QAADAD

AN ASSEMBLY SYSTEM FOR THE SOLID STATE II

BY D. E. KNUTH AND W. C. LYNCH

QADAAD ASSEMBLY PROGRAM

Brief Description

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"Quick-And-Dirty Assembler And Documentor"

QADAAD is an assembly program for a Univac Solid State II (80 column) configuration with at least 5000 words of drum memory. It assembles from tape to tape at a speed of about 600 lines per minute. QADAAD provides facilities for hand minimization of latency, for relocatable output, and a special optional feature which draws flow charts of the program as an extra pass.

Card Format

cols 1-5	A field:	location
col 6	AR field:	relocation code for A
cols 7-9	AH field:	latency code for A
cols 10-12	OP field:	symbolic operation code
col 13	IR field:	index register or sign
cols 14-18	M field:	M-address
col 19	MR field:	relocation code for M
cols 20-22	MH field:	latency code for M
cols 23-27	C field:	C-address
col 28	CR field:	relocation code for C
cols 29 - 31	CH field:	latency code for C
cols 32-80	remarks.	

Symbolic Fields

- The A, M, and C fields contain information of the following types: (here ∧ indicates a blank column)
- Regional addresses, Xnnnn where X is any non-blank character and n represents a digit, e.g. A0000. A0003 is three locations greater than A0000. A9999 is one location less than A0000.
- 2. Absolute addresses, ∧NNNN where N represents either a digit or one of the undigits A,B,C,F,G,H. This address is simply reproduced in

the output. Examples: 4078, B2AB, H66H.

- 3. Local addresses. In A these are of the form nAAAA, and in M or C they are of the form nFAAA or nBAAA (referring to the next or previous A of nAAAA respectively).
- 4. "Self": *AAAA refers to A address. If this address occurs in A or on a control operator card it refers to the previous A address.
- 5. Pair address, any address of the form &WWWW or -WWWW where W's are arbitrary, indicates two adjacent addresses as used in "c + 1 conditions." The - is the first of the pair, the & is the second or overflow address.
- 6. Blank address. In A field, indicates the address corresponding to a blank M or C address on the preceding instruction. In M if the address is ignored by the Op (e.g. ATL or HLT) this is equivalent to *. In C if the address is ignored (e.g. CLA) the C address is set equal to M. Otherwise the next A field must be blank and it refers to this next A address.
- 7. Symbolic address. This is any address other than the above, except the leftmost character must be nonblank and non-numeric. Several symbolic addresses are pre-defined: RA is 000A, RL is 000B, RX is 000C; RB4, RB5,..., RB9 are respectively B2AB, B3AB, B5AB, B6AB, B7AB, B8AB. If multiple assembly is used these predefined symbols disappear, however.

<u>H Fields</u>

The AH, MH, and CH fields are relevant only if the corresponding address is undefined, i.e. if it has not appeared earlier. When the address is undefined, the H field is used to control storage allocation as follows:

- 1. MAA or DAA Choose the best available place on the drum.
- 2. HAA Choose the best available place on the high-speed bands.
- 3. CAA Choose a place in the BOO1-B999 core storage area.
- 4. nnn Assign on this level on the drum.
- 5. Ann Assign on this level on the high-speed bands.
- 6. +nn Assign nn higher than the normal rules would say.

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R Fields

The AR, MR, and CR fields are merely transferred to the output area and are used to prepare FORTRAN subprograms. The FORTRAN assignment ignores the 4-bit and then has the following code:

	0 absolute	5 external reference					
	1 unique storage	6 in AR for ALF table entry					
۸	or 2 program storage	7 cross reference					
	3 common storage	8 special program storage					

IR Field

The IR field is normally used to specify index register modification. If blank there is no indexing, otherwise the digits 1-9 are used to indicate modification by RB1-RB9. (In the control operators CON NUM ZON and ALF which provide constants the IR field is used to denote the sign instead, and then the sign is 0 if IR is blank, the sign is 1-9 if the IR field is 1-9.)

A convenient literal constant feature is included, so that if the IR field contains the character # the M and C fields are filled with a ten-digit constant which is positive. This constant is assembled separately, then the M address of the instruction refers to this constant, and C address is treated as if it were blank.

Example: ADD#01234 56789

is equivalent to ADD CONST and the line

CONST CON 012340 567890

appearing later. (Notice that the MR and CR fields of the constant are set to 0; the AR field is set to the MR field of the instruction.) The MH and CH fields of the instruction retain their normal significance.

Warning

Do not attempt to assembly anything into location 0000; this location is sacred to QADAAD. As a matter of fact the loading routine supplied on the output tape goes into band 0000 so it is wise not to assemble into any of the locations 0000-0399.

OP Field

The operation code field is filled with either a three-letter mnemonic op-code as used in the S4 assembly system, or it is one of the control operators listed below. Operators which require a 4-bit in the sign digit are not provided.

ARA. This is used for comments only.

CPY. The card-to-tape pass uses this. M and C are absolute addresses which refer to line numbers on the previous assembly listing; lines M through C are copied from tape in place of this card.

FIN. This is the last card of a program, signalling the end for the card-to-tape pass.

END. This is the second-last card of a program, signalling the end of the program. It can also be used to separate programs in a multi-program assembly. The M address must be defined, and the loading routine will be set to halt and transfer to this address.

BLR and BLA. Reserve or unreserve drum locations starting at M and ending at C. M and C must be defined addresses on the drum. If the CH field is nonblank, take n to be CH mod 100, indicating that we reserve or unreserve every n-th word. If the CH field is blank it means take n to be 1. If A is nonblank an EQU operation is also performed (See EQU). If C is blank it is taken equal to M.

COR. Reserve core locations in the BOO1-B999 area. M must de defined. M locations in the core are reserved. If A is non-blank an EQU to the first address of this block is also performed (see EQU). For example, ROO01 COR 0030 reserves thirty core locations as the locations ROO01 through ROO30.

EQU. A must be an undefined regional, forward local, or symbolic address. $(\underline{Not} \ a \ pair \ address,)$ It is then defined to equal the equivalent of M which must be <u>defined</u>.

HHH. All succeeding blank H-fields are overridden by the contents of the MH field on this card. For example, HHH AAAAAAC is used to start assignment of undefined address in core. HHH AAAAAAH is used to assign in high-speed storage.

CON, NUM, ZON, ALF. These provide four types of constants. The A field is treated exactly like the A field of an instruction. The M and C fields specify a 10-digit constant, except with ALF only M is used. CON means treat the contents of M and C as digits or the undigits A,B,C,F,G,H. NUM and ZON mean to treat M and C as an alphabetic constant and take the numeric or zone word of this constant (expressed in machine code). ALF takes the M field and produces a constant zzzznnnn where zzzz is the zone part and nnnnn is the numeric part. On all four constant operators the IR field is used to specify the sign.

PAT. Print availability table. This causes 50 lines to be printed. Each line has five entries:

Level L L Level L+100 Level L+150

The availability word for a level contains 40 bits in machine code, corresponding to the status of the word in each band with the following band assignments:

00 02	18
20 22	38
40 42	58
60 62	78

As an example, at the beginning of assembly PAT would produce

BGGGG99999	GGGGG99999	000000000	GGGGG99999	GGGGG99999
GGGGG99999	GGGGG99999	0000010000	GGGGG99999	GGGGG99999

etc.

(locations 0001-4999 are available).

TYP, ON , OFF. Three control operators are provided in case a person wants to use a single assembly deck to produce several almost identical versions, without

going through the card-to-tape pass each time. TYP causes the computer to stop at assembly time and then RL should be set to 000000nnnn. The operators OFF and ON have a defined M address and are effective only¹ if the equivalent of M is nnnn matching a previous TYP. They temporarily shut assembly off or on. The FORTRAM system program is coded for five types:

> 8001:SS80, I 8002:SS80, II 9000:SS90, card 9001:SS90, tape I 9002:SS90, II

Note: The line B8AH NEW100001 00000 causes some assembly into the B00A-B99F area of core storage. Other uses of the control op NEW will not be listed here.

Operating Instructions

Loading	the	assemb	ler:		Loading	the	object	program:
rC	G2	0400	000A		rC	G2	0500 000	AC
rA	F6	B000	B000	1	rA	F6	7800 780	01

(The assembler is presently set up for an 8800 word drum.)

Mount QADAAD tape on unit 4, scratch tapes on units 3 and 5. If there is tape input, put it on unit 2. (The card-to-tape pass copies from cards and/or unit 2 to unit 3. If this phase is omitted, mount a good input tape on unit 3 and forget unit 2.) If flowcharting is desired, you will need scratch tapes on units 6 and 7 also.

As soon as the assembler has loaded itself, it stops. Depress RUN if the card-to-tape pass is to be used else depress M and RUN to skip this pass. When the second pass has loaded, it stops, and the operator should depress RUN. The computer should next stop with HLT. B9AH, but several other stops may occur.

Pass 1 (card-to-tape) stops:

0001 Comparison failure on card reader. The bad cards have been diverted to stacker 2. Reload them, depress RUN.

0002 Card reader off normal or empty. Fix, depress RUN.

- 0005 Error in writing tape. This is really unfortunate. If RUN is depressed the condition is ignored.
- 0006 Error in reading tape. If RUN is depressed reading is attempted again, in the opposite direction. Gain should be varied by the operator.
- Pass 2 (assembly) stops:
- 1111 Line numbers on input tape not sequential, indicates a bad situation on the input tape (unit 3)
- 2222 Tape error. Depressing RUN will cause the last tape instruction (read or write) to be tried again, with no backspacing provided. The last tape instruction appears in RA, and it is usually a G2 0300. In this case, restart by typing G2 0305 000A into rC, depress general clear and RUN.
- 3333 Printer off normal. Fix, depress RUN.
- 4444 Tape 3 not yet ready. Depress RUN.
- 8888 Parity error on tape buffer unload; depress RUN to ignore.
- 000A END card with undefined M address. Key into rA an address.
- 000B TYP card. Key into rL the type number, depress RUN.
- 1212 Normal stop for FLO. If flowcharting is to be bypassed, clear rA and RUN, elso just RUN.
- B9AH Normal stop at end of assembly. Depress M and RUN to continue assembly of another program. Otherwise, servo 5 contains a selfloading tape. Depress RUN to load it, or if flowcharting is used, to bring in pass 3.
- Pass 3 (flowcharting) stops:

A stop occurs after pass 3 has loaded itself; merely depress RUN.

- 1111 Tape reading error. Depressing RUN will try to read again, but depending on the tape error you may wish to execute a read backwards. RX contains the tape instruction, so type in G2 0605 000C or G2 0705 000C.
- 3333 Printer off normal. Fix, depress RUN.
- 4444 Parity error on tape buffer unload; depress RUN to ignore.
- 5555 Invalid flow chart. Depress RUN until program starts up again.

At the completion of Pass 3 a stop will occur. If you wish to load the object program depress RUN. Depress M and RUN to continue flowcharting of another program if you are doing multiple assemblies.

This list does not include stops encountered while QADAAD is loading itself from tape.

Errors on assembly listing

On the assembly listing any error indications will appear to the left of the line number. On the END card the error indication will be blank only if there were no errors detected earlier.

Errors are identified by the number of the field where the error occurred.

Number:	1	2	3	4	5	6	7	8
Field:	A	AH	М	MH	С	CH	OP	DK

There are two special error indications which are perhaps only apparent errors:

-indicates the latency time has come out pretty badly. This may or may not be significant.;indicates the H field instructions cannot be carried out; the next best alternative was tried. If the whole drum is unavailable 0000 is assembled.

The Automatic Documentation Feature

If automatic flowcharting is desired, the control operation FLO (not mentioned above) should appear at the head of the program. As soon as FLO appears, the remarks field of the listing is examined and has special significance. Three listings rathen than one are then produced:

- I. The assembly listing
- II. The algorithm listing
- III. The flow chart

Listings II and III are produced simultaneously as a second pass to the assembly.

The documentation is presented in a fixed format which has been designed for effective description. When flowcharting, the program is broken into logical segments called <u>sections</u>. For example, a subroutine constitutes a section. Each section is given a single alphabetic letter to identify it; thus there is section A, section B, etc. One flowchart is prepared for each section. In each section there are at most 99 subsections, each of which corresponds to a box on the chart. In section A, the subsections are A1, A2, ..., A10, All, etc. In each box of the flowchart appears the subsection number and a few "key words" which tell <u>generally</u> what is taking place. The key-words are usually rather vague until one is familiar with the algorithm listing which tells <u>specifically</u> what is going on. (The algorithm listing should if possible also tell <u>why</u> it is going on.)

The remarks field is split up between the listings, as follows:

- I The assembly listing contains the section names and with each subsection number the key words; and also special coding-oriented details.
- II The algorithm listing contains the section names and with each subsection number the specific algorithm descriptions.
- III The flow chart contains key words, condition branch names, appropriate boxes and connecting lines, and also line numbers from the assembly listing for cross reference.

The DK field

The first four columns of the remarks (cols 32-35) are of prime importance when flowcharting, and are called the DK or documentation key field. The following forms are used in DK fields:

1. GAAA Ordinarily the remarks are deleted from the assembly listing except section names and keywords, but G causes them to appear on the assembly listing, and not on the other listings. This is used to give coding-oriented details.

2. **AAAA** No special operation. The remainder of the remarks are part of the algorithm listing only.

3. $X_{\wedge\wedge\wedge}$ Same as $\wedge\wedge\wedge\wedge$ except this line is omitted from the assembly listing. This is used when remarks take more room than the machine language.

4. K. M Here K is the section letter. This line is the beginning of a new section, and causes a skip to next page on each listing.

5. Kn.A or Knn. Indicates a new subsection. The remarks field columns 36-56 contains the key words.

6. CODI These are the first four letters of "Coding Details" which is part of the format explained below.

7. TABL This indicates this is a title section, with no subsections, and it is part of the format explained below.

8. Anything else is a condition name, e.g. YES: or NO: This means

- a) If cols 36-40 are blank, this condition is to label the branch to the next box below.
- b) Otherwise this condition branches to the next name appearing in the rest of the remarks.

Within the remarks of the algorithm description, branches to any but the next subsection are indicated by prefixing the name by the symbol #. This symbol is deleted from the final listing. If a condition name preceded, it is a conditional branch to this place. If no condition name preceded on this subsection it is an unconditional branch.

These rules are best explained by example, and so a listing of a small sample program is attached to this report.

Format for algorithm listing:

P. NAME OF PROGRAM (used for long programs only) TABLE OF CONTENTS

A. NAME OF SECTION A

B. NAME OF SECTION B DESCRIPTION OF THE PROGRAM

- A. NAME OF SECTION A DETAILS ABOUT THE PURPOSE OF THIS SECTION AND GENERAL SUMMARY.
- A1. FIRST STEP WHAT HAPPENS AS THE FIRST STEP
- A3. NEXT STEP WHAT HAPPENS AS THE NEXT STEP. SUBSECTION NUMBERS MUST APPEAR IN ASCENDING ORDER, BUT NOT NECESSARILY SEQUENTIALLY

CODING DETAILS. AN OPTIONAL PART AT THE CLOSE OF A SECTION TELLS, FOR EXAMPLE, WHAT REGISTERS CONTAIN WHICH INPUTS AND OUTPUTS TO THIS SUBROUTINE.

1

INPUT AS PUNCHED ON CARDS

	FLO		٨.	SEPCH.
T0001	BLR 1000	1999		THIS SUBROUTINE SEARCHES THROUGH TABLE T
LOWER	EQU BOIA			TO SEE IF IT CAN FIND AN ENTRY MATCHING
UPPER	EQU BO2A			A GIVEN KEY.
KEY	EQU BO3A			
	ННН Н			·
SF RCH	STA KEY		A1.	INITIALIZE
	LDA#00000	10000		START OUT BY SETTING !LOWER! TO 1.
	STA LOWER			+UPPER+TO 1000.
	SHL 0300	1F.		THE TABLE IS TOOOL THROUGH TIUOD AND IS IN
1	STA UPPER			ASCENDING SEQUENCE.
•	ADD LOWER	3F		
. 3	ATL		Δ.2.	GET MIDPOINT
	MUL#00000	00045		SET MY TO (LOWER+UPPER)/2. MY WILL THUS
	LDX RA			APPROXIMATE THE MIDPOINT OF THE INTERVAL
	LDL LOWER			WHERE WE HAVE PINPOINTED THE SEARCH.
<i>a</i> .	TGR 2F		NO .	IF 'UPPER' IS LESS THAN 'LOWER', THE KEY
	TEQ 2F	NOT		ISTAN THE TABLE.
2	ADD	RA	OK:	
-	LDA TOODO	RA .		T(M) #KEY
	LDL KEY		H.J. •	COMPARE T(M) WITH THE SEARCH KEY.
	TEQ2 0000		50.	
	TGR 2F			IF EGUAL, WE#EXIT.
				IF GREATER. TO#A5.
	LDA RX	10000	LS:	
	ADD#00000	10000	A4 •	FIX LOWER
х.	STA LOWER			SET 'LOWER' TO M+1+AS T(M) IS TOO SMALL.
	ADD UPPER	38		T0#A2.
2	LDA RX		A5.	FIX UPPER
	SUB	18		SET 'UPPER' TO M-1, AS T(M) IS TOO BIG.
	CON 00000	10000		T0#A2.
			COD	ING DETAILS: AT ENTRY RB2 CONTAINS THE EXIT
			х	LOCATION AND RA CONTAINS THE KEYWORD.
			Х	IF FOUND. THE PLACE FOUND IS IN RX.
	ннн			IF NOT IN TABLE, EXIT OCCURS TO LOCATION'NOT!
NOT	HLT	*	Τ.	TEST.
TEST	LIR1 0000	T	Ţ1.	SET UP T
- T	IIR1 0001			FILL TABLE T: PUTTING 2I IN T(I).
	ADD RA			
	STA1T0000			
	IIR1 0000			
	ADD	-T		
	CON 99900	00000		
άT	LDA#00010	00000	T2.	SERCH 100.
	LIR2	SERCH		USE THE SEARCH ROUTINE TO SEE IF 100 IS IN.
	ADD#00000	10000	тз.	SERCH 101.
	LIR2&T	SERCH		SEARCH ALSO FOR 101 WHICH IS#NOT IN THE TABLE
	END TEST			
	FIN			

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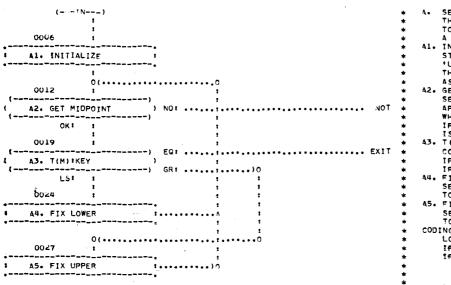
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0000						FL.0			A.	SERCH.
0001					T0001	91_R	1000	1999		
0002					LOWER	EQU	BO1A			
0003					UPPER	EQU	B02A			
0004					KEY	EQU	803A			
0 105						ннн	F	4		
(1006	4006	888 0 60	B03A	4010	SERCH	STA	KEY		41.	INITIALIZE
0007	4010	898 C 25	4012	4014		LDA#	00000	10000		
0008	4014	898 O 60	BOIA	4018		STA	LOWER			
0009	4018	898 O 37	0300	4024		SHL	0300	1F		
0010	4024	898 0 60	BO2A	4028	1	STA	UPPER			
0011	4028	898 0 70	BOIA	4033 .		ADD	LOWER	3F		
0012	4033	898 0 77	4033	4036	3	ATL			A2.	GET MIDPOINT
0013	4036	898 O 85	4039	4015		MUL#	00000	000A5		
0014	4015	BRB 0 05	000A	4019		LDX	RA	•		
0015	4019	B3B 0 30	BOIA	4023		LDL	LOWER			
0016	4023	838 0 87	4026	4226		TGR	2F			
0017	4226	838 0 82	4025	4029		TEQ	2F	NOT		
0018	4026	838 0 70	4229	000A	2	ADD		RA		
0019	4228	858 0 25	0999	4001		LDA	T0000		A3.	T (M) #KEY
0020	4001	838 0 30	B03A	4005		LDL	KEY			
0021	4005	BBB 1 82	0000	4009		TEQ2	0000			
0022	4009	858 0 87	4212	4412		TGR	2F			
0023	4412	838 0 25	000C	4016		LDA	RX			
0024	4016	89B 0 70	4218	4021		ADD#	00000	10000	Δ4.	FIX LOWER
0025	4021	898 0 60	BOIA	4025		STA	LOWER			
0026	4025	898 0 70	B024	4033		ADD	UPPER	38		
0027	4212	838 0 25	0000	4216	2	LDA	RX		Δ5.	FIX UPPER
0028	4216	838 0 75	4419	4024		SUB		18		
0029	4418	838 0 00	0001	0000		CON	00000	10000		
0600										
0033						HHH				

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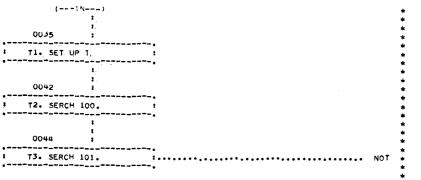
0034 0035 0036 0037	4029 0035 0038 0042	858 0 67 858 0 08 858 0 06 858 0 70	4029 0000 0001 0004	4029 0038 0042 0047	NOT TEST T	HLT LIR1: 0000 IIR1: 0001 ADD RA	* . -T	T. TEST. T1. SET UP T
0038 0039 0040	0047 0001 0005	838 0 64 838 0 0G 838 0 70	0999 0000 0007	0001 0005 0038		STA1 T0000 IIR1 0000 ADD	-1	
0041 0042 0043	0007 0039 0043	838 0 99 838 0 25 838 1 02	9000 0041 0247	0000 0043 4006	&T	CON 99900 LDA# 00010 LIR2	00000 00000 SERCH	T2. SERCH 100.
0044 0045 0046	0247 0052	898 0 70 898 1 02	0049 0039	0052 4006		ADD# 00000 LIR2 &T END TEST	10000 SERCH	T3. SERCH 101.

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- A. SERCH. THIS SUBROUTINE SEARCHES THROUGH TABLE T TO SEE IF IT CAN FIND AN ENTRY MATCHING A GIVEN KEY.
 AI. INITIALIZE START OUT BY SETTING 'LOWER' TO 1. 'UPPER'TO 1000. THE TABLE IS TOOOL THROUGH TLOOD AND IS IN ASCENTING SEQUENCE.
 A2. GET MIDPOINT SET 'M' TO (LOWER+UPPER)/2. 'M' WILL THUS APPROXIMATE THE MIDPOINT OF THE INTERVAL WHERE WE HAVE PINPOINTOF THE INTERVAL WHERE WE HAVE PINPOINTOF THE SEARCH. IF 'UPPER' IS LESS THAN 'LOWER', THE KEY IS NOT IN THE TABLE.
 A3. T(M):KEY COMPARE T(M) WITH THE SEARCH KEY. IF EQUAL: WE EXIT. IF GREATER, TO A5.
 A4. FIX LOWER SET 'LOWER' TO M+1.AS T(M) IS TOO SMALL. TO A2.
 CODING DETAILS: AT ENTRY RB2 CONTAINS THE EXIT LOCATION AND RA CONTAINS THE KEYWORD. IF NOT IN TABLE. EXIT CCCURS TC LOCATION'NOT'

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T. TEST.
T. SET UP T
FILL TABLE T. PUTTING 2I IN T(I).
T2. SERCH 100.
USE THE SEARCH ROJTINE TO SEE IF 100 IS IN.
T3. SERCH 101.
SEARCH ALSO FOR 101 WHICH IS NCT IN THE TABLE

OUTPUT OF SYMBOL-USAGE ROUTINE

-T	0036	0035	0040		
άT	0042	0045			
KEY	0004	0006	0020		
LOWER	0002	0008	0011	0015	0025
NOT	0034	0017			
RA	0014	0013	0037		
RX	0023	0027			
SERCH	0006	0043	0045		
T0000	0001	0019	0038		
TEST	0035	0046			