QAADAD

AN ASSEMBLY SYSTEM FOR THE SOLID STATE II

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"Quick-And-Dirty Assembler And Documentor"

QADAAD is an assembly program for a Univac Solid State II ( 80 column) configuration with at least 5000 words of drum memory. It assembies from tape to tape at a speed of about 600 lines per minute. QADAAD provides facilities for hand minimization of latency, for relocatable output, and a special optional feature which draws flow charts of the program as an extra pass.

## Card Format

| cols 1-5 | A field: location |
| :--- | :--- |
| col 6 | AR field: relocation code for A |
| cols 7-9 | AH field: latency code for A |
| cols 10-12 | OP field: symbolic operation code |
| col 13 | IR field: index register or sign |
| cols 14-18 | M field: M-address |
| col 19 | MR field: relocation code for M |
| cols 20-22 | MH field: latency code for M |
| cols 23-27 | C field: C-address |
| col 28 | CR field: relocation code for C |
| cols 29-31 | CH field: latency code for C |
| cols 32-80 | remarks. |

## Symbolic Fields

The A, M, and C fields contain information of the following types:
(here $\boldsymbol{\wedge}$ indicates a blank column)

1. Regional addresses, Xnnnn where $X$ is any non-blank character and n represents a digit, e.g, A0000. A0003 is three locations greater than A0000. A9999 is one location less than A0000.
2. Absolute addresses, $\wedge$ NNNN where $N$ represents either a digit or one of the undigits $A, B, C, F, G, H$. This address is simply reproduced in
the output. Examples: 4078, B2AB, H66H.
3. Local addresses. In $A$ these are of the form $n \wedge \wedge A \wedge$, and in $M$ or . $C$ they are of the form $n F A \wedge \wedge$ or $n B \wedge \wedge \wedge$ (referring to the next or previous $A$ of n^A^A respectively).
4. "Self": *aAAA refers to $A$ address. If this address occurs in A or on a control operator card it refers to the previous $A$ address.
5. Pair address, any address of the form \&WWWW or -WWWW where W's are arbitrary, indicates two adjacent addresses as used in "c + 1 conditions." The - is the first of the pair, the \& is the second or overflow address.
6. Blank address. In A field, indicates the address corresponding to a blank $M$ or $C$ address on the preceding instruction. In $M$ if the address is ignored by the Op (e.g. ATL or HLT) this is equivalent to *. In $C$ if the address is ignored (e.g. CLA) the $C$ address is set equal to $M$. Otherwise the next $A$ field must be blank and it refers to this next A address.
7. Symbolic address. This is any address other than the above, except the leftmost character must be nonblank and non-numeric. Several symbolic addresses are pre-defined: RA is 000 A , RL is 000 B , RX is 000 C ; $R B 4, R B 5, \ldots, R B 9$ are respectively $B 2 A B, B 3 A B, B 5 A B, B 6 A B, B 7 A B$, B8AB. If multiple assembly is used these predefined symbols disappear, however.

## H Fields

The $\mathrm{AH}, \mathrm{MH}$, and CH fields are relevant only if the corresponding address is undefined, i.e. if it has not appeared earlier. When the address is undefined, the $H$ field is used to control storage allocation as follows:

1. ^^^ or D^A Choose the best available place on the drum.
2. Han Choose the best available place on the high-speed bands.
3. C^1 Choose a place in the B001-B999 core storage area.
4. nnn Assign on this level on the drum.
5. Ann Assign on this level on the high-speed bands.
6. $+n n$ Assign $n n$ higher than the normal rules would say.

## R Fields

The $A R, M R$, and $C R$ fields are merely transferred to the output area and are used to prepare FORTRAN subprograms. The FORTRAN assignment ignores the 4-bit and then has the following code:

| 0 absolute | 5 external reference |
| :---: | :--- |
| 1 unique storage | 6 in AR for ALF table entry |
| ^ or 2 program storage | 7 cross reference |
| 3 common storage | 8 special program storage |

## IR Field

The IR field is normally used to specify index register modification. If blank there is no indexing, otherwise the digits 1-9 are used to indicate modification by RB1-RB9. (In the control operators CON NUM ZON and ALF which provide constants the IR field is used to denote the sign instead, and then the sign is 0 if $I R$ is blank, the sign is 1-9 if the IR field is 1-9.)

A convenient literal constant feature is included, so that if the IR field contains the character \# the $M$ and $C$ fields are filled with a ten-digit constant which is positive. This constant is assembled separately, then the $M$ address of the instruction refers to this constant, and $C$ address is treated as if it were blank.

Example: ADD\#01234 56789
is equivalent to ADD CONST and the line
CONST CON 012340567890
appearing later. (Notice that the $M R$ and $C R$ fields of the constant are set to 0; the AR field is set to the MR field of the instruction.) The MH and CH fields of the instruction retain their normal significance.

## Warning

Do not attempt to assembly anything into location 0000; this location is sacred to QADAAD. As a matter of fact the loading routine supplied on the output tape goes into band 0000 so it is wise not to assemble into any of the locations 0000-0399.

## OP Field

The operation code field is filled with either a three-letter mnemonic op-code as used in the $S 4$ assembly system, or it is one of the control operators listed below. Operators which require a 4-bit in the sign digit are not provided.
ann. This is used for comments only.

CPY. The card-to-tape pass uses this. $M$ and $C$ are absolute addresses which refer to line numbers on the previous assembly listing; lines $M$ through $C$ are copied from tape in place of this card.

FIN. This is the last card of a program, signalling the end for the card-to-tape pass.

END. This is the second-last card of a program, signalling the end of the program. It can also be used to separate programs in a multi-program assembly. The $M$ address must be defined, and the loading routine will be set to halt and transfer to this address.

BLR and BLA. Reserve or unreserve drum locations starting at $M$ and ending at $C$. $M$ and $C$ must be defined addresses on the drum. If the $C H$ field is nonblank, take $n$ to be CH mod 100, indicating that we reserve or unreserve every $n$-th word. If the $C H$ field is blank it means take $n$ to be 1. If $A$ is nonblank an EQU operation is also performed (See EQU). If $C$ is blank it is taken equal to $M$.

COR. Reserve core locations in the B001-B999 area. M must de defined. M locations in the core are reserved. If $A$ is non-blank an EQU to the first address of this block is also performed (see EQU). For example, R0001 COR 0030 reserves thirty core locations as the locations R0001 through R0030.

EQU. A must be an undefined regional, forward local, or symbolic address. (Not a pair address,) It is then defined to equal the equivalent of $M$ which must be defined.

HHH. All succeeding blank H-fields are overridden by the contents of the MH field on this card. For example, HHH a^AaA^C is used to start assignment of undefined address in core. HHHanaanaH is used to assign in high-speed storage.

CON, NUM, ZON, ALF. These provide faur types of constants. The A field is treated exactly like the $A$ field of an instruction. The $M$ and $C$ fields specify a 10-digit constant, except with ALF only $M$ is used. CON means treat the contents of $M$ and $C$ as digits or the undigits $A, B, C, F, G, H$. NUM and ZON mean to treat $M$ and $C$ as an alphabetic constant and take the numeric or zone word of thís constant (expressed in machine code). ALF takes the $M$ field and produces a constant zzzzznnnnn where zzzzz is the zone part and nnnnn is the numeric part. On all four constant operators the IR field is used to specify the sign.

PAT. Print availability table. This causes 50 lines to be printed. Each line has five entries:
Level L
L
Level L+100
Level L+150

The availability word for a level contains 40 bits in machine code, corresponding to the status of the word in each band with the following band assignments:

| 00 | 02 | $\ldots$ | 18 |
| :--- | :--- | :--- | :--- |
| 20 | 22 | $\ldots$ | 38 |
| 40 | 42 | $\ldots$ | 58 |
| 6062 | $\ldots$ | 78 |  |

As an example, at the beginning of assembly PAT would produce

$$
\text { BGGGG99999 GGGGG99999 } 0000000000 \text { GGGGG99999 GGGGG99999 }
$$

$$
\text { GGGGG99999 GGGGG99999 } 0000010000 \quad \text { GGGGG99999 GGGGG99999 }
$$

(locations 0001-4999 are available).
TYP, ON , OFF. Three control operators are provided in case a person wants to use a single assembly deck to produce several almost identical versions, without
going through the card-to-tape pass each time. TYP causes the computer to stop at assembly time and then $R L$ should be set to 000000 nnnn . The operators OFF and ON have a defined $M$ address and are effective only: if the equivalent of $M$ is nnnn matching a previous TYP. They temporarily shut assembly off or on. The FORTRAM system program is coded for five types:

$$
\begin{aligned}
& \text { 8001:SS80, I } \\
& \text { 8002:SS80, II } \\
& 9000: \mathrm{SS90}, \text { card } \\
& 9001: \mathrm{SS90}, \text { tape I } \\
& 9002: \mathrm{SS90}, \mathrm{II}
\end{aligned}
$$

Note: The line B8AH NEW100001 00000 causes some assembly into the B00A-B99F area of core storage. Other uses of the control op NEW will not be listed here.

## Operating Instructions

Loading the assembler: Loading the object program:

| rC | G2 | 0400 | $000 A$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $r A$ | $F 6$ | $B 000$ | B000 |$\quad$| rC | $G 2$ | 0500 |
| :--- | :--- | :--- |
|  | rA | $F 6$ |

(The assembler is presently set up for an 8800 word drum.)

Mount QADAAD tape on unit 4, scratch tapes on units 3 and 5. If there is tape input, put it on unit 2. (The card-to-tape pass copies from cards and/or unit 2 to unit 3. If this phase is omitted, mount a good input tape on unit 3 and forget unit 2.) If flowcharting is desired, you will need scratch tapes on units 6 and 7 also.

As soon as the assembler has loadeditself, it stops. Depress RUN if the card-to-tape pass is to be used else depress $M$ and RUN to skip this pass. Wher the second pass has loaded, it stops, and the operator shoula depress RUN. The computer should next stop with HLT. B9AH, but several other stops may occur.

Pass 1 (card-to-tape) stops:
0001 Comparison failure on card reader, The bad cards have been diverted to stacker 2. Reload them, depress RUN.

0002 Card reader off normal or empty. Fix, depress RUN.

0005 Error in writing tape. This is really unfortunate. If RUN is depressed the condition is ignored.

0006 Error in reading tape。 If RUN is depressed reading is attempted again, in the opposite direction. Gain should be varied by the operator.

Pass 2 (assembly) stops:
1111 Line numbers on input tape not sequential, indicates a bad situation on the input tape (unit 3)

2222 Tape error. Depressing RUN will cause the last tape instruction (read or write) to be tried again, with no backspacing provided. The last tape instruction appears in RA, and it is usually a G2 0300. In this case, restart by typing G2 0305 000A into rC , depress general clear and RUN.

3333 Printer off normal. Fix, depress RUN.
4444 Tape 3 not yet ready. Depress RUN.
8888 Parity error on tape buffer unload; depress RUN to ignore.
OOOA END card with undefined $M$ address. Key into rA an address.
OOOB TYP card. Key into rL the type number, depress RUN.
1212 Normal stop for FLO. If flowcharting is to be bypassed, clear rA and RUN, elso just RUN.

B9AF Normal stop at end of assembly. Depress $M$ and RUN to continue assembly of another program. Otherwise, servo 5 contains a selfloading tape. Depress RUN to load it, or if flowcharting is used, to bring in pass 3.

Pass 3 (flowcharting) stops:
A stop occurs after pass 3 has loaded itself; merely depress RUN.
1111 Tape reading error. Depressing RUN will try to read again, but depending on the tape error you may wish to execute a read backwards. RX contains the tape instruction, so type in G2 0605 000C or G2 0705 000C.

3333 Printer off normal. Fix, depress RUN.
4444 Parity error on tape buffer unload; depress RUN to ignore.
5555 Invalid flow chart. Depress RUN until program starts up again. At the completion of Pass 3 a stop will occur. If you wish to load the object program depress RUN. Depress M and RUN to continue flowcharting of another program if you are doing multiple assemblies. This list does not include stops encountered while QADAAD is loading itself from tape.

## Frrors on assembly Iisting

On the assembly listing any error indications will appear to the left of the line number. On the END card the error indication will be blank only if there were no errors detected earlier.

Errors are identified by the number of the field where the error occurred.

| Number: | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Field: | A | AH | M | MH | C | CH | OP | DK |

There are two special error indications which are perhaps only apparent errors:
-indicates the latency time has come out pretty badly. This may or may not be significant.
;indicates the $H$ field instructions cannot be carried out; the next best alternative was tried. If the whole drum is unavailable 0000 is assembled.

## The Automatic Documentation Feature

If automatic flowcharting is desired, the control operation FLO (not mentioned above) should appear at the head of the program. As soon as FLO appears, the remarks field of the listing is examined and has special significance. Three listings rathen than one are then produced:
I. The assembly listing
II. The algorithm listing
III. The flow chart

Listings II and III are produced simultaneously as a second pass to the assembly.

The documentation is presented in a fixed format which has been designed for effective description. When flowcharting, the program is broken into logical segments called sections. For example, a subroutine constitutes a section. Each section is given a single alphabetic letter to identify it; thus there is section $A$, section $B$, etc. One flowchart is prepared for each section.

In each section there are at most 99 subsections, each of which corresponds to a box on the chart. In section A, the subsections are A1, A2, ..., A10, All, etc. In each box of the flowchart appears the subsection number and a few "key words" which tell generally what is taking place. The key-words are usually rather vague until one is familiar with the algorithm listing which tells specifically what is going on. (The algorithm listing should if possible also tell why it is going on.)

The remarks field is split up between the listings, as follows:
I The assembly listing contains the section names and with each subsection number the key words; and also special coding-oriented details.

II The algorithm listing contains the section names and with each subsection number the specific algorithm descriptions.

III The flow chart contains key words, condition branch names, appropriate boxes and connecting lines, and also line numbers from the assembly listing for cross reference.

## The DK field

The first four columns of the remarks (cols 32-35) are of prime importance when flowcharting, and are called the DK or documentation key field. The following forms are used in DK fields:

1. G $\mathbf{1 1 \wedge}$ Ordinarily the remarks are deleted from the assembly listing except section names and keywords, but $G$ causes them to appear on the assembly listing, and not on the other listings. This is used to give coding-oriented details.
2. An^^ No special operation. The remainder of the remarks are part of the algorithm listing only.
3. X^^^ Same as ^ヘ^^n except this line is omitted from the assembly listing. This is used when remarks take more room than the machine language. 4. K.^^ Here $K$ is the section letter. This line is the beginning of a new section, and causes a skip to next page on each listing.
4. Kn.a or Knn. Indicates a new subsection. The remarks field columns 36-56 contains the key words.
5. CODI These are the first four letters of "Coding Details" which is part of the format explained below.
6. TABL This indicates this is a title section, with no subsections, and it is part of the format explained below.
7. Anything else is a condition name, e.g. YES: or NO: This means
a) If cols 36-40 are blank, this condition is to label the branch to the next box below.
b) Otherwise this condition branches to the next name appearing in the rest of the remarks.

Within the remarks of the algorithm description, branches to any but the next subsection are indicated by prefixing the name by the symbol \# 。 This symbol is deleted from the final listing. If a condition name preceded, it is a conditional branch to this place. If no condition name preceded on this subsection it is an unconditional branch.

These rules are best explained by example, and so a listing of a small sample program is attached to this report.

Format for algorithm listing:
P. NAME OF PROGRAM (used for long programs only)

TABLE OF CONTENTS
A. NAME OF SECTION A
B. NAME OF SECTION B

DESCRIPTION OF THE PROGRAM
A. NAME OF SECTION A

DETAILS ABOUT THE PURPOSE OF THIS SECTION
AND GENERAL SUMMARY。
A1. FIRST STEP
WHAT HAPPENS AS THE FIRST STEP
A3. NEXT STEP
WHAT HAPPENS AS THE NEXT STEP. SUBSECTION
NUMBERS MUST APPEAR IN ASCENDING ORDER, BUT
NOT NECESSARILY SEQUENTIALLY

CODING DETAILS. AN OPTIONAL PART AT THE CLOSE OF A SECTION TELLS, FOR EXAMPLE, WHAT REGISTERS CONTAIN WHICH INPUTS AND OUTPUTS TO THIS SUBROUTINE.


```
A. SERCH.
    THIS SIJGROUTINE SEARCHES THROUGH TABLE T
    TO SEE TF IT CAN FIND AN ENTRY MATCHING
    A GIVEN KEY.
    A1. INITIALIZE
    START OUT BY SETTING 'LOWER: TO 1,.
    - UPPER.TO 1000.
    THE TAOLE IS TOOO1 THROUGH TIUOO AND IS IN
    ASCENDTNG SEQUENCE.
A.2. GET MIDPOINT
    SET 'M' TO (LOWER+UPPER)/2. 'M' WHLL THUS
    APOROXTMATE THE MIDPOINT OF THE INTERVAL
    WHERE WF HAVE PINPOINTED THE SEARCH.
NO: IF 'UPDFR' IS LESS THAN 'LOWER', THE KEY
    IS#NOT IN THE TABLE.
OK:
A3. T(M):KEY
    COMPARE T(M) WITH THE SEARCH KEY.
EQ: IF EGUAL, WE#EXIT.
GR: IF GREATER. TOHA5.
LS:
A4. FIX LOWER
    SET !LOWER: TO M+1:AS T(M) IS TOO SMALL.
    TOHA2.
A5. FIX UPDER
    SET 'UPPER' TO M-1, AS T(M) IS TOO BIG.
    TOHA2.
CODING OETAILS: AT ENTRY RBZ CONTAINS THE EXIT
X LOCATION AND RA CONTAINS THE KEYWORD.
x IF FCUNO: THE PLACE FOUND IS IN RX.
    IF NOT IN TABLE, EXIT OCCURS TO LOCATION'NOT.
T. TEST.
T1. SET UP T
    FILL TABLE T, PUTTING 2I IN T(I).
    T2. SERCH 100.
        USE THE SEARCHI ROUTINE TO SEE IF 100 IS IN.
    T3. SERCH 101.
    SEARCH ALSO FOR 101 WHICH IS#NOT IN THE TARLE
```



| 0034 | 4029 | 338 | 0 | 67 | 4029 | 4029 | NOT | HLT |  | * | T. | TEST. |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0035 | 0035 | B98 | 0 | OB | 0000 | 0038 | TEST | LIR1 | 0000 | -T | T1. | SET UP | T |
| 0036 | 0038 | 898 | 0 | OG | 0001 | 0042 | - ${ }^{\text {T }}$ | IIR1 | 0001 |  |  |  |  |
| 0037 | 0042 | 88B | 0 | 70 | 0004 | 0047 |  | ADD | RA |  |  |  |  |
| 0038 | 0047 | 838 | 0 | 64 | 0999 | 0001 |  | STAI | T0000 |  |  |  |  |
| 0039 | 0001 | 838 | 0 | OG | 0000 | 0005 |  | IIRI | 0000 |  |  |  |  |
| 0040 | 0005 | B8B | 0 | 70 | 0007 | 0038 |  | ADD |  | - $T$ |  |  |  |
| 0041 | 0007 | 898 | 0 | 99 | 9000 | 0000 |  | CON | 9990 | 00000 |  |  |  |
| 0042 | 0039 | 398 | 0 | 25 | 0041 | 0043 | \& $T$ | LDAA | 000年0 | 00000 | T2. | SERCH | 100. |
| 0043 | 0043 | 898 | 1 | 02 | 0247 | 4006 |  | LIR2 |  | SERCH |  |  |  |
| 0044 | 0247 | в98 | 0 | 70 | 0049 | 0052 |  | ADD* | 00000 | 10000 | T3. | SERCH | 101. |
| $\begin{aligned} & 0045 \\ & 0046 \end{aligned}$ | 0052 | B8B | 1 | 02 | 0030 | 4006 |  | $\begin{aligned} & \text { LIRZ } \\ & \text { END } \end{aligned}$ | $\begin{aligned} & \& T \\ & T E S T \end{aligned}$ | SERCH |  |  |  |


A. SERCH. THIS SUBRONTIVE SEARCHES THROUGH TABLE ${ }^{\top}{ }^{\top}$
TO SEE IF IT CAN FIND AN ENTRY MATCHING TO SEE IF IF
A GIVEN KEY
41. INITIALIZE

START OUT BY SETTING 'LOWER' TO 1.

- UPPERITO 1000.

THE TABLE IS TOOO1 THROUGH T1000 AND IS IN ASCENTING SEQUENCE.
A2. GET MIDPOINT
SET MM TO (LOWER+UPPER)/2. 'MV WILL THUS APPROXIMATE THE MIDPOINT OF THE YNTERVAL WHERE WE HAVE PINPOINTED THE SEARCH. IF 'UPPER" IS LESS THAN 'LOWER'. THE KEY is not in the table.
43. T (M) :KEY

COMPARE T(M) WITH THE SEARCH KEY. IF EQUAL, WE EXIT.
IF GREATER, TO AS.
44. FIX LOWER

- EIX LOWER $\quad$ SET LOWER" $T O+1$ :AS T(M) IS TOO SMALL. TO A2.

45. EIX UTPER
SET UPPPER: TO M-1. AS T(M) IS TOO BIG.

CODING DETAILS: AT ENTRY RB2 CONTAINS THE EXIT LOCATION AND RA CONTAINS THE KEYWORD.
IF FOUND. THE PLACE FOUND IS IN $R X$. If HOT IN TABLE: EXIT OCCURS TC LOCATIONDNOT.


OUTPUT OF SYMBOL-USAGE ROUTINE

| -T | 0036 | 0035 | 0040 |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| KT | 0042 | 0045 |  | 015 | 0025 |
| KEY | 0004 | 0005 | 0020 |  |  |
| LOWER | 0002 | 0003 | 0011 | 003 |  |
| NOT | 0034 | 0017 |  |  |  |
| RA | 0014 | 0013 | 0037 |  |  |
| KX | 0023 | 0027 |  |  |  |
| SERCH | 0006 | 0043 | 0045 |  |  |
| TOOOO | 0001 | 0019 | 0038 |  |  |
| TEST | 0035 | 0046 |  |  |  |

