | 3 |
| AIF instruction | condition <br> test part <br> of operand | AIF | (NOT \& Bl OR 8 EQ 3).A |  |  |  |

The figure on the opposite page defines a logical expression.

NOTE: An arithmetic relation is two arithmetic expressions separated by a relational operator. A character relation is two character strings (for example, a character expression and a type attribute reference) separated ky a relational operator. The relational operators are:

| EQ | (equal) |
| :--- | :--- |
| NE | (not equal) |
| LE | (less than or equal) |
| LT | (less than) |
| GE | (greater than or equal) |
| GT | (greater than) |



RULES FOR CODING LOGICAL EXPRESSIONS:
The following is a summary of coding rules for logical expressions:

1. A logical expression must not contain two logical terms in succession.
2. A logical expression can kegin (1) with the logical cperator NOT.
3. A logical expression can contain two logical operators in succession; however, the only combinations
2
allowed are: OR NOT or ANL NOT. The two operators must be separated from each other by one or more blanks.
4. Any logical term, relation, or inner logical expression can be optionally enclosed in parentheses.
5. The relational and logical operators must be immediately preceded and followed by at least 3 one blank or other special character. (4)
6. A logical expression can contain up to 18 logical operators and up to 17 levels of parentheses.

Dos A logical expression can contain up to 18 logical operators and up to 5 levels of parentheses.

Note that the relational and other operators used by the arithmetic and character expressions in relations do not count toward this total.

Examples of Logical Expressions
(\&A GT 100 OR ' \&C' EQ F)

(NOT \&B OR \&A GE 10 AND \&A LE 0)

blank mandatory

EVALUATION CF LOGICAL EXPRESSIONS: The assembler evaluates logical expressions as follows:

1. It evaluates each logical term, which is given a binary value of 0 or 1.
2. If the logical term is an arithmetic or character relation, the assembler evaluates:
a. The arithmetic or character expression specified as values for comparison in these relations, and then
b. The arithmetic or character relation, and finally
c. The logical term, which is the result of the relation. If the relation is true, the logical term it represents is given a value of 1 ; if the relation is false, the term is given a value of 0 .

NOTE: If two comparands in a character relation have character values of unequal length, the assembler always takes the shorter character value to be less than the longer one.

('ABC' LT 'ABCD') $\begin{aligned} & \left.\begin{array}{l}\text { Always true } \\ \text { (given a value of 1) }\end{array}\right)\end{aligned}$
('ABC' LT 'ABCD') $\begin{aligned} & \text { Always true } \\ & \text { (given a value of 1) }\end{aligned}$

## 3

(\&B AND NOT (5 GT 3))

(\&B AND (NOT (5 GT 3)))
(\&B OR \&A AND ('\&C' EQ 'B'))
(\&B OR (\&A AND ('\&C' EQ 'B')))

$\left(\& B\right.$ OR $\& A$ AND ( $\& C$ B $\left.B^{\prime}\right)$

3. The assembler performs logical

operations from left to right. However: a. It performs logical NOTS before logical ANDs and ORs, and
(5) b. It performs logical ANDS before logical ORs.
4. In parenthesized logical expressions, the assembler evaluates the innermost expressions first and then considers them as logical terms in the next outer level of expressions. It continues this process until the outermost expression is evaluated.


## L5 -- Selecting Characters from a String

L5A -- SUESTRING NOTATION

## Purpose

The substring notation allows you to refer to one or more characters within a character string. You can therefore either select characters from the string and use them for substitution or testing, or scan through a complete string, inspecting each character. By concatenating substrings with other substrings or character strings, you can rearrange and build your own strings.

## Specifications

The substring notation can be used only in conditional assembly instructions as shown in the figure below.

| Can be Used in | Used as | Example |  |  | Value Assigned to SETC symbol |
| :---: | :---: | :---: | :---: | :---: | :---: |
| SETC <br> instruction <br> operand | operand | \& C 1 | SETC | ' ABC ' $(1,3)$ | ABC |
|  | part of operand | \&C2 | SETC | '\&Cl' $(1,2)$. DEF' | ABDEF |
| SETB or AIF instruction operand (logical - expression) | Character value in comparand of character relation |  | $\begin{aligned} & \text { AIF } \\ & \text { SETB } \end{aligned}$ | $\begin{aligned} & (' \& S T R I N G '(1,4) \text { EQ 'AREA').SEQ } \\ & (' \& S T R I N G '(1,4) . ' 9 ' ~ E Q ~ ' F U L L 9 ') ~ \end{aligned}$ |  |

The substring notation must be specified as shown in the figure to the right.


The character string is a character expression from which the substring is to be extracted. The first
 subscript indicates the first character that is to be extracted from the character string. The second subscript indicates the number of characters to be extracted from the character string, starting with the character indicated by the first subscript. Thus the second subscript specifies the length of the resulting substring.

| CHMDTCTEA STRT Examples | (01), (e) <br> Value of Variable Symbol | Character Value of Substring |
| :---: | :---: | :---: |
| 'ABCDE' ( 1,5 ) |  | ABCDE |
| ' $\operatorname{ABCDE}$ ' 2,3$)$ |  | BCD |
| '\&C' $(3,3)$ | ABCDE | CDE |
| ' \&PARAM' $(3,3)$ | - ( $(\mathrm{A}+3) * 10)$ | A+3 |
|  |  |  |

The character string must be a valid character expression with a length, N , in the range 1 through 255 characters.

The length of the resulting substring must be within the range $0-255$.

The subscripts, e1, and e2, must be arithmetic expressions. The substring notation is replaced by a value that depends on the three elements: $N$, e1, and e2, as summarized below:
(1) In the usual case, the assembler generates a correct substring of the specified length.

2 When e1 has a value of zero or a negative value, the assembler issues an error message.

3 When the value of e1 exceeds $N$, the assembler issues a warning message, and a null string is generated.

4 When $e 2$ has a value of 0 , the assembler generates the null character string. Note that if e2 is negative, the assembler issues an error message.

5 When e2 indexes past the end of the character expression (that is, $\mathrm{e} 1+\mathrm{e} 2$ is greater than $\mathrm{N}+1$ ), the assembler issues a warning message and generates a substring which includes only the characters up to the end of the character expression specified.


## L6 - Branching

## L6A -- The AIF INSTRUCTION

## Purpose

The AIF instruction allows you to branch according to the result of a condition test. You can thus alter the sequence in which your assembler language statements are processed.

The AIF instruction also provides loop control for conditional assembly processing, which allows you to control the sequence of statements to be generated.

It also allows you to check for error conditions and thereby to branch to the appropriate MNOTE instruction to issue an error message.

The AIF instruction statement must be specified as shown in the figure

## Specifications

 to the right.| Name | Operation | Operand |
| :--- | :--- | :--- |
| A sequence <br> Symbol or <br> Blank | AIF | (logical expression). SEQUENC |

The assembler evaluates the logical expression in the operand field expression is true (logical value $=1$ ). the next statement processed by
2 the assembler is the statement named by the sequence symbol. If it is
3 false (logical value $=0$ ), the next sequential statement is processed. 4


The sequence symbol in the operand field is a conditional assembly label that represents an address at pre-assembly time. It is the address of the statement to which a branch is taken if the logical expression preceding the sequence symbol is true.

The statement identified by the sequence symbol referred to in the AIF instruction can appear before or after the AIF instruction. However, the statement must appear
1 within the local scope of the sequence symbol. Thus, the statement identified ky the sequence symbol must appear:

2 - In open code, if the corresponding AIF instruction does or
(3) - In the same macro definition in which the corresponding AIF instruction appears.

The sequence symbols . EACK and . FORWARD are not multiply defined. No branch can be taken from open code into a macro definition or between macro definitions, regardless of nested calls to other macro definitions.


NOTE: For compatibility, the assemblers described in this manual will process the AIFB instruction (BOS/360) in the same way they process the AIF instruction.

L6B -- THE AGO INSTRUCTION

## Purpose

The AGO instruction allows you to kranch unconditionally. You can thus alter the sequence in which your assembler language statements are processed. This provides ycu with final exits from conditional assembly loops.

The AGO instruction statement must be specified as shown in the figure to the right.

| Name | Operation | Operand |
| :--- | :--- | :--- |
| A sequence <br> symbol or <br> blank | AGO | SEQUENC |

The statement identified by a sequence symbol referred to in the AGO instruction can appear before or after the AGO instruction. However, the statement must appear

## 1

 within the local scope of the sequence symbol. Thus, the statement identified ky the sequence symbol must appear2 - In open code, if the corresponding AGO instruction does or
(3) In the same macro definition in which the corresponding AGO instruction appears.

NOTE: For compatibility, the assemblers described in this manual will process the AGOB instruction (BOS/360) in the same way they process the AGO instruction.


L6C -- THE ACTR INSTRUCTION

## Purpose

The ACTR instruction allows you to set a conditional assembly loop counter either within a macro definition or in open code.

Each time the assembler processes an AIF or AGO branching instruction in a macro definition or in open code, the loop counter for that part of the program is decremented ky one. When the number of conditional assembly branches taken reaches the value assigned ky the ACTR instruction to the loop counter, the assembler exits from the macro definition or stops processing statements in open code.

By using the ACTR instruction, you avoid excessive looping during conditional assembly processing at pre-assembly time.

## Specifications

ACTR
The format of the ACTR instruction statement is given in the figure to the right.

The ACTR instruction can appear anywhere in open code or within a macro definition.

A conditional assembly loop counter
is set (or reset) to the value of the arithmetic expression in the operand field. The loop counter has a local scope; its value is decremented only by AGO and AIF instructions and reassigned only by ACTR instructions that appear within the same scope. Thus, the nesting of macros has no effect on the setting of individual loop counters.

The assembler sets its own internal loop counter both for open code and for each macro definition, if neither contains an ACTR instruction. The assembler assigns a standard value of 4096 to each of these internal loop counters.


LOOP COUNTER OPERATIONS: Within
the local scope of a particular
loop counter (including the internal
counters run by the assembler).
the following occurs:

1. Each time an AGO or AIF (also AGOB or AIFB) branch is executed, the assembler checks the loop counter for zero or a negative value.
2. If the count is not zero or negative, it is decremented by one.
(1) 3. If the count is zero, before decrementing, the assembler will take one of two actions:
a. If it is processing instructions in open code, the assembler will process the remainder of the instructions in the source module as comments. Errors discovered in these instructions during previous passes are flagged.
b. If it is processing instructions inside a macro definition, the assembler terminates the expansion of that macro definition and processes the next sequential instruction afte $\bar{r}$ the calling macro instruction. If the macro definition is called by an inner macro instruction, the assembler processes the next sequential instruction after this inner call, that is, continues processing at the next outer level of nested macros (for levels of nesting see K6A).


NOTE: The assembler halves the ACTR counter value when it encounters serious syntax errors in conditional assembly instructions.

## Purpose

You can specify a sequence symbol in the name field of an ANOP instruction, and use the symbol as a label for branching furposes.

The ANOP instruction performs no operation itself, but you can use it to branch to instructions that already have symbols in their name fields. For example, if you wanted to branch to a SETA, SETB, or SETC assignment instruction, which requires a variable symbol in the name field, you could insert a labeled ANOF instruction immediately before the assignment instruction. By branching to the ANOP instruction with an AIF or AGO instruction, you would, in effect, be branching to the assignment instruction.

## Specifications

The format of the ANOP instruction statement is given in the figure to the right.

No operation is performed by an ANOP instruction. Instead, if a branch is taken to the ANOP instruction, the assembler processes the next sequential instruction.


## L7 -- In Open Code

## L7A -- PURPOSE

Conditional assembly instructions in open code allow ycu:

1. To select at pre-assembly time statements or groups
of statements from the open code portion of a source rodule according to a pre-determined set of conditions. The assembler further processes the selected statements at assembly time.
2. To pass local variable information from open code through parameters into macro definitions.
3. To control the computation in and generation of macro definitions using global SET symbols.
4. To substitute values into the model statements in the open code of a source module and control the sequence of their generation.

All the conditional assembly elements and instructions can be specified in open code.

1. Conditional assembly instructions can appear anywhere in open code, but they must appear after any source macro definitions that are specified.

Dos The global and local declaration instructions (see L2) must appear first in open code; that is, they must follow any source macro definitions specified and precede the beginning of the first control section.


The specifications for the conditional assembly language described in L1 through L6 also apply in open code. However, the following restrictions apply:

1. To attributes in open code: For ordinary symbols, only references to the type, length, scaling, and integer attributes are allowed.

NOTE: References to the number attribute have no meaning in open code, because ESYSLIST is not allowed in open code and symbolic parameters have no meaning in open code.
2. To conditional assembly
expressions in open code, as shown in the figure to the right.

| Expression | Must not contain |
| :---: | :---: |
| Arithmetic (SETA) | - \&SYSLIST <br> - Symbolic parameters <br> - Any attribute references to symbolic parameters, or \&SYSLIST,\&SYSECT,\&SYSNDX |
| Character (SETC) | - \&SYSLIST,\&SYSECT,\&SYSNDX <br> - Attribute references to \&SYSLIST, \&SYSECT, \&SYSNDX, or to symbolic parameters <br> Symbolic parameters |
| Logical (SETB) | - Arithmetic expressions with the items listed above <br> - Character expressions with the items listed above |

```
Purpose
The listing options allow you to
print the conditional assembly
statements in the sequence they
are processed. You can thus follow
the conditional assembly logic in
open code or in the code within
any macro definition.
```


## Specifications

Conditional assembly statements in the open code of a source module or in a macro definition can be printed in the program listings in the order in which they are processed, including iterations. This must be requested by specifying the desired options in the PARM field of the EXEC statement for the assembler program (job control language), or by specifying the options in fields set up by a program that dynamically invokes the assembler. The options are listed in the figure to the right.

NOTE: For other listing options see J8.

| Option | Action |
| :--- | :--- |
| NOALOGIC | No conditional assembly statements in open code <br> are printed |
| ALOGIC | All conditional assembly statements in open code <br> that are processed are printed, including iterations |
| NOMLOGIC | No conditional assembly statements inside macro <br> definitions, called from your program, are printed. <br> NOTE: Conditional assembly statements in source <br> macro definitions are always printed along with the <br> rest of the code in a source module (assuming the <br> PRINT option LIST) |
| MLOGIC | All conditional assembly statements inside macro <br> definitions, that are processed when you call the <br> macro, are printed, including iterations |

## Appendix I: Character Codes

| 8-Bit <br> EBCDIC <br> code | Character set Punch Combination | Decimal | HexaDecimal | Printer Graphics |
| :---: | :---: | :---: | :---: | :---: |
| 00000000 | 12,0,9,8,1 | 0 | 00 |  |
| 00000001 | 12,9,1 | 1 | 01 |  |
| 00000010 | 12,9,2 | 2 | 02 |  |
| 00000011 | 12,9,3 | 3 | 03 |  |
| 00000100 | 12,9,4 | 4 | 04 |  |
| 00000101 | 12,9,5 | 5 | 05 |  |
| 00000110 | 12,9,6 | 6 | 06 |  |
| 00000111 | 12,9,7 | 7 | 07 |  |
| 00001000 | 12,9,8 | 8 | 08 |  |
| 00001001 | 12,9,8,1 | 9 | 09 |  |
| 00001010 | 12,9,8,2 | 10 | 0 A |  |
| 00001011 | 12,9,8,3 | 11 | OB |  |
| 00001100 | 12,9,8,4 | 12 | 0 C |  |
| 00001101 | 12,9,8,5 | 13 | OD |  |
| 00001110 | 12,9,8,6 | 14 | OE |  |
| 00001111 | 12,9,8,7 | 15 | 0 F |  |
| 00010000 | 12,11,9,8,1 | 16 | 10 |  |
| 00010001 | 11,9,1 | 17 | 11 |  |
| 00010010 | 11,9,2 | 18 | 12 |  |
| 00010011 | 11,9,3 | 19 | 13 |  |
| 00010100 | 11,9,4 | 20 | 14 |  |
| 00010101 | 11,9,5 | 21 | 15 |  |
| 00010110 | 11,9,6 | 22 | 16 |  |
| 00010111 | 11,9,7 | 23 | 17 |  |
| 00011000 | 11,9,8 | 24 | 18 |  |
| 00011001 | 11,9,8,1 | 25 | 19 |  |
| 00011010 | 11,9,8,2 | 26 | 1A |  |
| 00011011 | 11,9,8,3 | 27 | 1 B |  |
| 00011100 | 11,9,8,4 | 28 | 1 C |  |
| 00011101 | 11,9,8,5 | 29 | 1D |  |
| 00011110 | 11,9,8,6 | 30 | 1 E |  |
| 00011111 | 11,9,8,7 | 31 | 1 F |  |
| 00100000 | 11,0,9,8,1 | 32 | 20 |  |
| 00100001 | 0,9,1 | 33 | 21 |  |
| 00100010 | 0,9,2 | 34 | 22 |  |
| 00100011 | 0,9,3 | 35 | 23 |  |
| 00100100 | 0,9,4 | 36 | 24 |  |
| 00100101 | 0,9,5 | 37 | 25 |  |
| 00100110 | 0,9,6 | 38 | 26 |  |
| 00100111 | 0,9,7 | 39 | 27 |  |
| 00101000 | 0,9,8 | 40 | 28 |  |
| 00101001 | 0,9,8,1 | 41 | 29 |  |
| 00101010 | 0,9,8,2 | 42 | 2A |  |
| 00101011 | 0,9,8,3 | 43 | 2B |  |
| 00101100 | 0,9,8,4 | 44 | 2C |  |
| 00101101 | 0,9,8,5 | 45 | 2D |  |
| 00101110 | 0,9,8,6 | 46 | 2E |  |
| 00101111 | 0,9,8,7 | 47 | 2 F |  |
| 00110000 | 12,11,0,9,8,1 | 48 | 30 |  |
| 00110001 | 9,1 | 49 | 31 |  |
| 00110010 | 9,2 | 50 | 32 |  |


| 8-Bit <br> EBCDIC <br> Code | Character set Punch Combination | Decimal | HexaDecimal | Printer <br> Graphics |
| :---: | :---: | :---: | :---: | :---: |
| 00110011 | 9,3 | 51 | 33 |  |
| 00110100 | 9,4 | 52 | 34 |  |
| 00110101 | 9,5 | 53 | 35 |  |
| 00110110 | 9,6 | 54 | 36 |  |
| 00110111 | 9,7 | 55 | 37 |  |
| 00111000 | 9.8 | 56 | 38 |  |
| 00111001 | 9,8,1 | 57 | 39 |  |
| 00111010 | 9,8,2 | 58 | 3 A |  |
| 00111011 | 9,8,3 | 59 | 3 B |  |
| 00111100 | 9,8,4 | 60 | 3 C |  |
| 00111101 | 9,8,5 | 61 | 3 D |  |
| 00111110 | 9,8,6 | 62 | 3E |  |
| 00111111 | 9,8,7 | 63 | 3 F |  |
| 01000000 |  | 64 | 40 | blank |
| 01000001 | 12,0,9,1 | 65 | 41 |  |
| 01000010 | 12,0,9,2 | 66 | 42 |  |
| 01000011 | 12,0,9,3 | 67 | 43 |  |
| 01000100 | 12,0,9.4 | 68 | 44 |  |
| 01000101 | 12,0,9,5 | 69 | 45 |  |
| 01000110 | 12,0,9,6 | 70 | 46 |  |
| 01000111 | 12,0,9,7 | 71 | 47 |  |
| 01001000 | 12,0,9,8 | 72 | 48 |  |
| 01001001 | 12,8,1 | 73 | 49 |  |
| 01001010 | 12,8,2 | 74 | 4A |  |
| 01001011 | 12,8,3 | 75 | 4 B | - (period) |
| 01001100 | 12,8,4 | 76 | 4 C | $<$ |
| 01001101 | 12,8,5 | 77 | 4 D | ( |
| 01001110 | 12,8,6 | 78 | 4 E | + |
| 01001111 | 12,8,7 | 79 | 4 F |  |
| 01010000 | 12 | 80 | 50 | $\varepsilon$ |
| 01010001 | 12,11,9.1 | 81 | 51 |  |
| 01010010 | 12,11,9,2 | 82 | 52 |  |
| 01010011 | 12,11,9,3 | 83 | 53 |  |
| 01010100 | 12,11.9.4 | 84 | 54 |  |
| 01010101 | 12,11,9,5 | 85 | 55 |  |
| 01010110 | 12,11,9,6 | 86 | 56 |  |
| 01010111 | 12,11,9.7 | 87 | 57 |  |
| 01011000 | 12,11,9,8 | 88 | 58 |  |
| 01011001 | 11,8,1 | 89 | 59 |  |
| 01011010 | 11,8,2 | 90 | 5A |  |
| 01011011 | 11,8,3 | 91 | 5 B | \$ |
| 01011100 | 11,8,4 | 92 | 5 C | * |
| 01011101 | 11,8,5 | 93 | 5D | ) |
| 01011110 | 11,8,6 | 94 | 5 E |  |
| 01011111 | 11,8,7 | 95 | 5 F |  |
| 01100000 | 11 | 96 | 60 | - |
| 01100001 | 0.1 | 97 | 61 | 1 |
| 01100010 | 11,0,9,2 | 98 | 62 |  |
| 01100011 | 11,0,9,3 | 99 | 63 |  |
| 01100100 | 11,0,9,4 | 100 | 64 |  |
| 01100101 | 11,0,9,5 | 101 | 65 |  |
| 01100110 | 11,0,9,6 | 102 | 66 |  |
| 01100111 | 11,0,9,7 | 103 | 67 |  |
| 01101000 | 11,0,9,8 | 104 | 68 |  |
| 01101001 | 0,8,1 | 105 | 69 |  |
| 01101010 | 12,11 | 106 | 6 A |  |
| 01101011 | 0,8,3 | 107 | 6B | , (comma) |


| 8-Bit <br> EBCDIC <br> code | Character Set Punch Combination | Decimal | HexaDecimal | printer Graphics |
| :---: | :---: | :---: | :---: | :---: |
| 01101100 | 0,8,4 | 108 | 6 C | \% |
| 01101101 | 0,8,5 | 109 | 6D |  |
| 01101110 | 0,8,6 | 110 | 6 E |  |
| 01101111 | 0,8,7 | 111 | 6 F |  |
| \| 01110000 | 12,11,0 | 112 | 70 |  |
| - 01110001 | 12,11,0,9.1 | 113 | 71 |  |
| - 01110010 | 12,11,0,9,2 | 114 | 72 |  |
| - 01110011 | 12,11,0,9,3 | 115 | 73 |  |
| - 01110100 | 12,11,0,9,4 | 116 | 74 |  |
| 01110101 | 12,11,0,9,5 | 117 | 75 |  |
| 01110110 | 12,11,0,9,6 | 118 | 76 |  |
| 01110111 | 12,11,0,9,7 | 119 | 77 |  |
| 01111000 | 12,11,0,9,8 | 120 | 78 |  |
| - 01111001 | 8,1 | 121 | 79 |  |
| - 01111010 | 8,2 | 122 | 7 A |  |
| 01111011 | 8,3 | 123 | 7 B | \# |
| 01111100 | 8.4 | 124 | 7 C | a |
| 01111101 | 8,5 | 125 | 7D | - (apostrophe) |
| - 01111110 | 8,6 | 126 | 7 E | $=$. |
| - 01111111 | 8,7 | 127 | 7 F |  |
| - 10000000 | 12,0,8,1 | 128 | 80 |  |
| - 10000001 | 12,0,1 | 129 | 81 |  |
| \| 10000010 | 12,0,2 | 130 | 82 |  |
| - 10000011 | 12.0,3 | 131 | 83 |  |
| - 10000100 | 12,0,4 | 132 | 84 |  |
| - 10000101 | 12,0,5 | 133 | 85 |  |
| - 10000110 | 12,0,6 | 134 | 86 |  |
| 10000111 | 12,0,7 | 135 | 87 |  |
| 10001000 | 12,0,8 | 136 | 88 |  |
| - 10001001 | 12,0,9 | 137 | 89 |  |
| 10001010 | 12,0,8,2 | 138 | 8A |  |
| 10001011 | 12,0,8,3 | 139 | 8B |  |
| - 10001100 | 12,0,8,4 | 140 | 8 C |  |
| 10001101 | 12,0,8,5 | 141 | 8 D |  |
| - 10001110 | 12,0,8,6 | 142 | 8 E |  |
| - 10001111 | 12,0,8,7 | 143 | 8 F |  |
| - 10010000 | 12,11,8,1 | 144 | 90 |  |
| 10010001 | 12,11.1 | 145 | 91 |  |
| 10010010 | 12,11,2 | 146 | 92 |  |
| 10010011 | 12,11, 3 | 147 | 93 |  |
| 10010100 | 12,11,4 | 148 | 94 |  |
| - 10010101 | 12.11.5 | 149 | 95 |  |
| - 10010110 | 12.11.6 | 150 | 96 |  |
| 10010111 | 12,11,7 | 151 | 97 |  |
| - 10011000 | 12,11,8 | 152 | 98 |  |
| 10011001 | 12,11,9 | 153 | 99 |  |
| - 10011010 | 12,11,8,2 | 154 | 9A |  |
| 10011011 | 12,11,8,3 | 155 | 9 B |  |
| - 10011100 | 12,11,8,4 | 156 | 9 C |  |
| \| 10011101 | 12,11,8,5 | 157 | 9 D |  |
| 10011110 | 12,11,8,6 | 158 | 9 E |  |
| - 10011111 | 12,11,8,7 | 159 | 9 F |  |
| 10100000 | 11,0,8,1 | 160 | A0 |  |
| 10100001 | 11,0,1 | 161 | A1 |  |
| 10100010 | 11,0,2 | 162 | A2 |  |
| 10100011 | 11,0,3 | 163 | A 3 |  |
| 10100100 | 11,0,4 | 164 | A4 |  |


| $\begin{gathered} 8-\text { Bit } \\ \text { EBCDIC } \\ \text { Code } \end{gathered}$ | Character Set Punch Combination | Decimal | Hexa- Decimal | Printer <br> Graphics |
| :---: | :---: | :---: | :---: | :---: |
| 10100101 | 11,0,5 | 165 | A5 |  |
| 10100110 | 11, 0, 6 | 166 | A 6 |  |
| 10100111 | 11,0,7 | 167 | A7 |  |
| 10101000 | 11,0,8 | 168 | A8 |  |
| 10101001 | 11,0,9 | 169 | A9 |  |
| 10101010 | 11,0,8,2 | 170 | AA |  |
| 10101011 | 11,0,8,3 | 171 | AB |  |
| 10101100 | $11,0,8,4$ | 172 | AC |  |
| 10101101 | 11,0,8,5 | 173 | AD |  |
| 10101110 | $11,0,8,6$ | 174 | AE |  |
| 10101111 | 11,0,8,7 | 175 | AF |  |
| 10110000 | 12,11,0,8,1 | 176 | B0 |  |
| 10110001 | 12,11,0,1 | 177 | B1 |  |
| 10110010 | 12,11,0,2 | 178 | B2 |  |
| 10110011 | 12,11,0,3 | 179 | B3 |  |
| 10110100 | 12,11,0,4 | 180 | B4 |  |
| 10110101 | 12,11,0,5 | 181 | B5 |  |
| 10110110 | $12,11,0,6$ | 182 | B6 |  |
| 10110111 | 12,11,0,7 | 183 | B7 |  |
| 10111000 | 12,11;0,8 | 184 | B8 |  |
| 10111001 | 12,11,0,9 | 185 | B9 |  |
| 10111010 | 12,11,0,8,2 | 186 | BA |  |
| 10111011 | 12,11,0,8,3 | 187 | BB |  |
| 10111100 | $12,11,0,8,4$ | 188 | BC |  |
| 10111101 | 12,11,0,8,5 | 189 | BD |  |
| 10111110 | 12,11,0,8,6 | 190 | BE |  |
| 10111111 | 12,11;0,8,7 | 191 | BF |  |
| 11000000 | 12.0 | 192 | C0 |  |
| 11000001 | 12,1 | 193 | C1 | A |
| 11000010 | 12,2 | 194 | C2 | B |
| 11000011 | 12.3 | 195 | C3 | C |
| 11000100 | 12,4 | 196 | C4 | D |
| 11000101 | 12,5 | 197 | C5 | E |
| 11000110 | 12,6 | 198 | C6 | F |
| 11000111 | 12,7 | 199 | C7 | G |
| 11001000 | 12.8 | 200 | C8 | H |
| 11001001 | 12,9 | 201 | C 9 | I |
| 11001010 | 12,0,9,8,2 | 202 | CA |  |
| 11001011 | 12,0,9,8,3 | 203 | CB |  |
| 11001100 | 12,0,9,8,4 | 204 | CC |  |
| 11001101 | 12,0,9,8,5 | 205 | CD |  |
| 11001110 | $12,0,9,8,6$ | 206 | CE | : |
| 11001111 | 12,0,9,8.7 | 207 | CF |  |
| 11010000 | 11.0 | 208 | D0 |  |
| 11010001 | 11,1 | 209 | D1 | J |
| 11010010 | 11. 2 | 210 | D2 | K |
| 11010011 | 11.3 | 211 | D3 | L |
| 11010100 | 11,4 | 212 | D4 | M |
| 11010101 | 11,5 | 213 | D5 | N |
| 11010110 | 11.6 | 214 | D6 | 0 |
| 11010111 | 11.7 | 215 | D7 | p |
| 11011000 | 11.8 | 216 | D8 | Q |
| 11011001 | 11,9 | 217 | D9 | R |
| 11011010 | 12,11.9,8,2 | 218 | DA |  |
| 11011011 | 12,11,9,8,3 | 219 | DB |  |
| 11011100 | 12,11,9,8,4 | 220 | DC |  |
| 11011101 | $12,11,9,8,5$ | 221 | DD |  |


| 8-Bit | Character Set |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| EBCDIC | Punch |  | Hexa- | Printer |
| Code | Combination | Decimal | Decimal | Graphics |
| 11011110 | 12,11,9,8,6 | 222 | DE |  |
| 11011111 | 12,11,9,8,7 | 223 | DF |  |
| 11100000 | 0,8,2 | 224 | E0 |  |
| 11100001 | 11,0,9,1 | 225 | E1 |  |
| 11100010 | 0,2 | 226 | E2 | S |
| 11100011 | 0,3 | 227 | E3 | T |
| 11100100 | 0.4 | 228 | E4 | U |
| 11100101 | 0.5 | 229 | E5 | V |
| 11100110 | 0,6 | 230 | E6 | W |
| 11100111 | 0.7 | 231 | E7 | X |
| 11101000 | 0.8 | 232 | E8 | Y |
| 11101001 | 0,9 | 233 | E9 | 2 |
| 11101010 | 11,0,9,8,2 | 234 | EA |  |
| 11101011 | 11,0,9,8,3 | 235 | EB |  |
| 11101100 | 11,0,9,8,4 | 236 | EC |  |
| 11101101 | 11,0,9,8,5 | 237 | ED |  |
| 11101110 | 11, 0, 9,8,6 | 238 | EE |  |
| 11101111 | 11,0,9,8,7 | 239 | EF |  |
| 11110000 | 0 | 240 | F0 | 0 |
| 11110001 | 1 | 241 | F1 | 1 |
| 11110010 | 2 | 242 | F2 | 2 |
| 11110011 | 3 | 243 | F3 | 3 |
| 11110100 | 4 | 244 | F4 | 4 |
| 11110101 | 5 | 245 | F5 | 5 |
| 11110110 | 6 | 246 | F6 | 6 |
| 11110111 | 7 | 247 | F7 | 7 |
| 11111000 | 8 | 248 | F8 | 8 |
| 11111001 | 9 | 249 | F9 | 9 |
| 11111010 | 12,11,0,9,8,2 | 250 | FA |  |
| 11111011 | 12,11, 0,9,8,3 | 251 | FB |  |
| 11111100 | 12,11,0,9,8,4 | 252 | FC |  |
| 11111101 | $12,11,0,9,8,5$ | 253 | FD |  |
| 11111110 | 12,11,0,9,8,6 | 254 | FE |  |
| 11111111 | 12,11,0,9,8,7 | 255 | FF |  |



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## Appendix II: Hexadecimal-Decimal Conversion Table

The table in this appendix provides for direct conversion of decimal and hexadecimal numbers in these ranges:


Decimal numbers (0000-4095) are given within the 5-part table. The first two characters (high-order) of hexadecimal numbers ( $000-\mathrm{FFF}$ ) are given in the lefthand column of the table; the third character (x) is arranged across the top of each part of the table.

```
    To find the decimal equivalent of the hexadecimal number 0C9, look for OC in the left
colum, and across that row under the column for x = 9. The decimal number is 0201.
To convert from decimal to hexadecimal, look up the decimal number within the table and read the hexadecimal number by a combination of the hex characters in the left column, and the value for \(x\) at the top of the column containing the decimal number. For example, the decimal number 123 has the hexadecimal equivalent of 07B; the decimal number 1478 has the hexadecimal equivalent of 5C6.
For numbers outside the range of the table, add the following values to the table
```

Henadecimal

|  | $\mathbf{x}=0$ | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | A | B | C | D | E | F |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 00x | 0000 | 0001 | 0002 | 0003 | 0004 | 0005 | 0006 | 0007 | 0008 | 0009 | 0010 | 0011 | 0012 | 0013 | 0014 | 0015 |
| 01x | 0016 | 0017 | 0018 | 0019 | 0020 | 0021 | 0022 | 0023 | 0024 | 0025 | 0026 | 0027 | 0028 | 0029 | 0030 | 0031 |
| 02x | 0032 | 0033 | 0034 | 0035 | 0036 | 0037 | 0038 | 0039 | 0040 | 0041 | 0042 | 0043 | 0044 | 0045 | 0046 | 0047 |
| 03x | 0048 | 0049 | 0050 | 0051 | 0052 | 0053 | 0054 | 0055 | 0056 | 0057 | 0058 | 0059 | 0060 | 0061 | 0062 | 00\%. |
| 04x | 0064 | 0065 | 0066 | 0067 | 0068 | 0069 | 0070 | 0071 | 0072 | 0073 | 0074 | 0075 | 0076 | 0077 | 0078 | 0074 |
| 05x | 0080 | 0081 | 0082 | 0083 | 0084 | 0085 | 0086 | 0087 | 0088 | 0089 | 0090 | 0091 | 0092 | 0093 | 0094 | 0095 |
| 06x | 0096 | 0097 | 0098 | 0099 | 0100 | 0101 | 0102 | 0103 | 0104 | 0105 | 0106 | 0107 | 0108 | 0109 | 0110 | 0111 |
| 07x | 0112 | 0113 | 0114 | 0115 | 0116 | 0117 | 0118 | 0119 | 0120 | 0121 | 0122 | 0123 | 0124 | 0125 | 0126 | 0127 |
| 08x | 0128 | 0129 | 0130 | 0131 | 0132 | 0133 | 0134 | 0135 | 0136 | 0137 | 0138 | 0139 | 0140 | 0141 | 0142 | 0143 |
| 09x | 0144 | 0145 | 0146 | 0147 | 0148 | 0149 | 0150 | 0151 | 0152 | 0153 | 0154 | 0155 | 0156 | 0157 | 0158 | 0153 |
| 0Ax | 0160 | 0161 | 0162 | 0163 | 0164 | 0165 | 0166 | 0167 | 0168 | 0169 | 0170 | 0171 | 0172 | 0173 | 0174 | 0175 |
| OBx | 0176 | 0177 | 0178 | 0179 | 0180 | 0181 | 0182 | 0183 | 0184 | 0185 | 0186 | 0187 | 0188 | 0189 | 0190 | 0191 |
| 0 Cx | 0192 | 0193 | 0194 | 0195 | 0196 | 0197 | 0198 | 0199 | 0200 | 0201 | 0202 | 0203 | 0204 | 0205 | 0206 | 0207 |
| 0Dx | 0208 | 0209 | 0210 | 0211 | 0212 | 0213 | 0214 | 0215 | 0216 | 0217 | 0218 | 0219 | 0220 | 0221 | 0222 | 0223 |
| OEx | 0224 | 0225 | 0226 | 0227 | 0228 | 0229 | 0230 | 0231 | 0232 | 0233 | 0234 | 0235 | 0236 | 0237 | 0238 | 0239 |
| OFx | 0240 | 0241 | 0242 | 0243 | 0244 | 0245 | 0246 | 0247 | 0248 | 0249 | 0250 | 0251 | 0252 | 0253 | 0254 | 0255 |
| 10x | 0256 | 0257 | 0258 | 0259 | 0260 | 0261 | 0262 | 0263 | 0264 | 0265 | 0266 | 0267 | 0268 | 0269 | 0270 | 0271 |
| 11x | 0272 | 0273 | 0274 | 0275 | 0276 | 0277 | 0278 | 0279 | 0280 | 0281 | 0282 | 0283 | 0284 | 0285 | 0286 | 0287 |
| 12x | 0288 | 0289 | 0290 | 0291 | 0292 | 0293 | 0294 | 0295 | 0296 | 0297 | 0298 | 0299 | 0300 | 0301 | 0302 | 0303 |
| 13x | 0304 | 0305 | 0306 | 0307 | 0308 | 0309 | 0310 | 0311 | 0312 | 0313 | 0314 | 0315 | 0316 | 0317 | 0318 | 0319 |
| 14 x | 0320 | 0321 | 0322 | 0323 | 0324 | 0325 | 0326 | 0327 | 0328 | 0329 | 0330 | 0331 | 0332 | 0333 | 0334 | 0335 |
| 15x | 0336 | 0337 | 0338 | 0339 | 0340 | 0341 | 0342 | 0343 | 0344 | 0345 | 0346 | 0347 | 0348 | 0349 | 0350 | 0351 |
| 16x | 0352 | 0353 | 0354 | 0355 | 0356 | 0357 | 0358 | 0359 | 0360 | 0361 | 0362 | 0363 | 0364 | 0365 | 0366 | 0367 |
| 17x | 0368 | 0369 | 0370 | 0371 | 0372 | 0373 | 0374 | 0375 | 0376 | 0377 | 0378 | 0379 | 0380 | 0381 | 0382 | 0383 |
| 18x | 0384 | 0385 | 0386 | 0387 | 0388 | 0389 | 0390 | 0391 | 0392 | 0393 | 0394 | 0395 | 0396 | 0397 | 0398 | 0399 |
| 19x | 0400 | 0401 | 0402 | 0403 | 0404 | 0405 | 0406 | 0407 | 0408 | 0409 | 0410 | 0411 | 0412 | 0413 | 0414 | 0415 |
| 1Ax | 0416 | 0417 | 0418 | 0419 | 0420 | 0421 | 0422 | 0423 | 0424 | 0425 | 0426 | 0427 | 0428 | 0429 | 0430 | 0431 |
| 1Bx | 0432 | 0433 | 0434 | 0435 | 0436 | 0437 | 0438 | 0439 | 0440 | 0441 | 0442 | 0443 | 0444 | 0445 | 0446 | 0447 |
| 1Cx | 0448 | 0449 | 0450 | 0451 | 0452 | 0453 | 0454 | 0455 | 0456 | 0457 | 0458 | 0459 | 0460 | 0461 | 0462 | 046? |
| 1Dx | 0464 | 0465 | 0466 | 0467 | 0468 | 0469 | 0470 | 0471 | 9472 | 0473 | 0474 | 0475 | 0476 | 0477 | 0478 | 0479 |
| 1Ex | 0480 | 0481 | 0482 | 0483 | 0484 | 0485 | 0486 | 0487 | 0488 | 0489 | 0490 | 0491 | 0492 | 0493 | 0494 | 0495 |
| 1Fx | 0496 | 0497 | 0498 | 0499 | 0500 | 0501 | 0502 | 0503 | 0504 | 0505 | 0506 | 0507 | 0508 | 0509 | 0510 | 0511 |
| 20x | 0512 | 0513 | 0514 | 0515 | 0516 | 0517 | 0518 | 0519 | 0520 | 0521 | 0522 | 0523 | 0524 | 0525 | 0526 | 0527 |
| 21x | 0528 | 0529 | 0530 | 0531 | 0532 | 0533 | 0534 | 0535 | 0536 | 0537 | 0538 | 0539 | 0540 | 0541 | 0542 | 0543 |
| 22x | 0544 | 0545 | 0546 | 0547 | 0548 | 0549 | 0550 | 0551 | 0552 | 0553 | 0554 | 0555 | 0556 | 0557 | 0558 | 0559 |
| 23x | 0560 | 0561 | 0562 | 0563 | 0564 | 0565 | 0566 | 0567 | 0568 | 0569 | 0570 | 0571 | 0572 | 0573 | 0574 | 0575 |
| 24x | 0576 | 0577 | 0578 | 0579 | 0580 | 0581 | 0582 | 0583 | 0584 | 0585 | 0586 | 0587 | 0588 | 0589 | 0590 | 0591 |
| 25x | 0592 | 0593 | 0594 | 0595 | 0596 | 0597 | 0598 | 0599 | 0600 | 0601 | 0602 | 0603 | 0604 | 0605 | 0606 | 0607 |
| 26x | 0608 | 0609 | 0610 | 0611 | 0612 | 0613 | 0614 | 0615 | 0616 | 0617 | 0618 | 0619 | 0620 | 0621 | 0622 | 0623 |
| 27x | 0624 | 0625 | 0626 | 0627 | 0628 | 0629 | 0630 | 0631 | 0632 | 0633 | 0634 | 0635 | 0636 | 0637 | 0638 | 0639 |
| 28x | 0640 | 0641 | 0642 | 0643 | 0644 | 0645 | 0646 | 0647 | 0648 | 0649 | 0650 | 0651 | 0652 | 0653 | 0654 | 0655 |
| 29x | 0656 | 0657 | 0658 | 0659 | 0660 | 0661 | 0662 | 0663 | 0664 | 0665 | 0666 | 0667 | 0668 | 0669 | 0670 | 0671 |
| 2Ax | 0672 | 0673 | 0674 | 0675 | 0676 | 0677 | 0678 | 0679 | 0680 | 0681 | 0682 | 0683 | 0684 | 0685 | 0686 | 0687 |
| 2Bx | 0688 | 0689 | 0690 | 0691 | 0692 | 0693 | 0694 | 0695 | 0696 | 0697 | 0698 | 0699 | 0700 | 0701 | 0702 | 0703 |
| 2Cx | 0704 | 0705 | 0706 | 0707 | 0708 | 0709 | 0710 | 0711 | 0712 | 0713 | 0714 | 0715 | 0716 | 0717 | 0718 | 0719 |
| 2Dx | 0720 | 0721 | 0722 | 0723 | 0724 | 0725 | 0726 | 0727 | 0728 | 0729 | 0730 | 0731 | 0732 | 0733 | 0734 | 0735 |
| 2Ex | 0736 | 0737 | 0738 | 0739 | 0740 | 0741 | 0742 | 0743 | 0744 | 0745 | 0746 | 0747 | 0748 | 0749 | 0750 | 0751 |
| 2Fx | 0752 | 0753 | 0754 | 0755 | 0756 | 0757 | 0758 | 0759 | 0760 | 0761 | 0762 | 0763 | 0764 | 0765 | 0766 | 0767 |
| 30x | 0768 | 0769 | 0770 | 0771 | 0772 | 0773 | 0774 | 0775 | 0776 | 0777 | 0778 | 0779 | 0780 | 0781 | 0782 | 0783 |
| 31 x | 0784 | 0785 | 0786 | 0787 | 0788 | 0789 | 0790 | 0791 | 0792 | 0793 | 0794 | 0795 | 0796 | 0797 | 0798 | 0799 |
| 32x | 0800 | 0801 | 0802 | 0803 | 0804 | 0805 | 0806 | 0807 | 0808 | 0809 | 0810 | 0811 | 0812 | 0813 | 0814 | 0815 |
| 33x | 0816 | 0817 | 0818 | 0819 | 0820 | 0821 | 0822 | 0823 | 0824 | 0825 | 0826 | 0827 | 0828 | 0829 | 0830 | 0831 |
| 34 x | 0832 | 0833 | 0834 | 0835 | 0836 | 0837 | 0838 | 0839 | 0840 | 0841 | 0842 | 0843 | 0844 | 0845 | 0846 | 0847 |
| $35 x$ | 0848 | 0849 | 0850 | 0851 | 0852 | 0853 | 0854 | 0855 | 0856 | 0857 | 0858 | 0859 | 0860 | 0861 | 0862 | 0863 |
| $36 x$ | 0864 | 0865 | 0866 | 0867 | 0868 | 0869 | 0870 | 0871 | 0872 | 0873 | 0874 | 0875 | 0876 | 0877 | 0878 | 0879 |
| 37 x | 0880 | 0881 | 0882 | 0883 | 0884 | 0885 | 0886 | 0887 | 0888 | 0889 | 0890 | 0891 | 0892 | 0893 | 0894 | 0895 |
| 38x | 0896 | 0897 | 0898 | 0899 | 0900 | 0901 | 0902 | 0903 | 0904 | 0905 | 0906 | 0907 | 0908 | 0909 | 0910 | 0911 |
| 39x | 0912 | 0913 | 0914 | 0915 | 0916 | 0917 | 0918 | 0919 | 0920 | 0921 | 0922 | 0923 | 0924 | 0925 | 0926 | 0927 |
| 3Ax | 0928 | 0929 | 0930 | 0931 | 0932 | 0933 | 0934 | 0935 | 0936 | 0937 | 0938 | 0939 | 0940 | 0941 | 0942 | 0943 |
| 3 Bx | 0944 | 0945 | 0946 | 0947 | 0948 | 0949 | 0950 | 0951 | 0952 | 0953 | 0954 | 0955 | 0956 | 0957 | 0958 | 0959 |
| 3 Cx | 0960 | 0961 | 0962 | 0963 | 0964 | 0965 | $0966$ | $0967$ | $0968$ | 0969 | 0970 | 0971 | 0972 | 0973 | 0974 | 0975 |
| 3Dx | 0976 | 0977 | 0978 | 0979 | 0980 | 0981 | 0982 | 0983 | 0984 | 0985 | 0986 | 0987 | 0988 | 0989 | 0990 | 0991 |
| 3Ex | 0992 | 0993 | 0994 | 0995 | 0996 | 0997 | 0998 | 0999 | 1000 | 1001 | 1002 | 1003 | 1004 | 1005 | 1006 | 1007 |
| 3Fx | 1008 | 1009 | 1010 | 1011 | 1012 | 1013 | 1014 | 1015 | 1016 | 1017 | 1018 | 1019 | 1020 | 1021 | 1022 | 1023 |


|  | $x=0$ | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | A | B | C | D | E | F |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 40x | 1024 | 1025 | 1026 | 1027 | 1028 | 1029 | 1030 | 1031 | 1032 | 1033 | 1034 | 1035 | 1036 | 1037 | 1038 | 1039 |
| $41 x$ | 1040 | 1041 | 1042 | 1043 | 1044 | 1045 | 1046 | 1047 | 1048 | 1049 | 1050 | 1051 | 1052 | 1053 | 1054 | 1055 |
| 42 x | 1056 | 1057 | 1058 | 1059 | 1060 | 1061 | 1062 | 1063 | 1064 | 1065 | 1066 | 1067 | 1068 | 1069 | 1070 | 1071 |
| 43 x | 1072 | 1073 | 1074 | 1075 | 1076 | 1077 | 1078 | 1079 | 1080 | 1081 | 1082 | 1083 | 1084 | 1085 | 1086 | 1087 |
| $44 x$ | 1088 | 1089 | 1090 | 1091 | 1092 | 1093 | 1094 | 1095 | 1096 | 1097 | 1098 | 1099 | 1100 | 1101 | 1102 | 1103 |
| 45x | 1104 | 1105 | 1106 | 1107 | 1108 | 1109 | 1110 | 1111 | 1112 | 1113 | 1114 | 1115 | 1116 | 1117 | 1118 | 1119 |
| 46x | 1120 | . 1121 | 1122 | 1123 | 1124 | 1125 | 1126 | 1127 | 1128 | 1129 | 1130 | 1131 | 1132 | 1133 | 1134 | 1135 |
| 47 x | 1136 | 1137 | 1138 | 1139 | 1140 | 1141 | 1142 | 1143 | 1144 | 1145 | 1146 | 1147 | 1148 | 1149 | 1150 | 1151 |
| 48x | 1152 | 1153 | 1154 | 1155 | 1156 | 1157 | 1158 | 1159 | 1160 | 1161 | 1162 | 1163 | 1164 | 1165 | 1166 | 1167 |
| 49x | 1168 | 1169 | 1170 | 1171 | 1172 | 1173 | 1174 | 1175 | 1176 | 1177 | 1178 | 1179 | 1180 | 1181 | 1182 | 1183 |
| 4Ax | 1184 | 1185 | 1186 | 1187 | 1188 | 1189 | 1190 | 1191 | 1192 | 1193 | 1194 | 1195 | 1196 | 1197 | 1198 | 1199 |
| 4Bx | 1200 | 1201 | 1202 | 1203 | 1204 | 1205 | 1206 | 1207 | 1208 | 1209 | 1210 | 1211 | 1212 | 1213 | 1214 | 1215 |
| 4Cx | 1216 | 1217 | 1218 | 1219 | 1220 | 1221 | 1222 | 1223 | 1224 | 1225 | 1226 | 1227 | 1228 | 1229 | 1230 | 1231 |
| 4Dx | 1232 | 1233 | 1234 | 1235 | 1236 | 1237 | 1238 | 1239 | 1240 | 1241 | 1242 | 1243 | 1244 | 1245 | 1246 | 1247 |
| 4Ex | 1248 | 1249 | 1250 | 1251 | 1252 | 1253 | 1254 | 1255 | 1256 | 1257 | 1258 | 1259 | 1260 | 1261 | 1262 | 1263 |
| 4Fx | 1264 | 1265 | 1266 | 1267 | 1268 | 1269 | 1270 | 1271 | 1272 | 127.3 | 1274 | 1275 | 1276 | 1277 | 1278 | 1279 |
| 50x | 1280 | 1281 | 1282 | 1283 | 1284 | 1285 | 1286 | 1287 | 1288 | 1289 | 1290 | 1291 | 1292 | 1293 | 1294 | 1295 |
| 51x | 1296 | 1297 | 1298 | 1299 | 1300 | 1301 | 1302 | 1303 | 1304 | 1305 | 1306 | 1307 | 1308 | 1309 | 1310 | 1311 |
| 52x | 1312 | 1313 | 1314 | 1315 | 1316 | 1317 | 1318 | 1319 | 1320 | 1321 | 1322 | 1323 | 1324 | 1325 | 1326 | 1327 |
| 53x | 1328 | 1329 | 1330 | 1331 | 1332 | 1333 | 1334 | 1335 | 1336 | 1337 | 1338 | 1339 | 1340 | 1341 | 1342 | 1343 |
| 54 x | 1344 | 1345 | 1346 | 1347 | 1348 | 1349 | 1350 | 1351 | 1352 | 1353 | 1354 | 1355 | 1356 | 1357 | 1358 | 1359 |
| 55x | 1360 | 1361 | 1362 | 1363 | 1364 | 1365 | 1366 | 1367 | 1368 | 1369 | 1370 | 1371 | 1372 | 1373 | 1374 | 1375 |
| 56x | 1376 | 1377 | 1378 | 1379 | 1380 | 1381 | 1382 | 1383 | 1384 | 1385 | 1386 | 1387 | 1388 | 1389 | 1390 | 1391 |
| 57x | 1392 | 1393 | 1394 | 1395 | 1396 | 1397 | 1398 | 1399 | 1400 | 1401 | 1402 | 1403 | 1404 | 1405 | 1406 | 1407 |
| 58x | 1408 | 1409 | 1410 | 1411 | 1412 | 1413 | 1414 | 1415 | 1416 | 1417 | 1418 | 1419 | 1420 | 1421 | 1422 | 1423 |
| 59x | 1424 | 1425 | 1426 | 1427 | 1428 | 1429 | 1430 | 1431 | 1432 | 1433 | 1434 | 1435 | 1436 | 1437 | 1438 | 1439 |
| 5Ax | 1440 | 1441 | 1442 | 1443 | 1444 | 1445 | 1446 | 1447 | 1448 | 1449 | 1450 | 1451 | 1452 | 1453 | 1454 | 1455 |
| 5Bx | 1456 | 1457 | 1458 | 1459 | 1460 | 1461 | 1462 | 1463 | 1464 | 1465 | 1466 | 1467 | 1468 | 1469 | 1470 | 1471 |
| 5Cx | 1472 | 1473 | 1474 | 1475 | 1476 | 1477 | 1478 | 1479 | 1480 | 1481 | 1482 | 1483 | 1484 | 1485 | 1486 | 1487 |
| 5Dx | 1488 | 1489 | 1490 | 1491 | 1492 | 1493 | 1494 | 1495 | 1496 | 1497 | 1498 | 1499 | 1500 | 1501 | 1502 | 1503 |
| 5Ex | 1504 | 1505 | 1506 | 1507 | 1508 | 1509 | 1510 | 1511 | 1512 | 1513 | 1514 | 1515 | 1516 | 1517 | 1518 | 1519 |
| 5Fx | 1520 | 1521 | 1522 | 1523 | 1524 | 1525 | 1526 | 1527 | 1528 | 1529 | 1530 | 1531 | 1532 | 1533 | 1534 | 1535 |
| 60x | 1536 | 1537 | 1538 | 1539 | 1540 | 1541 | 1542 | 1543 | 1544 | 1545 | 1546 | 1547 | 1548 | 1549 | 1550 | 1551 |
| 61x | 1552 | 1553 | 1554 | 1555 | 1556 | 1557 | 1558 | 1559 | 1560 | 1561 | 1562 | 1563 | 1564 | 1565 | 1566 | 1567 |
| 62x | 1568 | 1569 | 1570 | 1571 | 1572 | 1573 | 1574 | 1575 | 1576 | 1577 | 1578 | 1579 | 1580 | 1581 | 1582 | 1583 |
| 63x | 1584 | 1585 | 1586 | 1587 | 1588 | 1589 | 1590 | 1591 | 1592 | 1593 | 1594 | 1595 | 1596 | 1597 | 1598 | 1599 |
| 64 x | 1600 | 1601 | 1602 | 1603 | 1604 | 1605 | 1606 | 1607 | 1608 | 1609 | 1610 | 1611 | 1612 | 1613 | 1614 | 1615 |
| $65 \times$ | 1616 | 1617 | 1618 | 1619 | 1620 | 1621 | 1622 | 1623 | 1624 | 1625 | 1626 | 1627 | 1628 | 1629 | 1630 | 1631 |
| 66x | 1632 | 1633 | 1634 | 1635 | 1636 | 1637 | 1638 | 1639 | 1640 | 1641 | 1642 | 1643 | 1644 | 1645 | 1646 | 1647 |
| $67 \times$ | 1648 | 1649 | 1650 | 1651 | 1652 | 1653 | 1654 | 1655 | 1656 | 1657 | 1658 | 1659 | 1660 | 1661 | 1662 | 1663 |
| 68x | 1664 | 1665 | 1666 | 1667 | 1668 | 1669 | 1670 | 1671 | 1672 | 1673 | 1674 | 1675 | 1676 | 1677 | . 1678 | 1679 |
| 69x | 1680 | 1681 | 1682 | 1683 | 1684 | 1685 | 1686 | 1687 | 1688 | 1689 | 1690 | 1691 | 1692 | 1693 | 1694 | 1695 |
| 6Ax | 1696 | 1697 | 1698 | 1699 | 1700 | 1701 | 1702 | 1703 | 1704 | 1705 | 1706 | 1707 | 1708 | 1709 | 1710 | 1711 |
| 6Bx | 1712 | 1713 | 1714 | 1715 | 1716 | 1717 | 1718 | 1719 | 1720 | 1721 | 1722 | 1723 | 1724 | 1725 | 1726 | 1727 |
| 6 Cx | 1728 | 1729 | 1730 | 1731 | 1732 | 1733 | 1734 | 1735 | 1736 | 1737 | 1738 | 1739 | 1740 | 1741 | 1742 | 1743 |
| 6Dx | 1744 | 1745 | 1746 | 1747 | 1748 | 1749 | 1750 | 1751 | 1752 | 1753 | 1754 | 1755 | 1756 | 1757 | 1758 | 1759 |
| 6Ex | 1760 | 1761 | 1762 | 1763 | 1764 | 1765 | 1766 | 1767 | 1768 | 1769 | 1770 | 1771 | 1772 | 1773 | 1774 | 1775 |
| 6Fx | 1776 | 1777 | 1778 | 1779 | 1780 | 1781 | 1782 | 1783 | 1784 | 1785 | 1786 | 1787 | 1788 | 1789 | 1790 | 1791 |
| 70x | 1792 | 1793 | 1794 | 1795 | 1796 | 1797 | 1798 | 1799 | 1800 | 1801 | 1802 | 1803 | 1804 | 1805 | 1806 | 1807 |
| 71x | 1808 | 1809 | 1810 | 1811 | 1812 | 1813 | 1814 | 1815 | 1816 | 1817 | 1818 | 1819 | 1820 | 1821 | 1822 | 1823 |
| 72x | 1824 | 1825 | 1826 | 1827 | 1828 | 1829 | 1830 | 1831 | 1832 | 1833 | 1834 | 1835 | 1836 | 1837 | 1838 | 1839 |
| 73x | 1840 | 1841 | 1842 | 1843 | 1844 | 1845 | 1846 | 1847 | 1848 | 1849 | 1850 | 1851 | 1852 | 1853 | 1854 | 1855 |
| 74 x | 1856 | 1857 | 1858 | 1859 | 1860 | 1861 | 1862 | 1863 | 1864 | 1865 | 1866 | 1867 | 1868 | 1869 | 1870 | 1871 |
| 75x | 1872 | 1873 | 1874 | 1875 | 1876 | 1877 | 1878 | 1879 | 1880 | 1881 | 1882 | 1883 | 1884 | 1885 | 1886 | 1887 |
| 76x | 1888 | 1889 | 1890 | 1891 | 1892 | 1893 | 1894 | 1895 | 1896 | 1897 | 1898 | 1899 | 1900 | 1901 | 1902 | 1903 |
| 77x | 1904 | 1905 | 1906 | 1907 | 1908 | 1909 | 1910 | 1911 | 1912 | 1913 | 1914 | 1915 | 1916 | 1917 | 1918 | 1919 |
| 78x | 1920 | 1921 | 1922 | 1923 | 1924 | 1925 | 1926 | 1927 | 1928 | 1929 | 1930 | 1931 | 1932 | 1933 | 1934 | 1935 |
| 79x | 1936 | 1937 | 1938 | 1939 | 1940 | 1941 | 1942 | 1943 | 1944 | 1945 | 1946 | 1947 | 1948 | 1949 | 1950 | 1951 |
| 7 Ax | 1952 | 1953 | 1954 | 1955 | 1956 | 1957 | 1958 | 1959 | 1960 | 1961 | 1962 | 1963 | 1964 | 1965 | 1966 | 1967 |
| 7Bx | 1968 | 1969 | 1970 | 1971 | 1972 | 1973 | 1974 | 1975 | 1976 | 1977 | 1978 | 1979 | 1980 | 198 ! | 1982 | 1983 |
| 7Cx | 1984 | 1985 | 1986 | 1987 | 1988 | 1989 | 1990 | 1991 | 1992 | 1993 | 1994 | 1995 | 1996 | 1997 | 1998 | 1999 |
| 7Dx | 2000 | 2001 | 2002 | 2003 | 2004 | 2005 | 2006 | 2007 | 2008 | 2009 | 2010 | 2011 | 2012 | 2013 | 2014 | 2015 |
| 7Ex | 2016 | 2017 | 2018 | 2019 | 2020 | 2021 | 2022 | 2023 | 2024 | 2025 | 2026 | 2027 | 2028 | 2029 | 2030 | 2031 |
| 7Fx | 2032 | 2033 | 2034 | 2035 | 2036 | 2037 | 2038 | 2039 | 2040 | 2041 | 2042 | 2043 | 2044 | 2045 | 2046 | 2047 |


|  | $x=0$ | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | A | B | C | D | E | $F$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 80x | 2048 | 2049 | 2050 | 2051 | 2052 | 2053 | 2054 | 2055 | 2056 | 2057 | 2058 | 2059 | 2060 | 2061 | 2062 | 2063 |
| $81 \times$ | 2064 | 2065 | 2066 | 2067 | 2068 | 2069 | 2070 | 2071 | 2072 | 2073 | 2074 | 2075 | 2076 | 2077 | 2078 | 2079 |
| 82x | 2080 | 2081 | 2082 | 2083 | 2084 | 2085 | 2086 | 2087 | 2088 | 2089 | 2090 | 2091 | 2092 | 2093 | 2094 | 2095 |
| 83x | 2096 | 2097 | 2098 | 2099 | 2100 | 2101 | 2102 | 2103 | 2104 | 2105 | 2106 | 2107 | 2108 | 2109 | 2110 | 2111 |
| $84 \times$ | 2112 | 2113 | 2114 | 2115 | 2116 | 2117 | 2118 | 2119 | 2120 | 2121 | 2122 | 2123 | 2124 | 2125 | 2126 | 2127 |
| 85 x | 2128 | 2129 | 2130 | 2131 | 2132 | 2133 | 2134 | 2135 | 2136 | 2137 | 2138 | 2139 | 2140 | 2141 | 2142 | 2143 |
| 86x | 2144 | 2145 | 2146 | 2147 | 2148 | 2149 | 2150 | 2151 | 2152 | 2153 | 2154 | 2155 | 2156 | 2157 | 2158 | 2159 |
| 87x | 2160 | 2161 | 2162 | 2163 | 2164 | 2165 | 2166 | 2167 | 2168 | 2169 | 2170 | 2171 | 2172 | 2173 | 2174 | 2175 |
| 88x | 2176 | 2177 | 2178 | 2179 | 2180 | 2181 | 2182 | 2183 | 2184 | 2185 | 2186 | 2187 | 2188 | 2189 | 2190 | 2191 |
| 89x | 2192 | 2193 | 2194 | 2195 | 2196 | 2197 | 2198 | 2199 | 2200 | 2201 | 2202 | 2203 | 2204 | 2205 | 2206 | 2207 |
| 8Ax | 2208 | 2209 | 2210 | 2211 | 2212 | 2213 | 2214 | 2215 | 2216 | 2217 | 2218 | 2219 | 2220 | 2221 | 2222 | 2223 |
| 8 Bx | 2224 | 2225 | 2226 | 2227 | 2228 | 2229 | 2230 | 2231 | 2232 | 2233 | 2234 | 2235 | 2236 | 2237 | 2238 | 2239 |
| 8 Cx | 2240 | 2241 | 2242 | 2243 | 2244 | 2245 | 2246 | 2247 | 2248 | 2249 | 2250 | 2251 | 2252 | 2253 | 2254 | 2255 |
| 8Dx | 2256 | 2257 | 2258 | 2259 | 2260 | 2261 | 2262 | 2263 | 2264 | 2265 | 2266 | 2267 | 2268 | 2269 | 2270 | 2271 |
| 8Ex | 2272 | 2273 | 2274 | 2275 | 2276 | 2277 | 2278 | 2279 | 2280 | 2281 | 2282 | 2283 | 2284 | 2285 | 2286 | 2287 |
| 8Fx | 2288 | 2289 | 2290 | 2291 | 2292 | 2293 | 2294 | 2295 | 2296 | 2297 | 2298 | 2299 | 2300 | 2301 | 2302 | 2303 |
| 90x | 2304 | 2305 | 2306 | 2307 | 2308 | 2309 | 2310 | 2311 | 2312 | 2313 | 2314 | 2315 | 2316 | 2317 | 2318 | 2319 |
| 91x | 2320 | 2321 | 2322 | 2323 | 2324 | 2325 | 2326 | 2327 | 2328 | 2329 | 2330 | 2331 | 2332 | 2333 | 2334 | 2335 |
| 92x | 2336 | 2337 | 2338 | 2339 | 2340 | 2341 | 2342 | 2343 | 2344 | 2345 | 2346 | 2347 | 2348 | 2349 | 2350 | 2351 |
| 93x | 2352 | 2353 | 2354 | 2355 | 2356 | 2357 | 2358 | 2359 | 2360 | 2361 | 2362 | 2363 | 2364 | 2365 | 2366 | 2367 |
| 94x | 2368 | 2369 | 2370 | 2371 | 2372 | 2373 | 2374 | 2375 | 2376 | 2377 | 2378 | 2379 | 2380 | 2381 | 2382 | 2383 |
| 95x | 2384 | 2385 | 2386 | 2387 | 2388 | 2389 | 2390 | 2391 | 2392 | 2393 | 2394 | 2395 | 2396 | 2397 | 2398 | 2399 |
| 96x | 2400 | 2401 | 2402 | 2403 | 2404 | 2405 | 2406 | 2407 | 2408 | 2409 | 2410 | 2411 | 2412 | 2413 | 2414 | 2415 |
| 97x | 2416 | 2417 | 2418 | 2419 | 2420 | 2421 | 2422 | 2423 | 2424 | 2425 | 2426 | 2427 | 2428 | 2429 | 2430 | 2431 |
| 98x | 2432 | 2433 | 2434 | 2435 | 2436 | 2437 | 2438 | 2439 | 2440 | 2441 | 2442 | 2443 | 2444 | 2445 | 2446 | 2447 |
| 99x | 2448 | 2449 | 2450 | 2451 | 2452 | 2453 | 2454 | 2455 | 2456 | 2457 | 2458 | 2459 | 2460 | 2461 | 2462 | 2463 |
| 9Ax | 2464 | 2465 | 2466 | 2467 | 2468 | 2469 | 2470 | 2471 | 2472 | 2473 | 2474 | 2475 | 2476 | 2477 | 2478 | 2479 |
| 98x | 2480 | 2481 | 2482 | 2483 | 2484 | 2485 | 2486 | 2487 | 2488 | 2489 | 2490 | 2491 | 2492 | 2493 | 2494 | 2495 |
| 9 Cx | 2496 | 2497 | 2498 | 2499 | 2500 | 2501 | 2502 | 2503 | 2504 | 2505 | 2506 | 2507 | 2508 | 2509 | 2510 | 2511 |
| 9Dx | 2512 | 2513 | 2514 | 2515 | 2516 | 2517 | 2518 | 2519 | 2520 | 2521 | 2522 | 2523 | 2524 | 2525 | 2526 | 2527 |
| 9Ex | 2528 | 2529 | 2530 | 2531 | 2532 | 2533 | 2534 | 2535 | 2536 | 2537 | 2538 | 2539 | 2540 | 2541 | 2542 | 2543 |
| 9Fx | 2544 | 2545 | 2546 | 2547 | 2548 | 2549 | 2550 | 2551 | 2552 | 2553 | 2554 | 2555 | 2556 | 2557 | 2558 | 2559 |
| A0x | 2560 | 2561 | 2562 | 2563 | 2564 | 2565 | 2566 | 2567 | 2568 | 2569 | 2570 | 2571 | 2572 | 2573 | 2574 | 2575 |
| A1x | 2576 | 2577 | 2578 | 2579 | 2580 | 2581 | 2582 | 2583 | 2584 | 2585 | 2586 | 2587 | 2588 | 2589 | 2590 | 2591 |
| A2x | 2592 | 2593 | 2594 | 2595 | 2596 | 2597 | 2598 | 2599 | 2600 | 2601 | 2602 | 2603 | 2604 | 2605 | 2606 | 2607 |
| A3x | 2608 | 2609 | 2610 | 2611 | 2612 | 2613 | 2614 | 2615 | 2616 | 2617 | 2618 | 2619 | 2620 | 2621 | 2622 | 2623 |
| A4 $x$ | 2624 | 2625 | 2626 | 2627 | 2628 | 2629 | 2630 | 2631 | 2632 | 2633 | 2634 | 2635 | 2636 | 2637 | 2638 | 2639 |
| A5x | 2640 | 2641 | 2642 | 2643 | 2644 | 2645 | 2646 | 2647 | 2648 | 2649 | 2650 | 2651 | 2652 | 2653 | 2654 | 2655 |
| A6x | 2656 | 2657 | 2658 | 2659 | 2660 | 2661 | 2662 | 2663 | 2664 | 2665 | 2666 | 2667 | 2668 | 2669 | 2670 | 2671 |
| A7x | 2672 | 2673 | 2674 | 2675 | 2676 | 2677 | 2678 | 2679 | 2680 | 2681 | 2682 | 2683 | 2684 | 2685 | 2686 | 2687 |
| A8x | 2688 | 2689 | 2690 | 2691 | 2692 | 2693 | 2694 | 2695 | 2696 | 2697 | 2698 | 2699 | 2700 | 2701 | 2702 | 2703 |
| A9x | 2704 | 2705 | 2706 | 2707 | 2708 | 2709 | 2710 | 2711 | 2712 | 2713 | 2714 | 2715 | 2716 | 2717 | 2718 | 2719 |
| AAx | 2720 | 2721 | 2722 | 2723 | 2724 | 2725 | 2726 | 2727 | 2728 | 2729 | 2730 | 2731 | 2732 | 2733 | 2734 | 2735 |
| ABx | 2736 | 2737 | 2738 | 2739 | 2740 | 2741 | 2742 | 2743 | 2744 | 2745 | 2746 | 2747 | 2748 | 2749 | 2750 | 2751 |
| $A^{\text {ACx }}$ | 2752 | 2753 | 2754 | 2755 | 2756 | 2757 | 2758 | 2759 | 2760 | 2761 | 2762 | 2763 | 276! | $\bigcirc 765$ | 2766 | 2767 |
| ADx | 2768 | 2769 | 2770 | 2771 | 2772 | 2773 | 2774 | 2775 | 2776 | 2777 | 2778 | 2779 | 2780 | 2781 | 2782 | 2783 |
| AEx | 2784 | 2785 | 2786 | 2787 | 2788 | 2789 | 2790 | 2791 | 2792 | 2793 | 2794 | 2795 | 2796 | 2797 | 2798 | 2799 |
| AFx | 2800 | 2801 | 2802 | 2803 | 2804 | -2805 | 2806 | 2807 | 2808 | 2809 | 2810 | 2811 | 2812 | 2813 | 2814 | 2815 |
| B0x | 2816 | 2817 | 2818 | 2819 | 2820 | 2821 | 2822 | 2823 | 2824 | 2825 | 2826 | 2827 | 2828 | 2829 | 2830 | 2831 |
| B1x | 2832 | 2833 | 2834 | 2835 | 2836 | 2837 | 2838 | 2839 | 2840 | 2841 | 2842 | 2843 | 2844 | 2845 | 2846 | 2847 |
| B2x | 2848 | 2849 | 2850 | 2851 | 2852 | 2853 | 2854 | 2855 | 2856 | 2857 | 2858 | 2859 | 2860 | 2861 | 2862 | 2863 |
| B3x | 2864 | 2865 | 2866 | 2867 | 2868 | 2869 | 2870 | 2871 | 2872 | 2873 | 2874 | 2875 | 2876 | 2877 | 2878 | 2879 |
| B4x | 2880 | 2881 | 2882 | 2883 | 2884 | 2885 | 2886 | 2887 | 2888 | 2889 | 2890 | 2891 | 2892 | 2893 | 2894 | 2895 |
| B5x | 2896 | 2897 | 2898 | 2899 | 2900 | 2901 | 2902 | 2903 | 2904 | 2905 | 2906 | 2907 | 2908 | 2909 | 2910 | 2911 |
| B6x | 2912 | 2913 | 2914 | 2915 | 2916 | 2917 | 2918 | 2919 | 2920 | 2921 | 2922 | 2923 | 2924 | 2925 | 2926 | 2927 |
| B7x | 2928 | 2929 | 2930 | 2931 | 2932 | 2933 | 2934 | 2935 | 2936 | 2937 | 2938 | 2939 | 2940 | 2941 | 2942 | 2943 |
| B8x | 2944 | 2945 | 2946 | 2947 | 2948 | 2949 | 2950 | 2951 | 2952 | 2953 | 2954 | 2955 | 2956 | 2957 | 2958 | 2959 |
| B9x | 2960 | 2961 | 2962 | 2963 | 2964 | 2965 | 2966 | 2967 | 2968 | 2969 | 2970 | 2971 | 2972 | 2973 | 2974 | 2975 |
| BAX | 2976 | 2977 | 2978 | 2979 | 2980 | 2981 | 2982 | 2983 | 2984 | 2985 | 2986 | 2987 | 2988 | 2989 | 2990 | 2991 |
| BBx | 2992 | 2993 | 2994 | 2995 | 2996 | 2997 | 2998 | 2999 | 3000 | 3001 | 3002 | 3003 | 3004 | 3005 | 3006 | 3007 |
| BCx | 3008 | 3009 | 3010 | 3011 | 3012 | 3013 | 3014 | 3015 | 3016 | 3017 | 3018 | 3019 | 3020 | 3021 | 302:. | 3023 |
| BDx | 3024 | 3025 | 3026 | 3027 | 3028 | 3029 | 3030 | 3031 | 3032 | 3033 | 3034 | 3035 | 3036 | 3037 | 3038 | 3039 |
| BEx | 3040 | 3041 | 3042 | 3043 | 3044 | 3045 | 3046 | 3047 | 3048 | 3049 | 3050 | 3051 | 3052 | 3053 | 3054 | 3055 |
| BFx | 3056 | 3057 | 3058 | 3059 | 3060 | 3061 | 3062 | 3063 | 3064 | 3065 | 3066 | 3067 | 3068 | 3069 | 3070 | 3071 |


|  | $x=0$ | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | A | B | C | D | E | $F$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| C0x | 3072 | 3073 | 3074 | 3075 | 3076 | 3077 | 3078 | 3079 | 3080 | 3081 | 3082 | 3083 | 3084 | suos | 3086 | 3087 |
| C1x | 3088 | 3089 | 3090 | 3091 | 3092 | 3093 | 3094 | 3095 | 3096 | 3097 | 3098 | 3099 | 3100 | 3101 | 3102 | 3103 |
| C2x | 3104 | 3105 | 3106 | 3107 | 3108 | 3109 | 3110 | 3111 | 3112 | 3113 | 3114 | 3115 | 3116 | 3117 | 3118 | 3119 |
| C3x | 3120 | 3121 | 3122 | 3123 | 3124 | 3125 | 3126 | 3127 | 3128 | 3129 | 3130 | 3131 | 3132 | 3133 | 3134 | 3135 |
| C4x | 3136 | 3137 | 3138 | 3139 | 3140 | 3141 | 3142 | 3143 | 3144 | 3145 | 3146 | 3147 | 3148 | 3149 | 3150 | 3151 |
| C5x | 3152 | 3153 | 3154 | 3155 | 3156 | 3157 | 3158 | 3159 | 3160 | 3161 | 3162 | 3163 | 3164 | 3165 | 3166 | 3167 |
| C6x | 3168 | 3169 | 3170 | 3171 | 3172 | 3173 | 3174 | 3175 | 3176 | 3177 | 3178 | 3179 | 3180 | 3181 | 3182 | 3183 |
| C7x | 3184 | 3185 | 3186 | 3187 | 3188 | 3189 | 3190 | 3191 | 3192 | 3193 | 3194 | 3195 | 3196 | 3197 | 3198 | 3199 |
| C8x | 3200 | 3201 | 3202 | 3203 | 3204 | 3205 | 3206 | 3207 | 3208 | 3209 | 3210 | 3211 | 3212 | 3213 | 3214 | 3215 |
| C9x | 3216 | 3217 | 3218 | 3219 | 3220 | 3221 | 3222 | 3223 | 3224 | 3225 | 3226 | 3227 | 3228 | 3229 | 3230 | 3231 |
| CAx | 3232 | 3233 | 3234 | 3235 | 3236 | 3237 | 3238 | 3239 | 3240 | 3241 | 3242 | 3243 | 3244 | 3245 | 3246 | 3247 |
| CBx | 3248 | 3249 | 3250 | 3251 | 3252 | 3253 | 3254 | 3255 | 3256 | 3257 | 3258 | 3259 | 3260 | 3261 | 3262 | 3263 |
| CCx | 3264 | 3265 | 3266 | 3267 | 3268 | 3269 | 3270 | 3271 | 3272 | 3273 | 3274 | 3275 | 3276 | 3277 | 3278 | 3279 |
| CDx | 3280 | 3281 | 3282 | 3283 | 3284 | 3285 | 3286 | 3287 | 3288 | 3289 | 3290 | 3291 | 3292 | 3293 | 3294 | 3295 |
| CEx | 3296 | 3297 | 3298 | 3299 | 3300 | 3301 | 3302 | 3303 | 3304 | 3305 | 3306 | 3307 | 3308 | 3309 | 3310 | 3311 |
| CFX | 3312 | 3313 | 3314 | 3315 | 3316 | 3317 | 3318 | 3319 | 3320 | 3321 | 3322 | 3323 | 3324 | 3325 | 3326 | 3327 |
| D0x | 3328 | 3329 | 3330 | 3331 | 3332 | 3333 | 3334 | 3335 | 3336 | 3337 | 3338 | 3339 | 3340 | 3341 | 3342 | 3343 |
| D1x | 3344 | 3345 | 3346 | 3347 | 3348 | 3349 | 3350 | 3351 | 3352 | 3353 | 3354 | 3355 | 3356 | 3357 | 3358 | 3359 |
| D2x | 3360 | 3361 | 3362 | 3363 | 3364 | 3365 | 3366 | 3367 | 3368 | 3369 | 3370 | 3371 | 3372 | 3373 | 3374 | 3375 |
| D3x | 3376 | 3377 | 3378 | 3379 | 3380 | 3381 | 3382 | 3383 | 3384 | 3385 | 3386 | 3387 | 3388 | 3389 | 3390 | 3391 |
| D4x | 3392 | 3393 | 3394 | 3395 | 3396 | 3397 | 3398 | 3399 | 3400 | 3401 | 3402 | 3403 | 3404 | 3405 | 3406 | 3407 |
| D5x | 3408 | 3409 | 3410 | 3411 | 3412 | 3413 | 3414 | 3415 | 3416 | 3417 | 3418 | 3419 | 3420 | 3421 | 3422 | 3423 |
| D6x | 3424 | 3425 | 3426 | 3427 | 3428 | 3429 | 3430 | 3431 | 3432 | 3433 | 3434 | 3435 | 3436 | 3437 | 3438 | 3439 |
| D7x | 3440 | 3441 | 3442 | 3443 | 3444 | 3445 | 3446 | 3447 | 3448 | 3449 | 3450 | 3451 | 3452 | 3453 | 3454 | 3455 |
| D8x | 3456 | 3457 | 3458 | 3459 | 3460 | 3461 | 3462 | 3463 | 3464 | 3465 | 3466 | 3467 | 3468 | 3469 | 3470 | 3471 |
| D9x | 3472 | 3473 | 3474 | 3475 | 3476 | 3477 | 3478 | 3479 | 3480 | 3481 | 3482 | 3483 | 3484 | 3485 | 3486 | 3487 |
| DAx | 3488 | 3489 | 3490 | 3491 | 3492 | 3493 | 3494 | 3495 | 3496 | 3497 | 3498 | 3499 | 3500 | 3501 | 3502 | 3503 |
| DBx | 3504 | 3505 | 3506 | 3507 | 3508 | 3509 | 3510 | 3511 | 3512 | 3513 | 3514 | 3515 | 3516 | 3517 | 3518 | 3519 |
| DCx | 3520 | 3521 | 3522 | 3523 | 3524 | 3525 | 3526 | 3527 | 3528 | 3529 | 3530 | 3531. | 3532 | 3533 | 3534 | 3535 |
| DDx | 3536 | 3537 | 3538 | 3539 | 3540 | 3541 | 3542 | 3543 | 3544 | 3545 | 3546 | 3547 | 3548 | 3549 | 3550 | 3551 |
| DEX | 3552 | 3553 | 3554 | 3555 | 3556 | 3557 | 3558 | 3559 | 3560 | 3561 | 3562 | 3563 | 3564 | 3565 | 3566 | 3567 |
| DFX | 3568 | 3569 | 3570 | 3571 | 3572 | 3573 | 3574 | 3575 | 3576 | 3577 | 3578 | 3579 | 3580 | 3581 | 3582 | 3583 |
| E0x | 3584 | 3585 | 3586 | 3587 | 3588 | 3589 | 3590 | 3591 | 3592 | 3593 | 3594 | 3595 | 3596 | 3597 | 3598 | 3599 |
| E1x | 3600 | 3601 | 3602 | 3603 | 3604 | 3605 | 3606 | 3607 | 3608 | 3609 | 3610 | 3611 | 3612 | 3613 | 3614 | 3615 |
| E2x | 3616 | 3617 | 3618 | 3619 | 3620 | 3621 | 3622 | 3623 | 3624 | 3625 | 3626 | 3627 | 3628 | 3629 | 3630 | 3631 |
| E3x | 3632 | 3633 | 3634 | 3635 | 3636 | 3637 | 3638 | 3639 | 3640 | 3641 | 3642 | 3643 | 3644 | 3645 | 3646 | 3647 |
| E4x | 3648 | 3649 | 3650 | 3651 | 3652 | 3653 | 3654 | 3655 | 3656 | 3657 | 3658 | 3659 | 3660 | 3661 | 3662 | 3663 |
| E5x | 3664 | 3665 | 3666 | 3667 | 3668 | 3669 | 3670 | 3671 | 3672 | 3673 | 3674 | 3675 | 3676 | 3677 | 3678 | 3679 |
| E6x | 3680 | 3681 | 3682 | 3683 | 3684 | 3685 | 3686 | 3687 | 3688 | 3689 | 3690 | 3691 | 3692 | 3693 | 3694 | 3695 |
| E7x | 3696 | 3697 | 3698 | 3699 | 3700 | 3701 | 3702 | 3703 | 3704 | 3705 | 3706 | 3707 | 3708 | 3709 | 3710 | 3711 |
| E8x | 3712 | 3713 | 3714 | 3715 | 3716 | 3717 | 3718 | 3719 | 3720 | 3721 | 3722 | 3723 | 3724 | 3725 | 3786 | 3727 |
| E9x | 3728 | 3729 | 3730 | 3731 | 3732 | 3733 | 3734 | 3735 | 3736 | 3737 | 3738 | 3739 | 3740 | 3741 | 3742 | 3743 |
| EAx | 3744 | 3745 | 3746 | 3747 | 3748 | 3749 | 3750 | 3751 | 3752 | 3753 | 3754 | 3755 | 3756 | 3757 | 3758 | 3759 |
| EBx | 3760 | 3761 | 3762 | 3763 | 3764 | 3765 | 3766 | 3767 | 3768 | 3769 | 3770 | 3771 | 3772 | 3773 | 3774 | 3775 |
| ECx | 3776 | 3777 | 3778 | 3779 | 3780 | 3781 | 3782 | 3783 | 3784 | 3785 | 3786 | 3787 | 3788 | 3789 | 3790 | 3791 |
| EDX | 3792 | 3793 | 3794 | 3795 | 3796 | 3797 | 3798 | 3799 | 3800 | 3801 | 3802 | 3803 | 3804 | 3805 | 3806 | 3807 |
| EEX | 3808 | 3809 | 3810 | 3811 | 3812 | 3813 | 3814 | 3815 | 3816 | 3817 | 3818 | 3819 | 3820 | 3821 | 3822 | 3823 |
| EFx | 3824 | 3825 | 3826 | 3827 | 3828 | 3829 | 3830 | 3831 | 3832 | 3833 | 3834 | 3835 | 3836 | 3837 | 3838 | 3839 |
| F0x | 3840 | 3841 | 3842 | 3843 | 3844 | 3845 | 3846 | 3847 | 3848 | 3849 | 3850 | 3851 | 3352 | 3853 | 3854 | 3855 |
| F1x | 3856 | 3857 | 3858 | 3859 | 3860 | 3861 | 3862 | 3863 | 3864 | 3865 | 3866 | 3867 | 3868 | 3869 | 3870 | 3871 |
| F2x | 3872 | 3873 | 3874 | 3875 | 3876 | 3877 | 3878 | 3879 | 3880 | 3881 | 3882 | 3883 | 3884 | 3885 | 3886 | 3887 |
| F3x | 3888 | 3889 | 3890 | 3891 | 3892 | 3893 | 3894 | 3895 | 3896 | 3897 | 3898 | 3899 | 3900 | 3901 | 3902 | 3903 |
| F4x | 3904 | 3905 | 3906 | 3907 | 3908 | 3909 | 3910 | 3911 | 3912 | 3913 | 3914 | 3915 | 3916 | 3917 | 3918 | 3919 |
| F5x | 3920 | 3921 | 3922 | 3923 | 3924 | 3925 | 3926 | 3927 | 3928 | 3929 | 3930 | 3931 | 3932 | 3933 | 3934 | 3935 |
| F6x | 3936 | 3937 | 3938 | 3939 | 3940 | 3941 | 3942 | 3943 | 3944 | 3945 | 3946 | 3947 | 3948 | 3949 | 3950 | 3951 |
| F7x | 3952 | 3953 | 3954 | 3955 | 3956 | 3957 | 3958 | 3959 | 3960 | 3961 | 3962 | 3963 | 3964 | 3965 | 3966 | 3967 |
| F8x | 3968 | 3969 | 3970 | 3971 | 3972 | 3973 | 3974 | 3975 | 3976 | 3977 | 3978 | 3979 | 3980 | 3981 | 3982 | 3983 |
| F9x | 3984 | 3985 | 3986 | 3987 | 3988 | 3989 | 3990 | 3991 | 3992 | 3993 | 3994 | 3995 | 3996 | 3997 | 3998 | 3999 |
| FAx | 4000 | 4001 | 4002 | 4003 | 4004 | 4005 | 4006 | 4007 | 4008 | 4009 | 4010 | 4011 | 4012 | 4013 | 4014 | 4015 |
| FBx | 4016 | 4017 | 4018 | 4019 | 4020 | 4021 | 4022 | 4023 | 4024 | 4025 | 4026 | 4027 | 4028 | 4029 | 4030 | 4031 |
| FCx | 4032 | 4033 | 4034 | 4035 | 4036 | 4037 | 4038 | 4039 | 4040 | 4041 | 4042 | 4043 | 4044 | 4045 | 4046 | 4047 |
| FDx | 4048 | 4049 | 4050 | 4051 | 4052 | 4053 | 4054 | 4055 | 4056 | 4057 | 4058 | 4059 | 4060 | 4061 | 4062 | 4063 |
| FEx | 4064 | 4065 | 4066 | 4067 | 4068 | 4069 | 4070 | 4071 | 4072 | 4073 | 4074 | 4075 | 4076 | 4077 | 4078 | 4079 |
| FFx | 4080 | 4081 | 4082 | 4083 | 4084 | 4085 | 4086 | 4087 | 4088 | 4089 | 4090 | 4091 | 4092 | 4093 | 4094 | 4095 |

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## Appendix III: Machine Instruction Format




## Notes for Appendix III:

1. R1, R2, and R3 are absolute expressions that specify general or floating-point registers. The general register numbers are 0 through l5; floating-point register numbers are $0,2,4$, and 6 .
2. D1 and D2 are absolute expressions that specify displacements. A value of $0-4095$ may be specified.
3. Bl and B2 are absolute expressions that specify base registers. Register numbers are 0-15.
4. X 2 is an absolute expression that specifies an index register. Register numbers are 0-15.
5. L, L1, and L2 are absolute expressions that specify field lengths. An L expression can specify a value of 1 - 256. Ll and $L 2$ expressions can specify a value of $1-16$ In all cases, the assembled value will be one less than the specified value.
6. I, I2, and I3 are absolute expressions that provide immediate data. The value of I and I2 may be 0 - 255. The value of I3 may be $0-9$.
7. S1 and S2 are absolute or relocatable expressions that specify an address.
8. RR; RS, and SI instruction fields that are blank under BASIC MACHINE FORMAT are not examined during instruction execution. The fields are not written in the symbolic operand, but are assembled as binary zeros.
9. MI and M3 specify a 4-bit mask.
10. In IBM System/370 the SIO, HIO, HDV and SIOF operation codes occupy one byte and the low order bit of the second byte. In all other systems the HIO and SIO operation codes occupy only the first byte of the instruction.
Appendix IV, pages 391-406, has been
removed. The following manuals are
the definitive publications for
machine instructions:
IBM System/370 Principles of
Operation
IBM 4300 Processors Principles of
Operation for ECPS:VSE Mode

## Appendix V: Assembler Instructions



| Operation | Name Entry | Operand Entry |
| :---: | :---: | :---: |
| EQU | An ordinary symbol or a variable symbol | One to three operands, separated by commas dos Only one operand |
| EXIRN | A sequence symbol or blank | One or more relocatable symbcls, separated by commas |
| GBLA | rust not be present | One or more variable symbols that are to be used as SEr symbols, separated by commas ${ }^{2}$ |
| GBT, B | Must not be present | One or more variable symbcls that are to be used as SET symbols, separated by commas ${ }^{2}$ |
| GBLC | Must not be present | One or more variable symbols that are to be used as SET symbols, separated by commas ${ }^{2}$ |
| IC'L | Must not be present | One to three decimal values, separated by commas |
| IS EQ | Must not be present | Two decimal values, separated by commas |
| LCI_A | Must not be present | One or more variable symbols that are to be used as SET symbols, separated by commas ${ }^{2}$ |
| LCLB | Must not be present | One or more variable symbols that are to be used as SET symbols, separated by commas ${ }^{2}$ |
| LCLC | Must not be present | One or more variable symbols that are to be used as SEI symbols, separated by commas ${ }^{2}$ |
| LTORG | Any symbol or blank | Not required |
| MACRO ${ }^{1}$ | Must not be present | Not required |
| MEND ${ }^{1}$ | A sequence symbol or blank | Not required |
| MEXIT ${ }^{1}$ | A sequence symbol or blank | Not required |
| MNOTE | A sequence symbol or blank | A severity code foll owed by a comma (this much is optional) followed by any combination of characters enclosed in apostrophes |


| Operation | Name Entry | Operand Entry |
| :---: | :---: | :---: |
| OS OPIV OPSYN | An ordinary symbol | A machine instruction mnemonic code, an extended mnemonic code, a macro operation, an assembler operation, an operation code defined by a previous OPSYN instruction, or blank |
| or OPSYN | A machine instruction mnemonic code, an extended mnemonic code, an assembler operation, an operation code defined by a previous OPSYN instruction | B1ank ${ }_{\text {chel }}$ |
| ORG | OS Any symbol or only blank <br> DOS A sequence symbol only or blank | A relocatable expression or blank <br> A relocatable expression or blank |
| os POP | A sequence symbol or blank | One or more operands, separated by a comma |
| PRINT | A sequence symbol or blank | One to three operands |
| PUNCH | A sequence symbol or blank | One to eighty characters, enclosed in apostrophes |
| Only PUSH | A sequence symbol or blank | One or more operands, separated by a comma |
| REPRO | A sequence symbol or blank | Not required |
| SETA | A SETA symbol | An arithmetic expression |
| SETB | A SETB symbol | A 0 or a 1 , a SETB symbol, or a logical expression enclosed in parentheses |
| SETC | A SETC symbol | A type attribute, a character expression, a substring notation, or a concatenation of character expressions and substring notations. <br> OS A duplication factor (a SEIA only expression enclosed in parentheses) can precede the above if desired. |
| SPACE | A sequence symbol or blank | A decimal self-defining term or blank |
| START | Any symbol or blank | A self-defining term or blank |
| TITLE | A string of alphameric characters. A variable symbol. A combination of the above. <br> A sequence symbol. A blank | One to 100 characters, enclosed in apostrophes |


| Operation | Name_Entry | Operand Entry |
| :---: | :---: | :---: |
| USING | A sequence symbol or blank | An absolute or relocatable expression followed by 1 to 16 absolute expressions, separated by commas |
| WXTRN | A sequence symbol or blank | One or more relocatable symbols, separated by commas |
| Instruction | Name Entry | Operand_Entry |
| Model Statements ${ }^{3}$ | An ordinary symbol, a variable symbol, a sequence symbol, a combination of variable symbols and other characters that is equivalent to a symbol, or blank | Any combination of characters (including variable symbols) |
| Prototype Statement ${ }^{2}$ | A symbolic parameter or blank | Zero or more operands that are symbolic parameters, separated by commas |
| Macro-Instruction Statement ${ }^{2}$ | An ordinary symbol, a variable symbol, a sequence symbol, a combination of variable symbols and other characters that is equivalent to a symbol, ${ }^{2}$ or blank | zero or more positional operands and/or zero or more keyword operands separated by commas ${ }^{2}$ |
| Assembler Language Statement ${ }^{3}$ | An ordinary symbol, a variable symbol, a sequence symbol, a combination of variable symbols and other characters that is equivalent to a symbol, or blank | Any combination of characters (including variable symbols) |
| 1 Can only be used as part of a macro definition. |  |  |
| 2 Variable symbols appearing in a macro instruction are replaced by their values before the macro instruction is processed. |  |  |
| 3 Restrictions on the in statement fields for each individua for Model Statement | se of variable symbo included in the d tatement and in "Rul elds" (See J4B). | ions |

## Appendix VI: Summary of Constants



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# Appendix VII: Summary of Macro Facility 

The four charts in this Appendix summarize the macro facility described in Part IV of this publication.

Chart 1 indicates which macro language elements can be used in the name and operand entries of each statement.

Chart 2 is a summary of the expressions that can be used in macro instruction statements. Chart 3 is a summary of the attributes that may be used in each expression. Chart 4 is a summary of the variable symbols that can be used in each expression.


Chart 2. Conditional Assembly Expressions


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| Expression <br> Operations are | Arithmetic Expressions <br> +, - (unary and binary), *, and /; <br> parentheses per- <br> mitted | Character Expressions <br> Concatenation, with a period (.) | Logical Expressions <br> AND, OR, and NOT parentheses permitted |
| :---: | :---: | :---: | :---: |
| Range of values | $-2^{31}$ to +231-1 | 0 through 255 characters | 0 (false) or 1 (true) |
| May be used in | - SETA operands <br> - Arithmetic relations <br> - Subscripted SET symbols <br> - ESYSLIST subscript <br> - Substring notation <br> - Sublist notation | - SETC operands <br> - Character relations ${ }^{2}$ | - SETB operands <br> - AIF operands |
| 1 An arithmetic relation consists of two arithmetic expressions related by the operators GT, LT, EQ, NE, GE, or LE. <br> 2 A character relation consists of two character expressions related by the operator GT, LT, EQ, NE, GE, or LE. Type attribute notation and Substring notation may also be used in character relations. The maximum size of the character expressions that can be compared is 255 characters. If the two character expressions are of unequal size, the the smaller one will always compare less than the larger. |  |  |  |

Chart 3. Attributes

| Attribute | Notation | Can be used with: | Can be used only if type attribute is: | Can be used in |
| :---: | :---: | :---: | :---: | :---: |
| Type | $T^{\prime}$ | Ordinary Symbols defined in open code; symbolic parameters inside macro definitions; ESYSLIST (m), ESYSL IST $(m, n)$, <br> SET symbols; ESYSTIME, ESYSPARM, ESYSDATE, ESYSECT, ESYSNDX | (May always be used) | 1. SETC operand fields <br> 2. Character relations |
| Length | L' | Ordinary Symbols defined in open code; symbolic parameters inside macro definitions: ESYSLIST (m) and ESYSLIST $(m, n)$ inside macro definitions | Any letter except $\mathrm{M}, \mathrm{N}, \mathrm{O}, \mathrm{T}$ and U | Arithmetic expressions |
| Scaling | $S^{\prime}$ | Ordinary Symbols defined in open code; symbolic parameters inside macro definitions; ESYSLIST (m) , and ESYSLIST ( $m, n$ ) inside macro definitions | $\begin{aligned} & \text { H, F, G, D, E, L, K, P, } \\ & \text { and } Z \end{aligned}$ | Arithmetic expressions |
| Integer | I' | Ordinary Symbols defined in open code; symbolic parameters inside macro definitions; ESYSLIST (m) , and ESYSLIST ( $\mathrm{m}, \mathrm{n}$ ) inside macro definitions | $\begin{aligned} & \text { H, F, G, D, E, L, K, P, } \\ & \text { and } Z \end{aligned}$ | Arithmetic expressions |
| Count | $K^{\prime}$ OS on | Symbolic parameters, ESYSLIST (m) and ESYSLIST ( $\mathrm{m}, \mathrm{n}$ ) inside macro definitions SET symbols; all system variable symbols | Any letter | Arithmetic expressions |
| Number | $\mathrm{N}^{\prime}$ | Symbolic parameters, ESYSLIST and GSYSLIST (m) inside macro definitions | Any letter | Arithmetic expressions |

NOTE: There are definite restrictions in the use of these attributes ( see L1B).

Chart 4. Variable Symbols (Part 1 of 2)

| Variable Symbol | Declared by: | Initialized, or set to: | Value changed by: | May be used in: |
| :---: | :---: | :---: | :---: | :---: |
| Symbolic ${ }^{1}$ parameter | Prototype statement | Corresponding macro instruction operand | (Constant throughout definition) | - Arithmetic expressions if operand is decimal self-defining term <br> - Character expressions |
| SETA | LCLA or GBLA instruction | 0 | $\begin{aligned} & \text { SETA } \\ & \text { instruction } \end{aligned}$ | - Arithmetic expressions <br> - Character expressions |
| SETB | LCLB or GBLB instruction | 0 | $\begin{aligned} & \text { SETB } \\ & \text { instruction } \end{aligned}$ | - Arithmetic expressions <br> - Character expressions <br> - Logical expressions |
| SETC | LCLC or GELC instruction | String of length 0 (null) | $\begin{aligned} & \text { SETC } \\ & \text { instruction } \end{aligned}$ | - Arithmetic expressions if value is decimal selfdefining term <br> - Character expressions |
| \&SYSNDX 1 | The assembler | Macro <br> instruction <br> index | (Constant throughout definition; unique for each macro instruction) | - Arithmetic expressions <br> - Character expressions |
| \&SYSECT ${ }^{1}$ | The assembler | Control section in which macro instruction appears | (Constant throughout definition; set by CSECT, DSECT, START, and COM) | - Character expressions |
| \&SYSLIST ${ }^{1}$ | The assembler | Not applicable | Not applicable | - N'ESYSLIST in arithmetic expressions |
| $\begin{aligned} & \text { \&SYSLIST } \\ & (n)^{1} \\ & \& S Y S L I S T \\ & (n, m)^{1} \end{aligned}$ | The assembler | Corresponding macro instruction operand | (Constant throughout definition) | - Arithmetic expressions if operand is decimal self-defining term |
|  |  |  |  | - Character expressions |

${ }^{1}$ Can be used only in macro definitions.

Chart 4. Variable Symbols cont. (Part 2 of 2)

| Variable Symbol | Declared by: | Initialized, or set to: | Value changed by: | May be used in: |
| :---: | :---: | :---: | :---: | :---: |
| ESYSPARM | PARM field | User defined or null | Constant throughout assembly | - Arithmetic expression if value is decimal selfdefining term <br> - Character expression |
| ESYSTIME | The assembler | System time | Constant throughout assembly | - Character expression |
| ESYSDATE | The assembler | System date | Constant throughout assembly | - Character expression |

${ }^{1}$ Can be used only in macro definitions.

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## Glossary

This glossary has three main types of definitions that afply:

- To the assembler language in particular (usually distinguished by reference to the words "assembler". "assembly", etc.)
- To programming in general
- To data processing as a whole

If you do not understand the meaning of a data processing term used in any of the definitions below, refer to the IEM Data Processing Glossary, Order No. GC20-1699.

IBM is grateful to the American National Standards Institute (ANSI) for permission to reprint its definitions from the American National Standard Vocabulary for Information Processing, which was prepared by Subcommittee X3 K5 on Terminology and Glossary of American National Standards Committee X3.

ANSI definitions are preceded by an asterisk ( $\star$ ).

* absolute address: A pattern of characters that identifies a unique storage location without further modification.
absolute expression: An assembly-time expression whose value is not affected by program relocation. An absolute expression can represent an absolute address.
absolute term: A term whose value is not affected ky relocation.


## *address:

1. An identification, as represented by a name, label, or number, for a register, location in storage, or any cther data source or destination such as the location of a station in a communication network.
2. Loosely, any part of an instruction that specifies the location of an operand for the instruction. Synonymous with address reference.
3. See absolute address, base address, explicit address, implicit address, symbolic address.
address constant: A value, or an expression representing a value, used in the calculation of storage addresses.
address reference: Same as address (2) •
alignment: The positioning of the beginning of a machine instruction, data constant, or area on a proper boundary in virtual storage.
alphabetic character: In assembler programming, the letters A through $Z$ and $\$$. \#, 2.
*alphameric: Same as alphanumeric.
*alphanumeric: Pertaining to a character set that contains letters, digits, and usually, other characters, such as punctuation marks. Synonymous with alphameric.
*AND: A logic operator having the property that if $P$ is a statement, $Q$ is a statement, $R$ is a statement,.... then the AND of $P$, $Q$, R.... is true if all statements are true, false if any statement is false.
arithmetic expression: A conditicnal assembly expression that is a combination of arithmetic terms, arithmetic cperators, and paired parentheses.

## arithmetic operator:

1. In assembler programming, an cperator that can be used in an absolute or relocatable expression, or in an arithmetic expression to indicate the
actions to be performed on the terms in the expression. The arithmetic operators allowed are: +, -, *, /.
2. See binary operator, unary operator.
arithmetic relation: Two arithmetic expressions separated by a relational operator.
*arithmetic shift:
3. A shift that does not affect the sign position.
4. A shift that is equivalent to the multiplication of a number ky a positive or negative integral power of the radix.
arithmetic term: A term that can be used only in an arithmeitc expression.
array: In assembler programming, a series of one or more values represented by a SET symbol.
*assemble: To prepare a machine language program from a symbolic language program by substituting absolute operation codes for symbolic operation codes and absolute or relocatable addresses for symbolic addresses.
*assembler: A computer program that assembles.
assembler instruction:
5. An assembler language statement that causes the assembler to perform a specific operation. Unlike the machine instructions, the assembler instructions are not translated into machine language.
6. See also conditional assembly instruction, macro processing instruction.
assembler lanquage: A source language that includes symbolic machine language statements in which there is a one-to-one correspondence with the instruction formats and data formats of the computer. The assembler language also includes statements that represent assembler instructions and macro instructions.
assembly time: The time at which the assembler translates the symbolic machine language statements into their object code form (machine instructions). The assembler also processes the assembler instructions at this time, with the exception of the conditional assembly and macro processing instructions, which it processes at pre-assembly time.
attribute: A characteristic of the data defined in a source module. The assembler assigns the value of an attribute to the symbol or macro instruction operand that represents the data. Synonymous with data attribute.
*base:
7. A number that is multiplied by itself as many times as indicated by an exponent.
8. See floating-point base.
*base address: A given address from which an absolute address is derived by combination with a relative address. NOTE: In
assembler programming, the relative address is synonymous with displacement.
base register: A register that contains the base address.
*binary: Pertaining to the number representation system with a radix of two.
*binary digit: In binary notation, either of the characters, 0 or 1 .
binary operator: An arithmetic operator having two terms. The binary operators that can be used in absolute or relocatable expressions and arithmetic expressions are: addition ( + ), subtraction ( $(-)$, multiplication (*) : and division (). Contrast with unary operator.
*bit: A binary digit.
bit-length modifier: A subfield in the DC assembler instruction that determines the length in bits of the area into which the defined data constant is to be assembled.
bit string: A string of binary digits in which the position of each binary digit is considered as an independent unit.
blank: In assembler programming, the same as space character.
rblank character: Same as space character.
boundary: In assembler programming, a location in storage that marks the beginning of an area into which data is a ssembled. For example, a fullword boundary is a location in storage whose address is divisible by four. The other boundaries are: doubleword (location divisible by eight), halfword (location divisible by two), and byte (location can be any number). See also alignment.

* branch: Loosely, a conditicnal jump.
buffer: An area of storage that is temporarily reserved for use in performing an input/output operation, and into which data is read or from which data is written.
$\star$ buq: A mistake or malfunction.
byte:

1. A sequence of adjacent binary digits operated upon as a unit and usually shorter than a computer word.
2. The representation of a character; eight binary digits (bits) in System/370.
call:
$\star$ 1. To transfer control to a specified closed subroutine.
3. See also macro call.

* Character:

1. A letter, digit, or other symbol that is used as part of the organization, control, or representation of data. A character is often in the form of a spatial arrangement of adjacent or connected strokes.
2. See blank character, character set, special character.
character expression: A character string enclosed by apostrophes. It can re used only in conditional assembly instructions.
The enclosing apostrophes are not part of the value represented. Contrast with quoted string.
character relation: Two character strings separated by a relational operator.

Character set:
$\star$ 1. A set of unique representations called characters, for example, the 26 letters of the English alphabet, 0 and 1 of the Boolean alphabet, the set of signals in the Morse code alphabet, the 128 characters of the ASCII alphabet.
2. In assembler programming, the alphabetic characters A through $Z$ and \$, \#, 0 ; the digits, 0 through 9; and the special characters $+-* /,()=$ . ' E and the blank character.
$\star$ Character string: A string consisting solely of characters.
closed subroutine: A subroutine that can be stored at one place and can be linked to one or more calling routines. Contrast with open subroutine.

## * code:

1. A set of unambigous rules specifying the way in which data may be represented, for example, the set of correspondences in the standard code for information interchange.
2. In data processing, to represent data or a computer program in a symbolic form that can be accepted by a data processor.
3. To write a routine.
4. See condition code, object code, operation code.
*coding: See symbolic coding.
collating sequence: An ordering assigned to a set of items, such that any two sets in that assigned order can be collated.
*column: A vertical arrangement of characters or other expressions.
comments statement: A statement used to include information that may be helpful in running a job or reviewing an cutput
listing.

* Complement:

1. A number that can be derived from a specified number by subtracting it from a second specified number. For example, in radix notation, the second specified number may be given power of the radix or one less than the given power of the radix. The negative of the number is often represented by its complement.
2. See radix complement, twos complement.
complex relocatable expression: A
relocatable expression that contains two or more unpaired relocatable terms or an unpaired relocatakle term preceded by a minus sign, after all unary operators have been resolved. A complex relccatable expression is not fully evaluated until program fetch time.
*Computer program: A series of instructions or statements, in a form acceptable to a computer, prepared in order to achieve a certain result.
*Computer word: A sequence of bits or characters treated as a unit and capable of being stored in one computer location.
concatenation character: The period (.) that is used to separate character strings that are to be joined together in conditional assembly processing.
condition code: A code that reflects the result of a previous input/output, arithmetic, or logical operation.
conditional assembly: An assembler facility for altering at pre-assembly time the content and sequence of source statements that are to be assembled.
conditional assembly expression: An expression that the assembler evaluates at pre-assembly time.
conditional assembly instruction: An assembler instruction that performs a conditional assembly operation. Conditional assembly instructions are processed at pre-assembly time. They include: the LCLA, LCLB, LCLC, GBLA, GBLB, and the GBIC declaration instructions; the SETA, SETB, and SETC assignment instructions; the AIF, AGO, ANOP, and ACTR branching instructions.

* conditional jump: A jump that occurs if specified criteria are met.
* constant: See figurative constant.
continuation line: A line of a source statement into which characters are entered when the source statement cannot ke contained on the preceding line or lines.
control program: A program that is designed to schedule and supervise the performance of data processing work by a computing system.
control section: That part of a program specified by the programmer to be a relocatable unit, all elements of which are to be loaded into adjoining virtual storage locations. Abbreviated CSECT.
control statement: See linkage editor control statement:

Copy: To reproduce data in a new location or other destination, leaving the source data unchanged, although the physical form of the result may differ from that of the source. For example, to copy a deck of cards onto a magnetic tape.
count attribute ( $K^{\prime}$ ): An attribute that gives the number of char acters that would be required to represent the data as a character string.

## * counter:

1. A device such as a register or storage location used to represent the number of occurrences of an event.
2. See instruction counter, location counter.

CPU: Central processing unit.
CSECT: See control section.
data attribute: Same as attribute.
data constant: See figurative constant.
*debug: To detect, locate, and remove mistakes from a routine or malfunctions from a computer.
*decimal: Pertaining to the number representation system with a radix of ten.
declare: To identify the variable symbols to be used ly the assembler at pre-assembly time.
*delimiter: A flag that separates and organizes items of data.
*device: See storage device.
kdictionary: Sef external symbcl dictionary.
dimension: The maximum number of values that can be assigned to a SET symbol representing an array.
dimensioned SET symbol: A SET symbcl, representing an array, followed by a decimal number enclosed in parentheses. A dimensioned SET symbol must be declared in a glokal (GBLA, GBLB, or GBLC) or local
(LCLA, LCLB, LCLC) declaration instruction.

## displacement:

*1. Same as relative address.
2. In assembler programming, the difference in bytes between a symbolic address and a specified base address.
doubleword: A contiguous sequence cf bits or characters which comprises two computer words and is capable of being addressed as a unit.
NOTE: In assembler programming, the doubleword has a length of eight bytes and can be aligned on a doubleword boundary (a location whose address is divisible by eight). Contrast with fullword, halfword.
kdummy: Pertaining to the characteristic of having the appearance of a specified thing but not having the capacity tc function as such. For example, a dummy control section.
dummy control section: A control section that the assembler can use to format an
area of storage without producing any object code. Synonymous with dummy section.
dummy section: Same as dumny control
duplication factor: In assembler programming, a value that indicates the number of times that the data specified immediately following the duplication factor is to be generated. For example, the first sutfield of a DC or DS instruction is a duplication factor.
$\star$ dynamic storage allocation: A storage allocation technique in which the location of computer programs and data is determined by criteria applied at the moment of need.

EBCDIC: Extended binary coded decimal interchange code.
entry name: A name within a control section that defines an entry point and can be referred to by any control section.
*entry point: In a routine, any place to which control can be passed.
entry symbol:

1. An ordinary symbol that represents an entry name (identified by the ENTRY assembler instruction) or control section name (defined by the CSECT or START assembler instruction).
2. See also external symbcl.

EQ: (equal to) See relational operator.
*error message: An indication that an error has been detected. Contrast with warning message.

ESD: External symbol dicticnary.
excess sixty-four binary notation: In assembler programming, a binary notation in which each exponent of a floating-point number $E$ is represented by the binary equivalent of $E$ plus sixty-four.
execution time: The time at which the machine instructions in object code form are processed by the central processing unit of the computer.
explicit address: An address reference which is specified as two absolute expressions. One expression supplies the value of a base register and the other supplies the value cf a displacement. The assembler assembles both values into the object code of a machine instruction.
exponent:
*1. In a floating-point representation, the numeral, of a pair of numerals representing a number, that indicates the power to which the base is raised.
2. See also excess sixty-four binary notation.
exponent modifier: A subfield in the operand of the DC assembler instruction that indicates the power of ten by which a number is to ke multiplied before being assembled as a data constant.

## expression:

1. One or more operations represented by a combination of terms, and paired parentheses.
2. See absolute expression, arithmetic expression, complex relocatable expression, relocatable expression.
3. See also character expression.
extended binary coded decimal interchange code: A set of 256 characters, each represented by eight bits.
external name: A name that can be referred to by any control section or separately assembled module; that is, a control section name ór an entry name in another module.
external reference: A reference to a symbol
that is defined as an external name in another module.
external symbol:
4. An ordinary symbol that represents an external reference. An external symbol is identified in a source module by the EXTRN or WXTRN assembler instruction, or by the $v$-type address constant.
5. Loosely, a symbol contained in the external symbol dictionary.
6. See also entry symbol.
external symbol dictionary: Contrcl
information associated with an olject or load module which identifies the external symbols in the module. Abbreviated ESD.

## EXTRN: External reference.

## fetch:

$\star 1$. To locate and load a quantity of data from storage.
2. In the Operating System (OS), to obtain load modules from auxiliary storage and load them into virtual storage. See also loader (1).
3. In the Eisk Operating System (LOS), to bring a program phase into virtual storage from the core image likrary for immediate execution.
4. A control program routine that accomplishes (1), (2), or (3). See also loader (2).
5. The name of the system macro instruction (FETCH) used to accomplish
(1).
(2) , or
(3).

* fiqurative constant: A preassigned, fixed, character string with a preassigned, fixed, data name in a particular programming language.
NCTE: In assembler programming, the two
types of figurative constant are:
a. data and address constants defined by the DC assembler instruction.
b. symbols assigned values ky the EQU assembler instruction.
flaq:
$\star$ 1. Any of various types of indicators used for identification. For example, in assembler programming, the paired apostrophes that enclose a character expression of a quoted string.

2. In assembler programming, to indicate the occurrence of an error.

* floating-point_base: In floating-point representation, the fixed positive integer that is the base of the power. NOTE: In assembler programming, this base is 16.
fullword: A contiguous sequence of kits or characters which comprises a computer word and is capable of being addressed as a unit.
NCTE: In assembler programming, the fullword has a length of four bytes and can be aligned on a fullword koundary (a location whose address is divisible by four). Contrast with doukleword, halfword.

GE: (greater than or equal to) See relational operator.

## generate:

* 1. To produce a program by selection of subsets from a set of skeletal coding under the control of parameters.

2. In assembler programming, to produce assembler language statements from the model statements of a macro definition when the definition is called by a macro instruction.
qlobal scope: Pertaining to that part of an assembler program that includes the kody of any macro definition called from a source
module and the open code portion of the source module. Contrast with local scope.

## qlobal variable symbol:

1. A variable symbol that can be used to communicate values between macro
definitions and ketween a macrc
definition and open code.
2. Contrast with local variable symbol.

GT: (greater than) See relational operator.
$\star$ halfword: A contiguous sequence of bits or characters which comprises half a computer word and is capable of being addressed as a unit.
NOTE: In assembler programming, the
halfword has a length of two bytes and can be aligned on a halfword boundary (a location whose address is divisible by two). Contrast with doubleword, fullword.
hexadecimal: Pertaining to a number system with a radix of sixteen; valid digits range from 0 through $F$, where $F$ represents the highest units position (15).
immediate data: Data specified in an SI type machine instruction that represents a value to be assembled into the object code of the machine instruction.
implicit address: An address reference which is specified as one absclute or relocatable expression. An implicit address must be converted into its explicit
base-displacement form before it can be assembled into the object code of a machine instruction.
index register:
*1. A register whose content may be added to or subtracted from the operand address prior to or during the execution of a computer instruction.
2. In assembler programming, a register whose content is added to the cperand or absolute address derived frem a combination of a base address with a displacement.
inner macro instruction: A macro
instruction that is specified, that is, nested inside a macro definition. Contrast with outer macro instruction.

## *instruction:

1. A statement that specifies an cperation and the values or locations of its oper ands.
2. See assembler instruction, conditional assembly instruction, machine instruction, macro instruction.
$\star$ instruction counter: A counter that indicates the location of the next computer instruction to be interpreted.
instruction statement: See instruction (1).
integer attribute ( $I^{\prime}$ ): An attribute that indicates the number of digit positions occupied by the integer portion of fixed-point, decimal, and floating-point constants in their object code form.
*interrupt: To stop a process in such a way it can be resumed.

* I $\angle O$ : An abbreviation for input/output.
* jok control statement: A statement in a job that is used in identifying the job or describing its requirements to the operating system.
* jump:

1. A departure from the normal sequence of executing instructions in a computer.
2. See conditional jump.
keyword: In assembler frogramming, an ordinary symbol containing up to seven characters. A keyword is used to identify a parameter, called a keyword parameter, in a macro prototype statement and the corresponding macro instruction operand.
keyword operand; An operand in a macro instruction that assigns a value to the corresponding keyword parameter declared in the prototype statement of the called macro definition. Keyword operands can be specified in any order, because they identify the corresponding parameter ky keyword and not by their position. NCTE: In assembler programming, the specification of a keyword oper and has the format: a keyword followed by an equal sign which, in turn, is followed by the value to be assigned to the keyword parameter.
keyword_parameter: A symbolic parameter in which the symbol fcllowing the ampersand represents a keyword.
NOTE: In assembler programming, the declaration of keyword parameter has the format: a keyword parameter followed ky an equal sign which, in turn, is followed by a standard (default) value.
lakel:
$\star \overline{1}$. One or more characters used to identify a statement or an item of data in a computer program.
3. In assembler programming, the entry in the name field of an assembler language statement. The three main types of name entry are:
a. the ordinary symbol which represents a label at assembly time.
b. the sequence symbcl which represents a label at pre-assembly time and is used as a conditional assembly branching destination. c. the variable symbol that represents a pre-assembly time label for conditional assembly processing and from which ordinary symbols can be generated tc create assemkly-time labels.

* lanquage:

1. A set of representations, conventions, and rules used to convey infcrmation.
2. See machine language, object language, source language.

LE: (less than or equal to) See relational operator.
*length: See word length.
length attrikute (L'): An attribute that gives the number of bytes to be cccupied by the object code for the data represented, such as machine instructions, constants, or areas.
length field: The operand entry cr subentry in machine instructions that specifies the number of bytes at a specific address that are affected by the execution of the instruction.
length modifier: A subfield in the oferand of the DS or DC assembler instruction that determines the length in bytes of the area to be reserved or of the area intc which the data defined is to be assembled.
*level: The degree of subordination in a hierarchy.
library macro definition: A macro definition stored in a program library. The IBM-supplied supervisor and data management macro definitions (such as those called ky GET or PUT) are examples of library macro definitions. A library macro definition can be included at the beginning of a source module: it then becomes a source macro definition.

* linkage: In programming, coding that connects two separately coded routines.
linkage editor: A processing program that prepares the output of language translators for execution. It combines separately produced object or load mcdules; resolves symbolic cross references among them; replaces, deletes, and adds control sections, and generates overlay structures on request; and produces executakle code (a load module) that is ready to be fetched into virtual storage.
linkaqe editor control statement: An instruction for the linkage editor.
literal: A symbol or a quantity in a source program that is itself data, rather than a reference to data. Contrast with figurative constant.
literal pool: An area in storage into which the assembler assembles the values of the literals specified in a source module.
* load: In programming, to enter data into storage or working registers.
load module: The output of the linkage editor; a program in a format suitable for loading into virtual storage for execution.
loader:

1. Under the Operating System (OS). a processing program that combines the basic editing and loading functions of the linkage editor and program fetch in one job step. It accepts object modules and load modules created by the linkage editor and generates executable code directly in virtual storage. The loader does not produce load modules for program libraries.
2. Under the Lisk Operating System ([OS), a supervisor routine that retrieves program phases from the core image library and loads them into virtual storage.
local scope: Fertaining to that part of an assembler program that is either the kody of any macro definition called from a source module or the open code portion of the source module. Contrast with global scope.
local variable symbol:
3. A variable symbol that can be used to communicate values inside a macro definition or in the open code portion of a source module.
4. Contrast with global variable symbol.
'location: Any place in which data may be stored.
location counter: A counter whose value indicates the address of data assembled from a machine instruction or a constant, or the address of an area of reserved storage, relative to the beginning of a control section.
logic shift: A shift that affects all positions.
logical expression: A conditional assembly expression that is combination of lcgical terms, logical operators, and paired parentheses.
logical operator: In assembler programming, an operator or pair of operators that can be used in a logical expression tc indicate the action to be performed on the terms in the expression. The logical operatcrs allowed are: AND, OR, NOT, AND NOT, and OR NOT.
logical relation:
5. A logical term in which two expressions are separated by a relational operator. The relational operators allowed are: EQ, GE, GT, LE, LT, and NE.
6. See arithmetic relation, character relation.
logical term: A term that can be used only in a logical expression.

100p:

1. A sequence of instructions that is executed repeatedly until a terminal condition prevails.
2. See loop counter.
loop counter: In assembler prograrming, a counter to prevent excessive lcoping during conditional assembly processing.

LT: (less than) See relational operatcr.
machine code: An operation code that a machine is designed to recognize.
machine instruction:

1. An instruction that a machine can recognize and execute.
2. In assembler programming, (locsely) the symbolic machine language statements which the assembler translates into machine language instructions.
$\star$ machine lanquage: A language that is used directly by a machine.

## macro:

1. Loosely, a macro definition.
2. See also macro definition, macro generation, racro instruction, macro prototype statement.
macro call: Same as macro instruction.
macro definition: A set of assembler language statements that defines the name of, format of, and conditions for generating a sequence of assembler language statements from a single source statement.

* macro expansion: Same as macro generation.
macro generation: An operation in which the assembler produces a sequence of assembler language statements by processing a macro definition called by a macro instruction. Macro generation takes place at
pre-assembly time. Synonymous with macro expansion.


## macro_instruction:

1. An instruction in a source language that is equivalent to a specified sequence of machine instructions.
2. In assembler programming, an assemkler language statement that causes the assembler to process a predefined set of statements (called a macro definition). The staterrents normally produced from the macro definition replace the macro instruction in the source program. Synonymous with macro call.
macro instruction operand: An operand that supplies a value to be assigned to the corresponding symbolic parameter of the macro definition called by the macro instruction. This value is passed into the macro definition to be used in its processing.
macro library: See program library.
macro processing instruction: An assembler instruction that is used inside macro definitions and processed at pre-assemkly time. These instructions are: MACRO, MENL, MEXIT, and MNOTE.
macro prototype: Same as macro prototype statement.
macro prototype statement: An assemkler language statement that is used to give a name to a macro definition and to provide a model (prototype) for the macro instruction that is to call the macro definition.
main storage:
$\star$ 1. The general purpose storage cf a computer.
Usually, main storage can be accessed directly ky the operating registers.
3. See-also real storage, virtual storage.

* mask: A pattern of characters that is used to control the retention or elimination of portions of another pattern cf characters.
mnemonic operation code: An operation code consisting of mnemonic symbols that
indicate the nature of the operation to be performed, the type of data used, or the format of the instruction performing the operation.
mnemonic symbol:
* 1. A symbol chosen to assist the human memory, for example, an abbreviation such as "mpy" for "multiply".

2. See also mnemonic operation ccde.
model statement: A statement in the body of a macro definition or in open code from which an assembler language statement can be generated at pre-assembly time. Values can be sutstituted at one or more points in a model statement; one or more identical or different statements can be generated from the same model statement under the control of a conditional assembly loop.
$\star$ module:
3. A program unit that is discrete and identifiable with respect to compiling, combining with other units, and loading, for example, the infut to, or output from, an assembler, compiler, linkage editor, or executive rcutine.
4. See load module, object module, source module.
name :
5. A 1- to 8-character alphameric term that identifies a data set, a control statement, an instruction statement, a program, or a cataloged procedure. The first character of the name must be alphabetic.
6. See entry name, external name.
7. See also name entry, label.
name entry: Usually syncnymous with label (2). However, the name entry of a model statement can ke any string of characters at pre-assembly time.
name field parameter: A symbolic parameter that is declared in the name field of a macro prototype staterent. It is assigned a
value from the entry in the name field of the macro instruction that corresponds to the macro prototype statement.

NE: (not equal to) See relational operator

* nest: To imbed subroutines or data in othe sukroutines or data at a different hierarchical level such that the different levels of routines or data can be executed or accessed recursively.
nesting level: In assembler programming, the level at which a term (or
subexpression) appears in an expression, o the level at which a macro definition containing an inner macro instruction is processed by the assembler.
no OP: An instruction that specifically instructs the computer to do nothing. except to proceed to the next instruction in sequence.

NOT: A logic operator having the property that if $p$ is a statement, then the NOT of is true if $P$ is false, false if $F$ is true.

* null_character: A control character that serves to accomplish media fill or time fill, for example, in ASCII the all zeros character (not numeric zero). Null characters may be inserted into or removed from a sequence of characters without affecting the meaning of the sequence, but control of equipment or the format may ke affected. Abbreviated NUL. Contrast with space character.
null_character string: Same as null string null string:
* 1. The notion of a string depleted of its entities, or the notion of a string prior to establishing its entities.

2. In assembler programming, synonymous with the null character string.
number attribute (N'):
3. An attribute of a symbolic parameter that gives the number of suklist entries in the corresponding macro instruction operand.
4. An attribute that gives the number of positional operands in a macro instruction (sfecified as N'ESYSLIST) or an attribute that gives the number of sublist entries in a specific positional operand (specified as N'ESYSLIST (n)).
$\star$ object code: Output from an assembler whic is itself executable machine code or is
suitable for processing to produce executable machine code.

* okject lanquage: The language to which a statement is translated. The machine language for the IBM System/370 is an cbject language.
* object module: A module that is the output of an assembler or compiler and is input to a linkage editor.
* object program: A fully compiled cr assembled program that is ready tc be loaded into the computer. Contrast with source program.
open code: That portion of a source module that lies outside of and after any source macro definitions that may be specified.
ogen subroutine: A subroutine that is inserted into a routine at each place it is used. Contrast with closed subroutine. NOTE: In assembler programming, a macro definition is an open subroutine, kecause the statements generated from the definition are inserted into the source module at the point of call.
operand:

1. That which is operated upon.
2. See keyword operand, positional operand.

* oferating system: Software which controls the execution of computer programs and which may provide scheduling, debugging, input/output control, accounting, compilation. storage assignment, data management, and related services.
* operation code: A code that represents specific operations.


## * operator:

1. In the description of a process, that which indicates the action to be performed on the operands. NOTE: In assembler programming, operands are referred to as terms.
2. See arithmetic operator, binary operator, logical operator, unary operator.
3. See also concatenation character.
$\star$ OR: A logic operator having the property that if $P$ is a statement, $Q$ is a statement, $R$ is a statement,...., then the OR of $P, 2$, R... is true if at least one statement is true, false if all statements are false.
ordinary symbol: A symbol that represents an assembly-time value when used in the name or operand field of an instruction in
the assembler language. Ordinary symbols are also used to represent operation codes for assembler language instructions. An ordinary symbol has one alphabetic character followed by zero to seven alphameric characters.
outer macro instruction: A macro instruction that is specified in open code. Contrast with inner macro instruction.

* overflow: That portion of the result of an operation that exceeds the capacity of the intended unit of storage.
* overlay: The technique of repeatedly using the same blocks of internal storage during different stages of a program. When one routine is no longer needed in storage, another routine can replace all part of it.
* padding: A technique used to fill a klock with dummy data.
paired parentheses: A left parenthesis and a right parenthesis that belong to the same level of nesting in an expression; the left parenthesis must appear before its matching right parenthesis. If parentheses are nested within paired parentheses, the nested parentheses must be paired.
paired relocatable terms: T'wo relocatable terms in an expression with the same relocatability attribute that have different signs after all unary operations have keen performed. Paired relocatable terms have an absolute value.
parameter:

1. A variable that is given a constant value for a specific purpose or process.
2. See keyword parameter, name field parameter, positional parameter, symbolic parameter.
point of substitution: Any place in an assembler language statement, particularly a model statement, into which values can be substituted at pre-assembly time. Variakle symbols represent points of sukstitution.
positional operand: An operand in a macro instruction that assigns a value to the corresponding positional parameter declared in the prototype statement of the called macro definition.
positional parameter: A symbolic parameter that occupies a fixed position relative to the other positional parameters declared in the same macro prototype statement.
pre-assembly time: The time at which the assembler process macro definitions and performs conditional assembly operations.
private code: An unnamed control section.
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* proqram:
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1. A series of actions propcsed in order to achieve a certain result.
2. Loosely, a routine.
3. To design, write, and test a program as in (1).
4. Loosely, to write a routine.
5. See computer program, object program, source program.
program fetch time:
6. The time at which a program (in the form of load modules or chases) is loaded into virtual storage for execution.
7. See also fetch (2), fetch (3).

* Erogram library: A collection of available computer programs and routines.
programmer macro definition: Locsely, a source macro definition.
prototype statement: Sare as macrc prototype statement.
* pushdown list: A list that is constructed and maintained so that the next item to be retrieved and removed is the most recently stored item in the list, that is, last in, first out. Synonymous with pushdown stack.
pushdown stack: Same as pushdown list.
quoted string: A character string enclosed by apostrophes that is used in a macro instruction operand to represent a value that can include tlanks. The enclcsing apostrophes are part of the value represented. Contrast with character expression.
* radix: In positional representation, that integer, if it exists, by which the significance of the digit place must be multiplied to give the significance of the next higher digit place. For example, in decimal notation, the radix of each place is ten.
* radix complement: A complement obtained by suctracting each digit from one less than its radix, then adding one to the least significant digit, executing all carries
required. For example, tens complement in decimal notation, twos complement in binary notation.
read-only: A type of access to data that allows it to read but not modified.
real storage: The storage of a IBM
System/370 computer from which the central processing unit can directly obtain instructions and data and to which it can directly return results. Real storage can occupy all or part of main storage. Contrast with virtual storage.
recursive: Fertaining to a process in whicl each step makes use of the results of earlier steps.
NOTE: In assembler programming, the inner macro instruction that calls the macro definition within which it is nested performs a recursive call.
$\star$ register:

1. A device capable of storing a specifiec amount of data such as one word.
2. See base register, index register.
relation: The comparison of two expressions to see if the value cf one is equal to, less than, or greater than the value of the other.
relational operator: An operator that can be used in an arithmetic or character relation to indicate the comparison to be performed between the terms in the relation. The relational cperators are: EQ (equal), GE (greater than or equal to), GT (greater than), LE (less to or equal to). LT (less than). NE (not equal to).

* relative address: The number that specifies the difference between the absolute address and the base address. Synonymous with displacement.
relocatability attribute: An attribute that identifies the control section to which a relocatable expression belongs. Two relocatable expressions have the same relocatability attribute if the unpaired term in each of them belongs to the same control section.
relocatable expression: An assembly-time expression whose value is affected by program relocation. A relocatable expression can represent a relocatable address.
relocatable term: A term whose value is affected by program relocation.
* relocate: In computer programming, to move a routine from one portion of storage to another and to adjust the necessary address references so that the routine, in its new location, can ke executed.
relocation: The modification cf address constants to compensate for a change in origin of a module, program, or control section.
*rounding: Same as roundoff.
roundoff: To delete the least significant digit or digits of a numeral and to adjust the part retained in accordance with some rule.
*routine:

1. An ordered set of instructions that may have some general or frequent use.
2. See sukroutine.
scale modifier: A subfield in the cperand of the DC assembler instruction that indicates the number of digits in the object code to be occupied by the fractional portion of a fixed-point or floating-point constant.

Scaling attribute: An attribute that indicates the number of digit positions occupied by the fractional portion of fised-point, decimal, and floating-point constants in their object code form.
scope:

1. In assembler programming, that part of a source program in which a variable symbol can communicate its value.
2. See global scope, local scope.
self-defining term: An absolute term whose value is implicit in the specification of the term itself.

Sequence symbol: A symbol used as a branching lakel for conditional assembly instructions. It consists of a period followed by one to seven alphameric characters, the first of which must be alphatetic.

SET symbol: A variable symbol used tc communicate values during conditicnal assembly processing. It must be declared to have either a global or local scope.
severity code: A code assigned by the assembler to an error detected in a source module. A severity code can also be specified and assigned tc an error message generated by the MNOTE instruction.

* sign kit: A binary digit occupying the sign position.
sign position: A position, normally located at one end of a numeral, that contains an indication of the algebraic sign of the number.
* Siqnificant diqit: A digit that is needed for a certain purpose, particularly one that must be kept to preserve a specific accuracy or precision.
* source language: The language from which a statement is translated.
source macro definition: A macro definition included in a source module. A source macro definition can be entered into a program library; it then becomes a library macro definition.
source module: A sequence cf statements in the assembler language that constitutes the input to a single execution of the assembler.
* Source program: A computer program written in a source language. Contrast with object program.
* space_character: A normally nonprinting graphic character used to separate words. Synonymous with blank character. Contrast with null character.
* special_character: A graphic character that is neither a letter, nor a digit, nor a space character.
* statement:

1. In computer programming, a meaningful expression or generalized instruction in a source language.
2. See job control statement, linkage editor control statement, comments statement, model statement.

* storage:

1. Fertaining to a device into which data can be entered, in which they can be held, and from which they can ke retrieved at a later time.
2. Loosely, any device that can store data.
3. See main storage, real storage, virtual storage.
storage_allocaticn:
4. The assignment of blocks of data to specified blocks of storage.
5. See dynamic stcrage allocation.

* storaqe protection: An arrangement for preventing access to storage for either reading, or writing, or both.
storage stack: Loosely, a pushdown list.


## $\star$ string:

1. A linear sequence of entities such as characters or physical elements.
2. See bit string, character string, null string.
sublist: A macro instruction operand that contains one or more entries separated by commas and enclosed in parentheses.
$\star$ subroutine:
3. A routine that can be part of another routine.
4. See closed sukroutine, ofen subroutine.
subscript: One or more elements, enclosed in parentheses, that appear immediately after a variakle symbol or character expression. The value of a subscript indicates a position in the array or string of values represented by the variable symbol or character expressicn.
subscripted ESYSLIST: The system variable symbol ESYSLIST immediately followed by either one sutscript or two subscripts separated ky commas, and enclosed in parentheses. The value of the first sutscript indicates the position cf a positional operand in a macro instruction and the value of the second subscript indicates the position of the entry in the sublist of the positional operand indicated by the first sukscript.

## subscripted SET symbol:

1. A SET symbol that is immediately followed by a subscript. A subscripted SET symbol must be declared with an allowable dimension before it can be used. The value of the subscrift indicates the position of the value given to the subscripted symbcl in the array represented by the SET symbol.
2. See also dimensioned SET symbcl.
subscripted symbolic parameter: A symbolic parameter that is immediately followed by a subscript. The value of the subscript indicates the position of the entry in the sublist in the macro instruction cperand referred to by the symbolic parameter.
substitution: The action taken by the assembler when it replaces a variable symbol with a value, for example, during the expansion of a racro definition.

## substring:

1. A character string that has keen extracted from a character expression. 2. See also substring notation.
sukstring notation: A character expression immediately followed by two subscripts, separated by a comma, and enclosed in parentheses. It can be used only in conditional assembly instructions. The value of the first subscript indicates the position of the character within the character expression that begins the substring. The value of the second subscript represents the numker of characters to be extracted from the character expression.

* Switch: A device or programming technique for making a selection, for example, a conditional jump.
* Symbol:

1. A representation of something ky reason of relationship, association, or convention.
2. See mnemonic symbol, crdinary symbol, sequence symbol, SET symbol, variakle symbol.

* symbolic address: An address expressed in symbols convenient to the computer programmer.
* Symbolic coding: Coding that uses machine instructions with symbolic addresses. NCTE: In assembler programming, any instruction can contain symbolic addresses. In addition, any other portion of an instruction may be represented with symbols, for example, lakels, registers, lengths and immediate data.


## symbolic parameter:

1. A variabie symbol declared in the prototype statement of a macro definition. A symbolic parameter is usually assigned a value from the corresponding cFerand in the macro instruction that calls the macro definition.
2. See also keyword parameter, name field parameter, positional parameter.
system loader: See loader (2).
system macro definition: Loosely, a library macro definition supplied by IBM.
system_macro instruction: Loosely, a macro instruction that calls for the processing of an IBM-supplied library macro definition, for example, the ATTACH macro.
system variable symbol: A variable symbol that always begins with the characters

ESYS. The system variable symbols do not have to ke declared, because the assembler assigns them read-only values autcratically according to specific rules.
term:

1. The smallest part of an expression that can te assigned a value.
2. See absolute term, arithmetic term, logical term, relocatable term.
*translate: To transform statements from one language to another without significantly changing the meaning.
*truncate: To terminate a computaticnal process in accordance with sorre rule, for example, to end the evaluation of a power series at a specified term.
NOTE: In assembler programming, the object code for data constants can be truncated by the assembler.
*twos complement: The radix complement in binary notation.
type attribute (T): An attribute that disting uishes one form of data from another, for example, fixed-pcint constants from floating-point constants or machine instructions from macro instructicns.
unary operator: An arithretic operator having only one term. The unary oper ators that can be used in absolute cr relocatable, and arithmetic expressions are: positive (+) and negative ( - ).
unnamed control section: A control section that is initiated in one of the following three ways:
3. Ey an unnamed START instructicn.
4. By an unnamed CSECT instruction, if no unnamed START instruction appears before the CSECT instruction.
5. By any instruction that affects the setting of the location ccunter.

* variable: A quantity that can assume any of a given set of values.
variable symbol: In assembler programming, a symbol, used in macro and conditional assembly processing, that can assume any of a given set of values. It consists of an ampersand ( 8 ) followed by one to seven alphameric characters, the first of which must be alphabetic.
NOTE: All variable symbols must be declared except the system variable symbols.
virtual_storage: Address space appearing to the user as real storage from which instructions and data are mapped into real storage locations. The size of virtual storage is limited only by the addressing scheme of the computing system rather than by the actual number of real storage locations. Contrast with real storage.
warning message: An indication that a possible error has been detected. The assembler does not assign a severity code to this type of error. Contrast with error message.
word:
* 1. A character string or bit string considered as an entity.
$\star$ 2. See computer word.

3. See doubleword, fullword, halfword.
$\star$ word length: A measure of the size of a word, usually specified in units such as characters or binary digits.
NOTE: In assembler programming, the word, or fullword, contains 32 bits (binary digits) or 4 bytes.
wrap-around: Loosely, the cverflow of the location counter when the value assigned to it exceeds 224-1

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## Appendix IV: Machine Instruction Mnemonic Codes

This appendix contains two tables of the mnemonic operation codes for all machine instructions that can be represented in assembler language, including extended mnemonic operation codes.

The first table is in alphabetic order by instruction. The second table is in numeric order by operation code.

In the first table is indicated: both the mnemonic and machine operation codes, explicit and implicit operand formats, program interruptions possible, and condition code set.

The column headings in the first table and the information each column provides follow:

Instruction: This column contains the name of the instruction associated with the mnemonic operation code.

Mnemonic Operation Code: This column contains the mnemonic operation code for the instruction. This is written in the operation field when coding the instruction.

Machine operation code: This column contains the hexadecimal equivalent of the actual machine operation code. The operation code will apppear in this form in most storage dumps and when displayed on the system control panel. For extended mnemonics, this column also contains the mnemonic code of the instruction from which the extended mnemonic is derived.

Operand Format: This column shows the symbolic format of the operand field in both explicit and implicit form. For both forms, R1, R2, and R3 indicate general registers in operand one, two, and three respectively. $X 2$ indicates a general register used as an index register in the second operand. Instructions which require an index register (X2) but are not to be indexed are shown with a 0 replacing $\mathrm{X} 2 . \mathrm{L}, \mathrm{L} 1$, and L 2 indicate lengths for either operand, operand one, or operand two respectively. m1 and m3 indicate a 4 -bit mask in operands one and three respectively. I, I2, and I3 indicate immediate data eight bits long (I and I2) or four bits long (I3).

For the explicit format, D1 and D2 indicate a displacement and B1 and B2 indicate a base register for operands one and two.

For the implicit format, D1, B1, and D2, B2 are replaced by S1, and S2 which indicate a storage address in operands one and two.

Type of Instruction: This column gives the basic machine format of the instruction (RR, RX, SI, or SS). If an instruction is included in a special feature or is an extended mnemonic, this is also indicated.

Proqram Interruptions Possible: This column indicates the possible program interruptions for this instruction. The abreviations used are: A - Addressing, $S$ - Specification, Ov Overflow, P - Protection, Op - Operation (if feature is not installed), and Other - other interruptions which are listed. The type of overflow is indicated by: D - Decimal, EExponent, or F - Fixed Point.

Condition code set: The condition codes set as a result of this instruction are indicated in this column. (See legend following the table.)

|  | Instruction | Mnemonic <br> Operation Code | Machine Operation Code | Operand Format |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | Explicit | Implicit |
|  | Add | A | 5A | R1, D2( $\times 2, \mathrm{B2}$ ) or R1, D2(, B2) | R1, S2(X2) or R1, 52 |
|  | Add | AR | 1A | R1,R2 |  |
|  | Add Decimal | AP | FA | D1(L1, B1), D2(L2, B2) | S1(L1), S2(L2) or S1, 52 |
|  | Add Halfword | AH | 4A | R1, D2 ( $\times 2, \mathrm{B2} 2$ ) r R1, D2 $2, \mathrm{B2} 2)$ | R1, S2 ( $\times 2$ )or R1, S2 |
|  | Add Logical | AL | 5E | R1, D2 ( $\times 2, \mathrm{B2}$ ) or R1, D2 (, B2) | R1, S2 ( $\times 2$ ) or R1, 52 |
|  | - Add Logical | ALR | $1 E$ | R1,R2 |  |
|  | Add Normalized, Extended Add Normalized, Long | $\begin{aligned} & A X R \\ & A D \end{aligned}$ | 36 64 | R1, R2 R1, D2 ( $22, \mathrm{~B} 2$ ) or R1, D2 (, B2) |  |
|  | Add Normalized, Long | AD ADR | $\begin{aligned} & 6 A \\ & 2 A \end{aligned}$ | R1, D2 $2 \times 2, \mathrm{~B} 2$ ) or R1, D2 $2, \mathrm{~B} 2)$ $\mathrm{R} 1, \mathrm{R} 2$ | R1, S2(X2) or R1, 52 |
|  | Add Normalized, Short | AE | 7A | R1, D2 ( $\times 2, \mathrm{~B} 2$ ) ${ }^{\text {R R1, }}$ D2 (, B2) | R1, 2 ( X 2 ) or R1, 52 |
|  | Add Normalized, Short | AER | 3A | R1, R2 | R1, S2(X2)or R1, ${ }^{\text {2 }}$ |
|  | Add Unnormalized, Long | AW | 6 E | R1, D2 (X2, B2) or R1, D2 (, B2) | R1, S2(X2)or R1, S2 |
|  | Add Unnormalized, Long | AWR | 2 E | R1, R2 | R1, S2(X2) or R1, ${ }^{\text {2 }}$ |
|  | Add Unnormalized, Short | AU | 7 F | R1, D2 ( $\times 2, \mathrm{B2}$ ) or R1, D2 (, B2) | R1, 22 (X2) or R1, 52 |
|  | Add Unnormalized, Short | AUR | 3E | R1, R2 | R1, $22 \times 2$ or R1, ${ }^{\text {a }}$ |
|  | And Logical | N | 54 | R1, D2 ( $\times 2, \mathrm{~B} 2$ ) or R1, D2 $(, \mathrm{B2}$ ) | R1, S2(X2) or R1, S2 |
|  | And Logical | NC | D4 | D1( $\mathrm{L}, \mathrm{Bl}$ ), D2(B2) | $\mathrm{Sl}(\mathrm{L}), \mathrm{S} 2$ or $\mathrm{S} 1, \mathrm{~S} 2$ |
|  | And Logical | NR | 14 | R1,R2 |  |
|  | And Logical Immediate | NI | 94 | $\mathrm{DI}(\mathrm{Bl}), 12$ | S1,12 |
|  | Branch and Link Branch and Link | BAL | 45 |  | R1, S2(X2)or R1, S2 |
|  | Branch and Save | ${ }_{\text {BALS }}$ | 4D | $\mathrm{R} 1, \mathrm{R} 2$ $\mathrm{R1}, \mathrm{D} 2(\mathrm{X} 2, \mathrm{B2})$ or $\mathrm{R1} 1, \mathrm{D} 2(, B 2)$ |  |
|  | Branch and Save | BASR | OD | $\begin{aligned} & \text { R1,D2(X2,B2) or R1,D2(,B2) } \\ & \text { R1,R2 } \end{aligned}$ | R1, S2(X2) or R1, 52 |
|  | Branch on Condition Branch on Condition | $\begin{aligned} & B C \\ & B C R \end{aligned}$ | 47 07 | $\begin{aligned} & \mathrm{M} 1, \mathrm{D} 2(X 2, B 2) \text { or } \mathrm{M1}, \mathrm{D} 2(, \mathrm{B2}) \\ & \mathrm{M1}, \mathrm{R2} 2 \end{aligned}$ | M1, $52(\mathrm{X} 2$ ) or $\mathrm{M1} 1,52$ |
|  | Branch on Count | ${ }_{\text {BCT }}$ | 46 |  | R1, 2 (X2) or R1, S2 |
|  | Branch on Count | ${ }^{\text {BCTTR }}$ |  | ${ }^{\mathrm{R} 1}$, R22 |  |
|  | Branch on Equal Branch on Equal | $\begin{aligned} & \text { BE } \\ & \text { BER } \end{aligned}$ | $\begin{aligned} & 47(B C 8) \\ & 07(B C R) \end{aligned}$ | $\mathrm{D} 2(\mathrm{X} 2, \mathrm{~B} 2) \text { or D2 } 2(, \mathrm{~B} 2)$ | S2(X2) or S2 |
|  | Branch on High | BH | 47(BC 2) | D2( $\times 2, \mathrm{B2}$ ) or D2(, B2) | S2(X2) or 52 |
|  | Branch on High | BHR | 07 (BCR 2) |  |  |
|  | Branch on Index Kigh <br> Branch on Index Low or Equal | $\underset{\text { BXLE }}{\text { BX }}$ |  | $\mathrm{R} 1, \mathrm{R3}$ $\mathrm{R1}, \mathrm{D} 3, \mathrm{D} 2(\mathrm{B2} 2)$ (82) |  |
|  | Branch on Low | ${ }_{B L}$ | 47(BC 4) | D2 2 ( 2 2, B2) or D2 2 ( $\mathrm{B2}$ ) | S2 ( X 2 ) or 52 |
|  | Branch on Low | BLR | 07 (BCR 4) |  |  |
|  | Branch if Mixed Branch if Mixed | $\begin{aligned} & B M \\ & B M R \end{aligned}$ | $\begin{aligned} & 47(B C 4) \\ & 07(B C R 4) \end{aligned}$ | $\begin{aligned} & \mathrm{D} 2(\mathrm{X} 2, \mathrm{~B} 2) \text { or D2(, B2) } \\ & \mathrm{R} 2 \end{aligned}$ | S2(X2) or S2 |
|  | Branch on Minus | BM | 47(BC 4) | D2( $\times 2, \mathrm{B2}$ ) or D2(, $\mathrm{B2}$ ) | S2(X2) or S2 |
|  | Branch on Minus | BMR | $07(B C R 4)$ |  |  |
|  | Branch on Not Equal Branch on Not Equal | BNE BNER | 47(BC 7) 07 ( $B C R 7$ ) | $\mathrm{R}_{\mathrm{R} 2}(\mathrm{X} 2, \mathrm{~B} 2)$ or D2(, B2) | S2(X2) or S2 |
|  | Branch on Not High | BNH | 47(BC 13) | D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) or D2 2 (, B2) | S2(X2) or S2 |
|  | Branch on Not High | BNHR | 07(BCR 13) |  |  |
|  | Branch on Not Low | BNL | $47(\mathrm{BC} 11)$ | $\mathrm{D} 2(\mathrm{X} 2, \mathrm{~B} 2)$ or D2 $(, \mathrm{B2} 2)$ | S2(X2) or S2 |
|  | Branch on Not Low Branch if Not Mixed | BNLR | 07(BCR 11) $47(B C$ 11) | $\mathrm{D} 2(\mathrm{X} 2, \mathrm{~B} 2) \text { or } \mathrm{D} 21, \mathrm{~B} 2)$ | (X2) or S2 |
|  | Branch if Not Mixed | BNMR | 07(BCR 11) |  |  |
|  | Branch on Not Minus Branch on Not Minus | BNM BNMR | $47(B C 11)$ $07(B C R 11)$ | D2( $\mathrm{R} 2, \mathrm{B2}$ ) or D2(, B2) | $52(\mathrm{X} 2)$ or 52 |
|  | Branch on Not Minus | BNMR | 07(BCR 11) |  |  |
|  | Branch if Not Ones | BNO | 47(BC 14) | $\mathrm{D} 2(\mathrm{X} 2, \mathrm{B2}$ ) or D2(, B2) | S2(X2) or S2 |
|  | Branch if Not Ones | BNOR | $07(B C R 14)$ |  |  |
|  | Branch on No Overflow Branch on No Overflow | BNO BNOR | 471BC 14) $07(B C R$ | $\mathrm{D} 2(\times 2, B 2)$ or $\mathrm{D} 21, \mathrm{~B} 2)$ | S2(X2) or S2 |
|  | Branch on Not Plus | BNP | $47\left(\mathrm{BC} \mathrm{l}^{\text {13) }}\right.$ | ${ }_{\text {D2 }}(\times 2, B 2)$ or D2 2 (,B2) | S2(X2) or S2 |
|  | Branch on Not Plus | BNPR | 07 (BCR 13) |  | S2(x2) or 52 |
|  | Branch if Not Zeros | BNZ | $47(\mathrm{BC} 7)$ | $\mathrm{D} 2(\mathrm{X} 2, \mathrm{~B} 2)$ or $\mathrm{D} 2(, \mathrm{~B} 2)$ | S2(X2) or S2 |
|  | Branch if Not Zeros Branch on Not Zero | BNZR BNZ | 07(BCR 7) $47(\mathrm{BC} 7$ ) | ${ }_{\text {R2 }}^{\text {D2 }}$ ( $\times 2, B 2$ or D2 $2(, B 2)$ | S2(X2) or S2 |
|  | Branch on Not Zero | BNZR | 07 (BCR 7) |  | ) |
|  | Branch if Ones | BO | 47 (BC 1) | D2( $\mathrm{X} 2, \mathrm{B2} 2$ or D2(, B2) | S2(X2) or S 2 |
|  | Branch if Ones | BOR | $07(B C R 1)$ |  |  |
|  | Branch on Overflow | BO | $47(B C 1)$ | $\mathrm{D} 2(\mathrm{X} 2, \mathrm{~B} 2)$ or $\mathrm{D} 2(, \mathrm{~B} 2)$ | S2(X2) or S2 |
|  | Branch on Overfiow | BOR | 07(BCR 1) | R2 |  |
|  | Branch on Plus | BP | 47(BC 2) | $\mathrm{D} 2(\mathrm{X} 2, \mathrm{~B} 2)$ or $\mathrm{D} 2(, \mathrm{~B} 2)$ | S2(X2) or 52 |
|  | Branch on Plus | BPR | $07(\mathrm{BCR} 2)$ |  |  |
|  | Branch if Zeros | BZ | 47(BC 8) | D2( $\times 2, B 2$ or D2(, B2) | S2(X2) or 52 |
|  | Branch if Zeros Branch on Zero | ${ }^{\text {BZZR }}$ | $07(\mathrm{BCR} 8)$ 47 (BC 8) | R2 D2 $(\mathrm{X} 2, \mathrm{~B} 2)$ or D2 $2(, \mathrm{~B} 2)$ | S2(X2) or S2 |
|  | Branch on Zero | ${ }^{\text {BZZR}}$ | 07(BCR 8) | R2 2,82 or ${ }^{\text {2 }}$ (, 22 ) | S2( |
|  | Branch Unconditional Branch Unconditional | ${ }^{B} \mathrm{~B}$ | $47(B C 15)$ $07(B C R 15)$ | D2(X2, B2) or D2(, B2) | S2(X2) or S2 |
| \| DOS/VSE onlv |  |  |  |  |  |
|  | Clear Storage Poge | CLRP | B215 | D2(B2) | $\begin{aligned} & \$ 2 \\ & \text { S2 } \end{aligned}$ |
|  | Compare Algebraic | ${ }^{\text {c }}$ | 59 | R1, D2 ( $\mathrm{X} 2, \mathrm{B2}$ ) or R1, D2 (, B2) | R1, 2 (X2 or R1, 52 |
|  | Compare Algebraic | CR | 19 | R1,R2 |  |
|  | Compare and Swap | CS | BA | R1,R3, D2, (B2) | R1, R3, S2 |
|  | Compare Decimal | CP | F9 | D1(L1, B1), D2 (L2, B2) | S1(L1), $52(\mathrm{~L} 2$ ) or 51,52 |
|  | Compare Double and Swap | CDS | BB | R1, R3, D2(B2) | R1,R3, S2 |
|  | Compare Halfword | CH | 49 | R1, D2 ( $22, \mathrm{~B} 2)$ or R1, D2 $2, \mathrm{~B} 2)$ | R1, S2 (X2) or R1, S2 |
|  | Compare Logical | CL | 55 | R1, D2 ( $\times 2, \mathrm{~B} 2$ )or R1, D2 (, B2) | R1, S2(X2) or R1, $\mathrm{S}^{2}$ |
|  | Compare Logical | CLC | D5 | D1( $\mathrm{L}, \mathrm{B1}$ ), D2(B2) | S1(L), S2 or S1, 52 |

\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|c|c|}
\hline \& \multirow[t]{2}{*}{Instruction} \& \multirow[t]{2}{*}{Type of Instruction Instruction} \& \multicolumn{5}{|l|}{Program Interruption Possible} \& \multicolumn{4}{|c|}{Condition Code Set} \\
\hline \& \& \& A S \& Or \& P \& Op \& Other \& 00 \& 01 \& 10 \& 11 \\
\hline \& \begin{tabular}{l}
Add \\
Add \\
Add Decimal \\
Add Halfword \\
Add Logical \\
Add Logical
\end{tabular} \& \[
\begin{aligned}
\& R X \\
\& R R \\
\& S S, \text { Decimal } \\
\& R X \\
\& R X \\
\& R R
\end{aligned}
\] \& \(\pm \times \mathrm{x}\) \& F \(\begin{aligned} \& \text { F } \\ \& \text { F } \\ \& \text { D } \\ \& F \\ \& F\end{aligned}\) \& \(\times\) \& \(\times\) \& Datic \& Sum \(=0\)
Sum \(=0\)
Sum \(=0\)
Sum \(=0\)
Sum=0 \((\oplus)\)
Sum \(=0\) ( \()\) \& Sum<0
Sum \(<0\)
Sum \(<0\)
Sum \(<0\)
Sum \(0 \oplus(H)\)
Sum=0
l \& \[
\begin{aligned}
\& \text { Sum }>0 \\
\& \text { Sum }>0 \\
\& \text { Sum }>0 \\
\& \text { Sum }>0 \\
\& \text { Sum }=0 \text { (1) } \\
\& \text { Sum }=0 \text { (1) }
\end{aligned}
\] \& Overflow Overflow Overflow Overflow Sum 0 (1) Sum 0 (1) \\
\hline \& \begin{tabular}{l}
Add Normalized, Extended \\
Add Normalized, Long \\
Add Normalized, Long \\
Add Normalized, Short \\
Add Normalized, Short
\end{tabular} \& RR,Floating Pt. RX, Floating Pt . RR, Flooting Pr. RX, Floating Pt. RR, Floating Pt. \& ( \(\begin{aligned} \& x \\ \& x \\ \& x \\ \& x \\ \& x \\ \& x \\ \& x \\ \& x\end{aligned}\) \& \(\begin{array}{r}\text { P } \\ \times \\ \times \\ \hline \\ \hline\end{array}\) \& \& \(\times\)
\(\times\)
\(\times\)
\(\times\)
\(\times\)
\(\times\)
\(\times\)
\(\times\)
\(\times\) \& B, C
B,
B,
B,
B,
B \& \[
\begin{aligned}
\& R \\
\& R \\
\& R \\
\& R \\
\& R \\
\& R
\end{aligned}
\] \& L
L
L
L
L \& \[
\begin{aligned}
\& M \\
\& M \\
\& M \\
\& M \\
\& M
\end{aligned}
\] \& \\
\hline \& Add Unnormalized, Long Add Unnormalized, Long Add Unnormalized, Short Add Unnormalized, Short \& RX, Floating Pt . RR, Flooting Pr. RX, Floating Pt. RR, Floating Pt. \&  \& E \& \& \[
\begin{aligned}
\& \mathbf{x} \\
\& \mathbf{x} \\
\& \mathbf{x} \\
\& \mathbf{x}
\end{aligned}
\] \& \[
\begin{aligned}
\& C \\
\& C \\
\& C \\
\& C
\end{aligned}
\] \& \[
\begin{aligned}
\& R \\
\& R \\
\& R \\
\& R \\
\& R
\end{aligned}
\] \& \(L\)
\(L\)
\(L\)
\(L\) \& \[
\begin{aligned}
\& M \\
\& M \\
\& M \\
\& M
\end{aligned}
\] \& \\
\hline \& \begin{tabular}{l}
And Logical \\
And Logical \\
And Logical \\
And Logical Immediate
\end{tabular} \& \[
\begin{aligned}
\& \text { RX } \\
\& \text { SS } \\
\& \text { RR } \\
\& \text { SI }
\end{aligned}
\] \& \(x\)
\(\times\)
\(\times\)

$\times$ \& \& $\times$ \& \& \& \[
$$
\begin{aligned}
& \mathrm{J} \\
& \mathrm{~J} \\
& \mathrm{~J} \\
& \mathrm{~J}
\end{aligned}
$$

\] \& \[

$$
\begin{aligned}
& K \\
& K \\
& K \\
& K
\end{aligned}
$$
\] \& \& <br>

\hline \& Branch and Link Branch and Link Branch and Save Branch and Save \& $$
\begin{aligned}
& R X \\
& R R \\
& R X \\
& R R
\end{aligned}
$$ \& \& \& \& \[

$$
\begin{aligned}
& x \\
& x
\end{aligned}
$$

\] \& \& \[

$$
\begin{aligned}
& N \\
& N \\
& N \\
& N
\end{aligned}
$$

\] \& \[

$$
\begin{aligned}
& N \\
& N \\
& N \\
& N
\end{aligned}
$$

\] \& \[

$$
\begin{aligned}
& N \\
& N \\
& N \\
& N
\end{aligned}
$$

\] \& \[

$$
\begin{aligned}
& N \\
& N \\
& N \\
& N
\end{aligned}
$$
\] <br>

\hline \& Branch on Condition Branch on Condition Branch on Count Branch on Count Branch on Equal Branch on Equal \& RX
RR
$R X$
RR
RX, Ext. Mnemonic

RR, Ext.Mnemonic \& \& \& \& \& \& $$
\begin{aligned}
& N \\
& N \\
& N \\
& N \\
& N \\
& N \\
& N
\end{aligned}
$$ \& N

N
N
N
N
N \&  \& N
N
N
N
N
N <br>

\hline \& | Branch on High |
| :--- |
| Branch on High |
| Branch on Index High |
| Branch on Index Low or Equal |
| Branch on Low |
| Branch on Low |
| Branch if Mixed |
| Branch if Mixed | \& RX, Ext.Mnemonic RR, Ext. Mnemonic RS RS RX, Ext.Mnemonic RR, Ext. Mnemonic RX, Ext.Mnemonic RR, Ext. Mnemonic \& \& \& \& \& \& \[

$$
\begin{aligned}
& N \\
& N \\
& N \\
& N \\
& N \\
& N \\
& N \\
& N \\
& N
\end{aligned}
$$

\] \& | N |
| :--- |
| N |
| N |
| N |
| N |
| N |
| N N | \& N

N
N
N
N
N
N

N \& $$
\begin{aligned}
& N \\
& N \\
& N \\
& N \\
& N \\
& N \\
& N \\
& N \\
& N
\end{aligned}
$$ <br>

\hline 1 \& | Branch on Minus |
| :--- |
| Branch on Minus |
| Branch on Not Equal |
| Branch on Not Equal |
| Branch on Not High |
| Branch on Nor High |
| Branch on Not Low |
| Branch on Not Low |
| Branch if Not Mixed |
| Branch if Not Mixed |
| Branch on Not Minus |
| Branch on Not Minus | \& RX, Ext.Mnemonic RR, Ext.Mnemonic RX, Ext.Mnemonic RR, Ext.Mnemonic RX, Ext.Mnemonic RR, Ext. Mnemonic RX, Ext.Mnemonic RR, Ext.Mnemonic RX, Ext.Mnemonic RR,Ext.Mnemonic RX, Ext.Mnemonic RR, Ext.Mnemonic \& \& \& \& \& \& $N$

$N$
$N$
$N$
$N$
$N$
$N$
$N$
$N$
$N$
$N$
$N$ \& N

N \& | N |
| :--- |
| N |
| N |
| N |
| N |
| N |
| N |
| N |
| N |
| N |
| N N | \& N

N
N
N
N
N
N
N
N
N
N
N <br>

\hline 1 \& | Branch if Not Ones |
| :--- |
| Branch if Not Ones |
| Branch on No Overflow |
| Branch on No Overflow |
| Branch on Not Plus |
| Branch on Not Plus |
| Branch if Not Zeros |
| Branch if Not Zeros |
| Branch on Not Zeros |
| Branch on Not Zeros |
| Branch if Ones |
| Bronch if Ones. |
| Branch on Overflow |
| Branch on Overflow | \& RX, Ext.Mnemonic RR, Ext. Mnemonic RX, Ext.Mnemonic RR.Ext.Mnemonic RX, Ext. Mnemonic RR, Ext.Mnemonic RX, Ext.Mnemonic RR, Ext.Mnemonic RX, Ext.Mnemonic RR, Ext.Mnemonic RX, Ext.Mnemonic RR, Ext.Mnemonic RX, Ext. Mnemonic RR, Ext.Mnemonic \& \& \& \& \& \& \[

z z Z Z Z Z Z Z Z Z Z Z Z Z
\] \& N

N
N
N
N
N
N
N
N
N
N
N
N

N \& \[
$$
\begin{aligned}
& \mathrm{N} \\
& \mathrm{~N} \\
& \mathrm{~N} \\
& \mathrm{~N} \\
& \mathrm{~N} \\
& \mathrm{~N} \\
& \mathrm{~N} \\
& \mathrm{~N} \\
& \mathrm{~N} \\
& \mathrm{~N} \\
& \mathrm{~N} \\
& \mathrm{~N} \\
& \mathrm{~N}
\end{aligned}
$$

\] \& | N |
| :--- |
| N |
| N |
| N |
| $N$ |
| N |
| N N |
| N N |
| N |
| N |
| N |
| N |
| N | <br>


\hline \& | Branch on Plus |
| :--- |
| Branch on Plus |
| Branch if Zeros |
| Branch if Zeros |
| Branch on Zero |
| Branch on Zero |
| Branch Unconditional |
| Branch Unconditional | \& RX, Ext.Mnemonic RR, Ext.Mnemonic RX, Ext.Mnemonic RR, Ext.Mnemonic RX, Ext.Mnemonic RR, Ext. Mnemonic RX, Ext.Mnemonic RR, Ext.Mnemonic \& \& \& \& \& \& \[

$$
\begin{aligned}
& N \\
& N \\
& N \\
& N \\
& N \\
& N \\
& N \\
& N
\end{aligned}
$$

\] \& | N |
| :--- |
| N |
| N |
| N |
| N |
| N |
| N |
| N | \& N

N
N
N
N
N
N
N \& N
N
N
N
N
N
N
N <br>

\hline | DOSIVSE only \& | Clear 1/O |
| :--- |
| Clear Storage Page |
| Compare Algebraic | \& \[

$$
\begin{aligned}
& S^{s} \\
& S \\
& R X
\end{aligned}
$$

\] \& \[

$$
\begin{array}{|l|l}
x & \\
x & x
\end{array}
$$

\] \& \& $\times$ \& \& A ${ }_{\text {A,GB }}$ \& \[

$$
\begin{aligned}
& A A X \\
& Z
\end{aligned}
$$

\] \& | CC |
| :--- |
| AA | \& EE

BB \& KK <br>
\hline \& Compare Algebraic \& RR \& \& \& \& \& \& z \& AA \& BB \& <br>
\hline \& Compare and Swap \& RS \& $\times \times$ \& \& $x$ \& $\times$ \& \& z \& AAW \& \& <br>
\hline \& Compare Decimal \& SS, Decimal \& \& \& \& $\times$ \& Data \& $z$ \& AA \& BB \& <br>
\hline \& Compare Double and Swap
Compare Halfword \& RS \&  \& \& $\times$ \& $\times$ \& \& z \& AAW
AA \& BB \& <br>
\hline \& \& \& \& \& \& \& \& z \& AA \& BB \& <br>
\hline \& Compare Logical \& \& $\times \times$ \& \& \& \& \& z \& AA \& BB \& <br>
\hline
\end{tabular}

|  | Instruction | $\begin{aligned} & \text { Mnemonic } \\ & \text { Operation } \\ & \text { Code } \\ & \hline \end{aligned}$ |  | Explicit ${ }^{\text {Operand F }}$ | Implicit |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| \| Dosivse only | Compare Logical | CLR | 15 | R1,R2 |  |  |
|  | Compare Logical Characters under Mask | CLM | BD | R1, M3, D2(B2) | R1,M3, S2 |  |
|  | Compare Logical Immediate | CLI | 95 | D1 (B1),12 | S1,12 |  |
|  | Compare Logical Long | CLCL | OF | R1, R2 |  |  |
|  | Compare, Long Compare, Long | CD CDR | 69 29 | R1, D2(X2, B2) or R1, D2(, B2) | R1, S2(X2) or R1, S2 |  |
|  | Compare, Short | CE | 79 | R1, D2( $\times 2, \mathrm{~B} 2$ ) or R1, D2 (, B2) | R1, S2(X2)or R1, S2 |  |
|  | Compare, Short | CER | 39 | R1,R2 |  |  |
|  | Connect Page . | CTP | B0 | 81, D2 (B2) | R1, 22 |  |
|  | Convert to Binary Convert to Decimal | CVB CVD | 4F |  | $\begin{aligned} & \mathrm{R1}, \mathrm{~S} 2(\times 2) \text { ) R1, S2 } \\ & \mathrm{R1}, \mathrm{~S} 2(\times 2) \text { lor R1, } 22 \end{aligned}$ |  |
| $\left\|\begin{array}{l} \text { DOS/VSE only } \\ \text { DOS/VSE only } \end{array}\right\|$ |  |  |  |  |  |  |
|  | Deconfigure Page | DEP | B21B | D2(B2) | S2 |  |
|  | Disconnect Page | DCTP | B21C |  |  |  |
|  | 俍 ${ }^{\text {Divide }}$ Divide | D ${ }_{\text {D }}$ | 5D | $\mathrm{R} 1, \mathrm{D} 2(\mathrm{X} 2, \mathrm{B2})$ or R1, D2 $2, \mathrm{B2} 2$ $\mathrm{R} 1, \mathrm{R} 2$ | $R 1, S 2(X 2) \text { or } R 1, S 2$ |  |
|  | Divide Decimal | DP | FD | R1, (L1, B1), D2 (L2, B2) | $\begin{aligned} & \mathrm{S1}(\mathrm{LI}), \mathrm{S} 2(\mathrm{~L} 2) \text { or } \mathrm{S} 1, \mathrm{~S} 2 \\ & \mathrm{R} 1, \mathrm{~S} 2(\mathrm{X} 2) \text { or } \mathrm{R} 1, \mathrm{~S} 2 \end{aligned}$ |  |
|  | Divide, Long | DD | 6 D | R1, D2(X2, B2), or R1, D2 (, B2) |  |  |
|  | Divide, Long | DDR | 2 D | $R 1, R 2$ | R1, S2(X2) or R1, S2 |  |
|  | Divide, Short Divide, Short | DE | $7 D$ 30 | $\mathrm{R} 1, \mathrm{D} 2(\mathrm{X} 2, \mathrm{~B} 2)$ or R1, $\mathrm{D} 2(, \mathrm{B2})$ $\mathrm{R1}, \mathrm{R} 2$ |  |  |
|  | Edit | ED | DE | D1(L, B1), D2(B2) | $S 1(L), S 2$ or $S 1, S 2$ <br> $S 1(L), S 2$ or $S 1, S 2$ <br> $R 1, S 2(X 2)$ or R1,S2 <br> $S 1(L), S 2$ or $S 1, S 2$ |  |
|  | Edit and Mark | EDMK | DF | D1 $(\mathrm{L}, \mathrm{B1} 1), \mathrm{D} 2(\mathrm{B2})$ |  |  |
|  | Exclusive Or | $\times$ | 57 | R1, $\mathrm{D} 2(\times 2, \mathrm{B2}$ ) or $\mathrm{R} 1, \mathrm{D} 2(, \mathrm{B2})$ |  |  |
|  | Exclusive Or | XC | D7 | D1(L, B1), D2 (B2) |  |  |
|  | Exclusive Or | XR | 17 | R1, R2 |  |  |
|  | Exclusive Or Immediate | X1 | 97 | D1(B1), 12 | \$1,12 |  |
|  | Execute | EX | 44 | $\mathrm{R} 1, \mathrm{D} 2(\times 2,82)$ or R1, $\mathrm{D} 2(, \mathrm{B2})$ | R1, S2(X2) R1, 52 |  |
|  | Halve, Long Halve, Short | $\begin{aligned} & \text { HDR } \\ & \text { HER } \end{aligned}$ | $\begin{aligned} & 24 \\ & 34 \end{aligned}$ | $\begin{aligned} & R 1, R 2 \\ & R 1, R 2 \end{aligned}$ |  |  |
|  | Halt Device | HDV | 9E011 | D1, ${ }^{11}$ | S1 | ${ }^{1}$ See Note 1 at end of this appendix |
|  | Halt 1/O | HIO | 9E00 | O1(BI) |  |  |
| \| DOS/VSE only | Insert Poge Bits. | IPB | B4 | R1, D2 (B2) | R1, S2 |  |
|  | Insert Character | IC | 43 | R1, D2 ( $\times 2, \mathrm{B2}$ ) or R1, D2 (, B2) | $\begin{aligned} & \mathrm{R} 1,52(\mathrm{X} 2) \\ & \mathrm{R} 1 . \mathrm{M} 3 . \mathrm{S} 2 \end{aligned} \text { or } \mathrm{R1}, \mathrm{~S} 2$ |  |
|  | Insert Characters under Mask Insert PSW Key | $\begin{aligned} & \text { ICM } \\ & \text { IPK } \end{aligned}$ | BF <br> B20B | R1, M3, D2(B2) | R1, M3, S2 |  |
|  | Insert Storage Key | ISK | 09 | R1,R2 |  |  |
|  | Load | L | 58 | $\mathrm{R1}, \mathrm{D} 2(\times 2, \mathrm{~B} 2) \text { or } \mathrm{R} 1, \mathrm{D} 2(, \mathrm{~B} 2)$ | R1, S2(X2) or R1, S2 |  |
|  | Load Load Address | LR | 18 | $\begin{aligned} & R 1, R 2 \\ & R 1, D 2(X 2, B 2) \text { or } \mathrm{R} 1, \mathrm{D} 2(, B 2) \end{aligned}$ | R1, S2(X2) or R1, S2 |  |
|  | Load and Test | LTR | 12 | R1,R2 |  |  |
|  | Load and Test, Long | LTDR | 22 | R1,R2 |  |  |
|  | Load and Test, Short | LTER | 32 | R1, R2 |  |  |
|  | Load Complement | LCR | 13 | R1,R2 |  |  |
|  | Load Complement, Long | LCDR | 23 | R1,R2 |  |  |
|  | Load Complement, Short | LCER | 33 | R1, R2 |  |  |
|  | Load Control | LCTL | B7 | R1, R3, D2(B2) | R1, R3, S2 |  |
| \|DOS/VSE only | Load Frame Index | LFI | B8. | R1, D2(B2) | R1, S2 . |  |
|  | Load Halfword | LH | 48 | R1, D2 ( $\times 2, \mathrm{B2}$ ) or R1, D2 (, B2) | $\left\lvert\, \begin{array}{ll} R 1, S 2(\times 2) & \text { or } \mathrm{R1}, \mathrm{S2} \\ \mathrm{R} 1, \mathrm{S2}(\times 2) & \text { or } \mathrm{R1}, \mathrm{~S} 2 \end{array}\right.$ |  |
|  | Load, Long | LD | 68 | R1, D2 ( $\times 2,82$ ) or R1, D2 $(, B 2)$ |  |  |
|  | Load, Long | LDR | 28 | R1,R2 |  |  |
|  | Load Multiple | LM | 98 | R1,R3, D2(B2) | R1, R3, S2 |  |
|  | Load Negative | LNR | 11 | R1,R2 |  |  |
|  | Load Negative, Long Load Negative, Short | LNDR LNER | 21 31 | $\begin{aligned} & \text { R1,R2 } \\ & \text { R1,R2 } \end{aligned}$ |  |  |
|  | Load Negative, Short | LNER | 31 | R1, R2 |  |  |
|  | Load Positive | LPR | 10 | R1,R2 |  |  |
|  | Load Positive, Long | LPDR | 20 | R1,R2 |  |  |
|  | Load Positive, Short | LPER | 30 | R1,R2 |  |  |
|  | Load PSW <br> Load Real Address | LPSW LRA | 82 ${ }^{\text {B1 }}$ | D1 (B1) <br> R1, D2 (X2, B2) or R1, D2 (, B2) |  |  |
| INat Dosivse | Load Rounded, Extended to Long | LRDR | 25 | $(\mathrm{R} 1, \mathrm{R} 2$ | $R 1, S 2(X 2) \text { or R1,S2 }$ |  |
|  | Load Rounded, Long to Short | LRER | 35 | R1, R2 |  |  |
|  | Load, Short <br> Load, Short | LEE | 78 38 | $\mathrm{R} 1, \mathrm{D} 2(\times 2, \mathrm{~B} 2)$ or R1, D2(, B2) $\mathrm{R1}$, 22 | R1, $52(\times 2)$ or R1, 52 |  |
| $\left\lvert\, \begin{aligned} & \text { DOS/VSE only } \\ & \text { DOS/VSE only } \end{aligned}\right.$ | Make Addressable .. | MAD | 3210 |  | $\begin{aligned} & \mathrm{s} 2 \\ & \mathrm{~s} 2 \end{aligned}$ |  |
|  | Make Unaddressable | MUN | B21E | D2(B2) . |  |  |
|  | Monitor Call | MC | AF | D1 $(\mathrm{B} 1), 12$ | S1,12 |  |
|  | Move Characters | MVC | D2 | D1(L, B1), D2(B2) | S1(L), S2 or S1, S2 |  |
|  | Move Immediate | MVI | 92 | D1(B1), 12 | S1,12 |  |



|  | Instruction | Mnemonic Operation Code | Machine Operation Code | Explicit Operand | Format Implicit |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | Move Long <br> Move Numerics <br> Move with Offset | MVCL MVN MVO | $\begin{aligned} & \text { OE } \\ & \mathrm{DI} \\ & \mathrm{FI} \end{aligned}$ | $\begin{aligned} & \mathrm{R} 1, \mathrm{R} 2 \\ & \mathrm{D} 1(\mathrm{~L}, \mathrm{B1}), \mathrm{D} 2(\mathrm{~B} 2) \\ & \mathrm{D} 1(\mathrm{LI}, \mathrm{BI}), \mathrm{D} 2(\mathrm{~L} 2, \mathrm{~B} 2) \end{aligned}$ | $\begin{aligned} & \mathrm{S1}(\mathrm{~L}), \mathrm{S} 2 \text { or } \mathrm{S} 1, \mathrm{~S} 2 \\ & \mathrm{~S} 1(\mathrm{LI}), \mathrm{S} 2(\mathrm{~L} 2) \text { or } \mathrm{S} 1, \mathrm{~S} 2 \end{aligned}$ |  |
|  | Move Zones | MVZ | D3 |  | $S 1(L), S 2 \text { or } S 1, S 2$ |  |
|  | Multiply Multiply | $M$ $M R$ | $\begin{aligned} & 5 \mathrm{C} \\ & 1 \mathrm{C} \end{aligned}$ | $\begin{aligned} & \mathrm{R} 1, \mathrm{D} 2(\mathrm{X} 2, \mathrm{~B} 2) \text { or } \mathrm{R} 1, \mathrm{D} 2(, \mathrm{~B} 2) \\ & \mathrm{R} 1, \mathrm{R} 2 \end{aligned}$ | $\mathrm{R1}, \mathrm{~S} 2(\mathrm{X} 2) \text { or } \mathrm{R1}, \mathrm{~S} 2$ |  |
|  | Multiply Decimal Multiply Extended | MP $M \times R$ | FC | D1 ( $\mathrm{L} 1, \mathrm{~B} 1), \mathrm{D} 2(\mathrm{L2}, \mathrm{~B} 2)$ $\mathrm{R1}$ ) | S1(LI), $52(\mathrm{~L} 2)$ or S1, S2 |  |
|  | Multiply Halfword | MH | 4 C | R1, D2(X2, B2) or R1, D2 (, B2) | R1, $52(\mathrm{X} 2)$ or $\mathrm{R1}, \mathrm{~S} 2$ |  |
|  | Multiply, Long Multiply, Long | MD MDR | ${ }^{6 C}$ | $\begin{aligned} & \mathrm{R} 1, \mathrm{D}_{2}(\mathrm{X} 2, \mathrm{~B} 2) \text { or } \mathrm{R} 1, \mathrm{D} 2(, \mathrm{~B} 2) \\ & \mathrm{R}, \mathrm{R} 2 \end{aligned}$ | $\mathrm{R1}, \mathrm{S2}(\mathrm{X} 2)$ or $\mathrm{R} 1, \mathrm{~S} 2$ |  |
|  | Multiply, Long to | MXD | 67 | R1, D2 (X2, B2) or R1, D2 (, B2) | R1,S2(X2) or R1, S2 |  |
|  | Multiply, Long to Extended | MXDR | 27 | R1, R2 |  |  |
|  | Multiply, Short | $\mathrm{ME}_{\mathrm{MER}}$ | $7 C$ $3 C$ |  | R1, S2(X2) or R1, 52 |  |
|  | No Operation | NOP NOPR | 47(BC <br> $07(B C)$ <br> 8 C | D2 $2(X 2, B 2)$ or D2(, B2) R2 | S2(X2) or 52 |  |
|  | Or Logical | $\bigcirc$ | 56 | R1, D2 ( $\times 2, \mathrm{B2}$ ) or R1, D2 (, B2) | $\mathrm{R} 1, \mathrm{~S} 2(\mathrm{X} 2) \text { or } \mathrm{R} 1, \mathrm{~S} 2$ |  |
|  | Or Logical | OC | D6 16 | D1(L, B1), D2(B2) ${ }_{\text {R1, } 21}$ | S1(L), $\mathrm{S2}$ or $\mathrm{S} 1, \mathrm{~S} 2$ |  |
|  | Or Logical Immediate | Ol | 96 | D1'(B1), 12 | S1,12 |  |
|  | Pack | PACK | F2 | D1(L1, B1), D2 (L2, B2) | S1'(L) , S2(L2) or S1,S2 |  |
|  | Purge Translation Lookaside Buffer | PtLb | B20D | - | - $\quad \cdots$ |  |
| \| Not Dos/Vse | Read Direct Reset Reference Bit | $\begin{aligned} & \text { RDD } \\ & \text { RRB } \end{aligned}$ | $\begin{aligned} & 85 \\ & 8213 \end{aligned}$ | $D 1 \text { (B1), } 12$ DI(BI) | $\begin{aligned} & \mathrm{s}, 12 \\ & \mathrm{~s} 1 \end{aligned}$ |  |
| DOS/VSE only | Refrieve Status and Page | RSP | D8 | $\begin{gathered} \mathrm{D1}(, \mathrm{B1}), \mathrm{D} 2(\mathrm{~B} 2) \\ \mathrm{or} \\ \mathrm{D1}(\mathrm{B1}), \mathrm{D} 2(\mathrm{~B} 2) \end{gathered}$ | S1, S2 |  |
| DOS/VSE only | Set Page Bits Set Clock | $\begin{array}{r} \text { SPB } \\ \text { SCK } \end{array}$ | $\begin{aligned} & B 5 \\ & B 204 \end{aligned}$ | R1, D2(B2) D1(B1) |  |  |
|  | Set Clock Comparator | SCKC | B206 | D1(B1) | Si |  |
|  | Set CPU Timer | SPT | B208 | DI(B1) |  |  |
|  | Set Prefix | SPX SPM | B210 04 | D2(B2) RI |  |  |
|  | Set PSW Key from Address | SPKA | B20A | D, (B, ) | SI |  |
|  | Set Storage Key | SSK | 08 | R1, R2 |  |  |
|  | Set System Mask Shift and Round Decimal | SSM SRP | 80 | $\begin{aligned} & \mathrm{D1}(\mathrm{BI}), \mathrm{Bi}, \mathrm{D} 2(\mathrm{B2}), \mathrm{I} 3 \\ & \mathrm{D} 1(\mathrm{LI}, \mathrm{l} \end{aligned}$ | $\begin{aligned} & S 1 \\ & S 1(L 1), S 2,13 \text { or } S 1, S 2,13 \end{aligned}$ |  |
|  | Shift Left Double Algebraic | SLDA | 8 F | R1, D2(B2) | R1, S2 : |  |
|  | Shift Left Double Logical | SLDL | 8D | R1, D2(B2) R1, 2 (B2) | R1, ${ }^{\text {R2 }}$ |  |
|  | Shift Left Single Logical | SLL | 88 | R1,D2(82) R1,D2(B2) | R1, R1, 22 |  |
|  | Shift Right Double Algebraic | SRDA SRDL | 8 E | R1, D2(82) | R1, ${ }^{\text {R2 }}$ |  |
|  | Shift Right Double Logical | SRDL | 8C | R1,D2(B2) R1,D2(B2) | R1, R1, 22 |  |
|  | Shift Right Single Logical | SRL | 88 | R1, ${ }^{\text {d2(B2) }}$ | R1, ${ }^{\mathrm{R} 1} \mathrm{~S}^{\text {R }}$ |  |
|  | Signal Processor | SIGP | AE | R1, R3, D2(B2) | R1, R3, S2 |  |
|  | ```Start 1/O Start I/O Fast Release``` | $\begin{aligned} & \text { SIO } \\ & \text { SIOF } \end{aligned}$ | $\begin{aligned} & 9 \mathrm{COO} \\ & 9 \mathrm{CO1} \\ & \hline 1 \end{aligned}$ | DI(BI) DI(BI) | S1 |  |
| I dosivse only | Store <br> Store Capacity Counts | $\begin{aligned} & \text { ST } \\ & \text { STCAP } \end{aligned}$ | ${ }^{50}$ | R1, $22(\mathrm{X2}, \mathrm{B2)} \mathrm{or} \mathrm{R1} ,\mathrm{D} 2(, \mathrm{B2})$ $\mathrm{D2} 2$ (B2) | R1,S2(X2) or R1, S2 |  |
|  | Store Channel ID Store Character | STIDC | B203 42 | D1(B1) | $\begin{aligned} & S 2 \\ & S 1 \end{aligned}$ |  |
|  |  | STC | 42 | R1,D2(X2, B2) or R1, D2 (, B2) | $\mathrm{R1}, \mathrm{D} 2(\mathrm{X} 2) \text { or } \mathrm{R} 1, \mathrm{~S} 2$ |  |
|  | Store Characters under Mask | STCM | BE | R1, M3, D2(B2) | R1, M3, S2 |  |
|  | Store Clock | STCK | B205 | DI(B1) | S1 |  |
|  | Store Clock Comparator | STCKC | B207 | DI(BI) |  |  |
|  | Store Control <br> Store CPU address | STCTL | ${ }_{8212}{ }^{\text {B6 }}$ | R1,R3, D2(B2) | R1, R3, S2 |  |
|  | Store CPU ID | STIDP | ${ }^{\text {B21 }}$ | D1(B1) |  |  |
|  | Store CPU Timer | STPT | 8209 | D1( $\mathrm{BI}^{\text {) }}$ | 51 |  |
|  | Store Halfword Store Long | STH | 40 |  | R1, $\mathrm{S2}(\times 2)$ or R1, S2 |  |
|  | Store Multiple | STM | 60 90 | R1,D2( ${ }^{\text {R2, }}$ (12) ${ }^{\text {R2, }}$ (B2) | R1, S2(X2) R1, 2 , 20 or R1, S2 |  |
|  | Store Prefix | STPX | B211 | D2 (B2) |  |  |
|  | Store Short | STE | 70 | R1, D2 (X2, B2) or R1, D2 (, B2) | R1, $22(\mathrm{X} 2)$ or R1, $\mathrm{S}^{2}$ |  |
|  | Store Then AND System Mask | STNSM | $A C$ | D1(B1), 12 |  | ${ }^{1}$ See Note 2 at end of this appendix |
|  | Store Then OR System Mask Subtract | ${ }_{\mathrm{S}}^{\text {STOSM }}$ | AD | D1(B1), 12 <br> R1, D2 ( ${ }^{2}$ ) | $51,12$ |  |
|  |  |  | 5B | R1, D2 ( ${ }^{\text {2 }}$ ) | $\mathrm{R1}, \mathrm{~S} 2(\mathrm{X} 2) \text { or } \mathrm{R} 1, \mathrm{~S} 2$ |  |
|  | Subtract | SR | 1B | R1, R2 |  |  |
|  | Subtract Decimal | SP | FB | D1(L1, B1), D2(L2, B2) | S1(LI) S2(L2) or S1, $\mathrm{S}^{\text {2 }}$ |  |
|  | Subtract Halfword | SH | ${ }^{48}$ | R1, D2 ( $\times 2, \mathrm{~B} 2$ ) or $\mathrm{R1} 1, \mathrm{D} 2(, \mathrm{B2} 2)$ | R1, S2 (X2) or R1, ${ }^{\text {S2 }}$ |  |
|  | Subtract Logical | SLR | 5F | $\mathrm{R} 1, \mathrm{D} 2(\mathrm{X} 2, \mathrm{~B} 2)$ or R1, $\mathrm{D} 2(, \mathrm{~B} 2)$ $\mathrm{R} 1, \mathrm{R} 2$ | R1, S2(X2) or R1, 52 |  |



| Instruction | Mnemonic Operation Code | Machine Operation Code | Operand Format |  |
| :---: | :---: | :---: | :---: | :---: |
|  |  |  | Explicit | Implicit |
| Subtract Normalized, Extended | SXR | 37 | R1, R2 |  |
| Subtract Normalized, Long | SD | 6 B | $\mathrm{R1} 1, \mathrm{D} 2(\mathrm{X} 2, \mathrm{B2})$ or R1, $\mathrm{D} 2(, \mathrm{B2})$ | R1, $\mathrm{S} 2(\mathrm{X} 2)$ or R1, S 2 |
| Subtract Normalized, Long | SDR | 2 B | R1,R2 |  |
| Subtract Normalized, Short Subtract Normalized, | SE | $7 B$ 38 |  | R1, $\mathrm{S2}$ (X2) or R1, S2 |
| Subtract Unnormalized, Long | SW | 6F | R1, D2 ( $\mathrm{X} 2, \mathrm{~B} 2$ ) or R1, D2 (, B2) | R1,S2(X2) or R1, S2 |
| Subtract Unnormalized, Long | SWR | 2 F | R1, R2 |  |
| Subtract Unnormalized, Short | SU | 7 F | R1, $22(\mathrm{X} 2, \mathrm{~B} 2)$ or R1, D2 (, B2) | $\mathrm{R1}, \mathrm{~S} 2(\mathrm{X} 2)$ or $\mathrm{R} 1, \mathrm{~S} 2$ |
| Subtract Unnormalized, Short Supervisor Call | SUR SVC | 3F | R1, R2 |  |
| Test and Set | is | 93 | D1(B1) | 51 |
| Test Channel | TCH | 9 F | D1(B1) | 51 |
| Test 1/O | TIO | 9 D | D1(BI) | S1 |
| Test Under Mask | TM | 91 | DI(BI), 12 | S1,12 |
| Translate | TR | DC | D1(L, B1), D2(B2) |  |
| Translate and Test | TRT | DD | D1(L, B1), D2 (B2) | $\mathrm{S} 1(\mathrm{~L}), \mathrm{S} 2 \text { or } \mathrm{S} 1 ; \mathrm{S} 2$ |
| Unpack <br> Write Direct | UNPK WRD | F3 | $\begin{aligned} & D 1(L 1, B 1), D 2(L 2, B 2) \\ & D 1(B 1), 12 \end{aligned}$ | $\begin{aligned} & \mathrm{S} 1(\mathrm{~L} 1), \mathrm{S} 2(\mathrm{~L} 2) \text { or } \mathrm{S} 1, \mathrm{~S} 2 \\ & \mathrm{~S} 1,12 \end{aligned}$ |
| Zero and Add Decimal | ZAP | F8 | D1(L1, B1), D2 (L2, B2) | S1(L1), S2(L2) or S1,S2 |

\begin{tabular}{|c|c|c|c|c|c|c|c|c|c|c|c|c|c|c|}
\hline \multicolumn{2}{|r|}{\multirow[t]{2}{*}{Instruction}} \& \multirow[t]{2}{*}{Type of Instruction} \& \& \multicolumn{6}{|r|}{Program Interruption Possible} \& \multicolumn{5}{|c|}{Condition Code Set} \\
\hline \& \& \& \& A \& 5 \& Ov \& P \& Op \& Other \& 00 \& 01 \& \& 10 \& 11 \\
\hline \begin{tabular}{l}
Subtract Normalized, Extended \\
Subtract Normalized, Long \\
Subtract Normalized, Long \\
Subrract Normalized, Short \\
Subtract Normalized, Short \\
Subtract Unnormaliezd, Long \\
Subtract Unnormailezd, Long \\
Subtract Unnormalized, Short \\
Subtract Unnormalized, Short \\
Supervisor Call
\end{tabular} \& \begin{tabular}{l}
tract Normalized, Extended tract Normalized, Long tract Normalized, Long tract Normalized, Short tract Normalized, Short tract Unnormaliezd, Long tract Unnormailezd, Long tract Unnormalized, Short tract Unnormalized, Short ervisor Call \\
and Set \\
Channel \\
1/O \\
under Mask \\
nslate \\
slate and Test \\
pack \\
ite Direct \\
o and Add Decimal
\end{tabular} \& RR, Floating Pt. RX, Floating Pt. RR, Floating Pt. RX, Floating Pr. RR, Floating Pt. RX, Floating Pt. RR, Floating Pt. RX, Floating Pr. RR, Floating Pt. RR \& \& \[
\begin{aligned}
\& x \\
\& x \\
\& x \\
\& x \\
\& x \\
\& x \\
\& x \\
\& x \\
\& x \\
\& x \\
\& x \\
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\& x
\end{aligned}
\] \&  \&  \&  \& \[
\begin{array}{|l}
x \\
x \\
x \\
x \\
x \\
x \\
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\\
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x \\
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x
\end{array}
\] \& B, C
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\(B, C\)
\(B, C\)
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\(C\)

A

A \& | $R$ |
| :--- |
| $R$ |
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| $N$ |
| SS |
| $J J$ |
| $L L$ |
| $U U$ |
| $N$ |
| $P P$ |
|  | \& $L$

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$M$
$M$

$N$ \& | $Q$ $Q$ $Q$ $Q$ $Q$ $Q$ $Q$ $Q$ $N$ |
| :--- |
| HH |
| KK |
| ww |
| N |
| $N$ |
| N O | <br>

\hline \multicolumn{3}{|c|}{\[
Under Ov: $$
\begin{aligned}
D & =\text { Decimal } \\
E & =\text { Exponent } \\
F & =\text { Fixed Point }
\end{aligned}
$$

\]} \& \& \multicolumn{2}{|l|}{der Other:} \& r: \& \& \multicolumn{3}{|l|}{| Privileged Operation |
| :--- |
| Exponent Underflow |
| Significance |
| Decimal Divide |} \& \multicolumn{4}{|c|}{| E Floating Point Divide |
| :--- |
| F Fixed Point Divide |
| G Execute |
| GA Monitoring |} <br>

\hline \multicolumn{15}{|l|}{Condition Code Set} <br>

\hline \[
$$
\begin{aligned}
& \mathrm{H} \\
& \mathrm{I} \\
& \mathrm{~J} \\
& \mathrm{~K} \\
& \mathrm{~L} \\
& \mathrm{M} \\
& \mathrm{~N} \\
& \mathrm{O} \\
& \mathrm{P} \\
& \mathrm{Q} \\
& \mathrm{R} \\
& \mathrm{~S} \\
& \mathrm{~T} \\
& \mathrm{U} \\
& \mathrm{~V} \\
& \mathrm{~W} \\
& \mathrm{X} \\
& \mathrm{Y} \\
& \mathrm{Z} \\
& \mathrm{AA} \\
& B B \\
& \mathrm{CC} \\
& \mathrm{DD} \\
& \mathrm{EE} \\
& \mathrm{FF} \\
& \mathrm{~GB} \\
& \mathrm{GC} \\
& \mathrm{GG} \\
& \mathrm{HH} \\
& \mathrm{II} \\
& \mathrm{JJ} \\
& \mathrm{KK} \\
& \mathrm{LL} \\
& M M
\end{aligned}
$$

\] \& | No carry |
| :--- |
| Carry |
| Result $=0$ |
| Result is not equal to zero |
| Result is less than zero |
| Result is greater than zero |
| Not changed |
| Overflow |
| Result exponent underflows |
| Result exponent overflows |
| Result fraction $=0$ |
| Result field equals to zero |
| Result field is less than zero |
| Resulf field is greater than $z$ |
| Difference $=0$ |
| Difference is not euqal to zero |
| Difference is less than zero |
| Difference is greater than ze |
| First operand equals second |
| First operand is less than sec |
| First operand is greater than |
| CSW stored |
| Channel and subchannel not |
| Channel or subchannel busy |
| Channel operating in burst $m$ |
| Page state |
| Page transision |
| Burst operation terminated Channel not operational Interruption pending on chan Channel available |
| Not operational |
| Available |
| 1/O operation initiated and proceeding with its executio | \& | rand d operand cond operand rking |
| :--- |
| e |
| annel | \& | NN |
| :--- |
| 00 PP QQ |
| RR |
| SS Tr |
| UU |
| vV |
| WW |
| XX |
| YY |
| ZZ |
| AAA |
| AAB |
| AAC |
| AAD |
| AAE |
| AAF |
| AAG |
| AAH |
| AAI |
| AAJ |
| AAK |
| AAL |
| AAM |
| AAN |
| AAO |
| AAP |
| AAQ |
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| ock |
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| fere | \&  \& | uncti on by by ing to oade ng to cifi of of ts ar ts ar its ar tes ld d and d an |
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| d d nt $b$ set sec opera corr tivit is not orkin busy ror or pa entry bit $z$ bit | \& | ion by nd fie |
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| ytes tes ar obits d |
| obits ed by byte byte e all mix all are of firs firs nd se |
| ount ount becau |
| cure ation rectly ty pro valid nec ng wi $y$ or state age-t ry inv ero, zero, | \& | yte foun ild is ex are zero re zero 34 and |
| :--- |
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| is lower is highe se of de |
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| $y$ stored hibited |
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| able len alid (Ichange change | \& | before hausted |
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| $d=0$ |
| $d=1$ |
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| mask in |
| dis low |
| d is high |
| rand cou |
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| valid |
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| bit zero |
| bit one | \& | zeros |
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| Sav | \& | nce bit one, change bit zero ce bit one, change bit one t table entry invalid (I-bit one) ation available |
| :--- |
| and second hand operands equal, operand replaced by |
| ird operand |
| eration is in progress for |
| dressed device |
| code accepted |
| stored |
| sful, block was disconnected, returned |
| was already disconnected, returned |
| ccessful, index returned |
| was connected |
| was already disconnected |
| returned, page is addressable |
| returned, page is connected |
| not returned, page is |
| ected |
| returned, address is invalid |
| already addressable |
| addressable |
| already connected |
| valid |
| invalid | <br>

\hline
\end{tabular}

| RR Format |  |  |  |
| :---: | :---: | :---: | :---: |
| Operation <br> Code | Name | Mnemonic | Remarks |
| 00 |  |  |  |
| 01 |  |  |  |
| 02 |  |  |  |
| 03 |  |  |  |
| 04 Set Program Mask SPM |  |  |  |
| 05 Branch and Link BALR |  |  |  |
| 06 Branch on Count BCTR |  |  |  |
| 07 Branch on Condition BCR |  |  |  |
| 08 Set Storage Key SSK |  |  |  |
| 09 Insert Storage Key ISK |  |  |  |
| 0A | Supervisor Call | SVC |  |
| OB |  |  |  |
| OC |  |  |  |
| OE | Move Long | MVCL |  |
| OF | Compare Logical Long | CLCL |  |
| 10 | Load Positive | LPR |  |
| 11 | Load Negative | LNR |  |
| 12 | Load and Test | LTR |  |
| 13 | Load Complement | LCR |  |
| 14 | AND | NR |  |
| 15 | Compare Logical | CLR |  |
| 16 | OR | OR |  |
| 17 | Exclusive OR | XR |  |
| 18 | Load | LR |  |
| 19 | Compare | CR |  |
| 1 A | Add | AR |  |
| 1 B | Subtract | SR |  |
| 1 C | Multiply | MR |  |
| 1 D | Divide | DR |  |
| 1 E | Add Logical | ALR |  |
| 1 F | Subtract Logical | SLR |  |
| 20 |  | LPDR |  |
| 21 | Load Negative (Long) | LNDR |  |
| 22 | Load and Test (Long) | LTDR |  |
| 23 | Load Complement (Long) | LCDR |  |
| 24 | Halve (Long) | HDR |  |
| 25 | Load Rounded (Extended to Long) | LRDR |  |
| 26 | Multiply (Extended) | MXR |  |
| 27 | Multiply (Long to Extended) | MXDR |  |
| 28 | Load (Long) , | LDR | - |
| 29 | Compare (Long) | CDR |  |
| 2A | Add Normalized (Long) | ADR |  |
| 2 B | Subtract Normalized | SDR |  |
| 2 C | Multiply (Long) | MDR |  |
| 2D | Divide (Long) | DDR |  |
| 2 E | Add Unnormalized (Long) | AWR |  |
| 2 F | Subtract Unnormalized (Long) | SWR |  |
| 30 | Load Positive (Short) | LPER |  |
| 31 | Load Negative (Short) | LNER |  |
| 32 | Load and Test (Short) | LTER |  |
| 33 | Load Complement (Short) | LCER |  |
| 34 | Halve (Short) | HER |  |
| 35 | Load Rounded (Long or Short) | LRER |  |
| 36 | Add Normalized (Extended) | AXR |  |
| 37 | Subtract Normalized (Extended) | SXR |  |
| 38 | Load (Short) | LER |  |


| RR Format |  |  |  |
| :---: | :---: | :---: | :---: |
| Operation code | Name | Mnemonic | Remarks |
| $\begin{aligned} & 39 \\ & 3 A \\ & 3 B \\ & 3 C \\ & 3 D \\ & 3 E \\ & 3 F \end{aligned}$ | Compare (Short) <br> Add Normalized (Short) <br> Subtract Normalized (Short) <br> Multiply (Short) <br> Divide (Short) <br> Add Unnormalized (Short) <br> Subtract Unnormalized (Short) | CER <br> AER <br> SER <br> MER <br> DER <br> AUR <br> SUR |  |
| RX Format |  |  |  |
| 40 <br> 42 <br> 43 <br> 44 <br> 45 <br> 46 <br> 47 <br> 48 <br> 49 <br> 4A <br> 4B <br> 4C <br> 4 E <br> 4F <br> 50 <br> 51 <br> 52 <br> 53 <br> 54 <br> 55 <br> 56 <br> 57 <br> 58 <br> 59 <br> 5A <br> 5B <br> 5C <br> 5D <br> 5 E 5 F <br> 60 <br> 61 <br> 62 <br> 63 <br> 64 <br> 65 <br> 66 <br> 67 <br> 68 <br> 69 <br> 6A <br> 6B <br> 6C <br> 6D <br> 6 F | Store Halfword <br> Load Address <br> Store Character <br> Insert Character <br> Execute <br> Branch and Link <br> Branch on Count <br> Branch on Condition <br> Load Halfword <br> Compare Halfword <br> Add Halfword <br> Subtract Halfword <br> Multiply Halfword <br> Convert to Decimal <br> Convert to Binary <br> Store <br> AND <br> Compare Logical <br> OR <br> Exclusive OR <br> Load <br> Compare <br> Add <br> Subtract <br> Multiply <br> Divide <br> Add Logical <br> Subtract Logical <br> Store (Long) <br> Multiply (Long to Extended) <br> Load (Long) <br> Compare (Long) <br> Add Normalized (Long) <br> Subtract Normalized (Long) <br> Multiply (Long) <br> Divide (Long) <br> Add Unnormalized (Long) <br> Subtract Unnormalized (Long) | STH <br> LA <br> STC <br> IC <br> EX <br> BAL <br> BCT <br> BC <br> LH <br> CH <br> AH <br> SH <br> MH <br> CVD <br> CVB <br> ST <br> N <br> CL <br> 0 <br> X <br> L <br> C <br> A <br> S <br> M <br> D <br> AL <br> SL <br> STD <br> MXD <br> LD <br> CD <br> AD <br> SD <br> MD <br> DD <br> AW <br> SW |  |


| RX Format |  |  |  |
| :---: | :---: | :---: | :---: |
| Operation Code | Name | Mnemonic | Remarks |
| $\begin{aligned} & 70 \\ & 71 \\ & 72 \\ & 73 \\ & 74 \\ & 75 \\ & 76 \\ & 77 \\ & 78 \\ & 79 \\ & 7 \mathrm{~A} \\ & 7 \mathrm{~B} \\ & 7 \mathrm{C} \\ & 7 \mathrm{D} \\ & 7 \mathrm{E} \\ & 7 \mathrm{~F} \end{aligned}$ | Store (Short) <br> Load (Short) <br> Compare (Short) <br> Add Normalized (Short) <br> Subtract Normalized (Short) <br> Multiply (Short) <br> Divide (Short) <br> Add Unnormalized (Short) <br> Subtract Unnormalized (Short) | STE <br> LE <br> CE <br> AE <br> SE <br> ME <br> DE <br> AU <br> SU |  |
| RS,SI, and S Format |  |  |  |
|  | Set System Mask <br> Load PSW <br> Diagnose <br> Write Direct <br> Read Direct <br> Branch on Index High <br> Branch on Index Low or Equal <br> Shift Right Single Logical <br> Shift Left Single Logical <br> Shift Right Single <br> Shift Left Single <br> Shift Right Double Logical <br> Shift Left Double Logical <br> Shift Right Double <br> Shift Left Double <br> Store Multiple <br> Test under Mask <br> Move (Immediate) <br> Test and Set <br> AND (Immediate) <br> Compare Logical (Immediate) <br> OR (Immediate) <br> Exclusive OR (Immediate) <br> Load Multiple <br> Start I/O, Start I/O Fast Release <br> Test I/O <br> Halt I/O, Halt Device <br> Test Channel | SSM <br> LPSW <br> WRD <br> RDD <br> BXH <br> BXLE <br> SRL <br> SLL <br> SRA <br> SLA <br> SRDI <br> SLDL <br> SRDA <br> SLDA <br> STM <br> TM <br> MVI <br> TS <br> NI <br> CLI <br> OI <br> XI <br> LM <br> SIO,SIOF <br> TIO <br> HIO, HDV <br> TCH | See Note 2 <br> See Note 1 |


| Operation <br> Code | Name | Mnemonic | Remarks |
| :---: | :---: | :---: | :---: |
| A7 |  |  |  |
| A8 |  |  |  |
| A9 |  |  |  |
| AA |  |  |  |
| AB |  |  |  |
| AC |  |  |  |
| AD | Store Then OR System Mask | STOSM |  |
| AE |  |  |  |
| AF |  |  |  |
| B0 | Connect Page | CTP |  |
| B1 | Load Real Address | LRA |  |
| B2 | (First byte of two-byte operation cod |  |  |
| B3 |  |  |  |
| B4 | Insert Page Bits | IPB |  |
| B5 | Set Page Bits | SPB |  |
| B6 | Store Control | STCTL |  |
| B7 | Load Control | LCTL |  |
| B8 | Load Frame Index | LFI |  |
| B9 |  |  |  |
| BA |  |  |  |
| BB |  |  |  |
| BD | Compare Logical Characters under Mask | CLM |  |
| BE | Store Characters under Mask | STCM |  |
| BF | Insert Characters under Mask | ICM |  |
| SS Format |  |  |  |
| C0 |  |  |  |
| Cl |  |  |  |
| C2 |  |  |  |
| C3 |  |  |  |
| C4 |  |  |  |
| C5 |  |  |  |
| C6 |  |  |  |
| C7 |  |  |  |
| C8 |  |  |  |
| C9 |  |  |  |
| CA |  |  |  |
| CB |  |  |  |
| CC |  |  |  |
| CD |  |  |  |
| CE |  |  |  |
| CF |  |  |  |
| D0 |  |  |  |
|  |  |  |  |
| D2 $\mathrm{l}^{\text {a }}$ Move (Characters) ${ }^{\text {a }}$ ( MVC |  |  |  |
| D3 Move Zones MVZ |  |  |  |
| D4 AND (Characters) ${ }^{\text {D }}$ ( ${ }^{\text {a }}$ |  |  |  |
|  |  |  |  |
|  |  |  |  |
| D7 Exclusive OR (Characters) XC |  |  |  |
| D8 Retrieve Status and Page RSP |  |  |  |
|  |  |  |  |
| DA |  |  |  |
| DB DC |  |  |  |
| DC | Translate | TR |  |



## NOTES

1. Under the System/370 architecture, the machine operations for Halt Device and Halt I/O are as follows:

( X denotes an ignored bit position)
2. Under the System/370 architecture, the machine operations for Start I/O and Start I/O Fast Release are as follows:
10011100 XxXX XXX0 Start I/O SIO

10011100 XXXX XXX1 Start I/O Fast Release SIOF
( X denotes an ignored bit position)

| Operation <br> Code | Name | Mnemonic |
| :--- | :--- | :--- |
| AE | Signal Processor |  |
| BA | Compare and Swap | SIGP |
| BB | Compare Double and Swap | CS |
| 9D01 | Clear I/O ID | CDS |
| B202 | Store CPU ID | CLRIO |
| B203 | Store Channel ID | STIDP |
| B204 | Set Clock | STIDC |
| B205 | Store Clock | SCK |
| B206 | Set Clock Comparator | STCK |
| B207 | Store Clock Comparator | SCKC |
| B208 | Set CPUTimer | STCKC |
| B209 | Store CPUTimer | SPT |
| B20A | Set PSW Key from Address | STPT |
| B20B | Insert PSW Key | SPKA |
| B20D | Purge Translation | IPK |
|  |  |  |
| B210 | Set Prefix Lookaside Buffer | PTLB |
| B211 | Store Prefix | SPX |
| B212 | Store CPUAddress | STPX |
| B213 | Reset Reference Bit | STAP |
| B215 | Clear Page | RRB |
| B21B | Deconfigure Page | CLRP |
| B21C | Disconnect Page | DEP |
| B21D | Make Addressable | DCTP |
| B21E | Make Unaddressable | MAD |
| B21F | Store Capacity Counts | MUN |

This Newsletter No. GN20-9372
Date 28 December 1981

Base Publication No. GC33-4010-5
File No. S370-21

Prerequisite Newsletters None

## OS/VS-DOS/VSE-VM/370 Assembler Language

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This technical newsletter, a part of Release 4 of OS/VS1, Release 3 of OS/VS2, Release 2 of VM/370, and DOS/VSE, provides replacement pages for the subject publication. These replacement pages remain in effect for any subsequent releases unless specifically altered. Pages to be inserted and/or removed are:
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A change to the text or to an illustration is indicated by a vertical line to the left of the change.

## Summary of Amendments

This technical newsletter contains maintenance changes.

Note: Please file this cover letter at the back of the manual to provide a record of changes.

